## Reflection & Coxeter Groups

Shubkarman Walia

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### Overview

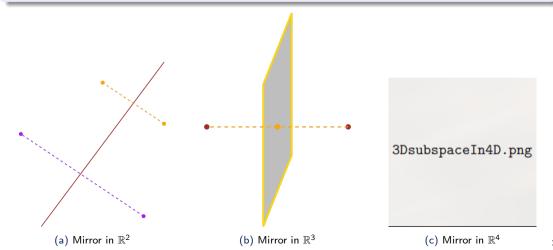
1. Mirrors and Reflections

2. Coxeter Groups

### Hyperplanes and Mirrors

### Hyperplane

In an n-dimensional vector space, an n-1 dimensional subspace is called a hyperplane.



### A precise definition

#### Reflections in a Vector space

Reflections are orthogonal transformations. Hence they preserve distance and angles. So if s is a reflection and  $\alpha, \beta \in \mathbb{R}^n$  then,  $||s(\alpha)|| = ||\alpha||$  and  $s(\alpha).s(\beta) = \alpha.\beta$ 

#### More Precisely

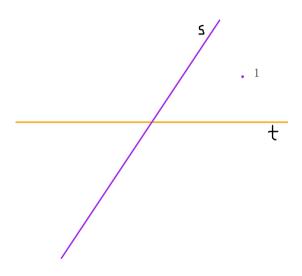
Reflection in a Real Euclidean Vector space endowed with a positive definite symmetric bilinear form  $(\lambda,\mu)$  is a linear operator that sends some non-zero vector  $\alpha$  (normal to the hyperplane  $H_{\alpha}$ ) to its negative while fixing everything in the hyperplane  $H_{\alpha}$ .

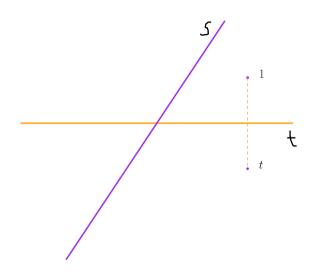
### Reflection as an orthogonal transformation

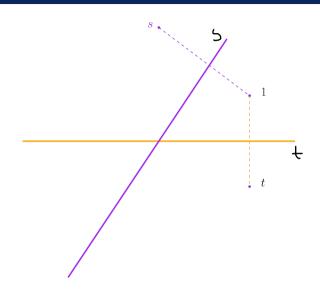
#### **Fact**

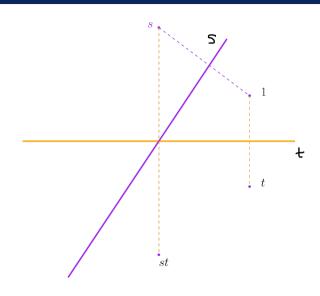
The set of all orthogonal transformations of a vector space V forms a group, O(V).

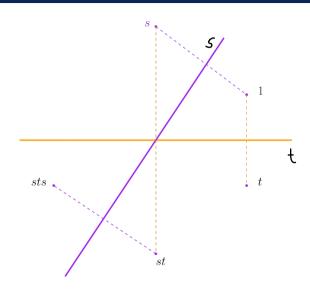
And so, a finite group generated by a system of mirrors (multiple reflections) is just a subgroup of O(V).

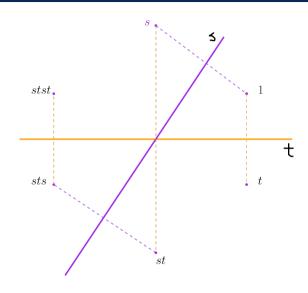


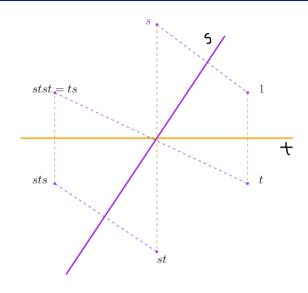


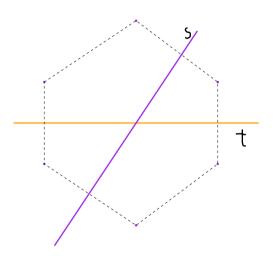












#### Conclusion

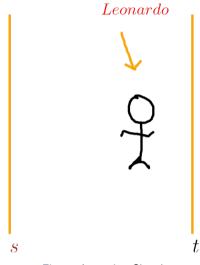
- $s^2 = t^2 = 1$  for any reflection s, t
- We get only 6 (finite) number of "images" or elements of the group.
- We get this mysterious relation: sts = tst
- We get the group:  $\langle s, t \mid s^2 = t^2 = 1, sts = tst \rangle$

## Another Example



Figure:  $\tilde{\mathcal{A}_2}$  in Inception

## Another Example



### How to find all the Finite Configurations?

#### Theorem

The group of all isometries of  $\mathbb{AR}^n$  which fix a point o coincides with the orthogonal group  $\mathbb{O}_n$ .

#### Consequence

The finite groups generated by reflections, i.e, Finite Reflection Groups, fix a point.

And so, all the hyperplanes in a system of mirrors must have a point in common in order to generate finite images.

### Enter: Coxeter Systems

#### Coxeter system

A Coxeter system is a group generated by a finite set of generators

$$S = \{s_1, s_2, ..., s_n\}$$

under the relations:

$$s_i^2 = 1$$
 for all  $1 \le i \le n$ 

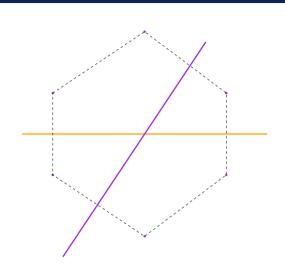
$$\underbrace{s_i s_j s_i \cdots}_{m_{ij} \text{ terms}} = \underbrace{s_i s_j s_i \cdots}_{m_{ij} \text{ terms}} \text{ for all } 1 \leq i, j \leq n$$

### An Example

$$S = \{s, t\}$$
 with  $m_{st} = 3$   
 $\implies W = \langle s, t | s^2 = t^2 = 1, sts = tst \rangle$ 

## An Example

And we know W =



### Equivalence

#### The Connection

Each finite reflection group can be represented as a coxeter system.

And so we have turned a geometrical problem to a combinatorial one!

### Coxeter Graphs

#### Representation of Coxeter systems

Each Coxeter system can be represented as a graph where each vertex is a generator of the corresponding Coxeter group, and there's an edge between vertices if  $m_{ij} \ge 3$  with edge label=  $m_{ii}$ .

### A combinatorial approach

We saw before how  $\langle s, t \rangle$  with  $m_{st} = 3$  was the finite group  $D_6$  What about?

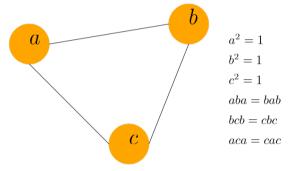


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Notice,  $(abc)^n \in W \ \forall n \in \mathbb{N}$  The group generated is infinite!

### A more elegant way

Let  $S = \{s_1, s_2, ..., s_n\}$  be the generator set of our Coxeter system. Encode the Coxeter graph into a Coxeter matrix, say A, known as the Cartan Matrix such that

$$(a_{ij})=(a_{ji})=-cos(\pi/m_{ij}) ext{ for } s_i,s_j\in S$$

#### Note

A is a symmetric matrix so all of its eigenvalues  $\in \mathbb{R}$ 

#### Theorem

If A is positive definite, i.e all eigenvalues of A are strictly positive then the Coxeter system associated to A is finite.

If A is positive semi-definite, i.e all eigenvalues of A are non-negative (zeroes are allowed) then the associated Coxeter System is affine.

### Computing

All that remains is to see all possible Coxeter systems that are possible and compute its eigenvalues...

#### 2.7 Classification of graphs of positive type

**Theorem** The graphs in Figure 1 of 2.4 and Figure 2 of 2.5 are the only connected Coxeter graphs of positive type.

*Proof.* Suppose there were a connected Coxeter graph  $\Gamma$  of positive type not pictured in either Figure 1 or Figure 2. We proceed in 20 easy steps

Figure: Humphreys: pg 36

### 20 Easy Steps

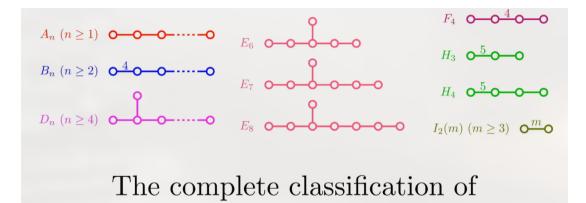
- All Coxeter graphs of rank 1 or 2 are clearly of positive type (A<sub>1</sub>, I<sub>2</sub>(m), A

  1, so we must have n ≥ 3.
- (2) Since  $\widetilde{A}_1$  cannot be a subgraph of  $\Gamma$ , we must have  $m < \infty$ .
- (3) Since An(n ≥ 2) cannot be a subgraph of Γ, Γ contains no circuits. Suppose for the moment that m = 3.
- (4)  $\Gamma$  must have a branch point, since  $\Gamma \neq A_n$ .
- (5)  $\Gamma$  contains no  $\widetilde{D_n}$ , n > 4, so it has a unique branch point.
- (6)  $\Gamma$  does not contain  $\widetilde{D_4}$ , so exactly three edges meet at the branch point (with  $a \le b \le c$  further vertices lying in these three directions).
- (7) Since  $\widetilde{\mathbf{E}_6}$  is not a subgraph of  $\Gamma$ , a=1.
- (8) Since  $\widetilde{E_7}$  is not a subgraph of  $\Gamma$ ,  $b \leq 2$ .
- (9) Since  $\Gamma \neq D_n$ , b cannot be 1, so b = 2.
- (10) Since  $\widetilde{E}_8$  is not a subgraph of  $\Gamma$ ,  $c \le 4$ .
- (11) Since  $\Gamma \neq E_n, E_7, E_8$ , the case m = 3 is impossible. Thus m > 4.

- (14) Since  $\Gamma \neq B_n$ , the two extreme edges of  $\Gamma$  are labelled 3.
- (15) Since  $\Gamma$  does not contain  $\widetilde{F}_4$ , n must be 4.
- (16) But  $\Gamma \neq F_4$ , so the case m = 4 is impossible. Thus  $m \geq 5$ .
- (17) Since  $\Gamma$  does not contain  $\widetilde{G}_2$ , we must have m = 5.
- (18)  $\Gamma$  does not contain the nonpositive graph Z<sub>4</sub> in 2.5, so the edge labelled 5 must be an extreme edge.
- (19)  $\Gamma$  does not contain the nonpositive graph  $Z_5$ , so  $n \leq 4$ .
- (20) Now Γ must be either H<sub>3</sub> or H<sub>4</sub>, which is absurd. So we have eliminated all possibilities. □

Figure: Humphreys: 20 Easy steps

### Finite Classification



finite irreducible Coxeter systems

Figure: All possible ways you can orient mirrors in n-dimensions so that you get finite images. Credits for the Image: Joseph Newton

#### References

- Alexandre V. Borovik, & Anna Borovik. (2010). *Mirrors and Reflections: The Geometry of Finite Reflection Groups*. Springer.
- James E. Humphreys. (1990). *Reflection Groups and Coxeter Groups*. Cambridge University Press.
- Newton, J. [Joseph Newton]. (n.d.). The Coxeter Classification [Video]. YouTube. https://www.youtube.com/watch?v=BV5mYjh8m4E&t=402s

# Thank you!