

# Integer sequences, algebraic series and differential operators<sup>1</sup>

PhD Defense

Sergey Yurkevich

University Paris-Saclay (Inria Saclay) and  
University of Vienna



universität  
wien

université  
PARIS-SACLAY

6th of July, 2023

---

<sup>1</sup>Supervised by Alin Bostan and Herwig Hauser

# Contents of the thesis I

**Chapter 1:** Introduction and summary of all chapters.

**Chapter 2:** “On a Class of Hypergeometric Diagonals”, with A. Bostan, 2022.  
In: *Proceedings of the American Mathematical Society*, vol 150, pp. 1071–1897.

**Chapter 3:** Joint work with A. Bostan and J.-A. Weil, and: “The art of algorithmic guessing in gfun”, 2022. In: *Maple Transactions*, vol 2, pp. 14421:1–14421:19.

**Chapter 4:** “A hypergeometric proof that Iso is bijective”, with A. Bostan, 2022.  
In: *Proceedings of the American Mathematical Society*, vol 150, pp. 2131–2136.

**Chapter 5:** “Fast Computation of the N-th Term of a q-Holonomic Sequence and Applications”, with A. Bostan, 2023. In *J. of Symbolic Comp.*, vol 115, pp. 96–123.

## Contents of the thesis II

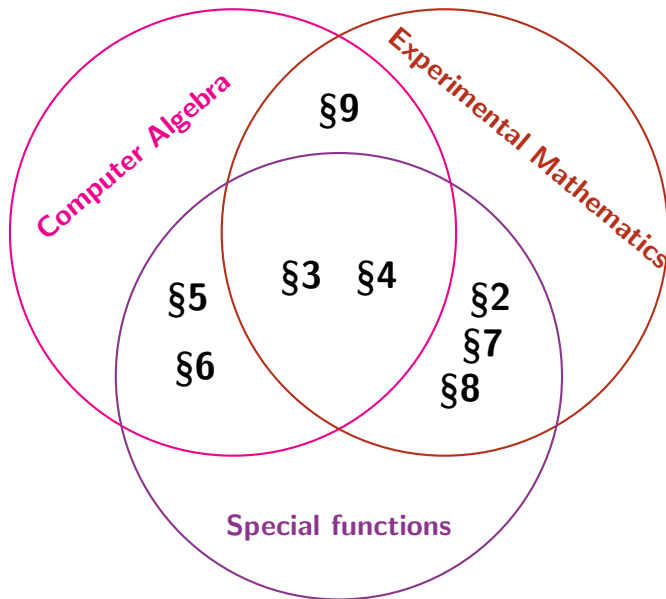
**Chapter 6:** “Beating binary powering for polynomial matrices”, with A. Bostan and V. Neiger, 2023. To appear in the Proceedings of *ISSAC'23*.

**Chapter 7:** “On the  $q$ -analogue of Pólya's Theorem”, with A. Bostan, 2023.  
In: *Electronic Journal of Combinatorics*, vol 30, pp. 2.9:1-9.

**Chapter 8:** “On the representability of sequences as constant terms”, with A. Bostan and A. Straub, 2023. To appear in *Journal of Number Theory*.

**Chapter 9:** “An algorithmic approach to Rupert's problem”, with J. Steininger, 2023,  
In: *Mathematics of Computation*, vol 92, pp. 1905–1929.

**Chapter 10:** A collection of 60 open problems and questions related to the thesis.



## Chapter 2: *Hypergeometric diagonals*

$$\text{Diag}((1+x_1)^{b_1} \cdots (1+x_1+\cdots+x_N)^{b_N}) = {}_M F_{M-1}([u]; [v]; (-N)^N t).$$

# Starting point

*"Guessing – that's the important beginning of solving any problem."*

- Starting point is the main identity from [\[Abdelaziz, Koutschan, Maillard, 2020\]](#):

$${}_3F_2\left(\left[\frac{2}{9}, \frac{5}{9}, \frac{8}{9}\right] ; \left[1, \frac{2}{3}\right] ; 27t\right) = \text{Diag}\left(\frac{(1-x-y)^{1/3}}{1-x-y-z}\right)$$

# Starting point

*"Guessing – that's the important beginning of solving any problem."*

- Starting point is the main identity from [Abdelaziz, Koutschan, Maillard, 2020]:

$${}_3F_2 \left( \left[ \frac{2}{9}, \frac{5}{9}, \frac{8}{9} \right] ; \left[ 1, \frac{2}{3} \right] ; 27t \right) = \text{Diag} \left( \frac{(1-x-y)^{1/3}}{1-x-y-z} \right)$$

- Left-hand side is a generalized *hypergeometric function*:

$${}_3F_2 \left( \left[ \frac{2}{9}, \frac{5}{9}, \frac{8}{9} \right] ; \left[ 1, \frac{2}{3} \right] ; 27t \right) := 1 + \frac{40}{9}t + \frac{5236}{81}t^2 + \cdots + a_n t^n + \cdots$$

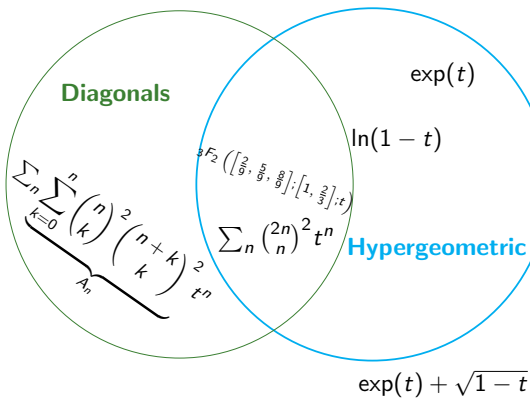
$$\frac{a_{n+1}}{a_n} = \frac{(9n+2)(9n+5)(9n+8)}{3(n+1)^2(9n+6)}$$

- Right-hand side is the diagonal of an *algebraic function*:

$$\frac{(1-x-y)^{1/3}}{1-x-y-z} = 1 + \frac{2}{3}x + \frac{2}{3}y + z + \frac{10}{9}xy + \frac{5}{3}xz + \cdots + \frac{40}{9}xyz + \cdots + \frac{5236}{81}x^2y^2z^2 + \cdots$$

# Setting

## P-finite/D-finite



A sequence  $(u_n)_{n \geq 0}$  is **P-finite** if it satisfies a linear recurrence with polynomial coefficients:

$$c_r(n)u_{n+r} + \cdots + c_0(n)u_n = 0.$$

$(u_n)_{n \geq 0}$  is **hypergeometric** if  $r = 1$ .

Let  $(\alpha)_n = \alpha \cdot (\alpha + 1) \cdots (\alpha + n - 1)$ .

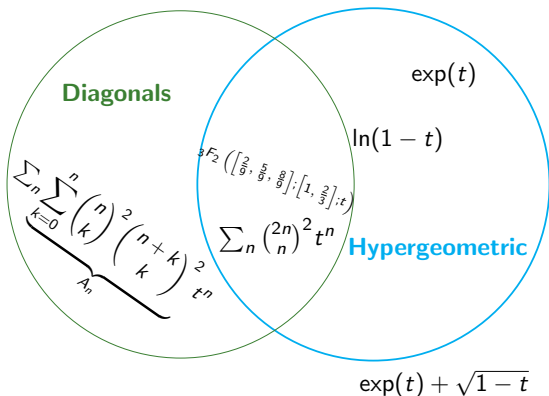
Then  $u_n = \frac{(a)_n \cdot (b)_n}{(c)_n \cdot n!}$  satisfies

$$(c + n)(n + 1)u_{n+1} - (a + n)(b + n)u_n = 0.$$



# Setting

## P-finite/D-finite



A series  $f(t) \in \mathbb{Q}[[t]]$  is **D-finite** if it satisfies a linear differential equation with polynomial coefficients:

$$p_r(t)f^{(r)}(t) + \cdots + p_0(t)f(t) = 0.$$

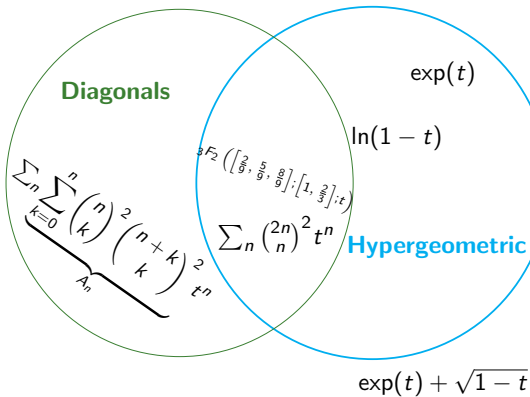
Let  $(\alpha)_n = \alpha \cdot (\alpha + 1) \cdots (\alpha + n - 1)$ .

Then  ${}_2F_1 \left[ \begin{smallmatrix} a & b \\ c \end{smallmatrix}; t \right] := \sum_{n \geq 0} \frac{(a)_n (b)_n}{(c)_n n!} t^n$  satisfies

$$t(1-t)f''(t) + (c - (a+b+1)t)f'(t) - abf(t) = 0.$$

# Setting

## P-finite/D-finite



For a multivariate power series

$$f(x_1, \dots, x_n) = \sum_{j_1, \dots, j_n} f_{j_1, \dots, j_n} x_1^{j_1} \cdots x_n^{j_n}$$

the **diagonal** is given by

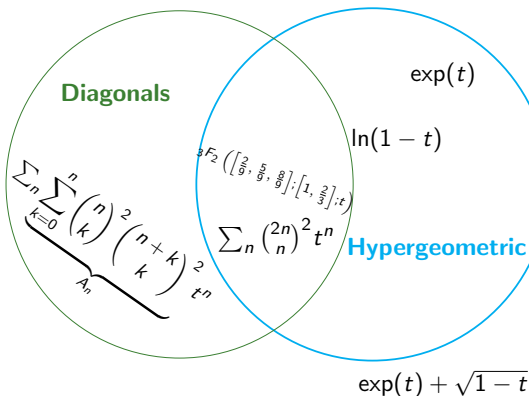
$$\text{Diag}(f) = \sum_j f_{j,j,\dots,j} t^j \in \mathbb{Q}[[t]].$$

**Diagonals** are series which can be written as diagonals of **multivariate algebraic** functions.

$$\text{Diag} \left( \frac{1}{1-x-y} \right) = \text{Diag} \sum_{i,j} \binom{i+j}{j} x^i y^j = \sum_n \binom{2n}{n} t^n = (1-4t)^{-1/2}$$

# Setting

## P-finite/D-finite



For a multivariate power series

$$f(x_1, \dots, x_n) = \sum_{j_1, \dots, j_n} f_{j_1, \dots, j_n} x_1^{j_1} \cdots x_n^{j_n}$$

the **diagonal** is given by

$$\text{Diag}(f) = \sum_j f_{j,j,\dots,j} t^j \in \mathbb{Q}[[t]].$$

**Diagonals** are series which can be written as diagonals of **multivariate algebraic** functions.

**Christol's Conjecture** [Christol, 1986]: Any convergent **D-finite** power series with integer coefficients is a **diagonal**. Specifically:  ${}_3F_2\left(\left[\frac{1}{9}, \frac{4}{9}, \frac{5}{9}\right]; \left[1, \frac{1}{3}\right], t\right) \in \text{Diagonals}$ .

# Main result **A**: Hypergeometric diagonals

*"First guess, then prove."*

*All great discoveries were made in this style."*

Theorem (Bostan, Y., 2022)

The **diagonal** of any finite product of algebraic functions of the form

$$(1 - x_1 - \cdots - x_n)^R, \quad R \in \mathbb{Q},$$

is a generalized **hypergeometric** function with explicitly determined parameters.

# Main result **A**: Hypergeometric diagonals

*"First guess, then prove."*

*All great discoveries were made in this style."*

Theorem (Bostan, Y., 2022)

The **diagonal** of any finite product of algebraic functions of the form

$$(1 - x_1 - \cdots - x_n)^R, \quad R \in \mathbb{Q},$$

is a generalized **hypergeometric** function with explicitly determined parameters.

- This vastly generalizes the main identity in [Abdelaziz, Koutschan, Maillard, 2020].
- We also settle down other memberships: E.g.  ${}_3F_2\left(\left[\frac{1}{4}, \frac{3}{8}, \frac{7}{8}\right]; \left[1, \frac{1}{3}\right], t\right) \in \mathbf{Diagonals}$ .

# Main result **A**: Hypergeometric diagonals

*"First guess, then prove."*

*All great discoveries were made in this style."*

Theorem (Bostan, Y., 2022)

The **diagonal** of any finite product of algebraic functions of the form

$$(1 - x_1 - \cdots - x_n)^R, \quad R \in \mathbb{Q},$$

is a generalized **hypergeometric** function with explicitly determined parameters.

- This vastly generalizes the main identity in [Abdelaziz, Koutschan, Maillard, 2020].
- We also settle down other memberships: E.g.  ${}_3F_2\left(\left[\frac{1}{4}, \frac{3}{8}, \frac{7}{8}\right]; \left[1, \frac{1}{3}\right], t\right) \in \mathbf{Diagonals}$ .
- **Main observation** for the proof:

$$\begin{aligned} & [x_1^{k_1} \cdots x_N^{k_N}] (1 + x_1)^{b_1} (1 + x_1 + x_2)^{b_2} \cdots (1 + x_1 + \cdots + x_N)^{b_N} \\ &= \binom{b_N}{k_N} \binom{b_{N-1} + b_N - k_N}{k_{N-1}} \cdots \binom{b_1 + \cdots + b_N - k_N \cdots - k_2}{k_1}. \end{aligned}$$

## Chapter 3: *Dubrovin-Yang-Zagier numbers and algebraicity of D-finite functions*

$$(a_n)_{n \geq 0} = (1, -48300, 7981725900, -1469166887370000, \dots)$$

$$(b_n)_{n \geq 0} = (1, -144900, 88464128725, -62270073456990000, \dots)$$

## Origin of $a_n$ and $b_n$

*"So this is a very mysterious example"*

- In [Arithmetic and Topology of Differential Equations, 2018](#) by Don Zagier:

$$\begin{aligned}
 &u_{n-3} + 20(4500n^2 - 18900n + 19739)u_{n-2} + 80352000n(5n-1)(5n-2)(5n-4)u_n + \\
 &\quad + 25(2592000n^4 - 16588800n^3 + 39118320n^2 - 39189168n + 14092603)u_{n-1} = 0, \\
 &\text{with initial terms } u_0 = 1, u_1 = -161/(2^{10} \cdot 3^5) \text{ and } u_2 = 26605753/(2^{23} \cdot 3^{12} \cdot 5^2).
 \end{aligned}$$



# Origin of $a_n$ and $b_n$

*"So this is a very mysterious example"*

- In [Arithmetic and Topology of Differential Equations, 2018](#) by Don Zagier:

$$u_{n-3} + 20(4500n^2 - 18900n + 19739)u_{n-2} + 80352000n(5n-1)(5n-2)(5n-4)u_n +$$

$$+ 25(2592000n^4 - 16588800n^3 + 39118320n^2 - 39189168n + 14092603)u_{n-1} = 0,$$

with initial terms  $u_0 = 1$ ,  $u_1 = -161/(2^{10} \cdot 3^5)$  and  $u_2 = 26605753/(2^{23} \cdot 3^{12} \cdot 5^2)$ .

## Problem (Zagier, 2018)

Find  $(\alpha, \beta) \in \mathbb{Q}^* \times \mathbb{Q}^*$  such that  $u_n \cdot (\alpha)_n \cdot (\beta)_n \cdot \gamma^n \in \mathbb{Z}$  for some  $\gamma \in \mathbb{Z}^*$ .  
 $(x)_n := x \cdot (x+1) \cdots (x+n-1)$ .

- [Yang and Zagier]:  $a_n = u_n \cdot (3/5)_n \cdot (4/5)_n \cdot (2^{10} \cdot 3^5 \cdot 5^4)^n \in \mathbb{Z}$ ,
- [Dubrovin and Yang]:  $b_n = u_n \cdot (2/5)_n \cdot (9/10)_n \cdot (2^{12} \cdot 3^5 \cdot 5^4)^n \in \mathbb{Z}$ .

# Origin of $a_n$ and $b_n$

*"So this is a very mysterious example"*

- In [Arithmetic and Topology of Differential Equations, 2018](#) by [Don Zagier](#):

$$u_{n-3} + 20(4500n^2 - 18900n + 19739)u_{n-2} + 80352000n(5n-1)(5n-2)(5n-4)u_n +$$

$$+ 25(2592000n^4 - 16588800n^3 + 39118320n^2 - 39189168n + 14092603)u_{n-1} = 0,$$

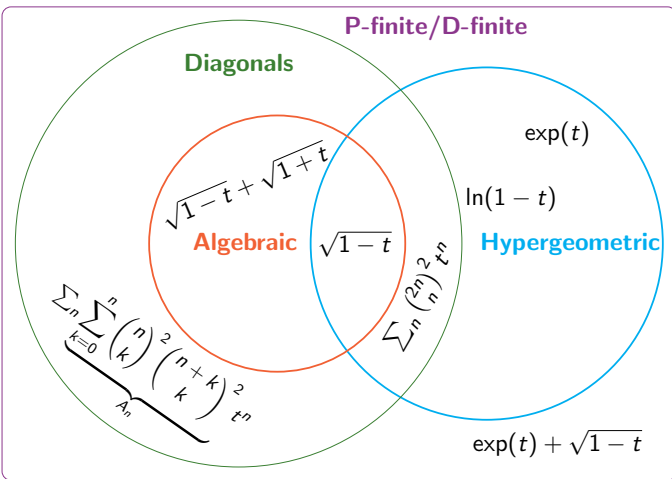
with initial terms  $u_0 = 1$ ,  $u_1 = -161/(2^{10} \cdot 3^5)$  and  $u_2 = 26605753/(2^{23} \cdot 3^{12} \cdot 5^2)$ .

## Problem (Zagier, 2018)

Find  $(\alpha, \beta) \in \mathbb{Q}^* \times \mathbb{Q}^*$  such that  $u_n \cdot (\alpha)_n \cdot (\beta)_n \cdot \gamma^n \in \mathbb{Z}$  for some  $\gamma \in \mathbb{Z}^*$ .  
 $(x)_n := x \cdot (x+1) \cdots (x+n-1)$ .

- [\[Yang and Zagier\]](#):  $a_n = u_n \cdot (3/5)_n \cdot (4/5)_n \cdot (2^{10} \cdot 3^5 \cdot 5^4)^n \in \mathbb{Z}$ ,
- [\[Dubrovin and Yang\]](#):  $b_n = u_n \cdot (2/5)_n \cdot (9/10)_n \cdot (2^{12} \cdot 3^5 \cdot 5^4)^n \in \mathbb{Z}$ .
- "Yang and I found a formula showing that the numbers  $a_n$  are integers [...]"  
 "Dubrovin and Yang found that the numbers  $b_n$  are also integral and that in this case the generating function [...] is actually **algebraic**!" [\[Zagier, 2018\]](#)

# Definitions and interactions



[Abel, 1827]: **Algebraic**  $\subseteq$  **D-finite**.

[Furstenberg, 1967]:  
**Algebraic**  $\subseteq$  **Diagonals**.

[Singer 1979, 2014]:  
**D-finite**  $f(t) \stackrel{?}{\in}$  **Algebraic**.

[Christol, 1984 and Lipshitz, 1988]:  
**Diagonals**  $\subseteq$  **D-finite**.

[Petkovsek 1992]:  
**D-finite**  $f(t) \stackrel{?}{\in}$  **Hypergeometric**.

[Beukers, Heckman, 1989]:  
**Algebraic**  $\cap$  **Hypergeometric**.

[Bostan, Lairez, Salvy, 2017]:  
**Diagonals** = **Multiple binomial sums**.

**André-Christol Conjecture** [André, 2004]:

**D-finite**  $f(t) \in \mathbb{Z}[[t]]$  convergent & minimal ODE ordinary in 0  $\Rightarrow f(t)$  **Algebraic**

## Main result B: Solving the mystery of $a_n$ and $b_n$

*"So this is a very mysterious example."*

- "Yang and I found a formula showing that the numbers  $a_n$  are integers [...]"  
"Dubrovin and Yang found that the numbers  $b_n$  are *also* integral and that in this case the generating function [...] is actually **algebraic!**"
- "My presumed arithmetic intuition [...] was entirely broken" – [Wadim Zudilin]

### Problem

Investigate the nature of  $(a_n)_{n \geq 0}$ ,  $(b_n)_{n \geq 0}$  and similar sequences.

## Main result B: Solving the mystery of $a_n$ and $b_n$

*"So this is a very mysterious example."*

- "Yang and I found a formula showing that the numbers  $a_n$  are integers [...]"  
"Dubrovin and Yang found that the numbers  $b_n$  are *also* integral and that in this case the generating function [...] is actually **algebraic**!"
- "My presumed arithmetic intuition [...] was entirely broken" – [Wadim Zudilin]

### Problem

Investigate the nature of  $(a_n)_{n \geq 0}$ ,  $(b_n)_{n \geq 0}$  and similar sequences.

### Theorem (Bostan, Weil, Y.)

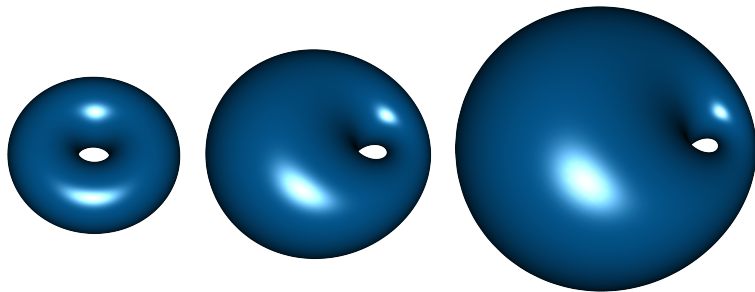
*The generating functions of both  $(a_n)_{n \geq 0}$  and  $(b_n)_{n \geq 0}$  are **algebraic**.*

### Theorem (Bostan, Weil, Y.)

*Seven more solutions to Zagier's problem:  $(c_n)_{n \geq 0}, \dots, (i_n)_{n \geq 0} \in \mathbb{Z}$ .*

## Chapter 4:

*On the reduced volume of conformal transformations of tori*

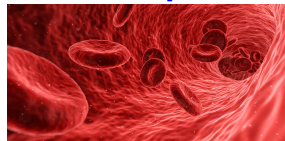


# Motivation and Introduction

*"Why do all humans have the same biconcave shaped red blood cells?"*

- *Canham model* predicts shape of biomembranes like blood cells [Canham, 1970].
- The model asks to minimize the *Willmore energy*

$$W(S) := \int_S H^2 dA, \quad (H \text{ is mean the curvature})$$



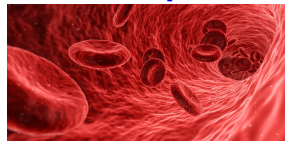
over orientable closed surfaces  $S \subseteq \mathbb{R}^3$  with genus  $g$ , area  $A_0$  and volume  $V_0$ .

# Motivation and Introduction

*"Why do all humans have the same biconcave shaped red blood cells?"*

- *Canham model* predicts shape of biomembranes like blood cells [Canham, 1970].
- The model asks to minimize the *Willmore energy*

$$W(S) := \int_S H^2 dA, \quad (H \text{ is mean the curvature})$$



over orientable closed surfaces  $S \subseteq \mathbb{R}^3$  with genus  $g$ , area  $A_0$  and volume  $V_0$ .

- [Willmore, 1965]: For a torus  $T = T(R, r)$  the Willmore energy is:

$$W(T) = \frac{\pi^2 R^2}{r\sqrt{R^2 - r^2}} \rightsquigarrow \text{minimal for } R/r = \sqrt{2}.$$

Theorem (Willmore 1964 (conjectured); Marques, Neves, 2014)

*Across all closed surfaces in  $\mathbb{R}^3$  of genus  $g \geq 1$  the Willmore energy is minimal for  $T_{\sqrt{2}}$ .*

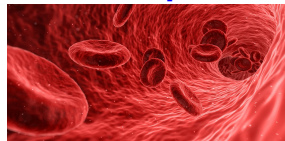


# Motivation and Introduction

*"Why do all humans have the same biconcave shaped red blood cells?"*

- *Canham model* predicts shape of biomembranes like blood cells [Canham, 1970].
- The model asks to minimize the *Willmore energy*

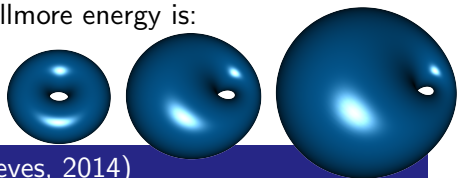
$$W(S) := \int_S H^2 dA, \quad (H \text{ is mean the curvature})$$



over orientable closed surfaces  $S \subseteq \mathbb{R}^3$  with genus  $g$ , area  $A_0$  and volume  $V_0$ .

- [Willmore, 1965]: For a torus  $T = T(R, r)$  the Willmore energy is:

$$W(T) = \frac{\pi^2 R^2}{r\sqrt{R^2 - r^2}} \rightsquigarrow \text{minimal for } R/r = \sqrt{2}.$$



Theorem (Willmore 1964 (conjectured); Marques, Neves, 2014)

*Across all closed surfaces in  $\mathbb{R}^3$  of genus  $g \geq 1$  the Willmore energy is minimal for  $T_{\sqrt{2}}$ .*

- $W(S)$  is invariant under Möbius transformations  $\Rightarrow$  no uniqueness of the shape.

# Main result C: Iso is bijective

*"Nature is not generic."*

- In Canham's model, instead of  $A_0$  and  $V_0$  rather prescribe the *isoperimetric ratio*:

$$\iota_0 := \pi^{1/6} \frac{\sqrt[3]{6V_0}}{\sqrt{A_0}} \in (0, 1].$$

## Question

Is the minimizer of  $W(S)$  with prescribed genus  $g$  and isoperimetric ratio  $\iota_0$  unique?

# Main result C: Iso is bijective

*"Nature is not generic."*

- In Canham's model, instead of  $A_0$  and  $V_0$  rather prescribe the *isoperimetric ratio*:

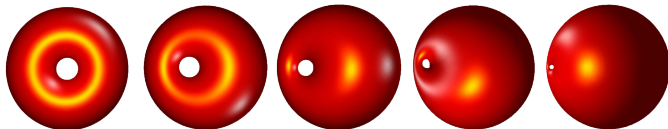
$$\iota_0 := \pi^{1/6} \frac{\sqrt[3]{6V_0}}{\sqrt{A_0}} \in (0, 1].$$

## Question

Is the minimizer of  $W(S)$  with prescribed genus  $g$  and isoperimetric ratio  $\iota_0$  unique?

Theorem (Yu, Chen, 21; Melczer, Mezzarobba, 21; Bostan, Y., 22)

*The shape of the projection of the Clifford torus to  $\mathbb{R}^3$  is uniquely determined by  $\iota_0$ . Thus, if  $g = 1$  and  $\iota_0^3 \in [3/(2^{5/4}\sqrt{\pi}), 1]$  then Canham's model has a unique solution.*



# Main result C': Iso is bijective

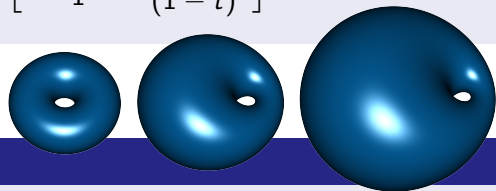
*"I could never resist a definite integral."*

Proposition (Bostan, Y., 2022)

The surface area  $\sqrt{2}\pi^2 A(t^2)$  and volume  $\sqrt{2}\pi^2 V(t^2)$  of  $i_{(t,0,0)}(T_{\sqrt{2}})$  are given by

$$A(t) = \frac{4(1-t^2)}{(t^2-6t+1)^2} \cdot {}_2F_1\left[\begin{matrix} -\frac{1}{2} & -\frac{1}{2} \\ 1 \end{matrix}; \frac{4t}{(1-t)^2}\right],$$

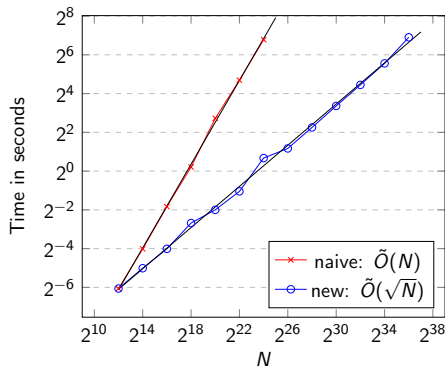
$$V(t) = \frac{2(1-t)^3}{(t^2-6t+1)^3} \cdot {}_2F_1\left[\begin{matrix} -\frac{3}{2} & -\frac{3}{2} \\ 1 \end{matrix}; \frac{4t}{(1-t)^2}\right].$$



Theorem (Bostan, Y., 2022)

The function  $\text{Iso}(t)^2 = 36\pi \frac{V(t^2)^2}{A(t^2)^3}$  is increasing on  $t \in (0, \sqrt{2}-1)$ .

## Chapter 5: Computing terms in $q$ -holonomic sequences



## Main result **D**: Sublinear algorithm for *q*-holonomic sequences

- A sequence  $(u_n)_{n \geq 0} \in \mathbb{K}$  is **holonomic/P-finite** if it satisfies

$$c_r(n)u_{n+r} + \cdots + c_0(n)u_n = 0 \quad n \geq 0, \quad c_0(x), \dots, c_r(x) \in \mathbb{K}[x].$$

Theorem (Strassen, 1977; Chudnovsky<sup>2</sup>, 1988)

Given  $N \in \mathbb{N}$ , one can compute  $u_N$  in  $\tilde{O}(\sqrt{N})$  arithmetic operations.

*Naive:*  $O(N)$

# Main result **D**: Sublinear algorithm for $q$ -holonomic sequences

- A sequence  $(u_n)_{n \geq 0} \in \mathbb{K}$  is **holonomic/P-finite** if it satisfies

$$c_r(n)u_{n+r} + \cdots + c_0(n)u_n = 0 \quad n \geq 0, \quad c_0(x), \dots, c_r(x) \in \mathbb{K}[x].$$

Theorem (Strassen, 1977; Chudnovsky<sup>2</sup>, 1988)

Given  $N \in \mathbb{N}$ , one can compute  $u_N$  in  $\tilde{O}(\sqrt{N})$  arithmetic operations. *Naive:  $O(N)$*

- A sequence  $(u_n(q))_{n \geq 0} \in \mathbb{K}$  is called  **$q$ -holonomic** if for some  $q \in \mathbb{K}$  it satisfies

$$c_r(q, q^n)u_{n+r} + \cdots + c_0(q, q^n)u_n = 0 \quad n \geq 0, \quad c_0(x, y), \dots, c_r(x, y) \in \mathbb{K}[x, y].$$

Theorem (Bostan, Y., 2023)

Given  $N \in \mathbb{N}$ , one can compute  $u_N(q)$  in  $\tilde{O}(\sqrt{N})$  arithmetic operations. *Naive:  $O(N)$*

**Idea:** For  $M(x) \in \mathbb{K}[x]^{r \times r}$  compute  $M(q^{N-1}) \cdots M(q)M(1)$  using baby-steps/giant-steps.

## Application: Evaluation of polynomials

*"Do not waste a factor of two!"*

- **Task:** Given a polynomial  $P(x) \in \mathbb{K}[x]$  and  $q \in \mathbb{K}$ , deduce  $P(q) \in \mathbb{K}$  fast.



## Application: Evaluation of polynomials

*"Do not waste a factor of two!"*

- **Task:** Given a polynomial  $P(x) \in \mathbb{K}[x]$  and  $q \in \mathbb{K}$ , deduce  $P(q) \in \mathbb{K}$  fast.
- Generically, Horner's rule needs  $O(\deg P)$  operations.
- Our results imply that one can do better for large families of polynomials.

# Application: Evaluation of polynomials

*"Do not waste a factor of two!"*

- **Task:** Given a polynomial  $P(x) \in \mathbb{K}[x]$  and  $q \in \mathbb{K}$ , deduce  $P(q) \in \mathbb{K}$  fast.
- Generically, Horner's rule needs  $O(\deg P)$  operations.
- Our results imply that one can do better for large families of polynomials.
- For example, the truncated Jacobi theta function

$$\vartheta_N(x) := 1 + x + x^4 + x^9 + \cdots + x^{N^2}$$

evaluated at  $q \in \mathbb{K}$  in  $\tilde{O}(\sqrt{N})$  operations [Nogneng, Schost, 2018], [Bostan, Y., 2023].

# Application: Evaluation of polynomials

*"Do not waste a factor of two!"*

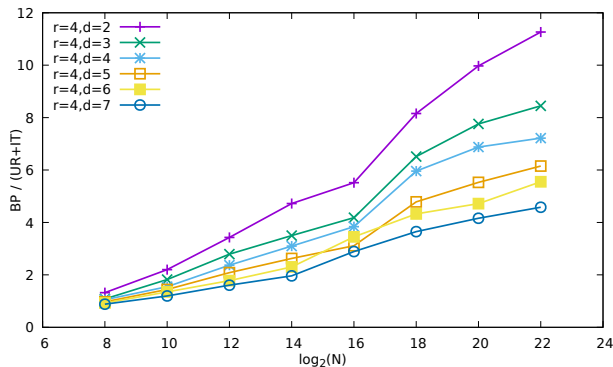
- **Task:** Given a polynomial  $P(x) \in \mathbb{K}[x]$  and  $q \in \mathbb{K}$ , deduce  $P(q) \in \mathbb{K}$  fast.
- Generically, Horner's rule needs  $O(\deg P)$  operations.
- Our results imply that one can do better for large families of polynomials.
- For example, the truncated Jacobi theta function

$$\vartheta_N(x) := 1 + x + x^4 + x^9 + \cdots + x^{N^2}$$

evaluated at  $q \in \mathbb{K}$  in  $\tilde{O}(\sqrt{N})$  operations [Nogneng, Schost, 2018], [Bostan, Y., 2023].

- **Method:**  $\vartheta_N(q) = u_N$ , where  $u_n = \sum_{k=0}^n q^{k^2}$  is  $q$ -holonomic.
- [Bostan, Y., 2023]: Same complexity via unified algorithm for  $\prod_{i=0}^N (x - a^i)$ , or  $q$ -Hermite polynomials, or  $\prod_{i=1}^{\infty} (1 - x^i)^3 \bmod x^N$ , etc.

## Chapter 6: *Computing terms in polynomial C-finite sequences*



## Polynomial C-finite sequences: Example

- Fibonacci polynomials:  $F_0(x) = 0$ ,  $F_1(x) = 1$  and  $F_{n+2}(x) = xF_{n+1}(x) + F_n(x)$   
 $F_9(x) = 1 + 10x^2 + 15x^4 + 7x^6 + x^8$  and  $F_{10}(x) = 5x + 20x^3 + 21x^5 + 8x^7 + x^9$ .

## Polynomial C-finite sequences: Example

- Fibonacci polynomials:  $F_0(x) = 0$ ,  $F_1(x) = 1$  and  $F_{n+2}(x) = xF_{n+1}(x) + F_n(x)$   
 $F_9(x) = 1 + 10x^2 + 15x^4 + 7x^6 + x^8$  and  $F_{10}(x) = 5x + 20x^3 + 21x^5 + 8x^7 + x^9$ .
- Compute using the definition:  $F_{n+2}(x) = xF_{n+1}(x) + F_n(x)$ .  $O(N^2)$

## Polynomial C-finite sequences: Example

- Fibonacci polynomials:  $F_0(x) = 0$ ,  $F_1(x) = 1$  and  $F_{n+2}(x) = xF_{n+1}(x) + F_n(x)$   
 $F_9(x) = 1 + 10x^2 + 15x^4 + 7x^6 + x^8$  and  $F_{10}(x) = 5x + 20x^3 + 21x^5 + 8x^7 + x^9$ .
- Compute using the definition:  $F_{n+2}(x) = xF_{n+1}(x) + F_n(x)$ .  $O(N^2)$
- [Folklore]: Use binary powering to compute  $M_N$ , where  $M_n(x) = \begin{pmatrix} x & 1 \\ 1 & 0 \end{pmatrix}^n$ :

$$M_n(x) = \begin{cases} M_{n/2}(x)^2 & \text{if } n \text{ even,} \\ M(x) \cdot M_{\frac{n-1}{2}}(x)^2 & \text{if } n \text{ odd.} \end{cases} \quad \text{color: red; } O(N \log(N))$$

## Polynomial C-finite sequences: Example

- Fibonacci polynomials:  $F_0(x) = 0$ ,  $F_1(x) = 1$  and  $F_{n+2}(x) = xF_{n+1}(x) + F_n(x)$   
 $F_9(x) = 1 + 10x^2 + 15x^4 + 7x^6 + x^8$  and  $F_{10}(x) = 5x + 20x^3 + 21x^5 + 8x^7 + x^9$ .
- Compute using the definition:  $F_{n+2}(x) = xF_{n+1}(x) + F_n(x)$ .  $O(N^2)$
- [Folklore]: Use binary powering to compute  $M_N$ , where  $M_n(x) = \begin{pmatrix} x & 1 \\ 1 & 0 \end{pmatrix}^n$ :

$$M_n(x) = \begin{cases} M_{n/2}(x)^2 & \text{if } n \text{ even,} \\ M(x) \cdot M_{\frac{n-1}{2}}(x)^2 & \text{if } n \text{ odd.} \end{cases} \quad \text{span style="color: red;"> $O(N \log(N))$$$

- **Idea:** Write  $F_N(x) = f_0 + f_1x + \dots + f_Nx^N$ . Then  $(f_k)_{k \geq 0}$  is **P-finite**:

$$f_{k+2} = \frac{(N+k+1)(N-k-1)}{4(k+1)(k+2)} f_k \quad \text{for } k \geq 0,$$

with  $(f_0, f_1) = (1, 0)$  for odd  $N$  and  $(f_0, f_1) = (0, N/2)$  for even  $N$ .

$O(N)$



## Main result E: Beating binary powering

*"The development of fast algorithms is slow!"*

A **polynomial C-finite sequence**  $(u_n(x))_{n \geq 0} \in \mathbb{K}[x]^{\mathbb{N}}$  satisfies a recurrence

$$u_{n+r}(x) = c_{r-1}(x)u_{n+r-1}(x) + \cdots + c_0(x)u_n(x),$$

for some polynomials  $c_0(x), \dots, c_{r-1}(x) \in \mathbb{K}[x]$ .

Theorem (Bostan, Neiger, Y., 2023)

Given a **polynomial C-finite sequence**  $(u_n(x))_{n \geq 0}$ , one can compute  $u_N(x)$  in  $O(N)$  operations in  $K$ .

Corollary

Given a polynomial matrix  $M(x)$ , one can compute  $M(x)^N$  in  $O(N)$  field operations.

## Chapter 6:

### *On the $q$ -analogue of Pólya's Theorem*

Specifically, if  $n, k, a, b$  satisfy the conditions stated earlier, is the function

$$F(x, q) = \sum_{t=0}^{\infty} \begin{bmatrix} n + at \\ k + bt \end{bmatrix}_q x^t$$

algebraic? That is, does there exist a nonzero polynomial  $P(x, y, z)$  whose coefficients are constants (say, complex numbers) such that  $P(x, q, F(x, q)) = 0$ , for all  $x$  and  $q$ ?

[Aissen, 1979]

# Main result F: A $q$ -analogue of Pólya's theorem

*"In mathematics often  
the simplest is the best."*

- Consequence of Pólya's theorem [Pólya, 1922]:

Theorem (Pólya, 1922)

For admissible  $n, k, a, b$ , the function  $F(x) := \sum_{j \geq 0} \binom{n+aj}{k+bj} x^j$  is **algebraic** over  $\mathbb{Q}(x)$ .

# Main result F: A $q$ -analogue of Pólya's theorem

*"In mathematics often the simplest is the best."*

- Consequence of Pólya's theorem [Pólya, 1922]:

## Theorem (Pólya, 1922)

For admissible  $n, k, a, b$ , the function  $F(x) := \sum_{j \geq 0} \binom{n+aj}{k+bj} x^j$  is **algebraic** over  $\mathbb{Q}(x)$ .

- Aissen asked whether a  $q$ -analogue holds [Aissen, 1979]. We prove:

## Theorem (Bostan, Y., 2022)

For admissible  $n, k, a, b$ , the function

$$F(x, q) := \sum_{j \geq 0} \left[ \begin{matrix} n+aj \\ k+bj \end{matrix} \right]_q x^j \in \mathbb{C}[q][[x]]$$

is **never algebraic** over  $\mathbb{Q}(q, x)$ . If  $q \in \mathbb{C}$ , then  $F(x, q)$  is algebraic iff  $q$  is root of unity.

# Main result F: A $q$ -analogue of Pólya's theorem

*"In mathematics often the simplest is the best."*

- Consequence of Pólya's theorem [Pólya, 1922]:

## Theorem (Pólya, 1922)

For admissible  $n, k, a, b$ , the function  $F(x) := \sum_{j \geq 0} \binom{n+aj}{k+bj} x^j$  is **algebraic** over  $\mathbb{Q}(x)$ .

- Aissen asked whether a  $q$ -analogue holds [Aissen, 1979]. We prove:

## Theorem (Bostan, Y., 2022)

For admissible  $n, k, a, b$ , the function

$$F(x, q) := \sum_{j \geq 0} \left[ \begin{matrix} n+aj \\ k+bj \end{matrix} \right]_q x^j \in \mathbb{C}[q][[x]]$$

is **never algebraic** over  $\mathbb{Q}(q, x)$ . If  $q \in \mathbb{C}$ , then  $F(x, q)$  is algebraic iff  $q$  is root of unity.

- $u_n(q) = \left[ \begin{matrix} n \\ k \end{matrix} \right]_q := \frac{[n]_q!}{[k]_q! [n-k]_q!}$ , where  $[n]_q! := (1+q) \cdots (1+q+\cdots+q^{n-1})$ .
- **Idea:** It holds that  $(u_n(q))_{n \geq 0}$  is  **$q$ -holonomic**.

## Chapter 8:

### *Representation of sequences as constant terms*

$$\sum_{k=0}^n \binom{n}{k}^2 \binom{n+k}{k}^2 = \text{ct} \left[ \left( \frac{(x+y)(z+1)(x+y+z)(y+x+1)}{xyz} \right)^n \right].$$

# Main result **G**: Describing **Constant terms** $\cap$ **C-finite sequences**

- A sequence  $A(n)$  is a **constant term** if it can be represented as

$$A(n) = \text{ct}[P(\mathbf{x})^n Q(\mathbf{x})],$$

where  $P, Q \in \mathbb{Q}[\mathbf{x}^{\pm 1}]$  are Laurent polynomials in  $\mathbf{x} = (x_1, \dots, x_d)$ .

# Main result **G**: Describing **Constant terms** $\cap$ **C-finite sequences**

- A sequence  $A(n)$  is a **constant term** if it can be represented as

$$A(n) = \text{ct}[P(\mathbf{x})^n Q(\mathbf{x})],$$

where  $P, Q \in \mathbb{Q}[\mathbf{x}^{\pm 1}]$  are Laurent polynomials in  $\mathbf{x} = (x_1, \dots, x_d)$ .

Question (Zagier, 2018; Gorodetsky, 2021; Straub, 2022)

Which **P-finite sequences** are constant terms?

Specifically: Are the Fibonacci numbers a constant term sequence?



# Main result **G**: Describing **Constant terms** $\cap$ **C-finite sequences**

- A sequence  $A(n)$  is a **constant term** if it can be represented as

$$A(n) = \text{ct}[P(\mathbf{x})^n Q(\mathbf{x})],$$

where  $P, Q \in \mathbb{Q}[\mathbf{x}^{\pm 1}]$  are Laurent polynomials in  $\mathbf{x} = (x_1, \dots, x_d)$ .

Question (Zagier, 2018; Gorodetsky, 2021; Straub, 2022)

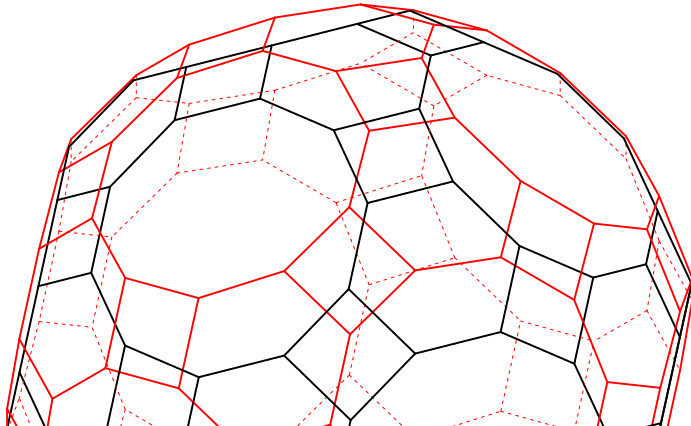
Which **P-finite sequences** are constant terms?

Specifically: Are the Fibonacci numbers a constant term sequence?

Theorem (Bostan, Straub, Y., 2023)

Let  $A(n)$  be a **C-finite sequence**.  $A(n)$  is a **constant term** if and only if it has a single characteristic root  $\lambda$  and  $\lambda \in \mathbb{Q}$ .

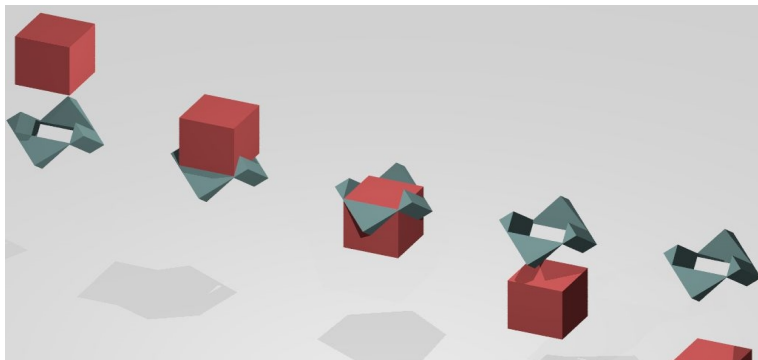
## Chapter 9: *On Rupert's problem*



# Summary and main result **H**: Deciding Rupertness *"It shows us 'what's out there'."*

## Definition

A convex polyhedron  $\mathbf{P} \subseteq \mathbb{R}^3$  is called **Rupert** if a hole with the shape of a straight tunnel can be cut into it such that a copy of  $\mathbf{P}$  can be moved through this hole.



# Summary and main result **H**: Deciding Rupertness *"It shows us 'what's out there'."*

## Definition

A convex polyhedron  $\mathbf{P} \subseteq \mathbb{R}^3$  is called **Rupert** if a hole with the shape of a straight tunnel can be cut into it such that a copy of  $\mathbf{P}$  can be moved through this hole.

Theorem (Prince Rupert; Nieuwland, 1816; Scriba, 1968; Jerrard, Wetzel, Yuan, 2017)

All **Platonic solids** are **Rupert**.

Theorem (Chai, Yuan, Zamfirescu, 18; Hoffmann, 18; Lavau, 19; Steininger, Y. 22)

At least 9 **Archimedean solids** are **Rupert**.



# Summary and main result **H**: Deciding Rupertness *"It shows us 'what's out there'."*

## Definition

A convex polyhedron  $\mathbf{P} \subseteq \mathbb{R}^3$  is called **Rupert** if a hole with the shape of a straight tunnel can be cut into it such that a copy of  $\mathbf{P}$  can be moved through this hole.

Theorem (Prince Rupert; Nieuwland, 1816; Scriba, 1968; Jerrard, Wetzel, Yuan, 2017)

All **Platonic solids** are **Rupert**.

Theorem (Chai, Yuan, Zamfirescu, 18; Hoffmann, 18; Lavau, 19; Steininger, Y. 22)

At least 9 **Archimedean solids** are **Rupert**.

- [Steininger, Y., 22]: Practical algorithm and proof of algorithmic decidability.

## Summary and conclusion

- A** **Diagonals** of products of  $(1 - x_1 - \cdots - x_n)^R$  are **hypergeometric** functions.
- B** The generating functions of the Dubrovin-Yang-Zagier numbers are **algebraic**.
- C**  $\text{Iso}_R(t)$  is a quotient of **hypergeometric** functions and increasing. Thus the shape of a projection of the Clifford torus is uniquely determined by its isoperimetric ratio.
- D** We can compute the  $N$ -th term of a  **$q$ -holonomic sequence** faster than previously.
- E** We can compute the  $N$ -th term of a **polynomial C-finite sequence** faster.
- F** The  $q$ -analogue of Pólya's theorem holds if and only if  $q$  is a root of unity.
- G** A **C-finite sequence** is a **constant term** iff it has 1 characteristic root  $\lambda$  and  $\lambda \in \mathbb{Q}$ .
- H** Rupertness is decidable and the truncated icosidodecahedron is Rupert.

## Perspectives and open questions

*"Curiouser and curiouser!"*

- A?** Describe **Diagonals** among **D-finite** functions.
- B?** Given a **D-finite** function, how to prove or disprove that it is **algebraic** in practice?
- C?** Given a **D-finite** function/**P-finite** sequence, how to prove that it is increasing?
- D?** Compute  $N$ -th terms in some **P-finite** sequences faster than in  $\tilde{O}(\sqrt{N})$  operations.
- E?** Compute the  $N$ -th term of an integer **C-finite sequence** in  $O(N)$  bit complexity.
- F?** Does there exist a suitable notion of " $q$ -algebraicity"?
- G?** Describe **Constant terms** among **Diagonals** or **P-finite** sequences.
- H?** Prove or disprove that the Rhombicosidodecahedron is Rupert.

*And many, many more...*

## Bonus: Definition of ${}_pF_q$ and algebraicity

The generalized **hypergeometric function** with parameters  $a_1, \dots, a_p$  and  $b_1, \dots, b_q$  is:

$${}_pF_q([a_1, \dots, a_p]; [b_1, \dots, b_q]; t) := \sum_{j \geq 0} \frac{(a_1)_j \cdots (a_p)_j}{(b_1)_j \cdots (b_q)_j} \frac{t^j}{j!},$$

where  $(x)_n := x \cdot (x+1) \cdots (x+n-1)$  is the rising factorial.

- [Fürnsinn, Y., 2023] Can also handle the case:  $a_j, b_k \notin \mathbb{Q}$  and  $a_j - b_k \in \mathbb{Z}$ .



## Bonus: Definition of ${}_pF_q$ and algebraicity

Theorem (Christol, 1986 and Beukers, Heckman, 1989)

Assume that the rational parameters  $\{a_1, \dots, a_p\}$  and  $\{b_1, \dots, b_{p-1}, b_p = 1\}$  are disjoint modulo  $\mathbb{Z}$ . Let  $N$  be their common denominator. Then

$${}_pF_{p-1}([a_1, \dots, a_p], [b_1, \dots, b_{p-1}]; t) \quad \text{is}$$

- **algebraic** if and only if for all  $1 \leq r < N$  with  $\gcd(r, N) = 1$  the numbers  $\{\exp(2\pi i r a_j), 1 \leq j \leq p\}$  and  $\{\exp(2\pi i r b_j), 1 \leq j \leq p\}$  interlace on the unit circle.
- **globally bounded** if and only if for all  $1 \leq r < N$  with  $\gcd(r, N) = 1$ , one encounters more numbers in  $\{\exp(2\pi i r a_j), 1 \leq j \leq p\}$  than in  $\{\exp(2\pi i r b_j), 1 \leq j \leq p\}$  when running through the unit circle from 1 to  $\exp(2\pi i)$ .
- [Fürnsinn, Y., 2023] Can also handle the case:  $a_j, b_k \notin \mathbb{Q}$  and  $a_j - b_k \in \mathbb{Z}$ .

## Bonus: DYZ-like numbers

### Zagier's problem

Find  $(\alpha, \beta) \in \mathbb{Q}^* \times \mathbb{Q}^*$  such that  $u_n \cdot (\alpha)_n \cdot (\beta)_n \cdot \gamma^n \in \mathbb{Z}$  for some  $\gamma \in \mathbb{Z}^*$ .  
 $(x)_n := x \cdot (x+1) \cdots (x+n-1)$ .

#	$u$	$v$	ODE order	degree	#	$u$	$v$	ODE order	degree
$a_n$	3/5	4/5	2	120	$f_n$	19/60	49/60	4	155520
$b_n$	2/5	9/10	4	120	$g_n$	19/60	59/60	4	46080
$c_n$	1/5	4/5	2	120	$h_n$	29/60	49/60	4	46080
$d_n$	7/30	9/10	4	155520	$i_n$	29/60	59/60	4	155520
$e_n$	9/10	17/30	4	155520					

### Theorem (Bostan, Weil, Y., 2023)

The sequences  $(a_n)_{n \geq 0}$ ,  $(b_n)_{n \geq 0}$ ,  $(c_n)_{n \geq 0}$ ,  $\dots$ ,  $(i_n)_{n \geq 0}$  are solutions to Zagier's problem.

- Estimates for degrees based on numerical monodromy group computations.
- Proof of **algebraicity**: Done:  $a_n, b_n, c_n$ . In progress:  $d_n, e_n, f_n, g_n, h_n, i_n$ .