Introduction to Vinyl

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Haskell records are nominally typed

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data R = R \{ x :: X \}
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data R = R \{ x :: X \}
data R' = R' \{ x :: X \} -- ^Frror
```

Structural Typing

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Sharing field names and accessors

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- Sharing field names and accessors
- Record types may be characterized structurally

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    (:&) :: !t → !(Rec rs) → Rec ((s ::: t) ': rs)
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class s \in (rs :: [*])
class ss \subseteq (rs :: [*]) where
  cast :: Rec rs → Rec ss
```

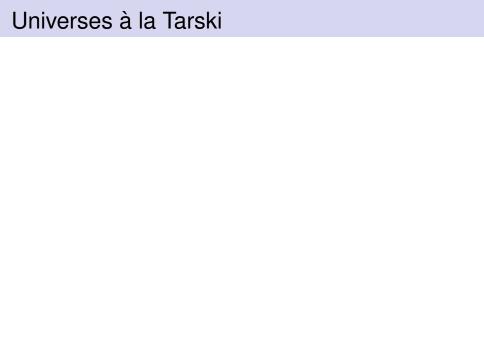
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lens : s ::: t \in rs \Rightarrow s ::: t \rightarrow Lens' (Rec rs) t
```

```
f :: Rec ("foo" ::: A ': rs)

→ Rec ("bar" ::: B ': "foo" ::: A ': rs)
```

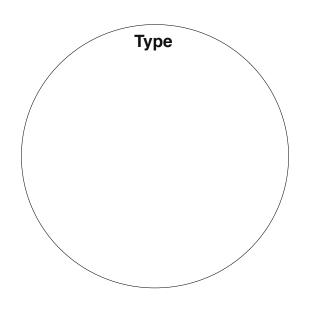


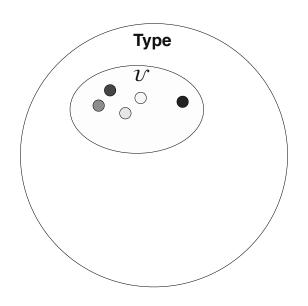
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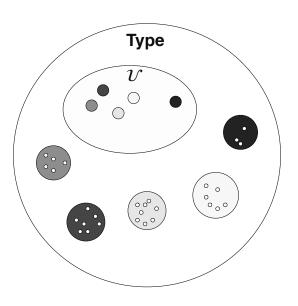
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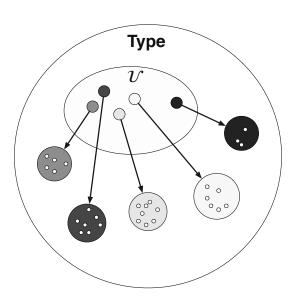
- ▶ A type \mathcal{U} of **codes** for types.
- ▶ Function $El_{\mathcal{U}}: \mathcal{U} \to \mathsf{Type}$.

$$\frac{\Gamma \vdash s : \mathcal{U}}{\Gamma \vdash El_{\mathcal{U}}(s) : \mathsf{Type}}$$









Records in Haskell

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data $\mbox{Rec} :: (\mathcal{U} \to *) \to [\ \mathcal{U}\] \to * \mbox{ where}$

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RNil :: Rec el_{\mathcal{U}} '[]
(:&) :: !(el_{\mathcal{U}} r) \to !(Rec el_{\mathcal{U}} rs) \to Rec el_{\mathcal{U}} (r ': rs)
```

 $\textbf{data} \; \mathsf{TyFun} :: * \to * \to *$

```
data TyFun :: * \rightarrow * \rightarrow * type family (f :: TyFun k I \rightarrow *) $ (x :: k) :: I
```

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data TyFun :: * \rightarrow * \rightarrow * type family (f :: TyFun k I \rightarrow *) $ (x :: k) :: I data Rec :: (TyFun \mathcal{U} * \rightarrow *) \rightarrow [ \mathcal{U} ] \rightarrow * where RNiI :: Rec el_{\mathcal{U}} '[] (:&) :: !(el_{\mathcal{U}} $ r) \rightarrow !(Rec el_{\mathcal{U}} rs) \rightarrow Rec el_{\mathcal{U}} (r ': rs)
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data Id :: (TyFun k k) \rightarrow * where type instance Id \$ x = x

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type HList rs = Rec ld rs

```
ex :: HList [\mathbb{Z}, Bool, String] ex = 34 :& True :& "vinyl" :& RNil
```

```
data Rec :: (TyFun \mathcal{U} * \to *) \to (* \to *) \to [\mathcal{U}] \to * where RNil :: Rec el<sub>\mathcal{U}</sub> f '[] (:&) :: !(f (el<sub>\mathcal{U}</sub> $ r)) \to !(Rec el<sub>\mathcal{U}</sub> f rs) \to Rec el<sub>\mathcal{U}</sub> f (r ': rs)
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data Rec :: (TyFun \mathcal{U} * \to *) \to (* \to *) \to [\mathcal{U}] \to * where RNil :: Rec el_{\mathcal{U}} f '[] (:&) :: !(f (el_{\mathcal{U}} \$ r)) \to !(Rec \ el_{\mathcal{U}} \ f \ rs) \to Rec \ el_{\mathcal{U}} \ f \ (r \ ': rs)
```

(=:) : Applicative $f \Rightarrow sing r \rightarrow el_{\mathcal{U}} \ r \rightarrow Rec \ el_{\mathcal{U}} \ f'[r]$

k =: x = pure x : & RNil

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data \operatorname{Rec} :: (\operatorname{TyFun} \mathcal{U} * \to *) \to (* \to *) \to [\mathcal{U}] \to * where \operatorname{RNil} :: \operatorname{Rec} \operatorname{el}_{\mathcal{U}} \operatorname{f}'[] (:\&) :: !(\operatorname{f} (\operatorname{el}_{\mathcal{U}} \$ \operatorname{r})) \to !(\operatorname{Rec} \operatorname{el}_{\mathcal{U}} \operatorname{f} \operatorname{rs}) \to \operatorname{Rec} \operatorname{el}_{\mathcal{U}} \operatorname{f} (\operatorname{r}' : \operatorname{rs}) (=:) : \operatorname{Applicative} \operatorname{f} \Rightarrow \operatorname{sing} \operatorname{r} \to \operatorname{el}_{\mathcal{U}} \$ \operatorname{r} \to \operatorname{Rec} \operatorname{el}_{\mathcal{U}} \operatorname{f}'[\operatorname{r}] \operatorname{k} =: \operatorname{x} = \operatorname{pure} \operatorname{x} : \& \operatorname{RNil} (\Leftarrow) : \operatorname{sing} \operatorname{r} \to \operatorname{f} (\operatorname{el}_{\mathcal{U}} \$ \operatorname{r}) \to \operatorname{Rec} \operatorname{el}_{\mathcal{U}} \operatorname{f}'[\operatorname{r}]
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data Rec :: (TyFun \mathcal{U} * \to *) \to (* \to *) \to [\mathcal{U}] \to * where RNil :: Rec el_{\mathcal{U}} f '[] (:&) :: !(f (el_{\mathcal{U}} \$ r)) \to !(Rec el_{\mathcal{U}} f rs) \to Rec el_{\mathcal{U}} f (r ': rs) (=:) : Applicative f \Rightarrow sing r \to el_{\mathcal{U}} $ r \to Rec el_{\mathcal{U}} f '[r] k =: x = pure x :& RNil (\Leftarrow): sing r \to f (el_{\mathcal{U}} \$ r) \to Rec el_{\mathcal{U}} f '[r] k \Leftarrow x = x :& RNil
```

newtype Lift o f g x = Lift { runLift :: f x 'o' g x }

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type Validator = Lift (\rightarrow) Identity (**Either** Error)

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type Validator = Lift (\rightarrow) Identity (Either Error) type JSONRenderer = Lift (\rightarrow) Identity (Const Aeson.Pair)
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 $\textbf{type} \ \, \text{Validator} = \text{Lift} \ \, (\rightarrow) \ \, \text{Identity} \ \, (\textbf{Either} \ \, \text{Error}) \\ \textbf{type} \ \, \text{JSONRenderer} = \text{Lift} \ \, (\rightarrow) \ \, \text{Identity} \ \, (\text{Const Aeson.Pair})$

 $(\textcircled{*}) :: \mathsf{Rec}_{\mathcal{U}} \; (\mathsf{Lift} \; (\to) \; \mathsf{f} \; \mathsf{g}) \; \mathsf{rs} \to \mathsf{Rec}_{\mathcal{U}} \; \mathsf{f} \; \mathsf{rs} \to \mathsf{Rec}_{\mathcal{U}} \; \mathsf{g} \; \mathsf{rs}$

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 $\begin{tabular}{ll} \textbf{type} \ \ Validator = Lift \ (\to) \ \ Identity \ (\textbf{Either} \ Error) \\ \begin{tabular}{ll} \textbf{type} \ \ JSONRenderer = Lift \ (\to) \ \ Identity \ \ (Const \ Aeson.Pair) \\ \end{tabular}$

 $(\textcircled{\$}) :: \mathsf{Rec}_\mathcal{U} \; (\mathsf{Lift} \; (\to) \; \mathsf{f} \; \mathsf{g}) \; \mathsf{rs} \to \mathsf{Rec}_\mathcal{U} \; \mathsf{f} \; \mathsf{rs} \to \mathsf{Rec}_\mathcal{U} \; \mathsf{g} \; \mathsf{rs}$

rdist :: Applicative $f \Rightarrow Rec_{\mathcal{U}} f rs \rightarrow f (Rec_{\mathcal{U}} Identity rs)$

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(\textcircled{k}) :: \mathsf{Rec}_{\mathcal{U}} (\mathsf{Lift} (\to) \mathsf{f} \mathsf{g}) \mathsf{rs} \to \mathsf{Rec}_{\mathcal{U}} \mathsf{f} \mathsf{rs} \to \mathsf{Rec}_{\mathcal{U}} \mathsf{g} \mathsf{rs}
rdist :: Applicative f \Rightarrow Rec_{\mathcal{U}} f rs \rightarrow f (Rec_{\mathcal{U}} Identity rs)
rtraverse ::
    Applicative h
        \Rightarrow (f \rightsquigarrow h \circ g)
        \rightarrow Rec<sub>14</sub> f rs
        \rightarrow h (Rec<sub>11</sub> g rs)
```

