

Q1:

Plaintext: VuNguyens4010423@student.

01010110 01110101 01001110 01100111 01110101 01111001 01100101 01101110
01110011 00110100 00110000 00110001 00110000 00110100 00110010 00110011
01000000 01110011 01110100 01110101 01100100 01100101 01101110 01110100
00101110

1. θ mapping

C[x] Operation

01010110	01110101	01001110	01100111	01110101
01111001	01100101	01101110	01110011	00110100
00110000	00110001	00110000	00110100	00110010
00110011	01000000	01110011	01110100	01110101
01100100	01100101	01101110	01110100	00101110

$$C[0] = A[0][0] \oplus A[0][1] \oplus A[0][2] \oplus A[0][3] \oplus A[0][4]$$

01010110 \oplus 01111001 = 00101111
00101111 \oplus 00110000 = 00011111
00011111 \oplus 00100111 = 00111000
00111000 \oplus 01100100 = 01011100

$$C[0] = 01011100$$

C[0]	C[1]	C[2]	C[3]	C[4]
01011100	00010100	00011101	00100000	00110111

D[x] = operation

$$\begin{aligned} C[x] &= A[x,0] \oplus A[x,1] \oplus A[x,2] \oplus A[x,3] \oplus A[x,4] \\ D[x] &= C[x-1] \oplus \text{rot}(C[x+1], 1) \\ A[x,y] &= A[x,y] \oplus D[x] \end{aligned}$$

D[0] = C[4] \oplus rot(C[1], 1)
D[0] = 00110111 \oplus rot(00010100, 1)
D[0] = 00110111 \oplus 00101000
D[0] = 00011111

D[0]	D[1]	D[2]	D[3]	D[4]
00011111	01100110	01010100	01110011	10011000

A[x,y] operation

$$A[x,y] = A[x,y] \oplus D[x]$$

01010110	01110101	01001110	01100111	01110101
01111001	01100101	01101110	01110011	00110100
00110000	00110001	00110000	00110100	00110010
00110011	01000000	01110011	01110100	01110101
01100100	01100101	01101110	01110100	00101110

D[0]	D[1]	D[2]	D[3]	D[4]
00011111	01100110	01010100	01110011	10011000

A[x,y]	0	1	2	3	4
0	01001001	00010011	00011010	00010100	11101101
1	01100110	00000011	00111010	00000000	10101100
2	00101111	01000111	01100100	00010111	10111010
3	00111000	00100110	00110111	00000110	11110010
4	01111011	00000011	00111010	00000111	10110110

2. q mapping

	$x = 3$	$x = 4$	$x = 0$	$x = 1$	$x = 2$
$y=2$	25	39	3	10	43
$y=1$	55	20	36	44	6
$y=0$	28	27	0	1	62
$y=4$	56	14	18	2	61
$y=3$	21	8	41	45	15

$A[x,y]$	0	1	2	3	4
0	01001001	00010011	00011010	00010100	11101101
1	01100110	00000011	00111010	00000000	10101100
2	00101111	01000111	01100100	00010111	10111010
3	00111000	00100110	00110111	00000110	11110010
4	01111011	00000011	00111010	00000111	10110110

$A[0,2] = 00101111$

$q[0,2] = 3$

$3 \bmod 8 = 3$

$\text{rotl}(00101111, 3) = 01111001$

$A[0,1] = 01100110$

$q[0,1] = 36$

$36 \bmod 8 = 4$

$\text{rotl}(01100110, 4) = 01100110$

01001001	00100110	10000110	01000001	01101111
01100110	00110000	10001110	00000000	11001010
01111001	00011101	00100011	00101110	01011101
01110000	11000100	10011011	11000000	11110010
11101101	00001100	01000111	00000111	10101101

3. π Mapping

$\varrho[x,y]$	0	1	2	3	4
0	01001001	00100110	10000110	01000001	01101111
1	01100110	00110000	10001110	00000000	11001010
2	01111001	00011101	00100011	00101110	01011101
3	01110000	11000100	10011011	11000000	11110010
4	11101101	00001100	01000111	00000111	10101101

$\varrho[x,y]$	0	1	2	3	4
0	$\varrho[0,0]$	$\varrho[1,0]$	$\varrho[2,0]$	$\varrho[3,0]$	$\varrho[4,0]$
1	$\varrho[0,1]$	$\varrho[1,1]$	$\varrho[2,1]$	$\varrho[3,1]$	$\varrho[4,1]$
2	$\varrho[0,2]$	$\varrho[1,2]$	$\varrho[2,2]$	$\varrho[3,2]$	$\varrho[4,2]$
3	$\varrho[0,3]$	$\varrho[1,3]$	$\varrho[2,3]$	$\varrho[3,3]$	$\varrho[4,3]$
4	$\varrho[0,4]$	$\varrho[1,4]$	$\varrho[2,4]$	$\varrho[3,4]$	$\varrho[4,4]$

$$\pi(x, y) = \varrho[y][(2x + 3y) \bmod 5]$$

$$\pi(0, 0) = \varrho[0][(2*0 + 3*0) \bmod 5]$$

$$\pi(0, 0) = \varrho(0,0)$$

$$\pi(1, 0) = \varrho[0][(2*1 + 3*0) \bmod 5]$$

$$\pi(1, 0) = \varrho(0,2)$$

$\pi[x,y]$	0	1	2	3	4
0	$q[0,0]$	$q[0,2]$	$q[0,4]$	$q[0,1]$	$q[0,3]$
1	$q[1,3]$	$q[1,0]$	$q[1,2]$	$q[1,4]$	$q[1,1]$
2	$q[2,1]$	$q[2,3]$	$q[2,0]$	$q[2,2]$	$q[2,4]$
3	$q[3,4]$	$q[3,1]$	$q[3,3]$	$q[3,0]$	$q[3,2]$
4	$q[4,2]$	$q[4,4]$	$q[4,1]$	$q[4,3]$	$q[4,0]$

$\pi[x,y]$	0	1	2	3	4
0	01001001	00110000	00100011	11000000	10101101
1	01000001	11001010	01111001	11000100	01000111
2	00100110	10001110	00101110	11110010	11101101
3	01101111	01100110	00011101	10011011	00000111
4	10000110	00000000	01011101	01110000	00001100

4. χ mapping

$$A[x,y] = B[x,y] \oplus ((\bar{B}[x+1,y]) \wedge B[x+2,y]) \quad , \quad x,y = 0,1,2,3,4$$

$\chi(0,0) = A[0][0] \oplus (\neg A[1][0] \wedge A[2][0])$
 $01001001 \oplus (\neg 00110000 \wedge 10000110)$
 $01001001 \oplus (11001111 \wedge 10000110)$
 $01001001 \oplus 10000110$
 $\chi(0,0) = 01001010$

$\chi(1,0) = A[1][0] \oplus (\neg A[2][0] \wedge A[3][0])$
 $00110000 \oplus (\neg 10000110 \wedge 11000000)$
 $00110000 \oplus (01111001 \wedge 11000000)$
 $00110000 \oplus 01000000$
 $\chi(1,0) = 11110000$

$\chi[x,y]$	0	1	2	3	4
0	01001010	11110000	00001110	10000000	10011101
1	01110000	01001110	01111010	11000100	11001101
2	00000110	01011110	00100011	11110000	01100101
3	01110110	11100100	00011001	11110011	00000111
4	11011011	00100000	01010001	11110010	00001100

5. ι mapping

$\iota[x,y]$	0	1	2	3	4
0	01001011	11110000	00001110	10000000	10011101
1	01110000	01001110	01111010	11000100	11001101
2	00000110	01011110	00100011	11110000	01100101
3	01110110	11100100	00011001	11110011	00000111
4	11011011	00100000	01010001	11110010	00001100

Q2:

K = 4010423/K7MDENG+bPxRfiCYEXAMPLEKEY =

343031303432332f4b374d44454e472b62507852666943594558414d504c454b4559000000
00

m =

AWS4-HMAC-SHA256

20250415M123600Z

20250415/us-east-1/iam/aws4_request

F536975d06c0309214f805bb90ccff089219ecd68b2577efef23edd43b7e1a59

415753342d484d41432d5348413235360d0a32303235303431354d3132333630305a0d0a3
2303235303431352f75732d656173742d312f69616d2f617773345f726571756573740d0a66
3533363937356430366330333039323134663830356262393063636666303839323139656
364363862323537376566656632336564643433623765316135390d0a

Q3:

Keys and timestamp:

KC = 7c87b7e19d52cd8663223cad6b9233ab

KS = 39ffe3021bcad0da66bc4efc46ef0360

nC = f04981feda3674db9223bb95eda3da2f

timestamp = 11/05/2025, 21:56:00

Ticket:

c524f78ebb017a214e31bf4b53033bd410b4823c7fd81005966d8f6467f02c35

Authenticator:

03705b3885e4d1a7737a13f147bd8c26940a8f0f4c027410d7a215899d5b27d29f030c238afb
88cfe9d33055c4c400230b323be14d3550543c6d53a480c9c8b4

Authenticator Decrypted

f04981feda3674db9223bb95eda3da2f11/05/2025, 21:56:00

Q5:

1)

e = s4010423 = 3D31B7

Ephemeral key =

B3D72C1F8EAD6B4AB9A39D531CE1EF2175A7623D0924BAF95358B028369D4ACF
5ED9D2C0A88D7E94E6CE41BE349B4D67E2C108D1AEAA3935371B7A3CF4C2A23C
1195D184B1B8FA32BFA2A75888EECFBD2D11E91A0C6ED79A1E6B8B730BCDA3FD
43C9A93BA2BB6974B81FC438C19D3AD4F6BB7DF31D4C8BA3F9913EF49C867B51

2)

Ephemeral key + e =

B3D72C1F8EAD6B4AB9A39D531CE1EF2175A7623D0924BAF95358B028369D4ACF5ED
9D2C0A88D7E94E6CE41BE349B4D67E2C108D1AEAA3935371B7A3CF4C2A23C1195D1
84B1B8FA32BFA2A75888EECFBD2D11E91A0C6ED79A1E6B8B730BCDA3FD43C9A93B
A2BB6974B81FC438C19D3AD4F6BB7DF31D4C8BA3F9913EF49C867B513D31B7

After hashed

Cfa64e36a57718a7ffa84c9227ec60f506286b94aaf455ace9fcbea0bf87029c

$a^b \bmod m =$

96aa0ab0f5eaf060e61f925bc8e713fa166b0ec9eff9c66d697079a68dca59b04ac16bfe760de
5451cf5bc7ba6312d63e1c8b018bc8329000ed45919989ae18a454f58418a3814b456c2aa68
172482bb7ed21e8e4fd41bd479f791e099109a9a1aa9eb4fb1f4068739dd756f740fe0862adef
fb6dd0fdff0706d7c209d0ab21cd55857ad1b3e67874b79179e2d92a75b6d7d2ec7979c2b229
de37273f70aacf3d6c45f80eabeda6b9d35387e6ef097d6b6b5ba273cf6f880bcc917a4b51e0a
275a4b69af2d5d6016abef5acdf12dbf77f957983bb76325c2fbb36a2df482121a1387345dfb68
28aa9615a4f287f6580f7815c494a070d417367104dd2d8e2dd6

3)Pre_master_secret= s4010423@student.rmit.edu.au =

5b6295b95611541897ef2591f488f94b083a892d9e75cf533e23636533e2c334b93ab469e8e3
c7b7f6586a9d27777959

Converted to decimal

1406545375955935357313680527750466447880708307956708036332991519758525758
0833924864938604378913021678854215105476953

Free and fast online Modular Exponentiation (ModPow) calculator. Just type in the base number, exponent and modulo, and click Calculate. This [Modular Exponentiation](#) calculator can handle big numbers, with any number of digits, as long as they are positive integers.

For a more comprehensive mathematical tool, see the [Big Number Calculator](#).

Calculate $a^b \bmod m$

Number (a)

7504664478807083079567080363
3299151975852575808339248649
3860437891302167885421510547
6053

☐ Use hexadecimal numbers

Exponent (b)

65537

Modulo (m)

2715381915028252110041696926
5020354495704323552394680394
9158074029099492628585443073

Calculate

Result

16646912904973800653869108633669548870197822014253601762359435685048034931664305141368479477051460
98002565524228277092645328159004228531710964392583181335682497369325131148760820909953551567535576
217983496868855032956088750264495519281226222768209028830772545492799504327031584416226701078976

1664691290497380065386910863366954887019782201425360176235943568504803493
1664305141368479477051460980025655242282770926453281590042285317109643925
831813356824973693251311487608209099535515675355762179834968688550329560
8875026449551928122622227682090288307725454927995043270315844162267010789
7615429458178938752032740157623977582403387583505970280944831050736762977
07672911999321856582477782016621519243110551503507681835277263499099965805
2157235848810352698994256337205238404006009756716818033388905958601016209
6770873038412640576153817043603601788745813670376883388054140605740492740
39025368428949645712489564678936

4) In the SSL handshake protocol, the authentication is achieved through using a digital certificate. When a user connects, it verifies the certificate and confirms the server's identity

5) Forward security ensures that if a private key of a server is breached and compromised, past communication will still be safe. This is possible through the usage of ephemeral key exchange