Exercise 4.1.3 Prove the following equality:

$$\operatorname{Tr} \{A\} = \langle \Gamma |_{RS} I_R \otimes A_S | \Gamma \rangle_{RS}, \tag{4.22}$$
 where A is a square operator acting on a Hilbert space \mathcal{H}_S , I_R is the identity operator acting

where A is a square operator acting on a Hilbert space \mathcal{H}_S , I_R is the identity operator acting on a Hilbert space \mathcal{H}_R isomorphic to \mathcal{H}_S and $|\Gamma\rangle_{RS}$ is the unnormalized maximally entangled vector from (3.233). This gives an alternate formula for the trace of a square operator A.

Exercise 4.1.16 (Union Bound) Prove a union bound for commuting projectors Π_1 and Π_2 where $0 < \Pi_1, \Pi_2 < I$ and for an arbitrary density operator ρ (not necessarily diagonal in the same basis as Π_1 and Π_2):

$$\Pi_2$$
 where $0 \leq \Pi_1, \Pi_2 \leq I$ and for an arbitrary density operator ρ (not necessarily diagonal in the same basis as Π_1 and Π_2):

 $\operatorname{Tr} \{ (I - \Pi_1 \Pi_2) \rho \} \leq \operatorname{Tr} \{ (I - \Pi_1) \rho \} + \operatorname{Tr} \{ (I - \Pi_2) \rho \}.$

(4.77)

Exercise 4.2.2 Suppose we have an ensemble $\{p_X(x), \rho_x\}$ of density operators and a POVM with elements $\{\Lambda_x\}$ that should identify the states ρ_x with high probability, i.e., we would like $\text{Tr}\{\Lambda_x\rho_x\}$ to be as high as possible. The expected success probability of the POVM is then

$$\sum_{x} p_X(x) \operatorname{Tr} \left\{ \Lambda_x \rho_x \right\}. \tag{4.99}$$

Suppose that there exists some operator τ such that

$$\tau \ge p_X(x)\rho_x\,,\tag{4.100}$$

where the condition $\tau \geq p_X(x)\rho_x$ is the same as $\tau - p_X(x)\rho_x \geq 0$ (i.e., that the operator $\tau - p_X(x)\rho_x$ is a positive semi-definite operator). Show that $\operatorname{Tr} \{\tau\}$ is an upper bound on the expected success probability of the POVM. After doing so, consider the case of encoding n bits into a d-dimensional subspace. By choosing states uniformly at random (in the case of the ensemble $\{2^{-n}, \rho_i\}_{i \in \{0,1\}^n}$), show that the expected success probability is bounded above by $d 2^{-n}$. Thus, it is not possible to store more than n classical bits in n qubits and have a perfect success probability of retrieval.

Exercise 4.3.1 Show that the purity $P(\rho_A)$ is equal to the following expression: $P(\rho_A) = \operatorname{Tr} \left\{ (\rho_A \otimes \rho_{A'}) F_{AA'} \right\}.$

where system A' has a Hilbert space structure isomorphic to that of system A and $F_{AA'}$ is the swap operator that has the following action on kets in A and A':

(4.106)

 $\forall x, y \qquad F_{AA'}|x\rangle_A|y\rangle_{A'} = |y\rangle_A|x\rangle_{A'}.$ (4.107)

(One can in fact show more generally that $\operatorname{Tr}\{f(\rho_A)\}=\operatorname{Tr}\{(f(\rho_A)\otimes I_{A'})F_{AA'}\}$ for any function f on the operators in system A.)

Exercise 4.3.6 Show that a parity measurement (defined in the previous exercise) of the state $|\Phi^{+}\rangle_{AB}$ returns an even parity result with probability one, and a parity measurement of the state $\pi_A \otimes \pi_B$ returns even or odd parity with equal probability. Thus, despite the fact that these states have the same local description, their global behavior is very different. Show

e states have the same local description, their global behavior is very different. Show same is true for the phase parity measurement, given by
$$\Pi_{\text{even}}^{X} \equiv \frac{1}{2} \left(I_A \otimes I_B + X_A \otimes X_B \right), \tag{4.134}$$

(4.135)

 $\Pi_{\text{odd}}^{X} \equiv \frac{1}{2} \left(I_A \otimes I_B - X_A \otimes X_B \right).$

that the same is true for the phase parity measurement, given by

Exercise 4.3.18 Show that performing a measurement with measurement operators $\{\Lambda_A^j\}$ on system A is the same as performing a measurement of the ensemble in (4.173). That is. show that $\operatorname{Tr}\{\rho_A\Lambda_A^j\} = \operatorname{Tr}\{\rho_{XA}(I_X \otimes \Lambda_A^j)\}$, where ρ_A is defined in (4.175).

fact that any positive semi-definite operator can be diagonalized, the fact that $id_R \otimes \mathcal{N}$ is linear, and use something similar to (4.202) –(4.205)).

Exercise 4.6.2 Prove that a quantum channel $\mathcal{N}_{A\to B}$ is entanglement-breaking if $(\mathrm{id}_R \otimes \mathcal{N}_{A\to B})$ (Φ_{RA}) is a separable state, where Φ_{RA} is a maximally entangled state. (Hint: You can use a trick similar to that which you used to solve Exercise 4.4.1. Alternatively, you can inspect the

proof of Theorem 4.6.1 below.)

Exercise 4.6.3 Show that both a classical-quantum channel and a quantum-classical channel are entanglement-breaking—i.e., if we input the A system of a bipartite state ρ_{RA} to

either of these channels, then the resulting state on systems RB is separable.