

# # Regular Expression

• Priority of operators of RE :

( ) ,  $a^*$  ,  $a^+$  ,  $\cdot$  ,  $+$

• Definition :

1.) RE is said to be valid iff it can be derived from RE by a finite no. of application of the rule  $x^*$ ,  $x^+$ ,  $x_1 \cdot x_2$ ,  $x_1 + x_2$ .

2.) If  $\Sigma$  is a given alphabet then,  $\Phi, \Sigma^+$ ,  $a \in \Sigma$  are REs.

$$\textcircled{1} a + \phi = a$$

$$\textcircled{2} a + \epsilon = \{a, \epsilon\}$$

$$\textcircled{3} u = a + b, d(u) = \{a, b\}$$

$$\textcircled{4} u = a \cdot b, d(u) = \{a, b\}$$

$$\textcircled{5} u = a + b + c, d(u) = \{a, b, c\}$$

$$\textcircled{6} u = (ab + a) \cdot b, d(u) = \{a, b, ab\}$$

$$\textcircled{7} u = a^+ , d(u) = \{a, aa, aaa, \dots\}$$

$$\textcircled{8} u = a^* , d(u) = \{\epsilon, a, aa, aaa, \dots\}$$

$$\textcircled{9} u = (a + ba)(b + a), d(u) = \{a, b, aa, bab, baa\}$$

$$\textcircled{10} u = (a + \epsilon)(\overbrace{b + a}^b), d(u) = \{ab, b\}$$

$$\textcircled{11} u = (a + b)(a + b), d(u) = \{aa, ab, ba, bb\}$$

$$\textcircled{12} u = (a + b)^* , d(u) = \{a, b\}$$

$\phi = \text{null}$   
 $\epsilon = \text{string of length 1}$

$= a^*$

$\{a, aa, aaa, \dots\}$

$R \rightarrow R^*$

$R_1 \Rightarrow R_2$

$R_1 \Rightarrow R_2$

$R_1 \Rightarrow R_2$

$R_1 \Rightarrow R_2$

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$$\begin{aligned} (ab)^3 &= (ababab) \\ (ab)^2 &= (ab)(ab) \\ (ab)^1 &= ab \\ (ab)^0 &= 1 \end{aligned}$$

(15)  $\mu = (ab)^* = \{ \varepsilon, ab, abab, ababab, \dots \}$

(14)  $\mu = a^* \cdot a^* = a^*$   
 $\{ \varepsilon, aa, aaaa, \dots \}$

(13)  $\mu = (a+b)^* \cdot (a+b)^+$   
 $\{ \varepsilon, \dots \} \cdot \{ a+b, (a+b)^2, (a+b)^3, \dots \}$   
 $= (a+b)^+$

$\therefore \{ \varepsilon, ab, abab, \dots \}$

$\{ \dots, \cancel{abab}, \dots \} = (a+b)^3$

$\{ aa, ab, ba, bb \} = (a+b)^2$

$\{ a, b \} = (a+b)^1$

$3 = (a+b)^0 \Rightarrow \mu$

(3)

eg 1 :  $u_1 = a^*$

$u_2 = a^* + (aa)^*$

(a)  $L(u_1) \subseteq L(u_2)$

(b)  $L(u_1) \supseteq L(u_2)$

✓ (c)  $L(u_1) = L(u_2)$

(d)  $L(u_1) \neq L(u_2)$

Ans 1

same language can be generated by one or more regular expressions.

ques 2:

$\phi$  = empty set / null set . = 0

$\epsilon$  = empty string.

(a)  $u = \epsilon^*$ ,  $L(u) = \{\epsilon^0, \epsilon^1, \epsilon^2, \dots\}$   
 $= \{\epsilon, \epsilon, \epsilon, \dots\}$

to remove  
 duplicacy  
 in set  
 $= \{\epsilon\}$

(b)  $u = \epsilon^+$ ,  $L(u) = \{\epsilon^1, \epsilon^2, \dots\} = \{\epsilon\}$

(5)

$$(c) \quad \Sigma = \phi^*, \Delta(\Sigma) = \{\phi^0, \phi^1, \phi^2, \dots\} = \{\epsilon\}$$

$$\{\epsilon, x, x, \dots\}$$

$$A = \{a, b\}$$

eg: -  $A^0 = \{\epsilon\}$  → string of length 0

$$A^1 = \{a, b\}$$

$$A^2 = \{aa, ab, ba, bb\}$$

$$(d) \quad \Sigma = \phi^+, \Delta(\Sigma) = \{\phi^1, \phi^2, \dots\} = \phi$$

Questions :

$$(1) \quad \Sigma^+ \cup \Sigma^* = \Sigma^*$$

$$(2) \quad \Sigma^+ \cap \Sigma^* = \Sigma^+$$

$$(3) \quad \Sigma^* \cdot \Sigma^+ = \Sigma^+$$

$$(4) \quad (\Sigma^*)^* = \{\epsilon^0, \epsilon^1, \epsilon^2, \dots\}^* = \{\epsilon, \epsilon, \epsilon\epsilon, \dots\}^*$$

$$\Rightarrow \{\epsilon^*, \epsilon^*, \epsilon\epsilon^*, \dots\} = \{\epsilon, \epsilon, \epsilon\epsilon, \dots\}$$

$$\Rightarrow \Sigma^*$$

$$(5) \quad (\Sigma^*)^+ = \Sigma^*$$



$$(6) (u^+)^* = u^*$$

$$(7) ((u^*)^+)^* \cdot u^+ = u^*$$

$$(8) (a+b)^* = (a^*+b)^* = (a+b^*)^*$$

$$(9) (a+b)^* = (a^*+b^*)^*$$

$$(10) (a+b)^* \neq (a \cdot b)^* \quad \begin{matrix} \{a\} \checkmark \\ \{b\} \checkmark \end{matrix} \quad \begin{matrix} \{a \cdot b\} \times \\ \{a\} \times \end{matrix}$$

$$(11) (a+b)^* \neq (a^* \cdot b)^* \neq (a \cdot b^*)^* = (a^* b^*)^* \quad \begin{matrix} \{a\} \checkmark \\ \{b\} \checkmark \end{matrix} \quad \begin{matrix} \{a\} \times \\ \{b\} \checkmark \end{matrix} \quad \begin{matrix} \{a \cdot b\} \times \\ \{b\} \times \end{matrix}$$

\* if you want ba from  $(a^* b^*)^*$

$$(a^* b^*)^2 \Rightarrow \begin{matrix} a^* & b^* & a^* & b^* \\ \downarrow & \downarrow & \downarrow & \downarrow \\ \epsilon & b & a & \epsilon \end{matrix}$$

missions

for rep  
algebraic

$\phi$   
symbol

regular

expres  
: R

: R, 1

$\epsilon$  :

$a^* = \{$

$a^+ = \{$

$a^+ \cdot a =$

# Regul# How to design RE from language (7)

followi

① start with ab:

$\{ab, aba, abb, abba, \dots\}$

$$\Rightarrow ab(a+b)^*$$

② start with bba:

$$\Rightarrow bba(a+b)^*$$

③ ends with abb:

$$\Rightarrow (a+b)^*abb$$

④ contains substring abb.

$$\Rightarrow (a+b)^*aab(a+b)^*$$

⑤ starts and ends with a

• if a is starting and ending

• if in starting + ending

$$\left. \begin{array}{l} a \\ a(a+b)^*a \end{array} \right\} \Rightarrow a + a(a+b)^*a$$

⑥ start and ends with same symbol  
 $\Rightarrow a + a(a+b)^*a + b + b(a+b)^*b.$

⑦ starts and end with different symbol.

$$\Rightarrow a(a+b)^*b + b(a+b)^*a.$$

⑧

$$|W| = 3$$

$$(a+b)(a+b)(a+b)$$

⑨

$$|W| \geq 3$$

$$(a+b)^3(a+b)^*$$

⑩

$$|W| \leq 3$$

$$\Rightarrow \{ \underbrace{\epsilon}_0, \underbrace{a, b}_1, aa, bb, ba, ab, \dots \}$$

$$\Rightarrow \underbrace{\epsilon}_0 + \underbrace{(a+b)}_1 + \underbrace{(a+b)^2}_2 + \underbrace{(a+b)^3}_3$$

$$\Rightarrow (a+b+\epsilon)$$



$$(1) |w|_a = 2$$

$$b^* a b^* a b^*$$

$$(2) |w|_a \geq 2$$

$$(a+b)^* a (a+b)^* a (a+b)^*$$

$$(3) |w|_a \leq 2$$

$$\epsilon + b^* a b^* + b^* a b^* a b^*$$

$$(4) \text{ Symbol from left end is } 1$$

$$(a+b)^2 b(a+b)^*$$

$$(5) |w| = 0 \pmod{3}$$

$$[(a+b)^3]^* \rightarrow \text{for } 0$$

in contribute  
of length 3

0  
3  
6  
9  
...

$$(6) |w| = 2 \pmod{3}$$

$$(a+b)^2 [(a+b)^3]^*$$

⑦  $|w|_b = 0 \pmod{2}$

$b =$   
0  
2  
4  
6  
8  
...

$a^* + (a^* b a^* b a^*)^*$

⑧  $|w|_a = 1 \pmod{3}$

✓  $b^* a b^* (b^* a b^* a b^* a b^*)^*$

⑨  $|w|_b = 2 \pmod{3}$

bb

⑦  $a^* (bb)^* a^*$   
↓ ↓  
ε 1

bb bb bb

⑧  $(\overbrace{b^* a (aaa)^* b^*}^{\text{consecutive}})^*$

4  
7

10

13

$a^* \underline{bba}^* \cdot (a^* b a^* b a^* b a^*)^*$