

# Bottom-up Parser

# Bottom-up Parser

- LR methods (Left-to-right, Rightmost derivation)
  - SLR
  - Canonical LR (CALR)
  - Look Ahead LR (LALR)
- Other special cases:
  - Shift-reduce parsing
  - Operator-precedence parsing

# Bottom-up parser

- Bottom-up parsers build a derivation by working from the input back toward the start symbol
  - Builds parse tree from leaves to root
  - Builds reverse rightmost derivation

# Handle

- Since Bottom-up parsers match the RHS of production with LHS, a concept called 'handle' is defined
- A *handle* is a substring of grammar symbols in a *right-sentential form* that matches a right-hand side of a production
- A handle's reduction to the non-terminal on the LHS represents one step along the reverse of a rightmost derivation
- This sub-string is a handle

# Handle - Example

- Expression Grammar Handles
  - id
  - $E * E$
  - $(E)$

# Shift Reduce Parser

- Simplest of the Bottom-up Parsers
- *Shift* input symbols until a handle is found.
- *Reduce* the substring to the non-terminal on the LHS of the corresponding production.

# Shift Reduce Parser

- A shift-reduce parser has 4 actions:
  - *Shift* the next input symbol is shifted onto the stack
  - *Reduce* the handle that is at top of stack
    - pop handle
    - push appropriate LHS symbol
  - *Accept* and stop parsing & report success
  - *Error* recovery routine is called

# Acceptance

- When the stack has the start symbol and the input is exhausted, the shift reduce parser goes to an accepting state



Consider a grammar

- 1.  $E \rightarrow E + E$
- 2.  $E \rightarrow E * E$
- 3.  $E \rightarrow \mathbf{id}$

| Stack            | Action           | Input           |
|------------------|------------------|-----------------|
| \$               | shift            | id + id * id \$ |
| \$ <u>id</u>     | Reduce by rule 3 | +id*id \$       |
| \$ E             | shift            |                 |
| \$ E+            | Shift            | id * id \$      |
| \$ E + <u>id</u> | Reduce by rule 3 | * id \$         |
| \$ <u>E + E</u>  | Reduce by Rule 1 | * id \$         |
| \$ E             | Shift            | * id \$         |
| \$ E *           | Shift            | id \$           |

# Parsing action

| Stack            | Action           | Input |
|------------------|------------------|-------|
| \$ E *           | Shift            | id \$ |
| \$ E * <u>id</u> | Reduce by rule 3 | \$    |
| \$ <u>E</u> * E  | Reduce by rule 2 | \$    |
| \$ E             | Accept           | \$    |

# Conflicts

- Shift-reduce and reduce-reduce conflicts are caused by
  - The limitations of the parsing method (even when the grammar is unambiguous)
  - Ambiguity of the grammar