

Q1 (A) #t

(cddr (list (list 2 3 4) 5))

First cdr : (list 5)

Second cdr : ()

$\frac{1}{2}$ for #t; $\frac{1}{2}$ for explanation
(cut $\frac{1}{2}$ for variations of #t)

(B) 147

Outer lambda applied to 3, 5

$\therefore x: 3$

$y: 5$

Inner (returned) lambda applied to 2, 7

$\therefore z: 2$

$w: 7$

$(\ast (+ 2 5) (\ast 7 21)) = 147$
 $(\ast 3 7))$

$\frac{1}{2} + \frac{1}{2}$

Any variation
of this explanation
is fine

Q2 (A) (define (fib n))

(define (fib-iter a b count)

(if (= count 0)

b

(fib-iter (+ a b) a (- count 1))))

(fib-iter 1 0 n))

(fib 5)

(fi 1 0 5)

(fi 1 1 4)

(fi 2 1 3)

(fi 3 2 2)

(fi 5 3 1)

(fi 8 5 0)

0	1	1	2	3	5	8
↑	↑	↑	↑	↑	↑	
0	1	2	3	4	5	

1 mark
(wrong index: $-\frac{1}{2}$)

(B) Tail-call optimization allows reusing the stack frame when recursive calls to a procedure occur at the tail position; i.e., without any computation left to be performed after unwinding the recursion. 1 mark for a reasonable explanation (above process not discussed: -1/2)
This happens with the above process too. Thus, with TCO, the six calls to "fi" can share the same stack frame.

0.5 marks for when TCO happens (function calls itself/another function with no added computation, etc)

0.5 marks for how the optimization works (best would be showing the stack frames and

noting that there is only 1 stack frame at each moment, earlier stack frames are reused) and hence space complexity reduces

Q3 (A) car (cons v w)

$$=_{\beta} \text{car } ((\lambda f. \lambda s. \lambda b. b \ f \ s) \ v \ w)$$

$$=_{\beta} \text{car } ((\lambda s. \lambda b. b \ v \ s) \ w)$$

$$=_{\beta} \text{car } (\lambda b. b \ v \ w)$$

$$= (\lambda p. p \ \text{true}) (\lambda b. b \ v \ w)$$

$$=_{\beta} (\lambda b. b \ v \ w) \ \text{true}$$

$$=_{\beta} \text{true} \ v \ w$$

$$= v$$

2 marks

(partial if jumped steps)

(B) cdr = $\lambda p. p \ \text{false}$, where $\text{false} \ x \ y = y$

1 mark if both

cdr & false are defined

0.5 marks if false not defined explicitly. (Implicitly assuming definition of false in derivation does not count)

Q4 (define (reverse l)

(if (null? l)

nil

; l also works

(append (reverse (cdr l)) (list (car l))))

□

2 marks if completely right

-1/2 without this

1 for missing base case. -0.5 for minor misplaced parentheses -0.5 for not having the arguments of append as lists

Bonus

(define (flip l)

(cons b a))

(define (reverse l)

(foldl flip nil l))

Put a ✖ in the
front page if done right.

Only fully correct solutions for a PC. Allowed cases where foldl was redefined to have different order or arguments for f.