

# CIE6020/MAT3350 HW Assignment 2

Due: 23:55, 18 Feb 2019

- Let  $X, Y, Z$  be three random variables with a joint probability mass function  $p(x, y, z)$ . The relative entropy between the joint distribution and the product of the marginals is

$$D(p(x, y, z) || p(x)p(y)p(z)) = \mathbb{E} \left[ \log \frac{p(x, y, z)}{p(x)p(y)p(z)} \right].$$

Expand this in terms of entropies. When is this quantity zero?

**Solution:**

$$\begin{aligned} D(p(x, y, z) || p(x)p(y)p(z)) &= \sum_{x,y,z} p(x, y, z) \log \frac{p(x, y, z)}{p(x)p(y)p(z)} \\ &= \sum_{x,y,z} p(x, y, z) \log p(x, y, z) \\ &\quad - \sum_x p(x) \log p(x) - \sum_y p(y) \log p(y) \\ &\quad - \sum_z p(z) \log p(z) \\ &= H(X) + H(Y) + H(Z) - H(X, Y, Z). \end{aligned}$$

Clearly, if we have  $p(x, y, z) = p(x)p(y)p(z)$ , the quantity is zero.

- Let the random variable  $X$  have three possible outcomes  $\{a, b, c\}$ . Consider two distributions on this random variable:

symbol	$p(x)$	$q(x)$
$a$	$\frac{1}{2}$	$\frac{1}{3}$
$b$	$\frac{1}{4}$	$\frac{1}{3}$
$c$	$\frac{1}{4}$	$\frac{1}{3}$

Calculate  $H(p)$ ,  $H(q)$ ,  $D(p||q)$  and  $D(q||p)$ . Verify that in this case,  $D(p||q) \neq D(q||p)$ .

**Solution:**

$$H(p) = -\frac{1}{2} \log \frac{1}{2} - 2 \times \frac{1}{4} \log \frac{1}{4} = \frac{3}{2}.$$

$$H(q) = -3 \times \frac{1}{3} \log \frac{1}{3} = \log 3.$$

$$D(p||q) = \frac{1}{2} \log \frac{1/2}{1/3} + 2 \frac{1}{4} \log \frac{1/4}{1/3} = \log 3 - \frac{3}{2}.$$

$$D(q||p) = \frac{1}{3} \log \frac{1/3}{1/2} + 2 \frac{1}{3} \log \frac{1/3}{1/4} = \frac{5}{3} - \log 3.$$

We see that  $D(p||q) = \log 3 - \frac{3}{2} \neq \frac{5}{3} - \log 3 = D(q||p)$ .

3. Show that  $\ln x \geq 1 - \frac{1}{x}$  for  $x > 0$ , where the equality holds when  $x = 1$ .

**Solution:** Let  $f(x) = \ln x - 1 + \frac{1}{x}$ . The derivative  $f'(x) = \frac{1}{x} - \frac{1}{x^2}$ . We have

- $f(1) = 0$ ;
- when  $x \in (0, 1)$ ,  $f'(x) < 0$ ;
- when  $x > 1$ ,  $f'(x) > 0$ .

Therefore, for  $x > 0$ ,  $f(x) \geq 0$ , i.e.,  $\ln x \geq 1 - \frac{1}{x}$ , where the equality holds when  $x = 1$ .

4. *Conditioning reduces entropy.* Show that  $H(Y|X) \leq H(Y)$  with equality iff  $X$  and  $Y$  are independent.

**Solution:**  $H(Y) - H(Y|X) = I(X;Y) \geq 0$  where the equality holds when  $X$  and  $Y$  are independent.

5. Show that  $I(X;Y|Z) \geq 0$  with equality iff  $X \rightarrow Z \rightarrow Y$ .

**Solution:**  $I(X;Y|Z) = \sum_z p(z) I(X;Y|Z = z) \geq 0$  as mutual information is nonnegative, where the equality holds iff for all  $z$ ,  $I(X;Y|Z = z) = 0$ , i.e., for all  $x, y$  and  $z$ ,  $p(x, y|z) = p(x|z)p(y|z)$ .

6. *Data processing.* Let  $X_1 \rightarrow X_2 \rightarrow X_3 \rightarrow \dots \rightarrow X_n$  form a Markov chain, i.e.,

$$p(x_1, x_2, \dots, x_n) = p(x_1)p(x_2|x_1) \cdots p(x_n|x_{n-1}).$$

Reduce  $I(X_1; X_2, \dots, X_n)$  to its simplest form.

**Solution:** First, the Markov chain implies  $p(x_2, \dots, x_n) = p(x_2)p(x_3|x_2) \cdots p(x_n|x_{n-1})$ . Then,

$$\begin{aligned}
 I(X_1; X_2, \dots, X_n) &= \sum_{x_1, \dots, x_n} p(x_1, \dots, x_n) \log \frac{p(x_1, \dots, x_n)}{p(x_1)p(x_2, \dots, x_n)} \\
 &= \sum_{x_1, \dots, x_n} p(x_1, \dots, x_n) \log \frac{p(x_1)p(x_2|x_1) \cdots p(x_n|x_{n-1})}{p(x_1)p(x_2)p(x_3|x_2) \cdots p(x_n|x_{n-1})} \\
 &= \sum_{x_1, x_2} p(x_1, x_2) \log \frac{p(x_2|x_1)}{p(x_1)} \\
 &= I(X_1, X_2).
 \end{aligned}$$

7. Let  $X$  and  $Y$  be two random variables and let  $Z$  be independent of  $(X, Y)$ . Show that  $I(X; Y) \geq I(X; g(Y, Z))$  for any function  $g$ .
8. *Bottleneck.* Suppose that a (nonstationary) Markov chain starts in one of  $n$  states, necks down to  $k < n$  states, and then fans back to  $m > k$  states. Thus,  $X_1 \rightarrow X_2 \rightarrow X_3$ , that is,  $p(x_1, x_2, x_3) = p(x_1)p(x_2|x_1)p(x_3|x_2)$ , for all  $x_1 \in \{1, 2, \dots, n\}, x_2 \in \{1, 2, \dots, k\}, x_3 \in \{1, 2, \dots, m\}$ .
  - (a) Show that the dependence of  $X_1$  and  $X_3$  is limited by the bottleneck by proving that  $I(X_1; X_3) \leq \log k$ .
  - (b) Evaluate  $I(X_1; X_3)$  for  $k = 1$ , and conclude that no dependence can survive such a bottleneck.

**Solution:**

- (a) Since  $X_1 \rightarrow X_2 \rightarrow X_3$ , by the data processing inequality,  $I(X_1; X_3) \leq I(X_1; X_2) \leq H(X_2) \leq \log k$ .
- (b) By (a), we have  $I(X_1, X_3) \leq \log 1 = 0$ . Since mutual information is non-negative, we have  $I(X_1, X_3) = 0$ , i.e.,  $X_1$  and  $X_3$  are independent.