

CS 261: Data Structures

Binary Tree Traversals

Binary Search Trees

Binary Tree Traversals

- How to examine nodes in a tree?
- A list is a simpler structure:
 - can be traversed either forward or backward
- What order do we visit nodes in a tree?
- Most common tree traversal orders:
 - Pre-order
 - In-order
 - Post-order

Binary Tree Traversals

- All traversal algorithms have to:
 - Process node
 - Process left subtree
 - Process right subtree

Traversal order is determined by the order these operations are done

Binary Tree Traversals – Most Common

1. Node, left, right → Pre-order

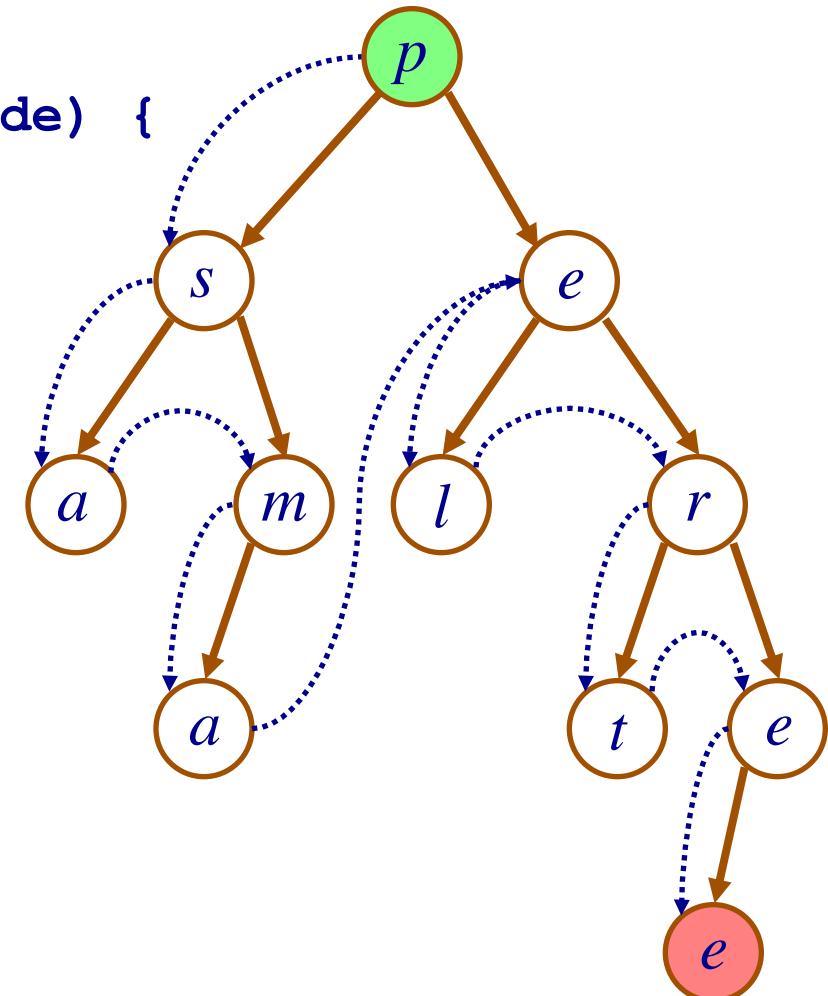
2. Left, node, right → In-order

3. Left, right, node → Post-order

Pre-Order Binary Tree Traversal

- Process order → Node, Left subtree, Right subtree

```
void preorder(struct Node *node) {  
    if (node != 0) {  
        process (node->val);  
        preorder (node->left);  
        preorder (node->right);  
    }  
}
```

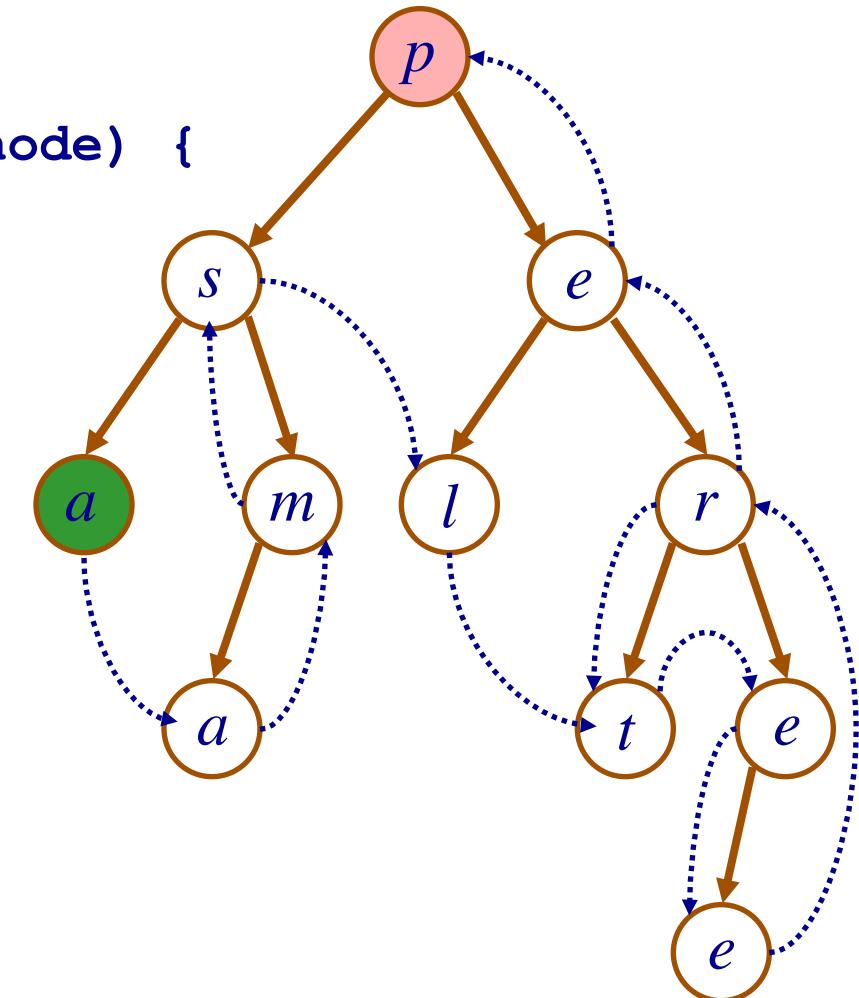


Example result: p s a m a e l r t e e

Post-Order Binary Tree Traversal

- Process order → Left subtree, Right subtree, Node

```
void postorder(struct Node *node) {  
    if (node != 0) {  
        postorder(node->left) ;  
        postorder(node->rght) ;  
        process (node->val) ;  
    }  
}
```

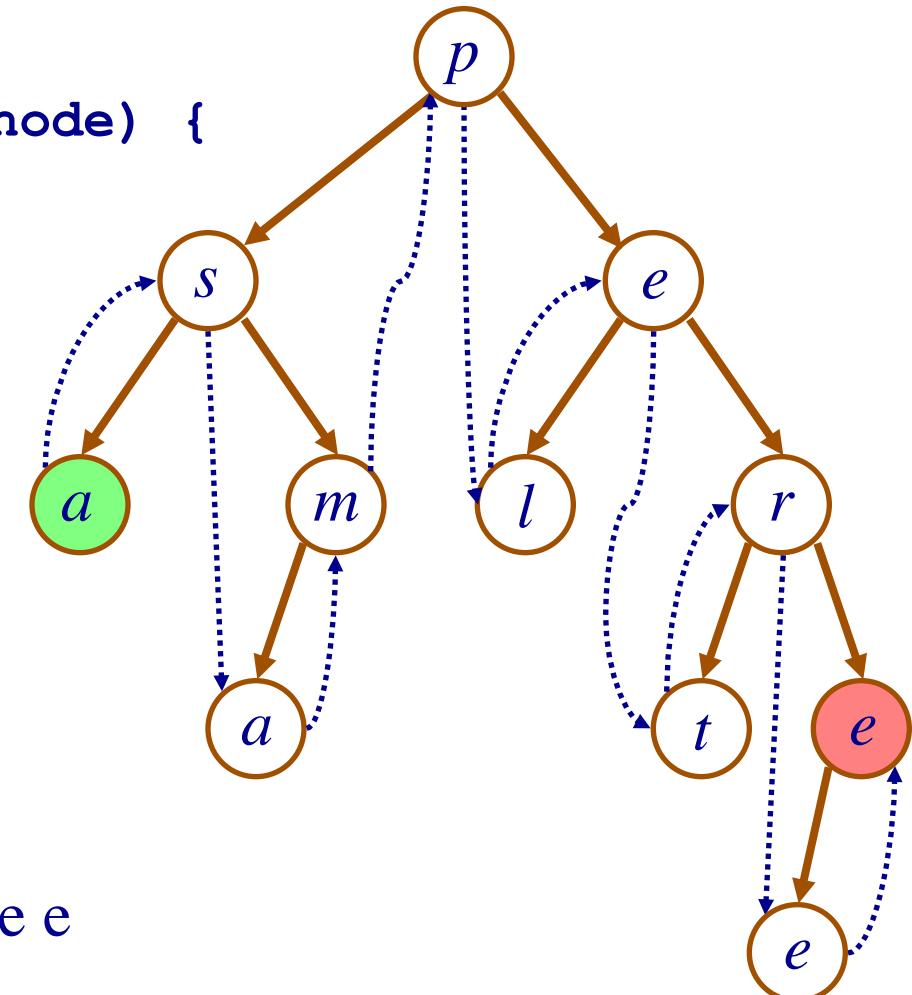


Example result: a a m s l t e e r e p

In-Order Binary Tree Traversal

- Process order → Left subtree, Node, Right subtree

```
void inorder(struct Node *node) {  
    if (node != 0) {  
        inorder(node->left) ;  
        process (node->val) ;  
        inorder(node->rght) ;  
    }  
}
```



Example result: a s a m p l e t r e e

Complexity of Binary Tree Traversals

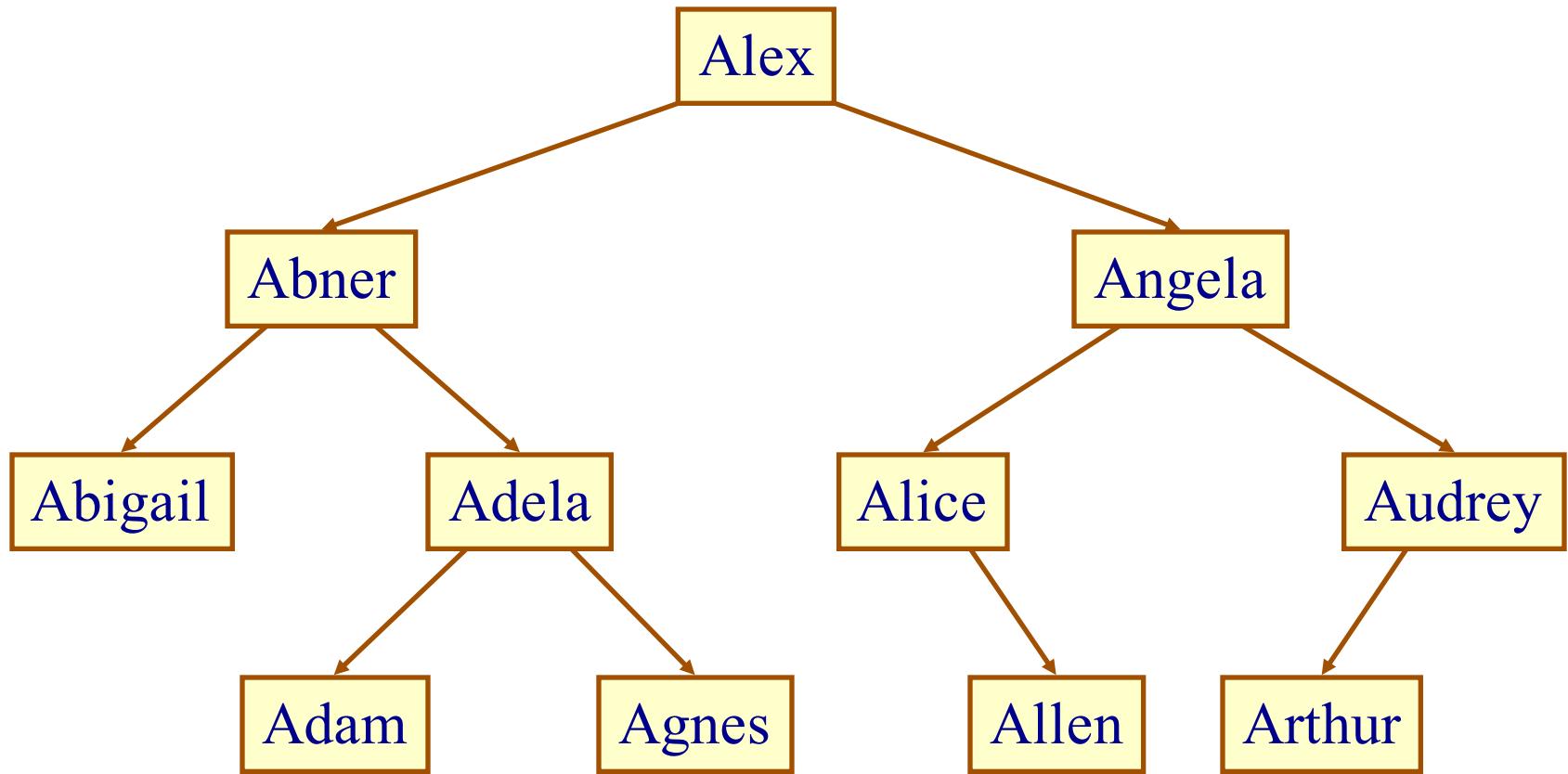
- Each traversal requires constant work at each node (not including recursive calls)
- Order not important
- Iterating over all n elements in a tree requires $O(n)$ time

Binary Search Trees

Binary Search Tree

- = Binary trees where every node value is:
 - **Greater than** all its **left** descendants
 - **Less than or equal to** all its **right** descendants
 - In-order traversal returns elements in sorted order
- If tree is reasonably full (well balanced), searching for an element is $O(\log n)$

Binary Search Tree: Example



A Node in BST

```
struct Node {  
    TYPE    val;  
    struct Node *left;     /* Left child */  
    struct Node *right;   /* Right child */  
};
```

Binary Search Tree (BST): Implementation

```
struct BST {  
    struct Node *root;  
    int size;  
};  
  
void initBST(struct BST *tree);  
int containsBST(struct BST *tree, TYPE val);  
void addBST(struct BST *tree, TYPE val);  
void removeBST(struct BST *tree, TYPE val);  
int sizeBST(struct BST *tree);
```

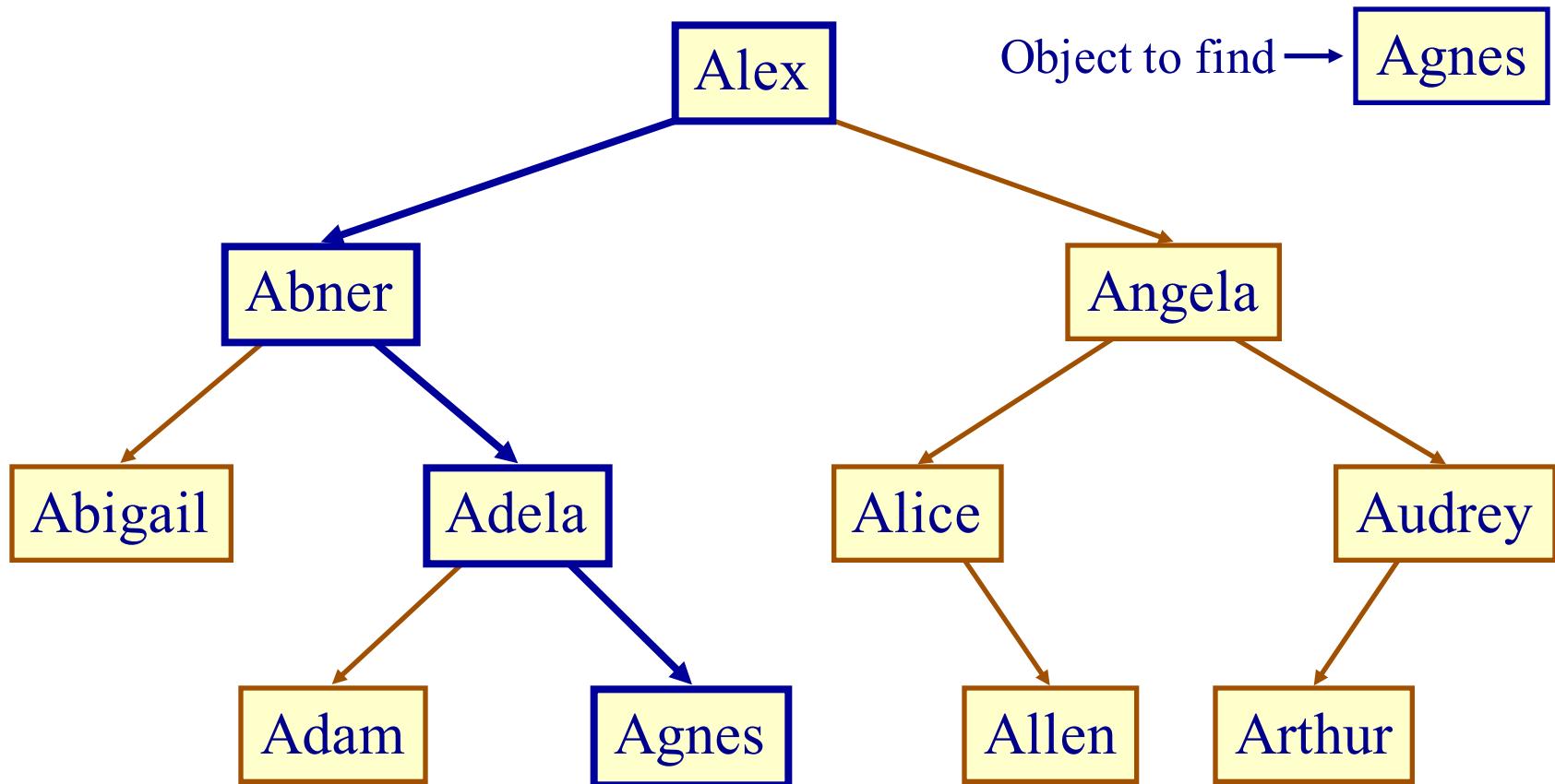
Sorted Bag Implementation: Contains

- Start at root
- At each node, compare value to node value:
 - Return true if match
 - If value is less than node value, go to left child (and repeat)
 - If value is greater than node value, go to right child (and repeat)
 - If node is null, return false

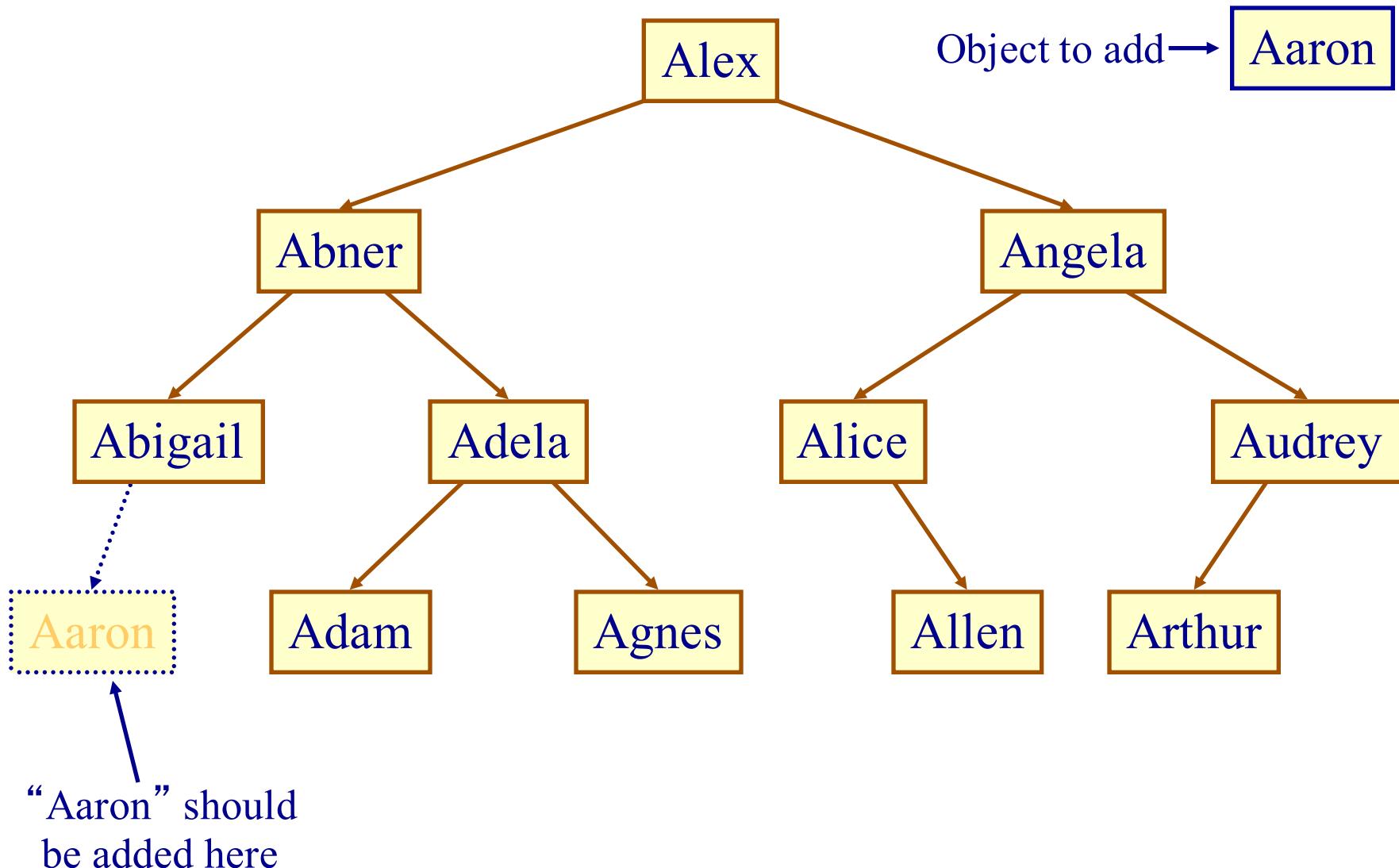
Bag Implementation: Contains

- Traverses a path from the root to the leaf
- Therefore, if tree is *reasonably full*,
execution time is $O(??)$

BST: Contains/Find Example



BST: Add Example



Add: Calls a Utility Routine

```
void addBST(struct BST *tree, TYPE val) {  
  
    assert(tree);  
  
    tree->root = _addNode(tree->root, val);  
  
    tree->size++;  
  
}
```

What is the complexity? $O(??)$

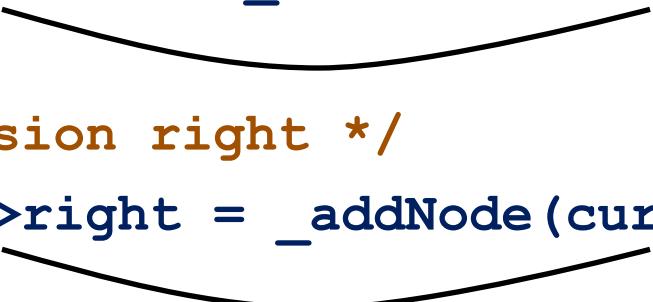
Useful Auxiliary Function -- Recursion

```
struct Node * _addNode(struct Node * current, TYPE val) {  
    if (current == 0)  
        /* always first check when to stop */  
  
    . . . /* make a new node and return the pointer to it */  
    }  
    else  
    {  
        /* call recursion left or right */  
    }  
    return current;  
}
```

return the same pointer

Useful Auxiliary Function -- Recursion

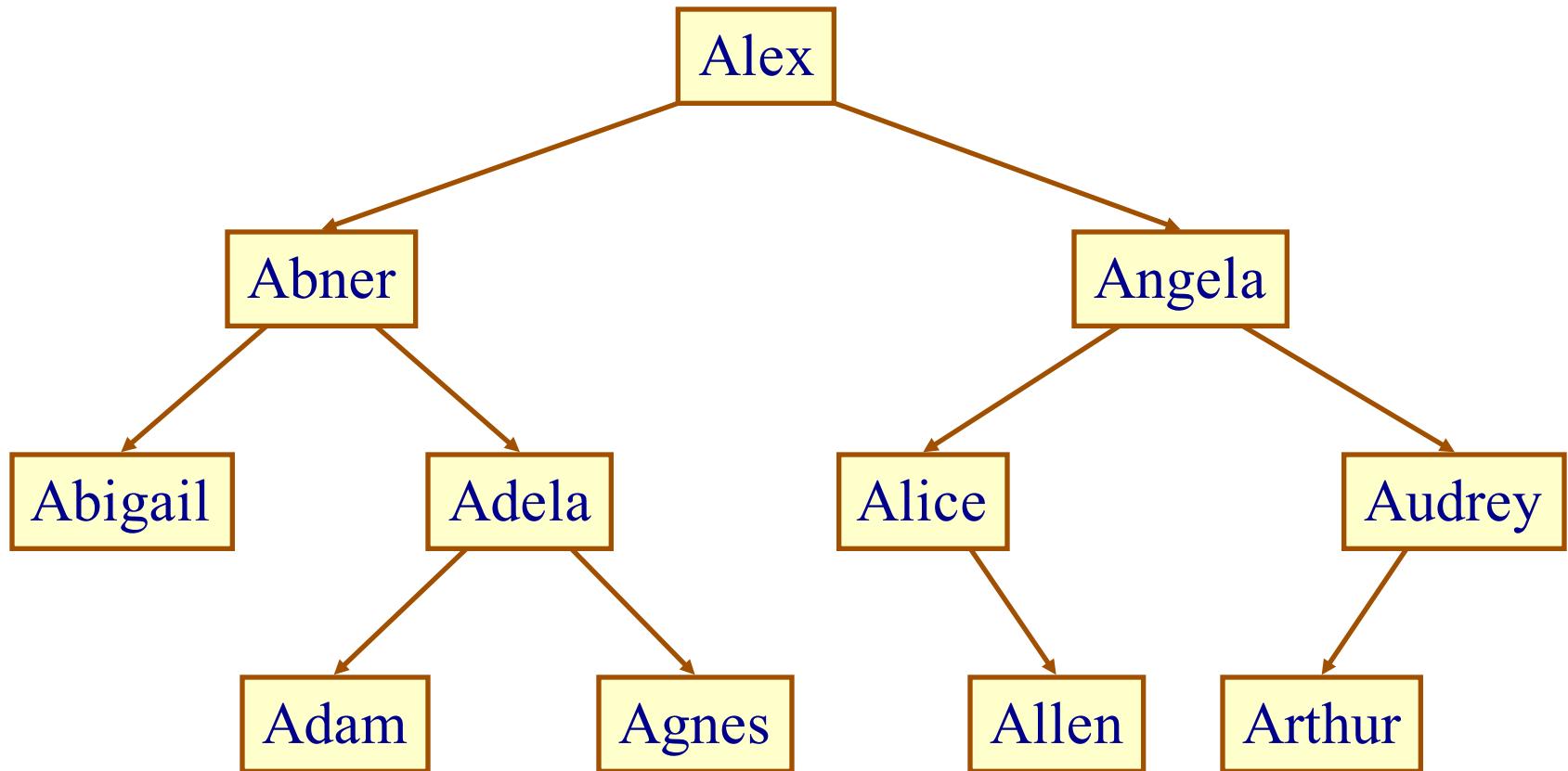
```
struct Node * _addNode(struct Node * current, TYPE val) {
    if (current == 0) {
        ... /*make a new node and return the pointer to it*/
    }else{
        if (val < current->val)
            /* recursion left */
            current->left = _addNode(current->left, val);
        else
            /* recursion right */
            current->right = _addNode(current->right, val);
    }
    return current;
}
```



Useful Auxiliary Function

```
struct Node * _addNode(struct Node * current, TYPE val) {
    if (current == 0) {
        struct Node * new =
            (struct Node *) malloc(sizeof(struct Node));
        assert(new != 0);
        new->value = val;
        new->left = new->right = 0;
        return new;
    } else
        if (val < current->val)
            current->left = _addNode(current->left, val);
        else
            current->right = _addNode(current->right, val);
    return current;
}
```

BST: Remove



How would you remove: Audrey? Angela?

Who Fills the Hole in BST?

- Answer:

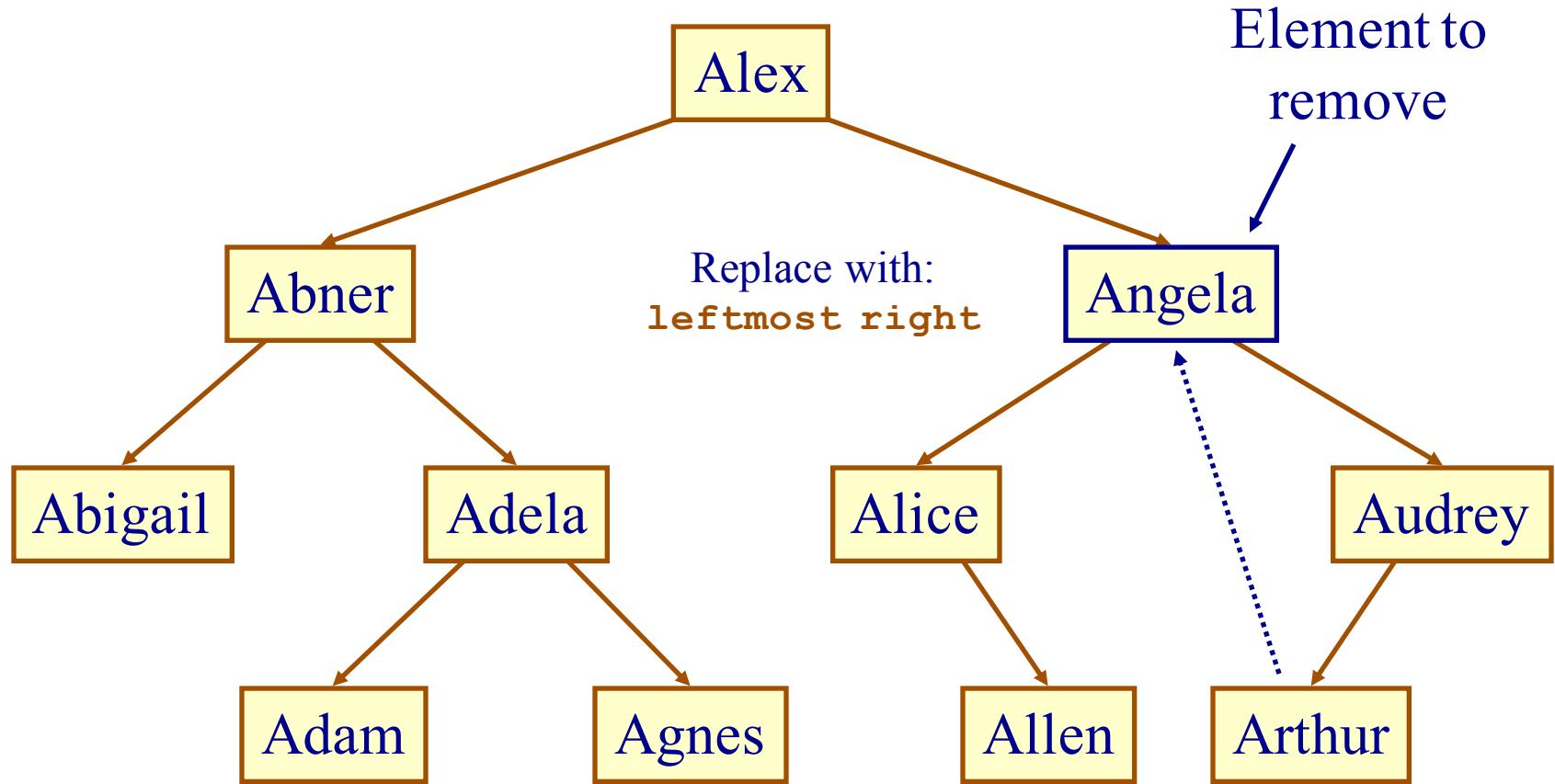
the leftmost descendant

of the right child

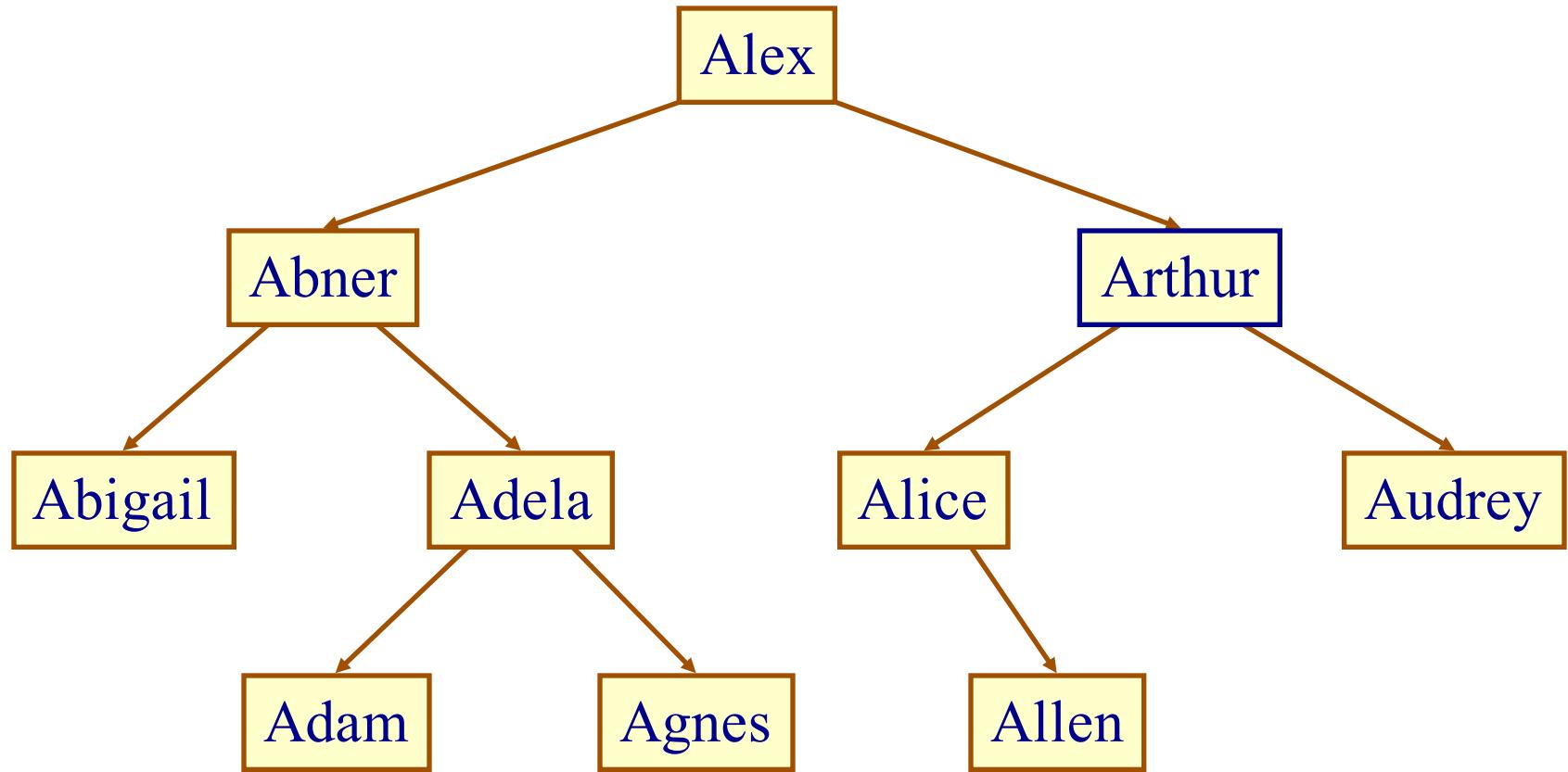
= smallest element in the right subtree

- Try this on a few values

BST: Remove Example



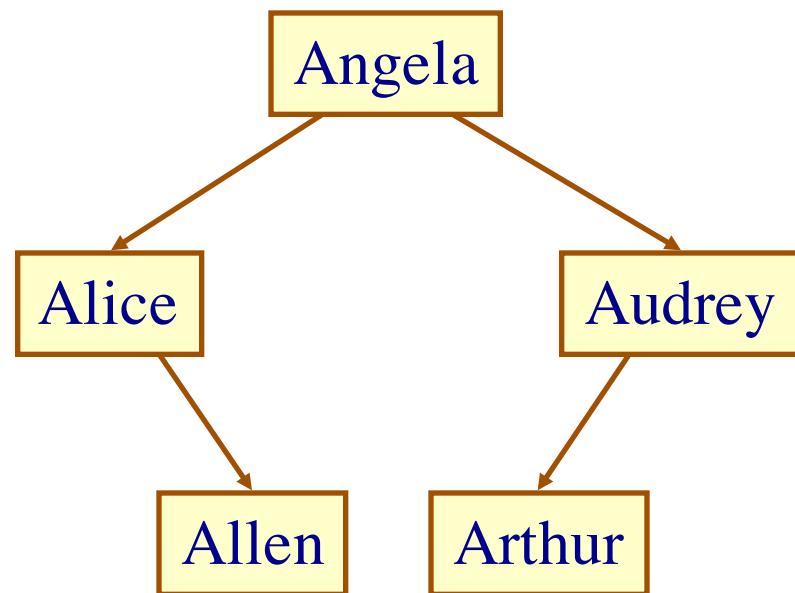
BST: Remove Example



After call to **remove**

Special Case

- What if you don't have the right child?
- Try removing "Audrey"
 - Simply, return the left child



Remove

```
void removeBST (struct BST *tree, TYPE e) {  
  
    if (containsBST(tree, e)) {  
  
        tree->root = _removeNode(tree->root, e);  
        tree->size--;  
  
    }  
  
}
```



Remove Node

```
struct Node *_removeNode(struct Node *current, TYPE e) {  
    if (current->val == e)  
        /* Found it, remove it */  
  
    else  
        /* not found, so recursive call */  
}
```

Remove Node

```
struct Node *_removeNode(struct Node *current, TYPE e) {
    struct Node *node;
    assert(current);
    if (current->val == e) /* Found it, remove it */
        if (current->right == NULL) {
            node = current->left;
            free(current);
            return node; /* returns the left child */
        }
    else
        /* replace current->val with the value of
         * the leftmost descendant of the right child */
        current->val = _leftMost(current->right);
        current->right = _removeLeftmost(current->right);
    else
        /* not found, so recursive call */
    ...
}
```

Remove Node

```
struct Node *_removeNode(struct Node *current, TYPE e) {  
    if (current->val == e) /* Found it, remove it */  
        ...  
    else /* not found, so recursive call */  
        if( e < current->val)  
            current->left = _removeNode(current->left, e);  
        else  
            current->right = _removeNode(current->right, e);  
  
    return current;  
}
```

Useful Routines: _leftMost

```
/*Returns value of leftmost child of current node*/  
  
TYPE _leftMost(struct Node *current) {  
  
    while(current->left != 0)  
  
        current = current->left;  
  
    return current->val;  
  
}
```

Useful Routines: _removeLeftmost

```
/* Removes the leftmost descendant of current */

struct Node *_removeLeftmost(struct Node *current) {

    struct Node *temp;

    if(current->left != 0) {

        current->left = _removeLeftmost(current->left);

        return current;
    }

    temp = current->right;
    free(current);
    return temp;
}
```

Complexity

- Running down a path from root to leaf
- What is the time complexity? $O(??)$