

1. $xs.map(id) = xs$

Base case: $xs = []$

Prove: $[] .map(id) = []$

LHS $= [] .map(id)$
 $= []$ (map definition)

RHS: $= [] = LHS$

Induction case: $xs = x:xs'$

IH: Assume $xs' .map(id) = xs'$ is true for any xs' .

LHS $= (x:xs') .map(id)$
 $= x:xs'$ (map definition)

RHS $= x:xs'$

LHS $= RHS$

Done

2. $(xs.append(ys)).map(f) = (xs.map(f)).append(ys.map(f))$

Base case: $xs = []$

Prove: $([] .append(ys)) .map(f) = ([] .map(f)) .append(ys.map(f))$

LHS $= ([] .append(ys)) .map(f)$
 $= ys$ (append, map definition)

RHS $= ([] .map(f)) .append(ys.map(f))$
 $= [] .append(ys.map(f))$
 $= [] .append(ys)$ (append, map definition)
 $= ys = LHS$

Induction case: $xs = x:xs'$

IH: assume $(xs' .append(ys)) .map(f) = (xs' .map(f)) .append(ys.map(f))$ is true for any xs' .

LHS $= (x:xs') .append(ys) .map(f)$
 $= x:(xs' .append(ys)) .map(f)$
 $= x:(xs' .append(ys) .map(f))$ (append, map definition)

RHS $= (x:xs') .map(f) .append(ys.map(f))$
 $= x:(xs.append(ys) .map(f))$ (append, map definition)

LHS $= RHS$

Done

3. $(xs.append(ys)).fold(g, a) = xs.fold(g, ys.fold(g, a))$

Base case: $xs = []$

Prove: $([].append(ys)).fold(g, a) = [].fold(g, ys.fold(g, a))$

$$\begin{aligned} \text{LHS} &= ([].append(ys)).fold(g, a) && \text{(append definition)} \\ &= ys.fold(g, a) && \text{(fold definition)} \\ \text{RHS} &= [].fold(g, ys.fold(g, a)) \\ &= ys.fold(g, a) && \text{(fold definition)} \\ &= ys = \text{LHS} \end{aligned}$$

Induction case: $xs = x:xs'$

IH: assume $(xs'.append(ys)).map(f) = (xs'.map(f)).append(ys.map(f))$ is true for any xy .

$$\begin{aligned} \text{LHS} &= ((x:xs').append(ys)).map(f) \\ &= x:(xs'.append(ys)).map(f) \\ &= x:(xs'.append(ys).map(f)) && \text{(append, map definition)} \\ \text{RHS} &= ((x:xs').map(f)).append(ys.map(f)) \\ &= x:(xs'.append(ys).map(f)) && \text{(append, map definition)} \\ \text{LHS} &= \text{RHS} \end{aligned}$$

Done

4. $(xs.append(ys)).length() = xs.length() + ys.length()$

Base case: $xs = []$

Prove: $([].append(ys)).length() = [].length() + ys.length()$

$$\begin{aligned} \text{LHS} &= ([].append(ys)).length() \\ &= ys.length() && \text{(append, length definition)} \\ \text{RHS} &= [].length() + ys.length() \\ &= ys.length() = \text{LHS} && \text{(length definition)} \end{aligned}$$

Induction: $xs = x:xs'$

IH: Assume $(xs'.append(ys)).length() = xs'.length() + ys.length()$ is true for any xy .

$$\begin{aligned} \text{LHS} &= (x:xs').append(ys).length() \\ &= x:(xs'.append(ys).length()) && \text{(append, length definition)} \\ &= 1 + xs'.length() + ys.length() \\ \text{RHS} &= (x:xs').length() + ys.length() \\ &= x:(xs'.append(ys)) + ys.length() && \text{(append, length definition)} \\ &= 1 + xs'.length() + ys.length() \\ \text{LHS} &= \text{RHS} \end{aligned}$$

Done

5. $(xs.reverseH\ ys).length() = xs.length() + ys.length()$

Base case: $xs = []$

Prove: $([].reverseH(ys)).length() = [].length() + ys.length()$

$$\begin{aligned} \text{LHS} &= ([].reverseH(ys)).length() \\ &= ys.length() && (\text{reverseH, length definition}) \\ \text{RHS} &= [].length() + ys.length() \\ &= ys.length() = \text{LHS} && (\text{length definition}) \end{aligned}$$

Induction: $xs = x:xs'$

IH: Assume $(xy.reverseH\ ys).length() = xy.length() + ys.length()$ is true for any xy .

$$\begin{aligned} \text{LHS} &= ((x:xs').reverseH(ys)).length() \\ &= xs'.length(x:ys).length() && (\text{reverseH definition}) \\ &= xs'.length() + (x:ys).length() \\ &= 1 + xs'.length() + ys.length() && (\text{length definition}) \\ \text{RHS} &= (x:xs').length() + ys.length() \\ &= 1 + xs'.length() + ys.length() && (\text{length definition}) \\ \text{LHS} &= \text{RHS} \end{aligned}$$

Done

6. $xs.length() = (xs.reverse2()).length()$

Prove: $xs.length() = (xs.reverse2()).length()$

$$\begin{aligned} \text{LHS} &= xs.length() \\ \\ \text{RHS} &= (xs.reverse2()).length() \\ &= xs.reverseH([]).length() && (\text{reverse2 definition}) \\ &= xs.length() + [].length() && (\text{reverseH definition}) \\ &= xs.length() = \text{LHS} && (\text{length definition}) \\ \text{LHS} &= \text{RHS} \end{aligned}$$

Done

7. $(xs.append(ys)).reverse() = ys.reverse().append(xs.reverse())$

Base case: $xs = []$

Prove: $([].append(ys)).reverse() = ys.reverse().append([].reverse())$

$$\begin{aligned} \text{LHS} &= ([] .append(ys)) .reverse() \\ &= ys.reverse() && (\text{append, reverse definition}) \\ \text{RHS} &= ys.reverse().append([].reverse()) \\ &= ys.reverse() && (\text{append, reverse definition}) \end{aligned}$$

Induction: $xs = x:xs'$

IH: Assume $(xy.append(ys)).reverse() = ys.reverse().append(xy.reverse())$ is true for any xy .

$$\begin{aligned} \text{LHS} &= (x:xs' .append(ys)) .reverse() \\ &= x:(xs' .append(ys) .reverse()) && (\text{append definition}) \\ &= xs' .append(ys) .reverse().append(x:[]) \\ &= ys.reverse().append(xs' .reverse()).append(x:[]) \end{aligned}$$

$$\begin{aligned} \text{RHS} &= ys.reverse().append(x:xs' .reverse()) && (\text{reverse definition}) \\ &= ys.reverse().append(xs' .reverse()).append(x:[]) \end{aligned}$$

$$\text{LHS} = \text{RHS}$$

Done

8. $(xs.reverse()).reverse() = xs$

Base case: $xs = []$

Prove: $([].reverse()).reverse() = []$

$$\begin{aligned} \text{LHS} &= ([] .reverse()).reverse() \\ &= [] .reverse() && (\text{reverse definition}) \end{aligned}$$

$$\text{RHS} = []$$

Induction: $xs = x:xs'$

IH: Assume $(xs' .reverse()).reverse() = xs'$ is true for any xs' .

$$\begin{aligned} \text{LHS} &= (x:xs' .reverse()).reverse() \\ &= x:(xs' .reverse().reverse()) && (\text{reverse definition}) \\ &= x:xs' \end{aligned}$$

$$\text{RHS} = x:xs'$$

$$\text{LHS} = \text{RHS}$$

Done

9. $xs.reverseH\ ys = (xs.reverse()).append(ys)$

Base case: $xs = []$

Prove: $[] .reverseH\ ys = ([] .reverse()).append(ys)$

LHS = $[] .reverseH\ ys$ (reverseH definition)

= ys

RHS = $([] .reverse()).append(ys)$

= $[] .append(ys)$ (append, reverse definition)

= ys

Induction: $xs = x:xs'$

IH: Assume $xs'.reverseH(ys) = (x:xs'.reverse()).append(ys)$ is true for any xs' .

LHS = $x:xs'.reverseH(ys)$

= $xs'.reverseH(x:ys)$ (reverseH definition)

= $(x:xs').reverse().append(ys)$ (append definition)

RHS = $(x:xs'.reverse()).append(ys)$

= $(xs'.reverse()).append(ys)$ (reverse definition)

LHS = RHS

Done

10. $(xs.reverseH\ ys).reverseH\ zs = ys.reverseH\ (xs.append(zs))$

Base case: $xs = []$

Prove: $([] .reverseH(ys)).reverseH\ zs = ys.reverseH([] .append(zs))$

LHS = $([] .reverseH(ys)).reverseH\ zs$

= $ys.reverseH(zs)$ (reverseH definition)

RHS = $ys.reverseH([] .append(zs))$

= $ys.reverseH(zs)$ (append definition)

Induction: $xs = x:xs'$

IH: Assume $(xs'.reverseH(ys)).reverseH(zs) = ys.reverseH(xs'.append(zs))$ is true for any xs' .

LHS = $((x:xs').reverseH\ ys).reverseH\ zs$

= $ys.reverseH((x:xs').append(zs))$ (reverseH definition)

RHS = $ys.reverseH(xs'.append(zs))$

LHS = RHS

Done

11. $xs.reverseH(ys.append(zs)) = (xs.reverseH(ys)).append(zs)$

Prove: $xs.reverseH(ys.append(zs)) = xs.reverseH(ys).append(zs)$

```
LHS    = (xs.reverse()).append(ys.append(zs))
        = ((xs.reverse()).append(ys)).append(zs) (append association)
        = (xs.reverseH(ys)).append(zs)
        = RHS
```

LHS = RHS

Done

12. $(xs.append(ys)).reverseH(zs) = ys.reverseH(xs.reverseH(zs))$

Base case: $xs = []$

Prove: $([].append(ys)).reverseH(zs) = ys.reverseH([], reverseH(zs))$

```
LHS    = ([].append(ys)).reverseH(zs)
        = ys.reverseH(zs) (append definition)
```

```
RHS    = ys.reverseH([], reverseH(zs))
        = ys.reverseH(zs) (reverseH definition)
```

Induction: $xs = x:xs'$

IH: Assume $(xs'.append(ys)).reverseH(zs) = ys.reverseH(xs'.reverseH(zs))$ is true for any xs' .

```
LHS    = ((x:xs').append(ys)).reverseH(zs)
        = (xs.append(ys).reverse()).append(zs) (solution 9)
        = (ys.reverse().append(xs.reverse())).append(zs) (solution 7)
        = ys.reverseH(xs.reverse().append(zs)) (solution 9)
        = ys.reverseH((x:xs').reverseH(zs)) (reverseH definition)
```

RHS = $ys.reverseH(xs.reverseH(zs))$

LHS = RHS

Done

13. $(xs.append(ys)).reverse2() = ys.reverse2().append(xs.reverse2())$

Base case: $xs=[]$

Prove: $([].append(ys)).reverse2() = ys.reverse2().append([].reverse2())$

LHS $= ([].append(ys)).reverse2()$

$= ys.reverse2()$

RHS $= ys.reverse2().append([].reverse2())$

$= ys.reverse2().append([])$

$= ys.reverse2()$

Inductive case: $xs=x:xs'$

IH:

Assume $(xs'.append(ys)).reverse2() = ys.reverse2().append(xs'.reverse2())$ is true for any xs' .

LHS $= (x:xs.append(ys)).reverse2()$

$= (x:xs.append(ys)).reverseH([])$

$= (xs.append(ys)).reverseH(x:[])$

$= ys.reverseH(xs.reverseH(x:[]))$ (solution 12)

RHS $= ys.reverse2().append((x:xs).reverse2())$

$= ys.reverseH([]).append(xs.reverseH(x:[]))$ (reverse2)

$= ys.reverseH(xs.reverseH(x:[]))$ (solution 11)

LHS $=$ RHS

Done

14. `(xs.reverse2()).reverse2() = xs`

Base case: `xs=[]`

Prove: `([].reverse2()).reverse2() = []`

```
LHS  = ( [].reverse2() ).reverse2()
      = [].reverseH([]).reverseH([])
      = []
RHS  = []
```

Inductive case: `xs=(x:xs')`

IH: `(xs'.reverse2()).reverse2() = xs'` is true for any `xs'`.

```
LHS  = ((x:xs).reverse2()).reverse2()
      = ((x:xs).reverseH([])).reverseH([])    (reverse2 definition)
      = xs.reverseH(x:[]).reverseH([])
      = (x:[]).reverseH(xs.append([]))          (solution 10)
      = x.reverseH(xs)
      = x.reverse().append(xs)                  (solution 9)
      = x.append(xs')
```

```
RHS  = x:xs'
      = x.append(xs')
```

```
LHS  = RHS
```

Done

15. $t.\text{flattenH}(xs) = t.\text{flatten}().\text{append}(xs)$

Base case: $t = d$

$$\text{LHS} = d.\text{flattenH}(xs) = d:xs$$

$$\begin{aligned}\text{RHS} &= d.\text{flatten}().\text{append}(xs) \\ &= (d:[]).\text{append}(xs) \\ &= d:[].\text{append}(xs) \\ &= \text{LHS}\end{aligned}$$

Inductive case: $t = N(d, t1, t2)$

$$xs.\text{append}(ys.\text{append}(zs)) = (xs.\text{append}(ys)).\text{append}(zs)$$

prove: $N(d, t1, t2).\text{flattenH}(xs) = N(d, t1, t2).\text{flatten}().\text{append}(xs)$

$$\text{IH1: } t1.\text{flattenH}(xs) = t1.\text{flatten}().\text{append}(xs)$$

$$\text{IH2: } t2.\text{flattenH}(xs) = t2.\text{flatten}().\text{append}(xs)$$

$$\begin{aligned}\text{LHS} &= N(d, t1, t2).\text{flattenH}(xs) \\ &= t1.\text{flattenH}(d:t2.\text{flattenH}(xs)) \\ &= t1.\text{flatten}().\text{append}(d:t2.\text{flattenH}(xs)) \\ &= t1.\text{flatten}().\text{append}(d:t2.\text{flatten}().\text{append}(xs))\end{aligned}$$

$$\begin{aligned}\text{RHS} &= N(d, t1, t2).\text{flatten}().\text{append}(xs) \\ &= (t1.\text{flatten}().\text{append}(d:t2.\text{flatten}().\text{append}(xs))).\text{append}(xs) \\ &= t1.\text{flatten}().\text{append}(d:t2.\text{flatten}().\text{append}(xs))\end{aligned}$$

$$\text{LHS} = \text{RHS}$$

Done

16. $t.\text{flatten2}() = t.\text{flatten}()$

Prove:

$$\begin{aligned}\text{LHS} &= t.\text{flatten2}() && (\text{flatten2 definition}) \\ &= t.\text{flattenH}([]) && (\text{flattenH definition}) \\ &= t.\text{flatten}().\text{append}([]) && (\text{append definition}) \\ &= t.\text{flatten}() \\ \text{LHS} &= \text{RHS}\end{aligned}$$

Done

17. $t.\text{map}(f1).\text{sum}() = t.\text{nodes}()$

Base case: $t = \text{leaf}(d)$

Prove: $\text{leaf}(d).\text{map}(f1).\text{sum}() = \text{leaf}(d).\text{nodes}()$

$$\begin{aligned}\text{LHS} &= \text{leaf}(d).\text{map}(f1).\text{sum}() \\ &= \text{leaf}(f1.\text{apply}(d)).\text{sum}() && (\text{definition of map}) \\ &= \text{leaf}(1).\text{sum}() && (f1 \text{ definition}) \\ &= 1 \\ \text{RHS} &= \text{leaf}(d).\text{nodes}() \\ &= 1\end{aligned}$$

Induction case: $T = N(d, t1, t2)$

Assume $t1.\text{map}(f1).\text{sum}() = t1.\text{nodes}()$

Assume $T2.\text{map}(f1).\text{sum}() = t2.\text{nodes}()$

$$\begin{aligned}\text{LHS} &= N(d, t1, t2).\text{map}(f1).\text{sum}() \\ &= N(f1.\text{apply}(d), t1.\text{map}(f), t2.\text{map}(f)).\text{sum}() \\ &= 1 + (t1.\text{map}(f), t2.\text{map}(f)).\text{sum}() && (f1 \text{ definition}) \\ &= 1 + t1.\text{map}(f).\text{sum}() + t2.\text{map}(f).\text{sum}() \\ &= 1 + t1.\text{nodes}() + t2.\text{nodes}() && (\text{IH Assumption}) \\ &= N(d, t1, t2).\text{nodes}()\end{aligned}$$

$$\text{RHS} = N(d, t1, t2).\text{nodes}()$$

$$\text{LHS} = \text{RHS}$$

Done

18. `t.nodes() = t.longestPath().length() + 1.`

Prove:

Base case: $t = \text{leaf}(d)$

```
LHS    = leaf(d).nodes()
        = 1
```

[illegible]

LHS not equal to RHS

Not a statement

Done .

19. For non-empty trees t , it is the case that $t.\text{internalNodes}() + 1 = t.\text{leaves}()$.

Base case:

Nodes = 1

Leaf = d

Internal nodes = 0

LHS = leaf(d).internalNodes() + 1

= 0 + 1

= 1

RHS = leaf(d).leaves()

= 1

Induction case: $T = N(d, t_1, t_2)$

Assume $t_1.\text{internalNodes}() + 1 = t_1.\text{leaves}()$

Assume $t_2.\text{internalNodes}() + 1 = t_2.\text{leaves}()$

LHS = $N(d, t_1, t_2).\text{internalNodes}() + 1$

= leaf(d).internalNodes() + $N(t_1, t_2).\text{internalNodes}() + 1$

= $t_1.\text{internalNodes}() + t_2.\text{internalNodes}() + 1$

(IH Assumption)

= $t_1.\text{internalNodes}() + t_2.\text{leaves}()$

RHS = $N(d, t_1, t_2).\text{leaves}()$

= $t_1.\text{leaves} + t_2.\text{leaves}()$

= $t_1.\text{internalNodes}() + 1 + t_2.\text{leaves}()$

(IH Assumption)

LHS not equal to RHS

Not a statement

Done.

20. A full m -ary with n nodes has $(n - 1)/m$ internal nodes and $((m - 1)n + 1)/m$ leaves.

Base case:

$$\text{Nodes} = 1$$

$$\text{Leaf} = d$$

$$\text{Internal nodes} = 0$$

$$(n - 1)/m = (1-1)/m$$

$$= 0$$

$$((m - 1)n + 1)/m = ((m-1)(1)+1)/m$$

$$= (m-1+1)/m$$

$$= 1$$

Induction case: $T = N(d, t_1, t_2, \dots, t_m)$

$$\text{LHS} = i_1 + i_2 + \dots + i_m + 1$$

$$= (n_1-1)/m + (n_2-1)/m + \dots + (n_m-1)/m + 1$$

$$= (n_1-1 + n_2-1 + \dots + n_m-1)/m + m/m$$

21. A full m -ary with i internal nodes has $mi + 1$ nodes and $(m - 1)i + 1$ leaves.

22. A full m -ary with l leaves has $(ml - 1)/(m - 1)$ nodes and $(l - 1)/(m - 1)$ internal nodes.

23. How many people have seen the letter, including the first person?

24. How many people sent out the letter?