

You may use up to 10 pages of single-side notes and the textbook. You may not use a computer, phone, calculator or any other electronic equipment except to compile and run exam questions. You may not communicate with any other student via in-person or via electronic or any other means.

- 1. Numeric conversions:
  - a. binary conversions ones complement, twos complement
  - b. Hex addition / subtraction
  - c. Big vs Little Endian
- 2. ASCII / Unicode conversion
- 3. Intel Assembly
  - a. Example 1:

### 4. Intel Assembly

What is the equivalent C code and what will be stored in var1



```
MOV EAX, 1
  MOV EBX, 100
  MOV ECX, 101
   SUB EBX, ECX
   JG next1
   INC EAX
next1:
  MOV var1, EAX ; what is stored in var1
   int eax, ebx, ecx, var1;
   eax = 1;
   ebx = 100;
   ecx = 101;
   ebx = ebx - ecx;
   if (ebx <= 0)
     eax += 1;
   var1 = eax; // var1 = 2
```

- 5. Cache: how lines are mapped from addresses, performance
- 6. Address translation page 391, figure 6.23
- 7. Disk
  - a. Access = seek time + rotational delay + transfer time
  - RAID = RAID 0 optimal performance, no redundancy
     RAID 1 mirror great performance, total redundancy, high cost
- 8. CPU scheduling / Context Switching / Interrupts FIFO, SJF, Priority, Round-Robin
- 9. I/O types programmed, memory mapped, DMA, interrupts
- 10. RISC vs CISC differences page 534, table 9.1
- 11. Intel 64, ARM 64, RISC-V basic architecture
  - a. Stack management EBP and ESP registers
  - b. CISC vs RISC
  - c. ISA of ARM, RISC-V (number of registers, instruction width, caching)
  - d. ARM-64: Data and instruction cache, can run 32- and 64-bit programs, MMU manages memory mapping
- 12. Networking
  - a. UDP vs TCP
  - b. TCP sequence numbers and Windowing



- c. IP addresses
- 13. Java VM
  - a. Bytecode
  - b. Garbage collection
- 14. Quantum Computing
  - a. Quantum mechanics and Qbit
  - b. Decoherence
- 15. Compilers

-99 (10) to hex

- a. Optimization
- b. Static vs shared libraries

### Examples:

```
99 / 2 = 49 r 1
49 / 2 = 24 r 1
24 / 2 = 12 r 0
12 / 2 = 6 r 0
6 / 2 = 3 r 0
3 / 2 = 1 r 1
1 / 2 = 0 r 1
0 \ 1 \ 1 \ 0 \ 0 \ 0 \ 1 \ 1 = 63 \ (16)
1 0 0 1 1 1 0 0 = one complement
              +1
1 \ 0 \ 0 \ 1 \ 1 \ 1 \ 0 \ 1 = 9D \ (16)
 1
9A4C
+16B2
____
BOFE
3210
11 * 16^3 + 0 * 16^2 + 15 * 16^1 + 14 * 16^0 = 45310 (10)
```



9 9A4C -16B2 ==== 839A

Big Endian

00 00 B0 FE

Little Endian

FE B0 00 00

### "789 a Σ

ASCII = 37 38 39 20 61 20 ??

Unicode = 0037 0038 0039 0020 0061 0020 03A3 Unicode (little) = 3700 3800 3900 2000 6100 2000 A303

int  $xx[3] = \{1, 2, 3\};$ 

1 2 3 00 00 00 01 00 00 02 00 00 00 03 (big endian) 01 00 00 00 00 00 00 03 00 00 (little endian)

#### Cache:

Each addresses tag + directMap + offset Offset bit = size of each cache lines directMap size = size of cache

cache line = 16 bytes cache size = 256 lines

address =  $12345678_{16}$  = 12345678 offset =  $8_{16}$  cache entry =  $67_{16}$ 

 $tag = 12345_{16}$ 



```
Virtual Address:
Page number + offset
offset size = size of each page frame
page size = 4096
virtual address space = 4096 pages
physical address space = 1024 pages
Page table
 0 - 01
 1 - 0D
 2 - 02
 3 - N/A
 4 - 04
 Virtual Address = XXX YYY<sub>16</sub>
                             001 \ 010_{16} \ -> \ 0D \ 010_{16}
                             002 0FF<sub>16</sub> -> 02 0FF<sub>16</sub>
                             003 100<sub>16</sub> \rightarrow page fault
Memory Layout:
address 0:
 code
 global data and constants
 heap (grows to higher addresses)
 stack (grows to lower addresses)
address N (highest)
```