

ASSIGNMENT - 3

1 The Relational Model

1.1

branch_no \rightarrow city, street_number, postcode

employee-ID \rightarrow employee-name

customer-ID \rightarrow customer-name

account-number \rightarrow balance

1.2

{customer-ID} + is the list of closures w.r.t customer-ID.

customer-ID \neq {customer-ID, Customer-name, account-number, balance, branch_no, street_number, city, postcode}

where

Customer-ID \rightarrow Customer-name (functional dependency from the Customer Data Attributes).

Therefore, customer-ID \neq {customer-name}

1.3

Relations:

BRANCH:

Branch (branch-no, street-number, city, postcode)

EMPLOYEE:

Employee (employee-ID, Employee-name)

CUSTOMER:

Customer (customer-ID, customer-name)

ACCOUNT:

Account (account-number, balance)

Relationships:

Employee to Branch:

Employee ("Works_In") Branch – Many-to-One relation.

An employee only works for one branch, but a branch can have many employees

Customer to Account:

Customer ("Has") Account – Many-to-Many relation.

A customer can have more than one account and an account may be owned by more than one customer.

New Relations:

BRANCH:

Branch (branch-no, street-number, city, postcode)

EMPLOYEE:

Employee (employee-ID, employee-name, branch-No*)

CUSTOMER:

Customer (customer-ID, customer-name)

ACCOUNT:

Account (account-number, balance)Works_In (employee-ID*, branch_no*)

employee-ID and branch_no are the foreign keys of Employee and Branch tables, respectively.

Has (customer-ID*, account-number*)

customer-ID and account-number are the foreign keys of Employee and Branch tables, respectively.

Or

A new table can be created instead of Has relationship. Therefore, new relations will be,

BRANCH:

Branch (branch-no, street-number, city, postcode)

EMPLOYEE:

Employee (employee-ID, Employee-name, branch-No*)

CUSTOMER:

Customer (customer-ID, customer-name)

ACCOUNT:

Account (account-number, balance)

CUSTOMER_ACCOUNT:

Cust_Acc (customer-ID*, account-number*)

2. Normalisation

2.1

Consider first two FDs,

- $\text{projID} \rightarrow \text{title, type, manager}$
- $\text{manager} \rightarrow \text{type}$

the above two relations imply the transitive dependency where the type can be deleted from the projID relation as the manager relation always implies to type.

Therefore, new FD is

- **$\text{projID} \rightarrow \text{title, manager}$**
- **$\text{manager} \rightarrow \text{type}$**

Similarly,

- $\text{projID, title, jobID} \rightarrow \text{contractNo}$

As projID determines title. Therefore, title can be removed

Therefore, it becomes

- $\text{projID, jobID} \rightarrow \text{contractNo}$

Also, jobID determines the projID. Therefore, the projID can be deleted.

New FD,

- $\text{jobID} \rightarrow \text{contractNo}$

And from

- $\text{jobID} \rightarrow \text{projID, start-date, end-date, contractor}$

Since, jobID determines contractNo we can write,

- $\text{jobID} \rightarrow \text{projID, contractNo, start-date, end-date, contractor}$

Now consider,

- **$\text{contractNo} \rightarrow \text{jobID, contractor, start-date, end-date}$**

Since, jobID determines contractNo we can write jobID by removing start-date, end-date, contractor. As they are already determined by contractNo.

Therefore, new FD is

- **$\text{jobID} \rightarrow \text{projID, contractNo}$**

Now consider the final FD

- $\text{jobID} \rightarrow \text{contractNo}$

We can remove this as jobID already determines contractNo

Minimal basis for given PNG functional dependencies:

- **$\text{projID} \rightarrow \text{title, manager}$**
- **$\text{manager} \rightarrow \text{type}$**
- **$\text{jobID} \rightarrow \text{projID, contractNo}$**
- **$\text{contractNo} \rightarrow \text{jobID, contractor, start-date, end-date}$**

2.2

- The relation PNG is not in BCNF or 3NF because it has transitive functional dependencies.
- The relation is not in 3NF because the functional dependencies are transitive dependencies which violates the 3NF rules.

Consider relations,

- **projID \rightarrow title, type, manager**
- **manager \rightarrow type**

projID implies both manager, type and manager imply type which is transitive functional dependency.

- Since, it is not 3NF it is not BCNF.

2.3

PMG relation into 3NF:

The relation is said to be in 3NF when it has the full functional dependency and should not have any transitive relations.

Therefore, FDs must be,

- **projID \rightarrow title, manager**
- **manager \rightarrow type**
- **jobID \rightarrow projID, contractNo**
- **contractNo \rightarrow jobID, contractor, start-date, end-date**

PROJECT:

Project (projID, title, manager*)

MANAGER:

Manager (manager, type)

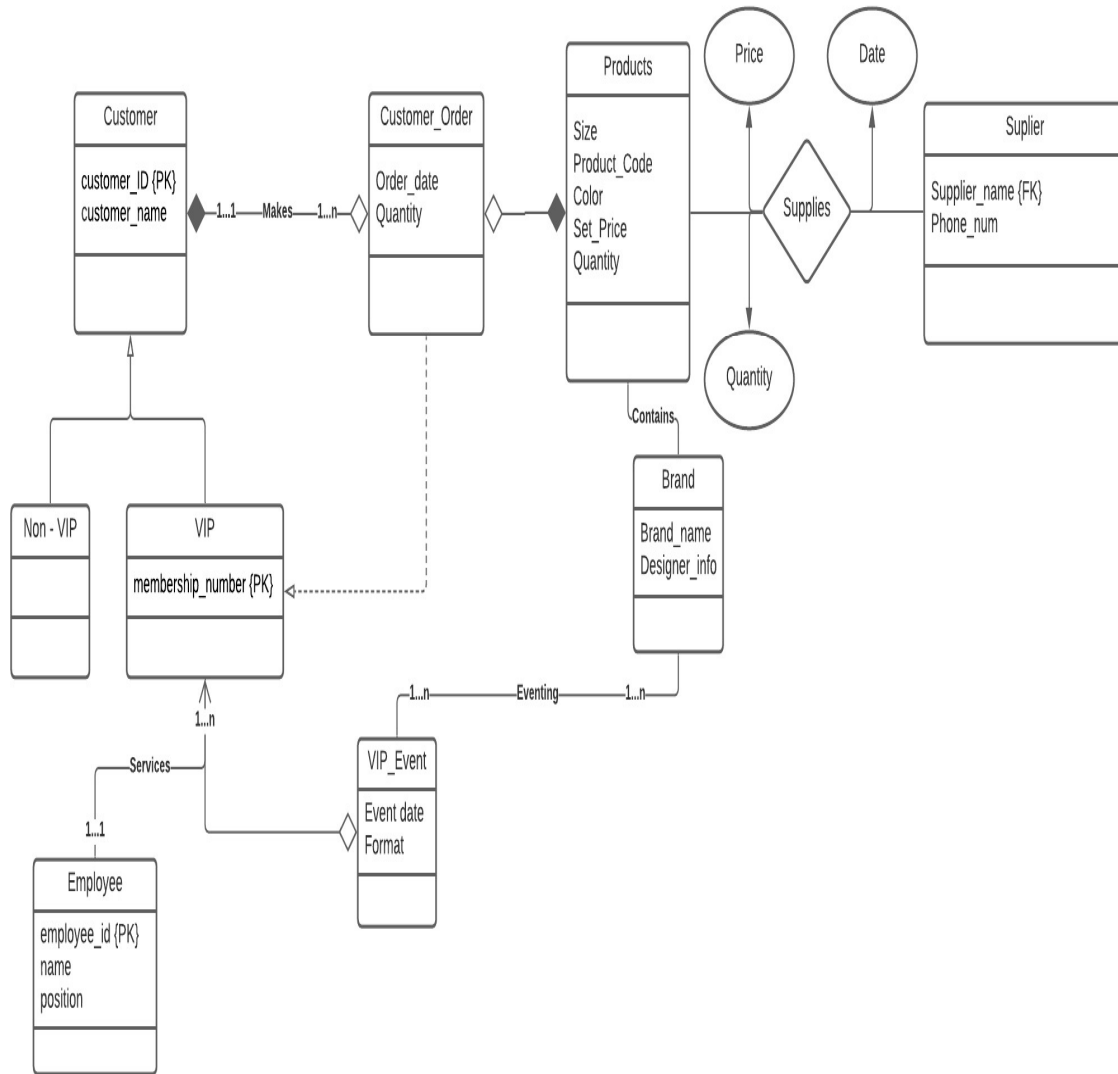
JOB:

Job (jobID, contractNo*, projID*)

CONTRACT:

Contract (contractNo, contractor, start-date, end-date, jobID*)

3 The ER Model



4 ER to Relational schema Mapping

STAFF:

Staff (empNo, givenname, surname, propertyNo*)

PROPERTY:

Property (propertyNo, streetnumber, street, suburb, postcode)

CLIENT:

Client (custNo, givenname, surname, phoneno, creditrating)

OPENSESSION:

OpenSession (sessionNo, opendate, opentime, propertyNo*)

LEASE:

Lease (leaseNo, start-date, end-date, propertyNo*, custNo*)

OWNER:

Owner (ownerNo, address, phoneno, propertyNo*)

IndividualOwner (ownerNo*, givenname, surname)

EntityOwner (ownerNo*, ABN, entityname)

Or

Owner (ownerNo, address, phoneno, givenname, surname, ABN, entityname, propertyNo*)

(Since, inheritance)