

# MAXIMUM SUM SUBARRAY

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## ABSTRACT

The project aims at finding the maximum sum subarray in the given 2D array.

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# INTRODUCTION

We are given a 2D array, our aim is to find the maximum sum subarray in it.

For solving this we have analyzed two different approaches:-

- Brute force.
- Dynamic Programming.

These algorithms are tested for different sample test cases for time complexity and space complexity. In our report, we tend to focus on establishing a rule or a relationship between the time and input data.

# ALGORITHM DESIGN

## Brute force(Naive method)

- ▶ In the brute force approach we try to check every possible rectangle in the given  $n \times m$  2D array (where  $n, m$  are number of rows and columns respectively).
- ▶ Set the position of the top-left and bottom-right corners of the sub-rectangle and adding the integers within it while iterating through all the rows sequentially.
- ▶ Parallely we try to find the maximum subarray sum value.

# PSEUDO CODE (Naive approach)

```
int A[101][101]
```

```
function MAIN()
```

```
    maxSum ← INT_MIN
```

```
    tempSum ← 0
```

```
    x ← 0, y ← 0, z ← 0, w ← 0, n ← 0, m ← 0
```

```
    input n, m
```

```
    for i ← 0 to n-1
```

```
        for j ← 0 to m-1
```

```
            input A[i][j]
```

```
        end
```

```
    end
```

```
    for i ← 0 to n-1
```

```
        for j ← 0 to m-1
```

```
            for k ← 0 to n-1
```

```
                for l ← 0 to m-1
```

```
                    tempSum = FINDSUM(i, j, k, l)
```

```
                    if tempSum > maxSum
```

```
                        x ← i
```

```
                        y ← j
```

```
                        z ← k
```

```
                        w ← l
```

```
                    maxSum = tempSum
```

```
                    end if
```

```
                end
```

```
            end
```

```
        end
```

```
    end
```

```
    print x, y, z, w, maxSum
```

## MAXSUM PSEUDO CODE

function FINDSUM(x, y, z, w)

    sum  $\leftarrow$  0

    for i  $\leftarrow$  x to z

        for j  $\leftarrow$  y to w

            sum  $\leftarrow$  sum + A[i][j]

        end

    end

return sum

# KADANE'S ALGORITHM FOR 1D ARRAY

- ▶ This is an efficient approach to find the sum of contiguous subarray within a one-dimensional array of numbers that has the largest sum.
- ▶ The idea of Kadane's algorithm is to look for all positive contiguous segments of the array.
- ▶ Keep track of maximum sum contiguous segment among all positive segments. Each time we get a positive-sum compare it with maximum sum so far and update it, if it is greater than maximum sum till then.



## EXAMPLE

0	1	2	3	4	5	6
-2	-3	4	-1	-2	1	5

In this example maximum contiguous array sum is  $= 4 + (-1) + (-2) + 1 + 5$   
 $= 7$

# ALGORITHM DESIGN

## Dynamic programming approach

- ▶ We use Kadane's algorithm and prefix sum to reduce time complexity to  $O(n^2 \times m)$ .
- ▶ The idea is to fix the top and bottom rows one by one and find the maximum sum contiguous columns for every top and bottom row pair.
- ▶ We basically find left and right column numbers (which have maximum sum) for every fixed up and bottom row pair. To find the left and right column numbers, calculate the sum of elements in every column from top to bottom and store these sums in an array say `temp[]`.
- ▶ `temp[i]` indicates sum of elements from top to bottom in column `i`. If we apply Kadane's 1D algorithm on `temp[]`, and get the maximum sum subarray of `temp`, this maximum sum would be the maximum possible sum with top and bottom as boundary rows. To get the overall maximum sum, we compare this sum with the maximum sum so far

## EXAMPLE

1	2	-1	-4	-20
-8	-3	4	2	1
3	8	10	1	3
-4	-1	1	7	-6

In the following 2D array, the maximum sum subarray is highlighted with red rectangle and sum of this subarray is  $= (-3)+4+2+8+10+1+(-1)+1+7$   
 $= 29$ .

# DYNAMIC PROGRAMMING PSEUDO CODE

function KADANE(V)

maxSum  $\leftarrow$  INT\_MIN, tempSum  $\leftarrow$  0

st  $\leftarrow$  1, end  $\leftarrow$  1, localSt  $\leftarrow$  0

for i  $\leftarrow$  0 to length[V]

tempSum  $\leftarrow$  tempSum + V[i]

if maxSum < tempSum

st  $\leftarrow$  localSt

end  $\leftarrow$  i

maxSum  $\leftarrow$  tempSum

end if

if tempSum < 0

localSt  $\leftarrow$  i+1

tempSum  $\leftarrow$  0

end if

end

vector<int> res  $\leftarrow$  {st, end, maxSum}

return res

# DYNAMIC PROGRAMMING PSEUDO CODE

```
function MAIN()
```

```
    r ← 0, c ← 0, maxSum ← INT_MIN
```

```
    x ← 0, y ← 0, z ← 0, w ← 0
```

```
    input r, c
```

```
    for i ← 0 to r-1
```

```
        for j ← 0 to c-1
```

```
            input A[i][j]
```

```
        end
```

```
    end
```

```
    for i ← 0 to r-1
```

```
        vector<int> sum(c)
```

```
        for j ← 0 to r-1
```

```
            for col ← 0 to c-1
```

```
                sum ← sum + A[j][col]
```

```
            end
```

```
        vector<int> res ← KADANE(sum)
```

```
        if maxSum < res[2]
```

```
            x ← i
```

```
            y ← res[0]
```

```
            z ← j
```

```
            w ← res[1]
```

```
            maxSum = res[2]
```

```
        end if
```

```
    end
```

```
end
```

```
print x, y, z, w, maxSum
```

# ALGORITHM ANALYSIS

## Brute force(Naive method)

In the brute force approach we try to check every possible rectangle in the given  $n \times m$  2D array (where  $n, m$  are number of rows and columns respectively).

This solution requires 6 nested loops :

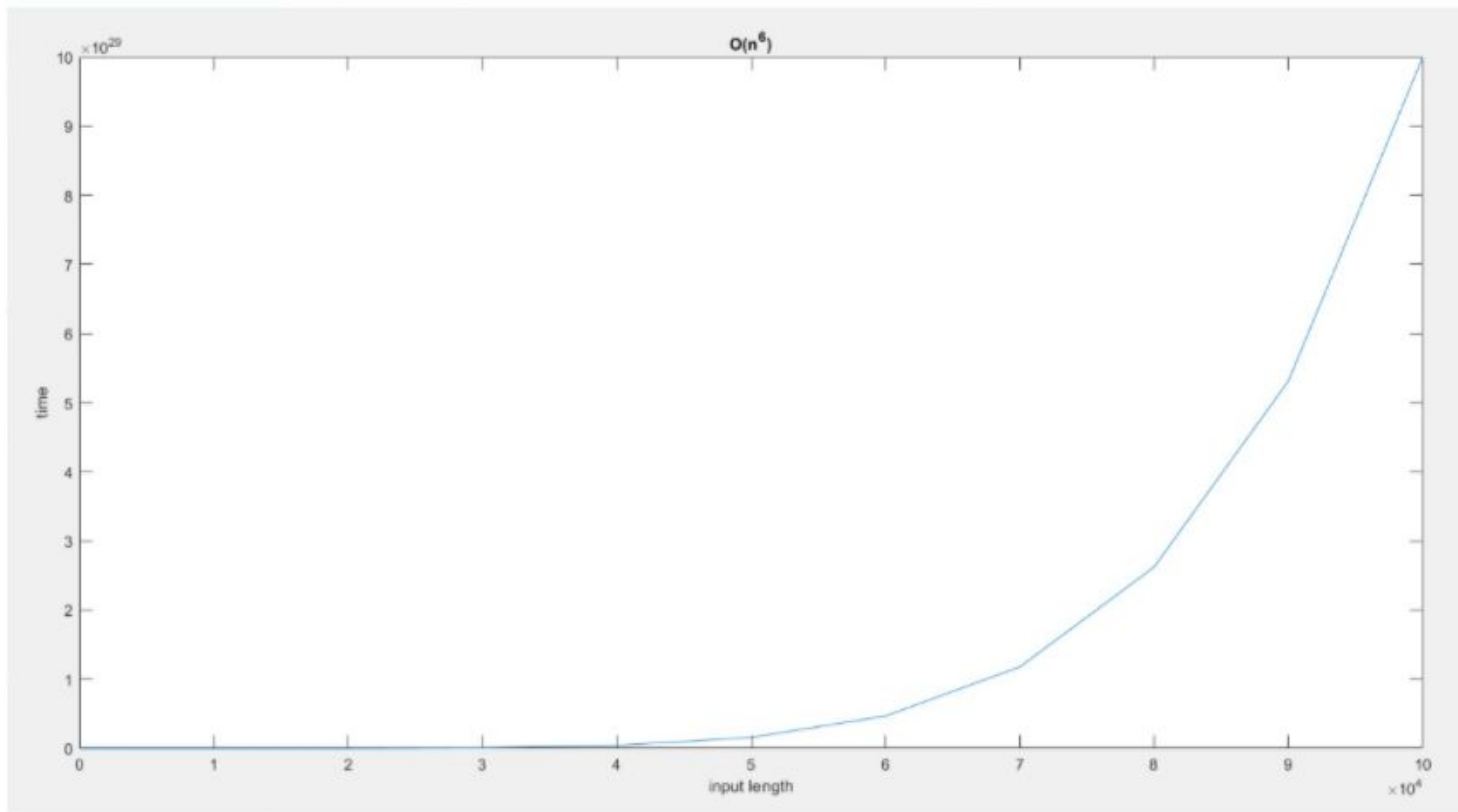
- ▶ 4 for start and end coordinate of the 2 axis  $O(n^2 \times m^2)$
- ▶ 2 for the summation of the sub-matrix  $O(n \times m)$
- ▶ The overall time complexity is  **$O(n^3 \times m^3)$**
- ▶ As we did not allocate extra space in this approach, space complexity is  $O(1)$ .

# ALGORITHM ANALYSIS

## DYNAMIC PROGRAMMING APPROACH

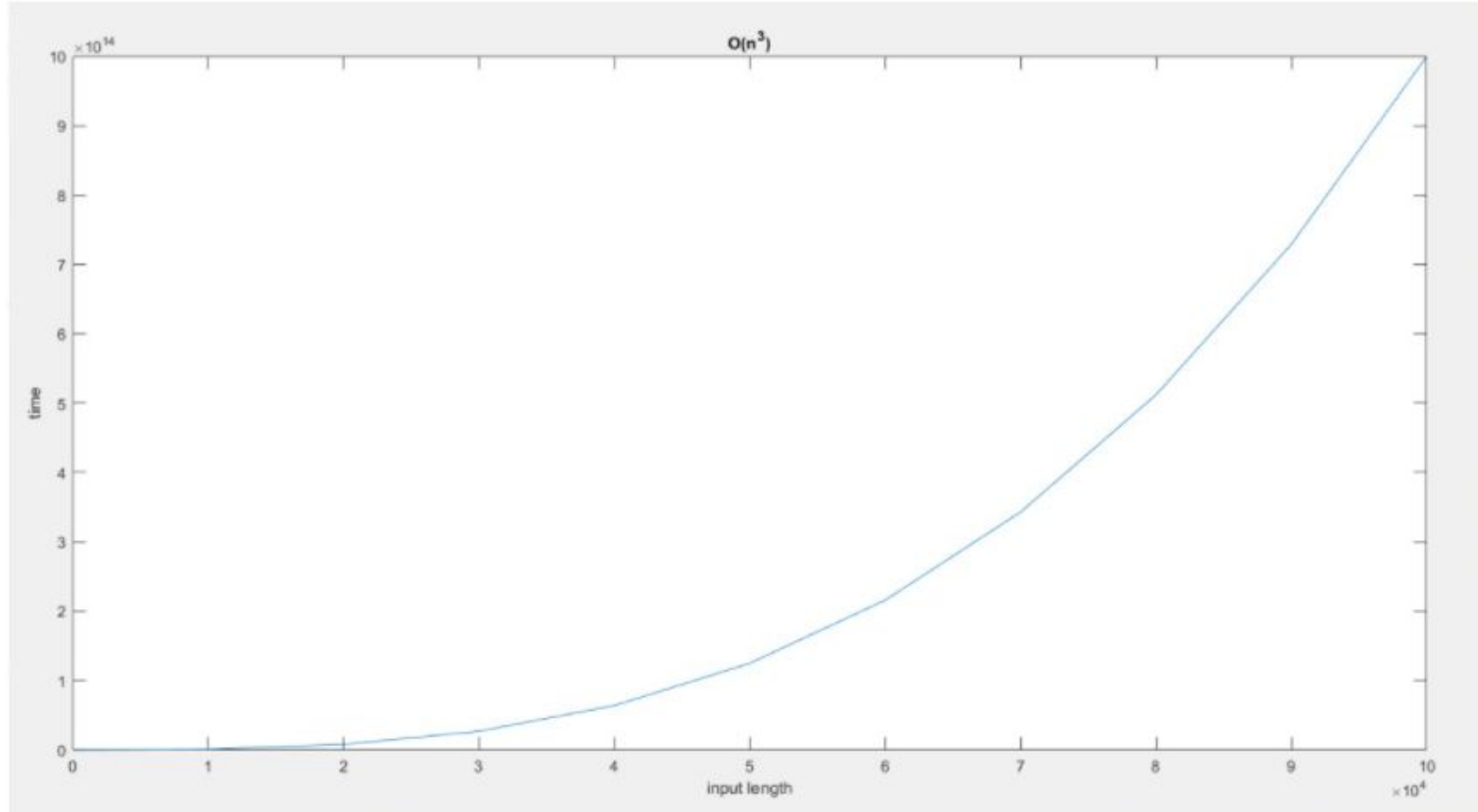
- ▶ Basically, kadane's dynamic algorithm (complexity:  $O(m)$  since we are doing it for column size) is used inside two loops of size  $n$  (complexity:  $O(n^2)$ ).
- ▶ This gives a total complexity of  $O(n^2 \times m)$
- ▶ In this approach we have used temp array, where it's size is equal to number of columns. Therefore space complexity is  $O(m)$

# EXPERIMENTAL STUDY (NAIVE APPROACH)





# EXPERIMENTAL STUDY(DP APPROACH)



## CONCLUSION

From the experimental study we concluded that the average running time of dynamic approach using kadane's algorithm is best, which can be observed from the mutual graph of kadane's dynamic programming algorithm and Brute Force algorithm as shown.

# IMPORTANT LINKS

Reference links:-

<https://www.geeksforgeeks.org/maximum-sum-rectangle-in-a-2d-matrix-dp-27/>

<https://www.geeksforgeeks.org/largest-sum-contiguous-subarray/>

Code link:-

Kadane's Dynamic programming approach: <https://ideone.com/Zw0Ape>

Brute force: <https://ideone.com/W0rVpg>

THANK YOU