

Module IV

③ Syntax Directed Translation is a generalisation of a context free grammar in which each grammar symbol has an associated set of attributes, and partition into 2 subsets called synthesized and inherited attributes of that grammar symbol.

Grammar + Semantic Rules.

Grammar:

$$E \rightarrow E + T \mid T$$

$$T \rightarrow T * F \mid F$$

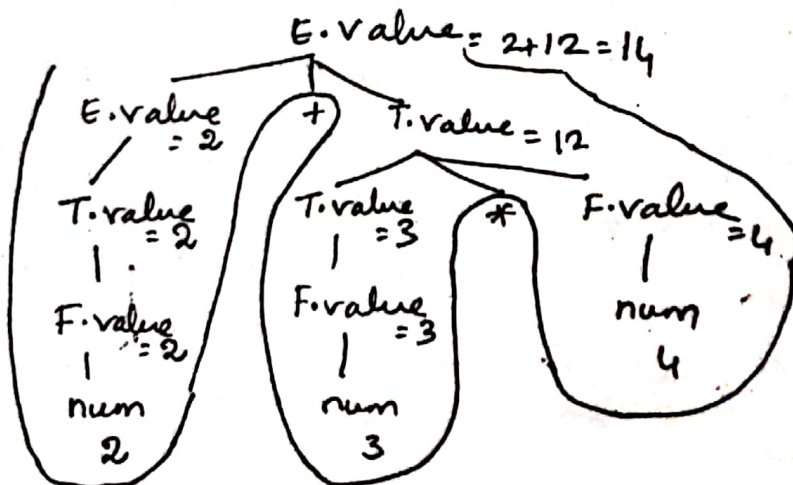
$$F \rightarrow \text{num}$$

Semantic Rules

$E \rightarrow E + T$	$\{ E.\text{Value} = E.\text{value} + T.\text{value} \}$
$E \rightarrow T$	$\{ E.\text{Value} = T.\text{value} \}$
$T \rightarrow T * F$	$\{ T.\text{Value} = T.\text{value} * F.\text{value} \}$
$T \rightarrow F$	$\{ T.\text{Value} = F.\text{value} \}$
$F \rightarrow \text{num}$	$\{ F.\text{Value} = \text{num}.\text{value} \}$

lexical value.

Consider the expression, $2 + 3 * 4$



Infix to postfix

$E \Rightarrow E + T$ { printf("+"); } ①

$E \Rightarrow T$ { } ②

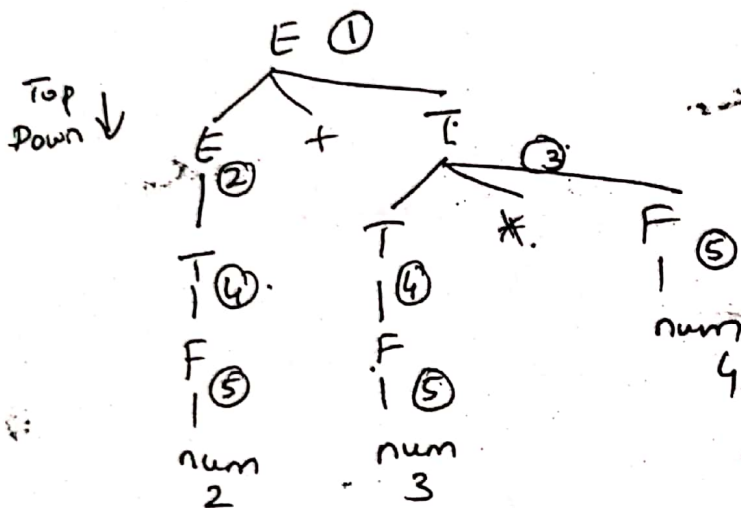
$T \Rightarrow T * F$ { printf("*"); } ③

$T \Rightarrow F$ { } ④

$F \Rightarrow \text{num}$ { printf(num = lval); } ⑤

Consider the expression $2 + 3 * 4$ ^{infix}. Convert into postfix

Top Down parsing



Postfix expression : $234 * +$

To build abstract Syntax tree

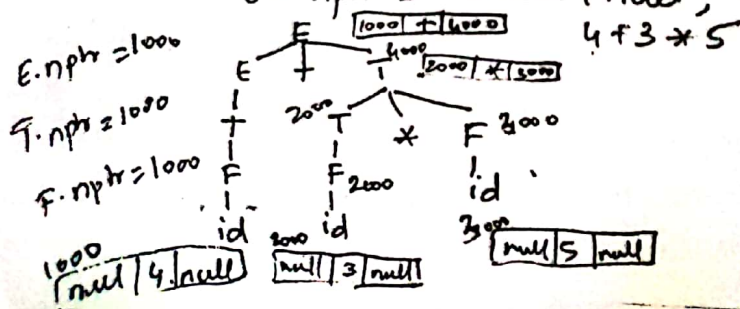
$E \Rightarrow E + T$ { E.nptr = mknode(E.nptr, '+', T.nptr); }

$E \Rightarrow T$ { E.nptr = T.nptr; }

$T \Rightarrow T * F$ { T.nptr = mknode(T.nptr, '*', F.nptr); }

$T \Rightarrow F$ { T.nptr = F.nptr; }

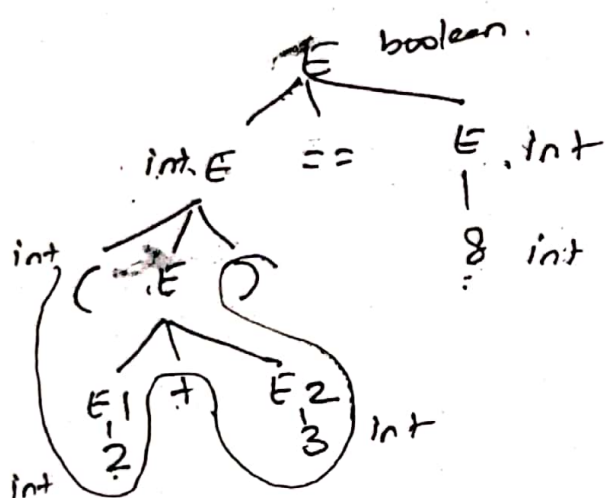
$F \Rightarrow \text{id}$ { F.nptr = mknode(null, ldbname, null); }



for type checks

$E \Rightarrow E_1 + E_2 / \{ \text{if } ((E_1 \text{ type} == E_2 \text{ type}) \ \&\& \ (E_1 \text{ type} = \text{int})) \text{ then } E \text{ type} = \text{int} \text{ else error} \}$
 $E_1 == E_2 / \{ \text{if } ((E_1 \text{ type} == E_2 \text{ type})) \ \&\& \ (E_1 \text{ type} = \text{int} / \text{boolean}) \text{ then } E \text{ type} = \text{boolean} \text{ else error} \}$
 $(E_1) / \{ E \text{ type} = E_1 \text{ type}; \}$
 $\text{num} / \{ E \text{ type} = \text{int}; \}$
 $\text{True} / \{ E \text{ type} = \text{boolean}; \}$
 $\text{False} / \{ E \text{ type} = \text{boolean}; \}$

Expression: $(2+3) == 8$



SAT for three address code generation.

$S \Rightarrow \text{id} = E \quad \{ \text{gen}(\text{id.name} = E \text{ place}); \}$

$E \Rightarrow E_1 + T / \{ E \text{ place} = \text{new Temp}(); \}$
 $\quad \quad \quad \text{gen}(E \text{ place} = E_1 \text{ place} + T \text{ place}); \}$

$T \quad \quad \{ E \text{ place} = T \text{ place}; \}$

$T \Rightarrow T_1 * F / \{ T \text{ place} = \text{new Temp}(); \}$
 $\quad \quad \quad \text{gen}(T \text{ place} = T_1 \text{ place} * F \text{ place}); \}$

$F \quad \quad \{ T \text{ place} = F \text{ place}; \}$

$F \Rightarrow \text{id} \quad \{ F \text{ place} = \text{id.name}; \}$

