

## \* Greedy Method:- (is a Technique)

- Simplest & straight forward approach.
- The decision is taken on basis of current available info.<sup>n</sup>. without worrying about the effect of current decision in future.
- feasible sol.<sup>n</sup> that may or may not be optimal.

feasible sol.<sup>n</sup>  $\rightarrow$  any subset that satisfy the condition.

optimal sol.<sup>n</sup>  $\Rightarrow$  best & most favorable sol.<sup>n</sup>.

### Characteristic & Features:-

- ① To construct the sol.<sup>n</sup> in an optimal way, Algorithm maintains 2 sets:-
  - one contains chosen items &
  - other " rejected items.
- ② Greedy algorithm make good local choices.
  - an optimal sol.<sup>n</sup>
  - feasible sol.<sup>n</sup>

## Components of Greedy Algorithm :-

- ① A Candidate Set :-  
A sol.<sup>n</sup> is created from this set.
- ② A Selection fun.<sup>c</sup> :-  
used to choose the best candidate to be added to the sol.<sup>n</sup>.
- ③ A feasibility fun.<sup>c</sup> :-  
used to determine whether a candidate can be used to contribute to the sol.<sup>n</sup>.
- ④ A objective fun.<sup>c</sup> :-  
used to assign value to a solution or partial sol.<sup>n</sup>.
- ⑤ A solution fun.<sup>c</sup> :-  
used to indicate whether a complete sol.<sup>n</sup> has been reached.

## Applications of Greedy :-

- ① finding shortest path
- ② finding Minimum Spanning Tree
- ③ Job sequencing with deadline / activity-selection problem.
- ④ fractional knapsack problem

⑤ Huffman coding.

Pseudocode for Greedy Algorithm:-

```
Algorithm Greedy(a, n)
{
    solution := 0;
    for i = 1 to n do
    {
        x := Select(a);
        if feasible(solution, x) then
            solution := Union(solution, x);
    }
    return solution;
}
```

for those Problem which are divisible.  
Example :- Fractional Knapsack using Greedy Method.

$n=7$   
 $W=15$

Object (O) $\Rightarrow$	1	2	3	4	5	6	7
Profit (P) $\Rightarrow$	10	5	15	7	6	18	3
Weights (W) $\Rightarrow$	2	(3)	5	7	1	4	1

Right Selection Procedure for the object.

$\Rightarrow \frac{P}{W} \Rightarrow$

5	(1.3)	3	1	6	4.5	3
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$(0 \leq x \leq 1)$

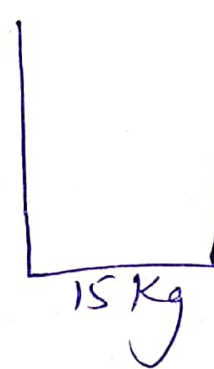
$x \left( \begin{matrix} 1 & 2/3 & 1 & 0 & 1 & 1 & 1 \\ x_1 & x_2 & x_3 & x_4 & x_5 & x_6 & x_7 \end{matrix} \right)$

$\uparrow$   
 4th object Not Included  $\therefore$  zero

How to include object/ITEM in Bag to get Max. Profit:-

Idea

- 1) Those Element/object give take that object first.  
 E.g. - (18) Profit take Complete object



$$\begin{aligned} 15 - 1 &= 14 \\ 14 - 2 &= 12 \\ 12 - 4 &= 8 \\ 8 - 5 &= 3 \\ 3 - 1 &= 2 \\ 2 - 2 &= 0 \end{aligned}$$

(x5)  
 (x6)  
 (x4)  
 (x3)  
 (x7)  
 (x2)  
 Take only 2 Kg

Idea

- 2) Or we may fill the bag with smaller things. So that I can add more & more things in the bag & can make more Profit.  
 i.e. More Thing we carry more will be the Profit

Right Method (3)

$\rightarrow \frac{P}{W}$  Then add item accordingly.

$$\begin{aligned} \sum x_i w_i &= 1 \times 2 + \frac{2}{3} \times 3 + 1 \times 5 + 0 \times 7 + 1 \times 1 + 1 \times 4 + 1 \times 1 \\ &= 2 + 2 + 5 + 0 + 1 + 4 + 1 = 15 \text{ Kg} \end{aligned}$$

$$\begin{aligned} \sum x_i p_i &= 1 \times 10 + \frac{2}{3} \times 5 + 1 \times 15 + 1 \times 6 + 1 \times 18 + 1 \times 3 \\ &= 10 + 2 \times 1.3 + 15 + 6 + 18 + 3 \end{aligned}$$

~~10 + 1.3~~  
 $= 10 + 2 \times 1.3 + 15 + 6 + 18 + 3$   
 $= 54.6$  Total Profit



## \* Fractional Knapsack (Array V, Array W, int W)

- 1) for  $i=1$  to  $\text{Size}(V)$
- 2) do  $P[i] = V[i]/W[i]$
- 3) Sort - Descending ( $P$ )
- 4)  $i \leftarrow 1$
- 5) while ( $W > 0$ )
- 6) do amount =  $\min(W, W[i])$
- 7) solution  $[i] = \text{amount}$
- 8)  $W = W - \text{amount}$
- 9)  $i \leftarrow i+1$
- 10 return solution.

Constraint:-

$$\sum x_i w_i \leq m$$

$\nwarrow$   
Capacity  
of  
Bag

Objective:-

$$\max \sum x_i P_i$$

$O(n \log n)$

Q= Consider 5 items along their respective weights and values.

$$I = (I_1, I_2, I_3, I_4, I_5)$$

$$W = (5, 10, 20, 30, 40)$$

$$V = (30, 20, 100, 90, 160)$$

The Capacity of Knapsack  $W=60$ . Find the solution to the fractional knapsack problem.

Sol.<sup>n</sup>  $\Rightarrow$

⇒ Initially,

Item	(kg) weight $w_i$	(Price) $v_i$
$I_1$	5	30
$I_2$	10	20
$I_3$	20	100
$I_4$	30	90
$I_5$	40	160

Taking value per weight ratio i.e.  $p_i = v_i/w_i$

Item	$w_i$	$v_i$	$p_i = v_i/w_i$
$I_1$	5	30	6.0
$I_2$	10	20	2.0
$I_3$	20	100	5.0
$I_4$	30	90	3.0
$I_5$	40	160	4.0

Now arrange the value of ' $p_i$ ' in decreasing order.

Item	$w_i$	$v_i$	$p_i = v_i/w_i$
$I_1$	5	30	6.0
$I_3$	20	100	5.0
$I_5$	40	160	4.0
$I_4$	30	90	3.0
$I_2$	10	20	2.0

Now, fill the knapsack acc. to the decreasing value of  $p_i$ .

First, we choose item  $I_1$  whose weight is 5, then choose item  $I_3$  whose weight is 20. Now the total weight

in knapsack is  $5+20=25$ .

Now, the next item is I5 & its weight is 40, but we want only 35. So we choose fractional part of it i.e.

35
20
5

) (60)

The value of fractional part of I5 is

$$\frac{160}{40} \times 35 = 140$$

The max. value is  $\Rightarrow 30 + 100 + 140$

$$= 270 \text{ Ans.}$$