

# MODEL BY VOIDJACKLEE



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#### Header

```
#include <stdio.h>
#include <stdarg.h>
#include <stdlib.h>
#include <math.h>
#include <string.h>
#include <time.h>
#include <limits.h>
#include <vector>
#include <list>
#include <set>
#include <utility> // pair
#include <map>
#include <iostream>
#include <sstream>
#include <algorithm> // sort
#include <functional>
#include <string>
#include <stack>
#include <queue>
#include <fstream>
#include <bitset>
//#include <unordered map>
//#include <unordered_set>
using namespace std;
#define ll long long
#define lll __int128
#define uchar unsigned char
#define ushort unsigned short
#define uint unsigned int
#define ulong unsigned long
#define ull unsigned long long
#define INT INF 0x7fffffff
#define pi acos(-1)
#define mem(a,b) memset(a,b,sizeof(a))
#define memn(a,b,c,n) memset(a,b,sizeof(c)*(n))
#define fre(a) freopen(a,"r",stdin)
#define cio ios::sync_with_stdio(false); // Do not use it with "scanf"
and other c input!
#define pb push back
#define mpair make_pair
```

```
#define rep(i,a,b) for(int i=a;i<=b;i++)</pre>
#define pre(i,a,b) for(int i=a;i>=b;i--)
#define REP(i,a,b) for(int i=a;i<b;i++)</pre>
#define reada(a,s,n) rep(i,s,n)scanf("%d",a + i);
#define READa(a,s,n) REP(i,s,n)scanf("%d",a + i);
#define readall(a,s,n) rep(i,s,n)scanf("%lld",a + i);
#define READall(a,s,n) REP(i,s,n)scanf("%lld",a + i);
//template <typename _Tp> inline void read(_Tp&x) {
      char ch;bool flag=0;x=0;
//
//
      ch=getchar();
//
      while(!isdigit(ch)){if(ch=='-')flag=1;ch=getchar();}
//
      while(isdigit(ch))x=x*10+ch-'0',ch=getchar();
//
      if(flag)x=-x;
//}
//inline void print lll(lll x) {
     if(x<0) {x=-x;putchar('-');}</pre>
//
//
      if(x>9) print lll(x/10);
      putchar(x%10+'0');
//
//}
#define _T_(T) int T;scanf("%d",&T);while(T--)
#define T T (TTESTCASES)
#define _E_(T) while(~T)
#define C(a) cout << a << endl;</pre>
#define dsci(a) int a;scanf("%d",&a)
#define dscii(a,b) int a,b;scanf("%d%d",&a,&b)
#define dsciii(a,b,c) int a,b,c;scanf("%d%d%d",&a,&b,&c)
#define dscl(a) 11 a;scanf("%11d",&a)
#define dscll(a,b) ll a,b;scanf("%11d%11d",&a,&b)
#define dsclll(a,b,c) ll a,b,c;scanf("%lld%lld%lld",&a,&b,&c)
#define dscd(a) double a;scanf("%lf",&a)
#define dscdd(a,b) double a,b;scanf("%lf%lf",&a,&b)
#define dscddd(a,b,c) double a,b,c;scanf("%lf%lf%lf",&a,&b,&c)
#define sci(a) scanf("%d",&a)
#define scii(a,b) scanf("%d%d",&a,&b)
#define sciii(a,b,c) scanf("%d%d%d",&a,&b,&c)
#define scl(a) scanf("%lld",&a)
#define scll(a,b) scanf("%lld%lld",&a,&b)
#define sclll(a,b,c) scanf("%11d%11d%11d",&a,&b,&c)
#define scd(a) scanf("%lf",&a)
#define scdd(a,b) scanf("%lf%lf",&a,&b)
#define scddd(a,b,c) scanf("%lf%lf%lf",&a,&b,&c)
```

```
#define lowbit(x) ((x)&(-(x)))
#define Tprint(a,s,e) REP(i,s,e){if(i!=s)printf(" ");printf("%lld",a
[i]);}
#define endl '\n'
#define itn int
#define iny int
#define nit int
#define inr int
#define mian main
#define iman main
#define mina main
#define mian main
#define ednl endl
#define fro for
#define fir for
#define reutrn return
#define return return
#define reutnr return
#define re0 return 0
#define re1 return 1
```

# 差分&前缀和

### 差分

```
int p[1000] = {0};
int a[1000];
int n;

void pls(int l,int r,int k)
{
    p[1] += k;
    p[r + 1] -= k;
}

void init()
{
    REP(i,0,n) {
        pls(i,i,a[i]);
    }
}
int main()
```

```
{
    scanf("%d",&n);
    READ(a,0,n);
    init();
    Tprint(p,0,n);
    printf("\n");
    int q;
    scanf("%d",&q);
    int 1,r;
   while (q --) {
        scanf("%d%d",&1,&r);
        pls(1,r,1);
    }
    int s[1000];
    s[0] = p[0];
    rep(i,1,n - 1) {
        s[i] = s[i - 1] + p[i];
    }
    rep(i,0,n - 1) printf("%d ",s[i]);
    return 0;
}
二维差分
int p[1000][1000] = {0};
int a[1000][1000];
void pls(int a,int b,int x,int y,int k)
{
    p[a][b] += k;
    p[x + 1][y + 1] += k;
    p[x][y + 1] -= k;
    p[x + 1][y] -= k;
}
int main()
{
    int n,m;
    scanf("%d%d",&n,&m);
    rep(i,1,n) {
        rep(j,1,m) {
            scanf("%d",&a[i][j]);
        }
```

```
}
    rep(i,1,n) {
        rep(j,1,m) {
            pls(i,j,i,j,a[i][j]);
        }
    }
    rep(i,1,n) {
        rep(j,1,m) {
            printf("%d ",p[i][j]);
        printf("\n");
    printf("\n");
    int sum[1000][1000] = {0};
    sum[1][1] = p[1][1];
    rep(i,1,n) {
        rep(j,1,m) {
            sum[i][j] = sum[i - 1][j] + sum[i][j - 1] - sum[i - 1][j -
1] + p[i][j];
        }
    }
   rep(i,1,n) {
        rep(j,1,m) {
            printf("%d ",sum[i][j]);
        printf("\n");
    }
    return 0;
}
二维前缀和
int main()
{
    int arr[1000][1000] = {0};
    int sum[1000][1000] = \{0\};
    int n,m;
    scanf("%d%d",&n,&m);
    rep(i,1,n) {
        rep(j,1,m) {
            scanf("%d",&arr[i][j]);
        }
```

```
}
    rep(i,1,n) {
        rep(j,1,m) {
            printf("%d ",arr[i][j]);
        printf("\n");
    }
    puts("");
    sum[1][1] = arr[1][1];
     rep(i,1,m) sum[0][i] = sum[0][i - 1] + arr[0][i];
      rep(i,1,n) sum[i][0] = sum[i - 1][0] + arr[i][0];
    rep(i,1,n) {
        rep(j,1,m) {
            sum[i][j] = sum[i - 1][j] + sum[i][j - 1] - sum[i - 1][j -
1] + arr[i][j];
        }
    }
   // sum up
    rep(i,1,n) {
        rep(j,1,m) {
            printf("%d ",sum[i][j]);
       printf("\n");
    }
   // obtain
    int q;
    scanf("%d",&q);
    int a,b,x,y;
   while (q --) {
        scanf("%d%d%d%d",&a,&b,&x,&y);
       printf("%d\n",sum[x][y] - sum[a - 1][y] - sum[x][b - 1] + sum[a]
 - 1][b - 1]);
    }
    return 0;
}
```

# 树状数组

### 单点修改&区间查询

```
// 树状数组:单点修改,区间和查询
const int MAXN = 5e5 + 10;
11 a[MAXN];
11 c[MAXN];
int n,m;
void modify(int i,ll x)
{
   // 从叶子结点一路向上更新
   for (;i <= n;i += lowbit(i)) {</pre>
       c[i] += x;
   }
}
11 sum(int i)
   // 查询: 由于每个c结点相当于一小段前缀和,因此全+起来,最后求得的便是总
共的前缀和
   11 \text{ ans} = 0;
   for (;i > 0;i -= lowbit(i))
       ans += c[i];
   return ans;
}
int main()
   scii(n,m);
   rep(i,1,n) scl(a[i]);
   rep(i,1,n) modify(i, a[i]);
   int t,x,y;
   while (m --) {
       sciii(t,x,y);
       if (t == 1) {
           // modify
           modify(x, y);
       } else if (t == 2) {
           // query
           printf("%1ld\n",sum(y) - sum(x - 1));
       }
```

```
}
re0;
}
```

## (+差分) 区间修改 & 单点查询

```
// 树状数组+差分: 区间修改+单点查询
const int MAXN = 5e5 + 10;
int a[MAXN];
int c[MAXN];
int n,m;
void modify(int idx,ll x)
{
   // 从叶子结点一路向上更新
   for (int i = idx;i <= n;i += lowbit(i))</pre>
       c[i] += x;
   }
}
11 sum(int idx)
   // 查询: 由于每个c 结点相当于一小段前缀和, 因此全+起来, 最后求得的便是总
共的前缀和
   11 ans = 0;
   for (int i = idx;i > 0;i -= lowbit(i))
       ans += c[i];
   return ans;
}
void pls(int l,int r,int k)
   modify(l, k);
   modify(r + 1, -k);
}
void init()
{
   rep(i,1,n) {
       pls(i,i,a[i]);
}
```

```
int main()
{
   scii(n,m);
   rep(i,1,n) sci(a[i]);
   init();
   int t;
   int x,y,k;
   while (m --) {
       sci(t);
       if (t == 1) {
           sciii(x,y,k);
           // modify
           pls(x,y,k);
       } else if (t == 2) {
           sci(x);
           // get_ans: sum_up
           printf("%lld\n",sum(x));
       }
   }
   return 0;
}
求最值
int n;
const int MAXN = 1e5 + 10;
int a[MAXN]; // 原数组
11 c[MAXN]; // 求和树状数组
void modify(int i,ll x)
{
   // 从叶子结点一路向上更新
   for (;i <= n;i += lowbit(i)) {</pre>
       c[i] += x;
}
11 sum(int i)
   // 查询: 由于每个c结点相当于一小段前缀和,因此全+起来,最后求得的便是总
共的前缀和
   11 ans = 0;
   for (;i > 0;i -= lowbit(i))
```

```
ans += c[i];
    }
    return ans;
}
int mx[MAXN]; // 最大值数状数组
void modify_m(int i,int x)
{
    int low;
    a[i] = x; // 会直接修改数组的值
    for (;i <= n;i += lowbit(i)) {</pre>
        mx[i] = a[i];
        low = lowbit(i);
        for (int j = 1; j < low; j <<= 1) {</pre>
            mx[i] = max(mx[i], mx[i - j]);
        }
    }
}
int query_max(int l,int r)
{
    int ans = max(a[1],a[r]);
   while(true)
        ans = max(ans, a[r]);
        if (1 == r) break;
        r --;
        for (;r - 1 > lowbit(r);r -= lowbit(r))
            ans = max(ans, mx[r]);
        }
    }
    return ans;
}
int main()
    int m;
    scii(n,m);
    rep(i,1,n) {
        sci(a[i]);
        modify_m(i, a[i]); // 修改最大值
        modify(i, a[i]);
    }
    int t,x,y;
   while (m --) {
        sciii(t,x,y);
        if (t == 1) {
```

```
modify(x, y - a[x]);
           modify_m(x, y);
       } else if (t == 2) {
           printf("%lld\n",sum(y) - sum(x - 1));
       } else printf("%d\n",query_max(x, y));
   return 0;
}
求逆序数
// 树状数组求逆序数
const int MAXN = 5e4 + 10;
int a[MAXN];
int b[MAXN];
11 c[MAXN];
int n,m;
void modify(int idx,ll x)
{
   // 从叶子结点一路向上更新
   for (int i = idx;i <= n;i += lowbit(i)) {</pre>
       c[i] += x;
   }
}
11 sum(int idx)
   // 查询: 由于每个c结点相当于一小段前缀和,因此全+起来,最后求得的便是总
共的前缀和
   11 \text{ ans} = 0;
   for (int i = idx;i > 0;i -= lowbit(i))
       ans += c[i];
   return ans;
}
int main()
   int *x;
   ll ans;
   int idx;
   while (~sci(n)) {
       rep(i,1,n) {
```

```
sci(a[i]);
            b[i] = a[i];
            c[i] = 0;
        }
        sort(a + 1, a + 1 + n);
        x = unique(a + 1, a + 1 + n);
//
          for (int *i = a + 1;i != x;i ++) printf("%d ",*i);
//
          puts("");
        ans = 0;
        pre(i,n,1) {
            idx = (int) (lower_bound(a + 1, x, b[i]) - a);
            ans += sum(idx);
            modify(idx, 1);
        printf("%11d\n",ans);
    }
    re0;
}
```

# 线段树

```
// 线段树 - 二叉树,节点存的是一个 L,r,区间的内容 n
const int MAXN = 1e5 + 10;
struct Node {
    int 1,r;
    11 mx;
    11 mn;
    11 sum;
    int lazy;
    11 lzn;
} tree[MAXN << 2];</pre>
11 a[MAXN];
void push_up(int i)
{
    tree[i].sum = tree[i << 1].sum + tree[i << 1 | 1].sum;
    tree[i].mn = min(tree[i << 1].mn, tree[i << 1 | 1].mn);</pre>
    tree[i].mx = max(tree[i << 1].mx,tree[i << 1 | 1].mx);</pre>
}
void push_down(int i) //下推标记
{
    if (tree[i].lazy == 1) {
        tree[i << 1].sum += (tree[i << 1].r - tree[i << 1].l + 1) * tre
e[i].lzn;
        tree[i \leftrightarrow 1 \mid 1].sum += (tree[i \leftrightarrow 1 \mid 1].r - tree[i \leftrightarrow 1 \mid 1].
```

```
1 + 1) * tree[i].lzn;
        tree[i << 1].mx += tree[i].lzn;</pre>
        tree[i << 1 | 1].mx += tree[i].lzn;
        tree[i << 1].mn += tree[i].lzn;
        tree[i << 1 | 1].mn += tree[i].lzn;
        tree[i << 1].lzn += tree[i].lzn;</pre>
        tree[i << 1 | 1].lzn += tree[i].lzn;</pre>
        tree[i << 1].lazy = tree[i].lazy;</pre>
        tree[i << 1 | 1].lazy = tree[i].lazy;</pre>
        tree[i].lazy = 0;
        tree[i].lzn = 0;
    } else if (tree[i].lazy == 2) {
        tree[i << 1].sum = (tree[i << 1].r - tree[i << 1].l + 1) * tree
[i].lzn;
        tree[i << 1 | 1].sum = (tree[i << 1 | 1].r - tree[i << 1 | 1].l
+ 1) * tree[i].lzn;
        tree[i << 1].mx = tree[i].lzn;</pre>
        tree[i << 1 | 1].mx = tree[i].lzn;
        tree[i << 1].mn = tree[i].lzn;</pre>
        tree[i << 1 | 1].mn = tree[i].lzn;
        tree[i << 1].lzn = tree[i].lzn;</pre>
        tree[i << 1 | 1].lzn = tree[i].lzn;
        tree[i << 1].lazy = tree[i].lazy;</pre>
        tree[i << 1 | 1].lazy = tree[i].lazy;</pre>
        tree[i].lazy = 0;
        tree[i].lzn = 0;
    }
}
// i - 二叉树节点编号,调用时取1
// L, r 区间左右端下标,调用的时候取最大范围即可 build(1,n,1);
void build(int l,int r,int i)
{
    tree[i].l = l;
    tree[i].r = r;
    tree[i].lazy = 0;
    tree[i].lzn = 0;
    if (1 == r) {
```

```
tree[i].sum = a[1]; // a 原数组,把原来的数值给叶子结点
       tree[i].mx = a[1];
       tree[i].mn = a[1];
        return;
    }
   int m = (1 + r) >> 1;
   build(1,m,i << 1);
   build(m + 1,r,i << 1 | 1);
   push up(i);
}
void add(int l,int r,ll x,int i) // 将区间[l,r]整个加上x,调用(l,r,x,1)
{
   if (1 <= tree[i].1 && r >= tree[i].r) {
       tree[i].sum += (tree[i].r - tree[i].l + 1) * x;
       tree[i].mx += x;
       tree[i].mn += x;
       tree[i].lzn += x;
       tree[i].lazy = 1;
        return;
   }
   push_down(i);
   int m = (tree[i].l + tree[i].r) >> 1;
   if (1 <= m) add(1,r,x,i << 1);
   if (r > m) add(1,r,x,i << 1 | 1);
   push up(i);
}
void modify(int l,int r,ll x,int i) // 将区间[l,r]直接变成x,调用(l,r,x,
1)
{
   if (1 <= tree[i].1 && r >= tree[i].r) {
        tree[i].sum = (tree[i].r - tree[i].l + 1) * x;
       tree[i].mx = x;
       tree[i].mn = x;
       tree[i].lzn = x;
       tree[i].lazy = 2;
       return;
   }
   push_down(i);
   int m = (tree[i].l + tree[i].r) >> 1;
   if (1 <= m) modify(1,r,x,i << 1);</pre>
   if (r > m) modify(l,r,x,i << 1 | 1);
   push_up(i);
}
ll query(int l,int r,int i) //查询
```

```
{
    if (1 <= tree[i].1 && r >= tree[i].r){
        return tree[i].sum;
//
          return tree[i].mx;
//
          return tree[i].mn;
    push down(i);
    int m = (tree[i].l + tree[i].r) >> 1;
    11 sum = 0;
      LL mx = 0;
//
    LL mn = INT_INF;
    if (1 <= m) {
          sum += query(l,r,i << 1);</pre>
//
           mx = max(mx, query(l, r, i << 1));
//
           mn = min(mn, query(l, r, i << 1));</pre>
    if (r > m) {
        sum += query(l,r,i \ll 1 \mid 1);
//
          mx = max(mx, query(l, r, i << 1 | 1));
          mn = min(mn, query(l, r, i << 1 | 1));</pre>
//
    }
    return sum;
//
     return mx;
//
      return mn;
}
int main()
{
    dscii(n,m);
    rep(i,1,n) scl(a[i]);
    build(1, n, 1);
    int k,a,b;
    11 c;
    while (m --) {
        sciii(k,a,b);
        if (k == 1) {
             printf("%lld\n",query(a, b, 1));
        } else if (k == 2) {
             scl(c);
             modify(a, b, c, 1);
             rep(i,1,n) printf("a[%d]=%lld\n",i,query(i, i, 1));
        }
    }
    re0;
}
```

### **Utils**

#### gcd

```
11 gcd(11 a,11 b)
{
    if (a % b == 0) return b;
    return gcd(b, a % b);
}
exgcd
```

```
/// 欧几里得护展 exgcd
/// ax + by = gcd(a,b)
ll exgcd(ll a, ll b, ll &x, ll &y)
{
    if (b == 0)
    {
        x = 1;
        y = 0;
        return a;
    }
    ll r = exgcd(b, a % b, x, y);
    ll t = y;
    y = x - (a / b) * y;
    x = t;
    return r;
}
```

ax + by = gcd(a,b)的通解为:

$$\begin{cases} x = x_0 - c \frac{b}{\gcd(a, b)} \\ y = y_0 + c \frac{a}{\gcd(a, b)} \end{cases}$$
 c 是任意常数

#### quickpow

• 可以求 a 的逆元:quickpow(a % mod, mod - 2),注意 a 的范围,它也要%mod 过才不会爆精度。

```
const int mod = 1000000007;

ll quickpow(ll a, ll b)
{
    ll ans = 1;
    while (b)
```

```
if (b & 1) ans = a * ans % mod;
        a = a * a % mod;
       b >>= 1;
    return ans;
}
dfs 序
可以将树线段化。
const int MAXN = 1e5 + 10;
int dfn = 0;
int in[MAXN],out[MAXN];
void dfs(int k,int f)
    in[k] = ++ dfn;
    for (int i = g[k];~i;i = e[i].nxt)
        if (e[i].to != f) {
           dfs(e[i].to,k);
   out[k] = dfn;
}
```

# **Euler**

#### **Euler Function**

```
\varphi(x) = x(1 - \frac{1}{p_1})(1 - \frac{1}{p_2})...(1 - \frac{1}{p_n})
p_1, p_2, ..., p_n是 x 的所有质因子。

ll eular(ll x)
{
    ll ans = x;
    for (int i = 2;i * i <= x;i ++)
    {
        if (x % i == 0)
        }
```

```
ans = ans / i * (i - 1);
while (x % i == 0) x /= i;
}
if (x > 1) ans = ans / x * (x - 1);
return ans;
}

欧拉降幂(指数取余)

a^p \equiv a^{pmod\varphi(m)+\varphi(m)}(mod m)
const int mod = 1000000007;
ll e = eular(mod);
ll ans = quickpow(a % mod,b % e + e);
```

## **Matrix**

```
/// 矩阵模版 快速幂
const int mod = 1e9 + 7;
class Matrix
    const static int MAXN = 10;
    11 a[MAXN][MAXN];
public:
    int w;
    int h;
    Matrix(int h,int w):w(w),h(h)
    {
        mem(a,0);
    }
   Matrix(const char format[], ...) {
        va_list args;
        W = 0;
        h = 0;
        char buf[1000];
        va_start(args, format);
        vsprintf(buf, format, args);
        va_end(args);
```

```
stringstream ss(buf);
    stringstream num;
    11 n;
    string line;
    while (getline(ss, line)) {
        num.clear();
        num << line;</pre>
        W = 0;
        while (num >> n) {
            a[h + 1][++ w] = n;
        h ++;
    }
}
void E()
{
    if (w == h) {
        mem(a,0);
        rep(i,1,w) a[i][i] = 1;
}
void print()
{
    int f = 1;
    rep(i,1,h) {
        f = 1;
        rep(j,1,w) {
            if (f) f = 0;
            else printf(" ");
            printf("%11d",a[i][j]);
        puts("");
    }
}
void read_in()
{
    rep(i,1,h) {
        rep(j,1,w) {
            scanf("%11d",&a[i][j]);
        }
    }
}
Matrix operator* (const Matrix &B) const
```

```
if (w != B.h) return Matrix(0,0); // invalid
         Matrix ans(h,B.w);
         rep(i,1,h) {
             rep(j,1,B.w) {
                 rep(k,1,w) {
                        ans[i][j] = (ans[i][j] + a[i][k] * B[k][j]); // N
//
ot Moduled
                      ans[i][j] = (ans[i][j] + a[i][k] * B[k][j] % mod) %
 mod; // Moduled
                 }
         return ans;
    }
    const ll* operator[] (int i) const {
         return a[i];
    }
    11* operator[] (int i) {
         return a[i];
    }
};
Matrix quickpow(Matrix a, 11 b)
{
    if (a.h != a.w) return Matrix(0,0); // invalid
    Matrix ans(a.h,a.w);
    rep(i,1,ans.h) ans[i][i] = 1; // Set ans matrix to E
    while (b)
         if (b & 1) ans = a * ans;
         a = a * a;
        b >>= 1;
    return ans;
}
sample of use:
https://www.luogu.com.cn/problem/P1939
                         a_x = \begin{cases} 1, x = 1, 2, 3 \\ a_x = a_{x-1} + a_{x-3}, x \ge 4 \end{cases}
/// Sample of build a Matrix to solve num sequent.
int main()
```

```
{
    Matrix p("1 0 1\n1 0 0\n0 1 0"); // Construct a matrix
    Matrix base("1\n1\n1");
    int n;
    __T {
        scanf("%d",&n);
        if (n <= 3) {
            printf("1\n");
            continue;
        }
        printf("%lld\n",(quickpow(p, n - 3) * base)[1][1]);
    }
    return 0;
}</pre>
```

# **Forward Star**

#### No weight

```
const int MAXN = 1e5 + 10;
struct Edge {
    int to;
    int nxt;
} e[MAXN * 2];
int g[MAXN]; // Please call init() to memset it to -1!
int cnt = 0;
void init(int n)
{
    cnt = 0;
    memn(g,-1,int,n);
}
void add_edge(int u,int v)
{
    e[cnt] = \{v,g[u]\};
    g[u] = cnt ++;
With weight
const int MAXN = 1e5 + 10;
struct Edge {
    int to;
    11 w;
```

```
int nxt;
} e[MAXN * 2];
int g[MAXN]; // Please call init() to memset it to -1!
int cnt = 0;

void init(int n)
{
    cnt = 0;
    memn(g,-1,int,n);
}

void add_edge(int u,int v,ll w)
{
    e[cnt] = {v,w,g[u]};
    g[u] = cnt ++;
}
```

# Dijkstra

#### **Build Graphic from Vector**

```
const int MAXN = 1e5 + 10;
int dis[MAXN];
int vis[MAXN];
struct ST {
    int n;
    int w;
    bool operator< (const ST &other) const {</pre>
        return w > other.w;
    }
};
vector<ST> g[MAXN];
void dij(int s) {
    mem(dis,-1);
    mem(vis,⊘);
    priority_queue<ST> q;
    q.push(\{s,dis[s] = \emptyset\});
    ST current;
    int k;
    while (!q.empty()) {
```

```
current = q.top();
    q.pop();
    if (vis[current.n]) continue;
    vis[current.n] = 1;

    for (auto to : g[current.n]) {
        k = current.w + to.w;
        if (dis[to.n] == -1 || dis[to.n] > k) {
            q.push({to.n,dis[to.n] = k});
        }
    }
}
```

#### **Build Graphic from Forward Star**

```
const int MAXN = 1e5 + 10;
struct Edge {
    int to;
    int w;
    int nxt;
} e[MAXN * 2];
int g[MAXN]; // Please call init() to memset it to -1!
int cnt = 0;
void init(int n)
    cnt = 0;
    memn(g,-1,int,n);
}
void add_edge(int u,int v,int w)
{
    e[cnt] = \{v,w,g[u]\};
    g[u] = cnt ++;
}
int dis[MAXN];
int vis[MAXN];
struct ST {
    int n;
    int w;
    bool operator< (const ST &other) const {</pre>
        return w > other.w;
    }
```

```
};
void dij(int s) {
    mem(dis,-1);
    mem(vis,0);
    priority_queue<ST> q;
    q.push(\{s,dis[s] = \emptyset\});
    ST current;
    int k,to;
    while (!q.empty()) {
        current = q.top();
        q.pop();
        if (vis[current.n]) continue;
        vis[current.n] = 1;
        for (int i = g[current.n];~i;i = e[i].nxt) {
            to = e[i].to;
            k = dis[current.n] + e[i].w;
            if (dis[to] == -1 || dis[to] > k) {
                q.push({to,dis[to] = k});
            }
        }
    }
}
int main()
    int n,m;
    int u, v, w;
    int s,t;
    __T {
        scanf("%d%d",&n,&m);
        init(n + 5);
        while (m --) {
            scanf("%d%d%d",&u,&v,&w);
            add_edge(u,v,w);
            add_edge(v,u,w);
        scanf("%d%d",&s,&t);
        dij(s);
        printf("%d\n",dis[t]);
    return 0;
}
```

# dij 变式,多状态

```
const int MAXN = 1e5 + 10;
struct ST {
    int n;
    double w;
    int s;
    bool operator< (const ST &other) const {</pre>
        return w > other.w;
    }
};
double dis[MAXN][3];
int vis[MAXN][3];
vector<ST> g[MAXN];
void dij(int s) {
    rep(i,1,MAXN) dis[i][0] = dis[i][1] = dis[i][2] = -1;
    mem(vis,0);
    priority_queue<ST> q;
    q.push({s,dis[s][0] = 0,0});
    inr step;
    ST current;
    double k,cost;
   while (!q.empty()) {
        current = q.top();
        q.pop();
        if (vis[current.n][current.s]) continue;
        vis[current.n][current.s] = 1;
        for (auto to : g[current.n]) {
            cost = to.w;
            // update costs from step(mode)
            // REP(i,0,current.s % 3) {
            // cost = 1 / (1 - cost);
            // }
            // cost = fabs(cost);
```

# 并查集

#### 路径压缩

{

```
const int MAXN = 1e5 + 10;
int find_set[MAXN];
int find(int a)
{
   if (find_set[a] == a) return a;
   return find_set[a] = find(find_set[a]);
}
inline void bind(int a,int b)
   find_set[find(a)] = find(b);
}
void init(int n)
{
   rep(i,0,n) {
       find_set[i] = i; // 每个种类初始状态只有自己一个点
   }
}
路径压缩+按秩合并
const int MAXN = 1e5 + 10;
int find_set[MAXN];
int depth[MAXN];
int find(int a)
```

```
if (find set[a] == a) return a;
   return find_set[a] = find(find_set[a]);
}
inline void bind(int a,int b)
   int x = find(a), y = find(b);
   if (depth[x] >= depth[y]) { // 如果 a 的 根的子树深度 比 b 的 根的子树深
度 大, 那 a 的根继续做根
       find set[y] = x; // 改变b 节点的根的根为a 的根
       if (depth[x] == depth[y]) { // 俩根深度一样
          if (x != y) depth[x] ++; // 作为a 的根,自然子树的深度++
       }
   } else find_set[x] = y;
}
void init(int n)
   rep(i,0,n) {
       find_set[i] = i; // 每个种类初始状态只有自己一个点
       depth[i] = 1; // 初始化秩
   }
}
```

# ST 表

```
const int MAXN = 100010;
int st[MAXN][20];
int a[MAXN];
int n,m;
inline int read_int()
{
   int x = 0,f = 1;
   char ch = getchar();
   while (!isdigit(ch)) {
      if (ch == '-') f = -1;
      ch = getchar();
   }
   while (isdigit(ch))
   {
      x = x * 10 + ch - 48;
      ch = getchar();
   }
   return x * f;
```

```
}
void init() {
   // 定义 st[i][j] 是从 i 开始,到 i + 2^j 这一段,即[i, i + 2^j]这一段中的
最大/小值
   rep(i,1,n) st[i][0] = a[i];
   for (int j = 1;(1 << j) <= n;j ++) { // 遍历所有的j,j是一个很小的数
字,最大值=Log2(n)
       rep(i,1,n - (1 << j) + 1) { // 在[1,n]区间范围内,确定 j 的情况下,
把所有的i 都遍历求值一遍
           st[i][j] = max(st[i][j-1], st[i+(1 << (j-1))][j-1]);
// 套公式
       }
   }
}
int query(int 1, int r)
{
   int x = log_2(r - l + 1);
   return \max(st[1][x], st[r - (1 << x) + 1][x]);
}
int main()
{
   scanf("%d%d",&n,&m);
   rep(i,1,n) {
       a[i] = read_int();
   init();
   int 1,r;
   while (m --) {
       scanf("%d%d",&1,&r);
       printf("%d\n",query(1, r));
   }
   return 0;
}
```

# 求树上一条边两边点的个数

```
const int MAXN = 1e5 + 10;
int all = 0;
int nn;

ll ch[MAXN];
ll a[MAXN];
```

```
struct Edge {
    int to;
    int nxt;
} e[MAXN * 2];
int g[MAXN]; // Please memset it to -1!!!!!!!
int cnt = 0;
void init(int n)
{
    cnt = 0;
   memn(g,-1,int,n);
}
void add_edge(int u,int v)
{
    e[cnt] = \{v,g[u]\};
    g[u] = cnt ++;
}
// 方法: 记录每个节点的子节点数,在用总结点数减去该边子节点数即可
void dfs(int n, int f)
{
    ch[n] ++;
    int to;
    for (int i = g[n];~i;i = e[i].nxt) {
        to = e[i].to;
        if (to == f) continue;
        dfs(to, n);
        ch[n] += ch[to];
    }
    for (int i = g[n];~i;i = e[i].nxt) {
        to = e[i].to;
        if (to == f) continue;
        a[all ++] = ch[to] * (nn - ch[to]);
    }
}
int main()
    int u, v;
   scanf("%d",&nn);
    init(nn + 5);
    rep(i,1,nn - 1) {
        scanf("%d%d",&u,&v);
```

```
add_edge(u,v);
add_edge(v,u);
}
mem(ch,0);
dfs(1,-1);
sort(a, a + all);
nn --;
ll sum = 0;
int i = 0;
while (nn) {
    sum += (nn --) * a[i ++];
}
printf("%lld\n",sum);
return 0;
}
```

### **Net Flow**

#### **Edmond-Karp**

```
/// 网络流-最大流模板 EK 算法 前向星
/// https://www.luogu.com.cn/problem/P3376
const int MAXN = 5000 + 10;
struct Edge {
    int to;
    11 w;
    int nxt;
} e[MAXN * 2];
int g[MAXN];
int cnt = 0;
void init(int n)
   cnt = 0;
   memn(g,-1,int,n);
}
void add_edge(int u,int v,ll w)
{
    e[cnt] = \{v,w,g[u]\};
   g[u] = cnt ++;
}
```

int pre[MAXN]; // record the previous node index and indicating a node
whether has been visited.

```
11 flow[MAXN];
int n,m;
11 getAgtPath(int s,int t) // method to find an augmented path
{
    // bfs
    queue<int> q;
    q.push(s);
    int current;
    memn(pre,-1,int,n + 5);
    flow[s] = INT_INF;
    pre[s] = 0; // sourse node has been visited
    int to;
   while (!q.empty()) {
        current = q.front();
        q.pop();
        if (current == t) break;
        for (int i = g[current];~i;i = e[i].nxt) {
            to = e[i].to;
            if (pre[to] != -1 || e[i].w <= 0) continue;</pre>
            pre[to] = i; // record the index of current edge
            flow[to] = min(flow[current],e[i].w);
            q.push(to);
        }
    }
    if (pre[t] != -1) return flow[t]; // return the terminal's flow
    return -1; // 404 not found
}
11 EKmaxFlow(int s,int t) // source -> terminal
{
    if (s == t) return INT INF;
    11 flow = 0;
    11 agt flow;
    int cur;
    int edge_idx;
   while ((agt_flow = getAgtPath(s, t)) != -1) { // find augmented pat
```

```
h until all flows are gone.
        // modify the w of path
        cur = t;
        while (cur != s) {
            edge_idx = pre[cur];
            e[edge idx].w -= agt flow;
            e[edge_idx ^ 1].w += agt_flow; // xor can get the reverse e
dge swiftly
            cur = e[edge_idx ^ 1].to;
        }
        flow += agt_flow; // add to ans
    }
    return flow;
}
int vis[210][210];
int main()
{
    int s,t;
    scanf("%d%d%d%d",&n,&m,&s,&t);
    init(n + 5);
    mem(vis,-1);
    int u,v,w;
    rep(i,1,m) {
        scanf("%d%d%d",&u,&v,&w);
        if (vis[u][v] == -1) { // record the vised edge
            add_edge(u, v, w);
            add_edge(v, u, 0); // add reversed edge, which w=0;
            vis[u][v] = cnt - 2;
        else {
            e[vis[u][v]].w += w; // increase the weight
        }
    printf("%11d\n",EKmaxFlow(s, t));
    return 0;
}
Dinic
/// 网络流-最大流模板 Dinic 算法 前向星
/// https://www.luogu.com.cn/problem/P3376
const int MAXN = 5010;
```

```
struct Edge {
   int to:
   11 w;
   int nxt;
} e[MAXN * 2];
int g[MAXN]; // Please call init() to memset it to -1!
int cnt = 0;
void init(int n)
{
   cnt = 0;
   memn(g,-1,int,n);
}
void add_edge(int u,int v,ll w)
   e[cnt] = \{v,w,g[u]\};
   g[u] = cnt ++;
}
int s,t,n,m;
int dis[MAXN];
int cur[MAXN]; // 替代 g 数组,记住上次 dfs 最后跑到的地方,优化,减少 dfs 的
跑的次数
int bfs()
   memn(dis,-1,int,n + 5);
   queue<int> q;
   q.push(s);
   dis[s] = 0;
   int to,current,k;
   while (!q.empty()) {
       current = q.front();
       q.pop();
       for (int i = g[current];~i;i = e[i].nxt) {
           to = e[i].to;
           k = dis[current] + 1;
           if (dis[to] == -1 && e[i].w > 0) { // 只有没有访问过的,且该
通路可以走(w > 0)
               dis[to] = k;
               if (to == t) return 1;
               q.push(to);
           }
```

```
}
    }
    return 0;
}
11 dfs(int node, 11 flow)
{
    if (node == t) return flow;
    int to;
    11 d;
    for (int &i = cur[node];~i;i = e[i].nxt) { // 改变i 的同时, cur[node]
的值也会被改变
        to = e[i].to;
        if (dis[node] + 1 == dis[to] && e[i].w > 0) {
            d = dfs(to,min(e[i].w,flow));
            if (d > 0) {
                e[i].w -= d;
                e[i ^1].w += d;
                return d;
            }
        }
    }
    return 0;
}
11 dinic()
    ll ans = 0;
    11 d;
   while (bfs()) {
        rep(i,1,n) cur[i] = g[i];
        while ((d = dfs(s,INT_INF)))
            ans += d;
    return ans;
}
int main()
{
    scanf("%d%d%d%d",&n,&m,&s,&t);
    int u,v;
    11 w;
    init(n + 5);
    while (m --) {
        scanf("%d%d%lld",&u,&v,&w);
        add_edge(u, v, w);
        add_edge(v, u, ∅);
    }
```

```
printf("%lld\n",dinic());
  return 0;
}
```

## Manacher

### 最长回文串

```
string Manacher(string &s)
{
   //改造字符串
   int n = (int) s.size();
   string res = "$#";
   for (int i = 0;i < n;i ++)</pre>
   {
       res += s[i];
       res += "#";
   }
   //数组
   n = (int) res.size();
   vector<int> P(n,0);
   int mi = 0, right = 0; //mi 为最大回文串对应的中心点, right 为该回文串能
达到的最右端的值
   int maxLen = 0, maxPoint = 0; //maxLen 为最大回文串的长度, maxPoint 为
记录中心点
   for (int i = 1;i < n;i ++)</pre>
   {
       P[i] = right > i ? min(P[2 * mi - i], right - i) : 1; // 关键句,
文中对这句以详细讲解
       while (res[i + P[i]] == res[i - P[i]]) {
          P[i] ++;
       }
       if (right < i + P[i]) //超过之前的最右端,则改变中心点和对应的最右
端
       {
          right = i + P[i];
          mi = i;
       if (maxLen < P[i]) //更新最大回文串的长度,并记下此时的点
          maxLen = P[i];
          maxPoint = i;
       }
   }
```

```
return s.substr((maxPoint - maxLen) / 2,maxLen - 1);
}
回文串个数
11 Manacher n(string &s)
   //改造字符串
   int n = (int) s.size();
   string res = "$#";
   for (int i = 0; i < n; i ++)
       res += s[i];
       res += "#";
   }
   //数组
   n = (int) res.size();
   vector<ll> P(n,∅);
   11 mi = 0, right = 0; //mi 为最大回文串对应的中心点, right 为该回文串能达
到的最右端的值
   ll maxLen = 0, maxPoint = 0; //maxLen 为最大回文串的长度, maxPoint 为记
录中心点
   11 ans = 0;
   for (ll i = 1;i < n;i ++)</pre>
   {
       P[i] = right > i ? min(P[2 * mi - i], right - i) : 1; // 关键句,
文中对这句以详细讲解
       while (res[i + P[i]] == res[i - P[i]]) {
           P[i] ++;
       if (right < i + P[i]) //超过之前的最右端,则改变中心点和对应的最右
端
       {
           right = i + P[i];
           mi = i;
       if (maxLen < P[i]) //更新最大回文串的长度,并记下此时的点
           maxLen = P[i];
           maxPoint = i;
       ans += P[i] / 2;
   return ans;
}
```

#### Min25

#### 快速求[1,1010]的质数和

```
const int MAXN = 1000010;
namespace Min25 {
    int prime[MAXN], id1[MAXN], id2[MAXN], flag[MAXN], ncnt, m;
    11 g[MAXN], sum[MAXN], a[MAXN], T, n;
    inline int ID(ll x) {
        return x \leftarrow T ? id1[x] : id2[n / x];
    }
    inline ll calc(ll x) {
        return x * (x + 1) / 2 - 1;
    }
    inline 11 f(11 x) {
        return x;
    }
    inline void init() {
        rep(i,0,MAXN - 1) {
            prime[i] = id1[i] = id2[i] = flag[i] = 0;
            g[i] = sum[i] = a[i] = 0;
        }
        ncnt = 0;
        m = 0;
        T = sqrt(n + 0.5);
        for (int i = 2; i <= T; i++) {</pre>
            if (!flag[i]) {
                prime[++ncnt] = i;
                 sum[ncnt] = sum[ncnt - 1] + i;
            }
            for (int j = 1; j <= ncnt && i * prime[j] <= T; j++) {</pre>
                flag[i * prime[j]] = 1;
                 if (i % prime[j] == 0) break;
            }
        }
        for (11 1 = 1; 1 \le n; 1 = n / (n / 1) + 1) {
            a[++m] = n / 1;
            if (a[m] \leftarrow T) id1[a[m]] = m; else id2[n / a[m]] = m;
            g[m] = calc(a[m]);
        }
```

#### rope

奇淫怪巧: rope 是一种类似块状链表操作的东西,速度很快,底层是可持续化平衡树实现。

只有gnu 编译器可以使用, clang 不行。

### 头文件加入

#include <ext/rope>

```
using namespace __gnu_cxx;
操作

push_back(x); //在末尾添加 x

insert(pos,x); //在 pos 插入 x, 自然支持整个 char 数组的一次插入

erase(pos,x); //从 pos 开始删除 x 个

copy(pos,len,x); //从 pos 开始到 pos+Len 为止用 x 代替

replace(pos,x); //从 pos 开始换成 x

substr(pos,x); //提取 pos 开始 x 个

at(x); //访问第 x 个元素

[x] //访问第 x 个元素
```