

Lecture 6

Quantum Fourier Transform and Phase Estimation

- Quantum Fourier transform
- Phase Estimation

1 Quantum Fourier transform

Hadamard transform: $|x\rangle \xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^n}} \sum_{y \in \{0,1\}^n} (-1)^{x \cdot y} |y\rangle$ where x is an integer modulo 2^n .
This is a Fourier transform over $\underbrace{\mathbb{Z}_2 \otimes \cdots \otimes \mathbb{Z}_2}_n$.

How about Fourier transform over \mathbb{Z}_{2^n} ? That has the form:

$$|x\rangle \mapsto \frac{1}{\sqrt{2^n}} \sum_{y \in \mathbb{Z}_{2^n}} e^{\frac{2\pi i xy}{2^n}} |y\rangle := |\tilde{x}\rangle$$

where $x \in \mathbb{Z}_{2^n}$ represents an integer modulo 2^n .

For $n = 1$, the transform is

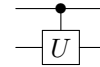
$$\begin{aligned} |0\rangle &\mapsto \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle), \quad |1\rangle \mapsto \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle) \\ |x\rangle &\xrightarrow{H} \frac{1}{\sqrt{2^n}} \sum_{y \in \{0,1\}} e^{\frac{2\pi i xy}{2^n}} \end{aligned}$$

In general, the $|\tilde{x}\rangle$ states form an orthonormal basis, the Fourier basis: $\langle \tilde{x} | \tilde{x}' \rangle = \delta_{x,x'}$.

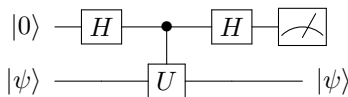
When do we need the quantum Fourier transform?

1.1 Phase estimation

- Given: Ability to implement a controlled unitary operator U and a quantum state $|\psi\rangle$ with $U|\psi\rangle = e^{i\theta}|\psi\rangle$.
- Problem: Learn θ .



1.2 Hadamard test



$$\begin{aligned} |0\rangle|\psi\rangle &\xrightarrow{H \otimes I} \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)|\psi\rangle \xrightarrow{\text{controlled-}U} \frac{1}{\sqrt{2}}(|0\rangle|\psi\rangle + |1\rangle U|\psi\rangle) \\ &= \frac{1}{\sqrt{2}}(|0\rangle + e^{i\theta}|1\rangle)|\psi\rangle \\ &\xrightarrow{H \otimes I} \frac{1}{2}[(|0\rangle + |1\rangle) + e^{i\theta}(|0\rangle - |1\rangle)]|\psi\rangle = \left(\frac{1+e^{i\theta}}{2}|0\rangle + \frac{1-e^{i\theta}}{2}|1\rangle\right)|\psi\rangle \end{aligned}$$

$$\begin{aligned}\Pr(0) &= \left| \frac{1 + e^{i\theta}}{2} \right|^2 = \frac{1}{4} [(1 + \cos \theta)^2 + \sin^2 \theta] & \Pr(1) &= \left| \frac{1 - e^{i\theta}}{2} \right|^2 = \sin^2 \frac{\theta}{2} \\ &= \frac{1}{4} [2 + 2 \cos \theta] = \cos^2 \frac{\theta}{2}\end{aligned}$$

If $\theta = 0$ or $\theta = \pi$, we learn θ perfectly: After we measure the state in the computational basis, we get $|0\rangle$ with probability 1 when $\theta = 0$ and $|1\rangle$ with probability 1 when $\theta = \pi$.

If $0 < \theta < \pi$, learn the probability distribution by samples to get information of θ .

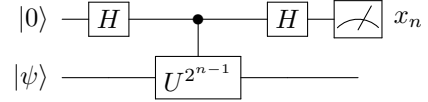
Now, consider $\theta = 2\pi \cdot \sum_{j=1}^n \frac{x_j}{2^j}$, $x_j \in \{0, 1\}$. i.e., $\theta = 2\pi \cdot 0.\overbrace{x_1 \dots x_n}^{n \text{ bits}}$

Consider what happens if we apply U^{2^k} .

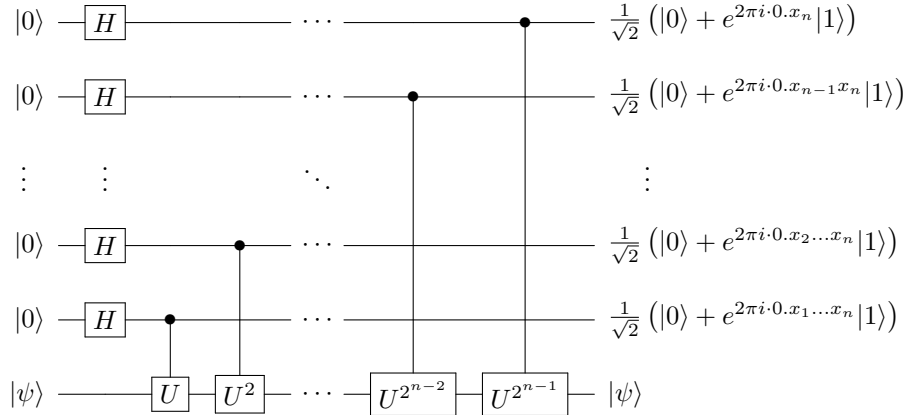
$$U^{2^k} |\psi\rangle = e^{i2^k \theta} |\psi\rangle = e^{2\pi i \cdot x_1 \dots x_k \cdot x_{k+1} \dots x_n} |\psi\rangle = e^{2\pi i \cdot 0.x_{k+1} \dots x_n} |\psi\rangle$$

In particular, $U^{2^{n-1}} |\psi\rangle = e^{2\pi i \cdot 0.x_n} |\psi\rangle$.

In other words, $U^{2^{n-1}} |\psi\rangle = \begin{cases} |\psi\rangle & x_n = 0 \\ -|\psi\rangle & x_n = 1 \end{cases}$



Idea: Combine n such experiments for different exponents of U :



The output state is:

$$\begin{aligned}& \bigotimes_{i=1}^n \frac{1}{\sqrt{2}} (|0\rangle + e^{2\pi i \cdot 0.x_{n+1-i} \dots x_n} |1\rangle) \\ &= \frac{1}{\sqrt{2^n}} \sum_{y=0}^{2^n-1} e^{\frac{2\pi i \cdot 2^{n-1} x y_{n-1}}{2^n}} |y_{n-1}\rangle e^{\frac{2\pi i \cdot 2^{n-2} x y_{n-2}}{2^n}} |y_{n-2}\rangle \dots e^{\frac{2\pi i x y_0}{2^n}} |y_0\rangle \\ &= \frac{1}{\sqrt{2^n}} \sum_{y=0}^{2^n-1} e^{\frac{2\pi i x y}{2^n}} |y\rangle = |\tilde{x}\rangle\end{aligned}$$

Conclusion: Phase estimation can be solved by **inverse** of quantum Fourier transform.

Note: $QFT : |x\rangle \longrightarrow |\tilde{x}\rangle = \frac{1}{\sqrt{2^n}} \sum_{y=0}^{2^n-1} e^{\frac{2\pi i x y}{2^n}} |y\rangle$ $x \in \{0, 1, \dots, 2^n - 1\}$

For phase estimation: Ideally, we want to have x_1, \dots, x_n directly as outputs. How do we implement this?

The first qubit is $\frac{|0\rangle + e^{2\pi i \cdot 0 \cdot x_n} |1\rangle}{\sqrt{2}} = H |x_n\rangle$. Therefore, applying H reveals x_n .

The second qubit is $\frac{|0\rangle + e^{2\pi i \cdot 0 \cdot x_{n-1} x_n} |1\rangle}{\sqrt{2}}$:

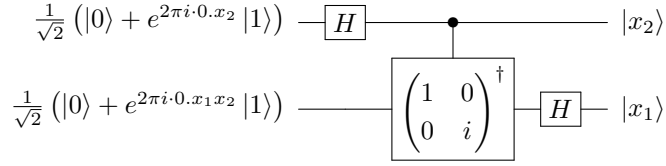
$$e^{2\pi i \cdot 0 \cdot x_{n-1} x_n} = e^{2\pi i \cdot 0 \cdot x_{n-1}} \cdot e^{2\pi i \cdot 0 \cdot 0 x_n} = e^{x_{n-1} \cdot \pi i} \cdot e^{x_n \cdot \frac{\pi i}{2}} = (-1)^{x_{n-1}} \cdot i^{x_n}.$$

We can remove the dependence on x_n :

If $x_n = 0$, do I ; if $x_n = 1$, do $\begin{pmatrix} 1 & 0 \\ 0 & i \end{pmatrix}^\dagger$.

Then the state becomes $\frac{1}{\sqrt{2}} (|0\rangle + (-1)^{x_{n-1}} |1\rangle)$. Then Hadamard gate H reveals x_{n-1} .

Inverse QFT with $n = 2$:

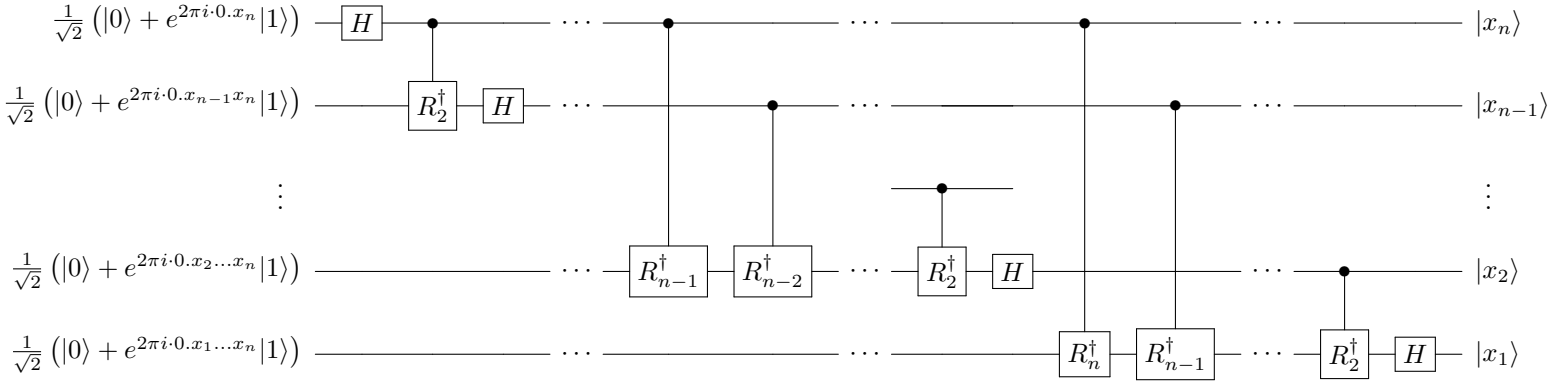


For convenience, denote $R_k = \begin{pmatrix} 1 & 0 \\ 0 & e^{\frac{2\pi i}{2^k}} \end{pmatrix}$.

Move generally, since $e^{2\pi i \cdot 0 \cdot x_{n-k+1} \dots x_n} = e^{2\pi i (\frac{x_{n-k+1}}{2^1} + \frac{x_{n-k+2}}{2^2} + \dots + \frac{x_n}{2^k})}$, the k^{th} qubit is

$$\frac{1}{\sqrt{2}} (|0\rangle + e^{2\pi i \cdot 0 \cdot x_{n-k+1} \dots x_n} |1\rangle) = R_k^{x_n} \dots R_3^{x_{n-k+3}} R_2^{x_{n-k+2}} H |x_{n-k+1}\rangle.$$

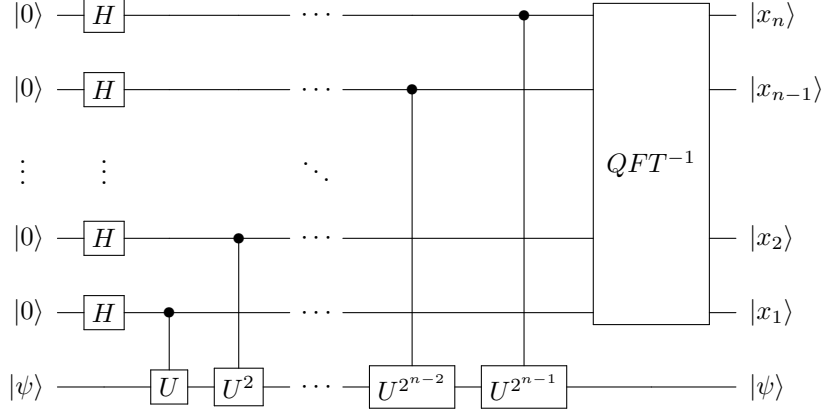
The circuit of inverse QFT:



Gate complexity of QFT^{-1} : $O(n^2)$

2 Phase Estimation

Big picture of phase estimation:



If $U|\psi\rangle = e^{2\pi i \cdot 0.x_1 \dots x_n} |\psi\rangle$, this works. How about $e^{i\varphi}$ for general $\varphi \in [0, 2\pi)$?

$$\begin{aligned}
 |0\rangle^{\otimes n} &\xrightarrow{H^{\otimes n}} \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |x\rangle & x = \sum_{i=1}^n x_i \cdot 2^{n-i} & \quad |x_i\rangle \xrightarrow{c-U^{2^{i-1}}} \begin{cases} |x_i\rangle & \text{if } x_i = 0 \\ e^{i\varphi 2^{i-1}} |x_i\rangle & \text{if } x_i = 1 \end{cases} \\
 &\xrightarrow{c-U^{\varphi}} \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} e^{i\varphi x} |x\rangle & |x_1 \dots x_n\rangle &\mapsto \prod_{i=1}^n e^{i\varphi x_i 2^{i-1}} |x_i\rangle
 \end{aligned}$$

The QFT^{-1} is $\sum_{x=0}^{2^n-1} |x\rangle \langle \tilde{x}| = \frac{1}{\sqrt{2^n}} \sum_{x,y=0}^{2^n-1} e^{-\frac{2\pi i xy}{2^n}} |x\rangle \langle y|$.

$$\begin{aligned}
 \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} e^{i\varphi x} |x\rangle &\xrightarrow{QFT^{-1}} \frac{1}{2^n} \sum_{x,y=0}^{2^n-1} e^{i\varphi x} e^{-\frac{2\pi i xy}{2^n}} |y\rangle \\
 &= \sum_{y=0}^{2^n-1} \alpha_y |y\rangle \quad \alpha_y := \frac{1}{2^n} \sum_{x=0}^{2^n-1} e^{i-\frac{2\pi y}{2^n}x} e^{i\varphi x} = \frac{1}{2^n} \sum_{x=0}^{2^n-1} e^{i(\varphi - \frac{2\pi y}{2^n})x}.
 \end{aligned}$$

where α_y is a geometric series. Denote $\tilde{\varphi} = \varphi - \frac{2\pi y}{2^n}$:

$$\sum_{x=0}^{2^n-1} e^{i\tilde{\varphi}x} = \frac{1 - e^{i\tilde{\varphi}2^n}}{1 - e^{i\tilde{\varphi}}} = \frac{e^{i\tilde{\varphi}2^{n-1}} (e^{-i\tilde{\varphi}2^{n-1}} - e^{i\tilde{\varphi}2^{n-1}})}{e^{i\tilde{\varphi}/2} (e^{-i\tilde{\varphi}/2} - e^{i\tilde{\varphi}/2})} \Rightarrow \left| \sum_{x=0}^{2^n-1} e^{i\tilde{\varphi}x} \right|^2 = \frac{\sin^2(\tilde{\varphi}2^{n-1})}{\sin^2(\tilde{\varphi}/2)}$$

Therefore: $\Pr(y) = |\alpha_y|^2 = \frac{1}{2^{2n}} \frac{\sin^2((\varphi - \frac{2\pi y}{2^n}) \cdot 2^{n-1})}{\sin^2(\varphi - \frac{2\pi y}{2^n}) \cdot \frac{1}{2}}$. $\frac{\sin mx}{\sin x} \rightarrow m$ when $x \rightarrow 0$.

This distribution is tightly peaked around those y for which $\frac{2\pi y}{2^n} \approx \varphi$.

If $\varphi = \frac{2\pi y}{2^n}$, $\Pr(y) = \frac{1}{2^{2n}} \cdot (2^{n-1} \cdot 2)^2 = 1$.

Claim: Let $\frac{2\pi k}{2^n} \leq \varphi \leq \frac{2\pi(k+1)}{2^n}$. Then the probability of outputting either k or $k+1$ is at least $\frac{8}{\pi^2}$.

Proof. The probability of success is $\Pr(k) + \Pr(k+1)$.

$$\begin{aligned}
\Pr(k) + \Pr(k+1) &= \frac{1}{2^{2n}} \left(\frac{\sin^2(2^{n-1}\varphi - \pi k)}{\sin^2\left(\frac{\varphi}{2} - \frac{\pi k}{2^n}\right)} + \frac{\sin^2(2^{n-1}\varphi - \pi(k+1))}{\sin^2\left(\frac{\varphi}{2} - \frac{\pi(k+1)}{2^n}\right)} \right) \\
\left(\min \text{ when } \varphi = \frac{2\pi\left(k + \frac{1}{2}\right)}{2^n} \right) &\geq \frac{2}{2^{2n}} \frac{\sin^2\left(\pi\left(k + \frac{1}{2}\right) - \pi k\right)}{\sin^2\left(\frac{\pi\left(k + \frac{1}{2}\right)}{2^n} - \frac{\pi k}{2^n}\right)} \\
&= \frac{1}{2^{2n-1}} \cdot \frac{1}{\sin^2\frac{\pi}{2^{n+1}}} \\
(\sin x \leq x) &\geq \frac{1}{2^{2n-1}} \cdot \frac{1}{\left(\frac{\pi}{2^{n+1}}\right)^2} = \frac{8}{\pi^2}.
\end{aligned}$$

□

Summary: Given $|\psi\rangle$ with $U|\psi\rangle = e^{i\varphi}|\psi\rangle$, we can produce an estimate of φ that differs from the true value by at most $\varepsilon\left(\frac{2\pi}{2^n}\right)$ with probability $\geq \frac{8}{\pi^2}$. This uses QFT^{-1} with **gate complexity** $O\left((\log \frac{1}{\varepsilon})^2\right)(n^2)$ and $O(1/\varepsilon)$ **controlled** $-U_s(2^n)$.