

Part I — Recognition (20)

1. A variable's **lifetime** is
 - a) How long its name is visible in the source
 - b) The time from allocation until deallocation of its storage
 - c) The time it stays in a register
 - d) The time it is cached by CPU
2. A variable's **scope** is
 - a) The textual region where its binding is visible
 - b) The OS process lifetime
 - c) The same as lifetime
 - d) The stack frame where it's stored
3. In C/C++, **global** objects (with external linkage) are allocated in
 - a) Stack segment
 - b) Static storage (data/BSS)
 - c) Heap
 - d) Register file
4. A **static local** (e.g., `static int z;`) has
 - a) Stack storage, lifetime = call
 - b) Static storage, lifetime = whole program
 - c) Heap storage, lifetime = call
 - d) Static storage, lifetime = function call only
5. An object from `malloc/new` has
 - a) Stack storage, lifetime = block
 - b) Heap storage, lifetime = until `free/delete`
 - c) Static storage, lifetime = program
 - d) Heap storage, lifetime = function call
6. A **dangling reference** happens when
 - a) Two names alias the same cell
 - b) A pointer refers to already-freed storage
 - c) A pointer is `NULL`
 - d) A global shadows a local
7. A **memory leak** happens when
 - a) You write past array bounds
 - b) You lose the last reference to allocated heap memory
 - c) You allocate on stack
 - d) You use `static`
8. An **activation record (stack frame)** typically contains
 - a) Only machine code
 - b) Return address, parameters, locals, saved registers, links
 - c) Only global variables
 - d) Only heap pointers
9. The **dynamic link** in a frame points to
 - a) Callee's frame

- b) Caller's frame (for return/unwind)
 - c) Global table
 - d) Heap block header
10. In lexically nested languages, the **static link/access link** enables
- a) Jump to caller
 - b) Access to variables in enclosing lexical scopes
 - c) Faster malloc
 - d) Garbage collection
11. **Static (lexical) scope** resolves names using
- a) The call chain at runtime
 - b) The source nesting at compile time
 - c) The loader's symbol map
 - d) OS environment variables
12. **Dynamic scope** resolves names using
- a) Source nesting only
 - b) The most recent active binding along the call chain
 - c) The global table only
 - d) The import list only
13. The **referencing environment** of a statement under static scope is
- a) Locals only
 - b) Locals + visible bindings in all enclosing blocks
 - c) Globals only
 - d) All active frames
14. The **referencing environment** under dynamic scope is
- a) Locals + all visible bindings in all active subprograms
 - b) Locals only
 - c) Globals only
 - d) Imports only
15. In C/C++, a pointer returned to a **stack variable's address** is
- a) Valid if `const`
 - b) Undefined (dangling)
 - c) Valid until next call
 - d) Valid for the program
16. In C/C++, a `static` data member of a class lives in
- a) Each object's subobject
 - b) Static storage (one per class)
 - c) Stack of each method call
 - d) Heap per object
17. In row-major languages, parameters are passed and a new AR is pushed
- a) When the ELF is linked
 - b) At each function call (on invocation)
 - c) At compile time
 - d) Only for recursive calls
18. **Early binding** vs **late binding** primarily distinguish
- a) Stack vs heap

- b) Compile-time vs run-time resolution
 - c) Static vs dynamic storage duration
 - d) Public vs private
19. In C, a function parameter `int y` typically has
- a) Static storage
 - b) Automatic (stack) storage
 - c) Heap storage
 - d) Register storage only
20. If a closure captures a local variable, its lifetime is
- a) Still limited to the call
 - b) Extended (boxed/heap cell) as long as closure exists
 - c) Moved to global section
 - d) Ignored by runtime
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Part II — Application (20) — Lifetime/Allocation & Outputs

21. Consider C++:

```
class MyClass { public: int a; static int b; }; int MyClass::b = 0; MyClass*
createObject() { MyClass obj1; MyClass* obj2 = new MyClass; static MyClass obj3; return
obj2; }
```

Which statement is correct about **memory regions**?

- a) obj1 heap, obj2 stack, obj3 data segment
- b) obj1 stack, obj2 stack, obj3 data segment
- c) obj1 stack, obj2 heap, obj3 data segment
- d) obj1 heap, obj2 heap, obj3 stack

22. Given C:

```
int x; void foo(int y){ static int z; int *t = malloc(sizeof(int)); // ... }
```

Which statement is **WRONG** about **lifetime**?

- a) `x` lives for the whole program
- b) `y` lives for the duration of `foo` call
- c) `z` lives for the duration of `foo` call
- d) `*t` lives until `free(t)` (or process ends)

23. With the same code, which statement is **WRONG** about **allocation**?

- a) `x` in static storage
- b) `y` in stack
- c) `z` in static storage
- d) `t` in heap and auto-freed at end of `foo`

24. C++:

```
void foo(int a, int *b, int &c, int *arr) { a += *b; *b *= 2; c += a; arr[0] = 42; arr = new int[2]{7,8}; } int main() { int x=5, y=3, z=10; int arr[2]={1,2}; int *par = arr; foo(x,&y,z,par); printf("%d %d %d | {%d,%d} | par-> {%d,%d}\n", x,y,z, arr[0],arr[1], par[0],par[1]); }
```

What prints?

- a) 5 6 18 | {42,2} | {7,8}
- b) 5 5 10 | {42,2} | {1,2}
- c) 5 3 18 | {1,2} | {42,2}
- d) 5 6 18 | {1,2} | {7,8}

25. C++:

```
int *p = new int(5); int *q = p; delete p; *q = 7; cout << *q;
```

Result is

- a) Prints 7
- b) Prints 5
- c) Undefined behavior (dangling write)
- d) Compile error

26. C++:

```
int *p = new int(1); p = nullptr;
```

Which is true?

- a) Safe; no leak occurs
- b) Leak (original address lost)
- c) Double free
- d) UB only if later dereferenced

27. C (assume typical calling convention):

```
int g = 1; void A(){ int g = 2; } void B(){ printf("%d", g); } int main(){ A(); B(); }
```

Output is

- a) 1 b) 2 c) UB d) compile error

28. Python (static scope):

```
x=1 def outer(): x=10 def inner(): return x+1 return inner() print(outer(), x)
```

Output

- a) 11 1 b) 11 10 c) 2 1 d) NameError

29. Python lifetime/escape:

```
def make(): bag=[] def add(v): bag.append(v) return sum(bag) return add f=make()
```

```
print(f(3), f(4), f(1))
```

a) 3 4 1 b) 3 7 8 c) 3 7 4 d) 8 7 3

30. C++:

```
int *bad(){ int x=5; return &x; } int *good(){ return new int(5); }
```

Which is correct?

- a) Both OK
- b) `bad` returns dangling; `good` OK (must delete)
- c) Both dangling
- d) Both leak always

31. C++:

```
const string& r = string("hi"); cout << r << '\n';
```

Which is correct?

- a) Always dangling
- b) Lifetime is extended for const lvalue ref to temporary → safe
- c) Moved to static storage
- d) UB at link time

32. C++:

```
int a=2,b=3; int &r=a, &s=r; s=b; cout<<a;
```

Output

- a) 2 b) 3 c) UB d) compile error

33. C++:

```
int *p=new int(5); delete p; delete p;
```

- a) Defined no-op
- b) Undefined behavior (double free)
- c) Compile error
- d) Silently leaks but defined

34. Python (import precedence last):

```
from math import sqrt sqrt = lambda z: z+1 x=3 def f(): return sqrt(x) print(f())
```

- a) 2.0 b) 4 c) NameError d) 3

35. Pseudo-Pascal — what prints under static vs dynamic scope?

```
var a: integer := 1; procedure P; var a: integer := 2; procedure Q; begin writeln(a) end;
```

```
begin Q end; begin P end.
```

- a) Static: 2 ; Dynamic: 1
- b) Static: 1 ; Dynamic: 2
- c) Both 2
- d) Both 1

36. Pseudo-Pascal:

```
var a: integer := 1; procedure P; var b: integer := 2; procedure Q; begin writeln(a+b)
end; begin Q end; begin P end.
```

Under **dynamic** scope (not real Pascal), Q prints

- a) 3 b) 2 c) 1 d) depends on call site

37. Pseudo-Pascal:

```
var x: integer := 5; procedure A; var x: integer := 7; procedure B; begin writeln(x) end;
begin B end; begin A end.
```

Under **static** scope \Rightarrow prints ... ; under **dynamic** scope \Rightarrow prints ...

- a) 7 ; 5 b) 5 ; 7 c) 7 ; 7 d) 5 ; 5

38. C++ (lifetime categories):

Pick the **incorrect** mapping.

- a) global \rightarrow static storage, lifetime=program
- b) local automatic \rightarrow stack, lifetime=call
- c) `new` object \rightarrow static storage, lifetime=program
- d) `static` local \rightarrow static storage, lifetime=program

39. C (frame content):

Which is **not** typically stored in an activation record?

- a) Return address
- b) Saved registers
- c) Heap free list
- d) Parameters

40. Python (closure escape & lifetime):

```
def counter(): n=0 def inc(): nonlocal n; n+=1; return n return inc c=counter()
print(c(), c(), c())
```

- a) 1 1 1 b) 1 2 3 c) 0 1 2 d) 2 3 4

Part III — Advanced Application (20) — Longer Case Studies & Static vs Dynamic Scope

41. Static vs Dynamic scope (sum updates):

```
var A,B: integer; procedure P(X: integer); var B: integer; procedure Q(Y: integer); begin
A := A + Y; B := B + X end; begin B := 2; Q(3); writeln(A, B) end; begin A := 1; B := 5;
P(4) end.
```

What prints inside P?

- a) Static: 4 6 ; Dynamic: 4 9
- b) Static: 4 9 ; Dynamic: 4 6
- c) Both 4 6
- d) Both 4 9

42. Output trace (C++, mixed parameters & heap):

```
void foo(int a, int* b, int& c, int*& par) { a += *b; *b += 1; c -= a; par = new int[3]
{a, *b, c}; } int main(){ int x=4, y=5, z=20; int arr[3]={0,0,0}; int* p=arr;
foo(x,&y,z,p); printf("x=%d y=%d z=%d | arr={%d,%d,%d} | p={%d,%d,%d}\n",
x,y,z,arr[0],arr[1],arr[2],p[0],p[1],p[2]); }
```

- a) x=4 y=6 z=11 | {0,0,0} | {9,6,11}
- b) x=4 y=6 z=11 | {9,6,11} | {9,6,11}
- c) x=5 y=5 z=?? (UB)
- d) x=4 y=6 z=16 | {0,0,0} | {10,6,16}

43. Dangling misuse after free (C++):

```
int *p=new int(5); int &r=*p; { int *q=p; *q=6; delete p; } cout<<r<<'\n';
```

- a) reliably prints 6
- b) undefined behavior (dangling)
- c) prints 5
- d) compile error

44. Python — deep closure/AR chain:

```
def outer(a): b=2 def mid(c): nonlocal b b += a + c def inner(d): return a + b + c + d
return inner f = mid(3) print(f(4))
```

Output

- a) 10 b) 12 c) 14 d) 16

45. Lifetime question (C++):

```
int* bad(){ int x=5; return &x; } int* good(){ return new int(5); }
```

Choose the best statement.

- a) bad ok; good ok
- b) bad dangling; good ok (must delete)
- c) both dangling
- d) both leak even if deleted

46. Static vs dynamic reading shadowed `x`:

```
var x: integer := 1; procedure A; var x: integer := 2; procedure B; begin writeln(x) end;
begin B end; begin A end.
```

- a) Static→2 ; Dynamic→1
- b) Static→1 ; Dynamic→2
- c) Both→2
- d) Both→1

47. C++ recursion + heap lifetime:

```
int sumUp(int n){ if(n<=1) return 1; int *p=new int(n); int r=*p + sumUp(n-1); delete p;
return r; } int main(){ cout<<sumUp(3); }
```

Prints

- a) 3 b) 4 c) 6 d) 7

48. Python — star args & kwargs inside nested default:

```
def K(a=1,*args,**kw): def M(b=2,*t,**m): return a + b + len(t) + len(m) return M(3, 7,
8, x=9, y=10) print(K())
```

- a) 5 b) 6 c) 7 d) 8

49. Dynamic scope thought experiment

In a dynamic-scoped language, with call chain `Main→S1→S2` and each declares `x`, a read of `x` inside `S2` binds to

- a) Main.x b) S1.x c) S2.x d) global.x

50. Python — closure list mutation persists (lifetime evidence):

```
def make(): bag=[] def add(v): bag.append(v) return bag[:] return add f=make()
print(f(2), f(5), f(7))
```

- a) [2] [5] [7]
- b) [2] [2,5] [2,5,7]
- c) [2] [5] [2,7]
- d) [2] [2] [2]

51. Union punning (C/C++):

```
union U{ unsigned i; float f; }; U u; u.i=0x3F800000; cout<<u.f;
```

Which is most accurate?

- a) Always 1.0 by the standard across languages
- b) Implementation/aliasing dependent; often 1.0 on many ABIs
- c) Always 0.0
- d) Compile-time error in C++

52. Python — global write then read via inner:


```
x=1 def f(): global x; x = x + 2 def g(): return x return g() print(f(), x)
```

a) 3 1 b) 3 3 c) 2 3 d) Error

53. Referencing environment summary

Pick the correct statement.

- a) Static RE = locals + lexically enclosing visibles; Dynamic RE = locals + all active frames' visibles
- b) Both REs = locals + globals only
- c) Static RE = activation record contents
- d) Dynamic RE ignores parameters

54. Python — nonlocal ladder shows lifetime extension:

```
def T(): a=0 def U(): nonlocal a; a+=1 def V(): nonlocal a; a+=2 def W(): nonlocal a; a+=3 return a return W() return V()+a return U()+a print(T())
```

a) 3 b) 4 c) 6 d) 7

55. C++ — heap without delete:

```
int* mk(int v){ return new int(v); } int main(){ int* q=mk(4); cout<<*q; /* no delete */ }
```

- a) Prints 4, leaks if not deleted
- b) Freed automatically at scope end
- c) Compile error
- d) Undefined

56. Name resolution order (Python): locals → nonlocals → globals → builtins/imports. For

```
from math import sqrt sqrt = lambda z: z*z x=3 def f(): return sqrt(x) print(f())
```

a) 9 b) 1.732... c) NameError d) 3

57. C++: global vs local with same name

```
int g=1; void A(){ int g=2; } void B(){ cout<<g; } int main(){ A(); B(); }
```

What concept explains the output?

- a) Dynamic scope
- b) Static scope & different bindings (local hides only inside A)
- c) Heap lifetime
- d) RAI

58. Python — higher-order closure shows lifetime:

```
def power(k): def p(x): return x**k return p sq=power(2); cu=power(3) print(sq(3), cu(2))
```

a) 6 8 b) 9 8 c) 9 6 d) 8 9

59. AR links

When inner function accesses an outer local, implementation typically uses

- a) Static link / closure environment to reach enclosing frame/cell
- b) Only dynamic link
- c) Only symbol table
- d) Only global register

60. Python — mixed nested with captured update (determine output):

```
x=10 def F(): x=1 def G(y): nonlocal x x+=y def H(): return x + y return H() return G(2)+x print(F(), x)
```

- a) 5 10 b) 3 10 c) 4 1 d) Error

Answer Key

- 1 b
- 2 a
- 3 b
- 4 b
- 5 b
- 6 b
- 7 b
- 8 b
- 9 b
- 10 b
- 11 b
- 12 b
- 13 b
- 14 a
- 15 b
- 16 b
- 17 b
- 18 b
- 19 b
- 20 b
- 21 c
- 22 c
- 23 d
- 24 a
- 25 c
- 26 b
- 27 a
- 28 a
- 29 b
- 30 b

31 b
32 b
33 b
34 b
35 a
36 a
37 a
38 c
39 c
40 b
41 a
42 a
43 b
44 c
45 b
46 a
47 d
48 d
49 b
50 b
51 b
52 b
53 a
54 d
55 a
56 a
57 b
58 b
59 a
60 a