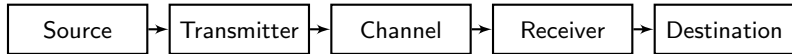


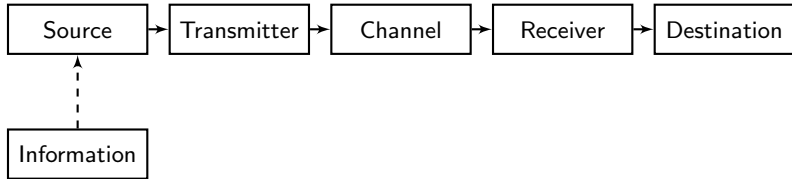


AVD623: Communication Systems-II  
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Dept. of Avionics  
**Lecture 2**

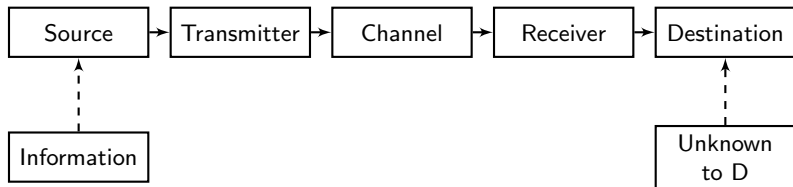
## A basic communication system



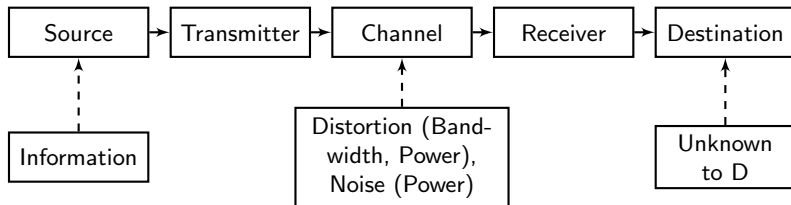
## A basic communication system



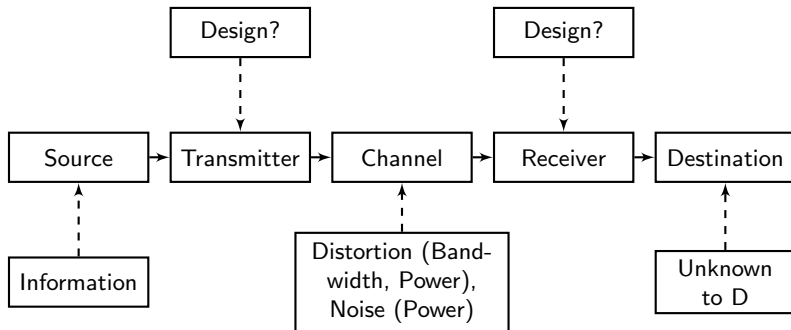
## A basic communication system



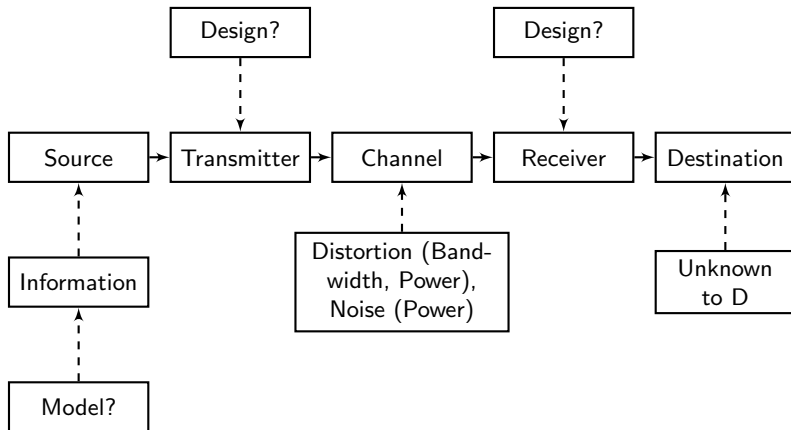
## A basic communication system



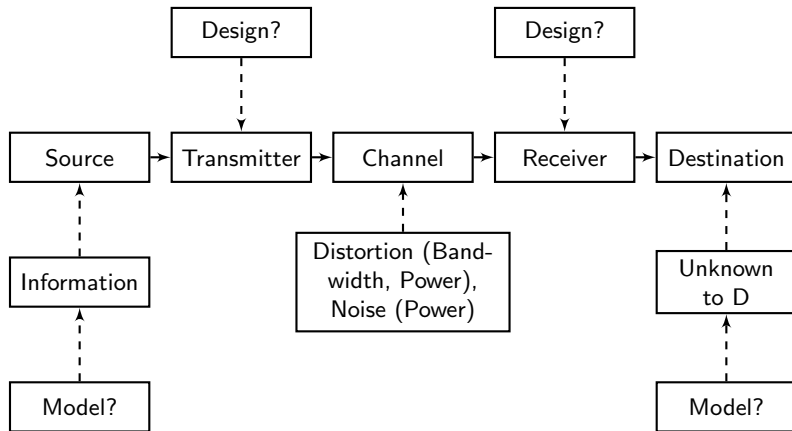
## A basic communication system



# A basic communication system

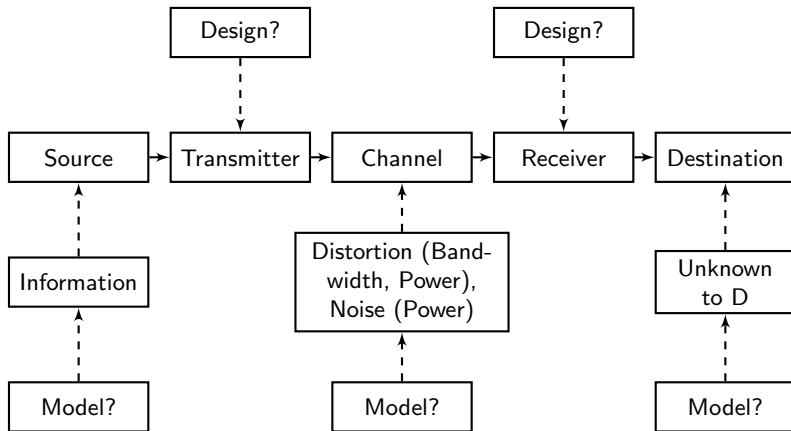


# A basic communication system

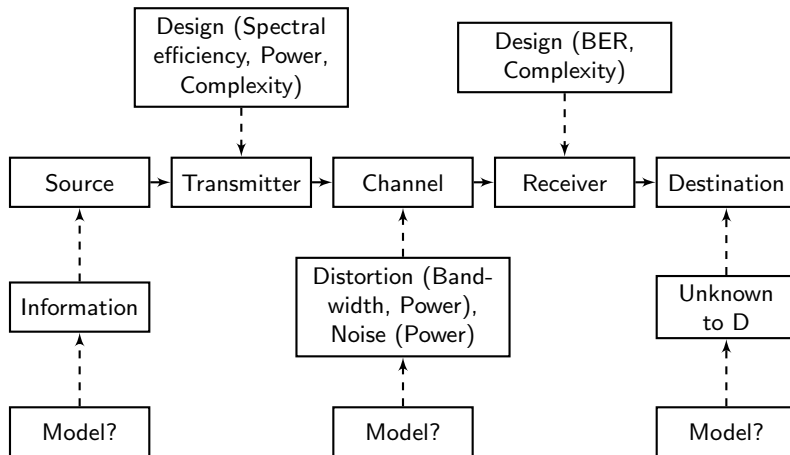




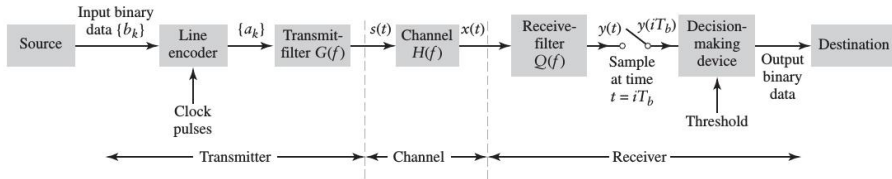
# A basic communication system



# A basic communication system



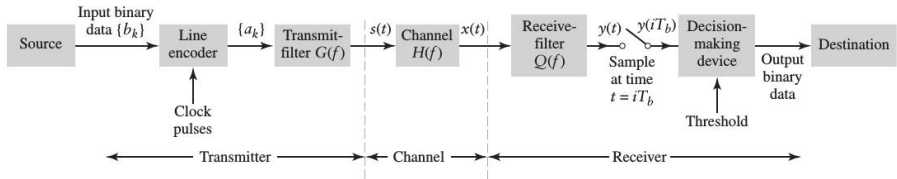
## Recall: Digital transmission system



- ▶ Source of digital information - characterized by bit duration  $T_b$  (or bit rate)
- ▶ Converted into a line code whose levels are represented by  $a_k$  (say  $-A$  and  $+A$ )
- ▶ Further transformation of  $a_k$  to “match” the signal to the channel (what if the channel were bandpass?)
- ▶ We obtain a continuous time signal  $s(t)$  which is transmitted over the channel
- ▶ At the receiver need to convert it into a digital signal - so synchronized sampling, usually at rate  $T_b$
- ▶ A decision device decides whether 0 or 1 was transmitted
- ▶ **Let us think about a baseband digital transmission system**

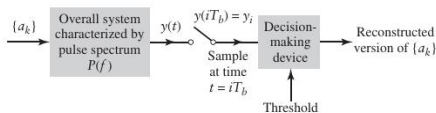


## An effective pulse shape $p(t)$



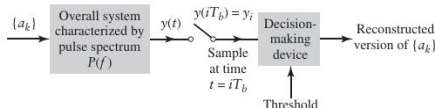
- ▶ Note that all filters and the channel are assumed to be LTI
- ▶ Let us think about  $a_k$ -s as an impulse train
- ▶ Then  $s(t) = \sum_{k=-\infty}^{\infty} a_k g(t - kT_b)$  since we are transmitting every  $T_b$  secs.
- ▶  $x(t) = s(t) \star h(t)$
- ▶  $y(t) = x(t) \star q(t)$
- ▶  $y(t) = \sum_{k=-\infty}^{\infty} a_k p(t - kT_b)$ , where  $p(t) = g(t) \star h(t) \star q(t)$
- ▶ Or  $P(f) = G(f)H(f)Q(f)$

## An effective pulse shape $p(t)$



►  $y(t) = \sum_{k=-\infty}^{\infty} a_k p(t - kT_b)$ , where  $p(t) = g(t) \star h(t) \star q(t)$

# Intersymbol interference problem



- ▶ At the sampling instants  $y(iT_b)$  we have  $y(iT_b) = \sum_{k=-\infty}^{\infty} a_k p((i - k)T_b)$  (notation: pulse is centered at zero)
- ▶ Suppose  $y_i = y(iT_b)$  and  $p_i = p(iT_b)$
- ▶  $y_i = \sum_{k=-\infty}^{\infty} a_k p_{i-k}$
- ▶ We need  $y_i = p_0 a_i$  for all  $i$ . Let us say that  $p_0 = \sqrt{E}$
- ▶ What we have is  $y_i = \sqrt{E} a_i + \sum_{k \neq i} a_k p_{i-k}$



- ▶ We have to design  $p(t)$  such that  $y_i = \sqrt{E} a_i$
- ▶  $p(t)$  has to be designed so that  $P(f)$  has minimum bandwidth
- ▶ Designing  $p(t)$  is called pulse shaping

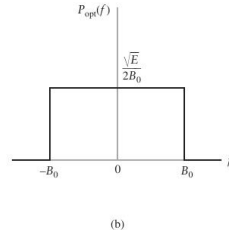
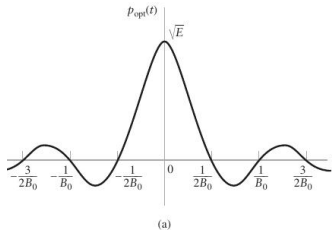


## Nyquist channel

- ▶ If  $y_i = p_0 a_i$  for every  $i$  then we require that

$$p_n = \begin{cases} \sqrt{E}, & \text{for } n = 0, \\ 0, & \text{otherwise.} \end{cases}$$

- ▶ Note that  $p_n = p(nT_b)$
- ▶ Is it possible to get  $P(f)$ ? Assuming that  $P(f)$  is bandlimited
- ▶ Consider the choice of  $p(t) = \text{sinc}\left(\frac{t}{T_b}\right)$
- ▶ With  $B_0 = \frac{1}{2T_b}$  we have the following optimal pulse shape  $p_{\text{opt}}(t)$



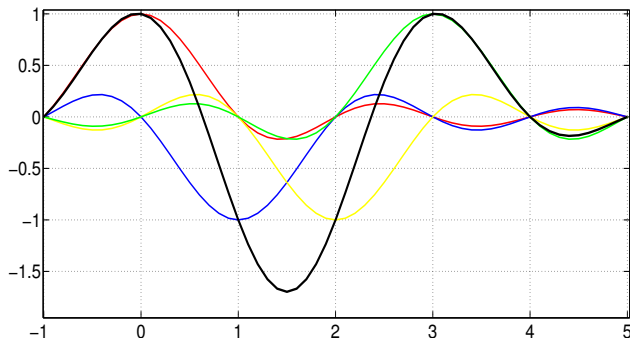
- ▶ The PAM system with  $P_{\text{opt}}(f)$  is called the **Nyquist channel**
- ▶ The bandwidth  $B_0$  is called the **Nyquist bandwidth**



## Nyquist channel pulse shaping - issues



- ▶ The transfer function  $P(f)$  is not realizable
- ▶ Issue of timing jitter



- ▶ Suppose sampling instants at which decoding is done has a *jitter*. Then is correct decoding possible?

## Raised cosine pulse shaping

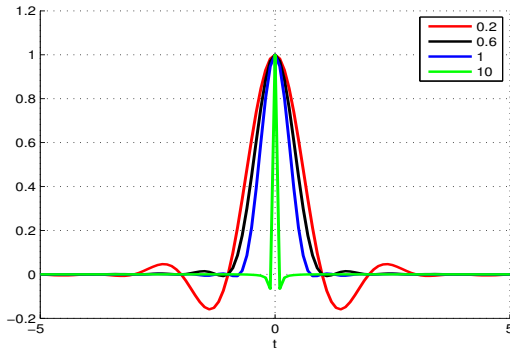


- ▶ The problem with sinc pulses -  $\frac{1}{t}$  decay
- ▶ How to increase the decay rate?

## Raised cosine pulse shaping

- ▶ The problem with sinc pulses -  $\frac{1}{t}$  decay
- ▶ **How to increase the decay rate?**
- ▶ *Damp* the sinc pulse using a window function
- ▶ Raised cosine pulse shape (actually damped sinc pulse shape)

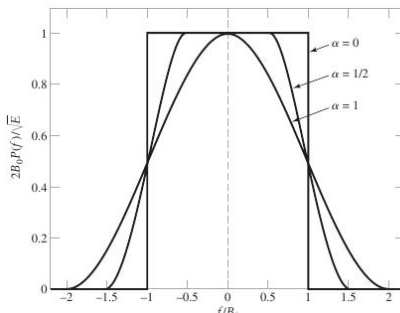
$$p(t) = \sqrt{E} \text{sinc}(2B_0 t) \frac{\cos(2\pi\alpha B_0 t)}{1 - (4\alpha B_0 t)^2}$$



- The F.T of  $p(t)$  is

$$P(f) = \begin{cases} \frac{\sqrt{E}}{2B_0}, & \text{for } |f| \leq f_1, \\ \frac{\sqrt{E}}{4B_0} \left[ 1 + \cos \left\{ \frac{\pi(|f| - f_1)}{2(B_0 - f_1)} \right\} \right], & \text{for } f_1 < |f| < 2B_0 - f_1, \\ 0, & \text{o/w.} \end{cases}$$

- $\alpha = 1 - \frac{f_1}{B_0}$ .  $\alpha$  is the roll-off factor.
- Bandwidth of the pulse is  $2B_0 - f_1$  or  $B_0(1 + \alpha)$





► Let  $r_b = \frac{1}{T_b}$

Scheme	Bandwidth	Power	Rate	Timing Jitter
Rectangular	$r_b$	95%	$r_b$	Robust
Sinc	$\frac{r_b}{2}$	100%	$r_b$	Weak
Raised cosine	$\frac{r_b}{2}(1 + \alpha)$	100%	$r_b$	less than Rect

► Read about square root raised cosine pulse shaping