# **RDBMS and Performance**

MariaDB is a drop-in replacement for MySQL, the most used open source database for online applications. MariaDB falls into the category of the NewSQL products,

i.e. a product that provides unique NoSQL features together with the typical features available in relational databases.

Therefore, aspects like transaction management, durability and consistency are available together with schema or schema-less modeling,

full text storage and analysis and integration with other NoSQL technologies.

MariaDB can be part of the database infrastructure for Big Data. It is not meant to be a replacement for Hadoop, but it can be a technology used in conjunction with it. Hadoop is used in batch, ad-hoc analysis. In projects that require the processing of Terabytes or Petabytes of data,

Hadoop is definitely a good fit. The results can be queried and reported via a standard MySQL/MariaDB interface, which is compatible with virtually

all the BI tools and development frameworks available today.

Let's take a detour to check out the differences.

#### **RDBMS**

So now lets take a detour and look at a RDBMS.

```
<img src="https://user-
images.githubusercontent.com/558905/40613898-7a6c70d6-624e-11e8-
9178-7bde851ac7bd.png" align="left" width="50" height="50" 
title="ToDo Logo" /> 
<h4>1. Log in to MySQL DB</h4>
```

Since it's loaded we'll use MySQL (may be MariaDB also):

[centos@ip-10-0-0-54 ~]\$ mysql -u root -p

Enter password:

Welcome to the MariaDB monitor. Commands end with ; or \g

.

Your MariaDB connection id is 207

Server version: 5.5.56-MariaDB MariaDB Server

Copyright (c) 2000, 2017, Oracle, MariaDB Corporation Ab a nd others.

Type 'help;' or '\h' for help. Type '\c' to clear the curr ent input statement.

MariaDB [(none)]>



```
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<h4>2. Now look at the databases</h4>
```

```
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9178-7bde851ac7bd.png" align="left" width="50" height="50"
title="ToDo Logo" />
<h4>2. And create some database, and use it</h4>
```

```
MariaDB [(none)]> create database class_test;
Query OK, 1 row affected (0.00 sec)
```

```
MariaDB [(none)]> use class_test;
Database changed
MariaDB [class_test]>

<img src="https://user-images.githubusercontent.com/558905/40613898-7a6c70d6-624e-11e8-9178-7bde851ac7bd.png" align="left" width="50" height="50" title="ToDo Logo" />
```

```
MariaDB [class_test]> create table users(id int, name char
(20), PRIMARY KEY(id));
Query OK, 0 rows affected (0.00 sec)
```

<h4>3. Let's create a table</h4>

And another table - this one will be related to the users table by the FOREIGN KEY user\_id:

```
MariaDB [class_test]> create table orders(id int, user_id
int, order_info text, PRIMARY KEY(id), FOREIGN KEY (user_i
d) REFERENCES users(id));
Query OK, 0 rows affected (0.00 sec)
```

```
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images.githubusercontent.com/558905/40613898-7a6c70d6-624e-11e8-
9178-7bde851ac7bd.png" align="left" width="50" height="50"
title="ToDo Logo" />
```

```
MariaDB [class_test]> insert into users (id, name) VALUES
  (1, 'bill');
Query OK, 1 row affected (0.00 sec)
```

Select the table:

```
MariaDB [class_test]> select * from users;
+---+---+
| id | name |
+---+---+
| 1 | bill |
+---+---+
1 row in set (0.00 sec)
```

And we see the where clause is supported and so forth:

```
MariaDB [class_test]> select * from users where id < 2;
+---+---+
| id | name |
+---+---+
| 1 | bill |
+---+---+
1 row in set (0.00 sec)</pre>
```

```
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images.githubusercontent.com/558905/40613898-7a6c70d6-624e-11e8-
9178-7bde851ac7bd.png" align="left" width="50" height="50" 
title="ToDo Logo" /> 
<h4>5. Create a row in the related table</h4>
```

```
MariaDB [class_test]> insert into orders (id, user_id, ord
er_info) VALUES (1,1,"something");
Query OK, 1 row affected (0.00 sec)
```

#### And check it:

```
MariaDB [class_test]> select * from orders;
+---+----+
| id | user_id | order_info |
+---+----+
| 1 | 1 | something |
+---+----+
1 row in set (0.00 sec)
```

```
MariaDB [class_test]> insert into orders (id, user_id, ord
er_info) VALUES (2,1,"something");
Query OK, 1 row affected (0.00 sec)
```

So you see the way 2 becomes another order number, no problem:

But what if we do something like this:

```
MariaDB [class_test]> insert into orders (id, user_id, ord
er_info) VALUES (1,2,"something");
```

What has happened?

## **RDBMS and Transactions**

```
<img src="https://user-
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9178-7bde851ac7bd.png" align="left" width="50" height="50" 
title="ToDo Logo" /> 
<h4>6. Do the same with an active Transaction</h4>
```

```
MariaDB [class_test]> start transaction;
Query OK, 0 rows affected (0.00 sec)
```

Now, delete an order:

```
MariaDB [class_test]> delete from orders where id=2;
Query OK, 1 row affected (0.01 sec)
```

And the order is now gone:

```
MariaDB [class_test]> select * from orders;
+---+---+
| id | user_id | order_info |
+---+----+
| 1 | 1 | something |
+---+----+
1 row in set (0.00 sec)
```

But we made a mistake, so roll back the trransaction:

```
MariaDB [class_test]> rollback;
Query OK, 0 rows affected (0.00 sec)
```

```
MariaDB [class_test]> select * from orders;
+---+----+
| id | user_id | order_info |
+---+----+
| 1 | something |
| 2 | 1 | something |
```

# **NOSql Performance Indicators**

So now let's do some performance testing. We have a .sql file called users.sql in the directory above. :

```
<img src="https://user-
images.githubusercontent.com/558905/40613898-7a6c70d6-624e-11e8-
9178-7bde851ac7bd.png" align="left" width="50" height="50" 
title="ToDo Logo" /> 
<h4>7. Import into MySQL</h4>
```

```
[centos@ip-10-0-0-54 data]$ mysql -u root -p < users.sql
```

Now go into MySQL and see the table:

And count the table:

```
MariaDB [user_data]> select count(*) from users;
+-----+
| count(*) |
+-----+
| 1000000 |
+-----+
1 row in set (0.14 sec)
```

Rather large table, isn't it?

```
<img src="https://user-
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9178-7bde851ac7bd.png" align="left" width="50" height="50" 
title="ToDo Logo" /> 
<h4>8. Look at the Data</h4>
```

```
MariaDB [user_data]> select * from users limit 1 \G
    *****************************
    id: 1
    name: Alexandre Sporer
address: 24594 Emmitt Locks, Greenfelderview, MT 48128
    dob: 2007-05-02
    phone: 07683017318
    state: MT
1 row in set (0.00 sec)
MariaDB [user_data]>
```



the '\G' just prints the data lengthwise

You see that we have name, address, date of birth, phone and state in this table for 1,000,000 users.

Now see if there are any indexes on the table:

```
MariaDB [user_data]> show indexes from users \G;
**************************

****

Table: users

Non_unique: 0

Key_name: PRIMARY

Seq_in_index: 1
```

Column\_name: id

Collation: A

Cardinality: 1002809

Sub\_part: NULL

Packed: NULL

Null:

Index\_type: BTREE

Comment:

Index comment:

1 row in set (0.00 sec)

ERROR: No query specified



if you see more than 1 index then drop the others:

```
MariaDB [user_data]> drop index state on users;
Query OK, 0 rows affected (0.03 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

Now lets get a count by an un-indexed column:

```
MariaDB [user_data]> select count(*), state from users gro
up by state \G
```

What were your results? We got:

#### So, 0.57 seconds. Now let's run an explain on the query:

```
MariaDB [user_data]> explain
    -> select count(*), state from users group by state \G
***************************

    id: 1
select_type: SIMPLE
    table: users
    type: ALL
possible_keys: NULL
    key: NULL
    key=len: NULL
    ref: NULL
    rows: 999129
Extra: Using temporary; Using filesort
```

```
1 row in set (0.00 sec)
```

So the size of the table means that Using temporary; Using
filesort is used by MySQL to run the query.

If you add some criteria on a column in the query, this is the **EXPLAIN**:

```
MariaDB [user data]> explain select count(*), state from u
sers where state = 'CA' group by state \G
***
         id: 1
 select type: SIMPLE
      table: users
       type: ALL
possible keys: NULL
        key: NULL
     key len: NULL
        ref: NULL
       rows: 999129
      Extra: Using where
1 row in set (0.00 sec)
```

Using temporary; Using filesort is gone. Now if you run the query:

```
MariaDB [user_data]> select count(*), state from users whe
re state = 'CA' group by state \G
```

Runtime is approximately 0.24 seconds.

```
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title="ToDo Logo" /> 
<h4>9. Index the Column</h4>
```

Now, what if we index that column:

```
MariaDB [user_data]> alter table users add index(state);
Query OK, 0 rows affected (2.30 sec)
Records: 0 Duplicates: 0 Warnings: 0
```

So now a b-tree structure is built to get only the rows matching the states we want.

Re-run the query:

```
MariaDB [user_data]> select count(*), state from users whe
re state = 'CA' group by state \G
```

What are the results now?

### Results

So we have seen where SQL databases (like MariaDB/MySQL) need to index a row by setting up an in-memory structure to make them perform better.

```
<button type="button"><a href="https://virtuant.github.io/hadoop-
overview-spark-hwx/">Go Back</a></button>
<br><br><br><br><br>
```