## **Lab: RDBMS and Performance**

MariaDB is a drop-in replacement for MySQL, the most used open source database for online applications. MariaDB falls into the category of the NewSQL products,

i.e. a product that provides unique NoSQL features together with the typical features available in relational databases.

Therefore, aspects like transaction management, durability and consistency are available together with schema or schema-less modeling,

full text storage and analysis and integration with other NoSQL technologies.

MariaDB can be part of the database infrastructure for Big Data. It is not meant to be a replacement for Hadoop, but it can be a technology used in conjunction with it. Hadoop is used in batch, ad-hoc analysis. In projects that require the processing of Terabytes or Petabytes of data,

Hadoop is definitely a good fit. The results can be queried and reported via a standard MySQL/MariaDB interface, which is compatible with virtually

all the BI tools and development frameworks available today.

Let's take a detour to check out the differences.

### **RDBMS**

So now lets take a detour and look at a RDBMS. Since it's loaded we'll use MySQL (may be MariaDB also):

```
[centos@ip-10-0-0-54 ~]$ mysql -u root -p
Enter password:
Welcome to the MariaDB monitor. Commands end with; or \g
.
Your MariaDB connection id is 207
Server version: 5.5.56-MariaDB MariaDB Server

Copyright (c) 2000, 2017, Oracle, MariaDB Corporation Ab a nd others.

Type 'help;' or '\h' for help. Type '\c' to clear the curr ent input statement.
MariaDB [(none)]>
```

Note: password should be just [return]

Now look at the databases:

And create some database, and use it:

```
MariaDB [(none)]> create database hbase_test;
Query OK, 1 row affected (0.00 sec)

MariaDB [(none)]> use hbase_test;
Database changed
MariaDB [hbase_test]>
```

Let's create a table:

```
MariaDB [hbase_test]> create table users(id int, name char
(20), PRIMARY KEY(id));
Query OK, 0 rows affected (0.00 sec)
```

And another table - this one will be related to the users table by the FOREIGN KEY user\_id:

```
MariaDB [hbase_test]> create table orders(id int, user_id
int, order_info text, PRIMARY KEY(id), FOREIGN KEY (user_i
```

```
d) REFERENCES users(id));
Query OK, 0 rows affected (0.00 sec)
```

Now insert a record:

```
MariaDB [hbase_test]> insert into users (id, name) VALUES
(1, 'bill');
Query OK, 1 row affected (0.00 sec)
```

Select the table:

```
MariaDB [hbase_test]> select * from users;
+---+---+
| id | name |
+---+---+
| 1 | bill |
+---+---+
1 row in set (0.00 sec)
```

And we see the where clause is supported and so forth:

```
MariaDB [hbase_test]> select * from users where id < 2;
+---+
| id | name |
+---+
| 1 | bill |</pre>
```

```
+---+----+
1 row in set (0.00 sec)
```

So now, create a row in the related table:

```
MariaDB [hbase_test]> insert into orders (id, user_id, ord
er_info) VALUES (1,1,"something");
Query OK, 1 row affected (0.00 sec)
```

#### And check it:

```
MariaDB [hbase_test]> select * from orders;
+---+---+
| id | user_id | order_info |
+---+----+
| 1 | 1 | something |
+---+----+
1 row in set (0.00 sec)
```

```
MariaDB [hbase_test]> insert into orders (id, user_id, ord
er_info) VALUES (2,1,"something");
Query OK, 1 row affected (0.00 sec)
```

So you see the way 2 becomes another order number, no problem:

```
MariaDB [hbase_test]> select * from orders;
```

```
+---+
| id | user_id | order_info |
+---+
| 1 | 1 | something |
| 2 | 1 | something |
+---+
2 rows in set (0.00 sec)
```

But what if we do something like this:

```
MariaDB [hbase_test]> insert into orders (id, user_id, ord
er_info) VALUES (1,2,"something");
```

What has happened?

## **RDBMS and Transactions**

This time we will do the same with an active transaction:

```
MariaDB [hbase_test]> start transaction;
Query OK, 0 rows affected (0.00 sec)
```

Now, delete an order:

```
MariaDB [hbase_test]> delete from orders where id=2;
Query OK, 1 row affected (0.01 sec)
```

And the order is now gone:

```
MariaDB [hbase_test]> select * from orders;
+---+----+
| id | user_id | order_info |
+---+----+
| 1 | 1 | something |
+---+----+
1 row in set (0.00 sec)
```

But we made a mistake, so roll back the trransaction:

```
MariaDB [hbase_test]> rollback;
Query OK, 0 rows affected (0.00 sec)
```

# **NOSql Performance Indicators**

So now let's do some performance testing. We have a .sql file called users.sql in the directory above. Import into MySQL:

```
[centos@ip-10-0-0-54 data]$ mysql -u root -p < users.sql</pre>
```

Now go into MySQL and see the table:

```
[centos@ip-10-0-0-54 data]$ mysql -u root -p
Enter password:
MariaDB [(none)]> show databases;
 Database
 information schema |
| hbase test
 hive
| mysql
 performance schema |
 test
 user data
+----+
7 rows in set (0.00 \text{ sec})
MariaDB [(none)]> use user_data;
```

And count the table:

```
MariaDB [user_data]> select count(*) from users;
+-----+
| count(*) |
+-----+
| 1000000 |
+-----+
1 row in set (0.14 sec)
```

Rather large table, isn't it?

Look at the data:

```
MariaDB [user_data]> select * from users limit 1 \G

********************************

id: 1

name: Alexandre Sporer

address: 24594 Emmitt Locks, Greenfelderview, MT 48128

    dob: 2007-05-02

phone: 07683017318

state: MT

1 row in set (0.00 sec)

MariaDB [user_data]>
```

You see that we have name, address, date of birth, phone and state in this table for 1,000,000 users.

Now see if there are any indexes on the table:

```
MariaDB [user data]> show indexes from users \G;
****
      Table: users
  Non unique: 0
    Key name: PRIMARY
Seq in index: 1
 Column name: id
   Collation: A
 Cardinality: 1002809
    Sub part: NULL
      Packed: NULL
       Null:
  Index type: BTREE
     Comment:
Index comment:
1 row in set (0.00 sec)
ERROR: No query specified
```

Note: if you see more than 1 index then drop the others:

```
MariaDB [user_data]> drop index state on users;
Query OK, 0 rows affected (0.03 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

Now lets get a count by an un-indexed column:

```
MariaDB [user_data]> select count(*), state from users gro
up by state \G
```

What were your results? We got:

So, 0.57 seconds. Now let's run an explain on the query:

```
MariaDB [user data] > explain
   -> select count(*), state from users group by state \G
****
         id: 1
 select type: SIMPLE
      table: users
       type: ALL
possible keys: NULL
       key: NULL
     key len: NULL
        ref: NULL
       rows: 999129
      Extra: Using temporary; Using filesort
1 row in set (0.00 sec)
```

So the size of the table means that Using temporary; Using filesort is used by MySQL to run the query.

If you add some criteria on a column in the query, this is the **EXPLAIN**:

```
MariaDB [user_data]> explain select count(*), state from u
sers where state = 'CA' group by state \G
*************************

id: 1
select_type: SIMPLE
```

```
table: users

type: ALL

possible_keys: NULL

key: NULL

key_len: NULL

ref: NULL

rows: 999129

Extra: Using where

1 row in set (0.00 sec)
```

Using temporary; Using filesort is gone. Now if you run the query:

Runtime is approximately 0.24 seconds.

Now, what if we index that column:

```
MariaDB [user_data]> alter table users add index(state);
Query OK, 0 rows affected (2.30 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

So now a b-tree structure is built to get only the rows matching the states we want.

Re-run the query:

```
MariaDB [user_data]> select count(*), state from users whe
re state = 'CA' group by state \G

**************************

****

count(*): 15195

state: CA
```

What are the results now?

## **Summary**

So we have seen where SQL databases (like MariaDB/MySQL) need to index a row by setting up an in-memory structure to make them perform better. In the next lab we'll see how NoSQL does a similar function but this time a little differently.