

✓ IP V4 / IP Header / IP datagram :-

VER	HLEN	DS	Total length						
4 bits	4 bits	8 bits	16 bits						
Identification		Flags	Fragmentation offset						
16 bits		3 bits	13 bits						
Time to Live	Protocol	Header checksum							
8 bits	8 bits	X X 16 bits							
Source IP address									
Destination IP address									
(Host)	Internet	0000							
option	option	1000							
utilization	utilization	0100							
Implementation	Implementation	0010							
Protocol	Protocol	0001							
Data									

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1) VER (Version)

- It tells which version of IP is used.
- $VER = 4$ (IPV4)
- $VER = 6$ (IPV6)

2) HLEN

- It contains value which when multiplied by four gives header length.
- For e.g. $HLEN = 5$ bytes
- $\therefore \text{Header length} = 20 \text{ B } (5 \times 4)$

Minimum header length is of 20 bytes.
In presence of options, it varies between 20B to 60B.

3) DS (Differentiated Services)

It tells whether any special type of service is required to deliver the packet between sender & receiver.

D: Minimize delay	R: Maximize reliability
T: Maximize throughput	C: Minimize cost
X X X D T R C X ← Don't care	Precedence TOS Bits

Types of services

TOS Bits	Description
0000	Normal (default)
0001	Minimize cost
0010	Maximize reliability
0100	Maximize throughput
1000	Minimize delay

4) Total length

- It is responsible to give the total size of the packet.
 $\text{Total length} = \text{Header length} + \text{Data length}$.

5) Fragmentation fields

- ① Identification
- ② Flags } Refer fragmentation.
- ③ Fragmentation offset }

6) Time to Live (TTL)

- This field limits unnecessary IP packet wandering ie the packet is routing in the network without reaching the destination.
- Router accepts the packet, checks the value of TTL, if TTL is zero, discards the packet else decrements TTL by one and pass it to the next router.

7) Protocol

It tells which protocol is responsible to handle this data.

Protocols

Value	Protocol	Description
1	ICMP	Internet Control Management Protocol
2	IGMP	
6	TCP	
17	UDP	
89	OSPF	Open shortest path first

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8) Checksum

- It is used for error detection to find errors (noise) is added in the header.

9) Source IP and Destination IP

It is 32 bit IP address of sender and receiver.

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V.V.SMP FRAGMENTATION :-

- The process of dividing the packets into smaller known as fragments is known as fragmentation.

- MTU (Maximum Transfer Unit) tells whether the packet should be fragmented or not (At sending side).

- MTU knows the network carrying capacity and depending on the size of the packet the fragmentation decision is taken.

If packet size > network carrying capacity then do the fragmentation.

Tuesday

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Fragging Fragmentation fields in IPV4 Header

1) Identification :-

In network every packet is uniquely identified by identification number.

and every fragment of the same packet would be assigned with same identification number which would help for reassembly at receiving side

Note : Order of sending the fragment & receiving the fragment is not same as all fragments may follow different paths to reach the destination

Flags:

DF MF

DF : Do not fragment.

MF : More fragment.

X : Don't care

DF : It tells whether the packet can be fragmented (if required) or not.

DF = 1 not to fragment

DF = 0 can be fragmented

MF : It tells whether it is a packet or first fragment or middle-fragment or last fragment

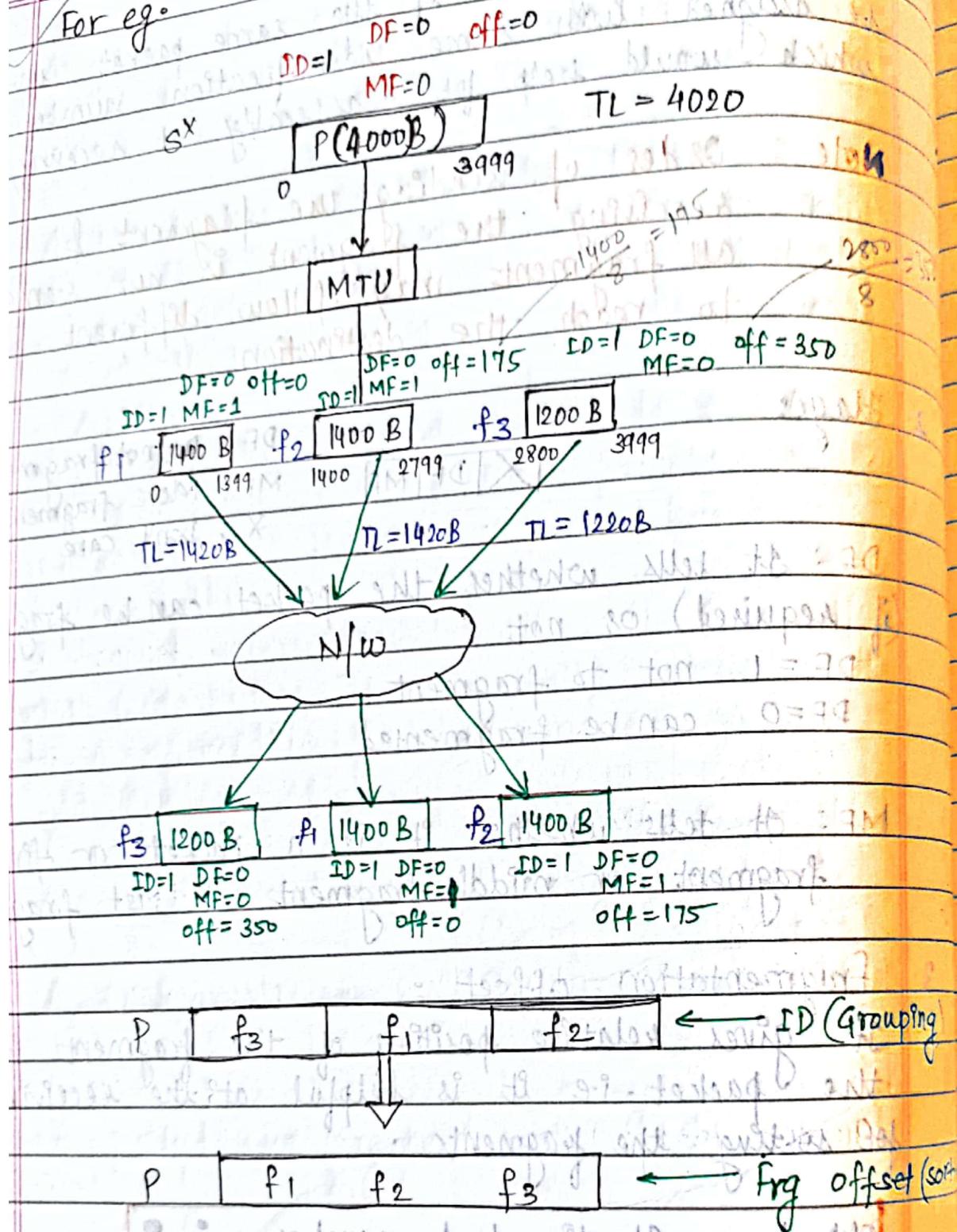
Fragmentation offset :

It gives relative position of the fragment in the packet i.e. it is helpful at the receiving for sorting the fragments.

Offset = Starting byte number $\div 8$.

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For e.g.



As shown in the diagram above, there is a packet of size 4000B which is divided into 3 fragments and the 3 fragments are sent sequentially but are received in any order.

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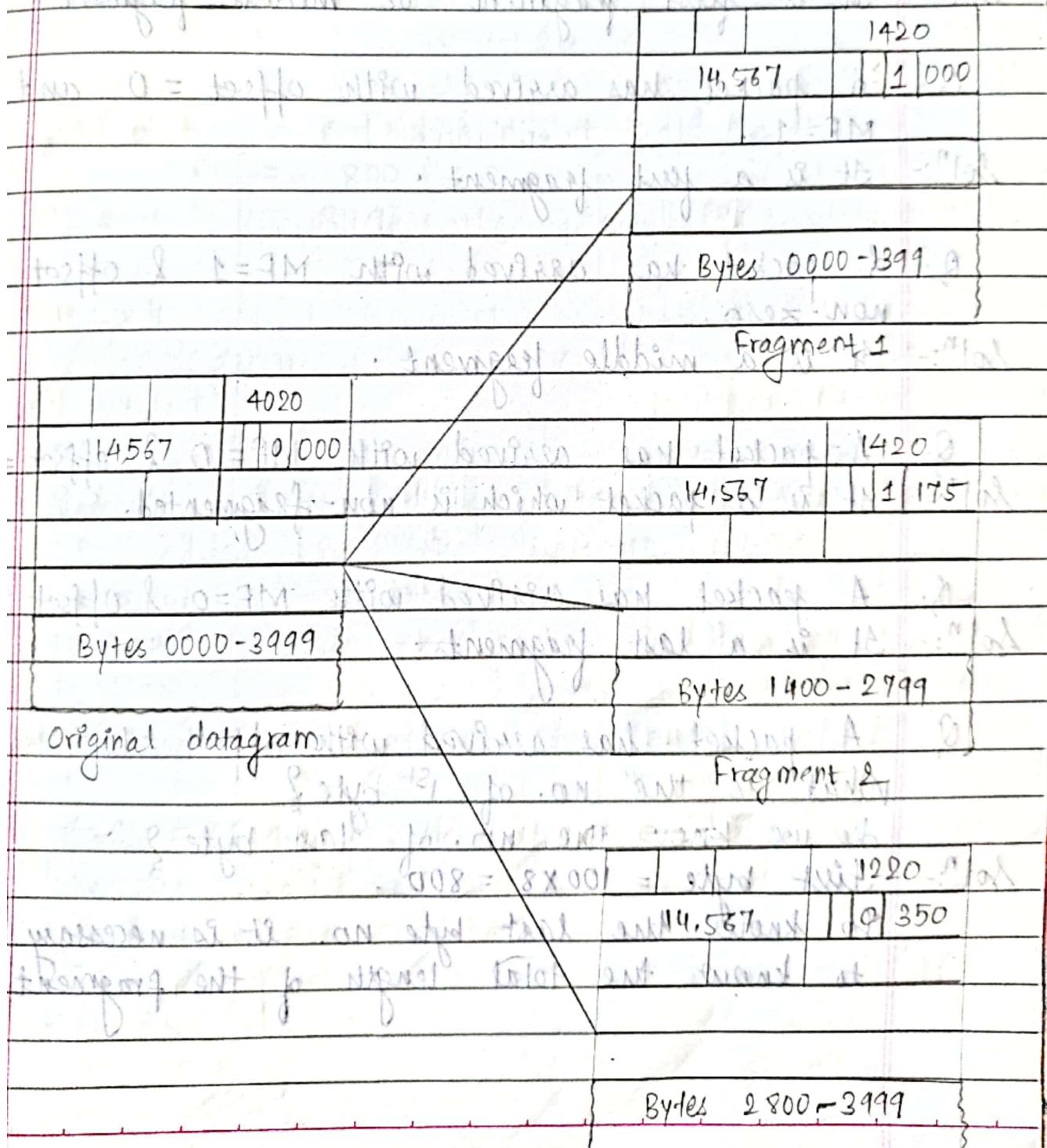
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At the receiving side, all the fragments are grouped into single packet with the help of identification no. & sorted w.r.t fragmentation offset.

For example : of detailed fragmentation w.r.t IP header.



Fragment 3

Scanned with CamScanner

Q. A packet has arrived with MF=0. Is it a packet or a fragment ?

Soln:- It is a packet or a last fragment.

Q. A packet has arrived with MF=1
Is it the first fragment, middle fragment
or the last fragment ?

Soln:- It is first fragment or middle fragment.

Q. A packet has arrived with offset = 0 and
MF=1.

Soln:- It is a first fragment.

Q. A packet has arrived with MF=1 & offset is
non-zero.

Soln:- It is a middle fragment.

Q. A packet has arrived with MF=0 & offset=0.

Soln:- It is a packet which is non-fragmented.

Q. A packet has arrived with MF=0 & offset is non-zero.

Soln:- It is a last fragment.

Q. A packet has arrived with offset = 100.

What is the no. of 1st byte ?

Do we know the no. of last byte ?

Soln:- First byte = $100 \times 8 = 800$.

To know the last byte no. it is necessary
to know the total length of the fragment.

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Q. A packet has arrived in which the offset value is 100, the value of HLEN is 5 and the value of the total

Soln:- Starting byte = $100 \times 8 = 800$

Header length = HLEN $\times 4 = 5 \times 4 = 20$

Total length = Data length + Header length

$$\therefore \text{Data Length} = \text{Total length} - \text{Header length}$$

$$= 100 - 20$$

$$= 80$$

$$\therefore \text{Last byte} = \text{first byte no.} + \text{Data length}$$

$$= 800 + 80$$

$$= 880$$

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V.IMP

(10 Marks)

Checksum

- It is the most commonly used error detection method that checks if there is any noise added in the network while traversing in the network.
- Receiver discards the packet in case of error.

Checksum Calculation at Sender Side.

Step-1: Divide IP header into k section each of n bits

Step-2: Set checksum field to zero

Step-3: Perform binary addition

Step-4: Take 1's complement of the sum, this is the checksum.

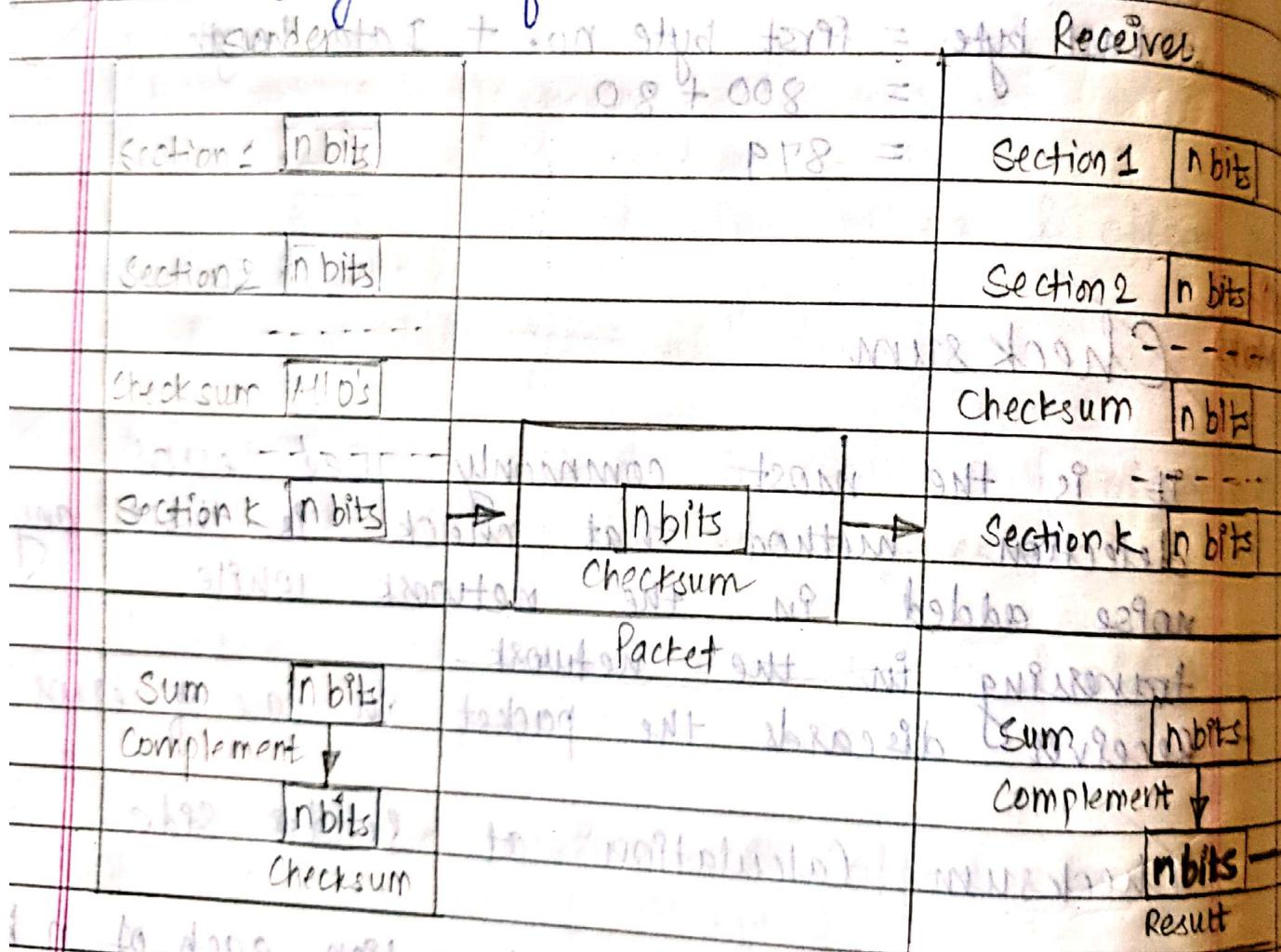
Checksum Calculation at the Receiver side.

- Step-1: Divide the IP header into k section each of n bits.
- Step-2: Do not set checksum to zero.
- Step-3: Perform binary addition $001 = 3$ bits
- Step-4: Take 1's complement of the sum, if the result is zero, accept it else discard it.

$001 + 001 = 010$

$010 - 001 = 001$

Block Diagram of checksum :-



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If the result is 0
keep otherwise

discard

for example :

4	5	00	0.09281	
	000110	0.	000	
4	17	0000		
	10.10.14.5			
	12.6.7.9			

4, 5, and 0 → 01000101 00000000
0028 → 00000000 00011100
0001 → 00000000 00000001
0 and 000 → 00000000 00000000
4 and 17 → 00000100 00010001
0000 → 00000000 00000000
10.12 → 00001010 00001100
14.5 → 00001110 00000101
12.6 → 00001100 00000110
7.9 → 00000111 00001001

sum → 01110100 01001110
checksum → 10001011 10110001

* even no of 1's - then ans is zero.

* odd no of 1's - then ans is 1 & add carry.

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