



Worksheet 6

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Branch: CSE **Section/Group:DWWC-43**

Semester:5 Date of

Performance: 10-01-

2023

Subject Name: Data Subject Code:

Structure

1. Aim/Overview of the practical:

i) Consider an infinite full binary tree (each node has two children except the leaf nodes) defined as follows. For a node labelled **v** its left child will be labelled **2*v** and its right child will be labelled **2*v+1**. The root is labelled as **1**.

You are given N queries of the form i j. For each query, you have to print the length of the shortest path between node labelled i and node labelled j.

You are given a rooted tree with NN nodes (numbered 11 through NN); node 11 is the root. Each node has a value; let's denote the value of node ii by A_iAi .

You may perform the following operation any number of times (including zero): choose any node which still exists in the tree and remove the whole subtree of this node including itself.

Let's define the *profit* as the sum of values of all nodes which currently exist in the tree minus $X \cdot k$, where k denotes the number of times this operation was performed. Find the maximum possible profit.

iii) You are given an undirected tree having NN nodes. The color of each node is either Black or White.

In one step, you can choose a node uu, toggle the color of the chosen node uu, as well as toggle the colors of all the nodes vv, such that $\{u, v\}\{u, v\}$ is an edge in the tree.

Can you find out the minimum number of steps required to make the color of all the nodes black?

Note: Since the input is large, prefer using fast input methods.



You are given a family tree with NN members numbered 1 ... N1...N. Every person ii, except for the founder of the family (root) has a parent denoted by P[i]P[i]. P[root]P[root] = -1 by definition. Person ii is a descendant of person jj if ii belongs to the subtree rooted at jj.

The net worth (adjusted for inflation) is given for all the members of the family. Your task is to find the maximum difference in net worth's of two members where one member is a descendant of the other.

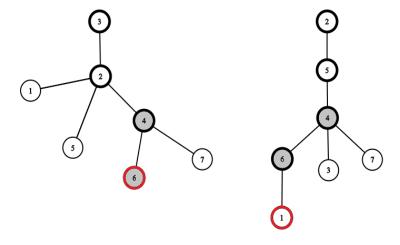
You can assume that no two members of this family married each other. So it remains a "family tree".

Ayush has two rooted trees. Each of them contains NN vertices (numbered 11 through NN); let's denote a vertex vv in tree 11 by v_1v_1 and a corresponding vertex vv in tree 22 by v_2v_2 .

For each tree $t\bar{t}$, any vertex on the path from the root to a vertex $v_t\bar{v}t$ (including $v_t\bar{v}t$) is an ancestor of $v_t\bar{v}t$ and the depth of $v_t\bar{v}t$ is the distance of $v_t\bar{v}t$ from the root of tree $t\bar{t}t$.

For any integer ww (1 \le W \le N1 \leq w \leq N) and vertices u_1u₁ and v_2v₂, ww is said to be a *common ancestor* of u_1u₁ and v_2v₂ if:

- in tree 11, w_1w_1 is an ancestor of u_1u_1
- in tree 22, w_2w_2 is an ancestor of v_2v_2
- the depths of vertices w_1w_1 and w_2w_2 (from their respective roots) are the same



For example, in the figure above, the pair $(u_1, v_2) = (6, 1)(u_1, v_2) = (6, 1)$ has two common ancestors — 44 and 66. Although vertex 22 is an ancestor for each of them, the depths of 2_{121} and 2_{222} are not the same.

Help Ayush find the maximum number of common ancestors for a pair of vertices, i.e. the largest integer cc such that there is a pair of vertices $(u_1, v_2)(u_1, v_2)$ (possibly u = vu = v) with cc common ancestors.

In a rooted tree, the (or LCA for short) of two vertices **u** and **v** is defined as the lowest vertex that is ancestor of both that two vertices.

Given a tree of N vertices, you need to answer the question of the form "r u v" which means if the root of the tree is at r then what is LCA of u and v.

Vii) Nearly 10^{9} + 7109+7 trillion kilometers far away from the Milkyway, there exists a parallel Universe of Light called the Cosmos of Solon where each planet is called a Temple. There are total NN temples. Unlike humans, the Solomons (Giant aliens of the Light World) can travel from one temple to another through Cosmic Bridges and the distance observed by them between any two temples is 00 units (Interesting? Well, they travel with more than the speed of light)



But the year 20202020, proved to be a curse for Solomons. This year, a Mega Star Urus struck their Universe and damaged many Cosmic Bridges. However the God of Light ensured that for any two distinct temples uu and vv there exists exactly one unique path through Cosmic Bridge(s), but the distance between two directly connected temples through a Bridge has now increased to 11 units.

To prevail the existence of this Universe, the God of Light carved himself into exactly NN Solomons which now rule the NN temples and are independent of each other. One by one, the i^{t} Solomon ($0 \le i \le N-1$) transmits Light pulses through the network of Cosmic Bridges and determines the sum of distances to the remaining N-1N-1 Solomons.

Can you help them find these quantities so they can reunite into God of Light again?

xiii) Abdullah has recently moved to Tolland. In Tolland, there are NN cities (numbered 11 through NN) connected by N-1N-1 bidirectional roads such that it is possible to visit any city from any other city. For each city ii, there is a toll value TL_iTLi . Abdullah lives in city XX. He has planned a tour of Tolland lasting for 2N2N days. For each ii (1 \le i \le $N1 \le i \le N$), on the 2i-12i-1-th day, he travels from city XX to city ii, and on the 2i2i-th day, he travels from city ii back to city XX. In Tolland, there are also special discount coupons. For each city cc and positive integer vv, there is a discount coupon for city cc with value vv. All discount coupons may be bought anywhere in Tolland; the cost of each discount coupon is twice its value. When travelling from any city ii to any (not necessarily different) city jj, Abdullah must pay a toll exactly once in each city on the path between them (including ii and jj); even if he passed through some cities before, he must pay the toll in these cities as well. For any city ii, the toll is TL_iTLi when not using any discount coupons; however, if Abdullah bought some discount coupons for city ii, he may use some or all of them to decrease this toll to \mathrm{max}\,\((0, TL_i-V))max(0, TLi-V), where VV is the sum of values of coupons for city ii used by Abdullah. Each coupon may be used at most once per day, but they do not expire when used and the same coupon may be used on multiple days. On each day, Abdullah may use any coupons he wants, but a coupon for city ii may only be used to decrease the toll in city ii.

Abdullah has a travel budget: on each day, he must not spend more than KK units of money on paying tolls. Before the start of the tour, he may choose to buy any number of discount coupons, for any cities and with any values (possibly different ones for different cities); the cost of coupons is not included in the budget for any day of the tour. Help Abdullah find the minimum total amount of money he must spend on discount coupons so that on each day during the tour, he will be able to stay within his budget by using the coupons he bought.

XIV) A tree is an undirected connected graph without any cycles. You are given a tree with **N** nodes and a parameter **K**. Each node in the node has a positive weight **W**[]. The tree needs to be partitioned into at most **K** subtrees by removing at most **K-1** edges such that the maximum weight of the resulting subtrees is as small as possible. Find that smallest value of the maximum subtree weight that can be obtained by removing at most **K-1** edges.

XV) The apocalyptic demon <u>Kali</u> is on a rampage. The ground shudders under his feet, trees shrivel and animals and birds scurry away from his path. In order to save the universe from devastation, all the devtas led by devraj <u>Indra</u> decide to meditate to please Lord <u>Vishnu</u> so that he appears in the form of <u>Kalki</u>, his last avatar, and kill the demon.

Each devta can sit in meditation at a particular place on <u>Swarglok</u> (the heaven). Considering Swarglok as a 2-dimensional plane, the position of each devta is a fixed point with integer coordinates. The positions of all the devtas are distinct from each other.



While meditating the devtas connect to each other by means of astral projections - metaphysical threads that connect one devta to another. They must do so in such a way to satisfy two criteria:

- Each devta must be able to connect to every other devta by following a path of astral projections.
- No subgroup of devtas may be connected by a cycle of astral projections.

In simple terms, the astral projections must form a tree connecting all devtas.

What needs to be optimized?

Once the devtas have taken their positions and are connected as a single group to start meditation, a *ring of influence* is formed around each devta. This ring is centered at the position of the devta, and has a radius equal to the Euclidean distance from this devta to the furthest devta directly connected via an astral projection.

Since different devtas have their own supernatural powers (Indra is the rain god, Agni is the fire god, Varuna is the water god, Vayu is the air god, etc), the influence ring of each devta has an adverse effect on all other devtas which lie within or on this ring. In other words, the efficiency of a devta to perform meditation decreases as the number of influence rings that include him increases. For each devta D_i , define C_i as the number of influence rings that include or touch D_i (including their own). Now devraj Indra wants the devtas to connect in such a manner so as to minimize the maximum value of C_i over all devtas. So he has sent a message to Bhulok (the earth) and he needs help from the best programmers on the planet.



2. <u>Steps for experiment/practical/Code:</u>

```
i)
#include <iostream>
using namespace std;
int main() {
     int t;
     cin>>t;
     while(t--) {
       int L, R;
       cin>>L>>R;
       int cnt=0;
       while(L!=R) {
         if(L>R){
            L=L/2;;
          }
          else {
            R=R/2;
          }
          cnt++;
       cout<<cnt<<endl;
     }
     return 0;
}
```

```
#include "bits/stdc++.h"
using namespace std;
typedef long long ll;
typedef long double ld;
typedef pair<ll,ll> pll;
typedef unsigned long long ull;
#define all(x) (x).begin(), (x).end()
const int MOD = 1000000007; // 1000000009, 1000000023, 1000000007, 998244353
const 11 INF = 11(1e18);
const ld PI = 4*atan((ld)1);
const char nl = '\n';
int main() {
 ios_base::sync_with_stdio(false);
 cin.tie(NULL);
 cout.precision(10);
 11 tt;
 cin >> tt;
 while(tt--){
  11 \text{ n,x; cin} >> n >> x;
  vector < vector < ll > g(n + 1, vector < ll > ());
  vector<ll>val(n);
  for(auto &x:val)cin >> x;
  for(11 i=0 ; i< n-1 ; i++){
   ll u,v; cin >> u >> v;
   g[u].push_back(v);
   g[v].push_back(u);
  vector < ll > dp(n+1);
```



```
auto dfs = [&](ll s, ll p , auto && dfs)->ll{
    ll res = val[s-1];
    for(auto x:g[s]){
        if(x != p)res += dfs(x , s , dfs);
    }
    return dp[s] = max(res , -x);
    };
    cout << dfs(1 , -1 , dfs) << nl;
}</pre>
```







iii)

```
#include "bits/stdc++.h"
using namespace std;
#define fast ios_base::sync_with_stdio(false);cin.tie(0);cout.tie(0);
#define tt int ct;cin>>ct;while(ct--)
#define MAX 100005
const int mod = 998244353;
typedef long long ll;
int n,a[MAX];
vector<int> ad[MAX];
int dp[MAX][2][2];
void dfs(int u,int par){
for(auto it : ad[u]){
if(it != par){
dfs(it,u);
for(int i=0; i<2; i++){
for(int j=0; j<2; j++){
int odd = a[u]^i;
int cc = j;
int dp1[2],dp2[2];
memset(dp2,0x3f,sizeof(dp2));
dp2[0] = 0;
for(auto v : ad[u]){
if(v != par)
swap(dp1[0],dp2[0]);
swap(dp1[1],dp2[1]);
memset(dp2,0x3f,sizeof(dp2));
dp2[0] = min(dp2[0], dp1[0] + dp[v][cc][0]);
dp2[1] = min(dp2[1],dp1[1]+dp[v][cc][0]);
dp2[0] = min(dp2[0], dp1[1] + dp[v][cc][1]);
dp2[1] = min(dp2[1],dp1[0]+dp[v][cc][1]);
dp[u][i][j] = dp2[odd]+j;
```



```
} }
return;
void solve(){
cin>>n;
for(int i=0;i<n;i++){
cin>>a[i];
for(int i=0;i<n;i++){
ad[i].clear();
for(int i=0;i< n-1;i++){
int u,v; cin>>u>>v; u--;v--;
ad[u].emplace_back(v);
ad[v].emplace_back(u);
dfs(0,-1);
int ans = min(dp[0][0][0],dp[0][0][1]);
if(ans > n){
printf("-1\n");
else{
printf("%d\n",ans);
return;
int32_t main() {
fast
#ifndef ONLINE_JUDGE
freopen("input.txt", "r", stdin);
freopen("output.txt","w",stdout);
#endif
tt{
solve();
return 0;
```



```
#include <iostream>
#include <bits/stdc++.h>
#include <cmath>
#include <algorithm>
#include <vector>
#include <deque>
#include <set>
#include <stack>
#include <string>
#include <map>
#include <queue>
//#pragma comment(linker, "/stack:200000000")
//#pragma GCC optimize("Ofast")
//#pragma GCC target("sse,sse2,sse3,sse4,popcnt,abm,mmx,avx,avx2,tune=native")
using namespace std;
#define int long long
#define 11 long long
#define vi vector<long long>
#define pb push_back
#define sz(s) (int)s.size()
#define all(v) v.begin(), v.end()
#define show(a) cerr << #a <<" -> "<< a <<"\n"
#define pp pair<int,int>
#define FF first
#define SS second
#define endl "\n"
#define ld long double
//typedef long double ld;
const int N = 1e6 + 2, N3 = 1e3 + 6, \inf = 1e9 + 7, LOG = 20;
map<char, int> md{{'N', 0}, {'E', 1}, {'S', 2}, {'W', 3}};
const int dx[] = \{-1, 0, 1, 0\};
const int dy[] = \{0, 1, 0, -1\};
string stepDir = "RDLU";
int n, w[N], dp[N], root, mn[N], mx[N], ans;
vi g[N];
```



```
void dfs(int v, int p) {
  mn[v] = w[v], mx[v] = w[v];
  for (auto to : g[v]) {
     if (to == p) continue;
     dfs(to, p);
     mn[v] = min(mn[v], mn[to]);
     mx[v] = max(mx[v], mx[to]);
   }
}
main() {
       ios_base::sync_with_stdio(0);
       cin.tie(0);
       cout.tie(0);
       freopen("triangles.in", "r", stdin);
//
       freopen("triangles.out", "w", stdout);
//
  cin >> n;
  for (int i = 1; i \le n; i++) {
     cin >> w[i];
  for (int i = 1; i \le n; i++) {
     int p;
     cin >> p;
     if (p == -1) {
        root = i;
        continue;
     g[p].pb(i);
  dfs(root, 0);
  for (int i = 1; i \le n; i++) {
     ans = max(ans, max(abs(w[i] - mn[i]), abs(w[i] - mx[i])));
   }
  cout << ans;
       return 0;
}
```







```
v)
```

```
#include <bits/stdc++.h>
using namespace std;
const int MX = 1e6 + 10;
int seg[4 * MX], lazy[4 * MX], beg[MX], fin[MX], dep[MX], cnt;
vector<int> adj[2][MX];
void updateNode(int idx, int v) {
  seg[idx] += v;
  lazy[idx] += v;
  return;
}
void shift(int idx, int st, int ed) {
  int lft = 2 * idx, rgt = lft + 1;
  if (lazy[idx]) {
     updateNode(lft, lazy[idx]);
     updateNode(rgt, lazy[idx]);
     lazy[idx] = 0;
  return;
}
void update(int s, int e, int v, int idx = 1, int st = 0, int ed = cnt - 1) {
  if (s > e \parallel e < st \parallel s > ed) return;
  if (s == st \&\& e == ed) {
     updateNode(idx, v);
```

```
return;
   }
  int lft = 2 * idx, rgt = lft + 1, mid = (st + ed) / 2;
  shift(idx, st, ed);
  update(s, min(e, mid), v, lft, st, mid), update(max(s, mid + 1), e, v, rgt, mid + 1, ed);
  seg[idx] = max(seg[lft], seg[rgt]);
  return;
}
void dfs0(int u, int d) {
  beg[u] = cnt++;
  dep[u] = d;
  for (auto v : adj[0][u]) dfs0(v, d + 1);
  fin[u] = cnt - 1;
  return;
}
void dfs1(int u, int d, int &ans) {
  if (dep[u] == d) update(beg[u], fin[u], 1);
  ans = max(ans, seg[1]);
  for (auto v : adj[1][u]) {
     dfs1(v, d + 1, ans);
   }
  if (dep[u] == d) update(beg[u], fin[u], -1);
  return;
}
int main() {
  ios::sync_with_stdio(false);
  cin.tie(0);
  cout.tie(0);
  int t;
  cin >> t;
```



```
while (t--) {
  cnt = 0;
  int n;
  cin >> n;
  for (int i = 0; i \le 4 * n; i++) {
     seg[i] = lazy[i] = 0;
     if (i \le n) {
        adj[0][i].clear();
        adj[1][i].clear();
   }
  for (int j = 0; j < 2; j++) {
     for (int i = 1; i \le n; i++) {
        int p;
        cin >> p;
        adj[j][p != -1 ? p : 0].push_back(i);
   }
  dfs0(0, 0);
  int ans = 0;
  dfs1(0, 0, ans);
  cout << ans - 1 << endl;
return 0;
```

```
#include<br/>
dits/stdc++.h>
using namespace std;
int sp[200100][18];
int depth[200100];int n;
vector <int> g[200100];
void dfs(int u,int pre){
  sp[u][0] = pre;
  for (int j = 1; j < 18; j ++){
     sp[u][j] = sp[sp[u][j-1]][j-1];
  depth[u] = depth[pre] + 1;
  for (auto v:g[u]){
     if (v == pre)continue;
     dfs(v,u);
   }
int lca(int x,int y){
  if (depth[x] < depth[y])swap(x,y);
  for (int j = 17; j >= 0; j --){
     if (depth[x] - (1 << j) >= depth[y]){
        x = sp[x][j];
     }
  for (int j = 17; j \ge 0; j --){
     if (sp[x][j] != sp[y][j]){
        x = sp[x][j];
        y = sp[y][j];
     }
  if (x != y){
     x = sp[x][0];
  return x;
int best(int x,int y){
  if (depth[x] > depth[y]) return x;
  return y;
int main(){
  ios_base::sync_with_stdio(0);cin.tie(nullptr);cout.tie(nullptr);
```

```
cin>>n;
  for (int i = 1; i < n; i ++){
     int u,v;
     cin>>u>>v;
     g[u].push_back(v);
     g[v].push_back(u);
  dfs(1,1);
  int q;
  cin>>q;
  while (q--){
     int r,u,v;
     cin>>r>>u>>v;
     cout << best(best(lca(r,u),lca(r,v)),lca(u,v)) << '\n';
  return 0;
vii)
#include <iostream>
#include <bits/stdc++.h>
using namespace std;
void bfs(int temp,int n,vector<list<int>>adj)
  vector <bool> visited(n,false);
  int c = 0;
  vector \langle int \rangle count(n,0);
  list<int> q;
  visited[temp] = true;
  count[temp]=0;
  q.push_back(temp);
  while(q.empty()==false)
     int sr = q.front();
     c = count[sr]+1;
     q.pop_front();
```



```
for(auto adjc : adj[sr])
        if(!visited[adjc])
          visited[adjc] = true;
          count[adjc] = c;
          q.push_back(adjc);
     }
  int s = 0;
  for(int i=0;i<n;i++)
     s = s + count[i];
  cout<<s<" ";
}
int main() {
      // your code goes here
       vector<list<int>>adj;
       int n;
       cin>>n;
       adj.resize(n);
       for(int i=0;i<n-1;i++)
         int v,u;
         cin>>v>>u;
         adj[v].push_back(u);
         adj[u].push_back(v);
       }
       for(int i=0;i<n;i++)
       bfs(i,n,adj);
      return 0;
}
```

xiii)

```
#include <bits/stdc++.h>
#define ll long long
using namespace std;
inline int read() {
int x = 0, f = 1; char ch = getchar();
while (ch < 0' \parallel ch > 9') { if (ch == '-') f = -1; ch = getchar(); }
while (ch >= '0' && ch <= '9') { x = x * 10 + ch - 48; ch = getchar(); }
return x * f;
}
inline ll readl() {
11 x = 0, f = 1; char ch = getchar();
while (ch < 0' \parallel ch > 9') { if (ch == '-') f = -1; ch = getchar(); }
while (ch >= '0' && ch <= '9') { x = x * 10 + ch - 48; ch = getchar(); }
return x * f:
}
const int N = 1e4 + 10;
vector<int> G[N];
bool vis[N];
ll dis[N], a[N], ans, k, n, p;
ll dfs1(int u, int fa) {
11 \text{ res} = 0;
for (auto v : G[u]) {
if (v != fa) res = max(res, dfs1(v, u));
return dis[u] = res + a[u];
void dfs2(int u, int fa) {
if (dis[u] \le k) return;
if (dis[u] - a[u] \le k) {
ans += dis[u] - k;
return;
}
ans += a[u];
for (auto v : G[u]) {
if (v != fa) dfs2(v, u);
}}
```



```
int main() {
int T = read();
while (T--) {
n = read(), p = read(), k = readl();
for (int i = 1; i \le n; i++) {
G[i].clear();
a[i] = readl();
}
for (int i = 1; i < n; i++) {
int u = read(), v = read();
G[u].emplace_back(v);
G[v].emplace_back(u);
dfs1(p, -1);
ans = 0;
dfs2(p, -1);
printf("%lld\n", ans * 2);
return 0;
```

xiv)

```
#include <bits/stdc++.h>
using namespace std;
using ll = long long;
11 \text{ n, k, sum} = 0;
ll\ cntk = 0,\ mxwt = -1;
vector<ll> wt_node;
vector<ll> wt_stree;
vector<vector<ll>> g;
void dfs(ll nn, ll par, ll x)
vector<ll> wt_child;
wt_stree[nn] = wt_node[nn];
for(auto it : g[nn])
if(it != par)
dfs(it, nn, x), wt_child.push_back(wt_stree[it]);
ll sz = wt_child.size();
sort(wt_child.begin(), wt_child.end());
for(ll i=0; i<sz; i++)
if(wt_stree[nn]+wt_child[i] <= x)</pre>
wt_stree[nn] += wt_child[i];
else
cntk += sz - i;
break;
}}}
bool check(ll x)
cntk = 0;
dfs(1, 0, x);
if(cntk \le k-1)
return 1;
return 0;
void solve()
```



```
cin >> n >> k;
wt_node.assign(n+1, 0);
g.assign(n+1, {});
wt_stree.assign(n+1, 0);
ll i, x, y;
for(i=1; i<=n; i++)
cin >> wt_node[i];
sum += wt_node[i];
mxwt = max(mxwt, wt_node[i]);
for(i=1; i<n; i++)
cin >> x >> y;
g[x].push_back(y);
g[y].push_back(x);
ll\ lo = mxwt,\ hi = sum;
ll mid, ans;
while(lo <= hi)
mid = (lo+hi)/2;
if(check(mid))
ans = mid;
hi = mid-1;
else
lo = mid+1;
cout << ans << "\n";
signed main()
ios_base::sync_with_stdio(0);
cin.tie(NULL);
cout.tie(NULL);
11 t = 1;
// cin >> t;
while(t--) solve();
return 0;}
```



```
xv)
```

```
#include < bits/stdc++.h>
using namespace std;
typedef long long LL;
void FastIO(void)
ios_base::sync_with_stdio(false);
cin.tie(NULL);
}
signed main()
FastIO();
LL T;
cin >> T;
while(T--)
LL N;
cin >> N;
vector<vector<LL>> coOrd;
for(LL i = 1; i \le N; ++i)
LL x, y;
cin >> x >> y;
coOrd.push_back({x, y, i});
sort(coOrd.begin(), coOrd.end());
for(LL i = 0; i < N-1; ++i)
cout << coOrd[i][2] << "\ " << coOrd[i+1][2] << endl;
```

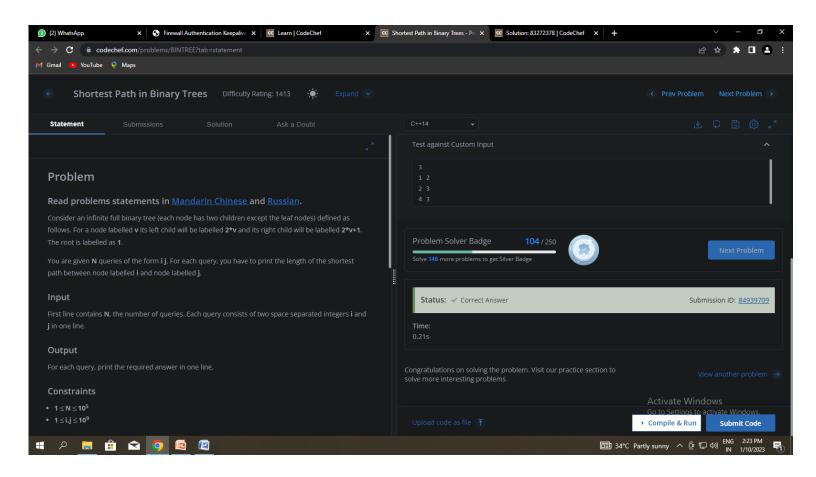




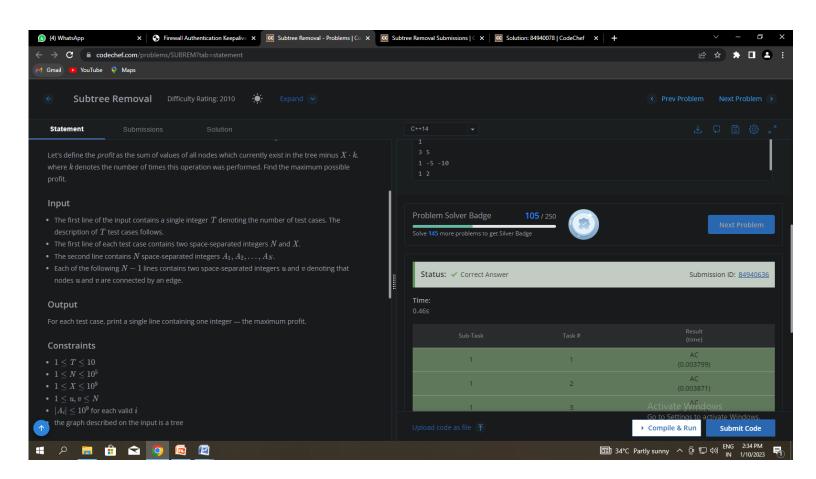


3. Observations/Discussions/ Complexity Analysis:

i)

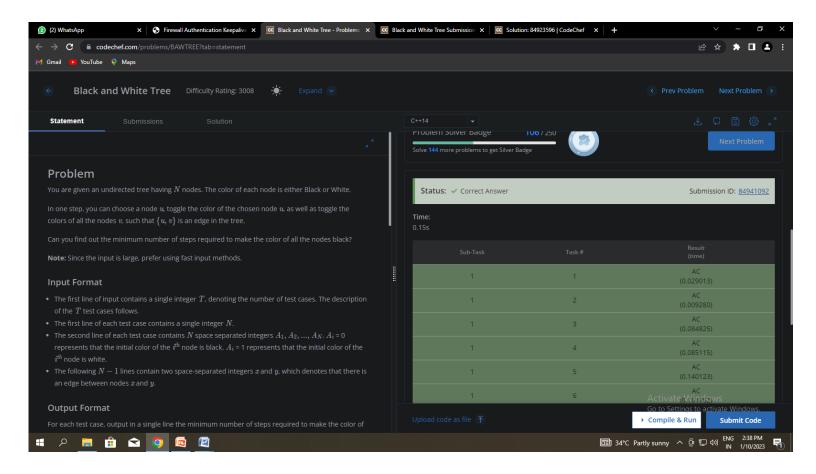






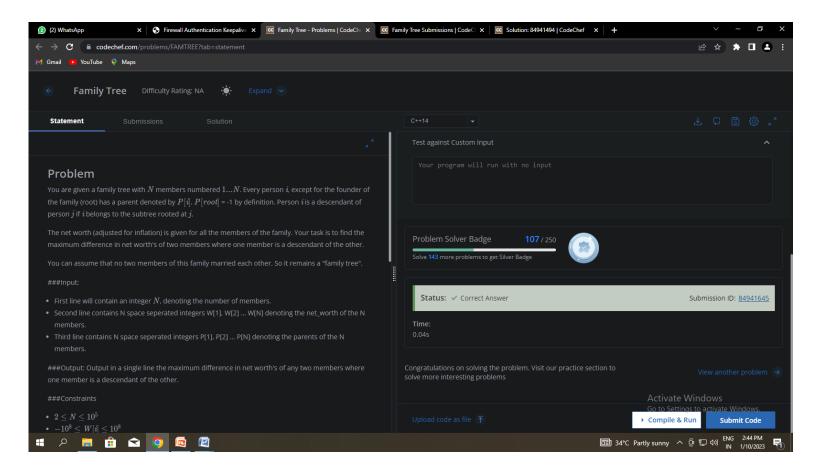


iii)



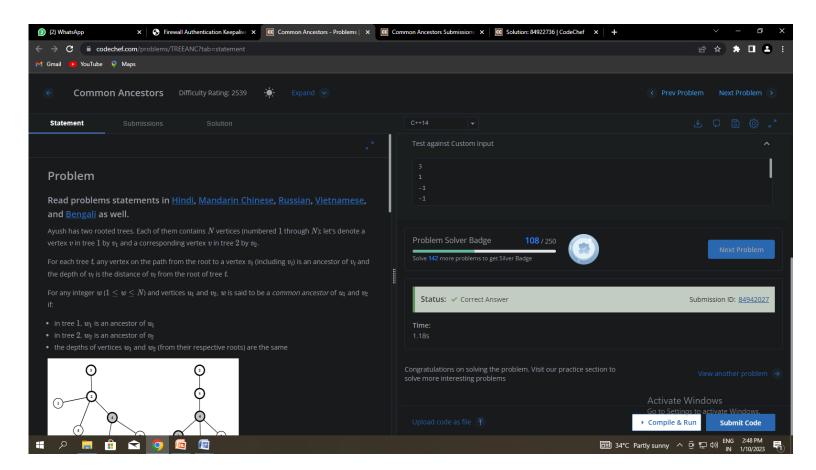


iv)

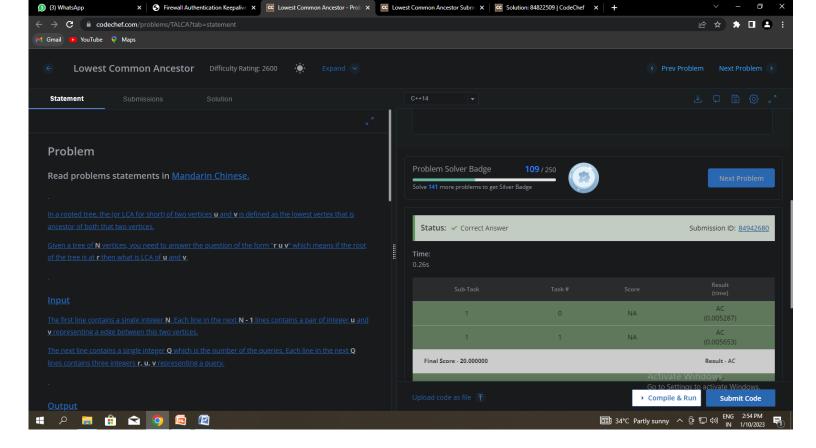




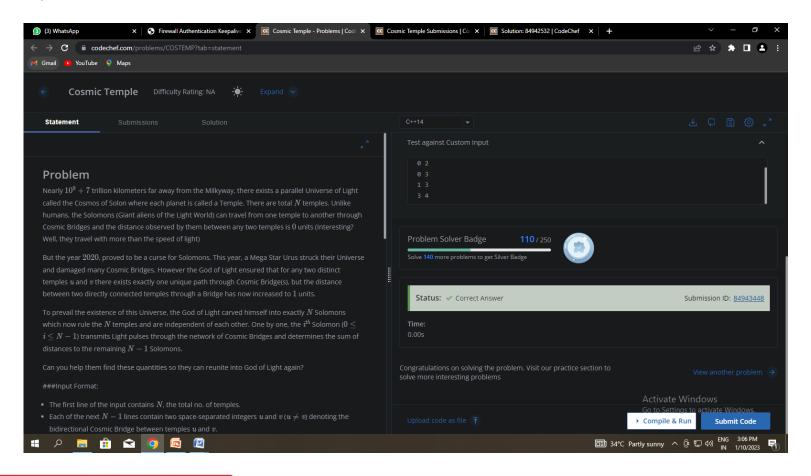
v)



vi)

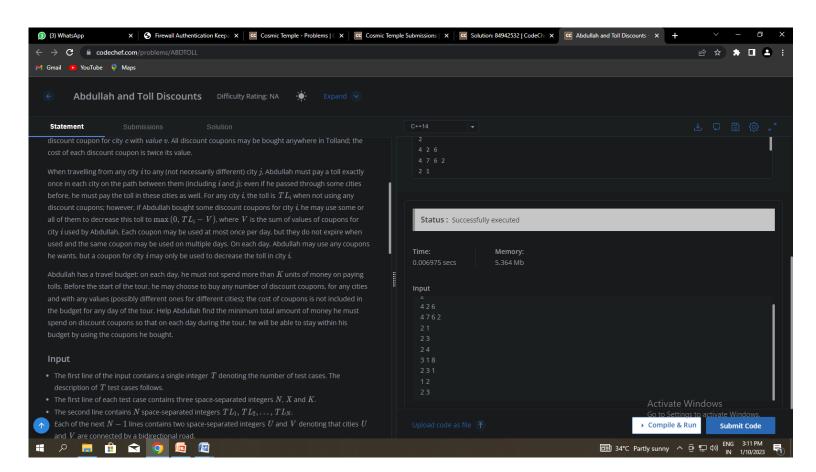


vii)



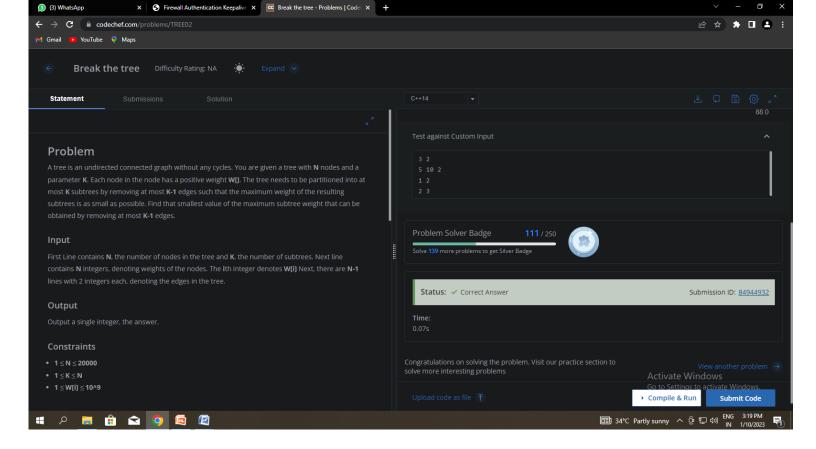


xiii)

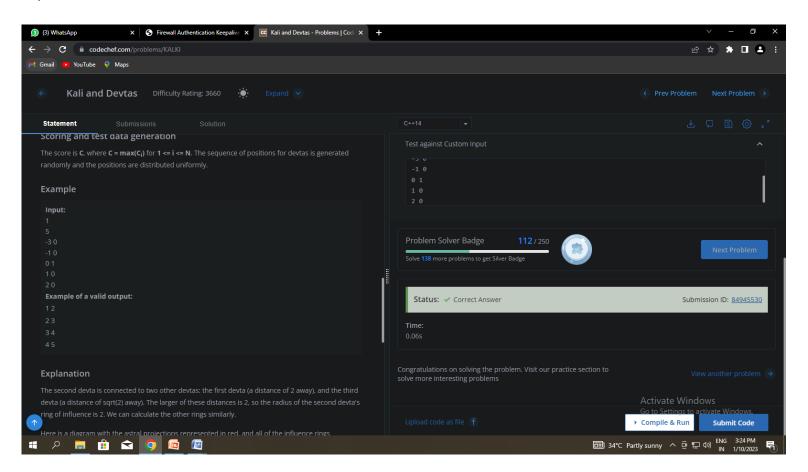


xiv)





xv)









Evaluation Grid (To be created as per the SOP and Assessment guidelines by the faculty):

Sr. No.	Parameters	Marks Obtained	Maximum Marks
1.			
2.			
3.			

