

## Trusted Computing

1. Student presentation
2. Questions from reviews:
  - a. Used in real world?
  - b. Mobile devices?
  - c. App-specific OS?
  - d. Device drivers?
  - e. Revocation – key compromise?
    - i. Example – DigiNotar
      1. Cert must be revoked by browsers; all certs it issued must be reissued
    - ii. Comodo – registration authority compromised
      1. Couldn't verify businesses asking for certs were the real business
  - f. Poor isolation in OS:
    - i. Shared resources in kernel (/proc, filesystem)
3. Comment:
  - a. Done in 2003; vmware just released in 1999 but not widely used
4. Public key crypto / protocols
  - a. General infrastructure: CA hierarchy
  - b. How do CA's protect key?
  - c. MS example:
    - i. locked room off campus
    - ii. Armed guard
    - iii. Tamper-proof hardware : chip surrounded by mercury that shorts out storage
5. Problem statement
  - a. What can applications trust today?
    - i. Can they trust the system to keep data private
    - ii. Can they games trust the user not to cheat
    - iii. Can banks trust client banking software not to disclose secrets?
  - b. Fact: today, physical access to a computer guarantees you full access to the information **on that computer**
    - i. Question: How?
      1. Debuggers
      2. Virtual Machines
      3. Hook audio / video path
    - ii. Any technique only makes it harder, but not impossible

- c. What if you have information so valuable you can't entrust it to the person using the computer?
    - i. Goal: Root Secure
    - ii. QUESTION: Examples?
      - 1. Digitally encoded songs / movies
        - a. Don't trust user not to give away / sell
      - 2. Sensitive legal documents
        - a. Don't trust user not to leak to competitors or press
      - 3. Networked applications – multiplayer games
        - a. Don't trust users not to cheat
      - 4. Secure video conferencing devices
    - d. Insecure OS allows one app to subvert all others
      - i. E.g. WinXP sp2 used in a botnet of 73,000 machines capable of sending 1 billion individual spam messages a day (bot runs anti-virus/anti-spyware to remove other bots)
- 6. Bad solution: closed platforms
  - a. Lock the machine
  - b. Obscure the hardware
  - c. Make it tamper-proof
    - i. Circuits that short-circuit if opened
    - ii. Example: Xbox
- 7. Solution: attestation / digital
  - a. Two General frameworks:
    - i. Layered code 0 (hardware) → Y (application)
    - ii. At run time, code at layer X checks signatures of layer 0 .. X-1 and X (check layers below)
      - 1. If correct, knows that is running on correct platform
      - 2. QUESTION: Why can you trust the signature of the layer below?
        - a. A: the hardware signatures are trusted
        - b. Can follow the chain of trust
    - iii. Code at layer X checksums / verifies layer X+1 (check layers above)
      - 1. Can verify next layer up is not corrupt if needed
      - 2. Can be used to make sure a computer only boots trusted code
    - iv. QUESTION: What is the difference?
      - 1. Windows/Chrome OS/IOS use option 2 – check layer above
- 8. Hardware support in TPM

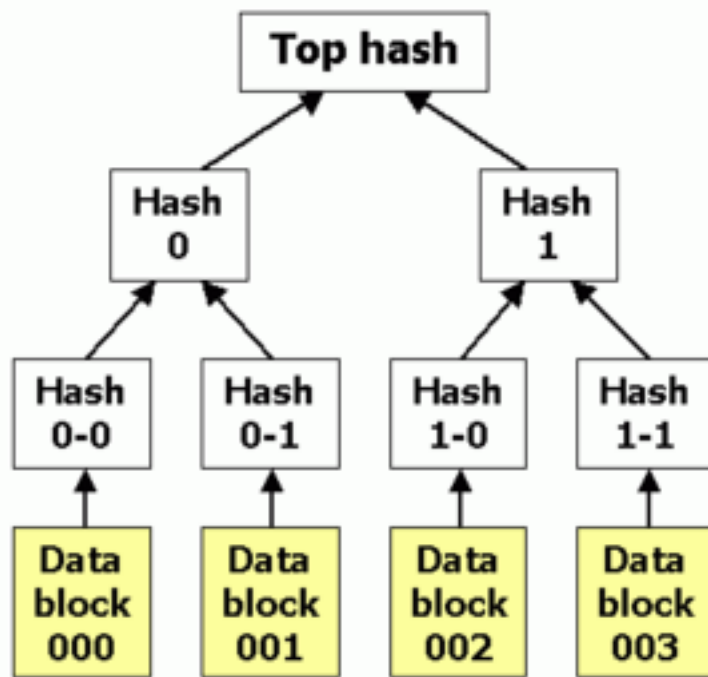
- a. Endorsement key – 2048 bit random private/public key; private key never leaves chip.
    - i. Used to sign a random number – show knowledge of endorsement key
  - b. Memory curtaining
    - i. Memory not available to OS – to store secure software
  - c. Sealed storage
    - i. Encrypted non-volatile storage that can only be decrypted if software with correct hash is running
    - ii. Need to present certified hash to decrypt the data
- 9. Implementation of Terra – using 1<sup>st</sup> framework of checking layers below
  - a. Problem with second: greatly restricts what you can run. How do you get flexibility and openness?
  - b. Vision: virtual appliances
    - i. Code can trust it runs in the right place
    - ii. Can have a customized OS for each applications
    - iii. QUESTION: Is this reasonable?
      - 1. Yes – change configuration of Linux, use BSD / WinXP Embedded, Flux OS kit
  - c. Layers
    - i. Hardware
      - 1. Provides secure storage, secret key
      - 2. HW must attest to booted system
        - a. Checksum & sign TVMM or boot loader
      - 3. Sealed storage – tamperproof storage controlled by coprocessor, only allows OS with same hash to unseal – for storing private keys
      - 4. Can buy today!
    - ii. Virtual machine
      - 1. Provides closed & open box environment
        - a. Close box cannot be inspected or manipulated by owner
      - 2. Management VM allows
        - a. Starting
        - b. Stopping
        - c. Installing devices
      - 3. Relatively small monitor – economy of mechanism, keep it small
    - iii. Multiple OSs
      - 1. Can run any complexity of OS you desire
      - 2. Communicate with virtual NIC – treat other VMs as external

- iv. Applications
- d. Attestation – provide a signature chain showing what code is running
  - i. 2 issues
    - 1. What code is running? Don't care what code, but want to get verifiable checksum of code
    - 2. Is the right code running? Can compare against known certificates
  - ii. Each layer is checksummed & signed by lower level
    - 1. Hardware checksums firmware
    - 2. Firmware checksums loader
    - 3. Loader checksums TVMM
    - 4. TVMM checksums VMs / OS
    - 5. NOTE: you can buy the hardware today + MS includes code that uses it (IBM Laptops have it)
  - iii. Checksum includes:
    - 1. Hash of all persistent state
    - 2. E.g. bios, executable code, constant data
    - 3. Does not include removable disks
  - iv. For attestation: a layer can call layer below for an attestation
    - 1. Layer X provides public key to layer X-1
    - 2. Layer X-1 signs checksum of Layer X state and key
    - 3. QUESTION: Why include the public key?
      - a. A: so can sign for the layer above – makes a chain
    - 4. QUESTION: What does this tell you?
      - a. Which code is running, not whether code is the right code or not
  - v. **Example:** Quicken
    - 1. On secure connection to server, pass list of attestations for all layers of software
      - a. HW signs bios, bios signs loader, loaders signs TVMM,
        - i. TVMM signs for VM – including OS plus quicken (a packaged appliance)
      - b. Send keys for each layer – public key for hw, bios, boot loader, TVMM, OS, etc.
      - c. Send certificate for HW from vendor
      - d. Verify

- i. Check certificate for HW against vendor (e.g. intel)
  - ii. Check hash of next level software & signature using public key provided
  - iii. Verify VM hash (OS+quicken) against certificate from vendor (obtained somewhere, could be the VM itself, embedded)
- vi. How to use:
  - 1. SW calls lower levels for attestation, pass certificate chain to other side of a channel
  - 2. Other side (server) has database of correct hashes, verifies that it is correct
  - 3. Must do within an authenticated channel to avoid man-in-the-middle attacks
- vii. Two certificate chains:
  - 1. One external: certify the software is known to be good (guest VM, TVMM, its bios)
    - a. CA certifies quicken, Quicken certifies a certain VM is good,
    - b. CA certifies vmware, vmware certifies TVMM
  - 2. One internal: certify what code is actually running (based on CPU key on up)
    - a. CA certifies HW – Intel, Intel certifies key in CPU (tamper resistant), HW certifies bootlaoder/TVMM, TVMM certifies guest VM
- viii. QUESTION: How do you know whether to trust the code that is running?
  - 1. Server keeps certificates of known good code
  - 2. Certificates come from vendors:
    - a. Intel for processor/ firmware
    - b. MS for OS
    - c. Quicken for application
- ix. QUESTION: Who needs to trust CA? If CA broken, who has to update? How do you securely update?
  - 1. A: Update servers that validate clients (or everybody in peer-to-peer)
  - 2. Problem: no way to bootstrap trust to new CA. Need to roll over to another, unbroken CA instead.

- x. QUESTION: What if HW key (e.g. Dell, Intel) is broken?
    - 1. Can no longer trust anything signed by that company; must put new keys/certificates in all hardware securely
  - xi. QUESTION: How much trust should you put in attestation?
    - 1. Can trust the right code was loaded
    - 2. QUESTION: can you tell it is still running?
      - a. Not really – could have executable data
      - b. Still have bugs in code that could be subverted
    - 3. Hardware may not be tamper proof
      - a. Someone could read secrets from DRAM, modify contents of DRAM externally (e.g. DMA from a device)
  - xii. QUESTION: Who does the checking?
    - 1. In Terra, a remote network machine
    - 2. QUESTION: How do it locally? Have to check at layer X-1
      - a. Need to install a key for app, or have a certificate chain for certifying apps
      - b. QUESTION: What are implications?
  - xiii. QUESTION: what about patching?
    - 1. A: must bundle in service packs
    - 2. A: reduced size OS means reduced patches
  - xiv. QUESTION: How limit consumer choice?
    - 1. Services can refuse to work with clients running non-approved software
      - a. E.g. Windows phones, Linux desktops, non-Chrome browsers
      - b. Example services: Netflix, Spotify –want to make sure you don't steal video/music when playing
- 10. Control over devices
  - a. TVMM/trustworthy computing separates out platform user from platform admin
    - i. Admin – boots TVMM, controls system booting
    - ii. User – choose what VMs to load
      - 1. TVMM just testifies/attests as to what runs, doesn't stop anything
  - b. "root secure" – hw is secure from root user or physical access

- i. no boot into single-user mode or enter bios commands to bypass security
  - c.
- 11. Implementation
  - a. Device access
    - i. QUESTION: do you need to trust devices? How much?
    - ii. Storage: 3 levels
      - 1. Encrypted with per-VM key
      - 2. Integrity checked: public but not secret data, e.g. binaries
      - 3. Unchecked – for sharing
      - 4. Choice depends on performance needs
    - iii. Network: rely on higher-level protocols for security
      - 1. E.g. ssl, ipsec
    - iv. Other devices
      - 1. Must prevent DMA into memory – need HW
        - a. IOMMU does remapping to prevent write to arbitrary memory
      - 2. Where do device drivers run? In TVMM or in VM?
        - a. If can partition devices, can run in VM
          - i. No need for sharing; just needs access to MMIO space + interrupt delivery
        - b. If can have untrusted drivers, can run in VM
          - i. Q: What do you need?
            - 1. Device encrypts, OS decrypts
          - ii. What does this protect?
            - 1. Sniffing, but not human eyes
  - 3.
- b. Performance
  - i. Problem: checksum of a large data item is slow
    - 1. E.g. entire virtual disk
    - 2. Observation: don't typically read entire disk at once
  - ii. Instead: want to checksum individual blocks for fine-grained performance
  - iii. Partial checksums of data at layer X
    - 1. Merkel Trees - [http://en.wikipedia.org/wiki/Hash\\_tree](http://en.wikipedia.org/wiki/Hash_tree)



- 2.
  3. Property: verify  $\log(N)$  checksums on block access
  - iv. Ahead-of-time attestation:
    1. Hash entire code for next stage of computation before executing
      - a. E.g. bootloader
  - v. Optimistic Attestation:
    1. VM Specifies desired checksum in advance
    2. Assume checksum correct
    3. Check blocks when first accessed, not when loading code
    4. Only fail if find a block with bad checksum
    5. QUESTION: Why does this work?
      - a. A: always fail before bad block is used
12. Hardware needed
- a. Attestation
    - i. Can checksum code & sign before running
  - b. Sealed storage
    - i. Encrypt data tied to hash of OS
      1. Only OS with same hash can decrypt
      2. Used by TVMM to store its private key – can only be retrieved by TVMM with same hash
  - c. HW virtualization
  - d. Secure IO
    - i. Devices that encrypt data



- ii. Simple encrypted interface, full untrusted interface  
(e.g. 2d graphics for TVMM, 3d for everyone else)
  - e. Device isolation
    - i. TVMM cannot be modified by a device
- 13. Political issues
  - a. Who designs systems like this?
    - i. Stanford – willing to play in commercial space
    - ii. Berkeley / MIT – more likely to build systems to defeat this
  - b. Who benefits from systems like this?
    - i. Music companies, movie companies
      - 1. Encode data such that needed trusted environment to play data
    - ii. Software companies
      - 1. Require trusted software to prevent piracy
    - iii. Large SW companies
      - 1. Can make file formats that can only be read by their code
      - 2. Can lock customers in
      - 3. If have DRM on mail, for legal reasons may need to transfer DRM – hard to move between vendors
      - 4. Content providers may tend to certify only a small number of large companies – lock out small companies
    - iv. Game companies / ethical gamers
      - 1. Fewer cheaters
    - v. Small companies:
      - 1. Can deploy their own platforms
      - 2. Less likely to be on existing lists of “trusted software/firmware/hardware”
    - vi. Users?
      - 1. Not clear what their benefit
      - 2. Claim: Content providers will provide more content if it can be made secure
      - 3.