資料庫管理 HW03

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1. (a) Left join all advisors(e) and their advisees(s), if someone has no advisee, then s would be NULL.

```
select e.id, e.name from employee as e
left join employee s on e.id=s.supervisor_id
where s.supervisor_id is null;
```

(b) Find the latest store id of each employee before 2025-01-05 and left join to the employee table.

```
select e.id as employee_id, h.store_id
from employee e
left join employee_store_history h
on e.id = h.employee_id
and h.start_date_time=(
select max(h2.start_date_time)
from employee_store_history h2
where h2.employee_id = e.id
and h2.start_date_time <= '2025-01-05'
);</pre>
```

(c) Using limit 1 to obtain the first store id and limit 1 offset 1 to obtain the second store id (after skipping the first one), then join them to produce the final result.

```
select e.id as employee_id,

(select h1.store_id from employee_store_history h1 where h1.employee_id=e.id order by

h1.start_date_time limit 1) as first_store_id,

(select h2.store_id from employee_store_history h2 where h2.employee_id=e.id order by

h2.start_date_time limit 1 offset 1) as second_store_id

from employee e;
```

(d) Calculating the total quantity purchased for each product, then ordering by total quantity and product_id, using limit 2 offset 3 to find the 4th and 5th products. Finally, joining with purchase_detail and purchase tables to get the required information.

```
with total_qty as(
select pd.product_id as product_id, sum(pd.qty) as total_qty, count(*) as purchase_count
from purchase_detail pd
```

```
group by pd.product_id
5),
6 target_product as(
       select product_id
       from total_qty
       order by total_qty desc, product_id asc
9
       limit 2 offset 3
10
11 )
  select p.id as product_id, p.name as product_name, pu.store_id as store_id, count(*) as

→ purchase_count, sum(pd.qty) as total_qty

13 from target_product tp
   join product p on tp.product_id = p.id
15  join purchase_detail pd on pd.product_id = p.id
16 join purchase pu on pd.purchase_no = pu.purchase_no
17 group by p.id, p.name, pu.store_id
  order by p.id, pu.store_id;
```

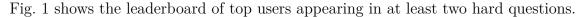
(e) First choose the target product ids by ranking the total purchased quantity, then cross join with store table to get all combinations of target products and stores. Finally, left join with purchase and purchase_detail tables to get the required information.

```
1 with total_qty as(
       select pd.product_id as product_id, row_number() over(order by sum(pd.qty) desc,
       → pd.product_id asc) as rnk
       from purchase_detail pd
3
       group by pd.product_id
5),
6
  target as(
       select product_id
8
       from total_qty
       where rnk >= 4 and rnk <= 5
9
10 )
11 select t.product_id as id, p.name as name, s.id as store_id, coalesce(sum(pd.qty), 0) as amount,
   \hookrightarrow coalesce(count(distinct pu.purchase_no), 0) as cnt
12 from target t --target store id
13 cross join store s
14 left join purchase pu on pu.store_id = s.id
15 left join purchase_detail pd on pu.purchase_no = pd.purchase_no and t.product_id = pd.product_id
   join product p on p.id = t.product_id
  group by s.id, t.product_id, p.name
   order by t.product_id, s.id;
```

(f) Calculate the total spending of each member in each store, then rank them within each store based on the total spending. Using rank() instead of row_number() to handle ties in spending amounts. While ranking, using partition by to separate rankings for each store.

```
6 group by sa.store_id, sa.member_id
7 order by sa.store_id, rnk;
```

2. (a) First, I identified the hard questions based on the number of total answers and correct answers. Then, for each hard question, I found each user's first correct submission and ranked the users by cost time and submission timestamp. Finally, I counted the number of hard questions where each user ranked in the top 3 and kept only those who appeared in at least two questions.



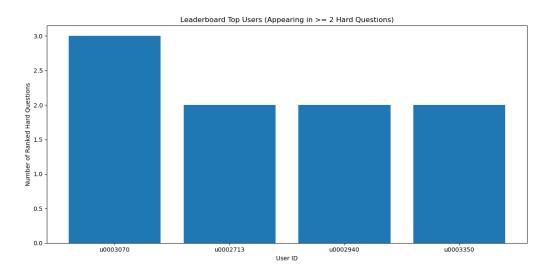


Figure 1: Leaderboard of Top Users Appearing in at Least Two Hard Questions

Following is my source code, I used DuckDB to execute the SQL query and Matplotlib to visualize the results.

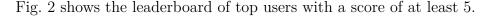
```
import duckdb
   import matplotlib.pyplot as plt
3
   con = duckdb.connect('my_database.db')
5
   df = con.sql("""
6
   -- list the hard questions
8
   with hard_questions as(
       select question_id as id, count(*) as total_ans, sum(is_correct) as total_ac
9
10
11
       group by question id
12
       having total_ans >= 1000 and total_ac <= 500
13
   ),-- list the rank for all users
   first_correct_pre as(
14
       select question_id, user_id, cost_time, created_at, rank() over(partition by question_id,

→ user_id order by created_at asc) as rnk

16
       from answers
17
       where is_correct = 1
   ).-- first correct for each user
18
   first_correct as(
       select question_id, user_id, cost_time, created_at
20
```

```
21
       from first_correct_pre
       where rnk = 1
23 ),--rank users for hard questions
       select question_id, user_id, rank() over(partition by question_id order by cost_time asc,
25
       from first_correct fc
       join hard_questions hq on hq.id = fc.question_id
27
28
   select user_id, count(distinct question_id) as appr_cnt
29
30
  from usr_rnk
   where rnk <= 3
   group by user_id
32
33 having appr_cnt >= 2
   order by appr_cnt desc, user_id asc;
   """).df()
35
36
   print(df)
37
38
  plt.figure(figsize=(12,6))
39 plt.bar(df["user_id"].astype(str), df["appr_cnt"])
40 plt.xlabel("User ID")
   plt.ylabel("Number of Ranked Hard Questions")
   plt.title("Leaderboard Top Users (Appearing in >= 2 Hard Questions)")
  plt.tight_layout()
44 plt.savefig("2a_leaderboard.png")
  plt.show()
```

(b) Using the same CTEs as in 2a to find the users who ranked in the top 3 for hard questions. Then I calculated their scores based on the given scoring system. Finally, I kept only those users with a total score of at least 5.



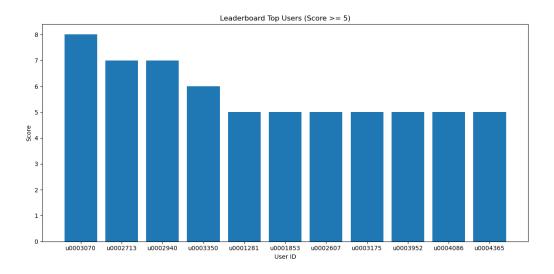


Figure 2: Leaderboard of Top Users with a Score of at Least 5

Following is my source code:

```
1 import duckdb
2 import matplotlib.pyplot as plt
```

```
con = duckdb.connect('my_database.db')
5
   df = con.sql("""
   -- list the hard questions
   with hard_questions as(
       select question_id as id, count(*) as total_ans, sum(is_correct) as total_ac
       from answers
10
11
       group by question_id
       having total_ans >= 1000 and total_ac <= 500
12
  ),-- list the rank for all users
13
   first_correct_pre as(
       select question_id, user_id, cost_time, created_at, rank() over(partition by question_id,
15
       \hookrightarrow user_id order by created_at asc) as rnk
16
       from answers
       where is_correct = 1
17
   ),-- first correct for each user
19
   first correct as(
20
       select question_id, user_id, cost_time, created_at
21
       from first_correct_pre
       where rnk = 1
22
23 ), -- rank users for hard questions
   usr_rnk as(
       select question_id, user_id, rank() over(partition by question_id order by cost_time asc,
25
       from first_correct fc
26
27
       join hard_questions hq on hq.id = fc.question_id
28
   select user_id,
29
30
          sum(
              cast(rnk = 1 as int) * 5 +
31
              cast(rnk = 2 as int) * 2 +
32
33
              cast(rnk = 3 as int) * 1
          ) as score
34
35 from usr_rnk
  where rnk <= 3
   group by user_id
37
   having score >= 5
   order by score desc, user_id asc;
39
   """).df()
   print(df)
41
42
   plt.figure(figsize=(12,6))
   plt.bar(df["user_id"].astype(str), df["score"])
  plt.xlabel("User ID")
  plt.ylabel("Score")
47 plt.title("Leaderboard Top Users (Score >= 5)")
   plt.tight_layout()
   plt.savefig("2b_leaderboard.png")
  plt.show()
```

3. (a) i. I'll use **Surrogate Key** as the primary key because Nature Key might be ambiguous. For example, multiple vehicles may be scheduled at the same time, and a composite Natural Key (such as vehicle ID, departure station, destination station, and time) would be too complex and hard to maintain.

I'll use Serial Number instead of UUID because dispatch record only exists

in a single table and does not require global uniqueness. Serial numbers are simpler, easier to manage, and provide better indexing performance for frequent insertions.

ii. Yes, I' ll add a "soft delete flag" to the dispatch record table. Dispatch records are important historical data that may be needed for future analysis or auditing, so they should not be physically deleted.

Soft deletion allows us to retain the record while marking it as inactive, preventing accidental data loss and preserving referential integrity with related tables. It also provides flexibility to restore records if needed.

iii. As Fig 3 shows:

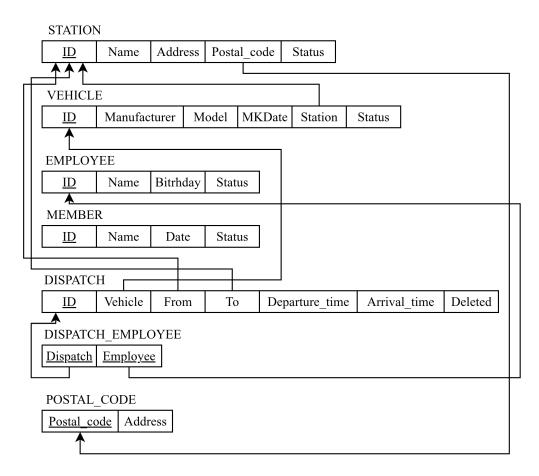


Figure 3: Relational Schema Diagram for Dispatch Record

(b) As Fig 4 shows:

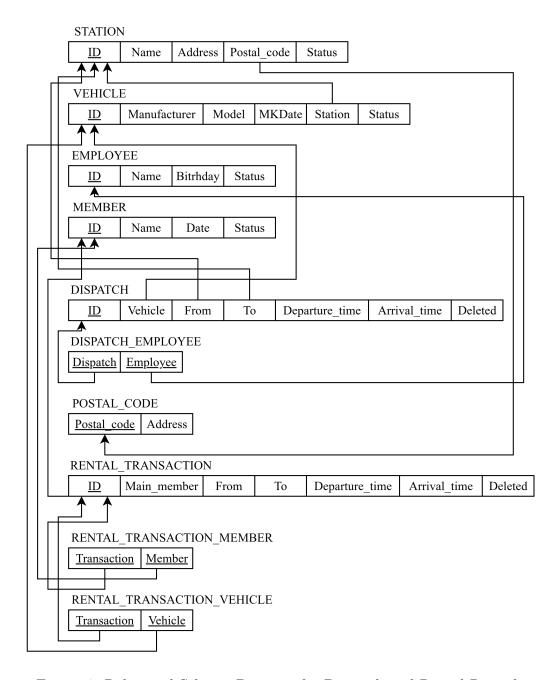


Figure 4: Relational Schema Diagram for Dispatch and Rental Record

4. As Table 1 shows. I know that if it violates a lower normal form, it must violate all higher normal forms. But this table only discusses each normal form separately.

| NF | Check | Reason |
|----|-------|--|
| 1 | No | It has multiple employees recorded in one table |
| 2 | No | ${\tt make_date}$ only related to the ${\tt v_id}$, but the primary key is com- |
| | | posite with v_id and from_datetime. |
| 3 | Yes | |
| BC | Yes | |
| 4 | Yes | |

Table 1: Normal Form Check for Vehicle Maintenance Record

5. As Table 2 shows. I know that if it violates a lower normal form, it must violate all higher normal forms. But this table only discusses each normal form separately.

| NF | Check | Reason |
|----|-------|---|
| 1 | Yes | |
| 2 | Yes | |
| 3 | No | In RELOCATION table, make_date is related to v_id and v_id is not |
| | | a primary key. |
| BC | Yes | |
| 4 | Yes | |

Table 2: Normal Form Check for Vehicle Maintenance Record