

Lecture 2 Notes

David Gieseke

October 14, 2020

1 Euclid's Knowledge

Let's first review the division algorithm.

Theorem. Let $a, b \in \mathbb{Z}$ and $b > 0$. Then, $\exists! q, r \in \mathbb{Z}$ such that $a = qb + r$ where $0 \leq r < b$.

This is the same as what we learned from elementary school. We know that q is the quotient and r is the remainder – and both are unique.

Proof. (existence) Let $S = \{n \in \mathbb{Z} \mid \exists x \in \mathbb{Z}, n = a - bx \geq 0\}$. Suppose that $S \neq \emptyset$ (the elements are potential q 's). Then, we can consider 3 different cases:

1. If $a \geq 0$, we take $x = 0$ and $a \in S$.
2. If $a < 0$, we take $x = a$. Then, $n = a - ab = a(1 - b)$. But since $b > 0$ and $b \in \mathbb{Z}$, $n = a(1 - b) \geq 0$. Therefore, $a \in S$.

And by the Well-Ordering Property, there is a least element r of S . If we denote its corresponding x to be q , then $r = a - bq \geq 0 \implies a = bq + r$.

Next, we show that $0 \leq r < b$. We showed $r \geq 0$ previously. Then assume that $r \geq b$, and thus $0 \leq r - b = a - b(q + 1) \in S$. However, r is assumed to be a least element of S , so $r \leq r - b$ and $b \leq 0$. This is a contradiction, so $r < b$. \square

Proof. (uniqueness) Assume that quotients and remainders are not unique. Then $a = qb + r = q'b + r'$, where $0 \leq r < b, 0 \leq r' < b$. If the remainder is not unique, then we can take $r > r'$ without loss of generality. Thus, we have $0 < r - r' \leq r < b$. And since $0 \neq q - q' \in \mathbb{Z}$, $|q - q'| \geq 1$. Then from the original equation, we have $r - r' = (q' - q)b \geq b$. This is a contradiction to the prior inequality. Thus, $r = r'$ must hold, and $q = q'$ follows as $b > 0$. \square

Definition. Let $d, m \in \mathbb{Z}$ where $d \neq 0$. Then d divides m if $\exists e \in \mathbb{Z}, m = ed$, notated by $d \mid m$.

Definition. Let $m, n \in \mathbb{Z}$. Then, $d \in \mathbb{Z}$ is a greatest common divisor, notated by $d = (m, n)$, of m and n if:

- (i) $d > 0$

(ii) $d \mid m, d \mid n$

(iii) if $e \in \mathbb{Z}$ and $e \mid m, e \mid n$, then $e \mid d$.

In other words, d is a divisor of m and n that divides any other common divisors.

Definition. $f \in \mathbb{Z}$ is a \mathbb{Z} linear combination of $m, n \in \mathbb{Z}$ if $\exists x, y \in \mathbb{Z}$ such that $f = xm + yn$.

Theorem. Let $m, n \in \mathbb{Z}$, and at least one of which is nonzero. Then, $d = (m, n)$ exists, is unique, and is a \mathbb{Z} linear combination of m and n .

Proof. Let's define $S = \{am + bn > 0 \mid a, b \in \mathbb{Z}\}$. Now, we first show that $d \mid m$. S is nonempty, since we can always take $a = \text{sgn}(m), b = 0$ and $am + bn \in S$ must hold. Thus, it must have a least element d . Then by the division algorithm, $\exists q, r \in \mathbb{Z}$ such that $m = qd + r$ and $0 \leq r < d$.

$$r = m - qd = m - (am + bn)q = (1 - aq)m - bqn \geq 0$$

Therefore, r is also a \mathbb{Z} linear combination of m and n . However $r < d$ and d is the least element of S , so $r = 0$ is the only possibility. Therefore, $m = qd \implies d \mid m$. Similarly, $d \mid n$ can be proven with a similar method.

Now, suppose that e is another common divisor of m and n . Thus $m = xe$ and $n = ye$, so $d = (ax + by)e$ and it follows that $e \mid d$. Therefore $d = (m, n)$ by definition, and it is a \mathbb{Z} linear combination of m and n as it is in S . \square

Proof. (uniqueness) Let d, d' be two GCD's of m and n . Then, $d \mid d'$ and $d' \mid d$ must both hold. This is true only when $d = \pm d'$, and since $d, d' > 0$ they must be equal. \square