# Category Theory - Lecture 5 (Notes)

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**Exercise 0.1.** Let C be a category.

Let  $f, g: A \longrightarrow B$  and let  $u: A' \longrightarrow A$  be an isomorphism. Then,

$$f = g \Leftrightarrow f \circ u = g \circ u$$

*Proof.*  $(\Rightarrow)$ : It is clear.

 $(\Leftarrow)$ : If  $u:A'\longrightarrow A$  is an isomorphism, then there exists  $u^{-1}:A\longrightarrow A'$  such that

$$fu = gu \Rightarrow fuu^{-1} = guu^{-1}$$
$$\Rightarrow f = g$$

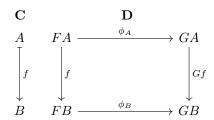
## 1 Natural Transformations

**Definition 1.1.** Let  $F, G : \mathbf{C} \longrightarrow \mathbf{D}$  be functors.

A natural transformation  $\phi: F \Rightarrow G$  is a family of maps:

$$\{\phi_A: FA \longrightarrow GA \mid A \in Ob(\mathbf{C})\}\$$

such that for all  $f: A \longrightarrow B$  in  $\mathbb{C}$ , we have



#### Example 1.1. Fix $n \in \mathbb{N}$

- ullet  $\mathbf{C} = CRing = category of commutative rings$
- $\mathbf{D} = category \ of \ monoids$
- $F: CRing \longrightarrow \underline{Mon}$  is the functor

$$F = M_n : CRing \longrightarrow \underline{Mon}$$

Here  $M_n(R)$  is  $n \times n$  matrix with entries in R. So  $\underline{Mon}$  is the category of monoids with  $n \times n$  matrices as objects.

$$R \longmapsto F \longrightarrow M_n(R)$$

$$\downarrow^f \qquad \qquad \downarrow^{M_n(f)}$$

$$S \longmapsto F \longrightarrow M_n(S)$$

•  $G: \underline{CRing} \longrightarrow \underline{Mon}$  is the functor

$$G = U : \underline{CRing} \longrightarrow \underline{Mon}$$

Here we have

$$(R, +, *, 0, 1) \longmapsto U \longrightarrow (R, +, 1)$$

$$\downarrow f \qquad \qquad \downarrow U(f)$$

$$(S, +, *, 0, 1) \longrightarrow U \longrightarrow (S, *, 1)$$

So U simply returns the underlying monoid of a given ring

Now we can define a natural transformation

$$det: M_n \Longrightarrow U$$

For  $R = (R, +, *, 0, 1) \in CRing$ , we need

$$det_R: FR \longrightarrow GR$$
  

$$det_R: M_n(R) \longrightarrow U(R)$$
  

$$det_R: M_n(R) \longrightarrow (R, *, 1)$$

Here  $det_R$  is a monoid homomorphism

This is monoid homomorphism because

$$\begin{cases} det_R(M*N) = det_R(M)*det_R(N) & (it preserves multiplication) \\ det_R(I_n) = 1_R & (it preserves identity) \end{cases}$$

So now need to check that this is a natural transformation. (i.e we need to check its naturality).

$$\frac{\mathbf{Cring}}{R} \qquad \frac{\mathbf{Mon}}{det_R} \longrightarrow R$$

$$\downarrow^f \qquad \downarrow^{M_n(f)} \qquad \downarrow^f$$

$$S \qquad M_n(S) \longrightarrow det_S \longrightarrow S$$

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \longmapsto det_R \longrightarrow ad - bc$$

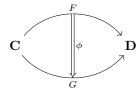
$$\downarrow^{M_n(f)} \qquad \qquad \downarrow^f$$

$$\begin{pmatrix} fa & fb \\ fc & fd \end{pmatrix} \longrightarrow f(ad - bc)$$

#### 1.1 Notation

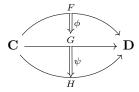
#### 1.1.1 Natural transformation diagram

For  $F,G:\mathbf{C}\longrightarrow\mathbf{D}$  and  $\phi:F\Rightarrow G$  as natural transformation, we write

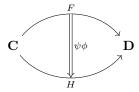


#### 1.1.2 Composition diagram

Given  $F, G, H : \mathbf{C} \longrightarrow \mathbf{D}$  and  $\phi, \psi : F \Rightarrow G$  as natural transformations, we write



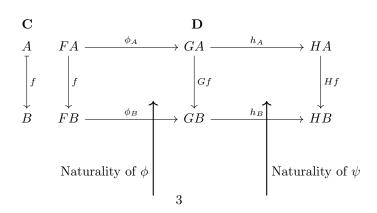
We define



Also

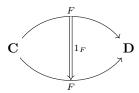
$$FA \xrightarrow{(\psi\phi)_A} HA \ =_{def} \ FA \xrightarrow{\phi_A} GA \xrightarrow{\psi_A} HA$$

So we can define composition



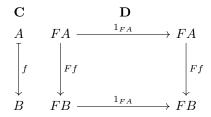
# 1.2 Identity natural transformation

For  $F: \mathbf{C} \longrightarrow \mathbf{D}$ , we can define



by 
$$FA \xrightarrow{1_{FA}} FA$$
 for all  $A \in \mathbf{C}$ .

Naturality of this is



### 1.3 Functor category

**Definition 1.2.** (Functor category)

For categories C, D, where C is a small category, we define the **functor category** [C, D] (sometimes denoted as  $D^C$ ) as having

• objects: functors  $F: \mathbf{C} \longrightarrow \mathbf{D}, G: \mathbf{C} \longrightarrow \mathbf{D}$ 

• maps: natural transformations  $\phi: F \Rightarrow G$ 

Exercise 1.1. Check the associativity and unit axioms for functor category.

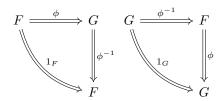
## 1.3.1 Isomorphisms in $D^{C}$ or [C, D]

Fix C, D.

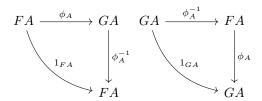
By definition, an isomorphism in  $[\mathbf{C}, \mathbf{D}]$  is  $\phi : F \Rightarrow G$  with an inverse

$$\phi^{-1}:G\Rightarrow F$$

This means that



This holds if and only if, for all  $A \in \mathbb{C}$ , we have



This implies  $\forall A \in \mathbf{C}, \ \phi_A : FA \longrightarrow GA$  is an isomorphism.

**Proposition 1.1.** Let  $\phi: F \Rightarrow G$  be a natural transformation.

Then  $\phi$  is an isomorphism if and only if for all  $A \in \mathbf{C}$ 

$$\phi_A: FA \longrightarrow GA$$

is an isomorphism.

Note 1.1. We need  $\phi$  to be natural in order to piece together all isomorphisms  $FA \longrightarrow GA$  for all  $A \in \mathbf{C}$