Lecture 21: Algorithms for Unconstrained Optimization

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Announcements

- ► Homework 7 due today
- ▶ Homework 8 posted, due next Wednesday (11/28)

Recap: Convexity/Convex Optimization Problem

We defined convex functions and convex optimization problems

- Can identify whether a given function is convex, concave or neither
- Can identify whether an optimization problem is a convex optimization problem
- Can convert problems into convex optimization problems

Implications of Convexity

Theorem

For a convex optimization problem, any local minimizer is also a global minimizer.

Theorem

For convex optimization problems, KKT conditions are sufficient for global optimality.



Recap: Duality for Nonlinear Optimization

There are dual problems for nonlinear optimization problems

- ► The dual problems are derived by considering the Lagrangian function and switching the order of max and min.
- Dual problems for nonlinear optimization do not necessarily have an explicit form.

Duality theorems:

- Weak duality always holds
- Strong duality holds when the problem is convex and the Slater's condition holds (otherwise, there could be a positive duality gap)

Upcoming Agenda

Discuss how to solve nonlinear optimization problems.

- ► We have shown that in many cases, KKT condition can be used to solve the optimization problem
- However, those are ad hoc situations. In most cases, we cannot directly find the optimal solution from the KKT conditions
- ▶ We want to have a robust procedure (an algorithm) that guarantees to solve the optimization problem.

Unconstrained Problems

We start with the unconstrained problem:

minimize_x
$$f(x)$$

We are going to study the following methods:

- Bisection search
- Golden section search
- Gradient descent method
- Newton's method

General Solution Idea

Typically, optimization algorithms are *iterative* procedures.

- ▶ Start from some point \mathbf{x}_0 , then generate a sequence of $\{\mathbf{x}_k\}$
- ► The sequence terminates when either no progress can be made or when we know that the current solution is already satisfactory
- ▶ Typically, we want to have $f(\mathbf{x}_{k+1}) < f(\mathbf{x}_k)$, i.e., each step we can improve the objective value.
- And hopefully, we want the sequence $\{x_k\}$ to *converge* to a local minimizer x^* (or global minimizer).

Recall the only algorithms we have studied: the simplex method and the interior point method. They both follow the above paradigm.



Some Useful Concepts: Convergent Sequences

Definition

Let $\{\mathbf{x}_k\}$ be a sequence of real vectors. Then $\{\mathbf{x}_k\}$ converges to \mathbf{x}^* if and only if for all real numbers $\epsilon>0$, there exists a positive integer K such that $||\mathbf{x}_k-\mathbf{x}^*||<\epsilon$ for all $k\geq K$.

In all our discussions, we assume $||\mathbf{x}||$ is the 2-norm of $\mathbf{x} = (x_1, ..., x_n)$, which means:

$$||\mathbf{x}|| = \sqrt{x_1^2 + x_2^2 + \dots + x_n^2}$$

Example of convergence:

- $> x_k = 1/k \rightarrow 0$
- $x_k = (1/2)^k \to 0$



Single Variable Problem

Assume f(x) is a single variable function.

Objective: find a local minimizer of f(x).

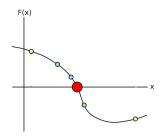
We introduce two methods:

- Bisection method
- Golden section method

Bisection Method

Bisection method uses the idea that the local minimizer must satisfy the FONC: f'(x) = 0.

Therefore, the problem becomes a root-finding problem for g(x) = f'(x).



Root Finding Algorithm: Bisection Method

Assume one can find x_l and x_r such that $g(x_l) < 0$ and $g(x_r) > 0$. By intermediate value theorem, if $g(\cdot)$ is continuous, there must exist a root of $g(\cdot)$ in $[x_l, x_r]$.

Bisection method:

- 1. Define $x_m = \frac{x_l + x_r}{2}$
- 2. If $g(x_m) = 0$, then output x_m
- 3. Otherwise
 - If $g(x_m) > 0$, then let $x_r = x_m$
 - If $g(x_m) < 0$, then let $x_l = x_m$
- 4. If $|x_r-x_l|<\epsilon$. stop and output $\frac{x_l+x_r}{2}$, otherwise go back to Step 1

One can also set the stop criterion based on g(x)



Bisection Method

In the bisection method, each iteration will divide the search interval to half.

Therefore, to find an ϵ approximation of x^* , we need at most $\log_2 \frac{x_r - x_l}{\epsilon}$ iterations

Applying bisection method to $f'(\cdot)$, one can find a point satisfying FONC (approximately), and if $f(\cdot)$ is convex, we can find the global minimizer of f(x) (approximately)

▶ Although simple, the bisection method is very useful in practice because it is easy to implement.

Example: Use bisection method to maximize

$$f(x) = \frac{xe^{-x}}{1 + e^{-x}}$$



Golden Section Method

One drawback of using the bisection method to solve (single variable, unconstrained) optimization problems is that it requires the knowledge (and computation) of f'(x).

Sometimes, we don't have f'(x) available. For example, f(x) sometimes is only a *black box*, which does not admit an analytical form (thus the derivative is hard to compute)

However, if we know that f(x) has a unique local minimum x^* in the range $[x_I, x_r]$, then we still have very efficient way to find it out

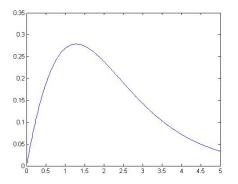
- ▶ We call such f unimodal on $[x_l, x_r]$
- Unimodal function has the property that the local minimum is global minimum (convex function is always unimodal)

In the context of maximization, we say a function is unimodal if it has a unique local maximum



Example of Unimodal Function

Consider $f(x) = \frac{xe^{-x}}{1+e^{-x}}$ (the function in last homework set):



This is a unimodal function, but not a concave function.



Golden Section Method

Assume we start with $[x_l, x_r]$. Assume $0 < \phi < 0.5$.

- 1. Set $x_I' = \phi x_r + (1 \phi)x_I$ and $x_r' = (1 \phi)x_r + \phi x_I$.
- 2. If $f(x'_l) < f(x'_r)$, then the minimizer must lie in $[x_l, x'_r]$, so set $x_r = x'_r$
- 3. Otherwise, the minimizer must lie in $[x'_l, x_r]$, so set $x_l = x'_l$.
- 4. If $x_r x_l < \epsilon$, output $\frac{x_l + x_r}{2}$, otherwise go back to Step 1

If we want to reuse the computation in the last step, we want to set

$$\frac{1-2\phi}{1-\phi} = \phi$$

i.e., $\phi=\frac{3-\sqrt{5}}{2},$ and $1-\phi=\frac{\sqrt{5}-1}{2}=0.618$ (where the name comes from)



For Maximization Problem

Both the bisection and golden section method can be easily adapted for maximization problem.

Just change the way the comparison works (or just change the interval to be cut off in each iteration).

Example Revisited: Use Golden section method to maximize:

$$f(x) = \frac{xe^{-x}}{1 + e^{-x}}$$

Higher Dimensional Problems

Next we consider the high-dimensional problem

minimize_x
$$f(x)$$

▶ There is not a clear bisection or golden section in that case

Solution idea:

- ▶ Each time, we first find a search direction
- ► Then we search for a good solution along that direction (which reduces to a one-dimensional problem)

General Framework for High Dimensional Search

From \mathbf{x}^0 , we generate a sequence of points:

$$\mathbf{x}^{k+1} = \mathbf{x}^k + \alpha_k \mathbf{d}^k.$$

We call \mathbf{d}^k the search direction (a vector) and α_k the step size (a scalar).

- ▶ Then the key is to choose proper \mathbf{d}^k at each iteration.
- $ightharpoonup \mathbf{d}^k$ typically depends on \mathbf{x}^k
- ▶ Then α_k may be chosen in accordance with some line (one-dimension) search rules.

We will study two such methods:

Gradient descent method and Newton's method



Gradient Descent Method

In the following discussion, we assume that $f(\mathbf{x})$ is continuously differentiable (differentiable and the derivative is continuous).

We know by Taylor expansion, for small α

$$f(\mathbf{x} + \alpha \mathbf{d}) \approx f(\mathbf{x}) + \alpha \nabla f(\mathbf{x})^T \mathbf{d}$$

Choosing $\mathbf{d} = -\nabla f(\mathbf{x})$ can decrease the objective value.

In the gradient descent method, when at point \mathbf{x}^k , we choose $\mathbf{d} = -\nabla f(\mathbf{x}^k)$



The Step Size

Now we choose the step size α_k .

 \blacktriangleright An intuitive idea is to choose α_k to achieve the largest descent

That is, to choose α_k such that

$$\alpha_k = \operatorname{argmin}_{\alpha} f(\mathbf{x}^k + \alpha \mathbf{d}^k) \tag{1}$$

- If we get the exact α_k in (??), we say we used an exact line search method to find the step size
- We can use the golden section method to perform the exact line search
- ightharpoonup In some situations, we can even find the exact lpha analytically



Example of Exact Line Search

Consider

$$f(\mathbf{x}) = \mathbf{c}^T \mathbf{x} + \frac{1}{2} \mathbf{x}^T Q \mathbf{x}$$
 (Q positive definite)

At \mathbf{x}^k , the gradient descent method will choose

$$\mathbf{d}^k = -\nabla f(\mathbf{x}^k) = -(\mathbf{c} + Q\mathbf{x}^k)$$

To choose the step size, notice that we can explicitly compute

$$f(\mathbf{x}^k + \alpha \mathbf{d}^k) = \mathbf{c}^T (\mathbf{x}^k + \alpha \mathbf{d}^k) + \frac{1}{2} (\mathbf{x}^k + \alpha \mathbf{d}^k)^T Q (\mathbf{x}^k + \alpha \mathbf{d}^k)$$
$$= \frac{1}{2} \alpha^2 (\mathbf{d}^k)^T Q \mathbf{d}^k + \alpha (\mathbf{c}^T \mathbf{d}^k + (\mathbf{x}^k)^T Q \mathbf{d}^k) + f(\mathbf{x}^k)$$

This is a quadratic function of α with positive second-order term. Thus we can find the optimal α that minimizes $f(\mathbf{x}^k + \alpha \mathbf{d}^k)$:

$$\alpha_k = \frac{(\mathbf{d}^k)^T \mathbf{d}^k}{(\mathbf{d}^k)^T Q \mathbf{d}^k}$$



Line Search Methods

However, in general, one cannot expect that

$$\alpha_k = \operatorname{argmin}_{\alpha} f(\mathbf{x}^k + \alpha \mathbf{d}^k) \tag{2}$$

can be solved explicitly. It is sometimes time-consuming.

Moreover, it is not clear how much benefit there is to solve (??) exactly. After all, it is just one iteration, it doesn't mean $\mathbf{x}^k + \alpha_k \mathbf{d}^k$ is optimal

Therefore, it might be good enough to get an approximately good point.

► There are multiple ways to do it, here we introduce the backtracking line search



Backtracking Line Search

Assume we have found \mathbf{d}^k and we want to choose step size α_k .

- 1. We first choose a small $\alpha \in (0,0.5)$. Also choose a constant $0 < \beta < 1$
- 2. Let t = 1
- 3. If $f(\mathbf{x}^k + t\mathbf{d}^k) \leq f(\mathbf{x}^k) + \alpha t \nabla f(\mathbf{x}^k)^T \mathbf{d}^k$, then choose $\alpha_k = t$. Otherwise, set $t = \beta t$ and repeat this step.

Why this works?

▶ We know by Taylor expansion, if t is sufficiently small, we must have

$$f(\mathbf{x}^k + t\mathbf{d}^k) \approx f(\mathbf{x}^k) + t\nabla f(\mathbf{x}^k)^T \mathbf{d}^k < f(\mathbf{x}^k) + \alpha t\nabla f(\mathbf{x}^k)^T \mathbf{d}^k$$

Therefore, as long as t is small enough, the condition in Step 3 must be satisfied (remember $\nabla f(\mathbf{x}^k)^T \mathbf{d}^k = -||\nabla f(\mathbf{x}^k)||^2 < 0$)



Stopping Criterion for Gradient Descent Method

Remember that for local optimality, we need

$$\nabla f(\mathbf{x}) = 0$$

Since we don't know the optimal value, we use the gradient as the stopping criterion:

▶ We stop when $||\nabla f(\mathbf{x})|| < \epsilon$ for a pre-chosen ϵ

Theorem

Suppose $f(\mathbf{x})$ is convex and the smallest eigenvalue of $\nabla^2 f(\mathbf{x})$ is m. Then

$$||\mathbf{x} - \mathbf{x}^*|| \leq \frac{2}{m} ||\nabla f(\mathbf{x})||$$

where \mathbf{x}^* is the global minimum of $f(\mathbf{x})$

▶ Therefore, when $||\nabla f(\mathbf{x})||$ is small enough, the solution is guaranteed to be close to the optimal solution

Gradient Descent Algorithm

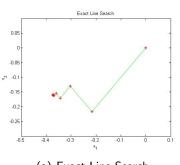
Start with any point \mathbf{x}^0 . Set k=0 and stopping criterion $\epsilon>0$

- 1. Check $||\nabla f(\mathbf{x}^k)||$. If $||\nabla f(\mathbf{x}^k)|| \le \epsilon$, stop and output \mathbf{x}^k . Otherwise, continue to Step 2
- 2. Let $\mathbf{d}^k = -\nabla f(\mathbf{x}^k)$
- 3. Use either exact line search or backtracking line search to find $\alpha_{\it k}$
- 4. Let $\mathbf{x}^{k+1} = \mathbf{x}^k + \alpha_k \mathbf{d}^k$, let k = k + 1. Go back to step 1.

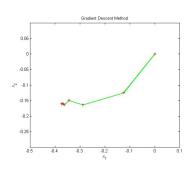


Illustration

Minimize $f(x) = \exp(x_1 + x_2) + x_1^2 + 3x_2^2 - x_1x_2$ using gradient method.



(a) Exact Line Search



(b) Backtracking Line Search