

MODULE 5

CONVERSION OF PDA TO CFG

$$Q) \quad M = (\overset{Q}{\{q_0, q_1\}}, \overset{\Sigma}{\{a, b\}}, \overset{\Gamma}{\{a, z_0\}}, \delta, q_0, z_0, \phi)$$

$$\delta(q_0, a, z_0) = (q_0, a z_0)$$

$$\delta(q_0, b, a) = (q_1, a)$$

$$\delta(q_1, a, a) = (q_1, \epsilon)$$

Soln:

Tuples of PDA :

$$M = (Q, \Sigma, \Gamma, \delta, q_0, z_0, F)$$

CFG is defined by four tuples :

$$G = \{V, T, S, P\}$$

where,

$V \rightarrow$ Set of non-terminals (A to Z)

Σ or $T \rightarrow$ Terminals (a-z, operators, Numbers, Special symbols)

$S \rightarrow$ Start symbol.

$P \rightarrow$ Production rule.

Variable, V is constructed using two rules.

1) Special symbol S

2) $[p x q]$ where p, q are states in Q & x is in Γ

Variable is constructed by combining all the possible combinations of variable in Γ with Q

$$(\overset{Q}{\{q_0, q_1\}} \times \overset{\Gamma}{\{a, z_0\}})$$

$$V = (S, [q_0 a q_0], [q_0 a q_1], [q_1 a q_0], [q_1 a q_1], [q_0 z_0 q_0], [q_0 z_0 q_1], [q_1 z_0 q_0], [q_1 z_0 q_1])$$

Variable is constructed.

Terminals, T

Write the terminal symbols given in Question directly.

$$T = \{a, b\}$$

Now, Terminals also constructed for CFG

Start symbol S

$S \rightarrow (\text{Initial state, Top symbol of stack, set of states})$

Here,

Initial state = q_0

Top symbol of stack = z_0

Set of states = $\{q_0, q_1\}$

$$S \rightarrow [q_0 z_0 q_0]$$

$$S \rightarrow [q_0 z_0 q_1]$$

} Start symbol.

Productions.

$$\delta(q_0, a, z_0) = (q_0, \underbrace{az_0}_2)$$

NOTE: There are two symbols in stack, so need to construct four production rules.

$$\delta(q_0, a, z_0) = (q_0, az_0)$$

$$[q_0 z_0 q_0] \rightarrow a [q_0 a q_0] [q_0 z_0 q_0]$$

$$[q_0 z_0 q_0] \rightarrow a [q_0 a q_1] [q_1 z_0 q_0]$$

$$[q_0 z_0 q_0] \rightarrow a [q_0 a q_0] [q_0 z_0 q_1]$$

$$[q_0 z_0 q_1] \rightarrow a [q_0 a q_1] [q_1 z_0 q_1]$$

(Since 2 symbols in the stack, Right side of P. rule contains 2 Brackets).

$$\delta(q_0, b, a) = (q_1, a)$$

(Stack contains only one symbol,
so 2 production rules can be
formed)

$$\delta(q_0, b, a) = (q_1, a)$$

$$[q_0 \ a \ q_0] \rightarrow b [q_1 \ a \ q_0]$$

$$[q_0 \ a \ q_1] \rightarrow b [q_1 \ a \ q_1]$$

(Since ^{stack} it contains only
one symbol, Right side
of rule has 1 bracket)

$$\delta(q_1, a, a) = (q_1, \epsilon)$$

$$[q_1 \ a \ q_1] \rightarrow a$$

(Since ^{stack} epsilon, only one
production)

The Tuples of CFG are,

$$G = \{ V, T, S, P \}$$

(After finding every tuples
separately, atlast (finally
write it like this))

(This part is Important)

$$V = \left(s, [q_0 \ a \ q_0], [q_0 \ a \ q_1], [q_1 \ a \ q_0], [q_1 \ a \ q_1], \right. \\ \left. [q_0 \ z_0 \ q_0], [q_0 \ z_0 \ q_1], [q_1 \ z_0 \ q_0], [q_1 \ z_0 \ q_1] \right)$$

$$T = \{ a, b \}$$

$$S \rightarrow [q_0 \ z_0 \ q_0]$$

$$S \rightarrow [q_0 \ z_0 \ q_1]$$

Production Rules : (P)

$$[q_0 \ z_0 \ q_0] \rightarrow a [q_0 \ a \ q_0] [q_0 \ z_0 \ q_0]$$

$$[q_0 \ z_0 \ q_0] \rightarrow a [q_0 \ a \ q_1] [q_1 \ z_0 \ q_0]$$

$$[q_0 \ z_0 \ q_1] \rightarrow a [q_0 \ a \ q_0] [q_0 \ z_0 \ q_1]$$

$$[q_0 \ z_0 \ q_1] \rightarrow a [q_0 \ a \ q_1] [q_1 \ z_0 \ q_1]$$

$$[q_0 \ a \ q_0] \rightarrow b [q_1 \ a \ q_0]$$

$$[q_0 \ a \ q_1] \rightarrow b [q_1 \ a \ q_1]$$

$$[q_1 \ a \ q_1] \rightarrow a$$