

Topic 38 – Graph Radio Coloring Parallelization

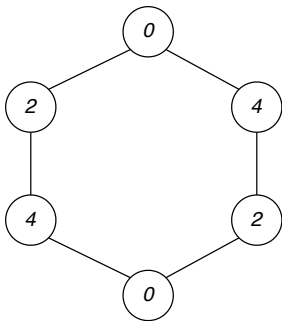
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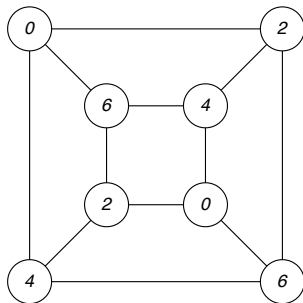
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17. December 2020



(a) 6-cycle graph radio coloring



(b) 3-dimensional cube graph radio coloring

Given an undirected graph $G = (V, E)$, the radiocoloring of G is a mapping $f : V \rightarrow N$, such that $\forall v, u \in V : |f(u) - f(v)| \geq k + 1 - d(u, v)$, where $d(u, v)$ is a distance between u and v in G .

```
# Step 1)
```

```
dist_matrix = floyd_warshall(adj_matrix)
```

```
C = Matrix(n, n)
```

```
# Step 2)
```

```
for i in (0 ... n):
```

```
    for j in (0 ... n):
```

```
        C(i)(j) = max(k+1 - dist_matrix(i)(j), 0)
```

```
    C(i)(i) = INF
```

```
labels = (0, INF, INF, ...)
```

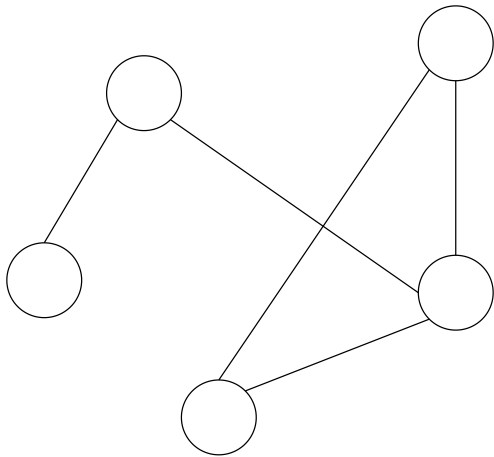
```
last = 0
```

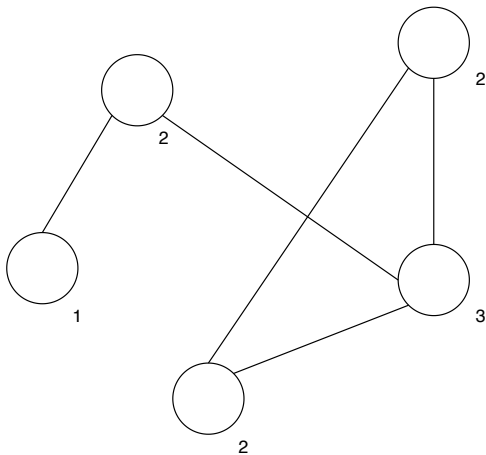
```
# Step 3)
repeat(n - 1):
    min_label = INF
    for j in (0 ... n):
        if min_label > C(last)(j):
            min_label = C(last)(j)
            p = j
    for j in (0 ... n):
        C(p)(j) += min_label
    for j in (0 ... n):
        if C(p)(j) < C(last)(j):
            C(p)(j) = C(last)(j)
    labels(p) = min_label
    last = p
```

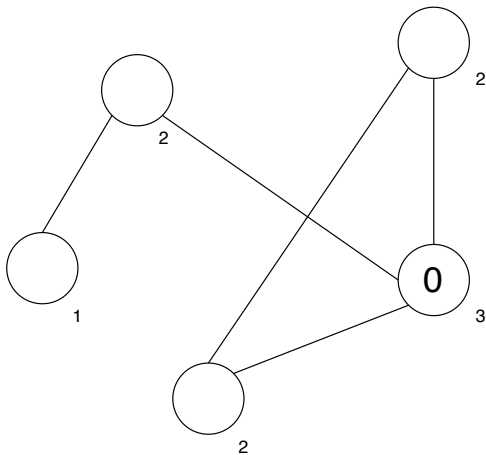
- Largest-degree first
- Algorithm
 - Sort the vertices in descending order by their degree.
 - Compute distance 2 binary matrix.
 - Initialize a $n \times 2n - 1$ binary matrix “forbidden” ($O(n^2)$).
 - Find the smallest non-conflicting color, assign it and update forbidden matrix.

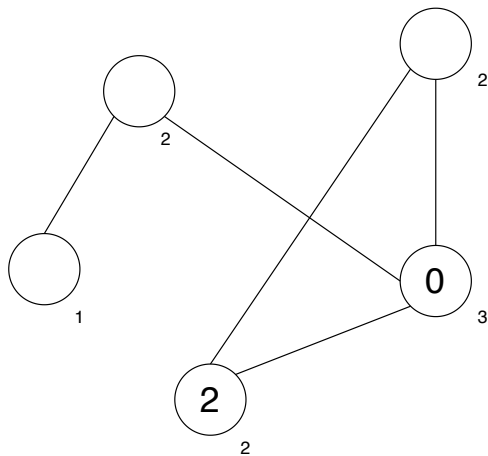
$$\text{DTIME} \approx O(n^2) + O(n \log n) + O(n^3) + O(n^3) \approx O(n^3)$$

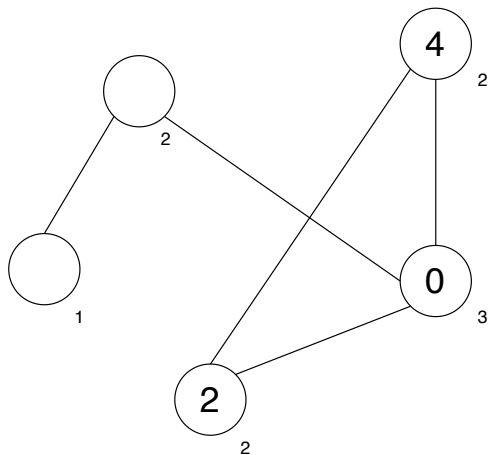
$$\text{DTIME}_n \approx O(n) + O(\log n) + O(n^2) + O(n \log n) \approx O(n^2)$$

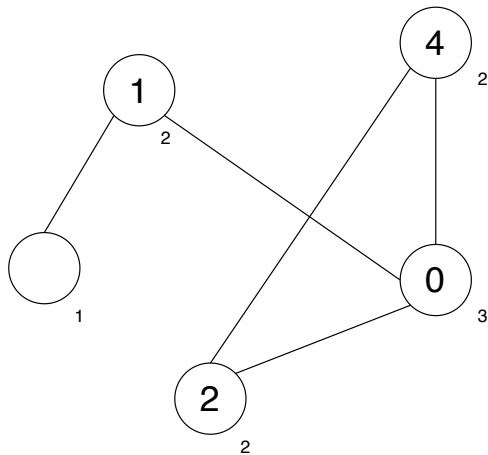


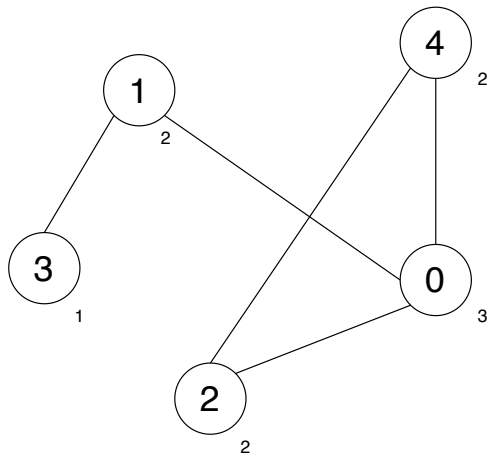


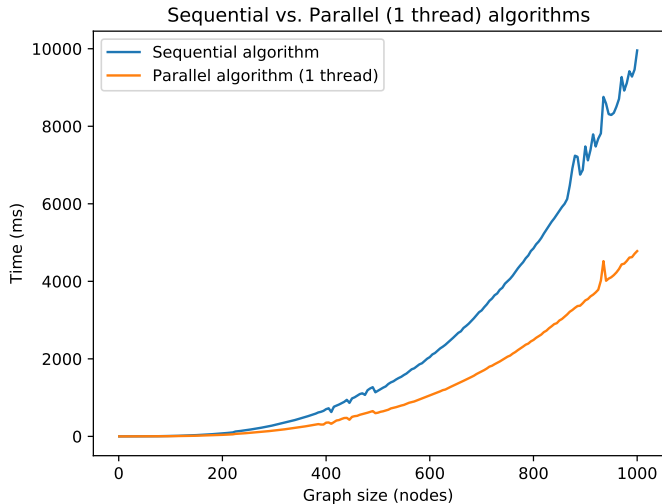


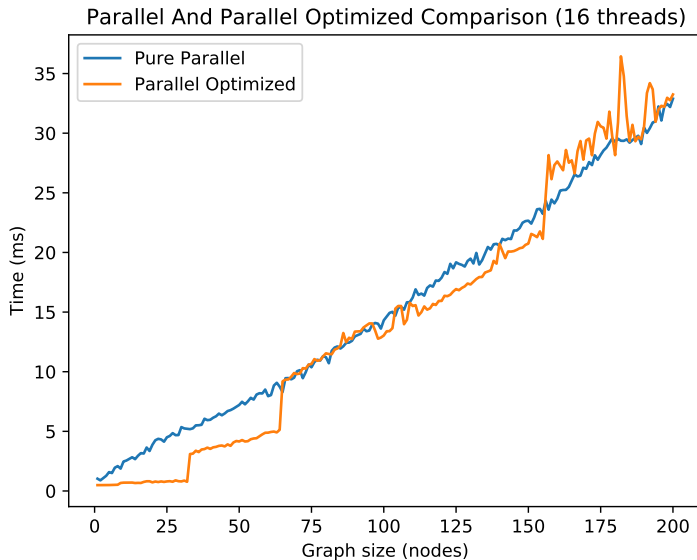




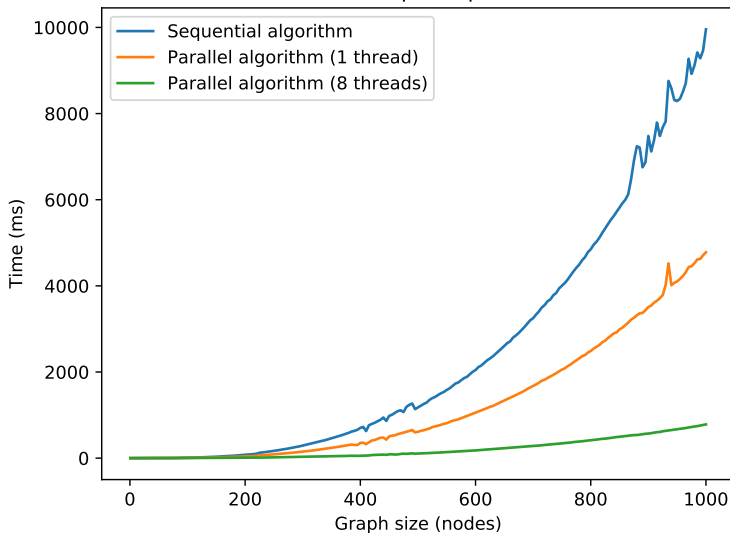








Parallelization speedup - 8 threads



Thank you for your attention