# 02-router

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This lab took me about 10 hours to do.

Implementation Explanation:

#### ARP报文处理

接收到的ARP报文分为请求报文和响应。对于请求报文,当其请求的目的IP为本端口地址时,回复一个ARP响应报文,并且把源的IP和MAC映射关系加入ARP缓存。对于响应报文,当其回复的目的地址为本端口地址时,说明发出的ARP请求得到了回复,此时把源端口的IP和MAC映射关系加入ARP缓存。

实现代码较长,在此不表

#### ARP表管理

首先是在ARP表中查找IP地址对应的MAC地址,我们遍历所有表项,如果有表项的IP和我们要查询的IP 相等并且表项有效,则选取该表项

```
int arpcache_lookup(u32 ip4, u8 mac[ETH_ALEN])
{
    pthread_mutex_lock(&arpcache.lock);

    for(int i=0;i<MAX_ARP_SIZE;i++)
    {
        if(arpcache.entries[i].ip4==ip4&&arpcache.entries[i].valid)
        {
            memcpy(mac, arpcache.entries[i].mac,ETH_ALEN);
            pthread_mutex_unlock(&arpcache.lock);
            return 1;
        }
    }
    pthread_mutex_unlock(&arpcache.lock);
    return 0;
}</pre>
```

如果在ARP表中查找不到相应条目,我们将该分组挂到ARPCache中的一个等待队列里,然后发出相应的ARP请求报文,直到我们收到这个IP地址对应的MAC地址,再释放这个分组

```
void arpcache_append_packet(iface_info_t *iface,u32 ip4,char *packet,int len)
{
        struct arp_req *req_entry=NULL, *req_q;
        struct cached_pkt *recv_pkt;
        recv_pkt=malloc(sizeof(struct cached_pkt));
        recv_pkt->len=len;
        recv_pkt->packet=packet;
        pthread_mutex_lock(&arpcache.lock);
        list_for_each_entry_safe(req_entry, req_q, &(arpcache.req_list), list)
        {
                if(req_entry->iface==iface&&req_entry->ip4==ip4)
                {
                        list_add_tail(&(recv_pkt->list),&(req_entry->cached_packets));
                        pthread_mutex_unlock(&arpcache.lock);
                        return;
                }
        }
        struct arp_req *added_req_list=(struct arp_req*)malloc(sizeof(struct arp_req));
        added_req_list->iface=iface;
        added_req_list->ip4=ip4;
        added_req_list->sent=time(NULL);
        added_req_list->retries=1;
        init_list_head(&(added_req_list->cached_packets));
        list_add_tail(&(recv_pkt -> list),&(added_req_list->cached_packets));
        list_add_tail(&(added_req_list->list),&(arpcache.req_list));
        arp_send_request(iface,ip4);
        pthread_mutex_unlock(&arpcache.lock);
}
```

第三是插入IP->MAC地址映射。我们插入映射时,需要找一个地方插入,如果有空闲表项(valid=0), 我们插入这个空闲表项,如果没有我们随机替换一个。所以我们首先需要遍历整个表寻找空闲表项,找 不到再随机替换。插入后需要释放正在等待该地址的分组,因为我们已经查到了他们所需要的MAC地址

```
void arpcache_insert(u32 ip4, u8 mac[ETH_ALEN])
{
        pthread_mutex_lock(&arpcache.lock);
        // Find the entry.
        int pos=-1;
        for(int i=0;i<MAX_ARP_SIZE;i++)</pre>
        {
                if(!arpcache.entries[i].valid)
                {
                         pos=i;
                         arpcache.entries[i].added=time(NULL);
                         arpcache.entries[i].ip4=ip4;
                         memcpy(arpcache.entries[i].mac, mac, ETH_ALEN);
                         arpcache.entries[i].valid=1;
                         break;
                }
        }
        if(pos==-1)
        {
                pos=time(NULL)%32;
                arpcache.entries[pos].added=time(NULL);
                arpcache.entries[pos].ip4=ip4;
                memcpy(arpcache.entries[pos].mac, mac, ETH_ALEN);
                arpcache.entries[pos].valid=1;
        }
        \ensuremath{//} Delete all the pending packets.
        struct arp_req *req_entry=NULL, *req_q;
        list_for_each_entry_safe(req_entry,req_q,&(arpcache.req_list),list)
        {
                if(req_entry->ip4==ip4)
                {
                         struct cached_pkt *pkt_entry=NULL,*pkt_q;
                         list_for_each_entry_safe(pkt_entry,pkt_q,&(req_entry->cached_pac
                         {
                                 struct ether_header *eth_hdr=(struct ether_header*)(pkt_
                                 memcpy(eth_hdr->ether_dhost, mac, ETH_ALEN);
                                 iface_send_packet(req_entry->iface,pkt_entry->packet,pkt
                                 list_delete_entry(&(pkt_entry->list));
```

```
free(pkt_entry);
}

list_delete_entry(&(req_entry->list));
free(req_entry);
}

pthread_mutex_unlock(&arpcache.lock);
}
```

最后是清理表项和分组缓存。这个主要有两部分,第一部分是清除存在时间过长的表项,第二部分是清 除请求次数过多的分组

```
void *arpcache_sweep(void *arg)
{
        while(1)
        {
                sleep(1);
                pthread_mutex_lock(&(arpcache.lock));
                struct arp_req *req_entry=NULL, *req_q;
                time_t now=time(NULL);
                for (int i=0;i<MAX_ARP_SIZE;i++)</pre>
                {
                         if(arpcache.entries[i].valid&&now-arpcache.entries[i].added>ARP_
                                 arpcache.entries[i].valid=0;
                }
                list_for_each_entry_safe(req_entry, req_q, &(arpcache.req_list), list)
                {
                         if(req_entry->retries>ARP_REQUEST_MAX_RETRIES)
                         {
                                 struct cached_pkt *pkt_entry=NULL,*pkt_q;
                                 list_for_each_entry_safe(pkt_entry,pkt_q,&(req_entry->ca
                                 {
                                         pthread_mutex_unlock(&(arpcache.lock));
                                         icmp_send_packet(pkt_entry->packet,pkt_entry->le
                                         pthread_mutex_lock(&(arpcache.lock));
                                         free(pkt_entry);
                                 }
                                 list_delete_entry(&(req_entry->list));
                                 free(req_entry);
                                 continue;
                         }
                         if(now-req_entry->sent>=1)
                         {
                                 arp_send_request(req_entry->iface, req_entry->ip4);
                                 req_entry->sent=now;
                                 req_entry->retries++;
                         }
                }
                pthread_mutex_unlock(&(arpcache.lock));
        }
        return NULL;
```

### 路由表查询与数据报转发

遍历整个路由表,当地址匹配且mask大于当前最大mask时,更新目标表项和最大mask,另外还要考虑 默认路由的掩码为全0的情况

#### 实现:

```
rt_entry_t *longest_prefix_match(u32 dst)
{
        rt_entry_t *pos, *res=NULL;
        u32 max_mask=0;
        list_for_each_entry(pos,&rtable,list)
        {
                if((pos->mask&pos->dest)==(pos->mask&dst))
                {
                         if(pos->mask>max_mask)
                         {
                                 res=pos;
                                 max_mask=pos->mask;
                         }
                }
        }
        return res;
}
```

在得到对应的路由器表项后,我们查看其网关地址,如果为0,意味着目的主机在同一子网内,此时直接向目的IP转发数据报;否则向下一跳网关转发数据报

```
void ip_send_packet(char *packet, int len)
{
        struct ether_header *eh=(struct ether_header*)packet;
        struct iphdr *ih=packet_to_ip_hdr(packet);
        rt_entry_t *find_rt=longest_prefix_match(ntohl(ih->daddr));
        if(find_rt==NULL)
        {
                free(packet);
                return;
        }
        u32 next_ip;
        if (find_rt->gw)
                next_ip=find_rt->gw;
        else
                next_ip=ntohl(ih->daddr);
        iface_send_packet_by_arp(find_rt->iface, next_ip, packet, len);
}
```

我们收到ip数据报后,首先查看目的地址和目前的iface的addr是不是一样,如果一样说明这是ICMP报文的请求回复报文,我们返回相应的reply;如果不一样则在路由表查找下一跳ip地址,如果查不到,说明我们无法到达,返回icmp unreachable异常,查的到的话再处理ttl,如果还能转发的话就将ip数据报转发出去

```
void handle_ip_packet(iface_info_t *iface, char *packet, int len)
{
        struct iphdr *ip_hdr=packet_to_ip_hdr(packet);
        u32 daddr=ntohl(ip_hdr -> daddr);
        // ICMP packet
        if(daddr==iface->ip)
        {
                struct iphdr *ip_hdr=packet_to_ip_hdr(packet);
                struct icmphdr *icmp_hdr=(struct icmphdr*)IP_DATA(ip_hdr);
                if(icmp_hdr->type==ICMP_ECHOREQUEST)
                        icmp_send_packet(packet,len,ICMP_ECHOREPLY,0);
                else
                        free(packet);
                return;
        }
        // Search daddr in router table.
        rt_entry_t *p_rt=longest_prefix_match(daddr);
        if(p_rt==NULL)
        {
                icmp_send_packet(packet, len, ICMP_DEST_UNREACH, ICMP_NET_UNREACH);
                return;
        }
        // ttl
        ip_hdr->ttl--;
        if(ip_hdr->ttl<=0)</pre>
        {
                icmp_send_packet(packet,len,ICMP_TIME_EXCEEDED,ICMP_EXC_TTL);
                return;
        ip_hdr->checksum=ip_checksum(ip_hdr);
        // Get the next jump.
        u32 next_jump=p_rt->gw?p_rt->gw:daddr;
        // forward packet by arp protocol.
        iface_send_packet_by_arp(p_rt->iface, next_jump, packet, len);
}
```

# ICMP数据报发送

考虑到ICMP数据报的格式,我们首先要确定其长度,主要是区分ping数据报和其他类型数据报;确定长度后就开始填充ICMP报文

```
void icmp_send_packet(const char *in_pkt,int len,u8 type,u8 code)
        struct iphdr *in_ip_hdr=packet_to_ip_hdr(in_pkt);
    int pkt_len=0;
    if (type==ICMP_ECHOREPLY)
        pkt_len=len;
        else
        pkt_len=ETHER_HDR_SIZE+IP_BASE_HDR_SIZE+ICMP_HDR_SIZE+IP_HDR_SIZE(in_ip_hdr)+8;
    char *sent_pkt=(char*)malloc(pkt_len);
    struct ether_header *eh=(struct ether_header*)sent_pkt;
    struct iphdr *ip_hdr=packet_to_ip_hdr(sent_pkt);
    struct icmphdr *icmp_hdr=(struct icmphdr*)(sent_pkt+ETHER_HDR_SIZE+IP_BASE_HDR_SIZE)
    eh -> ether_type=htons(ETH_P_IP);
    rt_entry_t *entry=longest_prefix_match(ntohl(in_ip_hdr->saddr));
    ip_init_hdr(ip_hdr,entry->iface->ip,ntohl(in_ip_hdr->saddr),pkt_len-ETHER_HDR_SIZE,1
    icmp_hdr->code=code;
    icmp_hdr->type=type;
    if(type==0)
        {
        memcpy(sent_pkt+ETHER_HDR_SIZE+IP_HDR_SIZE(ip_hdr)+4, \
        in_pkt+ETHER_HDR_SIZE+IP_HDR_SIZE(in_ip_hdr)+4,pkt_len-(ETHER_HDR_SIZE+IP_HDR_S]
    }
        else
        memset(sent_pkt+ETHER_HDR_SIZE+IP_HDR_SIZE(ip_hdr)+4,0,4);
                memcpy(sent_pkt+ETHER_HDR_SIZE+IP_HDR_SIZE(ip_hdr)+4+4,\
        in_ip_hdr, IP_HDR_SIZE(in_ip_hdr)+8);
    }
    icmp_hdr->checksum=icmp_checksum(icmp_hdr,pkt_len-ETHER_HDR_SIZE-IP_BASE_HDR_SIZE);
        ip_send_packet(sent_pkt,pkt_len);
}
```

Screenshots:

Host连通性测试:

#### h1 ping h2:

```
mininet> r1 ./router &
DEBUG: find the following interfaces: r1-eth0 r1-eth1 r1-eth2.
mininet> h1 ping h2
PING 10.0.2.22 (10.0.2.22) 56(84) bytes of data.
64 bytes from 10.0.2.22: icmp_seq=1 ttl=63 time=1.81 ms
64 bytes from 10.0.2.22: icmp_seq=2 ttl=63 time=0.313 ms
64 bytes from 10.0.2.22: icmp_seq=3 ttl=63 time=0.415 ms
64 bytes from 10.0.2.22: icmp_seq=4 ttl=63 time=0.456 ms
64 bytes from 10.0.2.22: icmp_seq=5 ttl=63 time=0.717 ms
64 bytes from 10.0.2.22: icmp_seq=5 ttl=63 time=0.638 ms
^C
--- 10.0.2.22 ping statistics ---
6 packets transmitted, 6 received, 0% packet loss, time 5054ms
rtt min/avg/max/mdev = 0.313/0.724/1.806/0.502 ms
```

### h1 ping h3:

```
mininet> h1 ping h3
PING 10.0.3.33 (10.0.3.33) 56(84) bytes of data.
64 bytes from 10.0.3.33: icmp_seq=1 ttl=63 time=0.413 ms
64 bytes from 10.0.3.33: icmp_seq=2 ttl=63 time=0.387 ms
64 bytes from 10.0.3.33: icmp_seq=3 ttl=63 time=0.673 ms
64 bytes from 10.0.3.33: icmp_seq=4 ttl=63 time=0.106 ms
64 bytes from 10.0.3.33: icmp_seq=5 ttl=63 time=0.148 ms
64 bytes from 10.0.3.33: icmp_seq=5 ttl=63 time=0.066 ms
^C
--- 10.0.3.33 ping statistics ---
6 packets transmitted, 6 received, 0% packet loss, time 5117ms
rtt min/avg/max/mdev = 0.066/0.298/0.673/0.214 ms
```

# h2 ping h1:

```
mininet> h2 ping h1

PING 10.0.1.11 (10.0.1.11) 56(84) bytes of data.

64 bytes from 10.0.1.11: icmp_seq=1 ttl=63 time=0.064 ms

64 bytes from 10.0.1.11: icmp_seq=2 ttl=63 time=0.270 ms

64 bytes from 10.0.1.11: icmp_seq=3 ttl=63 time=0.144 ms

64 bytes from 10.0.1.11: icmp_seq=4 ttl=63 time=0.118 ms

64 bytes from 10.0.1.11: icmp_seq=5 ttl=63 time=0.065 ms

64 bytes from 10.0.1.11: icmp_seq=5 ttl=63 time=0.122 ms

^C

--- 10.0.1.11 ping statistics ---

6 packets transmitted, 6 received, 0% packet loss, time 5097ms

rtt min/avg/max/mdev = 0.064/0.130/0.270/0.069 ms
```

### h2 ping h3:

```
mininet> h2 ping h3

PING 10.0.3.33 (10.0.3.33) 56(84) bytes of data.

64 bytes from 10.0.3.33: icmp_seq=1 ttl=63 time=12.0 ms

64 bytes from 10.0.3.33: icmp_seq=2 ttl=63 time=0.311 ms

64 bytes from 10.0.3.33: icmp_seq=3 ttl=63 time=0.319 ms

64 bytes from 10.0.3.33: icmp_seq=4 ttl=63 time=0.067 ms

64 bytes from 10.0.3.33: icmp_seq=5 ttl=63 time=0.100 ms

64 bytes from 10.0.3.33: icmp_seq=5 ttl=63 time=0.100 ms

64 bytes from 10.0.3.33: icmp_seq=6 ttl=63 time=0.089 ms

^C

--- 10.0.3.33 ping statistics ---

6 packets transmitted, 6 received, 0% packet loss, time 5096ms

rtt min/avg/max/mdev = 0.067/2.148/12.005/4.409 ms
```

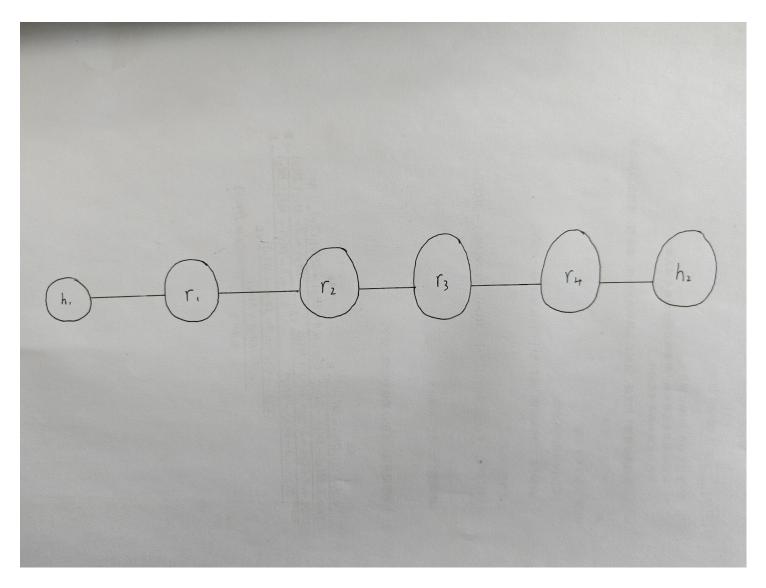
# h3 ping h1:

```
mininet> h3 ping h1
PING 10.0.1.11 (10.0.1.11) 56(84) bytes of data.
64 bytes from 10.0.1.11: icmp_seq=1 ttl=63 time=0.085 ms
64 bytes from 10.0.1.11: icmp_seq=2 ttl=63 time=0.207 ms
64 bytes from 10.0.1.11: icmp_seq=3 ttl=63 time=0.168 ms
64 bytes from 10.0.1.11: icmp_seq=4 ttl=63 time=0.106 ms
64 bytes from 10.0.1.11: icmp_seq=5 ttl=63 time=0.155 ms
64 bytes from 10.0.1.11: icmp_seq=5 ttl=63 time=0.221 ms
^C
--- 10.0.1.11 ping statistics ---
6 packets transmitted, 6 received, 0% packet loss, time 5179ms
rtt min/avg/max/mdev = 0.085/0.157/0.221/0.049 ms
```

### h3 ping h2:

```
mininet> h3 ping h2
PING 10.0.2.22 (10.0.2.22) 56(84) bytes of data.
64 bytes from 10.0.2.22: icmp_seq=1 ttl=63 time=0.272 ms
64 bytes from 10.0.2.22: icmp_seq=2 ttl=63 time=0.151 ms
64 bytes from 10.0.2.22: icmp_seq=3 ttl=63 time=0.135 ms
64 bytes from 10.0.2.22: icmp_seq=4 ttl=63 time=0.146 ms
64 bytes from 10.0.2.22: icmp_seq=5 ttl=63 time=0.146 ms
64 bytes from 10.0.2.22: icmp_seq=5 ttl=63 time=0.131 ms
64 bytes from 10.0.2.22: icmp_seq=6 ttl=63 time=0.131 ms
65 cc
67 cc
67 packets transmitted, 6 received, 0% packet loss, time 5110ms
67 rtt min/avg/max/mdev = 0.131/0.163/0.272/0.049 ms
```

#### 手动构建的拓扑网络:



```
mininet> dump

<Host h1: h1-eth0:10.0.0.1 pid=32174>

<Host h2: h2-eth0:10.0.0.2 pid=32176>

<Host r1: r1-eth0:10.0.0.3,r1-eth1:None pid=32178>

<Host r2: r2-eth0:10.0.0.4,r2-eth1:None pid=32180>

<Host r3: r3-eth0:10.0.0.5,r3-eth1:None pid=32182>

<Host r4: r4-eth0:10.0.0.6,r4-eth1:None pid=32184>
```

# 连通性测试:

```
mininet> h1 ping h2
PING 10.0.5.22 (10.0.5.22) 56(84) bytes of data.
64 bytes from 10.0.5.22: icmp_seq=1 ttl=60 time=2.97 ms
64 bytes from 10.0.5.22: icmp_seq=2 ttl=60 time=2.45 ms
64 bytes from 10.0.5.22: icmp_seq=3 ttl=60 time=2.03 ms
64 bytes from 10.0.5.22: icmp_seq=4 ttl=60 time=3.10 ms
64 bytes from 10.0.5.22: icmp_seq=5 ttl=60 time=2.35 ms
^C
--- 10.0.5.22 ping statistics ---
5 packets transmitted, 5 received, 0% packet loss, time 4007ms
rtt min/avg/max/mdev = 2.032/2.579/3.098/0.397 ms
```

```
mininet> h2 ping h1
PING 10.0.1.11 (10.0.1.11) 56(84) bytes of data.
64 bytes from 10.0.1.11: icmp_seq=1 ttl=60 time=0.961 ms
64 bytes from 10.0.1.11: icmp_seq=2 ttl=60 time=0.437 ms
64 bytes from 10.0.1.11: icmp_seq=3 ttl=60 time=1.12 ms
64 bytes from 10.0.1.11: icmp_seq=4 ttl=60 time=1.86 ms
64 bytes from 10.0.1.11: icmp_seq=5 ttl=60 time=1.86 ms
64 bytes from 10.0.1.11: icmp_seq=5 ttl=60 time=1.12 ms
^C
--- 10.0.1.11 ping statistics ---
5 packets transmitted, 5 received, 0% packet loss, time 4035ms
rtt min/avg/max/mdev = 0.437/1.099/1.859/0.455 ms
```

### traceout测试:

```
mininet> h1 traceroute h2
traceroute to 10.0.5.22 (10.0.5.22), 30 hops max, 60 byte packets
1 10.0.1.1 (10.0.1.1) 2.501 ms 0.078 ms 0.010 ms
2 10.0.2.2 (10.0.2.2) 0.071 ms 0.015 ms 0.115 ms
3 10.0.3.2 (10.0.3.2) 0.355 ms 0.345 ms 0.336 ms
4 10.0.4.2 (10.0.4.2) 0.548 ms 0.634 ms 0.516 ms
5 10.0.5.22 (10.0.5.22) 0.506 ms 0.496 ms 0.141 ms
```

# Remaining Bugs:

受制于能力,目前暂未发现