

# ICPC Templates For Grooming

WAKing @ NWPU

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# 1 数据结构

## 1.1 线段树

```

1 struct SegmentTree// 区间和,区间加
2 {
3     ll data[maxn<<2],lz[maxn<<2];
4     int L[maxn<<2],R[maxn<<2];
5     void pushup(int k) {
6         data[k]=data[k<<1]+data[k<<1|1];
7     }
8     void build(int l,int r,int k) {
9         L[k]=l;
10        R[k]=r;
11        data[k]=lz[k]=0;
12        if(l==r) {
13            return ;
14        }
15        int mid=(l+r)>>1;
16        build(l,mid,k<<1);
17        build(mid+1,r,k<<1|1);
18        pushup(k);
19    }
20    void lzadd(int k,ll x) {
21        lz[k]+=x;
22        data[k]+=(R[k]-L[k]+1)*x;
23    }
24    void pushdown(int k) {
25        if(lz[k]==0) return ;
26        lzadd(k<<1,lz[k]);
27        lzadd(k<<1|1,lz[k]);
28        lz[k]=0;
29    }
30    void change(int pos,int k,ll x) {
31        if(L[k]==R[k]) {
32            data[k]=x;
33            return ;
34        }
35        pushup(k);
36        int mid=(L[k]+R[k])>>1;
37        if(mid>=pos) change(pos,k<<1,x);
38        else change(pos,k<<1|1,x);
39        pushup(k);
40    }
41    void add(int l,int r,int k,ll x) {
42        if(L[k]>=l&&R[k]<=r) {
43            lzadd(k,x);
44            return ;
45        }
46        pushdown(k);
47        int mid=(L[k]+R[k])>>1;
48        if(mid>=l) add(l,r,k<<1,x);
49        if(mid<r) add(l,r,k<<1|1,x);

```

```

50     pushup(k);
51 }
52 ll query(int l,int r,int k) {
53     if(L[k]>=l&&R[k]<=r) return data[k];
54     pushdown(k);
55     ll ans=0;
56     int mid=(L[k]+R[k])>>1;
57     if(mid>=l) ans=ans+query(l,r,k<<1);
58     if(mid<r) ans=ans+query(l,r,k<<1|1);
59     return ans;
60 }
61 }tree;

```

## 1.2 线段树合并

```

1 struct SegmentTree {
2     int cnt=0,L[maxn<<5],R[maxn<<5],data[maxn<<5];// 空间巨大,merge函数可以优化
3     void pushup(int x) {
4         data[x]=data[L[x]]+data[R[x]];
5     }
6     void ins(int &now,int l,int r,int pos) {
7         if(!now) now=++cnt;
8         if(l==r) {
9             data[now]++;
10            return ;
11        }
12        int mid=(l+r)>>1;
13        if(pos<=mid) ins(L[now],l,mid,pos);
14        else ins(R[now],mid+1,r,pos);
15        pushup(now);
16    }
17    int merge(int u,int v,int l,int r) {// 保留u,v 构造一颗新的树
18        if(!u||!v) return u|v;
19        int mid=(l+r)>>1,now=++cnt;
20        if(l==r) data[now]=data[u]+data[v];
21        else {
22            L[now]=merge(L[u],L[v],l,mid);
23            R[now]=merge(R[u],R[v],mid+1,r);
24            pushup(now);
25        }
26        return now;
27    }
28    int merge(int u,int v,int l,int r) {// 将u合并到v, v不能再被使用.节约空间
29        if(!u||!v) return u|v;
30        int mid=(l+r)>>1;
31        if(l==r) {
32            data[u]+=data[v];
33        } else {
34            L[u]=merge(L[u],L[v],l,mid);
35            R[u]=merge(R[u],R[v],mid+1,r);
36            pushup(u);
37        }

```

```

38     return u;
39 }
40 int query(int now,int l,int r,int ql,int qr) { // 区间求和
41     if(ql<=l&&qr>=r) return data[now];
42     int mid=(l+r)>>1,ans=0;
43     if(mid>=ql) ans+=query(L[now],l,mid,ql,qr);
44     if(mid<qr) ans+=query(R[now],mid+1,r,ql,qr);
45     return ans;
46 }
47 };

```

### 1.3 主席树

```

1 struct ZXTree {
2     int cnt=0, T[maxn], L[maxn << 4], R[maxn << 4], sum[maxn << 4];
3     void init() { cnt=0; }
4     int build(int l, int r) {
5         int rt = ++cnt;
6         sum[rt] = 0;
7         if (l < r) {
8             int mid = (l + r) >> 1;
9             L[rt] = build(l, mid);
10            R[rt] = build(mid + 1, r);
11        }
12        return rt;
13    }
14    int update(int pre, int l, int r, int x) {
15        int rt = ++cnt;
16        L[rt] = L[pre];
17        R[rt] = R[pre];
18        sum[rt] = sum[pre] + 1;
19        if (l < r)
20        {
21            int mid = (l + r) >> 1;
22            if (x <= mid)
23                L[rt] = update(L[pre], l, mid, x);
24            else
25                R[rt] = update(R[pre], mid + 1, r, x);
26        }
27        return rt;
28    }
29    int query_kth(int u, int v, int l, int r, int k) { //第k小
30        if (l >= r)
31            return l;
32        int x = sum[L[v]] - sum[L[u]], mid = (l + r) >> 1;
33        if (x >= k)
34            return query_kth(L[u], L[v], l, mid, k);
35        else
36            return query_kth(R[u], R[v], mid + 1, r, k - x);
37    }
38    int query_cnt(int u,int v,int l,int r,int x) { //小于k的数目
39        if(r<=x) return sum[v]-sum[u];

```

```

40     int ans=0,mid=(l+r)>>1;
41     ans+=query_cnt(L[u],L[v],l,mid,x);
42     if(x>mid) ans+=query_cnt(R[u],R[v],mid+1,r,x);
43     return ans;
44 }
45 };

```

## 1.4 splay

```

1 struct Splay {
2     int rt, tot, fa[N], ch[N][2], val[N], cnt[N], sz[N];
3     void maintain(int x) { sz[x] = sz[ch[x][0]] + sz[ch[x][1]] + cnt[x]; }
4     bool get(int x) { return x == ch[fa[x]][1]; }
5     void clear(int x) {
6         ch[x][0] = ch[x][1] = fa[x] = val[x] = sz[x] = cnt[x] = 0;
7     }
8     void rotate(int x) {
9         int y = fa[x], z = fa[y], chk = get(x);
10        ch[y][chk] = ch[x][chk ^ 1];
11        fa[ch[x][chk ^ 1]] = y;
12        ch[x][chk ^ 1] = y;
13        fa[y] = x;
14        fa[x] = z;
15        if (z) ch[z][y == ch[z][1]] = x;
16        maintain(x);
17        maintain(y);
18    }
19    void splay(int x) {
20        for (int f = fa[x]; f = fa[x], f; rotate(x))
21            if (fa[f]) rotate(get(x) == get(f) ? f : x);
22        rt = x;
23    }
24    void ins(int k) {
25        if (!rt) {
26            val[++tot] = k;
27            cnt[tot]++;
28            rt = tot;
29            maintain(rt);
30            return;
31        }
32        int cnr = rt, f = 0;
33        while (1) {
34            if (val[cnr] == k) {
35                cnt[cnr]++;
36                maintain(cnr);
37                maintain(f);
38                splay(cnr);
39                break;
40            }
41            f = cnr;
42            cnr = ch[cnr][val[cnr] < k];
43            if (!cnr) {

```

```
44     val[++tot] = k;
45     cnt[tot]++;
46     fa[tot] = f;
47     ch[f][val[f] < k] = tot;
48     maintain(tot);
49     maintain(f);
50     splay(tot);
51     break;
52 }
53 }
54 }
55 int rk(int k) {
56     int res = 0, cnr = rt;
57     while (1) {
58         if (k < val[cnr]) {
59             cnr = ch[cnr][0];
60         } else {
61             res += sz[ch[cnr][0]];
62             if (k == val[cnr]) {
63                 splay(cnr);
64                 return res + 1;
65             }
66             res += cnt[cnr];
67             cnr = ch[cnr][1];
68         }
69     }
70 }
71 int kth(int k) {
72     int cnr = rt;
73     while (1) {
74         if (ch[cnr][0] && k <= sz[ch[cnr][0]]) {
75             cnr = ch[cnr][0];
76         } else {
77             k -= cnt[cnr] + sz[ch[cnr][0]];
78             if (k <= 0) {
79                 splay(cnr);
80                 return val[cnr];
81             }
82             cnr = ch[cnr][1];
83         }
84     }
85 }
86 int pre() {
87     int cnr = ch[rt][0];
88     while (ch[cnr][1]) cnr = ch[cnr][1];
89     splay(cnr);
90     return cnr;
91 }
92 int nxt() {
93     int cnr = ch[rt][1];
94     while (ch[cnr][0]) cnr = ch[cnr][0];
95     splay(cnr);
96     return cnr;
```

```

97     }
98     void del(int k) {
99         rk(k);
100         if (cnt[rt] > 1) {
101             cnt[rt]--;
102             maintain(rt);
103             return;
104         }
105         if (!ch[rt][0] && !ch[rt][1]) {
106             clear(rt);
107             rt = 0;
108             return;
109         }
110         if (!ch[rt][0]) {
111             int cnr = rt;
112             rt = ch[rt][1];
113             fa[rt] = 0;
114             clear(cnr);
115             return;
116         }
117         if (!ch[rt][1]) {
118             int cnr = rt;
119             rt = ch[rt][0];
120             fa[rt] = 0;
121             clear(cnr);
122             return;
123         }
124         int cnr = rt;
125         int x = pre();
126         fa[ch[cnr][1]] = x;
127         ch[x][1] = ch[cnr][1];
128         clear(cnr);
129         maintain(rt);
130     }
131 } tree;

```

## 1.5 LCT

```

1 namespace LCT{
2     const int N=maxn;
3     int ch[N][2],rev[N],fa[N],stk[N];
4     bool son(int x) {return ch[fa[x]][1] == x;}
5     bool isroot(int x) {return ch[fa[x]][1]!=x && ch[fa[x]][0]!=x;}
6     void reverse(int x) {swap(ch[x][1],ch[x][0]); rev[x] ^= 1;}
7     void pushup(int x) {
8         .....
9     }
10    void pushdown(int x) {if(rev[x])reverse(ch[x][0]) , reverse(ch[x][1]) , rev[x]
        ^= 1;}
11    void rot(int x) {
12        int y = fa[x] , z = fa[y] , c = son(x);
13        if(!isroot(y))ch[z][son(y)] = x; fa[x] = z;

```



```

14     ch[y][c] = ch[x][!c]; fa[ch[y][c]] = y;
15     ch[x][!c] = y; fa[y] = x; pushup(y);
16 }
17 void splay(int x) {
18     int top = 0; stk[++top] = x;
19     for( int i = x; !isroot(i) ; i = fa[i])stk[++top] = fa[i];
20     while(top)pushdown(stk[top--]);
21     for( int y = fa[x]; !isroot(x) ; rot(x) , y = fa[x])
22         if(!isroot(y))son(x) ^ son(y) ? rot(x) : rot(y);
23     pushup(x); return;
24 }
25 void access(int x) { for( int y = 0; x; y = x,x = fa[x])splay(x),ch[x][1] = y,
    pushup(x); }
26 void makeroot(int x) {access(x); splay(x); reverse(x);}
27 int findroot(int x) {access(x); splay(x); while(ch[x][0])x=ch[x][0]; return x;}
28 void split(int x,int y) {makeroot(x); access(y); splay(y);}
29 void link(int x,int y) {if((!x)||(!y))return; makeroot(x); fa[x] = y; }
30 void cut(int x,int y) {if((!x)||(!y))return; split(x,y); ch[y][0] = fa[x] = 0;
    pushup(y);}
31 bool adj(int x,int y) {
32     split(x,y);
33     return ch[y][0]==x&&ch[x][0]==0&&ch[x][1]==0;
34 }
35 }
36 /* 维护子树大小(不是splay的大小)
37 void pushup(int x) {
38     sz[x]=sz[ch[x][0]]+sz[ch[x][1]]+sz2[x]+1;
39 }
40 void access(int x) {
41     for( int y = 0; x; y = x,x = fa[x]) {
42         splay(x);
43         sz2[x]+=sz[ch[x][1]]-sz[y];
44         ch[x][1] = y,pushup(x);
45     }
46 }
47 void link(int x,int y) {
48     split(x,y);
49     fa[x]=y;sz2[y]+=sz[x];
50     pushup(y);
51 }
52 */

```

## 2 图论

### 2.1 迪杰斯特拉

```

1 vector<pair<ll,ll>> g[maxn];
2 ll dis[maxn];
3 void dij(int s) {
4     memset(dis,0x3f,sizeof(dis));
5     dis[s]=0;
6     priority_queue<pair<ll,ll>> q;

```

```

7   q.push({0,s});
8   while(!q.empty()) {
9       pair<ll,ll> now=q.top();q.pop();
10      now.fi=-now.fi;
11      for(auto x:g[now.se]) {
12          if(dis[x.fi]>dis[now.se]+x.se) {
13              dis[x.fi]=dis[now.se]+x.se;
14              q.push({-dis[x.fi],x.fi});
15          }
16      }
17  }
18 }

```

## 2.2 SPFA

```

1  struct edge {
2      int cost, to;
3      edge(int x = 0, int y = 0): to(x), cost(y) {}
4  };
5  vector<edge> g[maxn];
6  int dis[maxn], cnt[maxn], n, m;
7  bool vis[maxn];
8  queue<int> q;
9  bool spfa(int s) {
10     memset(vis, false, sizeof(vis));
11     memset(dis, 0x3f, sizeof(dis));
12     vis[s] = true;
13     dis[s] = 0;
14     while (!q.empty()) q.pop();
15     q.push(s);
16     while (!q.empty()) {
17         int now = q.front(); q.pop();
18         vis[now] = false;
19         for (int i = 0; i < g[now].size(); i++) {
20             int to = g[now][i].to;
21             if (dis[to] > dis[now] + g[now][i].cost) {
22                 dis[to] = dis[now] + g[now][i].cost;
23                 if (!vis[to]) {
24                     vis[to] = true;
25                     q.push(to);
26                     if (++cnt[to] > n) return false;
27                 }
28             }
29         }
30     }
31     return true;
32 }

```

## 2.3 tarjan

```

1  #include <bits/stdc++.h>

```

```

2 using namespace std;
3 typedef long long ll;
4 typedef unsigned long long ull;
5 #define met(s) memset(s, 0, sizeof(s))
6 const int mod = 1000000007;
7 const int maxn = 100000;
8 const int inf = 0x3f3f3f3f;
9 struct node {
10     int to, next;
11 } edge[maxn];
12 int head[maxn], tot;
13 void addedge(int u, int v) {
14     edge[++tot].to = v;
15     edge[tot].next = head[u];
16     head[u] = tot;
17 }
18 int Stack[maxn], top; //栈
19 int dfn[maxn], low[maxn], belong[maxn], scc, Index, num[maxn]; //belong[i]: i属于第i
    个强连通分量, num[i]:第i个强联通分量的个数
20 bool Instack[maxn];
21 void tarjan(int u) {
22     int v;
23     low[u] = dfn[u] = ++Index;
24     Stack[top++] = u;
25     Instack[u] = true;
26     for (int i = head[u]; i; i = edge[i].next) {
27         int v = edge[i].to;
28         if (!dfn[v]) {
29             tarjan(v);
30             if (low[u] > low[v]) low[u] = low[v];
31         }
32         else if (Instack[v] && low[u] > low[v]) low[u] = low[v];
33     }
34     if (low[u] == dfn[u]) {
35         scc++;
36         do {
37             v = Stack[--top];
38             Instack[v] = false;
39             belong[v] = scc;
40             num[scc]++;
41         } while (v != u);
42     }
43 }
44 void solve(int n) {
45     memset(dfn, 0, sizeof(dfn));
46     memset(Instack, false, sizeof(Instack));
47     memset(num, 0, sizeof(num));
48     Index = scc = top = 0;
49     for (int i = 1; i <= n; i++) {
50         if (!dfn[i]) tarjan(i);
51     }
52 }
53 void init() {

```

```

54     tot = 0;
55     memset(head, 0, sizeof(head));
56 }

```

## 2.4 LCA 倍增

```

1  vector<int> g[maxn];
2  int dep[maxn], father[maxn][maxbit], lg[maxn];
3  namespace LCA {
4      void init() {
5          lg[0] = -1;
6          for (int i = 1; i < maxn; i++) lg[i] = lg[i >> 1] + 1;
7      }
8      void dfs(int now, int fa) {
9          dep[now] = dep[fa] + 1;
10         father[now][0] = fa;
11         for (int i = 1; i <= lg[dep[now]]; i++)
12             father[now][i] = father[father[now][i - 1]][i - 1];
13         for (int i = 0; i < g[now].size(); i++) {
14             if (g[now][i] != fa) dfs(g[now][i], now);
15         }
16     }
17     int LCA(int u, int v) {
18         if (dep[u] < dep[v]) swap(u, v);
19         while (dep[u] != dep[v]) u = father[u][lg[dep[u]] - dep[v]];
20         if (u == v) return u;
21         for (int i = lg[dep[u]]; i >= 0; i--) {
22             if (father[u][i] != father[v][i]) {
23                 u = father[u][i];
24                 v = father[v][i];
25             }
26         }
27         return father[u][0];
28     }
29 }

```

## 2.5 虚树

```

1  vi g2[maxn];
2  namespace faketree {
3      int sta[maxn], top=0;
4      void insert(int x) {
5          if (top==0) {
6              sta[++top]=x;
7              return ;
8          }
9          int lca=LCA::LCA(sta[top], x);
10         if (lca==sta[top]) {
11             sta[++top]=x;
12             return ;
13         }

```

```

14     while(top>1&&dep[lca]<=dep[sta[top-1]]) {
15         g2[sta[top-1]].pb(sta[top]);
16         top--;
17     }
18     if(lca!=sta[top]) {
19         g2[lca].pb(sta[top]);
20         sta[top]=lca;
21     }
22     sta[++top]=x;
23 }
24 void end() {
25     while(top>1) {
26         g2[sta[top-1]].pb(sta[top]);
27         top--;
28     }
29     top=0;
30 }
31 }

```

## 2.6 树链剖分

```

1  int dep[maxn],son[maxn],siz[maxn],tp[maxn],fa[maxn];// tp链顶
2  int id[maxn],rk[maxn],cnt=0;// id dfs序, rk 反dfs序
3  vector<int> g[maxn];
4  void dfs1(int now,int f,int d) {
5      dep[now]=d,fa[now]=f,siz[now]=1;
6      for(int x:g[now]) {
7          if(x==f) continue;
8          dfs1(x,now,d+1);
9          siz[now]+=siz[x];
10         if(siz[x]>siz[son[now]]) son[now]=x;
11     }
12 }
13 void dfs2(int now,int t) {
14     tp[now]=t,id[now]=++cnt,rk[cnt]=now;
15     if(!son[now]) return ;
16     dfs2(son[now],t);
17     for(int x: g[now]) {
18         if(x==fa[now]||x==son[now]) continue;
19         dfs2(x,x);
20     }
21 }

```

## 2.7 长链剖分

```

1  int dep[maxn],height[maxn],son[maxn];
2  vector<int> g[maxn];
3  void dfs1(int now,int fa) {
4      height[now]=dep[now]=dep[fa]+1;
5      for(int x:g[now]) {
6          if(x==fa) continue;

```

```

7     dfs1(x,now);
8     if(height[x]>height[now]) {
9         height[now]=height[x];
10        son[now]=x;
11    }
12 }
13 }
14 // 求每个节点子树中,距离为k的节点数量,找到一个k,使节点数最多
15 int tmp[maxn],*dp[maxn],*p,ans[maxn];
16 void dfs2(int now,int fa) {
17     dp[now][0]=1;
18     if(son[now]) {
19         dp[son[now]]=dp[now]+1;
20         dfs2(son[now],now);
21         ans[now]=ans[son[now]]+1;
22     }
23     for(int x:g[now]) {
24         int len=height[x]-dep[x]+1;
25         if(x==fa || x==son[now]) continue;
26         dp[x]=p;
27         p+=len;
28         dfs2(x,now);
29         for(int j=0;j<len;j++) {
30             dp[now][j+1]+=dp[x][j];
31             if(dp[now][j+1]>dp[now][ans[now]] || (dp[now][j+1]==dp[now][ans[now]]&&ans[
32                 now>j+1)) {
33                 ans[now]=j+1;
34             }
35         }
36         if(dp[now][ans[now]]==1) ans[now]=0;
37     }
38 int main()
39 {
40     ios::sync_with_stdio(false);cin.tie(0);
41     int n;
42     cin>>n;
43     rep(i,1,n-1) {
44         int u,v;
45         cin>>u>>v;
46         g[u].pb(v);
47         g[v].pb(u);
48     }
49     dfs1(1,0);
50     p=tmp;
51     dp[1]=p;
52     p+=height[1];
53     dfs2(1,0);
54     rep(i,1,n) cout<<ans[i]<<'\\n';
55     return 0;
56 }

```

## 2.8 有向图欧拉回路

```

1 // 有向图具有欧拉回路的条件:
2 // 1.所有点出度等于入度
3 // 2.基图连通
4 vector<pair<int,bool>> g[maxn];
5 vector<int>ans;// need reverse
6 void dfs(int now) {
7     for(auto &x:g[now]) {
8         if(x.se) continue;
9         x.se=true;
10        dfs(x.fi);
11    }
12    ans.pb(now);
13 }

```

## 2.9 一般图匹配

```

1 int n,m;
2 // match[i], i所对应的匹配点
3 int match[maxn],pre[maxn],vis[maxn],fa[maxn],tim[maxn],idx,ans;
4 vi g[maxn];
5 queue<int>Q;
6 int find(int x){return x==fa[x]?x:fa[x]=find(fa[x]);}
7 int lca(int x,int y) {
8     for (++idx;;swap(x,y))
9         if (x) {
10            x=find(x);
11            if (tim[x]==idx) return x;
12            else tim[x]=idx,x=pre[match[x]];
13        }
14 }
15 void blossom(int x,int y,int p) {
16     while (find(x)!=p) {
17         pre[x]=y;y=match[x];
18         if (vis[y]==2) vis[y]=1,Q.push(y);
19         if (find(x)==x) fa[x]=p;
20         if (find(y)==y) fa[y]=p;
21         x=pre[y];
22     }
23 }
24 int Aug(int S) {
25     for (int i=1;i<=n;++i)
26         vis[i]=pre[i]=0,fa[i]=i;
27     while (!Q.empty()) Q.pop();
28     Q.push(S);vis[S]=1;
29     while (!Q.empty()) {
30         int u=Q.front();Q.pop();
31         for(int v:g[u]) {
32             if (find(u)==find(v)||vis[v]==2) continue;
33             if (!vis[v]) {
34                 vis[v]=2;pre[v]=u;

```

```

35         if (!match[v]) {
36             for (int x=v,lst;x;x=lst)
37                 lst=match[pre[x]],match[x]=pre[x],match[pre[x]]=x;
38             return 1;
39         }
40         vis[match[v]]=1,Q.push(match[v]);
41     }
42     else {
43         int gg=lca(u,v);
44         blossom(u,v,gg);blossom(v,u,gg);
45     }
46 }
47 }
48 return 0;
49 }
50 int MaxMatch() {
51     int ans=0;
52     rep(i,1,n) {
53         if(!match[i]) ans+=Aug(i);
54     }
55     return ans;
56 }

```

## 2.10 网络流

### 2.10.1 最大流 (Dinic)

```

1 struct E {
2     int to, cp;
3     E(int to, int cp): to(to), cp(cp) {}
4 };
5
6 struct Dinic {
7     static const int M = 1E5 * 5;
8     int m, s, t;
9     vector<E> edges;
10    vector<int> G[M];
11    int d[M];
12    int cur[M];
13    void init(int n, int s, int t) {
14        this->s = s; this->t = t;
15        for (int i = 0; i <= n; i++) G[i].clear();
16        edges.clear(); m = 0;
17    }
18    void addedge(int u, int v, int cap) {
19        edges.emplace_back(v, cap);
20        edges.emplace_back(u, 0);
21        G[u].push_back(m++);
22        G[v].push_back(m++);
23    }
24    bool BFS() {
25        memset(d, 0, sizeof d);

```



```

26     queue<int> Q;
27     Q.push(s); d[s] = 1;
28     while (!Q.empty()) {
29         int x = Q.front(); Q.pop();
30         for (int& i: G[x]) {
31             E &e = edges[i];
32             if (!d[e.to] && e.cp > 0) {
33                 d[e.to] = d[x] + 1;
34                 Q.push(e.to);
35             }
36         }
37     }
38     return d[t];
39 }
40 int DFS(int u, int cp) {
41     if (u == t || !cp) return cp;
42     int tmp = cp, f;
43     for (int& i = cur[u]; i < G[u].size(); i++) {
44         E& e = edges[G[u][i]];
45         if (d[u] + 1 == d[e.to]) {
46             f = DFS(e.to, min(cp, e.cp));
47             e.cp -= f;
48             edges[G[u][i] ^ 1].cp += f;
49             cp -= f;
50             if (!cp) break;
51         }
52     }
53     return tmp - cp;
54 }
55 int go() {
56     int flow = 0;
57     while (BFS()) {
58         memset(cur, 0, sizeof cur);
59         flow += DFS(s, inf);
60     }
61     return flow;
62 }
63 } DC;

```

### 2.10.2 上下界网络流

```

1 struct E {
2     int to, cp, id;
3     E(int to, int cp, int id): to(to), cp(cp), id(id) {}
4 };
5
6 struct Dinic {
7     static const int M = 1E5 * 5;
8     int m, s, t;
9     vector<E> edges;
10    vector<int> G[M];
11    int d[M];

```

```

12     int cur[M];
13     void change(int _s,int _t) {
14         s=_s,t=_t;
15     }
16     void init(int n, int s, int t) {
17         this->s = s; this->t = t;
18         for (int i = 0; i <= n; i++) G[i].clear();
19         edges.clear(); m = 0;
20     }
21     void addedge(int u, int v, int cap,int id) {
22         edges.emplace_back(v, cap,id);
23         edges.emplace_back(u, 0,0);
24         G[u].push_back(m++);
25         G[v].push_back(m++);
26     }
27     bool BFS() {
28         memset(d, 0, sizeof d);
29         queue<int> Q;
30         Q.push(s); d[s] = 1;
31         while (!Q.empty()) {
32             int x = Q.front(); Q.pop();
33             for (int& i: G[x]) {
34                 E &e = edges[i];
35                 if (!d[e.to] && e.cp > 0) {
36                     d[e.to] = d[x] + 1;
37                     Q.push(e.to);
38                 }
39             }
40         }
41         return d[t];
42     }
43     int DFS(int u, int cp) {
44         if (u == t || !cp) return cp;
45         int tmp = cp, f;
46         for (int& i = cur[u]; i < G[u].size(); i++) {
47             E& e = edges[G[u][i]];
48             if (d[u] + 1 == d[e.to]) {
49                 f = DFS(e.to, min(cp, e.cp));
50                 e.cp -= f;
51                 edges[G[u][i] ^ 1].cp += f;
52                 cp -= f;
53                 if (!cp) break;
54             }
55         }
56         return tmp - cp;
57     }
58     int go() {
59         int flow = 0;
60         while (BFS()) {
61             memset(cur, 0, sizeof cur);
62             flow += DFS(s, inf);
63         }
64         return flow;

```

```

65     }
66 } DC;
67 int flow[maxn];
68 // 有源汇最大流
69 int main()
70 {
71     ios::sync_with_stdio(false); cin.tie(nullptr);
72     int n,m,s,t;
73     cin>>n>>m>>s>>t;
74     DC.init(n+3,n+1,n+2);
75     DC.addedge(t,s,inf,0);
76     rep(i,1,m) {
77         int u,v,down,up;
78         cin>>u>>v>>down>>up;
79         DC.addedge(u,v,up-down,i);
80         flow[u]-=down;
81         flow[v]+=down;
82     }
83     rep(i,1,n) {
84         if(flow[i]>0) DC.addedge(n+1,i,flow[i],0);
85         if(flow[i]<0) DC.addedge(i,n+2,-flow[i],0);
86     }
87
88     DC.go();
89     bool f=true;
90     for(int x:DC.G[n+1]) {
91         if(DC.edges[x].cp>0) {
92             f=false;
93             break;
94         }
95     }
96     if(f) {
97         int ans=0;
98         for(int x:DC.G[t]) {
99             if(DC.edges[x].id==0&&DC.edges[x].to!=s) {
100                 ans+=DC.edges[x].cp;
101             }
102         }
103         rep(i,1,n+2) {
104             for(int x:DC.G[i]) {
105                 if(DC.edges[x].id==0) {
106                     DC.edges[x].cp=0;
107                 }
108             }
109         }
110         DC.change(s,t);
111         ans+=DC.go();
112         cout<<ans<<"\n";
113     }
114     else cout<<"please go home to sleep"<<"\n";
115     return 0;
116 }

```

## 2.10.3 费用流

```

1 struct E {
2     int from, to, cp, v;
3     E() {}
4     E(int f, int t, int cp, int v) : from(f), to(t), cp(cp), v(v) {}
5 };
6 struct MCMF {
7     int n, m, s, t;
8     vector<E> edges;
9     vector<int> G[M];
10    bool inq[M];
11    int d[M], p[M], a[M];
12    void init(int _n, int _s, int _t) {
13        n = _n; s = _s; t = _t;
14        rep(i, 0, n) G[i].clear();
15        edges.clear(); m = 0;
16    }
17    void addedge(int from, int to, int cap, int cost) {
18        edges.emplace_back(from, to, cap, cost);
19        edges.emplace_back(to, from, 0, -cost);
20        G[from].push_back(m++);
21        G[to].push_back(m++);
22    }
23    bool BellmanFord(int &flow, int &cost) {
24        rep(i, 0, n) d[i] = inf;
25        memset(inq, 0, sizeof inq);
26        d[s] = 0, a[s] = inf, inq[s] = true;
27        queue<int> Q; Q.push(s);
28        while (!Q.empty()) {
29            int u = Q.front(); Q.pop();
30            inq[u] = false;
31            for (int& idx: G[u]) {
32                E &e = edges[idx];
33                if (e.cp && d[e.to] > d[u] + e.v) {
34                    d[e.to] = d[u] + e.v;
35                    p[e.to] = idx;
36                    a[e.to] = min(a[u], e.cp);
37                    if (!inq[e.to]) {
38                        Q.push(e.to);
39                        inq[e.to] = true;
40                    }
41                }
42            }
43        }
44        if (d[t] == inf) return false;
45        flow += a[t];
46        cost += a[t] * d[t];
47        int u = t;
48        while (u != s) {
49            edges[p[u]].cp -= a[t];
50            edges[p[u] ^ 1].cp += a[t];
51            u = edges[p[u]].from;

```

```

52     }
53     return true;
54 }
55 int go() {
56     int flow = 0, cost = 0;
57     while (BellmanFord(flow, cost));
58     return cost;
59 }
60 } MM;

```

## 2.11 二分图

### 2.11.1 匈牙利

```

1  int book[505][505], k, m, n;
2  int cx[505], cy[505], vis[505];
3  void init() {
4      memset(cx, -1, sizeof(cx));
5      memset(cy, -1, sizeof(cy));
6      memset(vis, 0, sizeof(vis));
7      memset(book, 0, sizeof(book));
8  }
9  int line(int x) {
10     int i;
11     for (i = 1; i <= n; i++) {
12         if (book[x][i] && !vis[i]) {
13             vis[i] = 1;
14             if (cy[i] == -1 || line(cy[i])) {
15                 cx[x] = i;
16                 cy[i] = x;
17                 return 1;
18             }
19         }
20     }
21     return 0;
22 }
23 int maxmatch() {
24     int sum = 0, i;
25     for (i = 1; i <= m; i++) {
26         if (cx[i] == -1) {
27             memset(vis, 0, sizeof(vis));
28             sum += line(i);
29         }
30     }
31     return sum;
32 }

```

### 2.11.2 HK

```

1  namespace HK {
2      vector<int> G[maxn];
3      int uN;

```

```
4   int Mx[maxn],My[maxn];
5   int dx[maxn],dy[maxn];
6   int dis;
7   bool used[maxn];
8   bool searchp() {
9       queue<int> q;
10      dis=inf;
11      memset(dx,-1,sizeof(dx));
12      memset(dy,-1,sizeof(dy));
13      for(int i=0;i<uN;i++) {
14          if(Mx[i]==-1) {
15              q.push(i);
16              dx[i]=0;
17          }
18      }
19      while(!q.empty()) {
20          int u=q.front();q.pop();
21          if(dx[u]>dis) break;
22          for(int v:G[u]) {
23              if(dy[v]==-1) {
24                  dy[v]=dx[u]+1;
25                  if(My[v]==-1) dis=dy[v];
26                  else {
27                      dx[My[v]]=dy[v]+1;
28                      q.push(My[v]);
29                  }
30              }
31          }
32      }
33      return dis!=inf;
34  }
35  bool DFS(int u) {
36      for(int v:G[u]) {
37          if(!used[v]&&dy[v]==dx[u]+1) {
38              used[v]=true;
39              if(My[v]!=-1&&dy[v]==dis) continue;
40              if(My[v]==-1||DFS(My[v])) {
41                  My[v]=u;
42                  Mx[u]=v;
43                  return true;
44              }
45          }
46      }
47      return false;
48  }
49  int MaxMatch() {
50      int res=0;
51      memset(Mx,-1,sizeof(Mx));
52      memset(My,-1,sizeof(My));
53      while(searchp()) {
54          memset(used,false,sizeof(used));
55          for(int i=0;i<uN;i++) {
56              if(Mx[i]==-1&&DFS(i)) res++;
```

```

57     }
58 }
59 return res;
60 }
61
62 }

```

### 2.11.3 KM(DFS)

```

1  int nx, ny; //两边的点的数目
2  int g[maxn][maxn], linker[maxn]; //linker[i]:右边第i个点的匹配对象
3  int lx[maxn], ly[maxn], slack[maxn];
4  bool visx[maxn], visy[maxn]; //x,y是否被访问
5  bool dfs(int x) {
6      visx[x] = true;
7      for (int y = 0; y < ny; y++) {
8          if (visy[y]) continue;
9          int tmp = lx[x] + ly[y] - g[x][y];
10         if (tmp == 0) {
11             visy[y] = true;
12             if (linker[y] == -1 || dfs(linker[y])) {
13                 linker[y] = x;
14                 return true;
15             }
16         }
17         else if (slack[y] > tmp) slack[y] = tmp;
18     }
19     return false;
20 }
21 int km() {
22     memset(linker, -1, sizeof(linker));
23     memset(ly, 0, sizeof(ly));
24     for (int i = 0; i < nx; i++) {
25         lx[i] = -inf;
26         for (int j = 0; j < ny; j++) {
27             if (g[i][j] > lx[i]) lx[i] = g[i][j];
28         }
29     }
30     for (int x = 0; x < nx; x++) {
31         for (int i = 0; i < ny; i++) slack[i] = inf;
32         while (1) {
33             memset(visx, false, sizeof(visx));
34             memset(visy, false, sizeof(visy));
35             if (dfs(x)) break;
36             int d = inf;
37             for (int i = 0; i < ny; i++) {
38                 if (!visy[i] && d > slack[i]) d = slack[i];
39             }
40             for (int i = 0; i < nx; i++) {
41                 if (visx[i]) lx[i] -= d;
42             }
43             for (int i = 0; i < ny; i++) {

```

```

44         if (visy[i]) ly[i] += d;
45         else slack[i] -= d;
46     }
47 }
48 }
49 int res = 0;
50 for (int i = 0; i < ny; i++) {
51     if (linker[i] != -1) res += g[linker[i]][i];
52 }
53 return res;
54 }

```

#### 2.11.4 KM(BFS)

```

1  ll w[maxn][maxn];
2  ll lx[maxn],ly[maxn];
3  int linker[maxn]; //linker[i]:右边第i个点的匹配对象
4  ll slack[maxn];
5  int n;
6  bool visy[maxn];
7  int pre[maxn];
8  void bfs(int k) {
9      int x,y=0,yy=0,delta;
10     memset(pre,0,sizeof(pre));
11     for(int i=1;i<=n;i++) slack[i]=inf;
12     linker[y]=k;
13     while(1) {
14         x=linker[y];delta=inf;visy[y]=true;
15         for(int i=1;i<=n;i++) {
16             if(!visy[i]) {
17                 if(slack[i]>lx[x]+ly[i]-w[x][i]) {
18                     slack[i]=lx[x]+ly[i]-w[x][i];
19                     pre[i]=y;
20                 }
21                 if(slack[i]<delta) delta=slack[i],yy=i;
22             }
23         }
24         for(int i=0;i<=n;i++) {
25             if(visy[i])lx[linker[i]]-=delta,ly[i]+=delta;
26             else slack[i]-=delta;
27         }
28         y=yy;
29         if(linker[y]==-1) break;
30     }
31     while(y) linker[y]=linker[pre[y]],y=pre[y];
32 }
33 ll km() {
34     memset(lx,0,sizeof(lx));
35     memset(ly,0,sizeof(ly));
36     memset(linker,-1,sizeof(linker));
37     for(int i=1;i<=n;i++) {
38         memset(visy,false,sizeof(visy));

```



```

39     bfs(i);
40 }
41 ll ans=0;
42 for(int i=1;i<=n;i++) ans+=w[linker[i]][i];
43 return ans;
44 }

```

## 3 字符串

### 3.1 kmp

```

1  int Next[maxn];
2  void getNext(string a) {
3      int len = a.size(), k = -1, i = 0;
4      Next[0] = -1;
5      while (i < len)
6      {
7          while (k != -1 && a[i] != a[k]) k = Next[k];
8          Next[++i] = ++k;
9      }
10 }
11 int kmp(string a1,string a2) {
12     int i, j = 0;
13     for (i = 0; i < a1.size(); i++) {
14         while (j != -1 && a1[i] != a2[j]) j = Next[j];
15         j++;
16         if (j == a2.size()) return i - a2.size() + 2;
17     }
18     return -1;
19 }

```

### 3.2 exkmp

```

1  // next[i]: x[i...m-1]与x[0...m-1]的最长公共前缀
2  // extend[i]: y[i...n-1]与x[0...m-1]的最长公共前缀
3  void pre_exkmp(char x[],int m,int next[]) {
4      next[0]=m;
5      int j=0;
6      while(j+1<m&&x[j]==x[j+1]) j++;
7      next[1]=j;
8      int k=1;
9      for(int i=2;i<m;i++) {
10         int p=next[k]+k-1,L=next[i-k];
11         if(i+L<p+1) next[i]=L;
12         else {
13             j=max(0,p-i+1);
14             while(i+j<m&&x[i+j]==x[j]) j++;
15             next[i]=j;
16             k=i;
17         }
18     }

```

```

19 }
20 void EKMP(char x[],int m,char y[],int n,int next[],int extend[]) {
21     pre_exkmp(x,m,next);
22     int j=0;
23     while(j<n&& j<m&& x[j]==y[j]) j++;
24     extend[0]=j;
25     int k=0;
26     for(int i=1;i<n;i++) {
27         int p=extend[k]+k-1,L=next[i-k];
28         if(i+L<p+1) extend[i]=L;
29         else {
30             j=max(0,p-i+1);
31             while(i+j<n&& j<m&& y[i+j]==x[j]) j++;
32             extend[i]=j;
33             k=i;
34         }
35     }
36 }

```

### 3.3 最小表示法

```

1 int minpos(string a) {
2     int k=0,i=0,j=1,n=a.size();
3     while(k<n&& i<n&& j<n) {
4         if(a[(i+k)%n]==a[(j+k)%n]) k++;
5         else {
6             if(a[(i+k)%n]>a[(j+k)%n]) i=i+k+1;
7             else j=j+k+1;
8             if(i==j) i++;
9             k=0;
10        }
11    }
12    i=min(i,j);
13 }
14 string trans(string a,int pos) {
15     string ans;
16     rep(i,0,a.size()-1) ans+=a[(i+pos)%a.size()];
17     return ans;
18 }

```

### 3.4 AC 自动机

```

1 struct ACAM {
2     int next[maxn][26],fail[maxn],id[maxn],root,cnt;
3     int match[maxn];
4     int newnode() {
5         rep(i,0,25) next[cnt][i]=-1;
6         id[cnt++]=0;
7         return cnt-1;
8     }
9     void init() {

```

```

10     cnt=0;
11     root=newnode();
12 }
13 void ins(char a[],int x) {
14     int len=strlen(a),now=root;
15     rep(i,0,len-1) {
16         if(next[now][a[i]-'a']==-1)
17             next[now][a[i]-'a']=newnode();
18         now=next[now][a[i]-'a'];
19     }
20     id[now]=x;
21 }
22 void build() {
23     queue<int> q;
24     fail[root]=root;
25     rep(i,0,25) {
26         if(next[root][i]==-1) {
27             next[root][i]=root;
28         }
29         else {
30             fail[next[root][i]]=root;
31             match[next[root][i]]=root;
32             q.push(next[root][i]);
33         }
34     }
35     while(!q.empty()) {
36         int now=q.front();
37         q.pop();
38         rep(i,0,25) {
39             if(next[now][i]==-1) next[now][i]=next[fail[now]][i];
40             else {
41                 fail[next[now][i]]=next[fail[now]][i];
42                 q.push(next[now][i]);
43                 int tmp=fail[next[now][i]];
44                 if(id[tmp]>0) match[next[now][i]]=tmp;
45                 else match[next[now][i]]=match[tmp];
46             }
47         }
48     }
49 }
50 }AC;

```

### 3.5 SA

```

1 struct SuffixArray
2 {
3     int sa[maxn], rank[maxn], ws[maxn], wv[maxn], wa[maxn], wb[maxn], height[maxn],
        st[maxbit][maxn], N;
4     bool cmp(int *r, int a, int b, int l){return r[a]==r[b] and r[a+l]==r[b+l];}
5     void build(int *r, int n, int m)
6     {
7         N=n;

```

```

8      n++;
9      int i, j, k=0, p, *x=wa, *y=wb, *t;
10     for(i=0;i<m;i++)ws[i]=0;
11     for(i=0;i<n;i++)ws[x[i]=r[i]]++;
12     for(i=1;i<m;i++)ws[i]+=ws[i-1];
13     for(i=n-1;i>=0;i--)sa[--ws[x[i]]]=i;
14     for(p=j=1;p<n;j<=1,m=p)
15     {
16         for(p=0,i=n-j;i<n;i++)y[p++]=i;
17         for(i=0;i<n;i++)if(sa[i]>=j)y[p++]=sa[i]-j;
18         for(i=0;i<n;i++)wv[i]=x[y[i]];
19         for(i=0;i<m;i++)ws[i]=0;
20         for(i=0;i<n;i++)ws[wv[i]]++;
21         for(i=1;i<m;i++)ws[i]+=ws[i-1];
22         for(i=n-1;i>=0;i--)sa[--ws[wv[i]]]=y[i];
23         for(t=x,x=y,y=t,p=1,i=1,x[sa[0]]=0;i<n;i++)
24             x[sa[i]]=cmp(y,sa[i-1],sa[i],j)?p-1:p++;
25     }
26
27     for(i=0;i<n;i++)rank[sa[i]]=i;
28
29     for(i=0;i<n-1;height[rank[i+1]]=k)
30         for(k?k--:0,j=sa[rank[i]-1];r[i+k]==r[j+k];k++);
31 }
32 void build_st() //st表
33 {
34     int i, k;
35     for(i=1;i<=N;i++)st[0][i]=height[i];
36     for(k=1;k<=maxbit;k++)
37         for(i=1;i+(1<<k)-1<=N;i++)
38             st[k][i]=min(st[k-1][i],st[k-1][i+(1<<k)-1]);
39 }
40 int lcp(int x, int y) //最长公共前缀
41 {
42     int l=rank[x], r=rank[y];
43     if(l>r)swap(l,r);
44     if(l==r)return N-sa[l];
45     int t=log2(r-l);
46     return min(st[t][l+1],st[t][r-(1<<t)+1]);
47 }
48 }SA;

```

### 3.6 SAM

```

1 namespace sam {
2     const int M = maxn << 1;
3     int t[M][26], len[M] = {-1}, fa[M], sz = 2, last = 1;
4     ll endpos[M];
5     void init() { memset(t, 0, (sz + 10) * sizeof t[0]); sz = 2; last = 1; }
6     void ins(int ch) {
7         int p = last, np = last = sz++; endpos[np]=1;
8         len[np] = len[p] + 1;

```

```

9      for (; p && !t[p][ch]; p = fa[p]) t[p][ch] = np;
10     if (!p) { fa[np] = 1; return; }
11     int q = t[p][ch];
12     if (len[p] + 1 == len[q]) fa[np] = q;
13     else {
14         int nq = sz++; len[nq] = len[p] + 1;
15         memcpy(t[nq], t[q], sizeof t[0]);
16         fa[nq] = fa[q];
17         fa[np] = fa[q] = nq;
18         for (; t[p][ch] == q; p = fa[p]) t[p][ch] = nq;
19     }
20 }
21 //拓扑排序
22 int c[M] = {1}, a[M];
23 void rsort() {
24     rep (i, 1, sz-1) c[i] = 0;
25     rep (i, 1, sz-1) c[len[i]]++;
26     rep (i, 1, sz-1) c[i] += c[i - 1];
27     rep (i, 1, sz-1) a[--c[len[i]]] = i;
28 }
29 //计算endpos大小
30 void calendpos() {
31     rsort();
32     per(i,sz-1,2) {
33         endpos[fa[a[i]]]+=endpos[a[i]];
34     }
35 }
36 }

```

### 3.7 manacher

```

1  #include <bits/stdc++.h>
2  using namespace std;
3  typedef long long ll;
4  typedef unsigned long long ull;
5  #define met(s) memset(s, 0, sizeof(s))
6  const int MOD = 100000007;
7  const int MAXN = 1000000;
8  const int inf = 0x3f3f3f3f;
9  char str_new[220010], str[110010];
10 int p[220025];
11 int init() {
12     int len = strlen(str);
13     str_new[0] = '$';
14     str_new[1] = '#';
15     int t = 2;
16     for (int i = 0; i < len; i++) {
17         str_new[t++] = str[i];
18         str_new[t++] = '#';
19     }
20     return t;
21 }

```

```

22 int manacher() {
23     memset(str_new, 0, sizeof(str_new));
24     int len = init();
25     int id, mx = 0, Max = 0;
26     for (int i = 0; i < len; i++) {
27         if (i < mx) p[i] = min(p[2 * id - i], mx - i);
28         else p[i] = 1;
29         while (str_new[i - p[i]] == str_new[i + p[i]]) p[i]++;
30         if (mx < i + p[i]) {
31             mx = i + p[i];
32             id = i;
33         }
34         Max = max(Max, p[i] - 1);
35     }
36     for(int i=0;i<len;i++) p[i]--;
37     return Max;
38 }
39 int main() {
40     while (scanf("%s", str) != EOF) {
41         int ans = manacher();
42         printf("%d\n", ans );
43     }
44     return 0;
45 }

```

### 3.8 PAM

```

1 // cnt[i]: 表示该节点在原串中出现的次数....记得count()
2 namespace pam {
3     const int N=maxn<<1;
4     int t[N][26], fa[N], len[N], rs[N], cnt[N], num[N];
5     int sz, n, last;
6     int _new(int l) {
7         len[sz] = l; cnt[sz] = num[sz] = 0;
8         return sz++;
9     }
10    void init() {
11        memset(t, 0, sz * sizeof t[0]);
12        rs[n = sz = 0] = -1;
13        last = _new(0);
14        fa[last] = _new(-1);
15    }
16    int get_fa(int x) {
17        while (rs[n - 1 - len[x]] != rs[n]) x = fa[x];
18        return x;
19    }
20    void ins(int ch) {
21        rs[++n] = ch;
22        int p = get_fa(last);
23        if (!t[p][ch]) {
24            int np = _new(len[p] + 2);
25            num[np] = num[fa[np]] = t[get_fa(fa[p])][ch] + 1;

```

```

26         t[p][ch] = np;
27     }
28     ++cnt[last = t[p][ch]];
29     // len[last]; 以当前字符结尾的最长回文后缀
30     // num[last]; 以当前字符结尾的回文串数量(回文后缀)
31 }
32 void count() {
33     for(int i=sz-1;i>=2;i--)
34         cnt[fa[i]]+=cnt[i];
35 }
36 }

```

### 3.9 lyndon 分解

```

1 vector<string> duval(string const& s) {
2     int n = s.size(), i = 0;
3     vector<string> factorization;
4     while (i < n) {
5         int k = i + 1, j = i;
6         while (k < n && s[j] <= s[k]) {
7             if (s[j] < s[k])
8                 j = i;
9             else
10                 j++;
11             k++;
12         }
13         while (i <= j) {
14             factorization.push_back(s.substr(i, k - j));
15             i += k - j;
16         }
17     }
18     return factorization;
19 }

```

## 4 数学

### 4.1 素数表

```

1 int prime[maxn], p[maxn];
2 void init(int up) {
3     rep(i, 1, up) p[i] = i;
4     rep(i, 2, up) {
5         if (p[i] == i) prime[++prime[0]] = i;
6         for (int j = 1; j <= prime[0] && i <= up / prime[j]; j++) {
7             p[i * prime[j]] = prime[j];
8             if (i % prime[j] == 0) break;
9         }
10    }
11 }
12 vector<pii> smallfc(int x) {
13     vector<pii> ans;

```

```

14     while(x!=1) {
15         pii now;
16         now.fi=p[x];
17         while(x%now.fi==0) {
18             now.se++;
19             x/=now.fi;
20         }
21         ans.pb(now);
22     }
23     reverse(ans.begin(),ans.end());
24     return ans;
25 }
26 vector<pair<ll,int>> bigfc(ll x) {
27     vector<pair<ll,int>> ans;
28     rep(i,1,prime[0]) {
29         if(prime[i]>x/prime[i]) break;
30         if(x%prime[i]!=0) continue;
31         int cnt=0;
32         while(x%prime[i]==0) {
33             cnt++;
34             x/=prime[i];
35         }
36         ans.pb({prime[i],cnt});
37     }
38     return ans;
39 }

```

## 4.2 欧拉函数

```

1  //phi[1]=1, 1与1是互质的
2  int phi[maxn];
3  void euler(int n) {
4      memset(phi,0,sizeof(phi));
5      phi[1]=1;
6      for(int i=2;i<=n;i++) {
7          if(!phi[i]) {
8              for(int j=i;j<=n;j++) {
9                  if(!phi[j]) phi[j]=j;
10                 phi[j]=phi[j]/i*(i-1);
11             }
12         }
13     }
14 }

```

## 4.3 PollardRho

```

1  int t,s=20,cnt;
2  long long fac[1001];
3  long long ksc(long long x,long long y,long long mod){
4      long long res=0;
5      while(y) {

```



```
6     if(y&1)
7         res=(res+x)%mod;
8         x=(x<<1)%mod;
9         y>>=1;
10    }
11    return res;
12 }
13 long long ksm(long long x,long long y,long long mod) {
14     long long res=1;
15     while(y) {
16         if(y&1)
17             res=ksc(res,x,mod);
18             x=ksc(x,x,mod);
19             y>>=1;
20     }
21     return res;
22 }
23 int miller_rabin(long long n) {
24     if(n==2)
25         return 1;
26     if(n<2||!(n%2))
27         return 0;
28     long long u,pre,x;
29     int k=0;
30     u=n-1;
31     while(!(u&1)) {
32         ++k;
33         u>>=1;
34     }
35     for(int i=1;i<=s;++i) {
36         x=rand()%(n-2)+2;
37         x=ksm(x,u,n);
38         pre=x;
39         for(int j=1;j<=k;++j) {
40             x=ksc(x,x,n);
41             if(x==1&&pre!=1&&pre!=n-1)
42                 return 0;
43             pre=x;
44         }
45         if(x!=1)
46             return 0;
47     }
48     return 1;
49 }
50 long long gcd(long long a,long long b){//注意与一般的gcd不一样
51     if (a==0) return 1;//pollard_rho的需要
52     if (a<0) return gcd(-a,b);//可能有负数
53     while (b){
54         long long t=a%b; a=b; b=t;
55     }
56     return a;
57 }
58 long long pollard_rho(long long n,long long c) { //找因子
```

```

59
60     long long i=1,k=2;//用来判断是否形成了环
61     long long xx=rand()%n,y=xx;
62     while(1) {
63         i++;
64         xx=(ksc(xx,xx,n)+c)%n;
65         long long d=gcd(y-xx,n);
66         if(1<d&&d<n)//找到一个因数
67             return d;
68         if(y==xx)//出现循环,那么这次寻找失败
69             return n;
70         if(i==k){//相当于每次找连续k这么多次取模有没有得到相同余数
71             y=xx;
72             k<<=1;
73         }
74     }
75 }
76 void find(long long n) { //通过找因数来找质因子
77     if(miller_rabin(n)) {
78         fac[++cnt]=n;//记录质因子
79         return;
80     }
81     long long p=n;
82     while(p>=n)
83         p=pollard_rho(p,rand()%(n-1)+1);//如果转了一圈还是p,则继续while循环
84     //p是当前找到的一个因数(不一定是质因数),再分别找p和n/p的质因数
85     find(p);
86     find(n/p);
87 }
88 int main()
89 {
90     srand(time(0)+19260817);
91     scanf("%d",&t);
92     while(t-->0)
93     {
94         long long x;
95         scanf("%lld",&x);
96         if(miller_rabin(x))
97         {
98             printf("Prime\n");
99             continue;
100         }
101         cnt=0;
102         find(x);
103         sort(fac+1,fac+cnt+1);
104         for(int i=1;i<=cnt;i++) cout<<fac[i]<<" ";
105     }
106     return 0;
107 }

```

#### 4.4 线性求逆元

```
1 inv[1]=1;for(int i=2;i<maxn;i++)inv[i]=inv[mod%i]*(mod-mod/i)%mod;
```

## 4.5 组合数学

```
1 ll jc[maxn],inv[maxn];
2 void cinit(int n) {
3     jc[0]=inv[0]=1;
4     rep(i,1,n) jc[i]=jc[i-1]*i%mod;
5     inv[n]=qpow(jc[n],mod-2);
6     per(i,n-1,1) inv[i]=inv[i+1]*(i+1)%mod;
7 }
8 ll C(int n,int m,int p=mod) { return m>n?0:jc[n]*inv[m]%p*inv[n-m]%p; }
9 ll lucas(ll n,ll m,ll p=mod){ return m==0?1:lucas(n/p,m/p,p)*C(n%p,m%p,p)%p; }
```

## 4.6 FFT

```
1 const double PI=acos(-1.0);
2 struct Complex {
3     double x,y;
4     Complex(double _x=0.0,double _y=0.0){ x=_x,y=_y; }
5     Complex operator -(const Complex &b)const { return Complex(x-b.x,y-b.y); }
6     Complex operator +(const Complex &b)const { return Complex(x+b.x,y+b.y); }
7     Complex operator *(const Complex &b)const { return Complex(x*b.x-y*b.y,x*b.y+y*b
        .x); }
8 };
9 void change(Complex y[],int len) {
10     int i,j,k;
11     for(i=1,j=len/2;i<len-1;i++) {
12         if(i<j) swap(y[i],y[j]);
13         k=len/2;
14         while(j>=k) {
15             j-=k;
16             k/=2;
17         }
18         if(j<k) j+=k;
19     }
20 }
21 void fft(Complex y[],int len,int on) {
22     change(y,len);
23     for(int h=2;h<=len;h<=1) {
24         Complex wn(cos(-on*2*PI/h),sin(-on*2*PI/h));
25         for(int j=0;j<len;j+=h) {
26             Complex w(1,0);
27             for(int k=j;k<j+h/2;k++) {
28                 Complex u=y[k];
29                 Complex t=w*y[k+h/2];
30                 y[k]=u+t;
31                 y[k+h/2]=u-t;
32                 w=w*wn;
33             }
34         }
35     }
```

```

35     }
36     if(on==-1) {
37         for(int i=0;i<len;i++) y[i].x/=len;
38     }
39 }
40 void conv(Complex a[],Complex b[],int n) {
41     fft(a,n,1);
42     fft(b,n,1);
43     for(int i=0;i<n;i++) a[i]=a[i]*b[i];
44     fft(a,n,-1);
45 }

```

## 4.7 NTT

```

1  ll qpow(ll a,int b) {
2      ll ans=1;
3      while(b) {
4          if(b&1) ans=ans*a%mod;
5          a=a*a%mod;
6          b>>=1;
7      }
8      return ans;
9  }
10 const int G=3;
11 namespace NTT {
12     ll wn[maxn << 2], rev[maxn << 2];
13     int init(int n_) {
14         int step = 0; int n = 1;
15         for ( ; n < n_; n <= 1) ++step;
16         rep(i,1,n-1)
17             rev[i] = (rev[i >> 1] >> 1) | ((i & 1) << (step - 1));
18         int g = qpow(G, (mod - 1) / n);
19         wn[0] = 1;
20         for (int i = 1; i <= n; ++i)
21             wn[i] = wn[i - 1] * g % mod;
22         return n;
23     }
24     void NTT(ll a[], int n, int f) {
25         rep(i, 0, n-1)if(i < rev[i])
26             std::swap(a[i], a[rev[i]]);
27         for (int k = 1; k < n; k <= 1) {
28             for (int i = 0; i < n; i += (k << 1)) {
29                 int t = n / (k << 1);
30                 rep(j, 0, k-1) {
31                     ll w = f == 1 ? wn[t * j] : wn[n - t * j];
32                     ll x = a[i + j];
33                     ll y = a[i + j + k] * w % mod;
34                     a[i + j] = (x + y) % mod;
35                     a[i + j + k] = (x - y + mod) % mod;
36                 }
37             }
38         }

```

```

39     if (f == -1) {
40         ll ninv = qpow(n, mod-2);
41         rep (i, 0, n-1) a[i] = a[i] * ninv % mod;
42     }
43 }
44 ll tmp1[maxn<<2],tmp2[maxn<<2];
45 void mul(ll a[],ll b[],int lena,int lenb) { // lena为a的项数
46     int len=init(lena+lenb+2);
47     rep(i,0,lena-1) tmp1[i]=a[i];
48     rep(i,lena,len-1) tmp1[i]=0;
49     rep(i,0,lenb-1) tmp2[i]=b[i];
50     rep(i,lenb,len-1) tmp2[i]=0;
51     NTT(tmp1,len,1); NTT(tmp2,len,1);
52     rep(i,0,len-1) a[i]=tmp1[i]*tmp2[i]%mod;
53     NTT(a,len,-1);
54 }
55 void mul(ll a[],ll b[],ll c[],int lena,int lenb) {
56     int len=init(lena+lenb+2);
57     rep(i,0,lena-1) tmp1[i]=a[i];
58     rep(i,lena,len-1) tmp1[i]=0;
59     rep(i,0,lenb-1) tmp2[i]=b[i];
60     rep(i,lenb,len-1) tmp2[i]=0;
61     NTT(tmp1,len,1); NTT(tmp2,len,1);
62     rep(i,0,len-1) c[i]=tmp1[i]*tmp2[i]%mod;
63     NTT(c,len,-1);
64 }
65 ll P[20][maxn<<2];
66 void solve(ll *a,int l,int r,int dep) { //(1+a[0]*x)(1+a[1]*x)(1+a[2]*x)
67     if(l==r) {
68         P[dep][2*l]=1;
69         P[dep][2*l+1]=a[l];
70         return ;
71     }
72     int mid=(l+r)>>1;
73     solve(a,l,mid,dep+1);
74     solve(a,mid+1,r,dep+1);
75     int lenl=mid-l+1,lenr=r-mid;
76     mul(P[dep+1]+2*l,P[dep+1]+2*mid+2,P[dep]+1*2,lenl+1,lenr+1);
77 }
78 ll tmp[maxn << 2];
79 void inv(ll *a,ll *b,int len) { //lena为a的项数
80     if(len==1) { b[0]=qpow(a[0],mod-2); return ;}
81     inv(a,b,(len+1)>>1);
82     int l=init(len*2+1);
83     rep(i,0,len-1) tmp[i]=a[i];
84     rep(i,len,l-1) tmp[i]=0;
85     NTT(tmp,l,1);
86     NTT(b,l,1);
87     rep(i,0,l-1) b[i]=(2-b[i]*tmp[i]%mod+mod)%mod*b[i]%mod;
88     NTT(b,l,-1);
89     rep(i,len,l-1) b[i]=0;
90 }
91 ll pa[maxn<<2],pb[maxn<<2],pb_inv[maxn<<2];

```

```

92 void div(ll *a,ll *b,int lena,int lenb,ll *Q,ll *R) {//lena为a的最高次数
93     rep(i,0,lena) pa[i]=a[lena-i];
94     rep(i,0,lenb) pb[i]=b[lenb-i];
95     rep(i,lena-lenb+1,lenb) pb[i]=0;
96     inv(pb,pb_inv,lena-lenb+1);
97     mul(pa,pb_inv,lena+1,lena-lenb+1);
98     rep(i,0,lena-lenb) Q[i]=pa[lena-lenb-i];
99     rep(i,0,lenb) pb[i]=b[i];
100    mul(pb,Q,lenb+1,lena-lenb+1);
101    rep(i,0,lenb-1) R[i]=(a[i]-pb[i]+mod)%mod;
102 }
103 ll invb[maxn<<2],D[maxn<<2],inv2=qpow(2,mod-2);
104 void sqrt(ll *a,ll *b,int len) {//要保证b[0]=1,否则需要二次剩余求
105     if(len==1) {b[0]=1;return ;}
106     sqrt(a,b,(len+1)>>1);
107     rep(i,0,len<<1) invb[i]=0;
108     inv(b,invb,len);
109     int l=init(2*len);
110     rep(i,0,len-1) D[i]=a[i];
111     rep(i,len,l-1) D[i]=0;
112     NTT(D,l,1);NTT(b,l,1);NTT(invb,l,1);
113     rep(i,0,l-1) b[i]=(b[i]+D[i]*invb[i]%mod)%mod*inv2%mod;
114     NTT(b,l,-1);
115     rep(i,len,l-1) b[i]=0;
116 }
117 void dao(ll *a,ll *b,int len) {
118     rep(i,0,len-1) b[i-1]=i*a[i]%mod;b[len-1]=0;
119 }
120 void jifen(ll *a,ll *b,int len) {
121     rep(i,0,len-1) b[i]=a[i-1]*qpow(i,mod-2)%mod;b[0]=0;
122 }
123 ll A[maxn<<2],B[maxn<<2];
124 void ln(ll *a,ll *b,int len) {
125     int n;for(n=1;n<=len;n<<=1);
126     rep(i,0,n<<1) A[i]=B[i]=0;
127     dao(a,A,n);inv(a,B,n);
128     int l=init(n*2);
129     NTT(A,l,1),NTT(B,l,1);
130     rep(i,0,l-1) A[i]=A[i]*B[i]%mod;
131     NTT(A,l,-1);jifen(A,b,n);
132 }
133 ll F[maxn<<2];
134 void exp(ll *a,ll *b,int len) {//保证a[0]=0;
135     if(len==1) { b[0]=1;return ;}
136     exp(a,b,(len+1)>>1);
137     ln(b,F,len);
138     int l=init(2*len+1);
139     rep(i,0,len-1) F[i]=(a[i]-F[i]+mod)%mod;
140     rep(i,len,l-1) F[i]=b[i]=0;
141     F[0]++;
142     NTT(F,l,1),NTT(b,l,1);
143     rep(i,0,l-1) b[i]=b[i]*F[i]%mod;
144     NTT(b,l,-1);

```

```

145     rep(i,len,l-1) b[i]=0;
146 }
147 }

```

## 4.8 FWT

```

1  template<typename T>
2  void fwt(ll a[], int n, T f) {
3      for (int d = 1; d < n; d *= 2)
4          for (int i = 0, t = d * 2; i < n; i += t)
5              rep (j, 0, d-1)
6                  f(a[i + j], a[i + j + d]);
7  }
8
9  void AND(ll& a, ll& b) { a =(a+b)%mod; }
10 void OR(ll& a, ll& b) { b =(a+b)%mod; }
11 void XOR (ll& a, ll& b) {
12     ll x = a, y = b;
13     a = (x + y) % mod;
14     b = (x - y + mod) % mod;
15 }
16 void rAND(ll& a, ll& b) { a =(a-b+mod)%mod; }
17 void rOR(ll& a, ll& b) { b =(b-a+mod)%mod; }
18 void rXOR(ll& a, ll& b) {
19     static ll INV2 = (mod + 1) / 2;
20     ll x = a, y = b;
21     a = (x + y) * INV2 % mod;
22     b = (x - y + mod) * INV2 % mod;
23 }

```

## 4.9 第二类斯特林数

```

1  ll sb[maxn<<2],sr[maxn<<2];
2  // 计算  $k! \cdot S(n,k), k=0,1,2,\dots,n$ , 复杂度  $(n \cdot \log n)$ 
3  void stirling_second(int n){
4      int len=NTT::init(2*n+2);
5      for(int i=1; i<=n; i++) sr[i]=qpow(i, n)*inv[i]%mod;
6      for(int i=0; i<=n; i++){
7          if(i&1) sb[i]=mod-inv[i];
8          else sb[i]=inv[i];
9      }
10     NTT::mul(sr,sb,len);
11 }
12 // 单独计算  $m! \cdot S(n,m)$  复杂度  $O(m \cdot \log n)$ 
13 ll calsr(ll n,ll m) {
14     ll ans=0;
15     rep(k,0,m) {
16         ll now=C(m,k)*qpow(m-k,n)%mod;
17         ans=(k&1)?(ans-now+mod)%mod:(ans+now)%mod;
18     }
19     return ans;

```

20 }

#### 4.10 BSGS

```

1 ll BSGS(ll a, ll b, ll p) { // a^x = b (mod p)
2     a %= p;
3     if (!a && !b) return 1;
4     if (!a) return -1;
5     static map<ll, ll> mp; mp.clear();
6     ll m = sqrt(p + 1.5);
7     ll v = 1;
8     rep (i, 1, m) {
9         v = v * a % p;
10        mp[v * b % p] = i;
11    }
12    ll vv = v;
13    rep (i, 1, m) {
14        auto it = mp.find(vv);
15        if (it != mp.end()) return i * m - it->second;
16        vv = vv * v % p;
17    }
18    return -1;
19 }

```

#### 4.11 高斯消元

```

1 typedef double LD;
2 const LD eps = 1E-10;
3 const int maxn = 2000 + 10;
4
5 int n, m;
6 LD a[maxn][maxn], x[maxn];
7 bool free_x[maxn];
8
9 inline int sgn(LD x) { return (x > eps) - (x < -eps); }
10
11 int gauss(LD a[maxn][maxn], int n, int m) {
12     memset(free_x, 1, sizeof free_x); memset(x, 0, sizeof x);
13     int r = 0, c = 0;
14     while (r < n && c < m) {
15         int m_r = r;
16         rep (i, r + 1, n-1)
17             if (fabs(a[i][c]) > fabs(a[m_r][c])) m_r = i;
18         if (m_r != r)
19             rep (j, c, m)
20                 swap(a[r][j], a[m_r][j]);
21         if (!sgn(a[r][c])) {
22             a[r][c] = 0;
23             ++c;
24             continue;
25         }

```



```

26     rep (i, r + 1, n-1)
27         if (a[i][c]) {
28             LD t = a[i][c] / a[r][c];
29             rep (j, c, m + 1) a[i][j] -= a[r][j] * t;
30         }
31     ++r; ++c;
32 }
33 rep (i, r, n-1)
34     if (sgn(a[i][m])) return -1;
35 if (r < m) {
36     per (i, r - 1, 0) {
37         int f_cnt = 0, k = -1;
38         rep (j, 0, m-1)
39             if (sgn(a[i][j]) && free_x[j]) {
40                 ++f_cnt;
41                 k = j;
42             }
43         if(f_cnt > 0) continue;
44         LD s = a[i][m];
45         rep (j, 0, m-1)
46             if (j != k) s -= a[i][j] * x[j];
47         x[k] = s / a[i][k];
48         free_x[k] = 0;
49     }
50     return m - r;
51 }
52 per(i, m - 1, 0) {
53     LD s = a[i][m];
54     rep (j, i + 1, m-1)
55         s -= a[i][j] * x[j];
56     x[i] = s / a[i][i];
57 }
58 return 0;
59 }

```

#### 4.12 BM 线性递推

```

1 namespace linear_seq
2 {
3     #define SZ(x) ((int)x.size())
4     using VI=vector<int>;
5     const int N=10010;
6     ll res[N],base[N],_c[N],_md[N];
7     vector<int> Md;
8     void mul(ll *a,ll *b,int k) {
9         for(int i = 0 ; i < k + k ; ++i)
10             _c[i]=0;
11         for(int i = 0 ; i < k ; ++i)
12             if (a[i])
13                 for(int j = 0 ; j < k ; ++ j)
14                     _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;
15         for (int i=k+k-1;i>=k;i--)

```

```

16         if (_c[i])
17             for(int j = 0 ; j<(int ) Md.size() ; ++ j)
18                 _c[i-k+Md[j]]=( _c[i-k+Md[j]]-_c[i]*_md[Md[j]])%mod;
19     for(int i =0 ; i< k ; ++i)
20         a[i]=_c[i];
21 }
22 int solve(ll n,VI a,VI b) {
23     ll ans=0,pnt=0;
24     int k=SZ(a);
25     assert( SZ(a) == SZ(b) );
26     for(int i = 0 ; i < k ; ++ i)
27         _md[k-1-i] = -a[i] ; _md[k] = 1 ;
28     Md.clear() ;
29     for(int i =0 ; i < k ; ++ i)
30         if (_md[i]!=0)
31             Md.push_back(i);
32     for(int i = 0; i< k ;++ i)
33         res[i]=base[i]=0;
34     res[0]=1;
35     while ((1ll<<pnt)<=n)
36         pnt++;
37     for (int p=pnt;p>=0;p--) {
38         mul(res,res,k);
39         if ((n>>p)&1) {
40             for (int i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
41             for(int j = 0 ; j < (int)Md.size() ; ++ j)
42                 res[ Md[j] ]=(res[ Md[j] ]-res[k]*_md[Md[j]])%mod;
43         }
44     }
45     rep(i,0,k-1) ans=(ans+res[i]*b[i])%mod;
46     if (ans<0) ans+=mod;
47     return ans;
48 }
49 VI BM(VI s) {
50     VI C(1,1),B(1,1);
51     int L=0,m=1,b=1;
52     for(int n= 0 ; n < (int)s.size(); ++ n ) {
53         ll d=0;
54         for(int i =0 ; i < L +1 ;++ i)
55             d=(d+(1ll)C[i]*s[n-i])%mod;
56         if (d==0) ++m;
57         else if (2*L<=n) {
58             VI T=C;
59             ll c=mod-d*qpow(b,mod-2)%mod;
60             while (SZ(C)<SZ(B)+m)
61                 C.push_back(0);
62             for(int i =0 ; i < (int)B.size(); ++ i)
63                 C[i+m]=(C[i+m]+c*B[i])%mod;
64             L=n+1-L; B=T; b=d; m=1;
65         } else {
66             ll c=mod-d*qpow(b,mod-2)%mod;
67             while (SZ(C)<SZ(B)+m)
68                 C.push_back(0);

```

```

69         for(int i = 0 ; i < (int) B.size() ; ++ i)
70             C[i+m]=(C[i+m]+c*B[i])%mod;
71         ++m;
72     }
73 }
74 return C;
75 }
76 ll gao(VI a,ll n) {
77     VI c=BM(a);
78     c.erase(c.begin());
79     for( int i = 0 ; i < (int)c.size( );++i )
80         c[i]=(mod-c[i])%mod;
81     return (ll)solve(n,c,VI(a.begin(),a.begin()+SZ(c)));
82 }
83 };

```

### 4.13 O(1) 快速乘

```

1 inline ll mul(ll a,ll b,ll mod){
2     return (a*b-(ll)((long double)a/mod*b)*mod+mod)%mod;
3 }

```

## 5 其他

### 5.1 vimrc

```

1 set nu sw=4 ts=4 bs=2 ai smd
2 syntax on
3 set cindent
4 set noswapfile
5 set nobackup
6 set mouse =a
7 set showmatch
8 set autowrite
9 imap kj <esc>l
10 nmap <c-a> ggVG"+y
11 imap {<CR> {<CR><Esc>O
12 imap {<s-cr> {<CR>
13 map<F6> :call CR()<CR>
14 func! CR()
15     exec "w"
16     exec "!g++ % -std=c++11 -o %<.exe"
17     exec "! %<.exe "
18 endfunc

```