

leetcode Question: Linked List Cycle II

Linked List Cycle II

Given a linked list, return the node where the cycle begins. If there is no cycle, return null.

Follow up:

Can you solve it without using extra space?

Analysis:

This problem can be viewed as two major steps:

- (1) Detect whether the loop exists in the linked list.
- (2) Find the loop start node if loop exists.

The (1) step can be easily solved using the slow&fast pointers (see [Linked List Cycle](#))

How to deal with step (2)?

Firstly let us assume the slow pointer (S) and fast pointer (F) start at the same place in a **n node** circle. S run **t steps** while F can run **2t steps**, we want to know what is t (where they meet) , then just solve: **$t \bmod n = 2t \bmod n$** , it is not hard to know that when $t = n$, they meet, that is the start of the circle.

For our problem, we can consider the time when S enter the loop for the 1st time, which we assume k step from the head. At this time, the F is already in k step ahead in the loop. When will they meet next time?

Still solve the function: **$t \bmod n = (k + 2t) \bmod n$**

Finally, we can find out: **$t = (n-k)$** , S and F will meet, this is k steps before the start of the loop.

So, the Step (2) can be easily solved after Step(1):

Note that here S and F are not slow or fast pointers, but only regular pointers.

1. Set S to the head.
2. $S = S \rightarrow next$, $F = F \rightarrow next$
3. output either one of the pointer until they meet.

Here I show some figure for better illustration:

Note that, from step 5, p (slow pointer) firstly enter the loop, it will go (n-k) steps, and meet with fast pointer q. The rest of step is $n-(n-k) = k$, steps from current meet point, back to the start of the loop.



