

(1)

Regular Expressions

- Regular Expressions are used for representing certain sets of strings in an algebraic fashion.

Rules

→ $[a, b, c, \dots, \wedge, \phi]$

- Any terminal symbol i.e symbols ϵ, Σ , including \wedge and ϕ are regular expressions.
 $\{\epsilon, \Sigma, \wedge, \phi\}$
- The union of two regular expressions is also a regular expression
 $R_1, R_2 : R_1 \cup R_2 \Rightarrow (R_1 + R_2)$
- Concatenation is also RE. $: R_1, R_2 \Rightarrow (R_1 \cdot R_2)$
- Closure of RE is RE $: R \rightarrow R^*$

$$a^* = \{\epsilon, a, aa, aaa, \dots\}$$

$$a^+ = \{a, aa, aaa, \dots\}$$

$$a^+ \cdot a = a^*$$

NOVEN'S THEOREM ($FA \rightarrow RE$) (over Σ)

examples :

describe the following sets as Regular Expressions:

1) $\{0, 1, 2\}$: 0 or 1 or 2

$$R = 0 + 1 + 2$$

2) $\{\Lambda, ab\}$

$$R = \underbrace{\Lambda}_{\substack{\uparrow \\ \text{this is present}}} ab$$

*)

3) $\{abb, a, b, bba\}$: abb or a or b or bba

$$R = abb + a + b + bba$$

4) $\{\Lambda, 0, 00, 000, \dots\}$: $R = 0^*$ all strings that can be formed with 0 including empty symbol.

5) $\{1, 11, 111, 1111, \dots\}$: $R = 1^+$

$[P]^2$

Properties of Regular Expression

$$1) \phi + R = R$$

ϕ = empty set $\therefore \phi \cup R = R$

$$2) \phi.R + R.\phi = \phi$$

$$3) \epsilon R \doteq R \epsilon = R$$

$$(\wedge = \epsilon)$$

$$4) \epsilon^* = \epsilon \text{ and } \phi^* = \epsilon$$

$$5) R + R = R$$

$$6) R^*.R^* = R^*$$

$$7) RR^* = R^*R$$

$$8) (R^*)^* = R^*$$

$$9) \epsilon + RR^* = \epsilon + R^*R = R^*$$

$$10) (PQ^*)^* = P(QP)^*$$

$$11) (P+Q)^* = (P^*Q^*)^* = (P^*+Q^*)^*$$

$$12) R(P+Q) = RP + RQ \text{ and } (P+Q)R = PR + QR$$

THEOREM (Pumping Lemma)

ver Σ ,

design RE for following language over.

1) language accepting string of length exactly 2.

$$L_1 = \{aa, ab, ba, bb\}$$

$$R_1 = aa + ab + ba + bb$$

$$a(a+b) + b(a+b) \Rightarrow \boxed{(a+b)(a+b)}$$

2) language accepting of length atleast 2.

$$L_2 = \{aa, ab, ba, bb, aaa, bbb, abab, \dots\}$$

$$R_2 = \underbrace{(a+b)(a+b)}_{\text{atleast 2}} (a+b)^*$$

↑
anything more
than length 2.

ϵ^*

abba

3) language accepting string of length atleast 2.

$$L_3 = \{\epsilon, a, b, aa, bb, ab, ba\}$$

$$R_3 = \epsilon + a + b + aa + ab + bb + ba$$

$$= (\epsilon + a + b)(\epsilon + a + b)$$

1+RP

①

$$\frac{1}{2} \quad \frac{1}{1} \quad \frac{a}{a, b, c \rightarrow 0, 1, 2}$$

ab ba

$$(\check{a} + \check{b} + c)^* \underline{a (a + b + c)^* b (a + b + c)^*} + (a + b + c)^* b ($$

$$(0+1)^* \underline{1 \textcircled{01} 01} \text{ --- --- ---}$$

$$(0+1)^* \underline{1 (0+1) (0+1)^a} \text{ --- ---}$$

$$(0+1)^9$$

101

Σ^1
 Σ^2
 Σ^3

$$\underline{(0^* 1^* 00)^* 0^* 1^*}$$

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$$

$$\underline{\epsilon, 0, 000} \quad \epsilon, 0.$$

$$1 \quad \underline{1101} \quad \underline{0100}$$

$$0^* = \{ \epsilon \}$$

$$\{ \underline{a}^n \underline{b^{2m+1}} : n \geq 0 \text{ \& } m \geq 0 \}$$

P1

$$(0^* 1 \textcircled{00})^*$$

$$\epsilon, 0100, 01000100, 0$$

$$\begin{matrix} a^0 b^{2 \times 0 + 1} \\ a^0 b^1 \\ a^1 b^3 \\ a^2 b^5 \\ a^3 b^7 \\ a^4 b^9 \end{matrix}$$

$$\begin{matrix} (aa)^* & (bb)^* b. \\ 0, aa, aaaa & \underline{bb} b \\ & \underline{bbbb} b. \\ & \underline{bbbbbb} b. \end{matrix}$$

ab ba

$(a+b)^*$

a, a, b, aa, bbb, ab

r^*R

$P +$