

Definition for reference:

```
let rec plus a b =  
  match b with  
  | Zero -> a  
  | S c -> plus (S a) c
```

Here is a useful lemma that is used throughout:

$\text{plus } (S \ a) \ b = S \ (\text{plus } a \ b)$

proof by induction on b. Case b = Zero:

```
plus (S a) b  
= {case}  
plus (S a) Zero  
= {def}  
match Zero with  
| Zero -> S a  
| S c -> plus (S (S a)) c  
= {match}  
S a  
plus (S a) Zero  
S a  
= {match}  
S (match Zero with  
| Zero -> a  
| S c -> plus (S a) c)  
= {def}  
S (plus a Zero)  
= {case}  
S (plus a b)
```

Case b = S c, IH: $\text{plus } (S \ a) \ c = S \ (\text{plus } a \ c)$

```
plus (S a) b  
= {case}  
plus (S a) (S c)  
= {def}  
match S c with  
| Zero -> S a  
| S c -> plus (S (S a)) c  
= {match}  
plus (S (S a)) c  
= {IH}  
S (plus (S a) c)  
= {match}  
= {def}  
S (plus a b)
```

Thus we have proved that ' $\text{plus } (S \ a) \ b = S \ (\text{plus } a \ b)$ '.

I'll call this lemma 'S and plus commute in their first argument'.

P1) Prove $\text{plus Zero } b = b$

This proof would be trivial if we used P2, but we cannot prove P2 without P1.

Here's a proof by induction on b : Case $b = \text{Zero}$:

```
plus Zero b
= {case}
  plus Zero Zero
= {def}
  match Zero with
    | Zero -> Zero
    | S c -> plus (S Zero) c
= {match}
  Zero
= {case}
  b
```

Case $b = S \ d$, with IH: $\text{plus Zero } d = d$

```
plus Zero b
= {case}
  plus Zero (S d)
= {def}
  match S d with
    | Zero -> Zero
    | S c -> plus (S Zero) c
= {match}
  plus (S Zero) d
= {lemma 'S and plus commute in their first argument'}
  S (plus Zero d)
= {IH}
  S d
= {case}
  b
```

Thus we have proved that ' $\text{plus Zero } b = b$ '.
I'll call this lemma 'P1'.

P2) Prove $\text{plus } a \ b = \text{plus } b \ a$

We'll prove this by induction on b

Base case: $b = \text{Zero}$

```
plus a b
= {case}
  plus a Zero
```

```

= {def}
  match Zero with
  | Zero -> a
  | S c -> plus (S a) c
= {match}
  a
= {lemma 'P1'}
  plus Zero a
= {case}
  plus b a

```

Inductive case: $b = S\ c$, with IH: $plus\ a\ c = plus\ c\ a$

```

  plus a b
= {case}
  plus a (S c)
= {def}
  match S c with
  | Zero -> a
  | S c -> plus (S a) c
= {match}
  plus (S a) c
= {lemma 'S and plus commute in their first argument'}
  S (plus a c)
= {IH}
  S (plus c a)
= {lemma 'S and plus commute in their first argument'}
  plus (S c) a
= {case}
  plus b a

```

Thus we have proved that ' $plus\ a\ b = plus\ b\ a$ '.
I'll call this lemma 'P2'.

P3) Prove $plus\ a\ (plus\ b\ c) = plus\ (plus\ a\ b)\ c$

With P2 under our belt, we can move arguments to plus around all we like, so it really doesn't matter on what argument we do induction. However, I can use lemma 'S and plus commute in their first argument', since that avoids me having to use definitions (which is two steps rather than one).

For that reason I give a proof by induction on a here.

Here's a proof by induction on a: case $a = Zero$

```

  plus a (plus b c)
= {case}
  plus Zero (plus b c)
= {lemma 'P2'}
  plus (plus b c) Zero
= {lemma 'P1'}
  plus b c

```

```

= {lemma 'P1'}
  plus (plus b Zero) c
= {lemma 'P2'}
  plus (plus Zero b) c
= {case}
  plus (plus a b) c

case a = S d, IH: plus d (plus b c) = plus (plus d b) c
  plus a (plus b c)
= {case}
  plus (S d) (plus b c)
= {lemma 'S and plus commute in their first argument'}
  S (plus d (plus b c))
= {IH}
  S (plus (plus d b) c)
= {lemma 'S and plus commute in their first argument'}
  plus (S (plus d b)) c
= {lemma 'S and plus commute in their first argument'}
  plus (plus (S d) b) c
= {case}
  plus (plus a b) c

```

Thus we have proved that 'plus a (plus b c) = plus (plus a b) c' by induction.