

```

function Floyd( $L[1..n, 1..n]$ ): array  $[1..n, 1..n]$ 
    array  $D[1..n, 1..n]$ 
     $D \leftarrow L$ 
    for  $k \leftarrow 1$  to  $n$  do
        for  $i \leftarrow 1$  to  $n$  do
            for  $j \leftarrow 1$  to  $n$  do
                 $D[i, j] \leftarrow \min(D[i, j], D[i, k] + D[k, j])$ 
    return  $D$ 

```

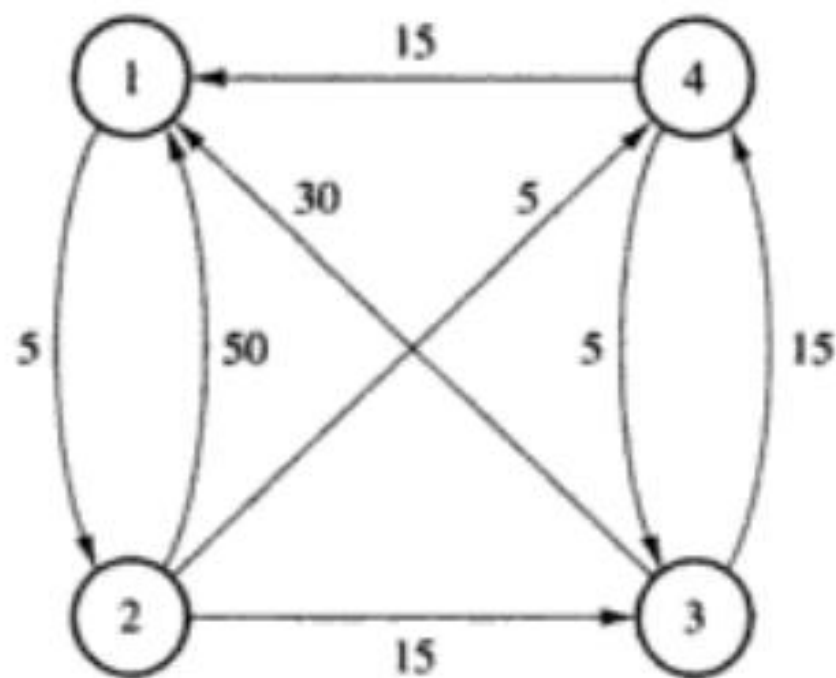
$$D_0 = L = \begin{pmatrix} 0 & 5 & \infty & \infty \\ 50 & 0 & 15 & 5 \\ 30 & \infty & 0 & 15 \\ 15 & \infty & 5 & 0 \end{pmatrix}$$

$$D_1 = \begin{pmatrix} 0 & 5 & \infty & \infty \\ 50 & 0 & 15 & 5 \\ 30 & 35 & 0 & 15 \\ 15 & 20 & 5 & 0 \end{pmatrix}$$

$$D_2 = \begin{pmatrix} 0 & 5 & 20 & 10 \\ 50 & 0 & 15 & 5 \\ 30 & 35 & 0 & 15 \\ 15 & 20 & 5 & 0 \end{pmatrix}$$

$$D_3 = \begin{pmatrix} 0 & 5 & 20 & 10 \\ 45 & 0 & 15 & 5 \\ 30 & 35 & 0 & 15 \\ 15 & 20 & 5 & 0 \end{pmatrix}$$

$$D_4 = \begin{pmatrix} 0 & 5 & 15 & 10 \\ 20 & 0 & 10 & 5 \\ 30 & 35 & 0 & 15 \\ 15 & 20 & 5 & 0 \end{pmatrix}$$



if $D[i, k] + D[k, j] < D[i, j]$ then $D[i, j] \leftarrow D[i, k] + D[k, j]$
 $P[i, j] \leftarrow k$

•

$$P = \begin{pmatrix} 0 & 0 & 4 & 2 \\ 4 & 0 & 4 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix}.$$

$$c_{ij} = \sum_{k=1}^q a_{ik}b_{kj}, \quad 1 \leq i \leq p, 1 \leq j \leq r.$$

```

for  $i \leftarrow 1$  to  $p$  do
  for  $j \leftarrow 1$  to  $r$  do
     $C[i, j] \leftarrow 0$ 
    for  $k \leftarrow 1$  to  $q$  do
       $C[i, j] \leftarrow C[i, j] + A[i, k]B[k, j]$ 

```

$$M = M_1 M_2 \cdots M_n$$

$$\begin{aligned} M &= (\cdots ((M_1 M_2) M_3) \cdots M_n) \\ &= (M_1 (M_2 (M_3 \cdots (M_{n-1} M_n) \cdots))) \\ &= (\cdots ((M_1 M_2) (M_3 M_4)) \cdots), \end{aligned}$$

$((AB)C)D$	10 582
$(AB)(CD)$	54 201
$(A(BC))D$	2 856
$A((BC)D)$	4 055
$A(B(CD))$	26 418

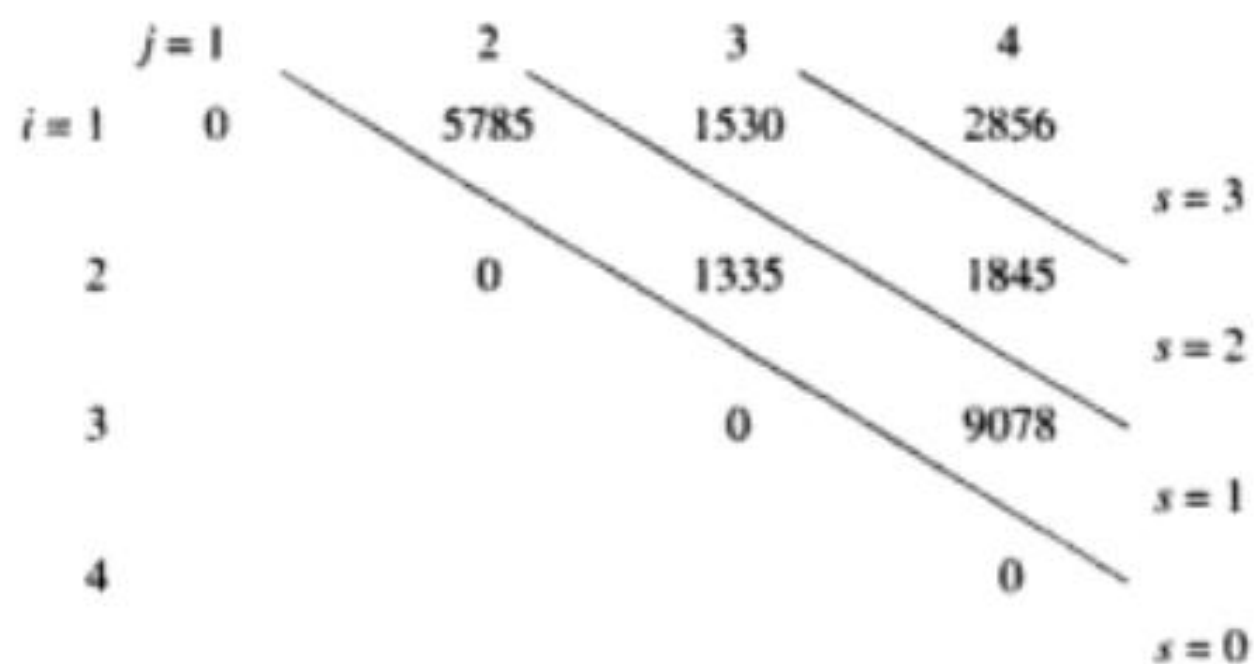
$$M = (M_1 M_2 \cdots M_i)(M_{i+1} M_{i+2} \cdots M_n).$$

$$T(n) = \sum_{i=1}^{n-1} T(i)T(n-i), \qquad T(n) \text{ is in } \Omega(4^n/n^2)$$

n	1	2	3	4	5	10	15
$T(n)$	1	1	2	5	14	4862	2674440

$$\begin{array}{llll}
s = 0 : & m_{ii} & = 0 & i = 1, 2, \dots, n \\
s = 1 : & m_{i,i+1} & = d_{i-1}d_id_{i+1} & i = 1, 2, \dots, n-1 \\
1 < s < n : & m_{i,i+s} & = \min_{i \leq k < i+s} (m_{ik} + m_{k+1,i+s} + d_{i-1}d_kd_{i+s}) & i = 1, 2, \dots, n-s
\end{array}$$

(13, 5, 89, 3, 34).



$$m_{13} = \min(m_{11} + m_{23} + 13 \times 5 \times 3, m_{12} + m_{33} + 13 \times 89 \times 3) \\ = \min(1530, 9256) = 1530$$

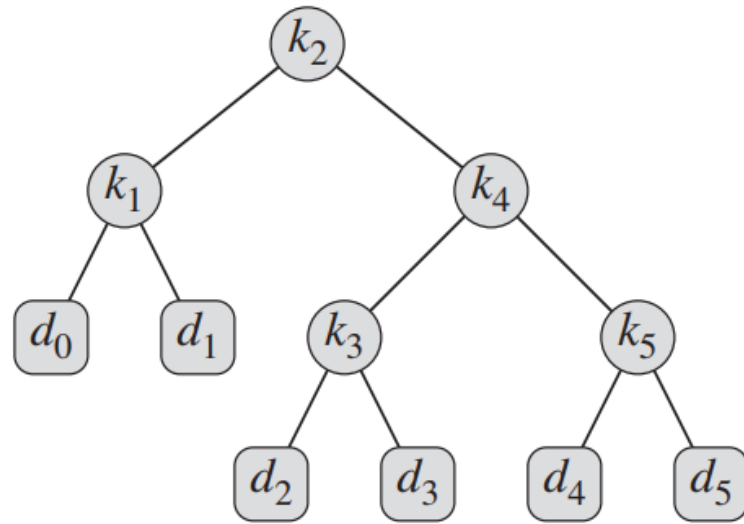
$$m_{24} = \min(m_{22} + m_{34} + 5 \times 89 \times 34, m_{23} + m_{44} + 5 \times 3 \times 34) \\ = \min(24208, 1845) = 1845.$$

Finally for $s = 3$

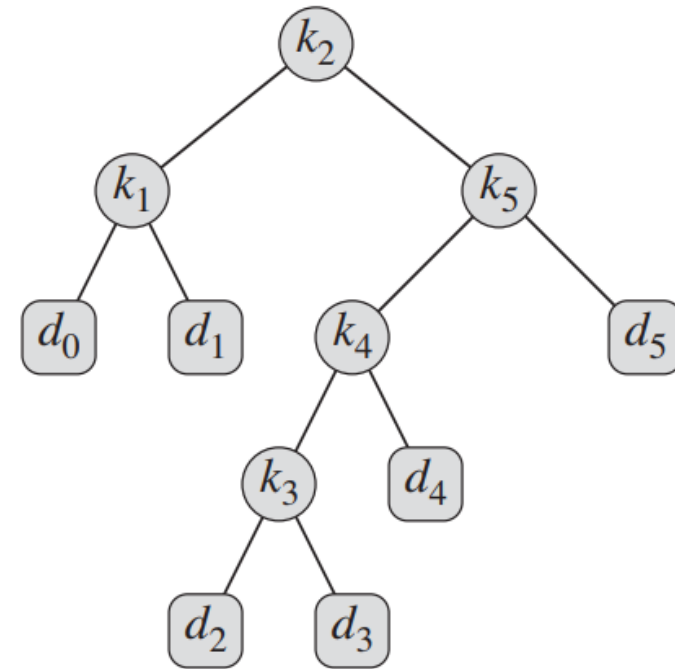
$$m_{14} = \min(m_{11} + m_{24} + 13 \times 5 \times 34, \quad \{k = 1\} \\ m_{12} + m_{34} + 13 \times 89 \times 34, \quad \{k = 2\} \\ m_{13} + m_{44} + 13 \times 3 \times 34) \quad \{k = 3\} \\ = \min(4055, 54201, 2856) = 2856.$$

$$\begin{aligned}
\sum_{s=1}^{n-1} (n-s)s &= n \sum_{s=1}^{n-1} s - \sum_{s=1}^{n-1} s^2 \\
&= n^2(n-1)/2 - n(n-1)(2n-1)/6 \\
&= (n^3 - n)/6,
\end{aligned}$$

		j	0	1	2	3	4	5	6
		i	y_j	B	D	C	A	B	A
0	x_i		0	0	0	0	0	0	0
1	A		0	↑ 0	↑ 0	↑ 0	↖1 1	←1 1	↖1 1
2	B		0	↖1 1	←1 1	←1 1	↑1 1	↖2 2	←2 2
3	C		0	↑1 1	↑1 1	↖2 2	←2 2	↑2 2	↑2 2
4	B		0	↖1 1	↑1 1	↑2 2	↑2 2	↖3 3	←3 3
5	D		0	↑1 1	↖2 2	↑2 2	↑2 2	↑3 3	↑3 3
6	A		0	↑1 1	↑2 2	↑2 2	↖3 3	↑3 3	↖4 4
7	B		0	↖1 1	↑2 2	↑2 2	↑3 3	↖4 4	↑4 4



(a)



(b)

Figure 15.9 Two binary search trees for a set of $n = 5$ keys with the following probabilities:

i	0	1	2	3	4	5
p_i		0.15	0.10	0.05	0.10	0.20
q_i	0.05	0.10	0.05	0.05	0.05	0.10

(a) A binary search tree with expected search cost 2.80. **(b)** A binary search tree with expected search cost 2.75. This tree is optimal.

