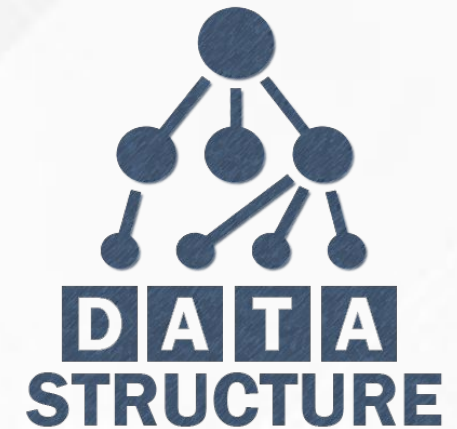


## Unit-2

# Linear Data Structure (Queue)



# Queue

- Imagine a line of people waiting at a ticket counter. People are **served in the order they come**, that is, people who **come first are served first** whereas people who come later are served after that.
- Let the action of someone joining a queue be called **enqueue** and someone being serve and getting out of the queue be called **dequeue**.
- A type of structure, similar to the example of the line of people, can be represented as a data structure. Such a data structure is known as a queue.

# Types of Queue

## ❑ SIMPLE QUEUE :

- ❑ A simple queue is a type of queue where **insertion is at the end** of the queue and **removal is at the front**.

## ❑ CIRCULAR QUEUE :

- ❑ A circular queue is a type of queue where the **last element is connected to the first**. An element is added to the end of the queue and removed from the front.

## ❑ PRIORITY QUEUE :

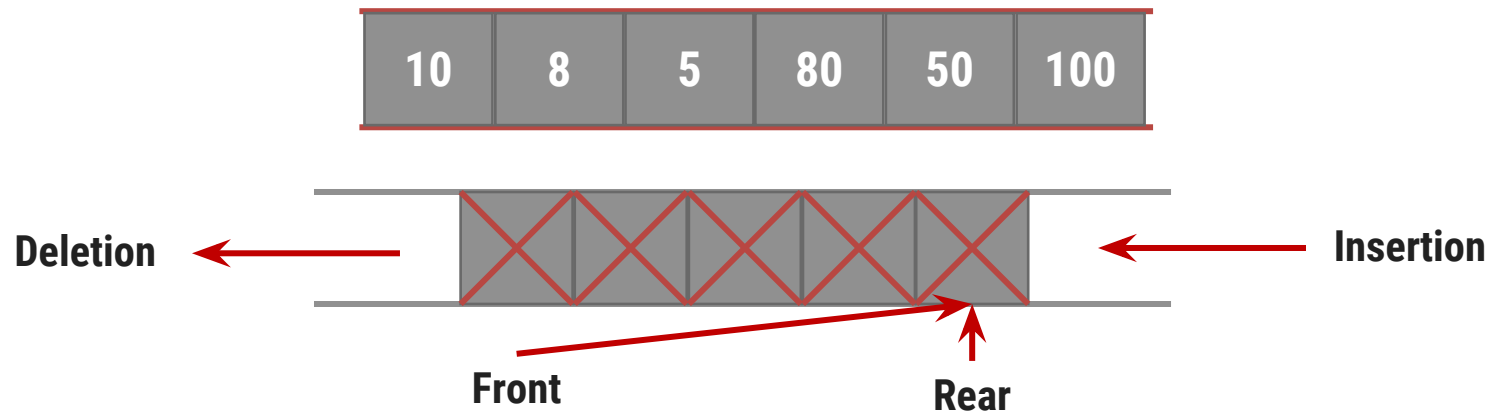
- ❑ A priority queue is a type of queue where the **elements are arranged in some priority order**. An element with the highest priority is removed first and insertion occurs according to the priority.

## ❑ DOUBLE ENDED QUEUE :

- ❑ A double ended queue is a type of queue where **insertion and deletion can happen at both ends of the queue**.

# Queue

- A linear list which permits **deletion** to be performed **at one** end of the list and **insertion at the other end** is called **queue**.
- The information in such a list is processed **FIFO (first in first out) or FCFS (first come first served)** manner.
- **Front** is the end of queue from that **deletion** is to be performed.
- **Rear** is the end of queue at which new element is to be **inserted**.
- Insertion operation is called **Enqueue** & deletion operation is called **Dequeue**.



# Applications of Queue

- ❑ Queue of people at any service point such as ticketing etc.
- ❑ Queue of air planes waiting for landing instructions.
- ❑ **Queue of processes** in OS.
- ❑ Queue is also used by Operating systems for **Job Scheduling**.
- ❑ When a **resource is shared** among multiple consumers. E.g., in case of printers the first one to be entered is the first to be processed.
- ❑ When **data is transferred asynchronously** (data not necessarily received at same rate as sent) between two processes. Examples include IO Buffers, pipes, file IO, etc.
- ❑ Queue is used in **BFS (Breadth First Search)** algorithm. It helps in traversing a tree or graph.
- ❑ Queue is used in networking to **handle congestion**.

# Simple Queue

- ❑ **Linear** Data Structure
- ❑ Perform Operation on **two ends**.
- ❑ Can be **implemented** using **Array** or **Linked List**.
- ❑ The simple queue is a **normal queue**
- ❑ **insertion** takes place at the **REAR END** of the queue - **enqueue**
- ❑ **deletion** takes place at the **FRONT END** of the queue – **dequeue**
- ❑ follows **FIFO – First In First Out** or **FCFS – First Come First In** principle.
- ❑ It is also called as **“Buffer”**.

# Simple Queue

## ❑ enqueue() - QINSERT








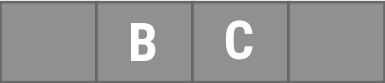
1. Check if the queue is full.
2. If the queue is full, then print "**Queue overflow**".
3. Else increment **REAR** by 1.
4. Assign `QUEUE [ REAR ] = ELEMENT`.

## ❑ dequeue() - QDELETE

1. Check if the queue is empty.
2. If the queue is empty, the print "**Queue underflow**".
3. Else Copy the element at the front of the queue to some temporary variable, `TEMP = QUEUE[ FRONT ]`.
4. Increment **FRONT** by 1.
5. Print temp and delete it.

# Example of Simple Queue Insert / Delete

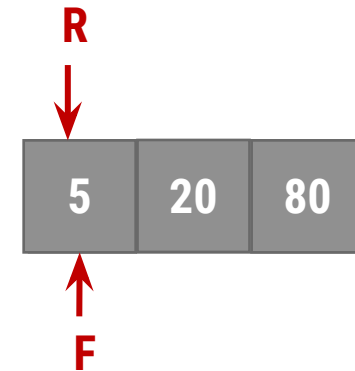
Perform following operations on queue with size 4 & draw queue after each operation  
Insert 'A' | Insert 'B' | Insert 'C' | Delete 'A' | Delete 'B' | Insert 'D' | Insert 'E'

<p><b>Empty Queue</b></p> <p>0 0</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Insert 'C'</b></p> <p>R=3 F=1</p> <p>A B C</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Insert 'D'</b></p> <p>R=4 F=3</p> <p>C D</p> <p>↑ ↑</p> <p>F R</p> 
<p><b>Insert 'A'</b></p> <p>R=1 F=1</p> <p>A</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Delete 'A'</b></p> <p>R=3 F=2</p> <p>A B C</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Insert 'E'</b></p> <p>R=4 F=3</p> <p>C D</p> <p>↑ ↑</p> <p>F R</p> 
<p><b>Insert 'B'</b></p> <p>R=2 F=1</p> <p>A B</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Delete 'B'</b></p> <p>R=3 F=3</p> <p>B C</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>(R=4) &gt;= (N=4) (Size of Queue)</b></p> <p><b>Queue Overflow</b></p> <p>Queue Overflow, but space is there with Queue, this leads to the memory wastage</p>



# Procedure: QINSERT(Q, F, R, N,Y) – Simple Queue [Enqueue]

- This procedure inserts **Y** at rear end of Queue.
- **Queue** is represented by a vector **Q** containing **N** elements.
- **F** is pointer to the front element of a queue.
- **R** is pointer to the rear element of a queue.



## 1. [Check for Queue Overflow]

```
If R >= N
Then write ('Queue Overflow')
Return
```

## 2. [Increment REAR pointer]

```
R = R + 1
```

## 3. [Insert element]

```
Q[R] = Y
```

## 4. [Is front pointer properly set?]

```
IF F=0
Then F = 1
Return
```

**N=3, R=0, F=0**

**F = 0**

**R = 0**

Enqueue (Q, F, R, N=3, **Y=5**)

Enqueue (Q, F, R, N=3, **Y=20**)

Enqueue (Q, F, R, N=3, **Y=80**)

Enqueue (Q, F, R, N=3, **Y=3**)

**Queue Overflow**

# Function: QDELETE(Q, F, R) – Simple Queue [Dequeue]

- This function **deletes & returns** an element **from front end** of the Queue.
- **Queue** is represented by a vector **Q** containing **N** elements.
- **F** is pointer to the **front** element of a queue.
- **R** is pointer to the **rear** element of a queue.

## 1. [Check for Queue Underflow]

```
If F = 0  
Then write ('Queue Underflow')  
Return(0)
```

## 2. [Delete element]

```
Y ← Q[F]
```

## 3. [Is Queue Empty?]

```
If F = R  
Then F ← R ← 0  
Else F ← F + 1
```

## 4. [Return Element]

```
Return (Y)
```

Case No 1:

F=0, R=0



Queue Underflow

Case No 2:

F=3, R=3

F R  
↓ ↓

F=0, R=0



Case No 3:

F=1, R=3

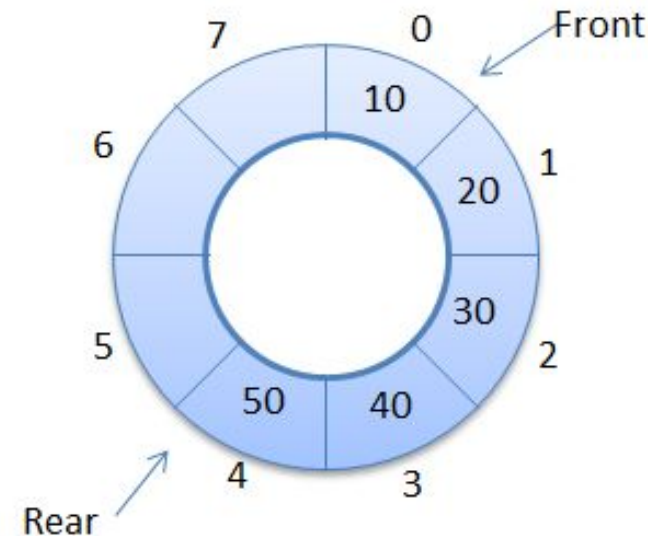
F R  
↓ ↓

F=2, R=3











# Circular Queue

- A more suitable method of representing simple queue which prevents an excessive use of memory is to **arrange the elements**  $Q[1], Q[2], \dots, Q[n]$  **in a circular fashion** with  $Q[1]$  following  $Q[n]$ , this is called **circular queue**.
- In circular queue the last node is connected back to the first node to make a circle.
- Circular queue is a linear data structure. It follows **FIFO** principle.
- It is also called as **"Ring buffer"**.



# Example of CQueue Insert / Delete

Perform following operations on Circular queue with size 4 & draw queue after each operation  
Insert 'A' | Insert 'B' | Insert 'C' | Delete 'A' | Delete 'B' | Insert 'D' | Insert 'E'

<p><b>Empty Queue</b></p> <p>0 0</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Insert 'C'</b></p> <p>R=3 F=1</p> <p>A B C</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Insert 'D'</b></p> <p>R=4 F=3</p> <p>C D</p> <p>↑ ↑</p> <p>F R</p> 
<p><b>Insert 'A'</b></p> <p>R=1 F=1</p> <p>A</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Delete 'A'</b></p> <p>R=3 F=2</p> <p>A B C</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Insert 'E'</b></p> <p>R=1 F=3</p> <p>E C D</p> <p>↑ ↑</p> <p>F R</p> 
<p><b>Insert 'B'</b></p> <p>R=2 F=1</p> <p>A B</p> <p>↑ ↑</p> <p>F R</p> 	<p><b>Delete 'B'</b></p> <p>R=3 F=3</p> <p>B C</p> <p>↑ ↑</p> <p>F R</p> 	

# Procedure: CQINSERT (F, R, Q, N, Y)

- This procedure inserts **Y** at rear end of the Circular Queue.
- **Queue** is represented by a vector **Q** containing **N** elements.
- **F** is pointer to the front element of a queue.
- **R** is pointer to the rear element of a queue.

## 1. [Reset Rear Pointer]

```
If      R = N  
Then    R ← 1  
Else    R ← R + 1
```

## 2. [Overflow]

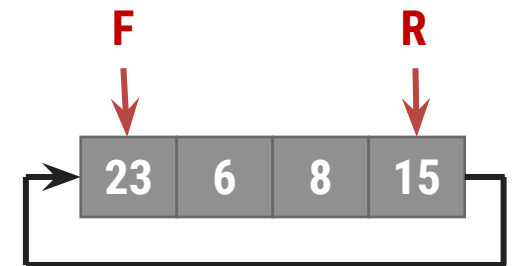
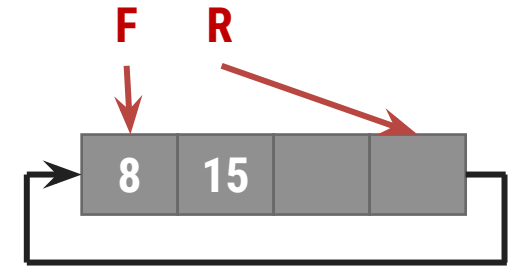
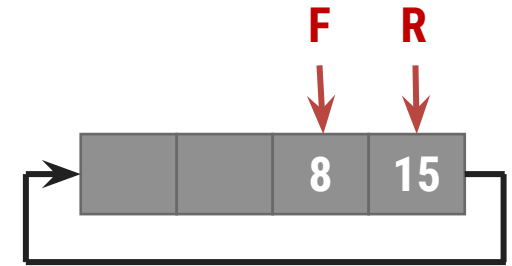
```
If      F=R  
Then    Write('Overflow')  
Return
```

## 3. [Insert element]

```
Q[R] ← Y
```

## 4. [Is front pointer properly set?]

```
IF F=0  
Then  F ← 1  
Return
```



# Function: CQDELETE (F, R, Q, N)

- This function **deletes & returns** an element **from front end** of the Circular Queue.
- **Queue** is represented by a vector **Q** containing **N** elements.
- **F** is pointer to the **front** element of a queue.
- **R** is pointer to the **rear** element of a queue.

## 1. [Underflow?]

```
If      F = 0  
Then Write('Underflow')  
      Return(0)
```

## 2. [Delete Element]

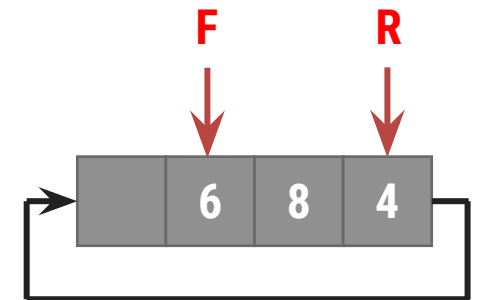
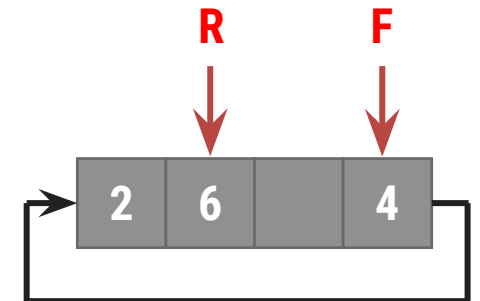
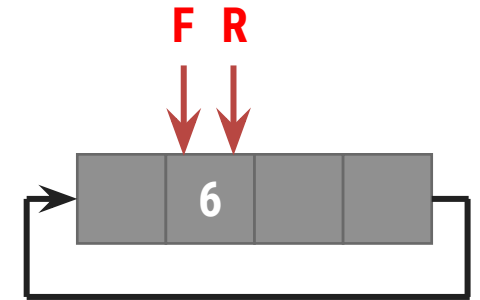
```
Y  $\leftarrow$  Q[F]
```

## 3. [Queue Empty?]

```
If      F = R  
Then F  $\leftarrow$  R  $\leftarrow$  0  
      Return(Y)
```

## 4. Increment Front Pointer]

```
IF F = N  
Then F  $\leftarrow$  1  
Else F  $\leftarrow$  F + 1  
Return(Y)
```

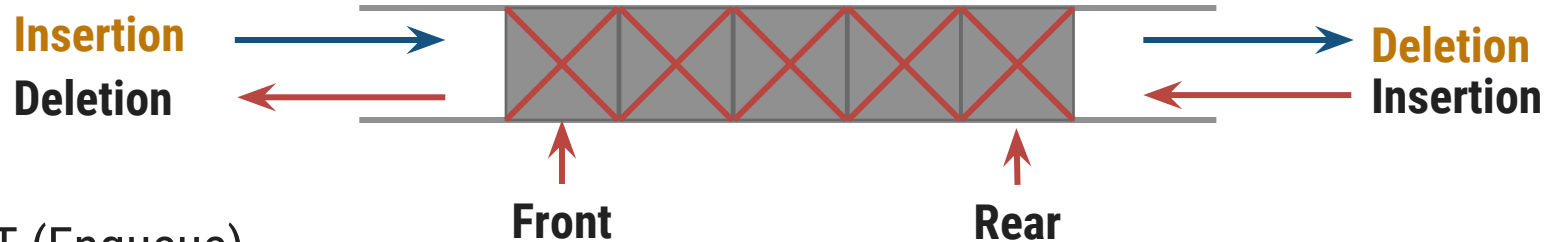


# DQueue

- A **DQueue (double ended queue)** is a linear list in which insertion and deletion are performed **from the either end of the structure**.
- There are two variations of Dqueue
  - **Input restricted dqueue** – allows insertion at only one end
  - **Output restricted dqueue** – allows deletion from only one end

## □ Dqueue Algorithms

- DQINSERT\_REAR is same as QINSERT (Enqueue)
- DQDELETE\_FRONT is same as QDELETE (Dequeue)
- DQINSERT\_FRONT
- DQDELETE\_REAR



# Procedure: DQINSERT\_FRONT (Q,F,R,N,Y)

- This procedure **inserts Y** at **front** end of the Circular Queue.
- Queue is represented by a vector **Q** containing **N** elements.
- **F** is pointer to the **front** element of a queue.
- **R** is pointer to the **rear** element of a queue.

## 1. [Overflow?]

```
If      F = 0
Then    Write('Empty')
        Return

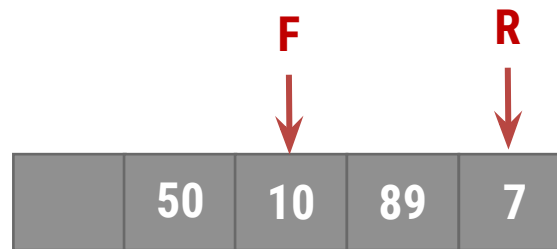
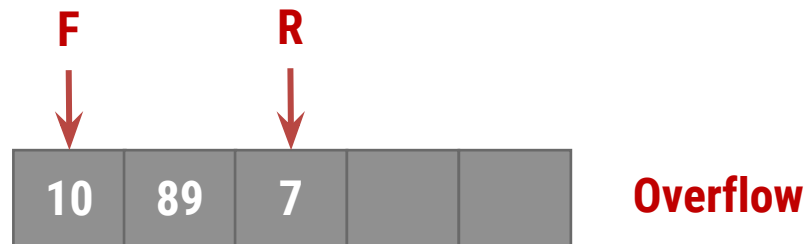
If      F = 1
Then    Write('Overflow')
        Return
```

## 2. [Decrement front Pointer]

```
F ← F - 1
```

## 3. [Insert Element?]

```
Q[F] ← Y
Return
```





# Function: DQDELETE\_REAR(Q,F,R)

- This function **deletes & returns** an element from **rear end** of the Queue.
- **Queue** is represented by a vector **Q** containing **N** elements.
- **F** is pointer to the **front** element of a queue.
- **R** is pointer to the **rear** element of a queue.

## 1. [Underflow?]

```
If      R = 0  
Then  Write('Underflow')  
      Return(0)
```

## 2. [Delete Element]

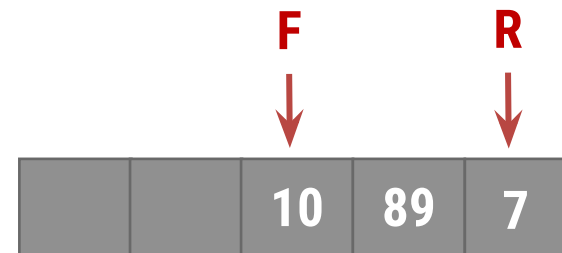
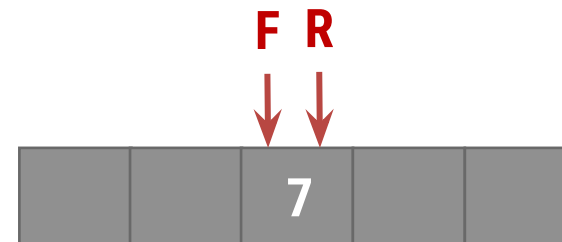
```
Y ← Q[R]
```

## 3. [Queue Empty?]

```
IF      R = F  
Then  R ← F ← 0  
Else  R ← R - 1
```

## 4. [Return Element]

```
Return(Y)
```



# Priority Queue

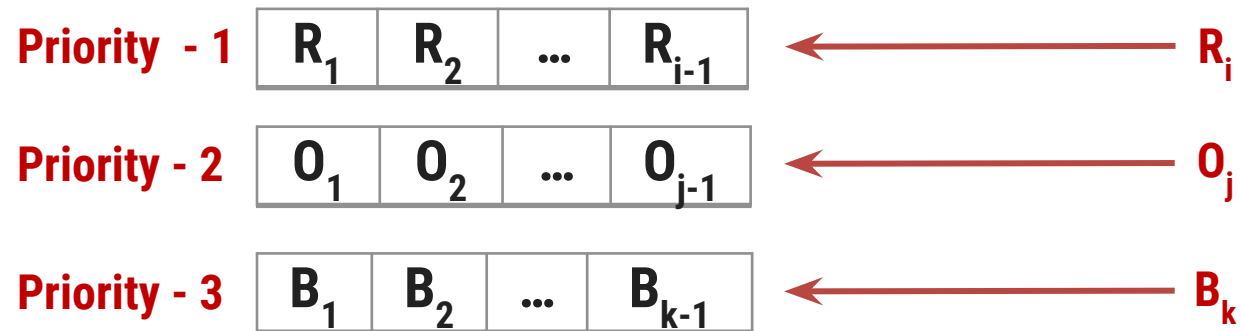
- A queue in which we are able to **insert & remove items** from **any position based on** some property (such as **priority** of the task to be processed) is often referred as **priority queue**.
- Below fig. represent a priority queue of jobs waiting to use a computer.
- Priorities are attached with each Job
  - **Priority 1** indicates **Real Time Job**
  - **Priority 2** indicates **Online Job**
  - **Priority 3** indicates **Batch Processing Job**
- Therefore if a job is initiated with priority  $i$ , it is inserted immediately at the end of list of other jobs with priorities  $i$ .
- Here jobs are always removed from the front of queue.

# Priority Queue Cont...

<b>Task</b>	$R_1$	$R_2$	...	$R_{i-1}$	$O_1$	$O_2$	...	$O_{j-1}$	$B_1$	$B_2$	...	$B_{k-1}$	...
<b>Priority</b>	1	1	...	1	2	2	...	2	3	3	...	3	...

$R_i$                        $O_j$                        $B_k$

Priority Queue viewed as a single queue with insertion allowed at any position



Priority Queue viewed as a Viewed as a set of queue

***Thank  
You***

