

Binary Relations Homework. Dmitry Semenov, M3100, ISU 409537

1. (a) $M_1 = \mathbb{R}$

$$xR_1y \leftrightarrow |x - y| \leq 1$$

(b) $M_2 = P(\{a, b, c\}) = \{\emptyset, \{a\}, \{b\}, \{c\}, \{a, b\}, \{a, c\}, \{b, c\}, \{a, b, c\}\}$

$$R_2 = " \subseteq "$$

(c) $M_3 = \{a, b, c, d\}$

$$||R_3|| = \begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

(d) $M_4 = \{"rock", "scissors", "paper"\}$

$$R_4 = \{\langle x, y \rangle | x \text{ beats } y\}$$

	R1	R2	R3	R4
Reflexive	<i>true</i>	<i>true</i>	<i>false : (x) = (a)</i>	<i>false : (x) = ("rock")</i>
Irreflexive	<i>false : (x) = (1)</i>	<i>false : (x) = ({a})</i>	<i>true</i>	<i>true</i>
Coreflexive	<i>false : (x, y) = (1, 2)</i>	<i>false : (x, y) = ({a}, {a, b})</i>	<i>false : (x, y) = (a, b)</i>	<i>false : (x, y) = ("paper", "rock")</i>
Symmetric	<i>true</i>	<i>false : (x, y) = ({a}, {a, b})</i>	<i>false : (x, y) = (a, d)</i>	<i>false : (x, y) = ("paper", "rock")</i>
Antisymmetric	<i>false : (x, y) = (1, 2)</i>	<i>true</i>	<i>true</i>	<i>true</i>
Asymmetric	<i>false : (x, y) = (1, 2)</i>	<i>false : (x, y) = ({a}, {a})</i>	<i>true</i>	<i>true</i>
Transitive	<i>false : (x, y, z) = (0, 1, 2)</i>	<i>true</i>	<i>false : (x, y, z) = (c, a, d)</i>	<i>false : (x, y, z) = ("rock", "scissors", "paper")</i>
Antitransitive	<i>false : (x, y, z) = (0, 0.5, 1)</i>	<i>false : (x, y, z) = ({\emptyset}, {a}, {a, b})</i>	<i>false : (x, y, z) = (a, b, d)</i>	<i>true</i>
Semiconnex	<i>false : (x, y) = (0, 2)</i>	<i>false : (x, y) = ({a}, {c})</i>	<i>true</i>	<i>true</i>
Connex	<i>false : (x, y) = (0, 2)</i>	<i>false : (x, y) = ({a}, {c})</i>	<i>false : (x, y) = (a, b)</i>	<i>false : (x, y) = ("rock", "rock")</i>
Left Euclidean	<i>false : (x, y, z) = (1, 0, 2)</i>	<i>false : ({a, b}, {a}, {b})</i>	<i>false : (x, y, z) = (a, b, c)</i>	<i>false : (x, y, z) = ("scissors", "rock", "rock")</i>
Right Euclidean	<i>false : (x, y, z) = (1, 0, 2)</i>	<i>false : (x, y, z) = ({a}, {a, b}, {a, c})</i>	<i>false : (x, y, z) = (c, b, a)</i>	<i>false : (x, y, z) = ("rock", "scissors", "scissors")</i>

**true* – valid for all x, y, z values from M

2. $R \subseteq M^2$ and $S \subseteq M^2$

(a) If R and S are reflexive, then $R \cap S$ is so.

Let $a \in M$.

$(a, a) \in R \wedge (a, a) \in S \Rightarrow$
 $(a, a) \in R \cap S \Rightarrow R \cap S$ is reflexive

(b) If R and S are symmetric, then $R \cap S$ is so.

Let $(a, b) \in R \cap S$.

$(a, b) \in R \cap S \Rightarrow$
 $(a, b) \in R \wedge (a, b) \in S \Rightarrow$
 $(b, a) \in R \wedge (b, a) \in S \Rightarrow$
 $(b, a) \in R \cap S \Rightarrow R \cap S$ is symmetric

(c) If R and S are transitive, then $R \cap S$ is so.

Let $(a, b) \in R \cap S \wedge (b, c) \in R \cap S$.

$(a, b) \in R \cap S \wedge (b, c) \in R \cap S \Rightarrow$
 $(a, b), (b, c) \in R \wedge (a, b), (b, c) \in S \Rightarrow$
 $(a, c) \in R \wedge (a, c) \in S \Rightarrow$
 $(a, c) \in R \cap S \Rightarrow R \cap S$ is transitive

(d) If R and S are reflexive, then $R \cup S$ is so.

Let $a \in M$.

$(a, a) \in R \vee (a, a) \in S \Rightarrow$
 $(a, a) \in R \cup S \Rightarrow R \cup S$ is reflexive

(e) If R and S are symmetric, then $R \cup S$ is so.

Let $(a, b) \in R \cup S$.

$(a, b) \in R \cup S \Rightarrow$
 $(a, b) \in R \vee (a, b) \in S \Rightarrow$
 $(b, a) \in R \vee (b, a) \in S \Rightarrow$
 $(b, a) \in R \cup S \Rightarrow R \cup S$ is symmetric

(f) If R and S are transitive, then $R \cup S$ is so.

Counter-example:

Let $(e, a) \in R \wedge (a, b) \in R \wedge (e, b) \in R$
 $(b, c) \in S \wedge (c, d) \in S \wedge (b, d) \in S$

So, $R \cup S$ will contain $(a, b), (b, c)$, but not (a, c) .

And hence, $R \cup C$ is not transitive.

3. An equinumerosity relation \sim over sets is defined as follows: $A \sim B \leftrightarrow |A| = |B|$.

(a) Show that \sim is an equivalence relation over finite sets.

- $(A, A) \in R \forall (A \in M) \leftrightarrow |A| = |A|$ - reflexive
- $(\forall A, B \in M : (A, B) \in R \leftrightarrow (B, A) \in R) \leftrightarrow ((|A| = |B|) \leftrightarrow (|B| = |A|))$ - symmetric
- $(\forall A, B, C \in M : (A, B) \in R \wedge (B, C) \in R \rightarrow (A, C) \in R) \leftrightarrow (|A| = |B| \wedge |B| = |C| \rightarrow |A| = |C|)$ - transitive

(b) Show that \sim is an equivalence relation over infinite sets.

For infinite sets, $|A| = |B|$ means there is a bijection between A and B .

Thus, in order to prove the equivalence of the equinumerosity relation, it is necessary to specify a bijection between sets for each of the properties described above.

- The identical mapping $id : A \rightarrow A$ is bijection
- If $f : A \rightarrow B$ is a bijection, then the inverse mapping $f^{-1} : B \rightarrow A$ is a bijection.
- If $f : A \rightarrow B$ and $g : B \rightarrow C$ are bijections, then their composition is $f \circ g : A \rightarrow C$ is bijection.

(c) Find the quotient set of $P(\{a, b, c, d\})$ by \sim .

Let $W = P(a, b, c, d) =$

$\{\emptyset, \{a\}, \{b\}, \{c\}, \{d\}, \{a, b\}, \{a, c\}, \{a, d\}, \{b, c\}, \{b, d\}, \{c, d\}, \{a, b, c\}, \{a, b, d\}, \{a, c, d\}, \{b, c, d\}, \{a, b, c, d\}\}$

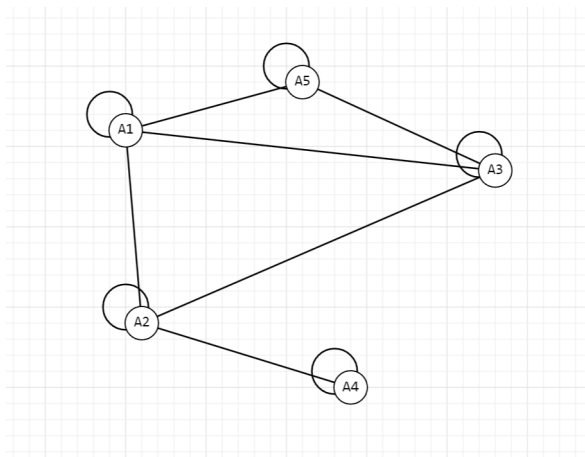
- $[\emptyset]_R = \emptyset$
- $[\{a\}]_R = [\{b\}]_R = [\{c\}]_R = [\{d\}]_R = \{\{a\}, \{b\}, \{c\}, \{d\}\}$
- $[\{a, b\}]_R = [\{a, c\}]_R = [\{a, d\}]_R = [\{b, c\}]_R = [\{b, d\}]_R = [\{c, d\}]_R = \{\{a, b\}, \{a, c\}, \{a, d\}, \{b, c\}, \{b, d\}, \{c, d\}\}$
- $[\{a, b, c\}]_R = [\{a, b, d\}]_R = [\{a, c, d\}]_R = [\{b, c, d\}]_R = \{\{a, b, c\}, \{a, b, d\}, \{a, c, d\}, \{b, c, d\}\}$
- $[\{a, b, c, d\}]_R = \{\{a, b, c, d\}\}$

$W/R = [W]_R = \{[x]_R : x \in W\} = \{\{a\}, \{a, b\}, \{a, b, c\}, \{a, b, c, d\}\}$

4. Let R_θ be a relation of θ -similarity (clearly, $\theta \in [0; 1] \subseteq R$) of finite non-empty sets defined as follows: a set A is said to be θ -similar to B iff the Jaccard index $Jac(A, B) = \frac{|A \cap B|}{|A \cup B|}$ for these sets is at least θ , i.e. $\langle A, B \rangle \in R_\theta \leftrightarrow Jac(A, B) \geq \theta$.

(a) Draw the graph of a relation $R_\theta \subseteq \{A_i\}^2$, where $\theta = 0.25$, $A_1 = \{1, 2, 5, 6\}$, $A_2 = \{2, 3, 4, 5, 7, 9\}$, $A_3 = \{1, 4, 5, 6\}$, $A_4 = \{3, 7, 9\}$, $A_5 = \{1, 5, 6, 8, 9\}$.

- $Jac(A_1, A_2) = \frac{|\{2, 5\}|}{|\{1, 2, 3, 4, 5, 6, 7, 9\}|} = 0.25 \Rightarrow \langle A_1, A_2 \rangle \in R_\theta$
- $Jac(A_1, A_3) = \frac{|\{1, 5, 6\}|}{|\{1, 2, 4, 5, 6\}|} = 0.6 \Rightarrow \langle A_1, A_3 \rangle \in R_\theta$
- $Jac(A_1, A_4) = \frac{|\{ \emptyset \}|}{|\{1, 2, 3, 5, 6, 7, 9\}|} = 0 \Rightarrow \langle A_1, A_4 \rangle \notin R_\theta$
- $Jac(A_1, A_5) = \frac{|\{1, 5, 6\}|}{|\{1, 2, 5, 6, 8, 9\}|} = 0.5 \Rightarrow \langle A_1, A_5 \rangle \in R_\theta$
- $Jac(A_2, A_3) = \frac{|\{4, 5\}|}{|\{1, 2, 3, 4, 5, 6, 7, 9\}|} = 0.25 \Rightarrow \langle A_2, A_3 \rangle \in R_\theta$
- $Jac(A_2, A_4) = \frac{|\{3, 7, 9\}|}{|\{2, 3, 4, 5, 7, 9\}|} = 0.5 \Rightarrow \langle A_2, A_4 \rangle \in R_\theta$
- $Jac(A_2, A_5) = \frac{|\{5, 9\}|}{|\{1, 2, 3, 4, 5, 6, 7, 8, 9\}|} = 0.2 \Rightarrow \langle A_2, A_5 \rangle \notin R_\theta$
- $Jac(A_3, A_4) = \frac{|\{ \emptyset \}|}{|\{1, 3, 4, 5, 6, 7, 9\}|} = 0 \Rightarrow \langle A_3, A_4 \rangle \notin R_\theta$
- $Jac(A_3, A_5) = \frac{|\{1, 5, 6\}|}{|\{1, 4, 5, 6, 8, 9\}|} = 0.5 \Rightarrow \langle A_3, A_5 \rangle \in R_\theta$
- $Jac(A_4, A_5) = \frac{|\{9\}|}{|\{1, 3, 5, 6, 7, 8, 9\}|} = 0.14 \Rightarrow \langle A_4, A_5 \rangle \notin R_\theta$
- $Jac(A_i, A_i) = \frac{|A_i|}{|A_i|} = 1 \Rightarrow \langle A_i, A_i \rangle \in R_\theta \forall i \in [1; 5], i \in \mathbb{N}$



(b) Determine whether θ -similarity is a tolerance relation.

- $Jac(A, A) = \frac{|A \cap A|}{|A \cup A|} = \frac{|A|}{|A|} = 1 \Leftrightarrow \forall \theta \in [0; 1] \quad Jac(A, A) \geq \theta \Leftrightarrow \text{reflexive}$
- Let $(A, B) \in R_\theta \Rightarrow Jac(A, B) \geq \theta \Rightarrow \frac{|A \cap B|}{|A \cup B|} \geq \theta \Rightarrow \frac{|B \cap A|}{|B \cup A|} \geq \theta \Rightarrow Jac(B, A) \geq \theta \Rightarrow (B, A) \in R_\theta \Rightarrow \text{symmetric}$

\Rightarrow tolerance relation

(c) Determine whether θ -similarity is an equivalence relation.

- Based on the previous paragraph, we can conclude about reflexive and symmetric properties
- Counterexample to the transitivity property: $A = \{1\}, B = \{1, 2\}, C = \{2\}, \theta = 0.239. Jac(A, B) = 0.5 \Rightarrow (A, B) \in R_\theta. Jac(B, C) = 0.5 \Rightarrow (B, C) \in R_\theta. Jac(A, C) = 0 \Rightarrow (A, C) \notin R_\theta \Rightarrow \text{not transitive}$
- It is worth noting that θ -similarity can be transitive when $\theta \in \{0\} \cup (0.5; 1]$.

\Rightarrow not always equivalence relation

5. The characteristic function f_S of a set S is defined as follows:

$$f_S(x) = \begin{cases} 1 & \text{if } x \in S \\ 0 & \text{if } x \notin S \end{cases}$$

Let A and B be finite sets. Show that for all $x \in \mathbb{U}$:

(a) $f_{\bar{A}}(x) = 1 - f_A(x)$

- Let $x \in A \Rightarrow x \notin \bar{A}. f_{\bar{A}}(x) = 0 = (1 - f_A(x) = 1 - 1 = 0)$
- Let $x \notin A \Rightarrow x \in \bar{A}. f_{\bar{A}}(x) = 1 = (1 - f_A(x) = 1 - 0 = 1)$

(b) $f_{A \cap B}(x) = f_A(x) * f_B(x)$

- Let $x \in A \wedge x \in B \Rightarrow x \in (A \cap B). f_{A \cap B} = 1 = (f_A(x) * f_B(x) = 1 * 1 = 1)$
- Let $x \in A \wedge x \notin B \Rightarrow x \notin (A \cap B). f_{A \cap B} = 0 = (f_A(x) * f_B(x) = 1 * 0 = 0)$
- Let $x \notin A \wedge x \notin B \Rightarrow x \notin (A \cap B). f_{A \cap B} = 0 = (f_A(x) * f_B(x) = 0 * 0 = 0)$
- Let $x \notin A \wedge x \in B \Rightarrow x \notin (A \cap B). f_{A \cap B} = 0 = (f_A(x) * f_B(x) = 0 * 1 = 0)$

(c) $f_{A \cup B}(x) = f_A(x) + f_B(x) - f_A(x) * f_B(x)$

- Let $x \in A \wedge x \in B \Rightarrow x \in (A \cup B). f_{A \cup B} = 1 = (f_A(x) + f_B(x) - f_A(x) * f_B(x) = 1 + 1 - 1 = 1)$
- Let $x \in A \wedge x \notin B \Rightarrow x \in (A \cup B). f_{A \cup B} = 1 = (f_A(x) + f_B(x) - f_A(x) * f_B(x) = 1 + 0 - 0 = 1)$
- Let $x \notin A \wedge x \in B \Rightarrow x \in (A \cup B). f_{A \cup B} = 1 = (f_A(x) + f_B(x) - f_A(x) * f_B(x) = 0 + 1 - 0 = 1)$
- Let $x \notin A \wedge x \notin B \Rightarrow x \notin (A \cup B). f_{A \cup B} = 0 = (f_A(x) + f_B(x) - f_A(x) * f_B(x) = 0 + 0 - 0 = 0)$

(d) $f_{A \oplus B}(x) = f_A(x) + f_B(x) - 2 * f_A(x) * f_B(x)$

- Let $x \in A \wedge x \in B \Rightarrow x \notin (A \oplus B). f_{A \oplus B} = 0 = (f_A(x) + f_B(x) - 2 * f_A(x) * f_B(x) = 1 + 1 - 2 = 0)$
- Let $x \in A \wedge x \notin B \Rightarrow x \in (A \oplus B). f_{A \oplus B} = 1 = (f_A(x) + f_B(x) - 2 * f_A(x) * f_B(x) = 1 + 0 - 0 = 1)$
- Let $x \notin A \wedge x \in B \Rightarrow x \in (A \oplus B). f_{A \oplus B} = 1 = (f_A(x) + f_B(x) - 2 * f_A(x) * f_B(x) = 0 + 1 - 0 = 1)$
- Let $x \notin A \wedge x \notin B \Rightarrow x \notin (A \oplus B). f_{A \oplus B} = 0 = (f_A(x) + f_B(x) - 2 * f_A(x) * f_B(x) = 0 + 0 - 0 = 0)$

6. Find the error in the "proof" of the following "theorem".

"Theorem": Let R be a relation on a set A that is symmetric and transitive. Then R is reflexive.

"Proof ": Let $a \in A$. Take an element $b \in A$ such that $\langle a, b \rangle \in R$. Because R is symmetric, we also have $\langle b, a \rangle \in R$. Now using the transitive property, we can conclude that $\langle a, a \rangle \in R$ because $\langle a, b \rangle \in R$ and $\langle b, a \rangle \in R$.

The above proof is incorrect for the following reason.

The case that does not take into account the proof is an empty relation, $\forall x, y \in A : \langle x, y \rangle \notin R$.

Then the relation R is symmetric, since the left part of the implication is always false: $\forall x, y \in A : (xRy) \rightarrow (yRx)$.
 Similarly, the relation R is transitive: $\forall x, y, z \in A : (xRy) \wedge (yRz) \rightarrow (xRz)$.

At the same time, the relation R is not reflexive, since $\forall x \in A : \neg(xRx)$

7. Give an example of a relation R on the set $\{a, b, c\}$ such that the symmetric closure of the reflexive closure of the transitive closure of R is not transitive.

Let $R = \{\langle a, b \rangle, \langle a, c \rangle\}$

$R^t = \{\langle a, b \rangle, \langle a, c \rangle\}$

$(R^t)^r = \{\langle a, b \rangle, \langle a, c \rangle, \langle a, a \rangle, \langle b, b \rangle, \langle c, c \rangle\}$

$((R^t)^r)^s = \{\langle a, b \rangle, \langle a, c \rangle, \langle a, a \rangle, \langle b, b \rangle, \langle c, c \rangle, \langle b, a \rangle, \langle c, a \rangle\}$

$((R^t)^r)^s$ is not transitive. Counterexample to the transitivity property: $\langle b, a \rangle \in ((R^t)^r)^s \wedge \langle a, c \rangle \in ((R^t)^r)^s$, but $\langle b, c \rangle \notin ((R^t)^r)^s$.

8. Prove or disprove the following statements about the functions f and g :

(a) If f and g are injections, then $g \circ f$ is also an injection.

- Let $f : A \rightarrow B, g : B \rightarrow C, g \circ f = g(f(x)) : A \rightarrow C$
- $\forall a, b : f(a) = f(b) \Leftrightarrow a = b$
- $\forall a, b : g(a) = g(b) \Leftrightarrow a = b$
- $\forall a, b : g(f(a)) = g(f(b)) \Rightarrow f(a) = f(b) \Rightarrow a = b \Rightarrow g \circ f \text{ is injection}$

(b) If f and g are surjections, then $g \circ f$ is also a surjection.

- Let $f : A \rightarrow B, g : B \rightarrow C, g \circ f = g(f(x)) : A \rightarrow C$
- $\forall c \in C \exists b \in B : g(b) = c$
- $\forall b \in B \exists a \in A : f(a) = b$
- $\forall c \in C \forall b \in B \exists a \in A : g(f(a)) = g(b) = c \Rightarrow g \circ f \text{ is surjection}$

(c) If f and $f \circ g$ are injections, then g is also an injection.

- Let $g : A \rightarrow B, f : B \rightarrow C, f \circ g = f(g(x)) : A \rightarrow C$
- $\forall a, b : f(a) = f(b) \Leftrightarrow a = b$
- $\forall a, b : f(g(a)) = f(g(b)) \Leftrightarrow g(a) = g(b)$
- g must be injective, otherwise, it would contradict the fact that $g \circ f$ is injective.

(d) If f and $f \circ g$ are surjections, then g is also a surjection

- Let $g : A \rightarrow B, f : B \rightarrow C, f \circ g = f(g(x)) : A \rightarrow C$
- $\forall c \in C \exists b \in B : f(b) = c$
- A counterexample about the surjectivity of g : $A = \{1, 2\}, B = \{1, 2\}, C = \{0\}$.

$g(1) = g(2) = f(1) = f(2) = 0$. Since, f and $f \circ g$ are surjections, but g is not.

9. Let $H = \{1, 2, 4, 5, 10, 12, 20\}$. Consider a divisibility relation $R \subseteq H^2$ defined as follows: $xRy \Leftrightarrow y \vdots x$.

(a) Sort R (as a set of pairs) lexicographically.

$R = \{\langle 1, 1 \rangle, \langle 1, 2 \rangle, \langle 1, 4 \rangle, \langle 1, 5 \rangle, \langle 1, 10 \rangle, \langle 1, 12 \rangle, \langle 1, 20 \rangle, \langle 2, 2 \rangle, \langle 2, 4 \rangle, \langle 2, 10 \rangle, \langle 2, 12 \rangle, \langle 2, 20 \rangle, \langle 4, 4 \rangle, \langle 4, 12 \rangle, \langle 4, 20 \rangle, \langle 5, 5 \rangle, \langle 5, 10 \rangle, \langle 5, 20 \rangle, \langle 10, 10 \rangle, \langle 10, 20 \rangle, \langle 12, 12 \rangle, \langle 20, 20 \rangle\}$

(b) Show that R is a partial order.

Using the divisibility properties of real numbers:

- $\forall x \in H : (x, x) \in R \Leftrightarrow (x \dot{:} x \Leftrightarrow \frac{x}{x} = 1) \Rightarrow \text{reflexive}$
- $\forall x, y \in H : (x, y) \in R \wedge (y, x) \in R \Rightarrow (x \dot{:} y \wedge y \dot{:} x) \Rightarrow (x = y) \Rightarrow \text{antisymmetric}$
- $\forall x, y, z \in H : (x, y) \in R \wedge (y, z) \in R \Rightarrow (y \dot{:} x \wedge z \dot{:} y)$

$y \dot{:} x \Rightarrow$ all exponents appearing in x prime factorization contains in y prime factorization

$z \dot{:} y \Rightarrow$ all exponents appearing in y prime factorization contains in z prime factorization

\Rightarrow all exponents appearing in x prime factorization contains in z prime factorization $\Rightarrow z \dot{:} x \Rightarrow (x, z) \in R \Rightarrow \text{transitive}$

$\Rightarrow R$ is a partial order

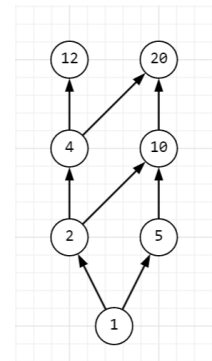
(c) Determine whether R is a linear (total) order.

- Counterexample to the connex property: $x = 10, y = 12 : (x, y) \notin R \wedge (y, x) \notin R$

$\Rightarrow R$ is not a total order

(d) Draw the Hasse diagram for a graded poset $\langle H, R, \rho \rangle$, where $\rho : H \rightarrow \mathbb{N}_0$ is a grading function which maps a number $n \in H$ to the sum of all exponents appearing in its prime factorization, e.g., $\rho(20) = \rho(2^2 * 5^1) = 2 + 1 = 3$, so the number 20 would have the 3rd rank (bottom-up).

- $\rho(1) = \rho(1^0) = 0$
- $\rho(2) = \rho(2^1) = 1$
- $\rho(5) = \rho(5^1) = 1$
- $\rho(4) = \rho(2^2) = 2$
- $\rho(10) = \rho(2^1 * 5^1) = 2$
- $\rho(12) = \rho(2^2 * 3^1) = 3$
- $\rho(20) = \rho(2^2 * 5^1) = 3$



Hasse diagram for $\langle H, R, \rho \rangle$

(e) Find the minimal, minimum (least), maximal and maximum (greatest) elements in the poset $\langle H, R \rangle$. If there are multiple or none, explain why.

- $1 \in H$ is minimal $\leftrightarrow \forall b \neq 1 : (b, 1) \notin R$
- $12 \in H$ is maximal $\leftrightarrow \forall b \neq 12 : (12, b) \notin R$, at the same time, $20 \in H$ is maximal $\leftrightarrow \forall b \neq 20 : (20, b) \notin R$
- greatest element doesn't exist, because $(12, 20) \notin R$
- $1 \in H$ is least $\leftrightarrow \forall b : (1, b) \in R$

(f) Perform a topological sort of the poset $\langle H, R \rangle$

$H_{\text{topological_sorted}} = \{1, 5, 2, 10, 4, 20, 12\}$

10. Prove that the transitive closure R^+ is in fact transitive.

Definition. $R^+ = \bigcup_{n \in \mathbb{N}^+} R^n$ is a transitive closure of $R \subseteq M^2$, where

- $R^{k+1} = R^k \circ R$ is a compositional (functional) power,

- $R^1 = R$,
- $S \circ R = \{\langle x, y \rangle \mid \exists z : (xRz) \wedge (zSy)\}$ is a composition (relative product) of relations R and S .

Let $(a, b) \in R^+ \wedge (b, c) \in R^+$. Then exist are $j, k \in [1, n] \in \mathbb{N} \mid (a, b) \in R^j, (b, c) \in R^k$. Hence $(a, c) \in R^j \circ R^k = R^{j+k}$. Since $R^{j+k} \subseteq R^+$, $(a, c) \in R^+ \Rightarrow R^+$ is transitive.

11. Prove that a set S is infinite if and only if there is a proper subset $A \subset S$ such that there is a one to-one correspondence (bijection) between A and S .

- Let S in infinite set. Since all infinite sets have counable infinite subset, we can construct countable infinite subset of S .
 $A = \{a_1, a_2, a_3, \dots\}$.
- Let's define two partitions of A as follows: $A_1 = \{a_2, a_4, a_6, a_8, \dots\}$, $A_2 = \{a_1, a_3, a_5, a_7, \dots\}$
- We can establish a one-to-one correspondence between A and A_1 like that : $a_n = a_{2n}$.
- Then we can assert that exist a bijection between $A \cup (S \setminus A) = S$ and $A_1 \cup (S \setminus A) = S \setminus A_2$ (mapping each element in $S \setminus A$ to itself). Since $S \setminus A_2 \subset S$, there is bijection between infinite set and it's proper subset.
- Let $A \subset S$, $f : A \rightarrow S$ is bijection. Using theorem about no bijection between finite set and proper subset, S must be infinite.

12. Given a set S and two partitions P_1 and P_2 of S , we define the relation $P_1 \preceq P_2$ as follows: partition P_1 is considered a refinement of the partition P_2 if every set in P_1 is a subset of one of the sets in P_2 . Show that the set of all partitions of a set S with the refinement relation \preceq is a lattice.

- To show that the set of all partitions of a set S with the refinement relation \preceq is a lattice, we need to show that it has a least upper bound and a greatest lower bound for every pair of partitions.
- Least Upper Bound: Let P_1 and P_2 are two partitions of S . To find the supremum of P_1 and P_2 , we need to find a partition P such that $P_1 \preceq P$ and $P_2 \preceq P$, and for any other partition Q , if $P_1 \preceq Q$ and $P_2 \preceq Q$, then $P \preceq Q$.
 - We can construct P by taking the union of each pair of corresponding sets from P_1 and P_2 . For every set A in P_1 and every set B in P_2 , we take the union of A and B and add it to P . \Rightarrow every set in P_1 and P_2 is a subset of some set in P .
 - To show that $P_1 \preceq P$ and $P_2 \preceq P$, we can observe that every set in P_1 is a subset of itself and is also a subset of the union of corresponding sets in P_1 and P_2 . Every set in P_2 is a subset of itself and is also a subset of the union of corresponding sets in P_1 and P_2 .
 - To show that if $P_1 \preceq Q$ and $P_2 \preceq Q$, then $P \preceq Q$ for any other partition Q , we can observe that every set in P_1 and P_2 is a subset of some set in Q . Since P is the union of corresponding sets from P_1 and P_2 , every set in P is also a subset of some set in Q . Hence, $P \preceq Q$.
 - \Rightarrow the supremum of P_1 and P_2 exists.
- Greatest Lower Bound: Let P_1 and P_2 are two partitions of S . To find the infimum of P_1 and P_2 , we need to find a partition P such that $P \preceq P_1$ and $P \preceq P_2$, and for any other partition Q , if $Q \preceq P_1$ and $Q \preceq P_2$, then $Q \preceq P$.
 - We can construct P by taking the intersection of each pair of corresponding sets from P_1 and P_2 . For every set A in P_1 and every set B in P_2 , we take the intersection of A and B and add it to P . \Rightarrow every set in P is a subset of some set in P_1 and P_2 .
 - To show $P \preceq P_1$ and $P \preceq P_2$, we can observe that every set in P is a subset of the corresponding sets in P_1 and P_2 , since it is the intersection of those sets.
 - To show that if $Q \preceq P_1$ and $Q \preceq P_2$, then $Q \preceq P$ for any other partition Q , we can observe that since $Q \preceq P_1$ and $Q \preceq P_2$, every set in Q is a subset of some set in both P_1 and P_2 . Since P is the intersection of corresponding sets from P_1 and P_2 , every set in P is also a subset of some set in Q . Hence, $Q \preceq P$.
 - \Rightarrow the infimum of P_1 and P_2 exists.
- \Rightarrow the set of all partitions of a set S with the refinement relation \preceq is a lattice.

