

## CS 61C Fall 2014 Discussion 2 – MIPS

C	MIPS
<pre>// \$s0 -&gt; a, \$s1 -&gt; b // \$s2 -&gt; c, \$s3 -&gt; z  int a = 4, b = 5, c = 6, z; z = a + b + c + 10;</pre>	<pre>addiu \$s0, \$0, 4 addiu \$s1, \$0, 5 addiu \$s2, \$0, 6 addu \$s3, \$s0, \$s1 addu \$s3, \$s3, \$s2 addiu \$s3, \$s3, 10</pre>
<pre>// \$s0 -&gt; int * p = intArr; // \$s1 -&gt; a;  *p = 0; int a = 2; p[1] = p[a] = a;</pre>	<pre>sw \$0, 0(\$s0) addiu \$s1, \$0, 2 sw \$s1, 4(\$s0) sll \$t0, \$s1, 2 add \$t0, \$t0, \$s0 sw \$s1, 0(\$t0)</pre>
<pre>// \$s0 -&gt; a, \$s1 -&gt; b  int a = 5, b = 10; if(a + a == b) {     a = 0; } else {     b = a - 1; }</pre>	<pre>addiu \$s0, \$0, 5 addiu \$s1, \$0, 10 addu \$t0, \$s0, \$s0 bne \$t0, \$s1, else xor \$s0, \$0, \$0 j exit  else:     addiu \$s1, \$s0, -1  exit:</pre>
<pre>// computes s1 = 2<sup>30</sup>  s1 = 1; for(s0=0;s0&lt;30;s++) {     s1 *= 2; }</pre>	<pre>addiu \$s0, \$0, 0 addiu \$s1, \$0, 1 addiu \$t0, \$0, 30  loop:     beq \$s0, \$t0, exit     addu \$s1, \$s1, \$s1     addiu \$s0, \$s0, 1     j loop  exit:</pre>
<pre>int sum(int n) {     int sum;     for(sum=0;n&gt;0;sum+=n--);     return sum; }</pre>	<pre>sum:     xor \$v0, \$0, \$0  loop:     blez \$a0, exit     addu \$v0, \$v0, \$a0     addiu \$a0, \$a0, -1     j loop  exit:     jr \$ra</pre>

Implement `streq`, which sets `$v0` to true only when its two character pointer arguments (`$a0` and `$a1`) point to equal strings, first in C, then in MIPS.

C	MIPS
<pre>int streq(char * s1, char * s2) {     do {         if(*s1 != *s2) {             return 0;         }         s2++;     } while(*s1++);     return 1; }</pre>	<pre>streq:     lb \$t0, 0(\$a0)     lb \$t1, 0(\$a1)     bne \$t0, \$t1, false     beq \$t0, \$0, true     addiu \$a0, \$a0, 1     addiu \$a1, \$a1, 1     j streq false:     xor \$v0, \$0, 0     jr \$ra true:     addiu \$v0, \$0, 1     jr \$ra</pre>

What are the instructions to branch on each of the following conditions?

<code>\$s0 &lt; \$s1</code>	<code>\$s0 &lt;= \$s1</code>	<code>\$s0 &gt; 1</code>	<code>\$s0 &gt;= 1</code>
<pre>slt \$t0, \$s0, \$s1 bne \$t0, \$0, label</pre>	<pre>slt \$t0, \$s1, \$s0 beq \$t0, \$0, label</pre>	<pre>sltiu \$t0, \$s0, 2 beq \$t0, \$0, label</pre>	<pre>bgtz \$s0, label</pre>

There are a few different meanings that the term “unsigned” takes in MIPS. Describe all three meanings, and give example instructions.

	Unsigned Meaning	Examples
1	Do not sign extend loaded data	<code>lbu, lhu</code>
2	No errors on <b>signed</b> overflow	<code>addu, subu, addiu</code>
3	Perform unsigned operations	<code>multu, divu, sltu, sltiu</code>

What is the distinction between zero extension and sign extension?

Sign extension and zero extension both extend a data type of a smaller bit width  $M$  to a data type of a larger bit width  $N$ . To zero extend from  $M$  to  $N$ , the lower  $M$  bits are kept unchanged, while the top  $N - M$  bits are 0. To sign extend from  $M$  to  $N$ , the lower  $M$  bits are again unchanged, but the top  $N - M$  bits are filled with the bit value of the bit at location  $M - 1$ .

What is the purpose of sign extension, and when is zero extension used instead of sign extension in MIPS?

Sign extension preserves the sign of a data type undergoing extension. The instructions `andi`, `ori`, and `xori` specifically zero extend their immediate values. The instructions `lbu` and `lhu` specifically zero extend their data loaded from memory. Also, jump addresses and shift amounts are not sign-extended in MIPS but are not exactly zero extended either.