

OPERATING SYSTEMS

CE-C14

ASSIGNMENT-1

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COE-2
SEM-4

Ques 1 A system has 6 processes as follows:

PROCESS	ARRIVAL TIME (msec)	CPU BURST TIME (msec)
P ₁	0	8
P ₂	4	50
P ₃	8	70
P ₄	28	60
P ₅	30	20
P ₆	32	

calculate average waiting time for processes if round robin scheduling is used with time quantum of:

i) 80 msec

Ready queue: P₁ P₂ P₃ P₄ P₅ P₆ P₁

Gantt chart:

P1	P2	P3	P4	P5	P6	P1
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0 80 88 138 208 268 288 298

Waiting time:

- P₁ → (0-0) + (288-80) = 208 msec
- P₂ → (80-4) = 76 msec
- P₃ → (88-8) = 80 msec
- P₄ → (138-28) = 110 msec
- P₅ → (208-30) = 178 msec
- P₆ → (268-32) = 236 msec

Avg waiting time = $\frac{148 \text{ msec} (208 + 76 + 80 + 110 + 178 + 236)}{6}$

ii) 20 msec

Ready queue: P₁ P₂ P₃ P₄ P₅ P₆ P₃ P₁ P₄ P₅ P₃ P₁ P₄ P₅ P₁ P₄

Gantt chart:

P1	P2	P3	P1	P4	P5	P6	P3	P1	P4	P5	P3	P1	P4	P5	P1	P4
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0 20 28 48 68 88 108 128 148 168 188 208 218 238 258 278 288 298

Waiting time:

- P₁ → (0-0) + (148-20) + (148-68) + (218-168) + (278-238) = 198
- P₂ → (20-4) = 16
- P₃ → (28-8) + (128-48) + (208-148) = 160

$$P_6 \rightarrow (108 - 32) = 76$$

$P_6 \rightarrow (108 - 32) = 76$
 Avg waiting time $\Rightarrow 139.66 \text{ msec}$ $\left(\frac{76 + 188 + 200 + 16 + 160 + 1}{6} \right)$

iii) 10 msec

Waiting time : $P_1 \rightarrow (0-0) + (28-10) + (78-32) + (128-88) + (168-136) + (208-176) + (238-216) + (268-248) + (298-278) = 208$

$(288-278)=208$
 $p_2 \rightarrow \dots \dots \dots (110-4)=6$

$$P_2 \rightarrow (18-8) + (38-28) + (88-48) + (138-98) + (178-148) + (218-198) = 130$$

$$P_4 \rightarrow (18-8) + (38-28) + (88-48) + (138-98) + (178-148) + (218-198) + (248-228) + (278-258) = 190$$

$$P_5 \rightarrow (58-30) + (108-68) + (158-118) + (198-168) + (228-208) + (258-238) = 178$$

$$p_6 \rightarrow (68 - 32) + (118 - 78) = 76$$

$P_6 \rightarrow (68-32) + (118-78) = 76$
 Avg waiting time $\rightarrow 131.33 \text{ msec } \left(\frac{208+6+130+190+178+76}{6} \right)$

Ques 2 Use a suitable deadlock algorithm to check if there is a deadlock in the system where
 available = $[2 \ 1 \ 2 \ 3 \ 1]$,
 Request = $[1 \ 0 \ 1 \ 0 \ 0]$

available = $[2 \ 1 \ 2 \ 3 \ 1]$,

Available = 1 2 1 2 3 4 5 7
Allocation = $\begin{bmatrix} 2 & 3 & 1 & 3 & 3 \\ 2 & 2 & 2 & 1 & 1 \\ 1 & 2 & 2 & 2 & 2 \\ 2 & 3 & 2 & 2 & 2 \\ 1 & 2 & 1 & 1 & 1 \\ 1 & 2 & 3 & 1 & 1 \end{bmatrix}$

$$\text{Request} = \begin{bmatrix} 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 & 0 \end{bmatrix}$$

Work = $[2 \ 1 \ 2 \ 3 \ 1]$

jurist = [F F F F F]

Since $\text{Finish}[i]$
= True for all
processes, there
is no deadlock

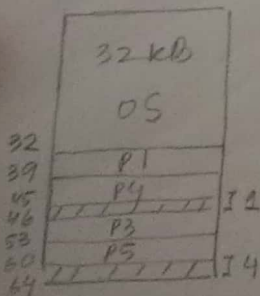
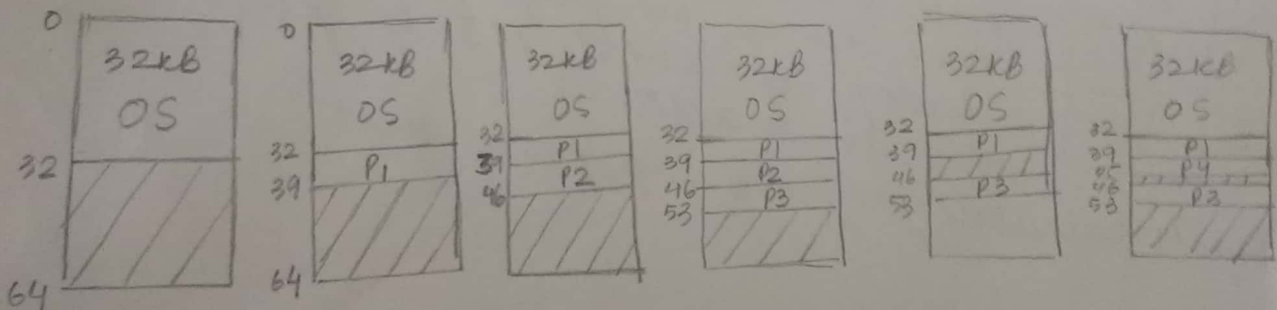
	Initially	Work	Finish
step ①	Select P1	2 1 2 3 1	F F F F F
step ②	Select P2	4 4 3 6 4	T F F F F
step ③	Select P3	5 6 5 7 5	T T F F F
step ④	Select P4	7 9 7 9 7	T T T F F
step ⑤	Select P5	8 11 8 11 9	T T T T F
step ⑥	Select P5	9 13 11 12 10	T T T T T

A system has 64KB RAM and uses contiguous memory allocation. The sizes of the OS and the user processes are as follows:

Process	Size (KB)
Operating system	32
P ₁	7
P ₂	7
P ₃	7
P ₄	6
P ₅	7
P ₆	5

The following sequence of events occur in the system, OS is loaded, then P₁, P₂, P₃ arrive, P₂ terminates, P₄, P₅, P₆ arrives. Determine if all the processes can be allocated space in the memory following:

i) FIRST FIT

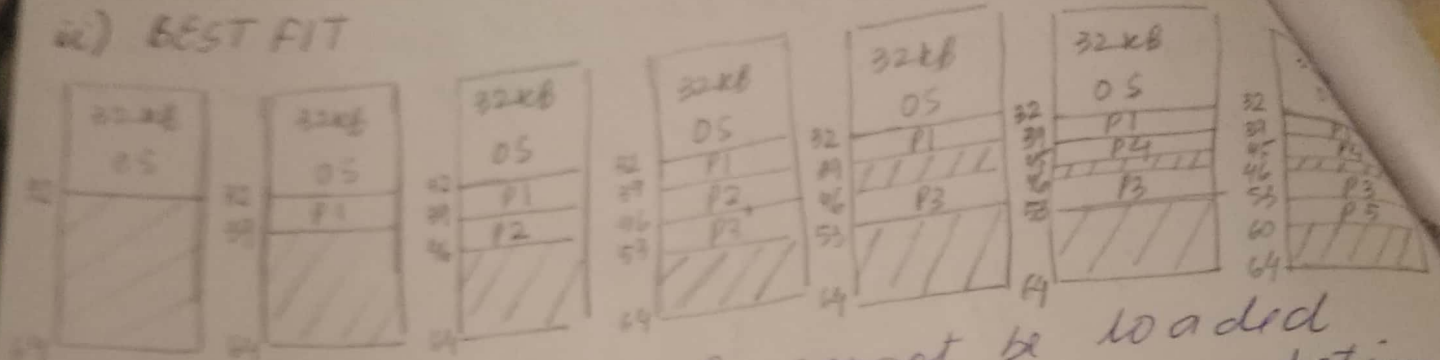


The total memory space still available = 1+4 = 5KB and the next incoming process P₆ has size = 5KB.

But since 5KB is available in different holes hence, we cannot load P₆ in memory even though the size of process is equal to the empty space available in the memory due to external fragmentation.

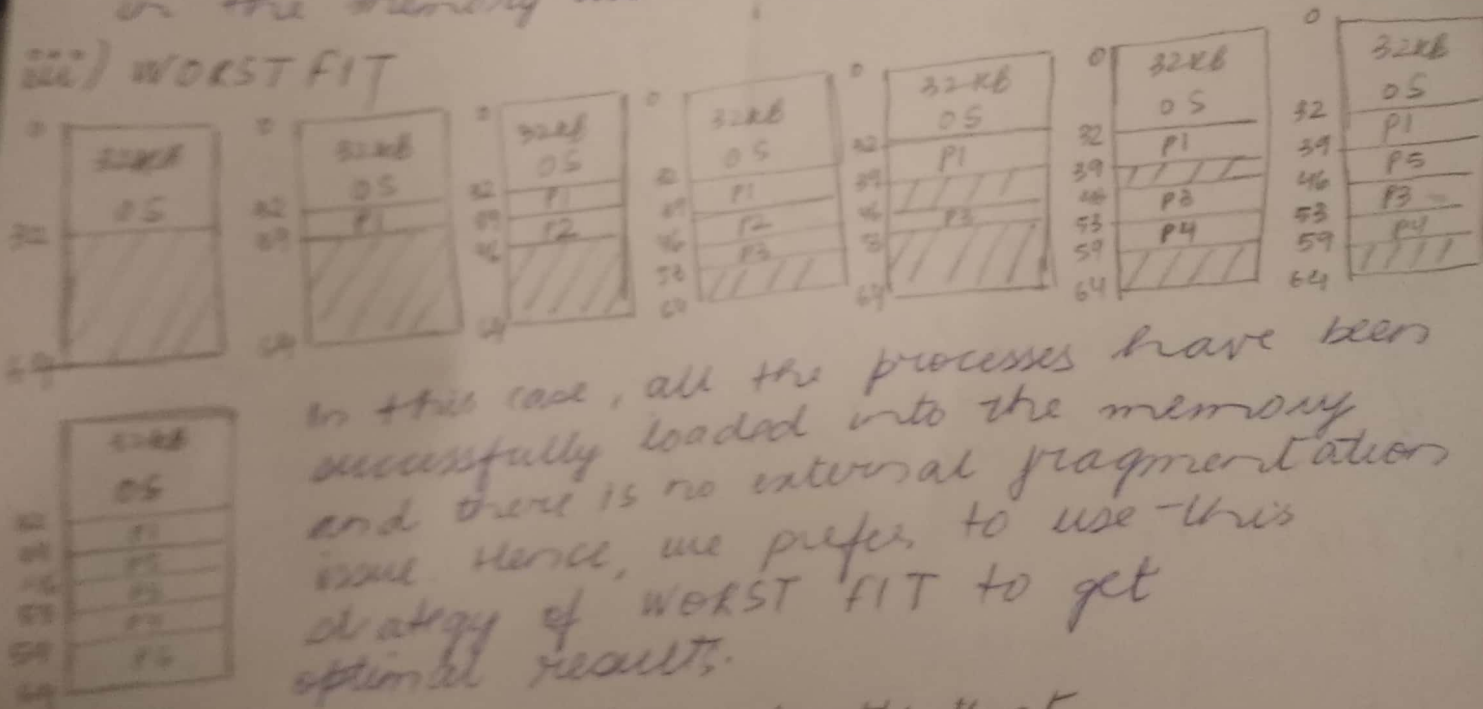
Ques 5 A hard disk has 1000 cylinders. The head is now at cylinder 615 and moves outwards. The disk access cylinders are (also):

ii) BEST FIT



→ In the above case, P_6 cannot be loaded in the memory due to external fragmentation.

iii) WORST FIT



In this case, all the processes have been successfully loaded into the memory and there is no external fragmentation issue. Hence, we prefer to use this strategy of WORST FIT to get optimal results.

Ques 4 Calculate the no. of page faults that

Ques 5 A hard disk has 1000 cylinders. The read/write head is now at cylinder 615 and moving outwards. The disk queue that requests to access cylinders 14, 918, 680, 183, 788 and 144. Calculate the number of cylinders that the read/write head will have to traverse to serve these requests if i) FCFS ii) SSTF iii) SCAN and iv) LOOK scheduling algorithm are used.

i) 0 14 144 183 615 680 788 918 999



$$\text{Number of cylinder traversed} = (615 - 14) + (918 - 14) + (918 - 680) + (680 - 183) + (788 - 183) + (788 - 144) = 3489$$

ii) 0 14 144 183 615 680 788 918 999



$$\text{Number of cylinder traversed} = (680 - 615) + (788 - 680) + (918 - 788) + (918 - 183) + (183 - 144) + (144 - 14) = 1207$$

iii) inward

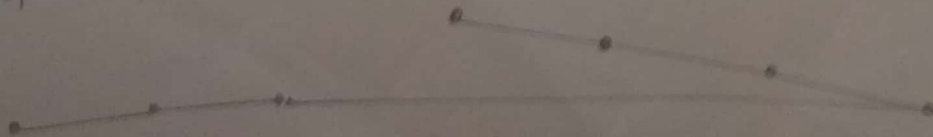
0 14 144 183 615 680 788 918 999



$$\text{Number of cylinder traversed} = (680 - 615) + (788 - 680) + (918 - 788) + (999 - 918) + (999 - 183) + (183 - 144) + (144 - 14) = 1369$$

iv) inward

0 14 144 183 615 680 788 918 999



$$\text{Number of cylinder traversed} = (680 - 615) + (788 - 680) + (918 - 788) + (918 - 183) + (183 - 144) + (144 - 14) = 1207$$