

Projekt: Kryptografie mit elliptischen Kurven

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Pseudocodes zu tatsächlicher Implementierung

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Algorithm 1: Modular Addition

input : integer a, integer b, natural number module
output: modular added values $[a + b \bmod \text{module}]$
1 return $((a \bmod \text{module}) + (b \bmod \text{module})) \bmod \text{module}$

Algorithm 2: Modular Subtraction

input : integer a, integer b, natural number module
output: modular subtracted values $[a - b \bmod \text{module}]$
1 return $((a \bmod \text{module}) - (b \bmod \text{module})) \bmod \text{module}$

Algorithm 3: Modular Multiplication

input : integer a, integer b, natural number module
output: modular multiplied values $[(a * b) \bmod \text{module}]$
1 $result \leftarrow 0$
2 for $i \leftarrow 0$ **to** $b - 1$ **do**
3 $result = \text{mod_add}(result, a, \text{module})$
4 $result = result \bmod \text{module}$
5 return $result$

Algorithm 4: Extended Euklidian Algororithm (recursive)

input : integer a, natural number b
output: Touple d,s,t ,result of the internally calculated table

```
1 if  $b = 0$  then  
2   return  $Touple(a, 1, 0)$  ; // return value without further compuattion  
  
3  $d1, s1, t1 \leftarrow \text{extended\_euklidian\_algorithm\_rec}(b, a \bmod b)$   
4  $d \leftarrow d1$   
5  $s \leftarrow t1$   
6  $t \leftarrow s1 - (a/b) * t1$   
7 return  $d, s, t$ 
```

Algorithm 5: Extended Euklidian Algorithm Wrapper for recursive function

input : integer a, natural number b
output: modular inverse of a in GF_2^b

```
1  $result \leftarrow \text{extended\_euklidian\_algorithm\_rec}(a, b)$   
2 return  $s$  from  $result$   $Touple \bmod b$ 
```

Algorithm 6: Square-And-Multiply

input : inetger base, int exp, natural number module
output: $base^{exp} \bmod \text{module}$

```
1  $exp\_bin \leftarrow$  exponent in binary form ; // binary exponent without righest 1  
2  $result \leftarrow base$  ; // initial value  
3 foreach  $integer$  of  $exp\_bin$  do  
4    $result \leftarrow result^2 \bmod \text{module}$   
5   if  $elem = 1$  then  
6      $result \leftarrow result * base \bmod \text{module}$   
7 return  $result$ 
```

Algorithm 7: Fermats Little Theorem

input : integer a, natural prime number p
output: modular inverse of a in GF_2^p

```
1 return  $sqaure\_and\_multiply(a, p - 2, p)$   
; // This function makes us of the frequently introduced function  
; //  $sqaure\_and\_multiply$  - fast exponentation is needed
```

Algorithm 8: Double Point

input : Point P, elliptic curve E, inversion mode
output: Point Q = P + P

- 1 $s \leftarrow inversion(2*y_P, mod_E)_{mode} * (3*(x_P)^2 + a_E \bmod mod_E)$
- 2 $x_3 \leftarrow s^2 - x_P - x_P \bmod mod_E$
- 3 $y_3 \leftarrow s * (x_P - x_3) - y_P \bmod mod_E$
- 4 **return** $Point(x_3, y_3)$
- 5

Algorithm 9: Add points

input : Point P, Point Q, elliptic curve E, inversion mode
output: Point R = P + Q

- 1 $s \leftarrow ((y_Q - y_P) \bmod mod_E * inversion(x_Q - x_P, mod_E)_{mode}) \bmod mod_E$
- 2 $x_3 \leftarrow s^2 - x_P - x_Q \bmod mod_E$
- 3 $y_3 \leftarrow s * (x_P - x_3) - y_P \bmod mod_E$
- 4 **return** $Point(x_3, y_3)$
- 5

Algorithm 10: Multiply Points

input : natural skalar factor k, Point P, elliptic curve E, inversion mode
output: T = k * P

- 1 $bin_exp \leftarrow$ binary representation of k without first 1
- 2 Point T $\leftarrow P$
- 3 **foreach** *integer of bin_exp* **do**
- 4 $T \leftarrow double_point(T, E, mode_{inversion})$
- 5 **if** *elem = 1* **then**
- 6 $T \leftarrow add_points(T, P, E, mode_{inversion})$
- 7 **return** T

Algorithm 11: NAF Point Multiplication

input : natural skalar factor k , Point P , elliptic curve E , inversion mode

output: Point $Q = k * P$

```
1  $Q \leftarrow PointP(0,0)$ 
2  $naf\_exp \leftarrow NAF(k)$ 
3 foreach integer of naf_exp do
4    $Q \leftarrow multiply\_points(2, Q, E, mode_{inversion})$ 
5   if  $elem = 1$  then
6      $Q \leftarrow add\_points(Q, P, E, mode_{inversion})$ 
7   if  $elem = -1$  then
8      $Q \leftarrow add\_points(Q, P^{-1}, E, mode_{inversion})$ 
9 return  $Q$ 
```

Algorithm 12: Calculate AES Key (res= 0x9a19f8e811c45299cb1e6625562f8505)

input : T_{AB}

output: AES key (base 16)

```
1  $T\_AB\_bin \leftarrow$  binary representation of  $T_{AB}$ 
2  $T\_AB\_bin\_1 \leftarrow$  bits  $b_0$  to  $b_{128}$  of  $T\_AB\_bin$ 
3  $T\_AB\_bin\_2 \leftarrow$  bits  $b_{128}$  to  $b_{256}$  of  $T\_AB\_bin$ 
4 return  $(T\_AB\_bin\_1 \oplus T\_AB\_bin\_2)_{16}$ 
```
