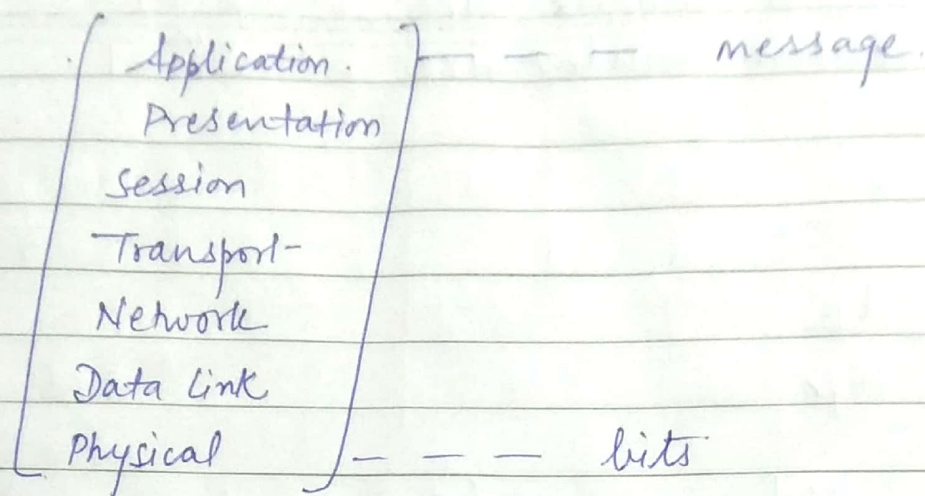


COMPUTER NETWORKS

OSI Model

— 7 layer model.



Noiseless channels

- (1) Simplest
- (2) Stop & wait.

Noisy channel

- (1) Stop & wait ARQ
- (2) Go back N
- (3) Selective Repeat ARQ.

Stop and wait ARQ :

uses mod-2 arithmetic as it has either 0 or 1 as data packets.

Go back N ARQ :

Sender window of size = $2^m - 1$
and receiver window size = 1

Selective Repeative ARQ

$$\text{Sender window size} = 2^m - 1$$

$$\text{Receiver window size} = 2^m - 1$$

basic formulas

$$\textcircled{1} \quad a = \frac{t_p}{t_{fr}}$$

$$\textcircled{2} \quad t_p = \frac{\text{Dist}}{sp} \quad (\text{time for propagation})$$

$$\textcircled{3} \quad t_{fr} = \frac{\text{msg}}{\text{Bandwidth}}$$

Utilizations

$$\Rightarrow \text{Stop \& wait} ; \quad U = \frac{1}{1+2a} ; \quad a \geq 1$$

$$\Rightarrow \text{Go back N ARQ} ; \quad U = \frac{W}{1+2a} ; \quad \begin{matrix} \leftarrow \\ W \geq 2a+1 \end{matrix}$$

$$\text{Throughput} = U \times \text{Bandwidth}$$

Ques: Consider a selective repeat sliding window protocol that uses a frame size of 1KB to send data on a 1.5 Mbps link with a one-way delay of 50 msec. To achieve a link utilization of 60%, the minimum no. of bits required to represent the sequence number field is _____?

Solⁿ:

$$\begin{aligned} M &= 1 \text{ KB} = 10^3 \times 8 \text{ Mb} \\ B &= 1.5 \text{ Mbps} \\ T_p &= 50 \text{ ms} \\ U &= 60\% \end{aligned}$$

For SR protocol $\Rightarrow U = \frac{W}{1+2a}$

$$a = \frac{t_p}{t_{fr}}$$

$$\frac{60}{100} = \frac{W}{1+2a}$$

$$t_{fr} = \frac{M}{B} = \frac{8 \times 10^3}{1.5 \times 10^6}$$

$$\frac{60}{100} = \frac{W}{1+2(9.43)}$$

$$t_{fr} = 5.3 \text{ ms}$$

$$\frac{60}{100} = \frac{W}{1+18.86}$$

$$a = \frac{5.3}{0.57} = 9.43$$

$$\frac{60}{100} = \frac{W}{19.86}$$

$$W = \frac{60}{100} \times \frac{19.86}{100}$$

$$W = \frac{1986 \times 6}{1000} = 11.9 \approx \underline{\underline{12}}$$

$$\therefore W = 2^{m-1}$$

$$\left[2^{m-1} = \frac{2^m}{2^1} \right]$$

$$\Rightarrow 12 = \frac{2^m}{2}$$

$$24 = 2^m$$

$$2^5 = 32$$

$$m = 5$$

Ques:- Suppose that a stop and wait protocol used bit rate on 64 Kbit/sec and 20 msec propagation delay. Assume that transmission time for the acknowledgement and processing time at nodes is negligible. Then min. frame size in bytes to achieve 50% utilization?

Solⁿ:-

$$t_p = 20 \text{ msec} = 20 \times 10^{-3} \text{ sec}$$

$$B = 64 \text{ Kbit/sec} = 64 \times 10^3 \text{ bit/sec}$$

$$U = 50\%$$

$$M = ?$$

for stop & wait protocol — $U = \frac{1}{1+2a}$

$$\frac{1}{2} \cdot \frac{50}{100} = \frac{1}{1+2\left(\frac{t_p}{t_{fr}}\right)}$$

$$\frac{1}{2} = \frac{1}{1 + 2 \left[\frac{20 \times 10^{-3}}{1} \right]}$$

$$a = \frac{t_p}{t_{fr}}$$

$$t_{fr} = \frac{M}{B}$$

$$\frac{1}{2} = \frac{1}{1 + 2 \left[\frac{20 \times 10^{-3} \times 64 \times 10^3}{M} \right]}$$

$$\frac{1}{2} = \frac{1}{1 + 2 \left[\frac{20 \times 64}{M} \right]}$$

$$\underline{\underline{M}} = 1 + 2 \left[\frac{20 \times 64}{M} \right] = 2$$

$$2 \left[\frac{20 \times 64}{M} \right] = 2 - 1$$

$$\frac{20 \times 64}{M} = \frac{1}{2} \Rightarrow M = 2560 \text{ bits}$$

$$\underline{\underline{M = 320 \text{ bytes}}}$$

Ques:- A link has transmission speed of 10^6 bit/sec.
It ~~reduces~~ uses data packets of size 1000 bytes.
Efficiency of stop & wait protocol in this set up
is exactly 25%. one way $t_p = \underline{\hspace{2cm}}$? in msec.

Ans:- $t_p = 12 \text{ msec}$

Ques:- Consider a network connecting two systems located
8000 km apart.

$$B = 500 \times 10^6 \text{ bits/sec}$$

$$\text{Speed} = 4 \times 10^6 \text{ m/sec.}$$

It is needed to design Go back N sliding window
protocol. Avg packet size is 10^7 bits.

Network is to be used to its full capacity.

Then min. size in bits of sequence number
fields is $\underline{\hspace{2cm}}$?

Ans:- 8