

Lecture 10 CSCI: Option Monad

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Outline

- Option monad
- Homework description



Pipelining Option-functions

- Let's implement overflow-safe addition and safe integer division

```
let plus x y
= let s = x + y in
  match (x >= 0, y >= 0, s >= 0) with
  | (true, true, true) -> Some s
  | (false, false, false) -> Some s
  | (true, false, _) -> Some s
  | (false, true, _) -> Some s
  | (_, _, _) -> None
```

```
let divide x
= function 0 -> None | y -> Some (x / y)
```



Pipelining Option-functions

- Let's do two more:

```
let double x = plus x x
```

```
let plus1 x = plus 1 x
```

- The point of these is not their implementation, it's their types:
- `plus : int -> int -> int option`
- `divide : int -> int -> int option`
- `double : int -> int option`
- `plus1 : int -> int option`



Pipelining Option-functions

- `plus : int -> int -> int option`
- `divide : int -> int -> int option`
- `double : int -> int option`
- `plus1 : int -> int option`
- We'd like to string together computations like so:
- `3 |> plus1 |> double = Some 8`
- Of course, this doesn't type-check...
- Solution: let's replace `|>` by a different function!



Pipelining Option-functions

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- `divide : int -> int -> int option`
- `double : int -> int option`
- `plus1 : int -> int option`
- We'd like to string together computations like so:
- `3 |> plus1 >>= double = Some 8`
- `>>=` is pronounced 'bind'
- What's the type of `>>=`?
- `(>>=) : 'type of 3 |> plus1' -> 'type of double' -> 'type of Some 8'`



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- `(>>=) : int option -> (int -> int option) -> int option`
- `let (|>) x f = f x`
- `let (>>=) x f = ...`



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- `(>>=) : int option -> (int -> int option) -> int option`
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`let (>>=) x f = match x with`
`...`



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- `let (|>) x f = f x`
- `let (>>=) x f = match x with`
 - `| None -> ...`
 - `| Some y ->`



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`let (>>=) x f = match x with`
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Pipelining Option-functions

- We'd like to string together computations like so:
- `3 |> plus1 >=> double = Some 8`

```
utop # 3 |> plus1 >=> double;;  
- : int option = Some 8
```

- `let (|>) x f = f x`
`let (>=>) x f = match x with`
 `| None -> None`
 `| Some y -> f y`



Pipelining Option-functions

- The mixing of `|>` and `>>=` is perhaps a bit ugly...
- instead we can write:

```
utop # Some 3 >>= plus1 >>= double;;  
- : int option = Some 8
```

- `let (|>) x f = f x`
`let (>>=) x f = match x with`
 `| None -> None`
 `| Some y -> f y`



Generalization aside

- The `>>=` operation in these slides is defined for the `int option` datatype.

- Actually, if we read `>>=` better, it's for 'a option:

```
utop # let (>>=) x f = match x with
```

```
    | None -> None
```

```
    | Some y -> f y;;
```

```
val ( >>= ) : 'a option -> ('a -> 'b option) -> 'b option
```

- In real life, people define `>>=` for a huge set of different data-types. We will only look at the 'a option one here.
- However, we'll try to do everything option related with just '`>>=`', so no more "match" statements after this.
- Reason: this way, this lecture applies to *any* `>>=` definition.



Pipelining Option-functions

- We can make these as long as we like...

```
utop # Some 3 >>= plus1 >>= double;;  
- : int option = Some 8
```

- Where `|>` was free
 `>>=` actually does a (very simple) computation for us



... functions with two arguments

- Let's calculate: $4 + 6 = 10$
- $6 \mid> (4 \mid> (+))$
- Now with our function...
- $6 \mid> (4 \mid> \text{plus}) = \text{Some } 10$
- ... but what if 6 and 4 are the result of another computation? (i.e. 'Some 6 / Some 4')
- $\text{Some } 6 >>= (4 \mid> \text{plus}) = \text{Some } 10 \dots$
- almost there... what can we write for '4 $\mid>$ plus'?



... functions with two arguments

```
utop # Some 4 |> plus;;
```

```
Error: This expression has type int -> int -> int option  
      but an expression was expected of type int option -> 'a  
      Type int is not compatible with type int option
```



... functions with two arguments

```
utop # Some 4 |> plus;;
```

```
Error: This expression has type int -> int -> int option  
      but an expression was expected of type int option -> 'a  
      Type int is not compatible with type int option
```

```
utop # Some 4 >>= plus;;
```

```
Error: This expression has type int -> int -> int option  
      but an expression was expected of type int -> 'a option  
      Type int -> int option is not compatible with type  
      'a option
```



... functions with 2 arguments

- Let's apply the currying idea!
- `Some 4 >>= (fun nr_four -> Some 6 >>= plus nr_four)`



... functions with 2 arguments

- Let's apply the currying idea!
- `Some 4 >>= (fun nr_four -> Some 6 >>= plus nr_four)`
- this part: `(fun nr_four -> Some 6 >>= plus nr_four)`
- it's a function that returns an 'int option'
- internally, it applies 'Some 6' to 'plus nr_four'
- here, 'plus nr_four' is also a function that returns an 'int option'
- this is writing `4 + 6` (in that order)



... functions with 2 arguments

- Let's apply the currying idea!
- `Some 4 >>= (fun nr_four -> Some 6 >>= plus nr_four)`
- `(fun nr_four -> Some 6 >>= plus nr_four)`
- should read like: `6 |> (+)`
- or: `(6 + ...)`
- I'll admit: it doesn't read nicely



Towards nicer notations

- If `(Some 6 >>= plus)` doesn't type-check,
... and should really read:
`(fun nr_four -> Some 6 >>= plus nr_four)`
- can't we define yet another `>>=`-like function?



Towards nicer notations

- If `(Some 6 >>= plus)` doesn't type-check,
... and should really read:
`(fun nr_four -> Some 6 >>= plus nr_four)`
- can't we define yet another `>>=`-like function?
- let `(>>==) x f = ...`



Towards nicer notations

- If `(Some 6 >>= plus)` doesn't type-check,
... and should really read:
`(fun nr_four -> Some 6 >>= plus nr_four)`
- ... should we define yet another `>>=`-like function?
- `let (>>==) x f = (fun nr_four -> x >>= f nr_four)`
- `let (>>==) x f nr_four = (x >>= f nr_four)`
- `let (>>==) x f a = (x >>= f a)`
- great.. now define one for each number of arguments?



Towards nicer notations

- It's not too common to introduce `>>==`
- People have invented tons of 'nice' notations for `>>=`
- It helps to realize that there is a pattern:
- Some `3 >>= plus1 >>= double >>= double`
- is the same as:
- Some `3 >>= (fun nr3
-> plus1 nr3 >>= (fun nr4
-> double nr4 >>= (fun nr8
-> double nr8)))`



Towards nicer notations

- It's not too common to introduce `>>==`
- People have invented tons of 'nice' notations for `>>=`
- What its use boils down to is a pattern:
- Some `3 >>= plus1 >>= double >>= double`

- is the same as:

- Some `3 >>= (fun nr3
-> plus1 nr3 >>= (fun nr4
-> double nr4 >>= (fun nr8
-> double nr8)))`

This doesn't look nice at first,
but it resembles "let" notation!!!
`let nr3 = Some 3 in
let nr4 = plus1 nr3 in
let nr8 = double nr4 in
double nr8`



Towards nicer notations

- This pattern works just fine for multiple argument functions!
- Some 3 >>= (fun nr3
-> plus1 nr3 >>= (fun nr4
-> Some 6 >>= (fun nr6
-> plus nr4 nr6)))
- (think: let nr3 be the result of Some3, let nr4 be ... etc)



Towards nicer notations

- This pattern also works when arguments are used twice!
- Some 3 >>= (fun nr3
-> plus1 nr3 >>= (fun nr4
-> plus nr3 nr4 >>= (fun nr7
-> plus nr4 nr7)))



Crucial observation

- If you look up the definition of $\gg=$ and try to fit it all into what the program is doing...
- ... the result is a very big complex program
- Instead, think of $e1 \gg= (\text{fun } x \rightarrow e2)$ as a special case of 'let $x = e1$ in $e2$ '
however, the ('a option) values are bound as 'a:
This means: $e1$ is 'a option, but x is 'a
(which makes this 'let' a bit special)



Towards nicer notations

- ocaml 4.08 allows us to define our own “let*”
 - some CSE machines are at 4.08 but I’m not sure if all of them are... Here’s what we could write:
-
- `Some 3 >>= (fun nr3
-> plus1 nr3 >>= (fun nr4
-> plus nr3 nr4 >>= (fun nr7
-> plus nr4 nr7)))`
 - `let (let*) = (>>=)`
 - `let* nr3 = Some 3 in
let* nr4 = plus1 nr3 in
let* nr7 = plus nr3 nr4 in
plus nr4 nr7`



Back to that generalization...

- Note that $>>=$ can be defined for other types..
- The program below still reminds us of the option monad because we used 'Some'
- People like to abstract away from that too, defining:
let return $x = \text{Some } x$
- $\text{Some } 3 >>= (\text{fun nr3}$
 $\rightarrow \text{plus1 } \underline{\text{nr3}} >>= (\text{fun nr4}$
 $\rightarrow \text{plus } \underline{\text{nr3}} \underline{\text{nr4}} >>= (\text{fun nr7}$
 $\rightarrow \text{plus } \underline{\text{nr4}} \text{ nr7 })))$
- $\text{return } 3 >>= (\text{fun nr3}$
 $\rightarrow \text{plus1 } \underline{\text{nr3}} >>= (\text{fun nr4}$
 $\rightarrow \text{plus } \underline{\text{nr3}} \underline{\text{nr4}} >>= (\text{fun nr7}$
 $\rightarrow \text{plus } \underline{\text{nr4}} \text{ nr7 })))$



What are monads?

- Monads are a variation on pipelining
- The main workhorse of monads is called 'bind', written $>>=$, the other part is 'return'.
- The option monad (aka maybe monad) defines 'bind' and 'return' like we did today.
- If you want to say you defined a monad, your functions bind and return should be pipeline-like enough
- Commonly, that means they should intend to satisfy:
- $(\text{return } r >>= f) = f \ r$
- $(v >>= \text{return}) = v$
- $((v >>= (\text{fun } x \rightarrow g \ x)) >>= h) = (v >>= (\text{fun } x \rightarrow (g \ x >>= h)))$



On Monads ...

- If you wish to define bind (`>>=`), go ahead!
- bind is not defined in the standard library, but some commonly used libraries provide them.
- There's one for option, and many more...
- Your TAs have not been taught 'monads'
 - I won't be creating homework that *requires* you to use them: you can do all midterms / assignments without `>>=`
 - I encourage you to try it out because it might help you (and maybe teach your TAs a thing or two)
 - ... and because I want my next TAs to know them :-)

