

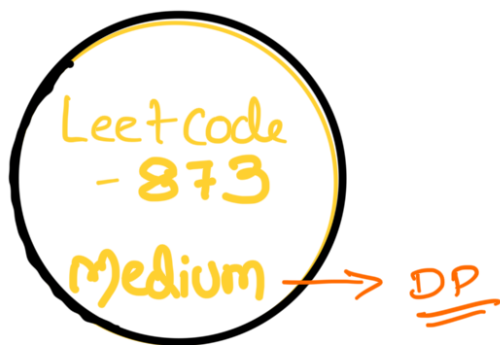
Dynamic

Video-107

Programming

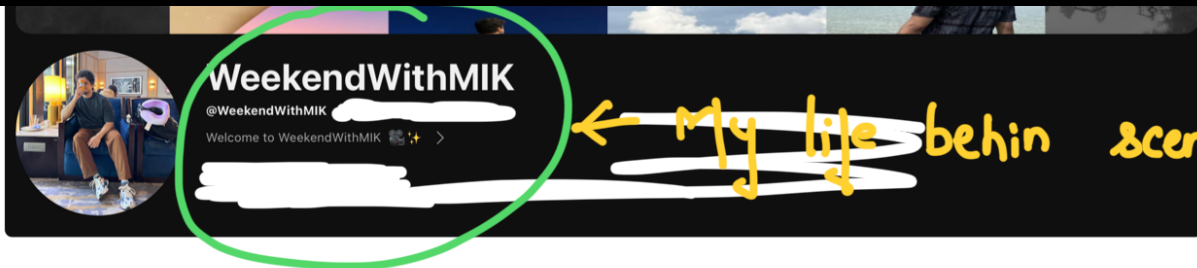


Note :- This playlist is on for explanation of Dns & solutions. See my "DP Concepts & Dns" playlist for understanding DP from scratch...



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Motivation :-

No one is perfect. Stop running after perfection. You cannot be 100%. Try for everything. You just have to keep moving in your own comfortable steady pace. Don't rush, it will all fall in right place in right time.

MIK...

873. Length of Longest Fibonacci Subsequence

Medium

Topics

Companies

A sequence x_1, x_2, \dots, x_n is Fibonacci-like if:

x_1, x_2, x_3, x_4, x_5

- $n \geq 3$
- $x_i + x_{i+1} == x_{i+2}$ for all $i + 2 \leq n$

Given a strictly increasing array `arr` of positive integers forming a sequence, return the length of the longest Fibonacci-like subsequence of `arr`. If one does not exist, return 0.

A subsequence is derived from another sequence `arr` by deleting any number of elements (including none) from `arr`, without changing the order of the remaining elements. For example, [3, 5, 8] is a subsequence of [3, 4, 5, 6, 7, 8].

F F F

Example :: arr = [1, 2, 3, 4, 5, 6, 7, 8]

Output = 5 ← {1, 2, 3, 5}

arr = [1, 3, 7, 11, 12, 14, 18]

Output = 3 ∈ {3, 11, 14} ✓
 {7, 11, 18} ✓
 {1, 11, 12} ✓

Thought Pro

“Let's think from a beginner po”

map = {1 → 0, 2 → 1, 3 → 2, 4 → 3, ...}

arr = [1, 2, 3, 4, 5, 6, 7, 8]
 _i _j _k



$$C = a + b \quad ; \Rightarrow C - b = a$$

$$\begin{aligned} \text{arr}[k] - \text{arr}[j] &= \\ 8 - 5 &= \textcircled{3} \end{aligned}$$

found 3 at index $i = 2$

$$i < j$$

$$\text{arr}[l] = \underset{5}{\text{arr}[j]} - \underset{-2}{(i)}$$

found 2 at
index = 1

// store (element \rightarrow idx) in the map

for ($j = 1$; $j < n$; $j++$) {

for ($k = j+1$; $k < n$; $k++$) {

length = solve($j, k, \text{arr}, \text{mp}$);

if (length ≥ 3) {

result = max(result, length);

}

return result;

```
int Solve(j, K, arr, mp) {
```

```
    target = arr[K] - arr[j];
```

```
    if (mp.Count(target) && mp[target] < j) {
```

```
        int i = mp[target];
```

```
        return Solve(i, j, arr, mp);
```

```
    }
```

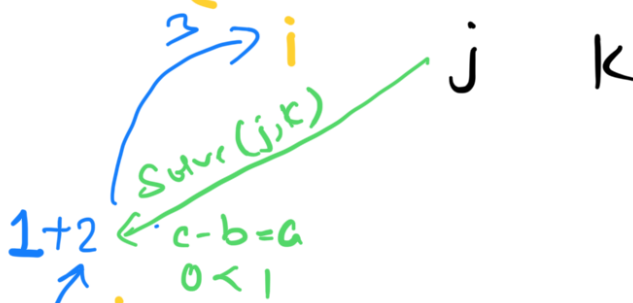
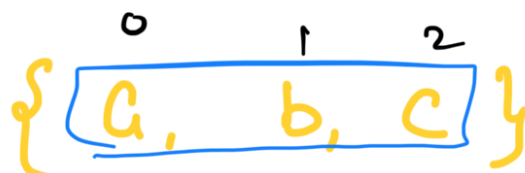
```
    return 2;
```



$arr[j] - arr[i]$



$f(j)$



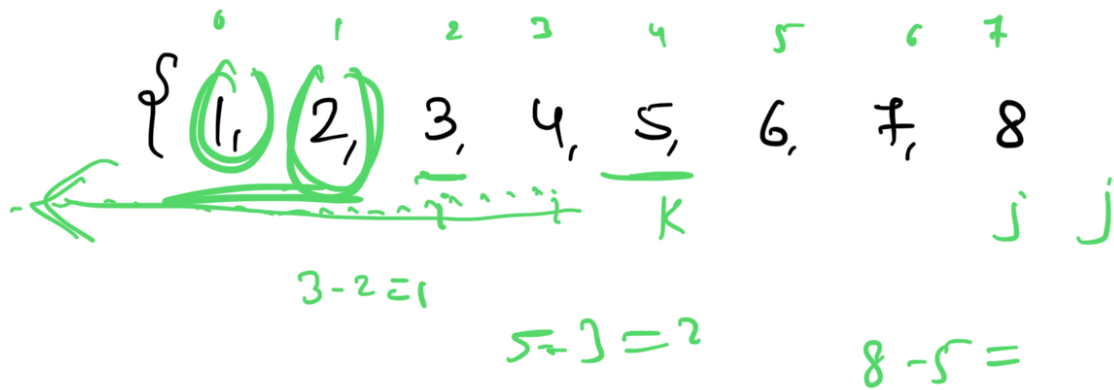
map

$a \rightarrow 0$

$b \rightarrow 1$

$c \rightarrow 2$

2 \downarrow solve(i, j)
 $b - a = z$



result = ~~3~~ ~~4~~ ~~5~~ ~~6~~ ~~7~~ ~~8~~

Recursion.

$$T.C = O(n^2 * n) = O(n^3)$$

$$S.C = O(n)$$

Why memoization :- ??

$\{1, 2, 3, 4, 5, 6, 7, 8\}$
 0 1 2 3 4 5 6 7
 i j k

$$3 + 1 = 4$$

3 Fibon.

$$f[j][k]$$

$$f[1][2] = 3 ;$$

↑ ↑

Bottom UP :-

Solve (i, j, arr, mp)

State Define:- $f[i][j]$ = Max length of Fibb'li' seq.
 ending at (i, j)

$\dots \{ \dots i, j \}$
 2

for (int j = 1 ; j < n ; j++) {

```
for (int k = j+1; k < n; k++) {
```

```
    int target = arr[k] - arr[j];
```

```
    if (mp.count(target) && mp[target] < j) {
```

```
        int i = mp[target];
```

```
        dp[j][k] = dp[i][j] + 1;
```

```
    }
```

```
    max-length = max(max-length, dp[j][k]);
```

```
}
```

```
}
```

```
return max-length >= 3 ? max-length : 0;
```

==

