

TOY1 Assembly Programming

Professor Kin K. Leung

kin.leung@imperial.ac.uk

www.commsp.ee.ic.ac.uk/~kkleung/

Heavily based on materials by Dr. Naranker Dulay

TOY1 Instruction Set

Opcode	Assembly Instruction		Action
0000	STOP		Stop Program Execution
0001	LOAD	Rn, [Addr]	$Rn = \text{Memory}[\text{Addr}]$
0010	STORE	Rn, [Addr]	$\text{Memory}[\text{Addr}] = Rn$
0011	ADD	Rn, [Addr]	$Rn = Rn + \text{Memory}[\text{Addr}]$
0100	SUB	Rn, [Addr]	$Rn = Rn - \text{Memory}[\text{Addr}]$
0101	GOTO	Addr	$PC = \text{Addr}$
0110	IFZER	Rn, Addr	IF $Rn == 0$ THEN $PC = \text{Addr}$
0111	IFNEG	Rn, Addr	IF $Rn < 0$ THEN $PC = \text{Addr}$

Example 1 - Multiplication

➤ Compute

$$A = B * C$$

$$B * C = \sum_{N=1}^C B$$

➤ Examples

$$12 * 3 = 12 + 12 + 12$$

$$12 * 1 = 12$$

$$12 * 0 = 0$$

A Solution

; Given: A, B, C

; Pre: C \geq 0

; Post: A = B * C

sum = 0 ; accumulate result in sum

n = C ; indicates how many additions remain

loop

exit when n \leq 0 ; no more additions remain

 sum = sum + B ; add another B

 n = n - 1 ; one less addition to do

end loop

A = sum

Variables

➤ We'll allocate variables to main memory as follows:

A to Memory [100H]

B to Memory [101H]

C to Memory [102H]

sum to R1

n to R2

sum = 0

- We need to perform **R1 = 0** but the closest instruction we have is **LOAD Rn, [Addr]**
- So we'll pre-set a memory location (e.g. memory word 200H) to 0.
- Now **sum = 0** simply translates to

LOAD R1, [200H] ; sum = 0

Memory Contents

- We'll fill main memory as follows:

Address	Contents
---------	----------

- | | |
|--------|---|
| ➤ 100H | A |
| 101H | B |
| 102H | C |

- | | |
|--------|---|
| ➤ 200H | 0 |
|--------|---|

- | | |
|--------|----------------------------|
| ➤ 080H | 1st Instruction of program |
| 081H | 2nd Instruction of program |
| | Etc. |

n = C

➤ This is easy to translate **LOAD R2, [102H]**

➤ We now have

080H	LOAD R1, [200H]	; sum = 0
081H	LOAD R2, [102H]	; n = C

...

100H	A	; holds A
101H	B	; holds B
102H	C	; holds C

...

200H	0	; constant 0
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exitloop when $n \leq 0$

- Let's consider a simpler statement **exitloop when $n = 0$**

Statement

```
loop
  exit when n = 0
  instructions
end loop
....
```

Address

L0

...

LY

Instructions

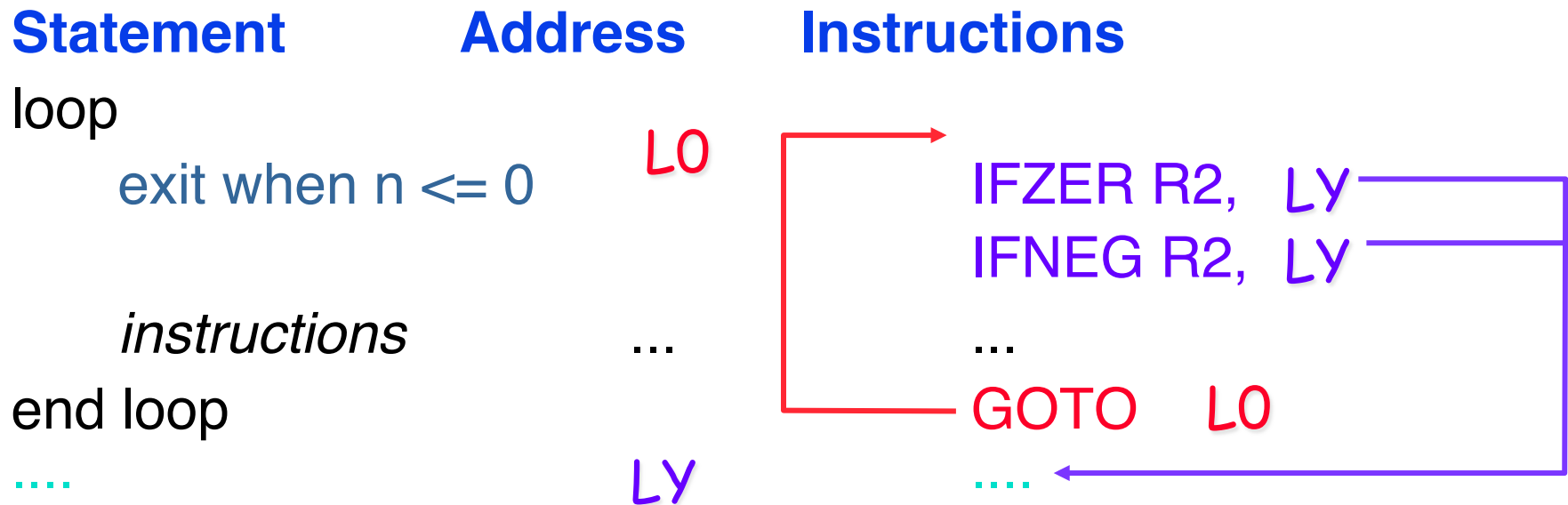
IFZER R2, **LY**

...


GOTO L0

....

exit when $n \leq 0$



Memory Contents

080H	LOAD R1, [200H]	; sum = 0	
081H	LOAD R2, [102H]	; n = C	
082H	IFZER R2, LY	; loop exit when n <= 0	
083H	IFNEG R2, LY		
...	...	; ...	
...	GOTO 082H	; end loop	
LY	...	; ...	
...	...	; ...	
100H	A	; holds A	
101H	B	; holds B	
102H	C	; holds C	
...			
200H	0	; constant 0	

sum = sum + B

- This simply translates to

ADD R1, [101H]

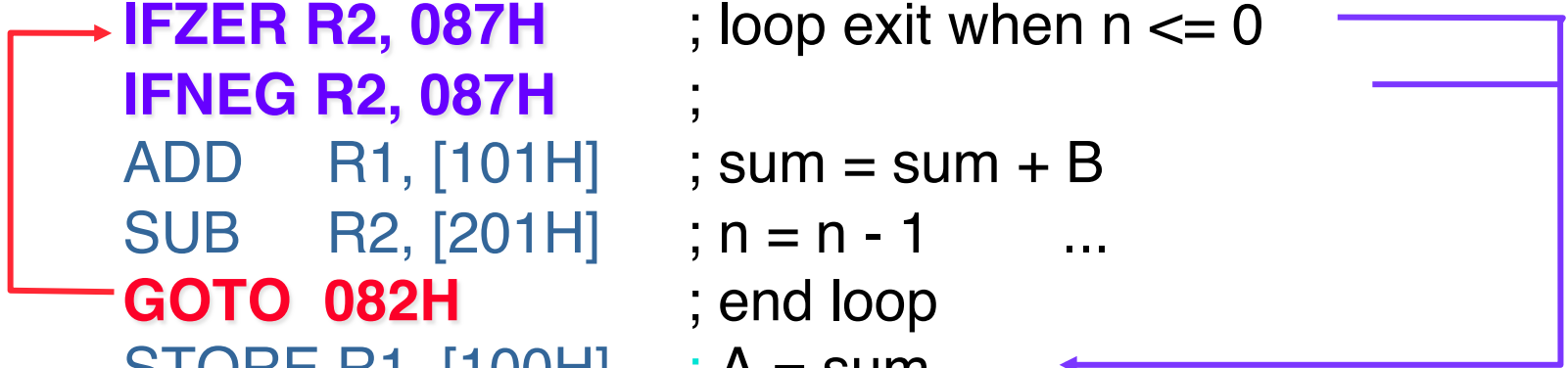
n = n - 1

- This is easy to translate if we pre-set another location (memory word 201H) with the constant 1

SUB R2, [201H]

Final Program

080H	LOAD R1, [200H]	; sum = 0
081H	LOAD R2, [102H]	; n = C
082H	IFZER R2, 087H	; loop exit when n <= 0
083H	IFNEG R2, 087H	;
084H	ADD R1, [101H]	; sum = sum + B
085H	SUB R2, [201H]	; n = n - 1 ...
086H	GOTO 082H	; end loop
087H	STORE R1, [100H]	; A = sum
088H	STOP	; That's all folks!
100H	A	; holds A
101H	B	; holds B
102H	C	; holds C
200H	0	; constant 0
201H	1	; constant 1



Example 2 - Vector Sum

➤ $\text{Sum} = \text{Memory}[200\text{H}] + \dots + \text{Memory}[200\text{H} + 99]$

<code>sum = 0</code>	<code>; accumulates result</code>
<code>n = 100</code>	<code>; no. of elements to add</code>
<code>addr = 200H</code>	<code>; address of 1st element</code>
loop	
<code>exit when n <= 0</code>	<code>; elements remaining</code>
<code>sum = sum + Memory [addr]</code>	<code>; add next element</code>
<code>addr = addr + 1</code>	<code>; advance to next element</code>
<code>n = n - 1</code>	<code>; one less element to add</code>
end loop	

TOY1 Instruction Set Continued

Opcode	Assembly Instruction	Action
1000	spare	
1001	LOAD Rn, [Rm]	$Rn = \text{Memory}[Rm]$
1010	STORE Rn, [Rm]	$\text{Memory}[Rm] = Rn$
1011	ADD Rn, [Rm]	$Rn = Rn + \text{Memory}[Rm]$
1100	SUB Rn, [Rm]	$Rn = Rn - \text{Memory}[Rm]$
1101	spare	
1110	spare	
1111	spare	

Instruction Format 2



➤ Example: Disassemble the instruction CEAB H

CEAB = 1100 1110 1010 1011

 ↓ ↓ ↓ ↓

 SUB R3 [R2] Unused

Memory Allocation

- We'll allocate variables to memories as follows:

Vector	to	Memory [200H] .. Memory [200H+99]
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➤ sum	to	R0
n	to	R1
addr	to	R2

➤ Constants	at	Memory [000H] onwards
Program	at	Memory [00FH] onwards

TOY1 Vector Sum Program

sum = 0	; accumulates result
0FH	LOAD R0, [00H]
00H	0
01H	1
n = 100	; no. of elements to add
10H	LOAD R1, [02H]
02H	100
addr = 200H	; address of 1st element
11H	LOAD R2, [03H]
03H	200H

TOY1 Vector Sum Program Contd.

loop exit when $n \leq 0$; elements remaining

12H IFZER R1, 18H

13H IFNEG R1, 18H

sum = sum + Memory [addr] ; add next element

14H ADD R0, [R2]

addr = addr + 1 % addr of next element

15H ADD R2, 01H

n = n - 1 % one less element to add

16H SUB R1, 01H

end loop

17H GOTO 12H

18H STOP

Think About

- Translating High-level Language statements to Assembly Language instructions (Compilation)
 - Allocating Variables to Registers and Main Memory
 - Branching and Looping
 - Indirect Addressing
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