# CS339: Abstractions and Paradigms for Programming

The Procedural Paradigm

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## Recall Substitution Model of Procedure Application

To apply a compound procedure to arguments, evaluate the body of the procedure with each formal parameter replaced by the corresponding argument.



# A "Procedural" Example

```
> (define (square x) (* x x))

> (define (sum-of-squares x y)
          (+ (square x) (square y)))

> (define (f a)
          (sum-of-squares (+ a 1) (* a 2)))
```



### Solve using substitution

```
(define (sum-of-squares x y)
  (+ (square x) (square y)))
(define (square x) (* x x))
(define (f a)
  (sum-of-squares (+ a 1) (* a 2)))
```

```
(f 5)
(sum-of-squares (+ 5 1) (* 5 2))
(+ (square (+ 5 1)) (square (* 5 2)))
(+ (* (+ 5 1) (+ 5 1)) (* (* 5 2) (* 5 2)))
(+ (* 6 6) (* 10 10))
(+ 36 100)
136
```



### Could we have substituted this way?

### Applicative Order of Evaluation

```
(f 5)
(sum-of-squares (+ 5 1) (* 5 2))
(sum-of-squares 6 10)
(+ (square 6) (square 10))
(+ (* 6 6) (* 10 10))
(+ 36 100)
136
```

```
(define (sum-of-squares x y)
  (+ (square x) (square y)))
(define (square x) (* x x))
(define (f a)
  (sum-of-squares (+ a 1) (* a 2)))
```

Evaluate arguments then apply



## Solve using substitution

#### Normal Order of Evaluation

```
(define (sum-of-squares x y)
  (+ (square x) (square y)))
(define (square x) (* x x))
(define (f a)
  (sum-of-squares (+ a 1) (* a 2)))
```

```
(f 5)
(sum-of-squares (+ 5 1) (* 5 2))
(+ (square (+ 5 1)) (square (* 5 2)))
(+ (* (+ 5 1) (+ 5 1)) (* (* 5 2) (* 5 2)))
(+ (* 6 6) (* 10 10))
(+ 36 100)
136
```

Fully substitute then reduce



### Applicative vs Normal Orders of Evaluation

- ➤ Applicative order
- Call by value

Call by name (need)

- > avoids redundant computation
- ➤ takes lesser memory

```
(f 5)
(sum-of-squares (+ 5 1) (* 5 2))
(sum-of-squares 6 10)
(+ (square 6) (square 10))
(+ (* 6 6) (* 10 10))
(+ 36 100)
136
```

```
(f 5)

(sum-of-squares (+ 5 1) (* 5 2))

(+ (square (+ 5 1)) (square (* 5 2)))

(+ (* (+ 5 1) (+ 5 1)) (* (* 5 2) (* 5 2)))

(+ (* 6 6) (* 10 10))

(+ 36 100)

136
```

➤ However, there are advantages of normal order too, and there are ways to make it more efficient (post mid-sem).

- ➤ Which one does Scheme use?
- ➤ How can you find out for any language?



### Conditionals

➤ Absolute value of x:

$$|x| = \begin{cases} x, & \text{if } x > 0 \\ 0, & \text{if } x = 0 \\ -x, & \text{if } x < 0 \end{cases}$$

**Action** 

$$|x| = \begin{cases} -x, & \text{if } x < 0 \\ x, & \text{otherwise} \end{cases}$$

OR



### Example: Newton's method for computing square roots

- ➤ Start with a guess and "improve" the guess until it is "good enough"
- ➤ Square root of 2:

Guess (y)	Quotient (x/y)	Average ((y+x/y)/2)
1	2/1 = 2	(1+2)/2 = 1.5
1.5	2/1.5 = 1.3333	1.4167
1.4167	1.4118	1.4142
1.4142		

➤ Say we stop when the square of the guess is equal to the number up to three decimal places.



# Example: Newton's square root [Cont.]

```
(define (sqrt-iter guess x)
 (if (good-enough? guess x)
      guess
      (sqrt-iter (improve guess x) x)))
(define (improve guess x)
 (average guess (/ x guess)))
(define (average x y)
 (/(+xy)2)
(define (good-enough? guess x)
 (< (abs (- (square guess) x)) 0.001))</pre>
(define (sqrt x)
 (sqrt-iter 1.0 x))
```

➤ Can you identify two "bad" things in this code?



### Example: Newton's square root [Cont.]

```
(define (average x y)
 (/(+ x y) 2))
(define (square x)
 (* X X)
(define (sqrt x)
  (define (improve guess)
    (average guess (/ x guess)))
  (define (good-enough? guess)
    (< (abs (- (square guess) x)) 0.001))
  (define (sqrt-iter guess)
    (if (good-enough? guess)
        guess
        (sqrt-iter (improve guess))))
  (sqrt-iter 1.0))
```

### Namespace Abstraction



Packaged together

#### Next class:

Why did we have sqrt-"iter"?

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