

TICKLAB

Documentary



MIPS RUNNER

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1 Idea

1.1 Initial Idea

Assembly is a low-level programming language that affects registers directly. The execution of an assembly program will affect the operating status of the computer. This leads to the need for a virtual machine to simulate assembly program behavior. The popular virtual machine is the Java virtual machine.

MIPS-Setting low-level machine language on C ++ that can translate and show users running each command.

1.2 Main idea

Write a compile program to check for code syntax errors.

Identify and save labels to perform address jump commands.

Simulate components with 32 bit memory cells, store values.

After compile, put the register data register on the screen and each command (10 sentences) for the user to manipulate to run each command line.

2 Research and development process

2.1 Knowledge for research:

2.1.1 Document

Computer organization and design 5th edition.

Mr Dat's report.

Web documents.

2.1.2 Instructor

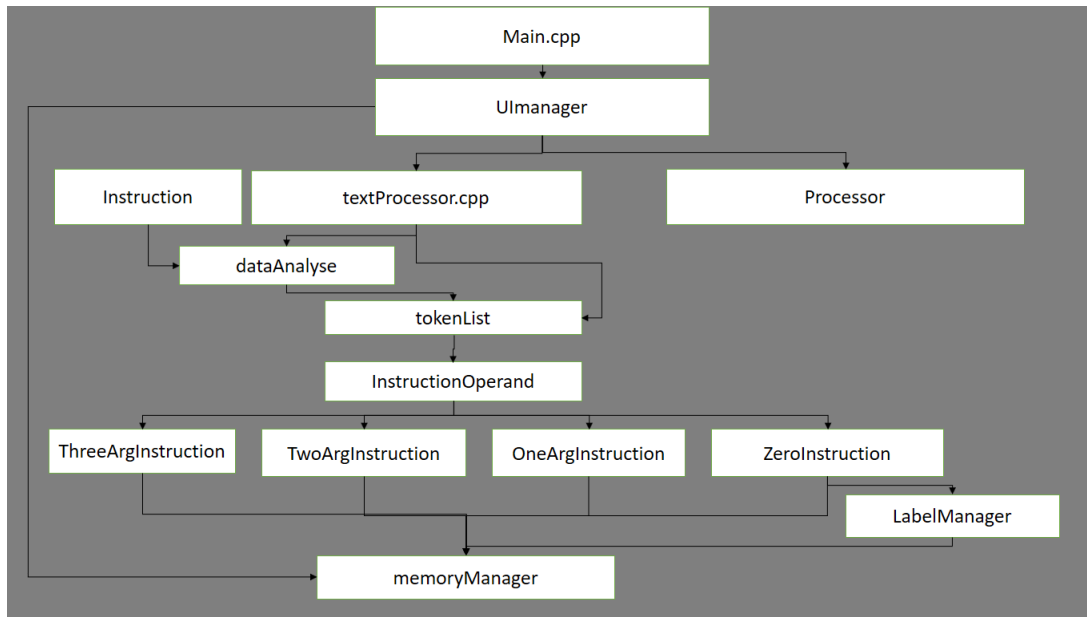
Huynh Hoang Kha.

2.1.3 The process of inquiry

Some noticeable things:

- Register.
- Instruction.
- The way computer executes variables.
- Initialize variable - load value into memory -Store in register - executes.
- How the computer processes loops by saving the address of the statement with a label.

2.1.4 Figure



2.2 Processing

2.2.1 First outline

Read the code file from a text file.

Split into lines in the textProcessor file.

Split each line into tokens in the tokenList file.

Identify the tokens (Instruction, label) in the InstructionOperand file separately the data envelope (.data) and the commandtext area: will be processed separately to store variable labels.

Instruction is divided into 4 types:

- ThreeArgInstruction(three argument instruction).
- TwoArgInstruction(two argument instruction).
- OneArgInstruction(one argument instruction).
- ZeroArgInstruction(zero argument instruction).

Each type of instruction has corresponding instructions which execute those instructions.

Register:

- Registers are stored in the memoryManager class.
- Simulate registers into two array functions of 35 elements and 31 elements (floating-point register) to store values.
- Particularly register value \$pc, is the order of lines commands are processed in the textProcessor class.

Difficulty:

- To recognize token, when handling each line, tokens that do not recognize it is a immediate value or register values.
- The idea is only suitable when the labels are declared before the jump or jal command to the labels.
- Before processing the line must have a compile.

2.2.2 Final idea

Right from the textProcessor file, identify the tokens as Immediate value, register, label, variable label.

Run the program before once to save the labels.

Add compile feature.

3 Product

3.1 Principle of operation

Read a text file (source.asm) from textProcessor class. Then separate these source code line by line, save these lines to src (char **) to save each char * corresponding to 1 statement.

Parse sourceCode into Instruction:

- Partition to declare data (.data) and write code (.text).
- Identify the label.
- Identify Instruction:

Use the register pc to save the order of the current statement in the char** array.

Use tokenList to split elements in the line into tokens, saved in token* array .

Counts the number of elements in tokenList * to identify which Instruction type (Three, Two, One or Zero instruction) is.

Compile: If the number of arguments in tokenList is more than 3, compile error.

Processing Instruction: e.g ThreeArgInstruction:

- Simulate 3 variables rd, rs, rt in class InstructionOperand.
- This Instruction Operand class is responsible for assigning identification 7 symbols:
 - + Variable label.
 - + Label.
 - + Register.
 - + Floating-point register.
 - + Integer.

- + Float.
- + Address.
- Identify the Instruction in ThreeArgInstruction:
 - + Use the variables rs, rt, rd to access memory if it is a register (include memoryManager into InstructionOperand, use the pointer to access array to store Register data).
 - + Compile: If in that Instruction, the identifiers of rs, rt cannot be exceeded (e.g. addi requires rt as immediate value = integer), an error will be displayed.
 - + If it is immediate value, simulate a new variable different from the register memory cell.
- Identify the Instruction in OneArgInstruction: the labels have been saved the original address, when jumping, just change pc (saving the current location).

Export values to the screen using UImanager class.

3.2 Demonstration

```
.data
prompt1: .asciiz "Input first real number: \n"
prompt2: .asciiz "Input second real number: \n"
prompt3: .asciiz "Sum: "
prompt4: .asciiz "Sub: "
.text
main:
    # Input the data
    la $a0, prompt1
    li $v0, 4
    syscall
    li $v0, 6
    syscall

    mfc1    $t0, $f0    # t0 contain the bits of first number

    la $a0, prompt2
    li $v0, 4
    syscall
    li $v0, 6
    syscall

    mfc1    $t1, $f0    # t1 contains the bits of second number
```

Assembly

```
-----SOURCE CODE-----
1:      .data
2:      prompt1:
3:      .asciiz "Input first real number: \n"
4:      prompt2:
5:      .asciiz "Input second real number: \n"
6:      prompt3:
7:      .asciiz "Sum: "
8:      prompt4:
9:      .asciiz "Sub: "
10:     .text
11:     main:
12:     la $a0,prompt1
13:     li $v0,4
14:     syscall
15:     li $v0,6
16:     syscall
17:     mfc1 $t0,$f0
18:     la $a0,prompt2
19:     li $v0,4
20:     syscall
21:     li $v0,6
22:     syscall
23:     mfc1 $t1,$f0
24:     F_SIGN_CHECKING:
25:     bgez $t0,F_GREATER_ZERO
26:     addi $t2,$zero,1
27:     sll $t2,$t2,31
28:     j END_F_SIGN_CHECKING
29:     F_GREATER_ZERO:
30:     add $t2,$zero,$zero
31:     END_F_SIGN_CHECKING:
32:     S_SIGN_CHECKING:
33:     bgez $t1,S_GREATER_ZERO
34:     addi $t3,$zero,1
```

Put Assembly on display

```
120:      or $t1,$t1,$t9
121:      mtc1 $t0,$f12
122:      la $a0,prompt3
123:      li $v0,4
124:      syscall
125:      li $v0,2
126:      syscall
127:      mtc1 $t1,$f12
128:      la $a0,prompt4
129:      li $v0,4
130:      syscall
131:      li $v0,$s0
132:      syscall
133:      j END
134:      STANDARDIZED:
135:      blt $v0,0x01000000,LESS_THAN
136:      GREATER_THAN:
137:      subi $v1,$v0,0x01000000
138:      blt $v1,0x01000000,END_STANDARDIZED
139:      srl $v0,$v0,1
140:      addi $s3,$s3,1
141:      j GREATER_THAN
142:      LESS_THAN:
143:      subi $v1,$v0,0x01000000
144:      bge $v1,$zero,END_STANDARDIZED
145:      sll $v0,$v0,1
146:      subi $s3,$s3,1
147:      j LESS_THAN
148:      END_STANDARDIZED:
149:      jr $ra
150:      END:
```

```
-----COMPILE ERROR-----
line 131:      "$s0" have to be an integer
-----
```

Compile error


```
-----RUNNING-----
11:      main:
12:      la $a0,prompt1
-> 13:      li $v0,4
14:      syscall
15:      li $v0,6
16:      syscall
17:      mfc1 $t0,$f0
-----
n
-----RUNNING-----
11:      main:
12:      la $a0,prompt1
13:      li $v0,4
-> 14:      syscall
15:      li $v0,6
16:      syscall
17:      mfc1 $t0,$f0
18:      la $a0,prompt2
-----
n
Input first real number:
-----RUNNING-----
11:      main:
12:      la $a0,prompt1
13:      li $v0,4
14:      syscall
-> 15:      li $v0,6
16:      syscall
17:      mfc1 $t0,$f0
18:      la $a0,prompt2
19:      li $v0,4
-----
```

press 'n' to execute next line.

press 'u' to show line are executing (not executing).

-----REGISTER-----				
\$0	\$zero	0	\$f0	0
\$1	\$at	0	\$f1	0
\$2	\$v0	4	\$f2	0
\$3	\$v1	0	\$f3	0
\$4	\$a0	13570112	\$f4	0
\$5	\$a1	0	\$f5	0
\$6	\$a2	0	\$f6	0
\$7	\$a3	0	\$f7	0
\$8	\$t0	0	\$f8	0
\$9	\$t1	0	\$f9	0
\$10	\$t2	0	\$f10	0
\$11	\$t3	0	\$f11	0
\$12	\$t4	0	\$f12	0
\$13	\$t5	0	\$f13	0
\$14	\$t6	0	\$f14	0
\$15	\$t7	0	\$f15	0
\$16	\$s0	0	\$f16	0
\$17	\$s1	0	\$f17	0
\$18	\$s2	0	\$f18	0
\$19	\$s3	0	\$f19	0
\$20	\$s4	0	\$f20	0
\$21	\$s5	0	\$f21	0
\$22	\$s6	0	\$f22	0
\$23	\$s7	0	\$f23	0
\$24	\$t8	0	\$f24	0
\$25	\$t9	0	\$f25	0
\$26	\$k0	0	\$f26	0
\$27	\$k1	0	\$f27	0
\$28	\$gp	0	\$f28	0
\$29	\$sp	11534223	\$f29	0
\$30	\$fp	0	\$f30	0
\$31	\$ra	\$ra		
\$32	\$hi	\$hi		
\$33	\$lo	\$lo		

press 'r' to show register memory.

[illegible]

press 'v' to show variable memory.

4 Conclusion

4.1 Summary

Completed most of comments of assembly.

Comprehended how to write an assembly.

Understand how a computer operated, organised and designed.

4.2 Evaluating working processing:

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Applying Git for operating project.

To aware the function of pare coding.

Frequent Interaction for planning idea,evaluating progress.

Comprehended the worthy value of using class,enhanced ability of. applying object oriented programming.

4.3 Lesson after project:

Drawing an sufficient picture clearly before starting to code.

Practicing enhancing ability of algorithm for dealing with problems fast, adaptive for fixing after.

References

- [1] Computer Organization and Design, Fifth Edition: The Hardware/Software Interface (The Morgan Kaufmann Series in Computer Architecture and Design).
- [2] <http://www.mrc.uidaho.edu/mrc/people/jff/digital/MIPSir.html>
- [3] <http://www.cs.uwm.edu/classes/cs315/Bacon/Lecture/HTML/ch05s04.html>
- [4] <https://people.cs.pitt.edu/~childers/CS0447/lectures/SlidesLab92Up.pdf>

This is the end of the report.