Master of Computer Applications

CAPOL403R01: Computer Organization & Architecture

Unit III: Lecture 3 – Part 2
Associative & Set associative mappings

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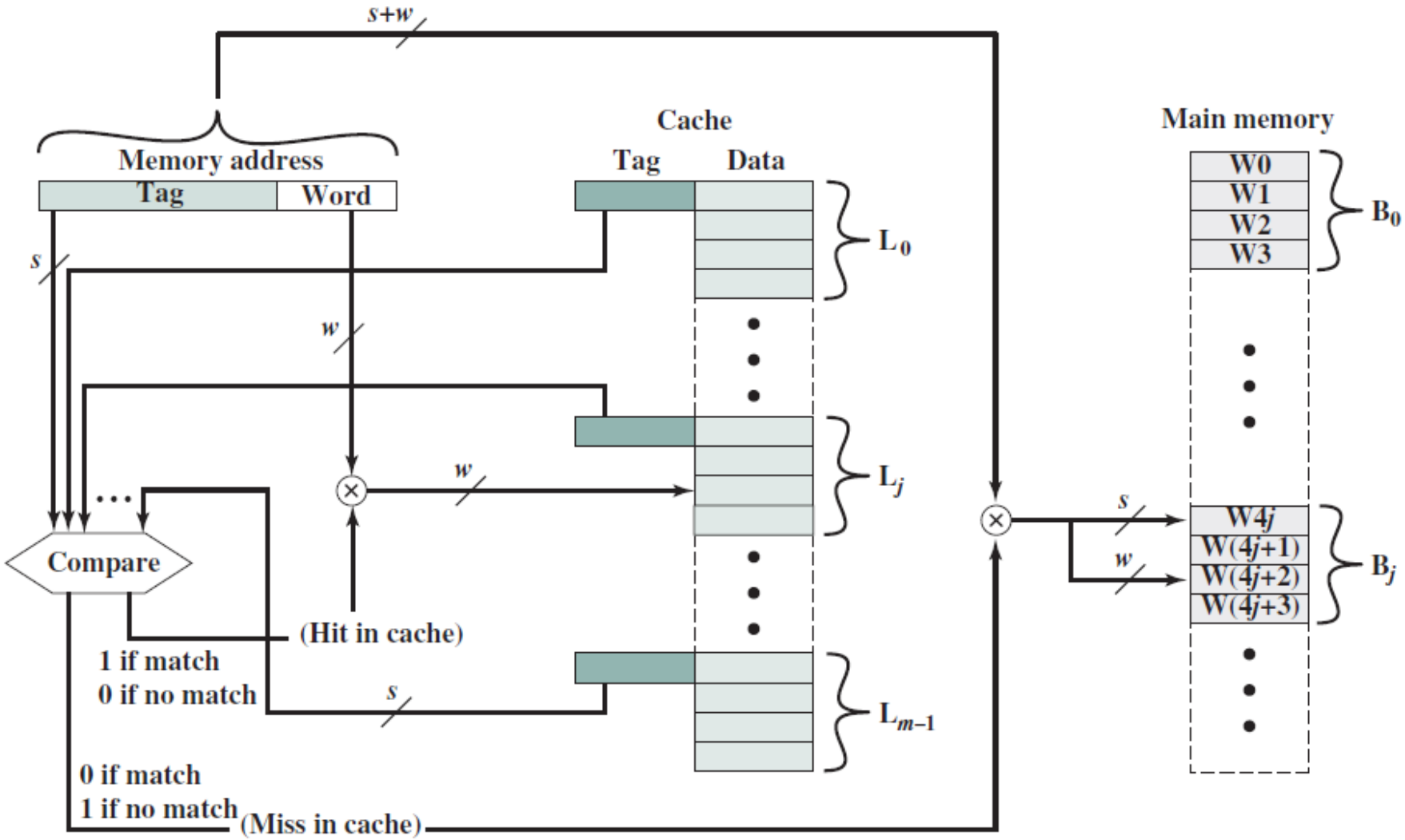
Associative mapping

- In this mapping, any main memory block is mapped on any cache line
- In this way it eliminates the “thrashing” problem
- But, to recognize the cache miss, all tags have to be checked
- To improve the performance, all tags are checked in parallel
  - It makes the design complex compared to direct mapping
- The address format is consider as given below

- A cache line has “k” words and “w” bits are allocated for word field
- When the address has “n” bits, “n-w” bits are allocated for tag field
  - The cache has any number of cache lines (2^L for some ‘L’ is not mandatory)

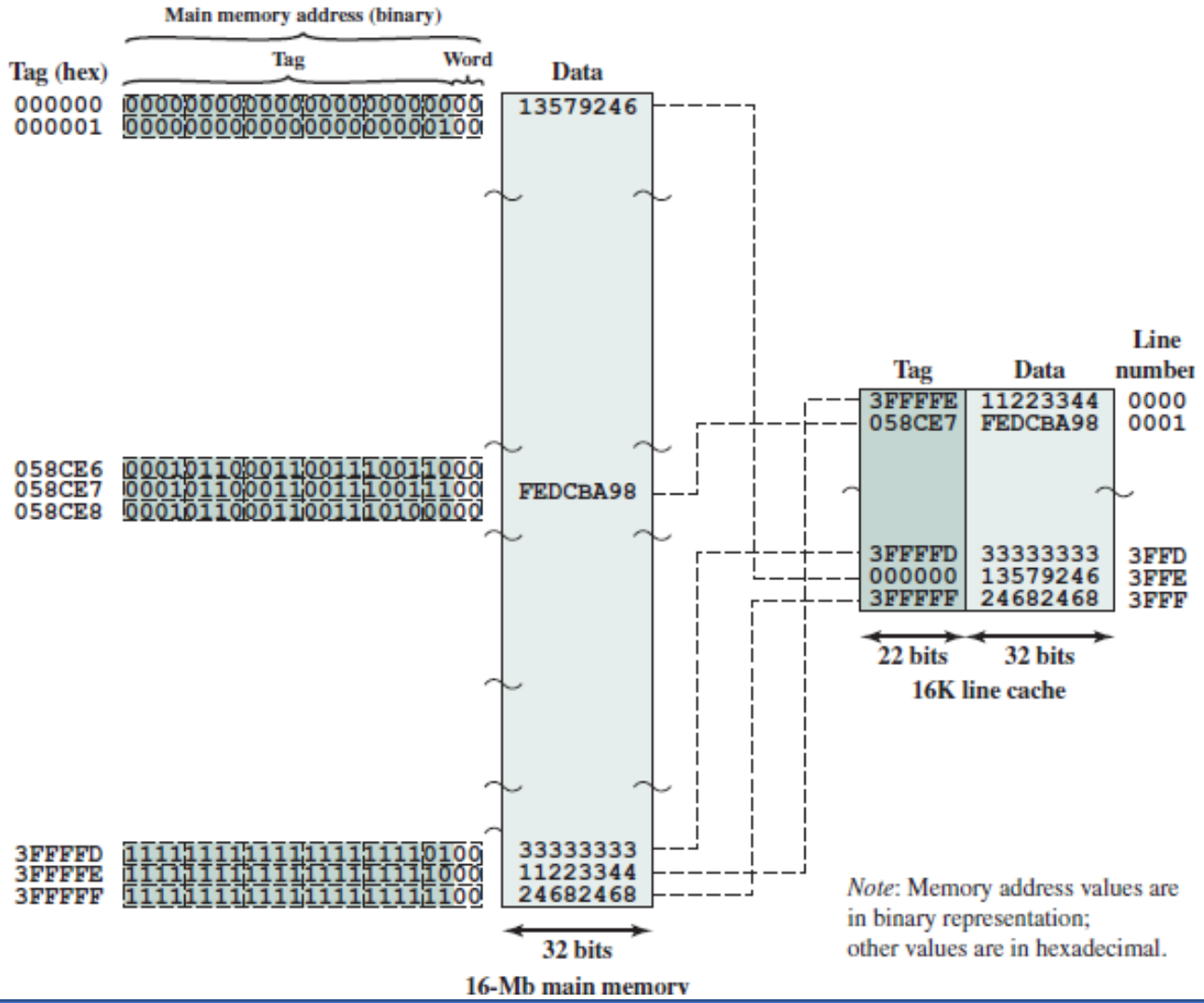
Associative mapping



A numerical example

- An associative cache can hold 64 KB words. Data are transferred between main memory and cache in blocks of 4 bytes each. The main memory size is 16 MB. Explain the address formats and mapping diagram
- Main memory size is 16 MB
  - Number of bits allocated for address is: $\log_2(16M)=\log_2(2^{24}) = 24$
- Block size is 4 bytes
  - Number of blocks in main memory = 16MB/4B = 4M = 2^{22} Blocks = 4 M Blocks
  - 22 bits are needed to identify a block – Tag size has to be 22 bits
  - Line size is 4 bytes – 2 bits are allocated to select a word from a line
  - Number of bits allocated for Tag field = (24 – 2) = 22
- Cache size is 64 KB = 2^{16} Bytes
  - Number of cache lines = $2^{16}/ 2^2 = 2^{14}$ = 16 K Lines
  - Any main memory block (among 2^{22}) is mapped any cache line (among 2^{14})

Fully associative mapping example



Set Associative mapping

- It has the strengths of direct and fully associative mapping techniques
- The cache has a number of sets
  - The number ‘v’ can be represented as 2^d for some positive integer ‘d’
- Each set has ‘k’ lines
  - K may be any positive number
- Size of cache is : $v*k*2^w = k*2^d*2^w = k*2^{d+w}$
  - 2^w is the number of words per line
- A main memory block is mapped to any line of “a particular set” only
  - Set number = block number % k
  - For example when k=8, the block number 32 is mapped only on **set 0**

    Reason: 0= 32%8
  - Any line from the set will be chosen

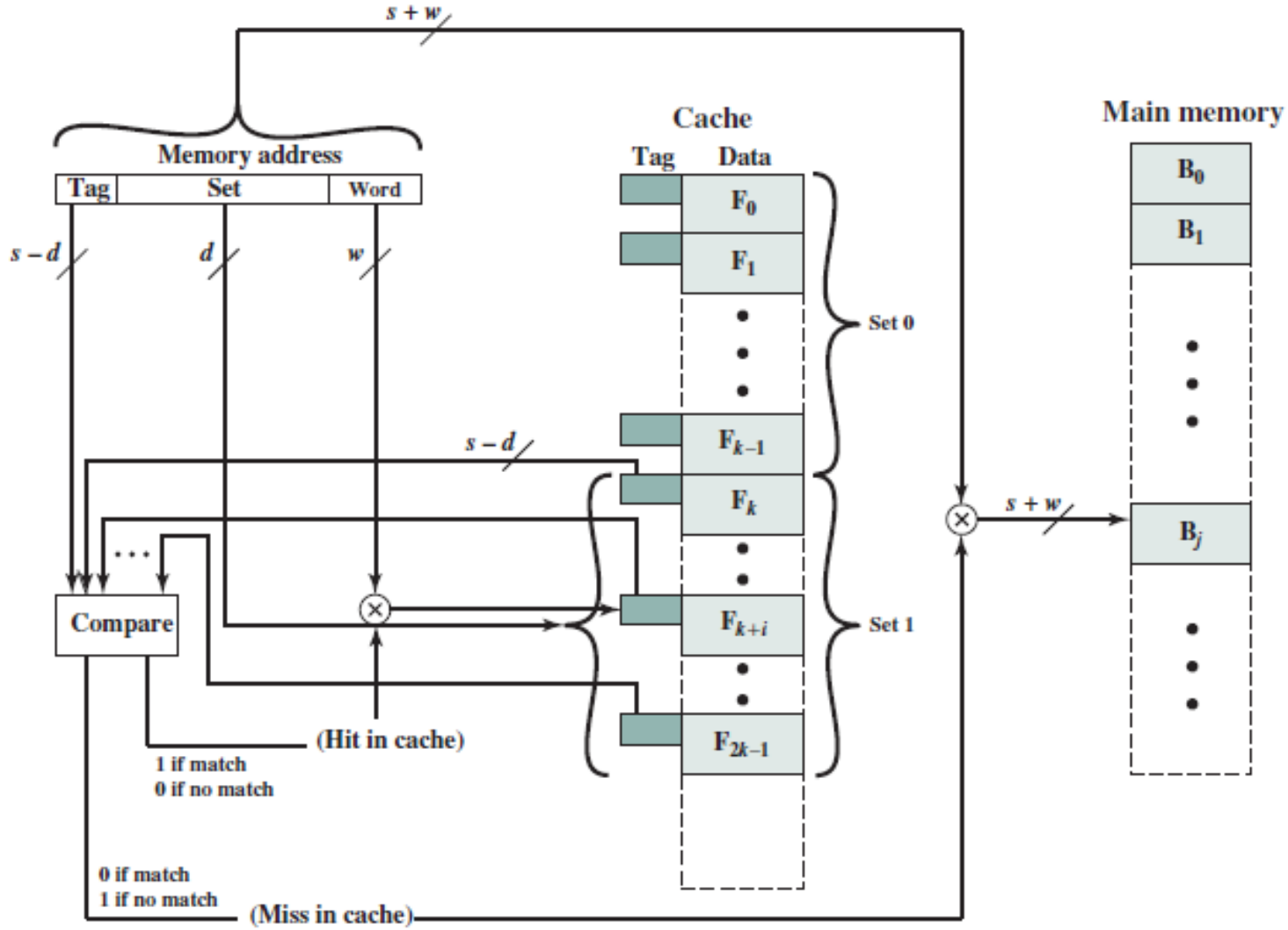
Set Associative mapping…

- A k-way set associative mapping has ‘k’ lines in a single set
- A k-way set associative mapping has (m/k) sets when the cache has ‘m’ lines
  - K has to be the divisor of m
- When k=1, it is direct mapping
- When k=m, it is fully-associative mapping
- A 2-way set associative mapping has better performance than direct mapping and the complexity is minimized
- Commonly used k values are:2 and 4
- Address format

Set Associative mapping…

- The address is divided into the above fields
- The appropriate set which is indicated by the bits of set field is chosen
- All tag bits of the lines of that set will be compared in parallel
- If there is a match in tag, the word field is used to access the word
- If not, then the word with the block is moved from main memory to any one lines of that set using replacement algorithms

K-way set associative cache



A numerical example

- A 2-way set associative cache can hold 64 KB words. Data are transferred between main memory and cache in blocks of 4 bytes each. The main memory size is 16 MB. Explain the address formats and mapping diagram
- Main memory size is 16 MB
  - Number of bits allocated for address is: $\log_2(16M)=\log_2(2^{24}) = 24$
- Block size is 4 bytes
  - Line size is 4 bytes – 2 bits are allocated to select a word from a line
- Cache size is 64 KB = 2^{16} Bytes
  - Number of cache lines = $2^{16}/ 2^2 = 2^{14}$
- Two-way set associative cache
  - Number of sets = $2^{14} / 2 = 2^{13}$
  - Number of bits allocated for a set = 13
  - Number of bits allocated for tag = 24 – (13+2) = 24 – 15 = 9

Set Associative mapping…

Tag (9 bits)	Set (13 bits)	Word (2 bits)

- For example, the address “0101_1010_1111_1000_1010_1110” is grouped as

0101_1010_1	1111_0001_0101_1	10

- The word is available in any one of the two cache lines of set number 1_1110_0010_1011 or equivalently on $(1E2B)_h$
- When any one of the line has the tag 0_1011_0101 or $(0B5)_h$, then cache hit
  - The third word (with word number 10 or $(2)_h$) of that line is accessed.
- If no lines has the tag, then the block 01_0110_1011_1110_0010_1011 or $(16BE2B)_h$ is copied from main memory and placed in anyone of the two cache lines of set $(1E2B)_h$

Thank you