

CSE 301 Concrete Mathematics Lecture Note Scribes

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Lect - 1 (Section A)

CSE 301 M.A. for C.S.

Part A: Tahmid Hasan

Textbook: Concrete Mathematics (Graham, Knuth, Patashnik)

- Syllabus:
- 1) Recurrence (Ch 1)
 - 2) Sums (Ch 2)
 - 3) Number Theory (Ch 4)
 - 4) Special Numbers (Ch 6)
 - 5) Generating Functions (Ch 7)

Recurrence

$$f_n = f_{n-1} + f_{n-2}$$

Tower of Hanoi

$$T_0 = 0$$

$$T_1 = 1$$

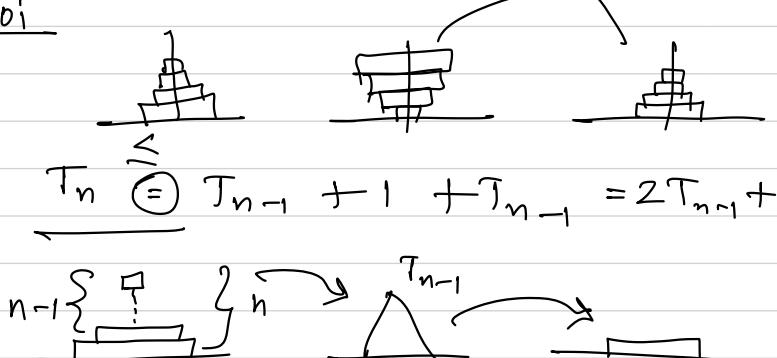
$$T_2 = 3$$

$$T_3 = 7$$

$$T_4 = 15$$

$$\overline{T_5 = 31}$$

$$\overline{T_6 = 63} \quad T_n \geq 1 + T_{n-1} + T_{n-1}$$



$$\boxed{\begin{array}{l} T_0 = 0 \\ T_n = 2T_{n-1} + 1 \end{array}}$$

$$T_n = 2^n - 1$$

Mathematical Induction:

(P(k))

(Sipser p. 29)

1) Basis: $\forall i \ P(i)$

2) Induction step: for any $i > 1$, if $P(i)$ is true, then $P(i+1)$ is true

$$\overbrace{i=1}^{\overbrace{i=2}^{\overbrace{i=3}^{\overbrace{i=4}{}_{\text{induction hypothesis}}}}}$$

$$\overbrace{i=i_0}^{\overbrace{i=i_0+1}^{\overbrace{i=i_0+2}^{\overbrace{i=i_0+3}{}_{\text{induction hypothesis}}}}}$$

1) Basis: $P(i_0)$

2) Induction step: for any $i > i_0$ if $P(i_0), P(i_0+1), \dots, P(i)$ is true, then $P(i+1)$ is true.

$$\begin{aligned} T_0 &= 0 \\ T_n &= 2T_{n-1} + 1 \end{aligned} \quad \left\{ \quad T_n = 2^n - 1 \right.$$

1) Basis: $T_0 = 2^0 - 1 = 0$

Hypo: given $P(i_0), P(i_0+1), \dots, P(i)$ true

$$\begin{aligned} T_{i+1} &= 2T_i + 1 = 2(2^i - 1) + 1 \\ &= 2^{i+1} - 2 + 1 = 2^{i+1} - 1 \end{aligned}$$

CS E 301: M.A. for C.S. Lect-1: Section B

Part A: Tahmid Hasan

Textbook: Concrete Mathematics (Graham, Knuth, Patashnik)

Syllabus: 1) Recurrence (Ch 1)

2) Sums (Ch 2)

3) Number Theory (Ch 4)

4) Special Numbers (Ch 6)

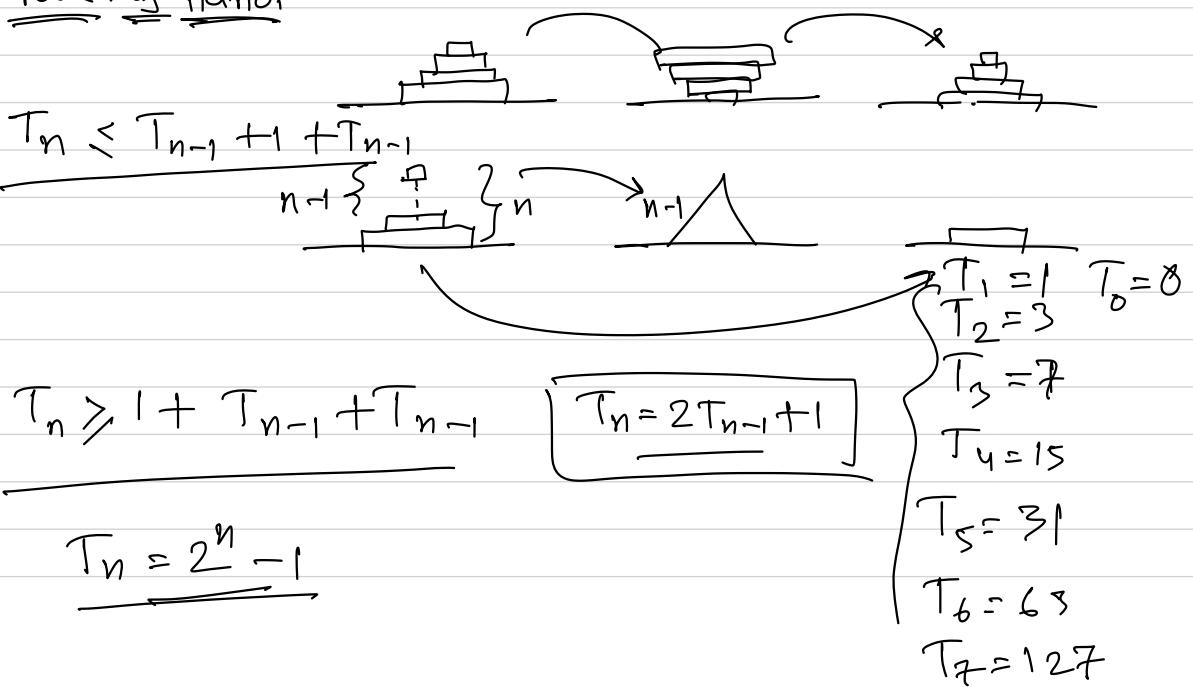
5) Generating functions (Ch 7)

Recurrence

$$S_n = n + S_{n-1}$$

$$f_n = f_{n-1} + f_{n-2}$$

Tower of Hanoi



← \mathbb{N}

Mathematical Induction: $P(k)$ (Sipser P.23)

Basis: $P(1)$ is true.

Induction: for any $i \geq 1$, if $P(i)$ is true, then $P(i+1)$ is true

$$\overline{k=4}$$

$$\overline{k=3}$$

$$\overline{k=2}$$

Basis: $P(i_0)$ is true.

$$\overline{k=i_0+1}$$

$$\overline{k=i_0+3}$$

$$\overline{k=i_0+2}$$

Induction: for any $i \geq i_0$, if $P(i_0), P(i_0+1), \dots, P(i)$ is true, then $P(i+1)$ is true.

$$\overline{T_n = 2^n - 1}$$

$$\overline{T_0 = 0}$$

$$\overline{T_n = 2T_{n-1} + 1}$$

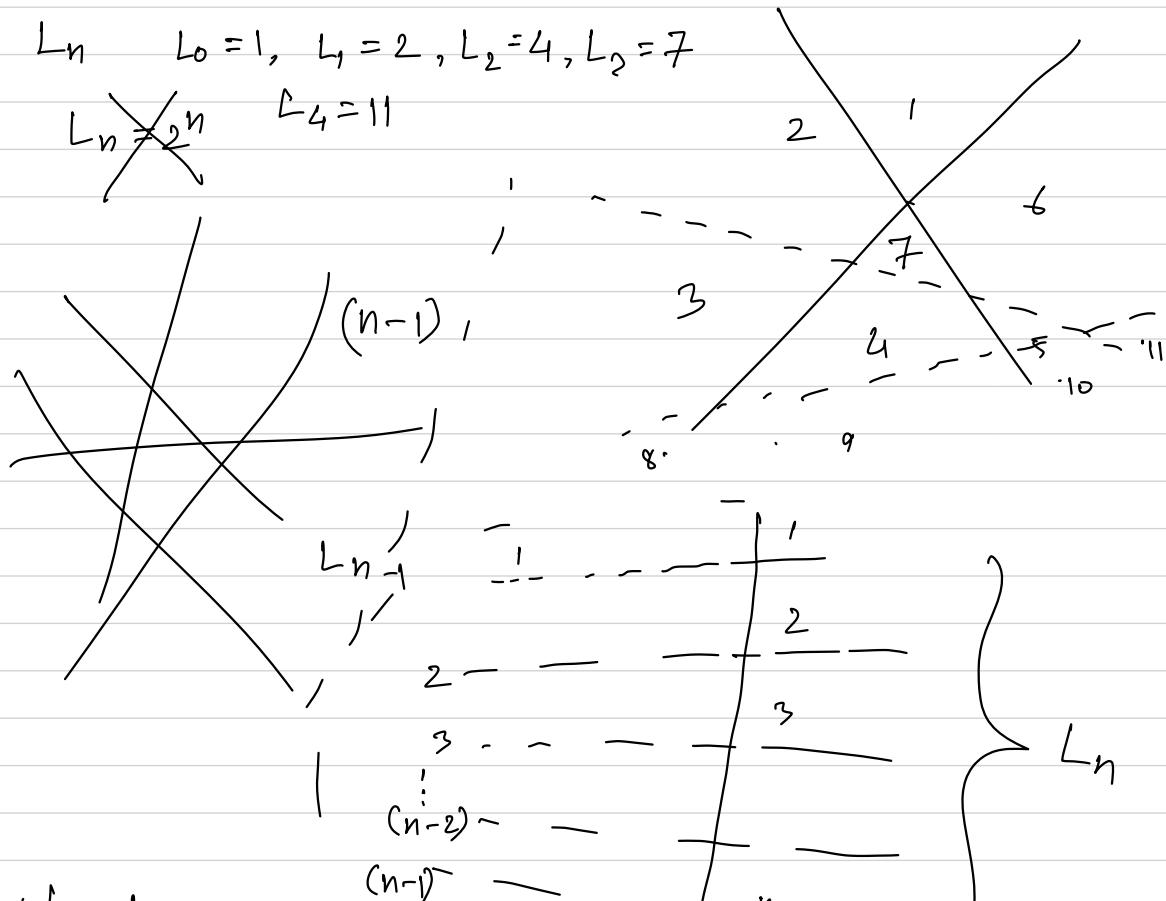
Basis: $i=0, T_0 = 2^0 - 1 = 0$

Induction: for any $i \geq 0$, if $T(i)$ is true, then we need to prove $T(i+1)$ is true

$$T_{i+1} = 2T_i + 1 = 2(2^i - 1) + 1 \\ = 2^{i+1} - 2 + 1 = 2^{i+1} - 1$$

Lect-2: Lines on a plane

Section: B



$$L_0 = 1$$

$$L_n = L_{n-1} + n \cancel{\times} v$$

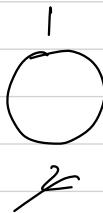
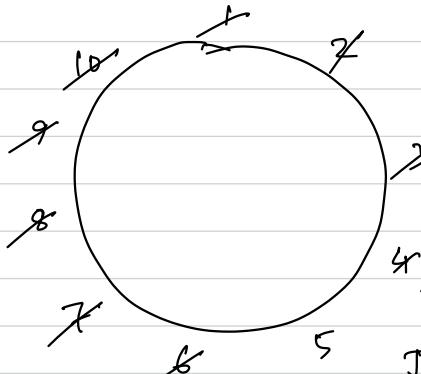
$$= L_{n-2} + (n-1) + n$$

$$= L_{n-3} + (n-2) + (n-1) + n$$

⋮

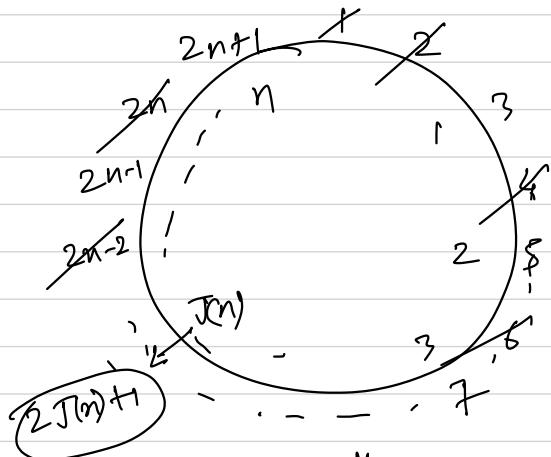
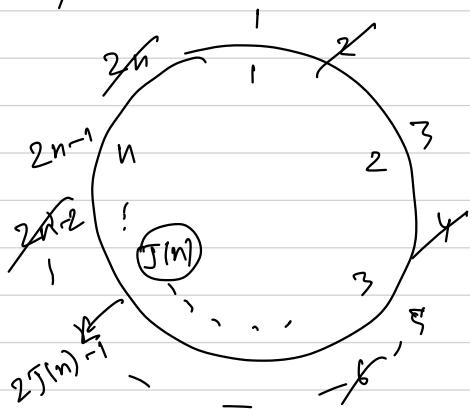
$$= \underbrace{L_0 + 1 + 2 + 3 + \dots + n}_{= 1 + \frac{n(n+1)}{2}}$$

Josephus Problem:



n	1	2	3	4	5	6
$J(n)$	1	1	3	1	3	5

$$\begin{aligned} J(2n) &= 2J(n) - 1 \\ J(2n+1) &= 2J(n) + 1 \end{aligned} \quad J(1) = 1 \quad \left\{ \begin{array}{l} O(\lg n) \end{array} \right.$$



n	2^0	2^1	2^2	2^3	2^4
$J(n)$	1	1	3	1	3

0	1	2	3	4	5	6	7
0	1	3	1	3	5	7	9

$$n = 2^m + l$$

$$J(n) = 2l + 1 \Rightarrow J(2^m + l) = 2l + 1 \quad m \geq 0$$

$$n < 2^{m+1} \Rightarrow 2^m + l < 2^{m+1} \Rightarrow l < 2^{m+1} - 2^m = 2^m \Rightarrow 0 \leq l < 2^m$$

Induction on m :

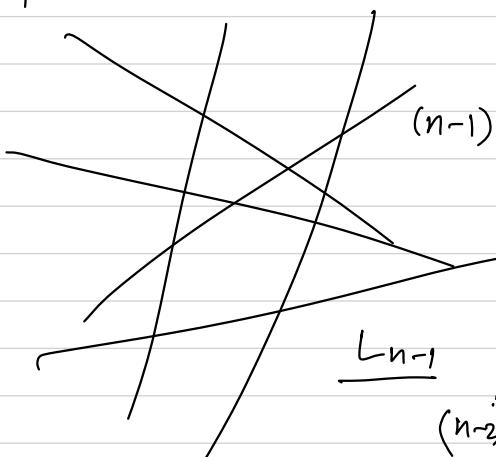
Lect 2: Lines on a plane

Sec. A

$$L_n \quad L_0 = 1, L_1 = 2, L_2 = 4, L_3 = 7, L_4 = 11$$

$$L_n = 2^n$$

$$L_{n-1}$$



$$L_0 = 1$$

$$L_n = L_{n-1} + n$$

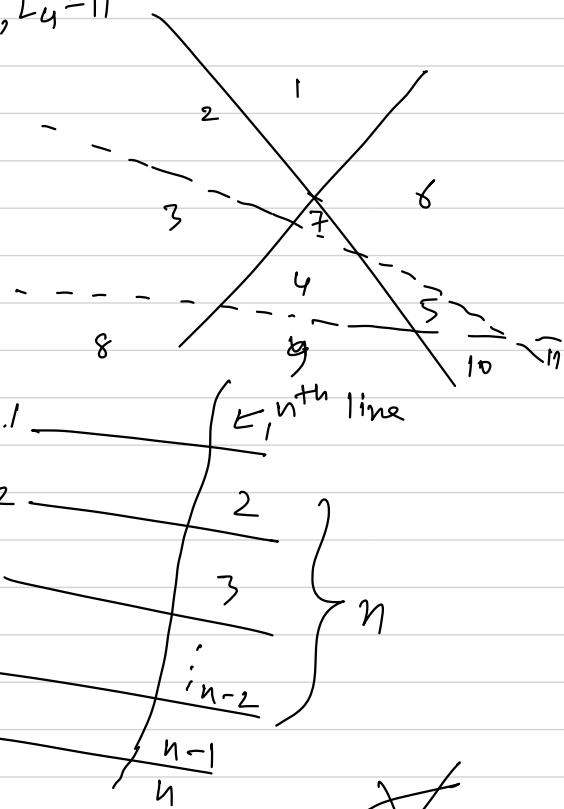
$$= L_{n-2} + (n-1) + n$$

$$= L_{n-3} + (n-2) + (n-1) + n$$

⋮

$$= L_0 + 1 + 2 + 3 + \dots + (n-2) + (n-1) + n$$

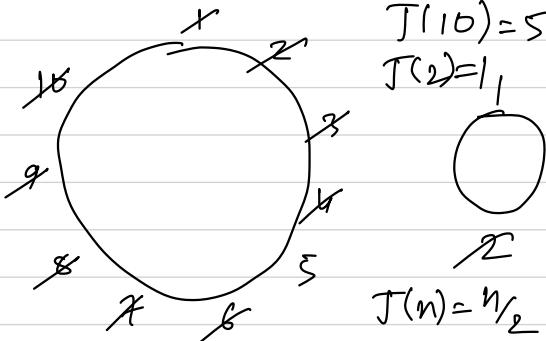
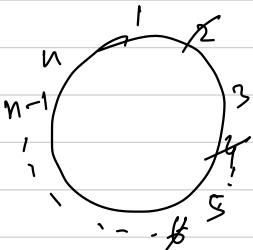
$$= \frac{1}{2} + \frac{n(n+1)}{2}$$



~~X~~

~~X~~

Josephus Problem:

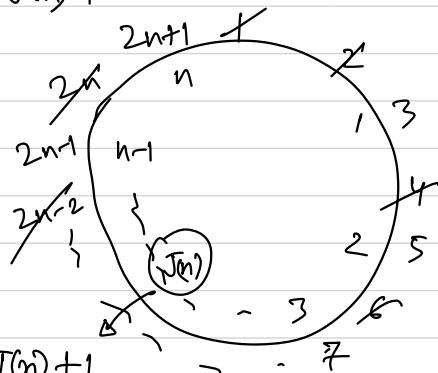
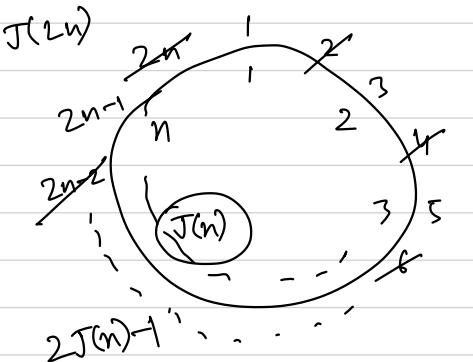


n	1	2	3	4	5	6
$J(n)$	1	1	3	1	3	5

$$J(1) = 1$$

$O(\lg n)$

$$J(2n) = 2J(n) - 1$$



$$J(2n+1) = 2J(n) + 1$$

n	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
$J(n)$	1	1	3	1	3	5	7	1	3	5	7	9	11	13	15	1

$$n = 2^m + l$$

$$J(n) = 2l + 1$$

$$15 = 2^3 + 7$$

$$25 = 2^4 + 9$$

$$n < 2^{m+1} \Rightarrow 2^m + l < 2^{m+1} \Rightarrow l < 2^{m+1} - 2^m = 2^m$$

$$0 \leq l < 2^m$$

$$\begin{aligned} J(1) &= 1 \\ J(2n) &= 2J(n) - 1 \\ J(2n+1) &= 2J(n) + 1 \end{aligned}$$

$$\left\{ \begin{array}{l} J(2^m + l) = 2l + 1; \quad m > 0, \quad 0 \leq l < 2^m \\ \text{Induction on } m \end{array} \right.$$

Lect-3: Josephus contd.

Sec. A

$$T(1) = 1 \quad n = 2^m + l; \quad m \geq 0, \quad 0 \leq l < 2^m$$

$$n = 2^m + l; m \geq 0, 0 \leq l < 2^m$$

$$T(2^m + l) = 2l + 1$$

$$0 \leq \lambda < 2^{\circ} \Rightarrow 0 \leq \lambda < 1$$

$$\begin{aligned}T(2n) &= 2T(n) - 1 \\T(2n+1) &= 2T(n) + 1\end{aligned}$$

$$\text{Basis: } m=0 \quad T(2^m + l) = T(2^0 + 0) = 2 \cdot 0 + 1 = 2l + 1 \\ \Rightarrow T(1) = 1$$

Induction:

Assume true for $0, \dots, M-1$

$$J\left(\frac{2^m + l}{2}\right) = J(2^{m-1} + \lfloor \frac{l}{2} \rfloor) - 1 = 2\left(\cancel{2} \cdot \lfloor \frac{l}{2} \rfloor + 1\right) - 1 = \cancel{2}l + 1$$

$$J\left(\frac{2n}{2^m+l}\right) = J\left(2^m + l - 1 + 1\right) = 2J\left(2^{m-1} + \frac{l-1}{2}\right) + 1$$

$$= 2 \left(2 \cdot \frac{l-1}{2} + 1 \right) + 1 = 2l + 1$$

$$n = (b_m b_{m-1} b_{m-2} \dots b_1 b_0)_2$$

$$n = b_m 2^m + b_{m-1} 2^{m-1} + \dots + b_1 x_2 + b_0$$

$$v = (1 \ b_{m-1} \ b_{m-2} \ \dots \ b_1 \ b_0)_2$$

$$n = 2^m + l \Rightarrow l = n - 2^m$$

$$2^m = (1 \ 0 \ 0 \ \dots 0)_2$$

$$l = (0 \ b_{m-1} \ b_{m-2} \dots \ b_1 \ b_0)_2$$

$$\mathfrak{I}((b_m b_{m-1} \cdots b_1 b_o)) = (b_{m-1} b_{m-2} \cdots b_1 b_o b_m)$$

$$2l = (b_m, b_{m-1}, \dots, b_1, b_0, 0)_2$$

4

$$ZHT = (b_{m-1}, b_{m-2}, \dots, b_1, b_0)_2$$

↳ b_m

$$J((1101)_2) = \underline{(1011)_2} \rightarrow (111)_L \rightarrow (111)_2$$

fixed point
 $f(n) = n$

$$T(n) = n/2$$

$$\Rightarrow 2\ell + 1 = (2^m + \ell)/2$$

$$\Rightarrow 4\ell + 2 = 2^m + \ell \Rightarrow \ell = \frac{1}{3}(2^m - 2)$$

$$m=0 \rightarrow \ell \times$$

$$m=1 \rightarrow \ell = 0$$

$$m=2 \rightarrow \ell \times$$

$$m=3 \rightarrow \ell = 2$$

//

<u>m</u>	<u>ℓ</u>	<u>$n = 2^m + \ell$</u>	<u>$J(n) = 2\ell + 1 = n/2$</u>	<u>n (binary)</u>
1	0	2	1	10
3	2	10	5	1010
5	10	42	21	101010
7	42	170	85	10101010

$$J(1) = 1$$

$$f(1) = \alpha$$

$$J(2n) = 2J(n) - 1$$

$$f(2n) = 2f(n) + f$$

$$J(2n+1) = 2J(n) + 1$$

$$f(2n+1) = 2f(n) + \gamma$$

$$n \quad f(n)$$

$$f(n) = A(n)\alpha + B(n)\beta + C(n)\gamma$$

$$1. \quad \alpha$$

$$2. \quad 2\alpha + \beta$$

$$3. \quad 2\alpha + \gamma$$

$$4. \quad 4\alpha + 3\beta$$

$$5. \quad 4\alpha + 2\beta + \gamma$$

$$6. \quad 4\alpha + \beta + 2\gamma$$

$$7. \quad 4\alpha + 3\beta$$

$$8. \quad 8\alpha + 7\beta$$

$$9. \quad 8\alpha + 6\beta + \gamma$$

$$n = 2^m + \ell$$

$$A(n) = 2^m$$

$$B(n) = 2^m - \ell - 1$$

$$C(n) = \ell$$

$$\frac{3n+2}{(n-1)(n-2)} = \frac{A}{(n-1)} + \frac{B}{(n-2)}$$

$$\frac{B(n) + C(n)}{\ell} = \frac{2^m - \ell - 1 + \ell}{\ell} = \frac{2^m - 1}{\ell}$$

Lect-3: Josephus cont.

Sec. B

$$J(1) = 1$$

$$n = 2^m + \ell, m \geq 0 \text{ & } 0 \leq \ell < 2^m$$

$$J(2n) = 2J(n) - 1$$

$$J(n) = 2\ell + 1 \quad \dots \quad (1)$$

$$J(2n+1) = 2J(n) + 1$$

$$0 \leq \ell < 2^0 \Rightarrow 0 \leq \ell < 1$$

$$\Leftrightarrow \ell = 0$$

Basis: $m=0$; $J(2^0 + \ell) = 2\ell + 1$

$$\Leftrightarrow J(2^0 + 0) = 2 \cdot 0 + 1$$

$$\Leftrightarrow J(1) = 1$$

Induction: Assume (1) is true for $0, 1, \dots, m-1$

$$J(\overbrace{2^m + \ell}^{2n}) = 2J(\overbrace{2^{m-1} + \ell/2}^0) - 1 = 2(2 \cdot \ell/2 + 1) - 1 = 2\ell + 1$$

$$J(\overbrace{2^m + \ell}^0) = J(\overbrace{2^m + \ell-1 + 1}^0) = 2J(\overbrace{2^{m-1} + \frac{\ell-1}{2}}^0) + 1 \\ = 2\left(2 \cdot \frac{\ell-1}{2} + 1\right) + 1 = 2\ell + 1$$

$$n = (b_m b_{m-1} \dots b_1 b_0)_2$$

$$n = b_m 2^m + b_{m-1} 2^{m-1} + \dots + 2b_1 + b_0$$

$$n = (1 b_{m-1} \dots b_1 b_0)_2$$

$$n = 2^m + \ell \Rightarrow \ell = n - 2^m$$

$$2^m = (1 0 \dots 0 0)_2$$

$$\ell = (0 b_{m-1} \dots b_1 b_0)_2$$

$$J((b_m b_{m-1} \dots b_1 b_0)_2) = (\underbrace{b_{m-1} b_{m-2} \dots b_1}_{111,111,111}, b_0 b_m)_2$$

$$2\ell = (b_{m-1} b_{m-2} \dots b_0 0)_2$$

$$2\ell + 1 = (b_{m-1} b_{m-2} \dots b_0 1)_2$$

$$(011)_2 \rightarrow (0111)_2 \rightarrow (111)_2 \rightarrow (1111)_2$$

$$\underbrace{10110010101101}_{111,111,111}$$

$$J(n) = n \rightarrow \text{fixed point}$$

$$f(n) = n$$

$$J(n) = \frac{n}{2}$$

$$n = 2^m + l \Rightarrow J(n) = 2l + 1$$

$$\Rightarrow 2l + 1 = (2^m + l)/2 \Rightarrow 4l + 2 = 2^m + l \Rightarrow l = \frac{1}{3}(2^m - 2)$$

<u>m</u>	<u>l</u>	<u>$n = 2^m + l$</u>	<u>$J(n) = \frac{n}{2} = 2l + 1$</u>	<u>n (binary)</u>
1	0	2	1	10
3	2	10	5	1010
5	10	42	21	101010
7	42	170	85	10101010

$$J(1) = 1$$

$$J(2n) = 2J(n) - 1$$

$$J(2n+1) = 2J(n) + 1$$

$$\left. \begin{array}{l} f(1) = \alpha \\ f(2n) = 2f(n) + \beta \\ f(2n+1) = 2f(n) + \gamma \end{array} \right\} f(n) = A(n)\alpha + B(n)\beta + C(n)\gamma$$

<u>n</u>	<u>f(n)</u>
1.	<u>α</u>
2.	$2\alpha + \beta$
3.	$2\alpha + \gamma$
4.	$4\alpha + 3\beta$
5.	$4\alpha + 2\beta + \gamma$
6.	$4\alpha + \beta + 2\gamma$
7.	<u>$4\alpha + 3\gamma$</u>
8.	$8\alpha + 7\beta$
9.	$8\alpha + 6\beta + \gamma$

$$n = 2^m + l$$

$$A(n) = \underline{\underline{2^m}}$$

$$C(n) = l$$

$$B(n) + C(n) = 2^m - 1$$

$$\Rightarrow B(n) + l = 2^m - 1$$

$$\Rightarrow B(n) = \underline{\underline{2^m - l - 1}}$$

Lect. 4: Generalized Josephus

sec. A

$$f(1) = \alpha$$

$$f(2n) = 2f(n) + \beta$$

$$f(2n+1) = 2f(n) + \gamma$$

$$f(n) = A(n)\alpha + B(n)\beta + C(n)\gamma \dots \quad (1)$$

$$A(n) = 2^m$$

$$B(n) = 2^m - l - 1$$

$$n = 2^m + l$$

$$C(n) = l$$

$$(\alpha, \beta, \gamma) = (1, 0, 0)$$

$$A(1) = 1$$

$$A(2n) = 2A(n)$$

$$A(2n+1) = 2A(n)$$

$$\frac{3n+2}{(x-1)(x-2)} = \frac{A}{x-1} + \frac{B}{x-2}$$

$$\begin{array}{ccccccccc} n & | & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 \\ A(n) & | & 1 & 2 & 2 & 4 & 4 & 4 & 4 & 8 & 8 \end{array}$$

$$f(n) = A(n)$$

$$\boxed{A(n) = 2^m} \quad \checkmark$$

$$f(n) = 1 \Rightarrow 1 = \alpha$$

$$1 = 2 \cdot 1 + \beta$$

$$1 = 2 \cdot 1 + \gamma$$

$$\boxed{A(n) - B(n) - C(n) = 1} \Rightarrow B(n) = A(n) - C(n) - 1 = \boxed{2^m - l - 1} \quad \checkmark$$

$$(\alpha, \beta, \gamma) = (1, -1, -1)$$

$$f(n) = n \Rightarrow 1 = \alpha$$

$$2n = 2 \cdot n + \beta$$

$$2n+1 = 2 \cdot n + \gamma$$

$$\boxed{A(n) + C(n) = n}$$

$$(\alpha, \beta, \gamma) = (\underline{1}, \underline{0}, \underline{1})$$

$$C(n) = n - 2^m = \underline{l} \quad \checkmark$$

Repertoire method: ex. 16 & 20

$$f((b_m b_{m-1} \dots b_1 b_0)_2) = (b_{m-1} \dots b_1 b_0 b_m)_2; \quad b_m = 1$$

$$f(1) = \alpha$$

$$f(\overline{2n+j}) = 2f(n) + \beta_j; \quad j \in \{0, 1\}, \quad \beta_0 = \beta, \quad \beta_1 = \gamma$$

$$d \mid \frac{a}{dq} \mid q$$

$$a = dq + r$$

$$2 \mid \frac{2n+j}{j} \mid n$$

$$\begin{aligned}
f((b_m b_{m-1} \dots b_1 b_0)_2) &= 2f((b_m b_{m-1} \dots b_1)_2) + \beta_{b_0} \\
&= 2(f((b_m b_{m-1} \dots b_2)_2) + \beta_{b_1}) + \beta_{b_0} \\
&= 4f((b_m b_{m-1} \dots b_2)_2) + 2\beta_{b_1} + \beta_{b_0} \\
&= 8f((b_m b_{m-1} \dots b_3)_2) + 4\beta_{b_2} + 2\beta_{b_1} + \beta_{b_0} \\
&\vdots \\
&= 2^m f((b_m)_2) + 2^{m-1} \beta_{b_{m-1}} + 2^{m-2} \beta_{b_{m-2}} + \dots + 2\beta_{b_1} + \beta_{b_0} \\
&= 2^m \alpha + 2^{m-1} \beta_{b_{m-1}} + 2^{m-2} \beta_{b_{m-2}} + \dots + 2\beta_{b_1} + \beta_{b_0} \\
&= (\alpha \beta_{b_{m-1}} \beta_{b_{m-2}} \dots \beta_{b_1} \beta_{b_0})_2 \text{ - notation abuse}
\end{aligned}$$

$$J(100) = J(2^6 + 36) = 2 \cdot 36 + 1 = 73$$

$$\alpha = 1$$

$$\beta = -1 = \beta_0$$

$$\gamma = 1 = \beta,$$

$$100 \rightarrow (1100100)_2$$

$$\begin{aligned}
&\frac{2^6 x_1 + 2^5 x_1 + 2^4 x(-1) + 2^3 x(-1) + 2^2 x_1 + 2^1 x(-1) + 2^0 x(-1)}{=} \\
&= \underline{\underline{73}}
\end{aligned}$$

$$f(j) = \alpha_j \quad \text{for } 1 \leq j < d$$

$$f(dn+j) = c f(n) + \beta_j \quad \text{for } 0 \leq j < d; \quad n \geq 1$$

$$f((b_m b_{m-1} \dots b_1 b_0)_d) = (\alpha_{b_m} \beta_{b_{m-1}} \beta_{b_{m-2}} \dots \beta_{b_1} \beta_{b_0})_c$$

Lect. 4: Generalized Josephus

Sec. B

$$f(1) = \alpha$$

$$f(n) = A(n)\alpha + B(n)\beta + C(n)\gamma \dots (1)$$

$$f(2n) = 2f(n) + \beta$$

$$A(n) = 2^m, B(n) = 2^m - l - 1, C(n) = l$$

$$f(2n+1) = 2f(n) + \gamma$$

$$\frac{3x+2}{(x-1)(x-2)} = \frac{A}{x-1} + \frac{B}{x-2}$$

$$(\alpha, \beta, \gamma) = (1, 0, 0) \quad A(1) = 1$$

$$f(m) = A(n)$$

$$A(2n) = 2A(n)$$

<u>n</u>	1	2	3	4	5	6	7	8	9
A(n)	1	2	2	4	4	4	8	8	

$$A(2n+1) = 2A(n)$$

$$n = 2^m + l \Rightarrow \boxed{A(n) = 2^m} \checkmark$$

$$f(n) = 1 \Rightarrow \begin{cases} 1 = \alpha \\ 1 = 2 \cdot 1 + \beta \\ 1 = 2 \cdot 1 + \gamma \end{cases} \left\} (\alpha, \beta, \gamma) = (1, -1, -1)$$

$$\boxed{A(n) - B(n) - C(n) = 1}$$

$$f(n) = n \Rightarrow \begin{cases} 1 = \alpha \\ 2n = 2 \cdot n + \beta \\ 2n+1 = 2 \cdot n + \gamma \end{cases} \left\} (\alpha, \beta, \gamma) = (1, 0, 1)$$

$$\boxed{A(n) + C(n) = n} \Rightarrow C(n) = n - 2^m = \underline{\underline{l}} \checkmark$$

$$A(n) - B(n) - C(n) = 1 \Rightarrow B(n) = A(n) - C(n) - 1 = 2^m - l - 1 \checkmark$$

Repertoire method: Ex. 16 & 20.

$$J((b_m b_{m-1} \dots b_1 b_0)_2) = (b_{m-1} \dots b_1 b_0 b_m)_2 \xrightarrow[b_m=1]{d}$$

$$\begin{array}{c|cc} a & | & q \\ \hline dq & & \\ \hline r & & \end{array}$$

$$f(1) = \alpha$$

$$f(2n+j) = 2f(n) + \beta j; \quad j \in \{0, 1\}, \quad \beta_0 = \beta, \quad \beta_1 = \gamma$$

$$2 \left[\frac{2n+j}{2n} \right] \mid n$$

$$a = dq + r, \quad 0 \leq r < d$$

$$\begin{aligned}
f((b_m b_{m-1} \dots b_1 b_0)_2) &= 2f((b_m b_{m-1} \dots b_1)_2) + \beta_{b_0} \\
&= 2(2f((b_m b_{m-1} \dots b_2)_2) + \beta_{b_1}) + \beta_{b_0} \\
&= 4f((b_m b_{m-1} \dots b_2)_2) + 2\beta_{b_1} + \beta_{b_0} \\
&= 8f((b_m b_{m-1} \dots b_3)_2) + 4\beta_{b_2} + 2\beta_{b_1} + \beta_{b_0} \\
&\vdots \\
&= 2^m f((b_m)_2) + 2^{m-1} \beta_{m-1} + \dots + 4\beta_{b_2} + 2\beta_{b_1} + \beta_{b_0} \\
&= 2^m \alpha + 2^{m-1} \beta_{m-1} + \dots + 4\beta_{b_2} + 2\beta_{b_1} + \beta_{b_0} \\
&= (\alpha \beta_{b_{m-1}} \dots \beta_{b_2} \beta_{b_1} \beta_{b_0})_2 \rightarrow \text{notation abuse!}
\end{aligned}$$

$$J(100) = J(2^6 + 36) = 36 \cdot 2 + 1 = 73$$

$$\underline{(b_m b_{m-1} b_{m-2} \dots b_1 b_0 \mid 1 0 0 \mid 0 0)_2}$$

$$\begin{aligned}
J(1) &= 1 & \alpha &= 1 \\
J(2n) &= 2J(n) - 1 & \beta &= -1 = \beta_0 \\
J(2n+1) &= 2J(n) + 1 & \gamma &= 1 = \beta_1
\end{aligned}$$

$$\begin{aligned}
2^6 x 1 + 2^5 x 1 + 2^4 x (-1) + 2^3 x (-1) + 2^2 x 1 + 2^1 x (-1) + 2^0 x (-1) \\
= 64 + 32 - 16 - 8 + 4 - 2 - 1 = 73
\end{aligned}$$

$$\begin{aligned}
f(j) &= \alpha_j & 1 \leq j < d \\
f(dn+j) &= c f(n) + \beta_j & 0 \leq j < d ; n \geq 1 \\
f((b_m b_{m-1} \dots b_1 b_0)_d) &= (\alpha_{b_m} \beta_{b_{m-1}} \beta_{b_{m-2}} \dots \beta_{b_1} \beta_{b_0})_d
\end{aligned}
\right| \quad
\begin{aligned}
f(1) &= 1 \\
f(2) &= 2 \\
f(3n) &= 4f(n) + 1 \\
f(3n+1) &= 4f(n) + 2 \\
f(3n+2) &= 4f(n) + 3
\end{aligned}$$

Lect 5: Sums

Sec. B

$$\text{Sum of first } n \text{ integers: } 1 + 2 + 3 + \dots + (n-1) + n$$

$$1 + 2 + \dots + n \quad \xrightarrow{\text{term}}$$

$$1 + \dots + n$$

$a_1 + a_2 + \dots + a_n \rightarrow \text{general term: } a_k$

$$1 + 2 + \dots + 2^{n-1} \xrightarrow{n} a_k = 2^{k-1} \mid 2^0 + 2^1 + \dots + 2^{n-1}$$

Sigma notation: $\sum_{k=1}^n a_k \rightarrow \text{delimited form}$

Generalized sigma notation: $\sum_{1 \leq k \leq n} a_k \rightarrow \sum_{p(k)} a_k ; \text{sum of all terms satisfying } p(k)$

$$\sum_{\substack{1 \leq k \leq 100 \\ k \text{ odd}}} k^2 \equiv \sum_{k=0}^{49} (2k+1)^2 \mid \sum_{\substack{1 \leq p \leq n \\ p \text{ prime}}} \frac{1}{p} = \sum_{k=1}^{\pi(n)} \frac{1}{p_k}$$

$$\sum_{1 \leq k \leq n} a_k = \sum_{1 \leq k+1 \leq n} a_{k+1} ;$$

$$\sum_{k=2}^{n-1} k(k-1)(n-k) = \sum_{k=0}^n k(k-1)(n-k)$$

$$\sum_k a_k [p(k)]$$

$$p(k) \equiv 1 \leq k \leq n$$

$$\sum_p \frac{1}{p} [\text{prime}] [p \leq n]$$

[statement] = $\begin{cases} 1 & \text{if statement is true} \\ 0 & \text{if statement is false} \end{cases}$

\hookrightarrow Inversion notation

$$\text{Sum} \rightarrow \text{Recurrence: } S_n = \sum_{k=0}^n a_k \rightarrow S_0 = a_0 ; S_n = \sum_{k=0}^{n-1} a_k + a_n$$

$$\Rightarrow S_n = S_{n-1} + a_n$$

$$a_k = \beta + \gamma k \quad \boxed{R_0 = \alpha, R_n = R_{n-1} + \beta + \gamma n}$$

$$R_1 = \alpha + \beta + \gamma, R_2 = \alpha + 2\beta + 3\gamma \rightarrow R_n = A(n)\alpha + B(n)\beta + C(n)\gamma$$

$$R_n = 1; \quad 1 = \alpha; \quad 1 = 1 + \beta + \gamma n \Rightarrow \overbrace{\beta + \gamma n}^{\sim \sim \sim} = 0 + 0 \cdot n \quad (\alpha, \beta, \gamma) = (1, 0, 0)$$

$$\boxed{A(n) = 1}$$

$$R_n = n; \quad 0 = \alpha, \quad n = n - 1 + \beta + \gamma n \Rightarrow \beta + \gamma n = 1 + 0 \cdot n \quad (\alpha, \beta, \gamma) = (0, 1, 0)$$

$$\boxed{B(n) = n}$$

$$R_n = n^2; \quad \alpha = 0, \quad n^2 = (n-1)^2 + \beta + \gamma n \Rightarrow \beta + \gamma n = 2n-1; \quad (\alpha, \beta, \gamma) = (0, -1, 2)$$

$$2C(n) - B(n) = n^2 \Rightarrow C(n) = \frac{(n^2+n)}{2}$$

$$R_n = \alpha + \beta n + \gamma(n^2+n)/2; \quad a_k = \beta + \gamma k; \quad a_0 = \alpha$$

$$\sum_{k=0}^n (a+bk)$$

Recurrence \rightarrow Sum

$$T_0 = 0, \quad T_n = 2T_{n-1} + 1 \rightarrow T_n/2^n = T_{n-1}/2^{n-1} + 1/2^n; \quad S_n = T_n/2^n$$

$$\hookrightarrow T_0/2^0 = 0/2^0;$$

$$\hookrightarrow S_n = S_{n-1} + 1/2^n$$

$$= S_{n-2} + 1/2^{n-1} + 1/2^n$$

$$= ;$$

$$= S_0 + \frac{1}{2} + \frac{1}{2^2} + \dots + \frac{1}{2^n}$$

$$= 0 + \frac{\frac{1}{2}(1 - \frac{1}{2^n})}{1 - \frac{1}{2}} = 1 - \frac{1}{2^n}$$

$$S_n = 1 - 1/2^n$$

$$\Rightarrow T_n/2^n = 1 - 1/2^n$$

$$\Rightarrow T_n = 2^n - 1$$

Lect. 5: Sums

Sum of first n integers: $1+2+3+\dots+(n-1)+n$

$$1+2+\dots+n \quad \xrightarrow{\text{term}}$$

$$1+\dots+n$$

$a_1+a_2+\dots+a_n$. General term: a_k

$1+2+\dots+2^{n-1}$; $a_k = 2^{k-1}$

$\sum_{k=1}^n a_k$ (sigma notation) $\sum_{1 \leq k \leq n} a_k$ (generalized sigma notation)

$$\sum_{\substack{1 \leq k \leq 100 \\ k \text{ odd}}} k^2$$

$$\sum_{k=0}^{49} (2k+1)^2$$

$$\left| \sum_{\substack{p \leq n \\ p \text{ prime}}} \frac{1}{p} \right|$$

$$\sum_{k=1}^{\pi(n)} \frac{1}{p_k}$$

$$\sum_{1 \leq k \leq n} a_k = \sum_{1 \leq k \leq n} a_{k+1}; \quad \sum_{k=1}^n a_k = \sum_{k=0}^{n-1} a_{k+1}$$

$\sum a_k$ sum over all k satisfying property $P(k)$

$$\sum_{k=2}^{n-1} k(k-1)(n-k) = \sum_{k=0}^n k(k-1)(n-k)$$

[statement] = $\begin{cases} 1, & \text{if statement is true} \\ 0, & \text{if statement is false} \end{cases}$ Inversion notation

$$\sum_K a_K [P(K)] \quad \sum_p \frac{1}{p} [p \leq n] [p \text{ prime}]$$

Sums \rightarrow Recurrence $S_n = \sum_{k=0}^n a_k$, $S_0 = a_0$; $S_n = \sum_{k=0}^{n-1} a_k + a_n$

$$S_n = S_{n-1} + a_n$$

$$a_n = \beta + \gamma n; \quad R_0 = \alpha, \quad R_n = R_{n-1} + \beta + \gamma n$$

$$R_1 = \alpha + \beta + \gamma, \quad R_2 = \alpha + 2\beta + 3\gamma \quad R_n = A(n)\alpha + B(n)\beta + C(n)\gamma$$

$$R_n = 1; \quad 1 = \alpha, \quad 1 = 1 + \beta + \gamma n \Rightarrow \underbrace{\beta + \gamma n}_{} = 0 + 0 \cdot n \quad (\alpha, \beta, \gamma) = (1, 0, 0)$$

$$[A(n) = 1]$$

$$R_n = n; \quad 0 = \alpha, \quad n = n - 1 + \beta + \gamma n \Rightarrow \underbrace{\beta + \gamma n}_{} = \underbrace{1 + 0 \cdot n}_{} \quad (\alpha, \beta, \gamma) = (0, 1, 0)$$

$$[B(n) = n]$$

$$R_n = n^2; \quad 0 = \alpha, \quad n^2 = (n-1)^2 + \beta + \gamma n \Rightarrow \beta + \gamma n = -1 + 2n \quad (\alpha, \beta, \gamma) = (0, -1, 2)$$

$$- B(n) + 2C(n) = n^2 \Rightarrow C(n) = (n^2 + n)/2$$

$$\sum_{k=0}^n (a + b k)$$

$$\begin{aligned} a &= \alpha \\ a &= \beta \\ b &= \gamma \end{aligned}$$

$$\begin{array}{c} R_0 = a \quad (\alpha) \\ R_n = \beta + \gamma n \\ \downarrow \quad \downarrow \\ b \end{array}$$

Recurrence \rightarrow Sum

$$T_0 = 0, \quad T_n = 2T_{n-1} + 1 \Rightarrow T_n/2^n = T_{n-1}/2^{n-1} + 1/2^n; \quad S_n = T_n/2^n$$

$$\hookrightarrow T_0/2^0 = 0/2^0 \qquad \qquad \qquad S_{n-1} = T_{n-1}/2^{n-1}$$

$$\begin{aligned} \Rightarrow S_n &= S_{n-1} + 1/2^n \\ &= S_{n-2} + 1/2^{n-1} + 1/2^n \\ &\vdots \\ &= S_0 + \underbrace{1/2 + 1/2^2 + \dots + 1/2^n}_{\text{...}} \end{aligned}$$

$$T_n/2^n = \frac{\frac{1}{2}(1 - \frac{1}{2^n})}{1 - \frac{1}{2}} = 1 - 1/2^n$$

$$\Rightarrow T_n = 2^n - 1$$

Lect 6: Summation factors and manipulation of sums

Sec. A

$$T_0 = 0, T_n = 2T_{n-1} + 1$$

$$a_n T_n = b_n T_{n-1} + c_n \Rightarrow s_n a_n T_n = s_n b_n T_{n-1} + s_n c_n$$

$$[s_n b_n = s_{n-1} a_{n-1}] \Rightarrow s_n a_n T_n = s_{n-1} a_{n-1} T_{n-1} + s_n c_n$$

$$[s_n a_n T_n = \tilde{s}_n] \Rightarrow \tilde{s}_n = s_{n-1} + s_n c_n$$

$$= \tilde{s}_{n-2} + s_{n-1} c_{n-1} + s_n c_n$$

$$\vdots \\ \Rightarrow s_n a_n T_n = \tilde{s}_0 + \sum_{k=1}^n s_k c_k = s_0 a_0 T_0 + \sum_{k=1}^n s_k c_k = s_0 b_0 T_0 + \sum_{k=1}^n s_k c_k$$

$$\Rightarrow T_n = \frac{1}{s_n a_n} (s_0 b_0 T_0 + \sum_{k=1}^n s_k c_k)$$

$$s_n b_n = s_{n-1} a_{n-1} \Rightarrow s_n = \frac{a_{n-1}}{b_n} s_{n-1} = \frac{a_{n-1}}{b_n} \times \frac{a_{n-2}}{b_{n-1}} \times s_{n-2} = \dots = \frac{a_{n-1} \dots a_1}{b_n \dots b_2} s_1$$

$$T_0 +: a_n = 1, b_n = 2, c_n = 1; s_n = \underbrace{\frac{1}{2} \times \frac{1}{2} \times \dots \times \frac{1}{2}}_{(n-1)} \times \frac{1}{2} = \frac{1}{2^n}$$

$$T_n = 2^n \times 1 \times \sum_{k=1}^n \frac{1}{2^k} \times 1 = 2^n - 1$$

Quicksort:

$a_1 a_2 a_3$	↓	$a_m a_n$
1 1 - - * - - 1		

$a_1 < a_k > a_j$

$$C_n = (n+1) + C_0 + C_{n-1}$$

$$C_n = (n+1) + C_1 + C_{n-2}$$

$$C_n = (n+1) + C_2 + C_{n-3}$$

:

$$\underline{C_n = (n+1) + C_{n-1} + C_0}$$

$$n C_n = n(n+1) + 2 \sum_{k=0}^{n-1} C_k \Rightarrow (n-1) C_{n-1} = n(n-1) + 2 \sum_{k=0}^{n-2} C_k$$

$$\Rightarrow n C_n - (n-1) C_{n-1} = 2n + 2C_{n-1}$$

$$\Rightarrow C_n = n+1 + \frac{2}{n} \sum_{k=0}^{n-1} C_k$$

$$\Rightarrow n C_n = (n+1) C_{n-1} + 2n$$

$$a_n T_n = \overbrace{b_n}^1 T_{n-1} + \overbrace{C_n}^1$$

$$s_n = \frac{a_{n-1}a_{n-2}\dots a_1}{b_n b_{n-1} \dots b_2} = \frac{(n-1)(n-2)\dots 2 \cdot 1}{(n+1) n (n-1) \dots 3} = \frac{2}{n(n+1)}$$

$$C_n = \ln(n+1)/2 \times \frac{1}{2} \sum_{k=1}^n \frac{2}{k(k+1)} \times 2k = 2(n+1) \sum_{k=1}^n \frac{1}{k+1}$$

$$H_n = 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} = \sum_{k=1}^n \frac{1}{k} \quad [\text{Harmonic number}]$$

$$\begin{aligned} \sum_{k=1}^n \frac{1}{k+1} &= \sum_{1 \leq k \leq n} \frac{1}{k+1} = \sum_{1 \leq k-1 \leq n} \frac{1}{k} = \sum_{2 \leq k \leq n+1} \frac{1}{k} = \sum_{2 \leq k \leq n} \frac{1}{k} + \frac{1}{n+1} \\ &= \sum_{1 \leq k \leq n} \frac{1}{k} + \frac{1}{n+1} - 1 \end{aligned}$$

$$\begin{aligned} C_n &= 2(n+1) \left(H_n - \frac{1}{n+1} \right) \\ &= 2(n+1)H_n - 2n \quad \lim_{n \rightarrow \infty} H_n = \ln n + \gamma \quad [0.577\dots, \text{a constant}] \\ &= O(n \lg n) \end{aligned}$$

Manipulation of sums: (\mathbb{K} is a finite set of integers)

$$1) \sum_{K \in \mathbb{K}} c a_K = c \sum_{K \in \mathbb{K}} a_K \quad (\text{distributive law})$$

$$2) \sum_{K \in \mathbb{K}} (a_K + b_K) = \sum_{K \in \mathbb{K}} a_K + \sum_{K \in \mathbb{K}} b_K \quad (\text{associative law})$$

$$3) \sum_{K \in \mathbb{K}} a_K = \sum_{P(K) \in \mathbb{K}} a_{P(K)} \quad (\text{commutative law}) \quad P(K) \text{ is a permutation over } \mathbb{Z}$$

$$\mathbb{K} = \{-1, 0, 1\}, P(K) = -K \quad \begin{matrix} -1 \rightarrow 1 \\ 0 \rightarrow 0 \\ 1 \rightarrow -1 \end{matrix}$$

$$a_{-1} + a_0 + a_1 = a_1 + a_0 + a_{-1} \quad P(K) = n-K \quad \begin{matrix} 0 \rightarrow n \\ 1 \rightarrow n-1 \\ \vdots \rightarrow 0 \end{matrix}$$

$$0 \leq n-K \leq n \Rightarrow \begin{matrix} K \leq n \\ 0 \leq K \\ \rightarrow 0 \leq K \leq n \end{matrix}$$

$$S = \sum_{0 \leq K \leq n} (a + bK)$$

$$= \sum_{0 \leq n-K \leq n} (a + b(n-K)) = \sum_{0 \leq K \leq n} (a + b(n-K))$$

$$2S = \sum_{0 \leq K \leq n} (a + bK + a + b(n-K))$$

$$= (2a + bn) \sum_{0 \leq K \leq n} 1 = (2a + bn)(n+1) \Rightarrow S = (a + \frac{1}{2}bn)(n+1)$$

$$= \sum_{0 \leq K \leq n} (2a + bn)$$

Lect. 6: Summation factors and manipulation of sums

Sec. B

$$T_0 = 0, T_n = 2T_{n-1} + 1$$

$$a_n T_n = b_n T_{n-1} + c_n \Rightarrow s_n a_n T_n = s_n b_n T_{n-1} + s_n c_n$$

$$[s_n b_n = s_{n-1} a_{n-1}] \Rightarrow s_n a_n T_n = s_{n-1} a_{n-1} T_{n-1} + s_n c_n$$

$$[s_n a_n T_n = \tilde{a}_n] \Rightarrow \tilde{a}_n = \tilde{a}_{n-1} + s_n c_n$$

$$= \tilde{a}_{n-2} + s_{n-1} c_{n-1} + s_n c_n$$

!

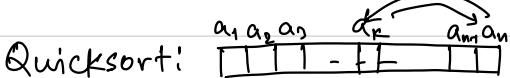
$$= \tilde{a}_0 + s_1 c_1 + s_2 c_2 + \dots + s_n c_n = \tilde{a}_0 + \sum_{k=1}^n s_k c_k$$

$$\begin{aligned} s_n a_n T_n &= s_0 a_0 T_0 + \sum_{k=1}^n s_k c_k = s_0 b_0 T_0 + \sum_{k=1}^n s_k c_k \\ \Rightarrow T_n &= \frac{1}{s_n a_n} (s_0 b_0 T_0 + \sum_{k=1}^n s_k c_k) \end{aligned}$$

$$s_n b_n = s_{n-1} a_{n-1} \Rightarrow s_n = \frac{a_{n-1}}{b_n} \times s_{n-1} = \frac{a_{n-1}}{b_n} \times \frac{a_{n-2}}{b_{n-1}} \times s_{n-2} = \dots = \frac{a_{n-1} \dots a_1}{b_n \dots b_2} s_1$$

$$\text{To H1: } T_n = 2T_{n-1} + 1 \quad a_n = 1 \quad s_n = \frac{1 \times 1 \dots \times 1}{2 \times 2 \dots \times 2} s_1 = \frac{1}{2^{n-1}} \times \frac{1}{2} = \frac{1}{2^n}$$

$$a_n T_n = b_n T_{n-1} + c_n \quad \begin{matrix} b_n = 2 \\ c_n = 1 \end{matrix}$$



$$C_0 = 0 \quad a_i < a_k > a_j$$

$$C_n = (n+1) + C_0 + C_{n-1}$$

$$C_n = (n+1) + C_1 + C_{n-2}$$

!

$$\underline{C_n = (n+1) + C_{n-1} + C_0}$$

$$n C_n = n(n+1) + 2 \sum_{k=0}^{n-1} C_k \longrightarrow (n-1) C_{n-1} = n(n-1) + 2 \sum_{k=0}^{n-2} C_k$$

$$\Rightarrow C_n = (n+1) + 2n \sum_{k=0}^{n-1} C_k \Rightarrow n C_n - (n-1) C_{n-1} = 2n + 2 C_{n-1}$$

$$\Rightarrow n C_n = (n+1) C_{n-1} + 2n$$

$$a_n^{\uparrow} T_n = b_n^{\uparrow} T_{n-1} + c_n^{\uparrow}$$

$$s_n = \frac{a_1 a_2 \dots a_n}{b_1 b_2 \dots b_n} = \frac{(n-1)(n-2) \dots 2 \cdot 1}{(n+1)n(n-1) \dots 3} = \frac{2}{n(n+1)}$$

$$C_n = \pi(n+1)/2 \times \sum_{k=1}^n \frac{2}{k(k+1)} \times 2 \cancel{\times} = 2(n+1) \sum_{k=1}^n \frac{1}{k+1}$$

$$H_n = 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} = \sum_{k=1}^n \frac{1}{k} \quad [\text{Harmonic number}]$$

$$\begin{aligned} \sum_{k=1}^n \frac{1}{k+1} &= \sum_{1 \leq k \leq n} \frac{1}{k+1} = \sum_{1 \leq k \leq n} \frac{1}{k+1} = \sum_{2 \leq k \leq n+1} \frac{1}{k} = \sum_{2 \leq k \leq n} \frac{1}{k} + \frac{1}{n+1} \\ &= \sum_{1 \leq k \leq n} \frac{1}{k} + \frac{1}{n+1} - 1 = H_n - \frac{n}{n+1} \quad \lim_{n \rightarrow \infty} H_n = \ln n + \gamma \begin{matrix} [0.572] \\ \text{a const} \end{matrix} \end{aligned}$$

$$C_n = 2(n+1)(H_n - \frac{n}{n+1}) = 2(n+1)H_n - 2n = O(n \lg n)$$

Manipulation of sums: \mathbb{K} is a finite set of integers

$$1) \sum_{k \in \mathbb{K}} c a_k = c \sum_{k \in \mathbb{K}} a_k \quad (\text{distributive law})$$

$$2) \sum_{k \in \mathbb{K}} (a_k + b_k) = \sum_{k \in \mathbb{K}} a_k + \sum_{k \in \mathbb{K}} b_k \quad (\text{associative law})$$

$$3) \sum_{k \in \mathbb{K}} a_k = \sum_{p(k) \in \mathbb{K}} a_{p(k)} \quad (\text{commutative law}) \quad p(k) \text{ is a permutation over } \mathbb{Z}$$

$$\mathbb{K} = \{-1, 0, 1\} \quad p(k) = -k \quad \begin{matrix} -1 \rightarrow 1 \\ 0 \rightarrow 0 \\ 1 \rightarrow -1 \end{matrix} \quad a_{-1} + a_0 + a_1 = a_1 + a_0 + a_{-1}$$

$$S = \sum_{0 \leq k \leq n} (a + b k) \quad p(k) = n - k \quad \begin{matrix} 0 \rightarrow n \\ \vdots \rightarrow n-1 \\ n \rightarrow 0 \end{matrix} \quad S = \sum_{0 \leq n-k \leq n} (a + b(n-k)) = \sum_{0 \leq k \leq n} (a + b(n-k))$$

$$2S = \sum_{0 \leq k \leq n} (a + b k) + \sum_{0 \leq k \leq n} (a + b(n-k)) \quad \begin{matrix} 0 \leq n-k \leq n \\ \Rightarrow k \leq n \& n-k \leq n \\ \Rightarrow 0 \leq k \leq n \end{matrix}$$

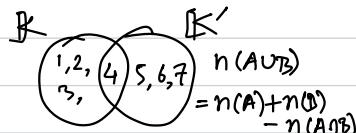
$$\begin{aligned} &= \sum_{0 \leq k \leq n} (a + b k + a + b n - b k) = \sum_{0 \leq k \leq n} (2a + bn) = (2a + bn) \sum_{0 \leq k \leq n} 1 \\ &= (2a + bn)(n+1) \end{aligned}$$

$$\Rightarrow S = (a + \frac{1}{2} bn)(n+1)$$

Lec. 7: Manipulation of Sums (contd.) & Multiple Sums

Sec. A

$$\sum_{K \in K} a_K + \sum_{K \in K'} a_K = \sum_{K \in K \cup K'} a_K + \sum_{K \in K \cap K'} a_K$$



$$\sum_{K=1}^m a_K + \sum_{K=m}^n a_K = \sum_{K=1}^n a_K + a_m$$

$$a_1 + a_2 + a_3 + a_4 = a_4 + a_1 + a_2 + a_3 \\ + a_4 + a_5 + a_6 + a_7 \quad + a_4 + a_5 + a_6 + a_7$$

$$S_n = \sum_{0 \leq k \leq n} a_k \rightarrow S_{n+1} = \sum_{0 \leq k \leq n+1} a_k = \sum_{0 \leq k \leq n} a_k + a_{n+1} = S_n + a_{n+1}$$

$$S_{n+1} = a_0 + \sum_{1 \leq k \leq n+1} a_k = a_0 + \sum_{1 \leq k+1 \leq n+1} a_{k+1} = a_0 + \sum_{0 \leq k \leq n} a_{k+1}$$

$$S_n + a_{n+1} = a_0 + \sum_{0 \leq k \leq n} a_{k+1} \quad [\text{Perturbation method}]$$

$$S_n = \sum_{0 \leq k \leq n} a_k x^k \Rightarrow S_n + a x^{n+1} = a + \sum_{0 \leq k \leq n} a x^{k+1} = a + x S_n$$

$$\Rightarrow S_n(1-x) = a(1-x^{n+1}) \Rightarrow S_n = \frac{a(1-x^{n+1})}{1-x}$$

=

$$S_n = \sum_{0 \leq k \leq n} k 2^k \Rightarrow S_n + (n+1) 2^{n+1} = \sum_{0 \leq k \leq n} (k+1) 2^{k+1} = \sum_{0 \leq k \leq n} k 2^{k+1} + \sum_{0 \leq k \leq n} 2^{k+1} \\ = 2S_n + \frac{2 \cdot (1-2^{n+1})}{1-2} = 2S_n + 2^{n+2} - 2$$

$$\Rightarrow S_n + (n+1) 2^{n+1} = 2S_n + 2^{n+2} - 2 \Rightarrow S_n = (n+1) 2^{n+1} - 2 \cdot 2^{n+1} + 2 \\ = (n-1) 2^{n+1} + 2$$

$$\sum_{K=0}^n K x^K = \frac{x - (n+1)x^{n+1} + nx^{n+2}}{(1-x)^2} \quad [x \neq 1]$$

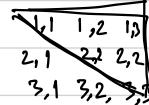
Multiple Sums: $\sum_{1 \leq j, k \leq 3} a_j b_k = a_1 b_1 + a_1 b_2 + a_1 b_3 \quad (a'_1)$
 $+ a_2 b_1 + a_2 b_2 + a_2 b_3 \quad (a'_2) \rightarrow (a_1 + a_2 + a_3)$
 $+ a_3 b_1 + a_3 b_2 + a_3 b_3 \quad (a'_3) \quad (b_1 + b_2 + b_3)$

$$\sum_{P(j,k)} a_{j,k} = \sum_{j,k} a_{j,k} [P(j,k)] \quad b'_1 \quad b'_2 \quad b'_3 \quad s$$

$$\sum_j \left(\sum_k a_{j,k} [P(j,k)] \right) = \sum_k \left(\sum_j a_{j,k} P[j,k] \right) \text{ [Interchanging order of summation]}$$

$$\begin{aligned} \sum_{1 \leq j, k \leq 3} a_j b_k &= \sum_{j, k} a_j b_k [1 \leq j, k \leq 3] = \sum_{j, k} a_j b_k [1 \leq j \leq 3] [1 \leq k \leq 3] \\ &= \sum_j \left(\sum_k a_j b_k [1 \leq j \leq 3] [1 \leq k \leq 3] \right) \\ &= \sum_j (a_j [1 \leq j \leq 3] \sum_k b_k [1 \leq k \leq 3]) \\ &= \sum_j a_j [1 \leq j \leq 3] (\sum_{k=1}^3 b_k) = \sum_{k=1}^3 b_k (\sum_j a_j [1 \leq j \leq 3]) \\ &= \sum_{k=1}^3 b_k \times \sum_{j=1}^3 a_j = (\sum_{j=1}^3 a_j) (\sum_{k=1}^3 b_k) \end{aligned}$$

$$\sum_{\substack{j \in J \\ k \in K}} = (\sum_{j \in J} a_j) (\sum_{k \in K} b_k)$$



$$\sum_{j \in J} \sum_{k \in K} a_{j,k} = \sum_{\substack{j \in J \\ k \in K}} a_{j,k} = \sum_{k \in K} \sum_{j \in J} a_{j,k}$$

$$\sum_{j \in J} \sum_{k \in K(j)} a_{j,k} = \sum_{k \in K} \sum_{j \in J'(k)} a_{j,k}$$

$$\begin{aligned} \Rightarrow \sum_j \sum_k a_{j,k} [j \in J] [k \in K(j)] &= \sum_k \sum_j a_{j,k} [k \in K'] [j \in J'(k)] \\ \Rightarrow [j \in J] [k \in K(j)] &= [k \in K'] [j \in J'(k)] \end{aligned}$$

Lect. 7: Manipulation of Sums (contd.) & Multiple Sums

$$\sum_{K \in K} a_K + \sum_{K \in K'} a_K = \sum_{K \in K \cup K'} a_K + \sum_{K \in K \cap K'}$$



$$(a_1 + a_2 + a_3 + a_4) + (a_4 + a_5 + a_6 + a_7) = a_4 + a_1 + a_2 + a_3 + \dots + a_7 \quad n(A \cup B) = n(A) + n(B) - n(A \cap B)$$

$$\Rightarrow n(A) + n(B) = n(A \cap B) + n(A \cup B)$$

$$\sum_{K=1}^m a_K + \sum_{K=m}^n a_K = \sum_{K=1}^n a_K + a_m$$

$$\sum_{0 \leq K \leq n} a_K = a_0 + \sum_{1 \leq K \leq n} a_K = \sum_{0 \leq K \leq n} a_K + a_n$$

$$S_n = \sum_{0 \leq K \leq n} a_K \Rightarrow S_{n+1} = \sum_{0 \leq K \leq n+1} a_K = \sum_{0 \leq K \leq n} a_K + a_{n+1} = S_n + a_{n+1}$$

$$= a_0 + \sum_{1 \leq K \leq n+1} a_K = a_0 + \sum_{1 \leq K+1 \leq n+1} a_{K+1} = a_0 + \sum_{0 \leq K \leq n} a_{K+1}$$

[Perturbation method]

$$\Rightarrow S_n + a x^{n+1} = a + \sum_{0 \leq K \leq n} a x^{K+1} = a + x \sum_{0 \leq K \leq n} a x^K = a + x S_n$$

$$\Rightarrow S_n(1-x) = a(1-x^{n+1}) \Rightarrow S_n = \frac{a(1-x^{n+1})}{(1-x)}$$

=

$$S_n = \sum_{0 \leq K \leq n} K 2^K \Rightarrow S_n + (n+1) 2^{n+1} = \sum_{0 \leq K \leq n} (K+1) 2^{K+1} = \sum_{0 \leq K \leq n} K 2^{K+1} + \sum_{0 \leq K \leq n} 2^{K+1}$$

$$= 2 \sum_{0 \leq K \leq n} K 2^K + \frac{2(1-2^{n+1})}{1-2} = \underline{2S_n + 2^{n+2} - 2}$$

$$\Rightarrow S_n = (n+1) 2^{n+1} - 2 \cdot 2^{n+1} + 2 = (n-1) 2^{n+1} + 2$$

$$\sum_{K=0}^n K x^K = \frac{x - (n+1)x^{n+1} + nx^{n+2}}{(1-x)^2}; \quad x \neq 1$$

$$\text{Multiple Sums: } \sum_{1 \leq j, k \leq 3} a_j b_k = a_1 b_1 + a_1 b_2 + a_1 b_3 \quad (a'_1)$$

$$+ a_2 b_1 + a_2 b_2 + a_2 b_3 \quad (a'_2) = (a_1 + a_2 + a_3)(b_1 + b_2 + b_3)$$

$$+ a_3 b_1 + a_3 b_2 + a_3 b_3 \quad (a'_3)$$

$$\sum_{P(j,k)} a_{j,k} = \sum_{j,k} a_{j,k} [P(j,k)] \quad (b'_1) \quad (b'_2) \quad (b'_3) \quad s$$

$$\sum_j \sum_k a_{j,k} [P(j,k)] = \sum_j \left(\sum_k a_{j,k} [P(j,k)] \right) = \sum_k \left(\sum_j a_{j,k} [P(j,k)] \right)$$

$$\sum_{1 \leq j, k \leq 3} a_j b_k = \sum_{j,k} a_j b_k [1 \leq j, k \leq 3] = \sum_{j,k} a_j b_k [1 \leq j \leq 3] [1 \leq k \leq 3]$$

$$= \sum_j \left(\sum_k a_j b_k [1 \leq j \leq 3] [1 \leq k \leq 3] \right)$$

$$= \sum_j a_j [1 \leq j \leq 3] \left(\sum_k b_k [1 \leq k \leq 3] \right) = \sum_j a_j [1 \leq j \leq 3] \left(\sum_{k=1}^3 b_k \right)$$

$$= \left(\sum_{k=1}^3 b_k \right) \left(\sum_j a_j [1 \leq j \leq 3] \right) = \left(\sum_{k=1}^3 b_k \right) \left(\sum_{j=1}^3 a_j \right) = \left(\sum_{j=1}^3 a_j \right) \left(\sum_{k=1}^3 b_k \right)$$

$$\sum_{\substack{j \in J \\ k \in K}} a_j b_k = \left(\sum_{j \in J} a_j \right) \left(\sum_{k \in K} b_k \right)$$

$$\sum_{j \in J} \sum_{k \in K} a_{j,k} = \sum_{j \in J} a_{j,k} = \sum_{k \in K} \sum_{j \in J} a_{j,k}$$

$$\sum_{j \in J} \sum_{k \in K(j)} a_{j,k} = \sum_{k \in K'} \sum_{j \in J'(k)} a_{j,k}$$

$$\Leftrightarrow \sum_j \sum_k a_{j,k} [j \in J] [k \in K(j)] = \sum_k \sum_j a_{j,k} [k \in K'] [j \in J'(k)]$$

$$\Leftrightarrow [j \in J] [k \in K(j)] = [k \in K'] [j \in J'(k)]$$

	1, 2	1, 3
2, 1	2, 2	2, 3
3, 1	3, 2	3, 3

Lect. 8: Multiple Sums contd.

Sec. B

$$[1 \leq j \leq n] [j \leq k \leq n] = [1 \leq j \leq k \leq n] = [1 \leq k \leq n] [1 \leq j \leq k]$$

$$\begin{bmatrix} a_1 a_1 & a_1 a_2 & \cdots & a_1 a_n \\ a_2 a_1 & a_2 a_2 & \cdots & a_2 a_n \\ \vdots & \vdots & \ddots & \vdots \\ a_n a_1 & a_n a_2 & \cdots & a_n a_n \end{bmatrix}$$

$$\int_0^1 \int_0^y f(x, y) dx dy$$

$$\int_0^1 \int_x^1 f(x, y) dy dx$$

$$S_{\nabla} = \sum_{1 \leq j \leq k \leq n} a_j a_k = \sum_{1 \leq k \leq j \leq n} a_k a_j = \sum_{1 \leq k \leq j \leq n} a_j a_k = S_{\Delta}$$

$$[1 \leq j \leq k \leq n] + [1 \leq k \leq j \leq n] = [1 \leq j, k \leq n] + [1 \leq j = k \leq n]$$

$$2S_{\nabla} = S_{\nabla} + S_{\Delta} = \sum_{1 \leq j, k \leq n} a_j a_k + \sum_{1 \leq j = k \leq n} a_j a_k$$

$$= (\sum_{j=1}^n a_j) (\sum_{k=1}^n a_k) + \sum_{1 \leq j \leq n} a_j^2$$

$$= (\sum_{j=1}^n a_j)^2 + \sum_{j=1}^n a_j^2$$

$$S = \sum_{1 \leq j < k \leq n} (a_k - a_j)(b_k - b_j) = \sum_{1 \leq k < j \leq n} (a_j - a_k)(b_j - b_k) = \sum_{1 \leq k < j \leq n} (a_k - a_j)(b_k - b_j)$$

$$[1 \leq j < k \leq n] + [1 \leq k < j \leq n] = [1 \leq j, k \leq n] - [1 \leq j = k \leq n]$$

$$2S = \sum_{1 \leq j, k \leq n} (a_k - a_j)(b_k - b_j) - \sum_{1 \leq j \neq k \leq n} (a_k - a_j)(b_k - b_j)$$

$$= \sum_{1 \leq j, k \leq n} a_k b_k - \sum_{1 \leq j, k \leq n} a_j b_k - \sum_{1 \leq j, k \leq n} a_k b_j + \sum_{1 \leq j, k \leq n} a_j b_j$$

$$= 2 \sum_{1 \leq j, k \leq n} a_k b_k - 2 \sum_{1 \leq j, k \leq n} a_j b_k$$

$$\sum_{1 \leq j, k \leq n} a_k b_k = \sum_{1 \leq k \leq n} \sum_{1 \leq j \leq n} a_k b_k = \sum_{1 \leq k \leq n} a_k b_k \sum_{1 \leq j \leq n} 1 = \sum_{1 \leq k \leq n} a_k b_k n = n \sum_{1 \leq k \leq n} a_k b_k$$

$$2s = 2n \sum_{1 \leq k \leq n} a_k b_k - 2 \left(\sum_{k=1}^n a_k \right) \left(\sum_{k=1}^n b_k \right)$$

$$\left(\sum_{k=1}^n a_k \right) \left(\sum_{k=1}^n b_k \right) = n \sum_{k=1}^n a_k b_k - \sum_{1 \leq j < k \leq n} (a_k - a_j)(b_k - b_j) \Rightarrow y - x = z \geq 0 \Rightarrow x \leq y$$

$$\left(\sum_{k=1}^n a_k \right) \left(\sum_{k=1}^n b_k \right) \leq n \sum_{k=1}^n a_k b_k ; \text{ if } a_1 \leq \dots \leq a_n ; b_1 \leq \dots \leq b_n$$

$$\left(\sum_{k=1}^n a_k \right) \left(\sum_{k=1}^n b_k \right) \geq n \sum_{k=1}^n a_k b_k ; \text{ if } a_1 \leq \dots \leq a_n ; b_1 \geq \dots \geq b_n$$

\hookrightarrow Chebyshew's monotonic inequalities.

$$y - x = z \leq 0 \Rightarrow x \geq y$$

$$S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j} [S_1 = 0, S_2 = 1, S_3 = S_2]$$

$$= \sum_{1 \leq k \leq n} \sum_{1 \leq j < k} \frac{1}{k-j}$$

$$= \sum_{1 \leq j \leq n} \sum_{j < k \leq n} \frac{1}{k-j}$$

$$= \sum_{1 \leq k \leq n} \sum_{1 \leq k-j < k} \frac{1}{j} [j < k-j] \quad [1 \leq k-j \Rightarrow j < k-1]$$

$$= \sum_{1 \leq j \leq n} \sum_{j < k \leq n} \frac{1}{k} [k < k+j]$$

$$= \sum_{1 \leq k \leq n} \sum_{0 < j \leq k} \frac{1}{j} \quad k-j < k \Rightarrow 0 < j$$

$$= \sum_{1 \leq j \leq n} \sum_{0 < k \leq n-j} \frac{1}{k}$$

$$= \sum_{1 \leq k \leq n} H_{k-1}$$

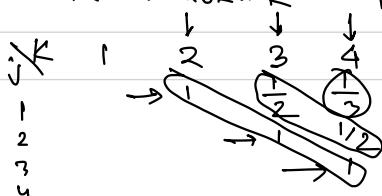
$$= \sum_{1 \leq j \leq n} H_{n-j}$$

$$= \sum_{0 \leq k \leq n} H_k [k < k+1]$$

$$= \sum_{1 \leq n-j \leq n} H_j = \sum_{0 \leq j \leq n} H_j$$

$$S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j} = \sum_{1 \leq j < k < j \leq n} \frac{1}{k} [k < k+j]$$

$$= \sum_{1 \leq k \leq n} \sum_{1 \leq j \leq n-k} \frac{1}{k} = \sum_{1 \leq k \leq n} \frac{n-k}{k} = \sum_{1 \leq k \leq n} \frac{n}{k} - \sum_{1 \leq k \leq n} 1 = nH_n - n = \sum_{0 \leq k \leq n} H_k$$



$$\hookrightarrow \frac{4-1}{1} + \frac{4-2}{2} + \frac{4-3}{3}$$

Lect. 8: Multiple Sums contd.

Sec A

$$[1 \leq j \leq n] [j \leq k \leq n] = [1 \leq j \leq k \leq n] = [1 \leq k \leq n] [1 \leq j \leq k]$$

$$\left[\begin{matrix} a_{1,1} & a_{1,2} & \dots & a_{1,n} \\ a_{2,1} & a_{2,2} & \dots & a_{2,n} \\ \vdots & \ddots & & \vdots \\ a_{n,1} & a_{n,2} & \dots & a_{n,n} \end{matrix} \right] \sum_{j=1}^n \sum_{k=j}^n a_{j,k} = \sum_{1 \leq j \leq n} \sum_{k=j}^n a_{j,k} \int_0^1 \int_0^y f(x,y) dx dy$$

$$\int_0^1 \int_x^1 f(x,y) dy dx$$

$$S_{\nabla} = \sum_{1 \leq j \leq k \leq n} a_{j,k} = \sum_{1 \leq k \leq j \leq n} a_{k,j} = \sum_{1 \leq k \leq j \leq n} a_{j,k} = S_{\Delta}$$

$$2S_{\nabla} = S_{\nabla} + S_{\Delta} = \sum_{1 \leq j, k \leq n} a_{j,k} + \sum_{1 \leq j=k \leq n} a_{j,k}$$

$$\sum_{1 \leq j, k \leq n} a_{j,k} = (\sum_{j=1}^n a_j) (\sum_{k=1}^n a_k) = (\sum_{j=1}^n a_j)^2, \sum_{1 \leq j \neq k \leq n} a_{j,k} = \sum_{k \neq j} a_j^2 = \sum_{j=1}^n a_j^2$$

$$S_{\nabla} = \frac{1}{2} ((\sum_{j=1}^n a_j)^2 + \sum_{j=1}^n a_j^2)$$

$$S = \sum_{1 \leq j < k \leq n} (a_k - a_j)(b_k - b_j) = \sum_{1 \leq k < j \leq n} (a_j - a_k)(b_j - b_k) = \sum_{1 \leq k < j \leq n} (a_k - a_j)(b_j - b_k)$$

$$2S = \sum_{1 \leq j, k \leq n} (a_k - a_j)(b_k - b_j) - \sum_{1 \leq j=k \leq n} (a_k - a_j)(b_k - b_j)$$

$$= \sum_{1 \leq j, k \leq n} a_k b_k - \sum_{1 \leq j, k \leq n} a_j b_k - \sum_{1 \leq j, k \leq n} a_k b_j + \sum_{1 \leq j, k \leq n} a_j b_j$$

$$= 2 \sum_{1 \leq j, k \leq n} a_k b_k - 2 \sum_{1 \leq j, k \leq n} a_j b_k$$

$$\sum_{1 \leq j, k \leq n} a_k b_k = \sum_{1 \leq k \leq n} \sum_{1 \leq j \leq n} a_k b_k = \sum_{1 \leq k \leq n} a_k b_k \sum_{1 \leq j \leq n} 1 = \sum_{1 \leq k \leq n} a_k b_k n = n \sum_{1 \leq k \leq n} a_k b_k$$

$$\Rightarrow S = n \sum_{k=1}^n a_k b_k - (\sum_{j=1}^n a_j)(\sum_{k=1}^n b_k)$$

$$\Rightarrow (\sum_{k=1}^n a_k)(\sum_{k=1}^n b_k) = n \sum_{k=1}^n a_k b_k - \sum_{1 \leq j < k \leq n} (a_k - a_j)(b_k - b_j)$$

$$x = y - z \Rightarrow y - x = z > 0 \Rightarrow y - x > 0 \Rightarrow x \leq y \quad (a_i, b_i \text{ monotonically increasing})$$

$$(\sum_{k=1}^n a_k)(\sum_{k=1}^n b_k) \leq n \sum_{k=1}^n a_k b_k; \quad a_1 \leq \dots \leq a_n, \quad b_1 \leq \dots \leq b_n$$

$$(\sum_{k=1}^n a_k)(\sum_{k=1}^n b_k) \geq n \sum_{k=1}^n a_k b_k; \quad a_1 \leq \dots \leq a_n, \quad b_1 \geq \dots \geq b_n$$

↳ Chebychev's monotonic inequalities

$$S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j} \quad [S_1 = 0, S_2 = 1, S_3 = \frac{5}{2}]$$

$$= \sum_{1 \leq k \leq n} \sum_{1 \leq j < k} \frac{1}{k-j}$$

$$= \sum_{1 \leq j \leq n} \sum_{j < k \leq n} \frac{1}{k-j}$$

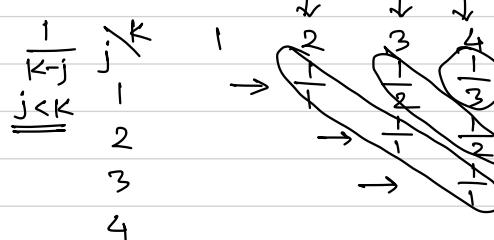
$$= \sum_{1 \leq k \leq n} \sum_{1 \leq k-j < k} \frac{1}{j} \quad [j < k-j] \quad = \sum_{1 \leq j \leq n} \sum_{j < k \leq n} \frac{1}{k} \quad [k < k+j]$$

$$= \sum_{1 \leq k \leq n} \sum_{0 < j \leq k-1} \frac{1}{j} \quad [k-j < k \Rightarrow 0 < j] \quad = \sum_{1 \leq j \leq n} \sum_{0 < k \leq n-j} \frac{1}{k} \quad [k+j \leq n \Rightarrow k \leq n-j]$$

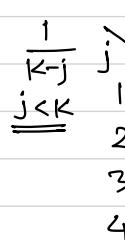
$$= \sum_{1 \leq k \leq n} H_{k-1} \quad = \sum_{1 \leq j \leq n} H_{n-j}$$

$$= \sum_{0 \leq k < n} H_k \quad = \sum_{1 \leq n-j \leq n} H_j \quad [j < n-j] = \sum_{0 \leq j < n} H_j$$

$$S_n = \sum_{1 \leq j < k \leq n} \frac{1}{k-j}$$



$$= \sum_{1 \leq j < k+j \leq n} \frac{1}{k} \quad [k < k+j]$$



$$= \sum_{1 \leq k \leq n} \sum_{1 \leq j \leq n-k} \frac{1}{k}$$

$$= \sum_{1 \leq k \leq n} \frac{1}{k} \sum_{1 \leq j \leq n-k} 1 = \sum_{1 \leq k \leq n} \frac{n-k}{k} = \sum_{1 \leq k \leq n} \frac{n}{k} - \sum_{1 \leq k \leq n} \frac{1}{k}$$

$$= n \sum_{1 \leq k \leq n} \frac{1}{k} - n = nH_n - n = \sum_{0 \leq k < n} H_k$$

$$\frac{4-1}{1} = \frac{3}{1} = 3 \times 1, \quad \frac{4-2}{2} = 2 \times \frac{1}{2}, \quad \frac{4-3}{3} = \frac{1}{3}$$

See. A

Lect. 9: General methods for sums

$$n \quad 0 \quad 1 \quad 2 \quad 3 \quad 4 \quad \dots$$

$$S_n = \sum_{0 \leq k \leq n} k^2; \quad S_n \quad 0 \quad 1 \quad 5 \quad 14 \quad 30 \quad \dots$$

Method 0: Look it up: $S_n = \frac{n(n+1)(2n+1)}{6} \rightarrow$

Method 1: Guess the answer; prove by induction

$$S_n = an^3 + bn^2 + cn + d; \quad 0 = d, \quad 1 = a+b+c$$

$$5 = 8a + 4b + 2c \quad (a, b, c) = \left(\frac{1}{3}, \frac{1}{2}, \frac{1}{6}\right)$$

$$S_n = \frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{1}{6}n$$

$$14 = 27a + 9b + 2c$$

$$= n/6(2n^2 + 3n + 1) = \frac{1}{6}(n+1)(2n+1) \rightarrow \text{prove by induction}$$

Method 2: Perturbation

$$\sum_{k=0}^{n+1} (k+1)^2 = 0 + \sum_{0 \leq k \leq n} (k+1)^2 = \sum_{0 \leq k \leq n} (k^2 + 2k + 1) = \sum_{0 \leq k \leq n} k^2 + 2 \sum_{0 \leq k \leq n} k + \sum_{0 \leq k \leq n} 1$$

$$\Rightarrow 2 \sum_{0 \leq k \leq n} k = (n+1)^2 - (n+1)$$

$$= S_n + 2 \sum_{0 \leq k \leq n} k + (n+1)$$

$$C_n = \sum_{0 \leq k \leq n} k^3; \quad C_n + (n+1)^3 = 0 + \sum_{0 \leq k \leq n} (k+1)^3 = \sum_{0 \leq k \leq n} (k^3 + 3k^2 + 3k + 1)$$

$$= \sum_{0 \leq k \leq n} k^3 + 3 \sum_{0 \leq k \leq n} k^2 + 3 \sum_{0 \leq k \leq n} k + \sum_{0 \leq k \leq n} 1$$

$$= C_n + 3S_n + 3 \cdot \frac{n(n+1)}{2} + (n+1)$$

Method 3: Repertoire Method

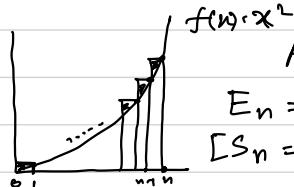
$$R_0 = \alpha$$

$$R_n = R_{n-1} + \beta + \gamma n + \delta n^2 \quad R(n) = A(n)\alpha + B(n)\beta + C(n)\gamma + D(n)\delta$$

$$R_n = 1, n, n^2, \underline{n^3}$$

Method 4: Replace sums by integrals

$$S_n = \sum_{0 \leq k \leq n} k^2 = 1^2 + 2^2 + 3^2 + \dots = 1 \cdot 1 + 1 \cdot 4 + 1 \cdot 9 + \dots + 1 \cdot n^2$$



$$\text{Area under the curve} = \int_0^n x^2 dx = \frac{1}{3} n^3$$

$$E_n = S_n - \frac{1}{3} n^3 = S_{n-1} + n^2 - \frac{1}{3} n^3 = E_{n-1} + \frac{1}{3} (n-1)^3 + n^2 - \frac{1}{3} n^3 \\ [S_n = E_n + \frac{1}{3} n^3] \rightarrow \\ = E_{n-1} + \frac{1}{3} n^3 - n^2 + n - \frac{1}{3} + \frac{n^2}{3}$$

$$S_n - \int_0^n x^2 dx = \sum_{k=1}^n k^2 - \sum_{k=1}^n \left(\int_{k-1}^k x^2 dx \right)$$

$$= \sum_{k=1}^n \left(k^2 - \int_{k-1}^k x^2 dx \right) = \sum_{k=1}^n \left(k^2 - \frac{k^3 - (k-1)^3}{3} \right) \\ = \sum_{k=1}^n \left(k - \frac{1}{3} \right)$$

Method 5: Expand and contract

$$[1 \leq j \leq n] [j \leq k \leq n] = [1 \leq j \leq k \leq n] = [1 \leq k \leq n]$$

$$S_n = \sum_{1 \leq k \leq n} k^2 = \sum_{1 \leq k \leq n} k \cdot k = \sum_{1 \leq k \leq n} k \cdot \sum_{1 \leq j \leq k} 1$$

$[1 \leq j \leq k]$

$$= \sum_{1 \leq k \leq n} \sum_{1 \leq j \leq k} k = \sum_{1 \leq j \leq n} k = \sum_{1 \leq j \leq n} \sum_{j \leq k \leq n} k$$

$$= \sum_{1 \leq j \leq n} \left(\sum_{k=1}^n k - \sum_{k=1}^{j-1} k \right)$$

$j + (j+1) + (j+2) + \dots + n$

$$= \sum_{1 \leq j \leq n} \left(\frac{n(n+1)}{2} - \frac{j(j-1)}{2} \right)$$

$$= 1 + 2 + \dots + (j-1) + j + \dots + n - (1 + 2 + \dots + (j-1))$$

$$= \frac{1}{2} n^2(n+1) - \frac{1}{2} \sum_{1 \leq j \leq n} (j^2 - j) = \frac{1}{2} n^2(n+1) - \frac{1}{2} \sum_{1 \leq j \leq n} j^2 + \frac{1}{2} \sum_{1 \leq j \leq n} j$$

$$= \frac{1}{2} n^2(n+1) - \frac{1}{2} S_n + \frac{1}{2} \times \frac{1}{2}(n+1)n$$

$$\Rightarrow \frac{3}{2} S_n = \frac{1}{2} n^2(n+1) + \frac{1}{4} n(n+1) \Rightarrow S_n = \frac{n(n+1)(2n+1)}{6}$$

Lect. 9: General methods for sums

Sec. 13

$$S_n = \sum_{0 \leq k \leq n} k^2 \quad \text{1st Strategy: } S_n \begin{array}{c} n \\ 0 & 1 & 2 & 3 & 4 \end{array} \quad \begin{array}{c} 0 & 1 & 5 & 14 & 30 \end{array}$$

Method 0: Look it up

Method 1: Guess the answer; prove by induction

$$S_n = an^3 + bn^2 + cn + d; \quad 0 = d; \quad a+b+c = 1 \quad (a, b, c) = \left(\frac{1}{3}, \frac{1}{2}, \frac{1}{6}\right)$$

$$8a+4b+2c = 5$$

$$27a+9b+3c = 14$$

$$S_n = \frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{1}{6}n$$

$$= \frac{1}{6}n(2n^2 + 3n + 1) = \frac{1}{6}n(n+1)(2n+1) \rightarrow \text{induction}$$

Method 2: Perturbation

$$S_n + (n+1)^2 = \sum_{0 \leq k \leq n} (k+1)^2 = \sum_{0 \leq k \leq n} (k^2 + 2k + 1) = \sum_{0 \leq k \leq n} k^2 + 2 \sum_{0 \leq k \leq n} k + \sum_{0 \leq k \leq n} 1$$

$$= S_n + 2 \sum_{0 \leq k \leq n} k + (n+1)$$

$$\Rightarrow 2 \sum_{0 \leq k \leq n} k = (n+1)^2 - (n+1)$$

$$C_n = \sum_{0 \leq k \leq n} k^3; \quad C_n + (n+1)^3 = \sum_{0 \leq k \leq n} (k+1)^3 = \sum_{0 \leq k \leq n} k^3 + 3 \sum_{0 \leq k \leq n} k^2 + 3 \sum_{0 \leq k \leq n} k + \sum_{0 \leq k \leq n} 1$$

$$= C_n + 3S_n + \frac{3n(n+1)}{2} + (n+1)$$

$$\Rightarrow S_n = n(n+1)(2n+1)/6$$

Method 3: Repertoire

$$R_0 = \alpha$$

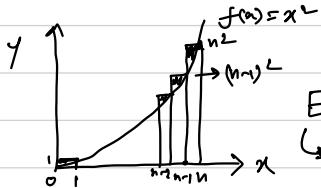
$$R_n = R_{n-1} + \underbrace{\beta}_{\alpha_n} + \gamma n + \delta n^2$$

$$R_n = A(n)\alpha + B(n)\beta + C(n)\gamma + D(n)\delta$$

$$R_n = 1, n, n^2, n^3$$

Method 4: Replace sums by integrals

$$S_n = \sum_{1 \leq k \leq n} k^2 = 1^2 + 2^2 + \dots + n^2 = 1 \cdot 1 + 1 \cdot 4 + 1 \cdot 9 + \dots + 1 \cdot (n-1)^2 + 1 \cdot n^2$$



$$\text{Area under curve} = \int_0^n x^2 dx = \frac{1}{3} n^3$$

$$E_n = S_n - \frac{1}{3} n^3 = S_{n-1} + n^2 - \frac{1}{3} n^3$$

$$\hookrightarrow S_n = E_n + \frac{1}{3} n^3 \quad | = E_{n-1} + \frac{1}{3} (n-1)^3 + n^2 - \frac{1}{3} n^3$$

$$= E_{n-1} + \frac{1}{3} n^3 - n^2 + n - \frac{1}{3} + n^2 - \frac{1}{3} n^3$$

$$= E_{n-1} + \left(n - \frac{1}{3} \right).$$

$$E_n = S_n - \int_0^n x^2 dx$$

$$= \sum_{k=1}^n k^2 - \int_0^n x^2 dx$$

$$= \sum_{k=1}^n k^2 - \sum_{k=1}^n \left(\int_{k-1}^k x^2 dx \right) = \sum_{k=1}^n \left(k^2 - \int_{k-1}^k x^2 dx \right)$$

$$= \sum_{k=1}^n \left(k^2 - \frac{k^3 - (k-1)^3}{3} \right)$$

$$= \sum_{k=1}^n \left(k - \frac{1}{3} \right)$$

Method 5: Expand and contract

$$S_n = \sum_{1 \leq k \leq n} k^2 = \sum_{1 \leq k \leq n} k \cdot k = \sum_{1 \leq k \leq n} k \cdot \sum_{1 \leq j \leq k} 1 = \sum_{1 \leq k \leq n} \sum_{1 \leq j \leq k} k$$

$$\hookrightarrow [1 \leq j \leq k \leq n] = [1 \leq k \leq n][1 \leq j \leq k] \quad | = \sum_{1 \leq j \leq k \leq n} k$$

$$\hookrightarrow [1 \leq j \leq n][j \leq k \leq n] \quad | = \sum_{1 \leq j \leq n} \sum_{j \leq k \leq n} k \quad \begin{cases} j+(j+1)+\dots+n \\ = 1+2+\dots+(j-1)+j+\dots+n \\ -(1+2+\dots+(j-1)) \end{cases}$$

$$= \sum_{1 \leq j \leq n} \left(\sum_{k=1}^n k - \sum_{k=1}^{j-1} k \right)$$

$$= \sum_{1 \leq j \leq n} \left(\frac{n(n+1)}{2} - \frac{j(j-1)}{2} \right) = \frac{n^2(n+1)}{2} - \frac{1}{2} \sum_{1 \leq j \leq n} (j^2 - j)$$

$$= \frac{n^2(n+1)}{2} - \frac{1}{2} \sum_{1 \leq j \leq n} j^2 + \frac{1}{2} \sum_{1 \leq j \leq n} j = \frac{n^2(n+1)}{2} - \frac{1}{2} S_n + \frac{1}{2} \cdot \frac{n(n+1)}{2}$$

$$\Rightarrow \frac{3}{2} S_n = \frac{n^2(n+1)}{2} - \frac{n(n+1)}{4} \Rightarrow S_n = \frac{n(n+1)(2n+1)}{6}$$

Lect. 10: Number Theory (Divisibility)

m divides n (or n is divisible by m) [2 divides 6; 6 is divisible by 2]
 if $m > 0$ and n/m is an integer.
 $m \mid n \Leftrightarrow m > 0 \ \& \ n = mk$ for some integer k .

gcd (greatest common divisor): gcd of m, n is the largest integer that divides both m and n .

$$\text{gcd}(m, n) = \max \{ k \mid k \mid m \ \& \ k \mid n \}$$

$$m=0, n>0 \Rightarrow \text{gcd}(m, n) = \text{gcd}(0, n) = n \quad \frac{1}{12} + \frac{1}{18} = \frac{2}{36} + \frac{2}{36} = \frac{2}{36}$$

lcm (least common multiple): $\text{lcm}(m, n) = \min \{ k \mid k > 0, m \mid k \ \& \ n \mid k \}$

Euclid's algorithm: $\text{gcd}(m, n)$, for $0 \leq m < n$

$$\text{gcd}(0, n) = n$$

$$\text{gcd}(m, n) = \text{gcd}(n \text{ mod } m, m)$$

$\begin{matrix} \text{divisor} \\ \downarrow d \\ \text{dividend} \\ \downarrow \end{matrix}$

$\begin{matrix} \text{quotient} \\ \downarrow q \\ \text{remainder} \\ \downarrow r \end{matrix}$

$\begin{matrix} 12 & | & 18 & | \\ & | & 12 & | \\ & 6 & | & 12 & | \\ & & 0 & & \end{matrix}$

$[0 \leq r < d]$

$$a = dq + r \rightarrow a \text{ mod } d$$

$$n = m \lfloor n/m \rfloor + n \text{ mod } m$$

$$\Rightarrow n \text{ mod } m = n - m \lfloor n/m \rfloor$$

Diophantine eqⁿ: $ax + by = c \quad x \in \mathbb{Z}, y \in \mathbb{Z}$

$$2x + 3y = 5$$

has solutions (x, y) iff $\text{gcd}(a, b) \mid c$

$$\begin{array}{r} m \mid n (\lfloor n/m \rfloor)^q \\ \hline n \text{ mod } m \end{array}$$

$$\frac{5}{2} = \underline{\underline{2}}\frac{1}{2}$$

Necessity: solution exists $\Rightarrow \text{gcd}(a, b) \mid c \quad \left| \begin{array}{l} \text{gcd}(a, b) = k \mid a, k \mid b \\ a = a'k, b = b'k \end{array} \right.$

$$\text{L.H.S.} = ax + by = a'kx + b'ky$$

$$= K(a'x + b'y) \Rightarrow K \mid \text{LHS} \Rightarrow K \mid \text{RHS} \Rightarrow K \mid c$$

$$\Rightarrow \text{gcd}(a, b) \mid c$$

Sufficiency (using book's notation):

Given m, n ; we need to compute $\underline{m'}, \underline{n'}$ such that

$$\underline{m'm} + \underline{n'n} = \gcd(m, n)$$

~~var^t~~ coeff ~~var^t~~ coeff

$$\gcd(0, n) = n$$

$$\gcd(m, n) = \gcd\left(\frac{n \bmod m}{r}, m\right)$$

$$\text{If } m=0, \text{ RHS} = \gcd(0, n) = \underline{n}$$

$$\text{L.H.S.} = \frac{\underline{m'm}}{0} + \frac{\underline{n'n}}{1} = n \quad \text{If } m=0, \text{ we pick } (m', n') = (0, 1)$$

$r = n \bmod m$, for r and m , we compute \bar{r} and \bar{m} such that
 $\bar{r}r + \bar{m}m = \gcd(r, m)$ [$\gcd(r, m) = \gcd(n \bmod m) = \gcd(m, n)$]

$$r = n - m \lfloor \frac{n}{m} \rfloor$$

$$\Rightarrow \bar{r}(n - m \lfloor \frac{n}{m} \rfloor) + \bar{m}m = \gcd(m, n)$$

$$\Rightarrow (\bar{m} - \bar{r} \lfloor \frac{n}{m} \rfloor)m + \bar{r}n = \gcd(m, n)$$

$$m' = \bar{m} - \lfloor \frac{n}{m} \rfloor \bar{r}, \quad n' = \frac{n}{r}$$

$$\begin{array}{r} 18 \\ \overline{) 30} \end{array} \begin{array}{r} |1 \\ 18 \\ \overline{) 12} \\ |18 \\ 12 \\ \overline{) 6} \\ |12 \\ 12 \\ \overline{) 0} \end{array}$$

$$m=18, n=12, \underline{m'}x18 + \underline{n'}x12 = \gcd(12, 18) = 6$$

$$b = 0 \times 0 + 1 \times 6 = \frac{(1-2,0)}{\bar{r}} 1 \times 6 + \frac{(0-1,1)}{\bar{m}m} 1 \times 12 = (-1) \times 12 + 1 \times 18 = 18 - 30$$

$$\underline{ax} + \underline{by} = c$$

linear combination of x & y . [Self-certifying solution]

$$k|m \& k|n \Leftrightarrow k|\gcd(m, n) \quad [\text{self-study}]$$

Leat. 11: Primes

Sec. A+B

A positive integer is prime iff it has only two divisors, 1 & p.

2, 3, 5, 7, 11, ...

[otherwise, composite]

uniquely

Any positive number can be expressed as a product of primes.

$$n = \underbrace{p_1 p_2 \cdots p_m}_{\substack{\text{may be empty}}} = \prod_{k=1}^m p_k, p_1 \leq p_2 \leq \cdots \leq p_m; [12 = 2 \times 2 \times 3, 18 = 2 \times 3 \times 3]$$

Fundamental thm of arithmetic.

Proof: Base case. $n = 1$ [empty set of product]

Prove for any $n > 1$ assuming true for all numbers $< n$.

Assume otherwise: $n = p_1 \cdots p_m = q_1' \cdots q_K'$, $p_1 \leq \cdots \leq p_m, q_1' \leq \cdots \leq q_K'$

We will prove $p_1 = q_1$. If not, say $p_1 < q_1 \Rightarrow p_1 \neq q_1 \Rightarrow \gcd(p_1, q_1) = 1$

$$\exists a, b \text{ s.t. } ap_1 + bq_1 = 1 \Rightarrow ap_1 q_2 \cdots q_K + bq_1 q_2 \cdots q_K = q_2 \cdots q_K$$

$$\Rightarrow ap_1 q_2 \cdots q_K + bn = q_2 \cdots q_K \Rightarrow p_1 \mid q_2 \cdots q_K$$

$p_1 < q_1 \leq q_2 \leq \cdots \leq q_K$, a contradiction.

$$p_2 \cdots p_m = q_2 \cdots q_K = n/p_1 < n$$

$\hookrightarrow p_2 \cdots p_m$ is eq. to $q_2 \cdots q_K$

$\Rightarrow p_1, p_2 \cdots p_m$ is eq. to $q_1, q_2 \cdots q_K$ [Proven by induction + contradiction]

$$12 = 2^2 \times 3, 18 = 2 \times 3^2, n = \prod_p p^{n_p}, n_p \geq 0 \quad [\text{Prime power factorization}]$$

$$12 = 2^2 \times 3^0 \times 5^0 \times 7^0 \times \dots$$

$$K = mn \Leftrightarrow k_p = m p + n p \quad \forall p \rightarrow 6 \mid 12$$

$$m \mid n \Leftrightarrow m_p \leq n_p$$

$$K = \gcd(m, n) \Leftrightarrow k_p = \min(m_p, n_p) \quad \forall p$$

$$K = \text{lcm}(m, n) \Leftrightarrow k_p = \max(m_p, n_p) \quad \forall p$$

$$\gcd(12, 18) = \gcd(2^2 \times 3, 2 \times 3^2) = 2^{\min(2,1)} \times 3^{\min(1,2)} = 2 \times 3 = 6$$

$$\text{lcm}(12, 18) = \text{lcm}(2^2 \times 3, 2 \times 3^2) = 2^{\max(2,1)} \times 3^{\max(1,2)} = 2^2 \times 3^2 = 36$$

Primes are infinite. Proof by contradiction.

Assume finite: P_1, P_2, \dots, P_k . $P = P_1 P_2 \dots P_k + 1$, contradiction

not divisible by any of
 P_1, P_2, \dots, P_k

Euclid number: $e_1 = 1, e_n = e_1 \dots e_{n-1} + 1. e_5 = 1807 = 13 \cdot 139$
 $\gcd(e_m, e_n) = 1; m \neq n$ [self-study]

$\hookrightarrow \gcd(1, e_m) = \gcd(0, 1) = 1$

Mersenne numbers: $2^p - 1$ is prime only if p is prime.

Proof by contradiction: $p = mK$. $a^n - b^n = (a - b)(a^{n-1} + a^{n-2}b + \dots + b^{n-1})$
 $2^{mK} - 1 = (2^m)^K - 1^K = (2^m - 1)(2^{m(K-1)} + 2^{m(K-2)} + \dots + 1) \Rightarrow$ composite

$$2^{11} - 1 = 2047 = 23 \cdot 89$$

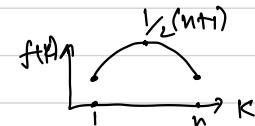
Composite number: $n! = 1 \cdot 2 \cdot 3 \cdots n = \prod_{k=1}^n k$

$$(n!)^2 = (1, 2, \dots, n)(n \cdots 2 \cdot 1) = \prod_{k=1}^n k(n+1-k)$$

$$f(k) = k(n+1-k)$$

$$n \leq k(n+1-k) \leq \frac{1}{4}(n+1)^2$$

$$\hookrightarrow n^n \leq \prod_{k=1}^n k(n+1-k) \leq \frac{1}{4^n}(n+1)^{2n}$$



$$\Rightarrow n^n \leq (n!)^2 \leq \frac{1}{4^n} (n+1)^{2n} \Rightarrow n^{n/2} \leq n! \leq \frac{1}{2^n} (n+1)^n$$

$$n! \sim \sqrt{2\pi n} (n/e)^n$$

[Stirling's approximation]

What is the highest power of \underline{p} in $\underline{n!}$? $5! = 120 = \underline{2} \times 15$

	1	2	3	4	5	6	7	8	9	$\cancel{10}$	$\cancel{10}$; power of 2
divisible by 2	X	X		X		X		X		X	$\lfloor \frac{10}{2} \rfloor = 5$
divisible by 4				X			X			X	$\lfloor \frac{10}{4} \rfloor = 2$
divisible by 8					X				X		$\lfloor \frac{10}{8} \rfloor = 1$

$$e_2(n!) = \lfloor \frac{n}{2} \rfloor + \lfloor \frac{n}{4} \rfloor + \dots = \sum_{k \geq 1} \lfloor \frac{n}{2^k} \rfloor$$

$$e_2(100) = 50 + 25 + 12 + 6 + 3 + 1 = 97 \quad \lfloor \frac{n}{2^{k+1}} \rfloor = \lfloor \lfloor \frac{n}{2^k} \rfloor / 2 \rfloor$$

$$\begin{aligned} 100 &= (1100100)_2 \\ \lfloor \frac{100}{2} \rfloor &= (110010)_2 \quad e_p(n!) = \lfloor \frac{n}{p} \rfloor + \lfloor \frac{n}{p^2} \rfloor + \dots = \sum_{k \geq 1} \lfloor \frac{n}{p^k} \rfloor \\ \lfloor \frac{100}{4} \rfloor &= (11001)_2 \end{aligned}$$

$\lfloor x \rfloor$ = largest integer not exceeding x
 $\Leftrightarrow \lfloor x \rfloor \leq x$

$$\begin{aligned} e_p(n!) &\leq \frac{n}{p} + \frac{n}{p^2} + \frac{n}{p^3} + \dots \\ &= n_p \left(1 + \frac{1}{p} + \frac{1}{p^2} + \dots\right) = \frac{n}{p} \times \frac{1}{1 - \frac{1}{p}} = \frac{n}{p} \times \frac{p}{p-1} = \frac{n}{p-1} \end{aligned}$$

$97 \leq 100$

HW: Calculate the number of trailing zeroes in $100!$

Lect. 12: Relative primality and Stern-Brocot tree

Sec. A

Two integers m, n are relatively prime (co-prime) when $\gcd(m, n) = 1$ [m 1n]
 $\frac{12}{18} = \frac{2}{3}$. A fraction $\frac{m}{n}$ is in its lowest term iff $m \perp n$ [irreducible fractions]

$$12, 18 : \gcd(12, 18) = 6, \quad \frac{12}{6} = 2, \quad \frac{18}{6} = 3. \quad \text{If } \overline{\gcd}(m, n) \perp \overline{\gcd}(m, n)$$

$$m = 2^2 \times 3^6 \times 5^6 \times 7^1 = 28; \quad n = 2^6 \times 3^1 \times 5^1 \times 7^6 = 15; \quad \gcd(m,n) = 1$$

$$m \perp n \iff \min(m_p, n_p) = 0 \iff m_p n_p = 0$$

$$k \perp m \text{ & } k \perp n \Leftrightarrow k \perp mn ; 2 \perp 3 \text{ & } 2 \perp 5 \Leftrightarrow 2 \perp 3 \times 5 = 15$$

Stern-Brocot Tree: Insert $\frac{m+m'}{n+n'}$ between two adjacent nodes $\frac{m}{n}$ & $\frac{m'}{n'}$

If $\frac{m}{n}$ & $\frac{m'}{n'}$ are consecutive fractions at any stage of the construction of the Stern-Brocot tree, then $m'n - mn' = 1$

Proof: By induction (on tree level). Base case: $\frac{m}{n} = \frac{0}{1} \Rightarrow m=0, n=1$

$$\frac{m'}{n'} = \frac{1}{0} \Rightarrow m' = 1, n' = 0$$

$$\Rightarrow m'n - mn' = 1 \cdot 1 - 0 \cdot 0 = 1$$

Hypothesis: True upto level K. \rightarrow consecutive fraction $\frac{m}{n}$ & $\frac{m'}{n'}$ [$m \neq m'$]

Induction step! Level $(k+1) \rightarrow \frac{m}{n} \rightarrow \frac{m+m'}{n+n'} \rightarrow \frac{m'}{n'}$

$$(m+m')n - m(n+n') = mn + m'n - mn - mn' = m'n - mn' = 1$$

$$m'(n+ni) - (m+ni)n = m'n + m'i^n - mn - ni^n = m'n - mn = 1$$

$$\gcd(m+m', n+n') = 1;$$

Let $\gcd(m+m', n+n') = d$. So, $d \mid m+m'$, $d \mid n+n'$

$$\Rightarrow d \mid (m+m')n, d \mid (n+n')m$$

$$\Rightarrow d \mid (m+m')n - (n+n')m$$

$$\Rightarrow d \mid mn + m'n - mn - mn' = mn' - mn$$

$$\Rightarrow d \mid mn' - mn = 1 \Rightarrow d \mid 1 \Rightarrow d = 1 \quad [\text{Proved}]$$

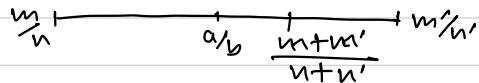
$$\frac{m}{n} < \frac{m+m'}{n+n'} < \frac{m'}{n}. \text{ Proof: } \frac{m}{n} < \frac{m+m'}{n+n'} \Leftrightarrow mn + mn' < mn + m'n \Leftrightarrow 0 < mn' - mn = 1$$

\hookrightarrow The construction preserves order; so no fraction occurs twice.

Claim: All non-negative fractions are present in the tree.

Proof (By contradiction): Assume $\frac{a}{b}$ doesn't exist in this tree. $\frac{m+m'}{n+n'}$

$$\frac{m}{n} = a_1 < \left(\frac{a}{b}\right) < \frac{1}{0} = \frac{m'}{n'}$$



if $\frac{a}{b} < \frac{m+m'}{n+n'}$; then $\frac{m'}{n'} < \frac{m+m'}{n+n'}$

if $\frac{a}{b} > \frac{m+m'}{n+n'}$; then $m/n < m+m'/n+n'$

if $\frac{a}{b} = \frac{m+m'}{n+n'}$, we are done.

$$\frac{a}{b} - \frac{m}{n} > 0$$

$$\& \frac{m'}{n'} - \frac{a}{b} > 0$$

$$\Rightarrow an - bm > 0$$

$$\& bm' - an' > 0$$

$$\Rightarrow an - bm \geq 1$$

$$\& bm' - an' \geq 1$$

$$\Rightarrow (m+n)(an - bm) > m+n \quad \& (m+n)(bm' - an') > m+n$$

$$\Rightarrow (m+n)(an - bm) + (m+n)(bm' - an') > m+n + m+n$$

$$\Rightarrow am'n + an'm - bmn' - bn'm + bm'n + bn'm - amn - an' > m+n + m+n$$

$$\Rightarrow a(m'n - mn') + b(m'n - mn') \geq m+n + m+n$$

$$\Rightarrow \underbrace{a+b}_{\text{fixed}} \geq \underbrace{m+n+m+n}_{\text{increase with every itr.}} \quad [\text{contradiction}]$$

Lect. 12: Relative Primality & Stern-Brocot tree

Sec. B

Two integers m, n are relative prime (co-prime) when $\gcd(m, n) = 1$
 $12/18 = 2/3$. A fraction $\frac{m}{n}$ is in lowest terms iff $m \perp n$ [relatively prime]

$$12/6 = 2, 18/6 = 3; \frac{m}{\gcd(m, n)} \perp \frac{n}{\gcd(m, n)}$$

$$m = 2^2 \times 3^0 \times 5^0 \times 7^1 = 28, n = 2^0 \times 3^1 \times 5^1 \times 7^0 = 15; \gcd(m, n) = 1$$

$$m \perp n \Leftrightarrow \min(m_p, n_p) = 0 \forall p. m \perp n \Leftrightarrow m_p n_p = 0 \forall p$$

$$K \perp m \text{ & } K \perp n \Leftrightarrow K \perp mn. 2 \perp 3 \text{ & } 2 \perp 5 \Leftrightarrow 2 \perp 15$$

Stern-Brocot tree: Insert $\frac{m+m'}{n+n'}$ between two adjacent fractions $\frac{m}{n}$ & $\frac{m'}{n'}$

$$\begin{array}{ccccccc} L_0 & \frac{0}{1} & - & - & - & - & \frac{1}{0} \\ L_1 & & & \frac{1}{1} & & & \\ L_2 & \frac{1}{2} & & & \frac{2}{1} & & \\ L_3 & \frac{1}{3} & \frac{2}{3} & \frac{3}{2} & \frac{3}{1} & & \end{array}$$

If $\frac{m}{n}$ & $\frac{m'}{n'}$ are two consecutive fractions at any stage of the construction of the Stern-Brocot tree, then $m'n - mn' = 1$ [Invariant]

Proof: By induction on the level K .

Base case, L_0 : $m'n - mn' = 1 \cdot 1 - 0 \cdot 0 = 1$

Hypo.: True up to L_K . $\Rightarrow \frac{m}{n}$ & $\frac{m'}{n'}$ at $L_K \Rightarrow m'n - mn' = 1$
 At $L_{K+1} \Rightarrow \frac{m}{n} \leftrightarrow \frac{m+m'}{n+n'} \leftrightarrow \frac{m'}{n'}$

$$\text{Now, } (m+m')n - m(n+n') = \cancel{mn} + m'n - \cancel{mn} - mn' = m'n - mn' = 1$$

$$\text{Again, } m'(n+n') - (m+m')n' = m'n + m'n' - mn' - \cancel{m'n'} = m'n - mn' = 1$$

$$\gcd(m+m', n+n') = 1,$$

Proof: Let $\gcd(m+m', n+n') = d$. So, $d \mid m+m'$, $d \mid n+n'$

$$\Rightarrow d \mid n(m+m')$$

$$\Rightarrow d \mid n(m+m') - m(n+n')$$

$$\Rightarrow d \mid mn + m'n - mn - mn'$$

$$\Rightarrow d \mid m'n - mn' = 1 \Rightarrow d \mid 1 \Rightarrow d = 1$$

$$\frac{m}{n} < \frac{m+m'}{n+n'} < \frac{m'}{n'}$$

$$\Leftrightarrow m(n+n') < n(m+m')$$

$$\Leftrightarrow mn + mn' < mn + m'n$$

$$\Leftrightarrow 0 < m'n - mn' = 1 \Leftrightarrow 0 < 1$$

The construction preserves order; so no fraction occurs twice.

Claim: All non-negative fractions are present in the tree.

Proof (By contradiction): Assume $\frac{a}{b}$ is absent from the tree.

$$\frac{m}{n} = \frac{0}{1} < \frac{a}{b} < \frac{m'}{n'} = \frac{1}{0}$$

if $\frac{a}{b} = \frac{m+m'}{n+n'}$, we are done.

$$\frac{m}{n} \xrightarrow{\quad} \frac{a}{b} \xrightarrow{\quad} \frac{m'}{n'}$$

if $\frac{a}{b} < \frac{m+m'}{n+n'}$, $\frac{m'}{n'} \leftarrow \frac{m+m'}{n+n'}$

if $\frac{a}{b} > \frac{m+m'}{n+n'}$, $\frac{m}{n} \rightarrow \frac{m+m'}{n+n'}$

$$\frac{a}{b} - \frac{m}{n} > 0$$

$$\& \quad \frac{m'}{n'} - \frac{a}{b} > 0$$

$$\Rightarrow an - bm > 0$$

$$\& \quad bm' - an' > 0$$

$$\Rightarrow an - bm > 1$$

$$\& \quad bm' - an' \geq 1$$

$$\Rightarrow (m+n')(an - bm) > m+n' \quad \& \quad (m+n)(bm' - an) > m+n$$

$$\Rightarrow (m+n')(an - bm) + (m+n)(bm' - an) > m+n' + m+n$$

$$\Rightarrow am'n + anm' - bmn' - bmn' + bm'n - anm' - amn' \geq m+n' + m+n$$

$$\Rightarrow a(m'n - mn') + b(m'n - mn') \geq m+n' + m+n$$

$$\Rightarrow \underbrace{a+b}_{\text{fixed}} \geq \underbrace{m+n'}_{\text{increase with every itr.}} + m+n \quad [\text{contradiction}]$$

Lect. 13: Congruences (Modular Arithmetic)

Sec. B

$$a = mq + r; \quad 0 \leq r < m \Rightarrow a = m \lfloor a/m \rfloor + a \text{ mod } m \quad (\%)$$

Defⁿ: $a \equiv b \pmod{m} \Leftrightarrow a \text{ mod } m = b \text{ mod } m$

~~#~~ $a \equiv b \pmod{m} \Leftrightarrow m \mid a - b$

1st part: $a = \lfloor a/m \rfloor m + a \text{ mod } m$

$$\stackrel{L \rightarrow R}{b = \lfloor b/m \rfloor m + b \text{ mod } m}$$

$$a - b = (\lfloor a/m \rfloor - \lfloor b/m \rfloor)m + a \text{ mod } m - b \text{ mod } m \Rightarrow m \mid a - b$$

2nd part: $m \mid a - b; \quad a = \lfloor a/m \rfloor m + r, \quad b = \lfloor b/m \rfloor m + r_2$

$$\stackrel{r \rightarrow l}{\Rightarrow a - b = (\lfloor a/m \rfloor - \lfloor b/m \rfloor)m + r_1 - r_2 \quad [0 \leq r, r_2 < m]}$$

$$\Rightarrow m \mid r_1 - r_2; \quad \max(r_1 - r_2) = m-1; \quad \min(r_1 - r_2) = -(m-1)$$

Multiples of m : ... - $3m, -2m, -m, 0, m, 2m, 3m, \dots$

$$r_1 - r_2 = 0 \Rightarrow r_1 = r_2 \Rightarrow a \text{ mod } m = b \text{ mod } m \Rightarrow a \equiv b \pmod{m}$$

- $a \equiv a \pmod{m}$ [Reflexive law] $LHS = RHS \Rightarrow LHS \equiv RHS \pmod{m}$

- $a \equiv b \Rightarrow b \equiv a \pmod{m}$ [Symmetric law]

- $a \equiv b \& b \equiv c \Rightarrow a \equiv c \pmod{m}$ [Transitive law]

$a \equiv b \& c \equiv d \Rightarrow a + c \equiv b + d \pmod{m}$

$$a - b = mk, c - d = ml \Rightarrow a - b + c - d = m(k+l) \Rightarrow (a+c) - (b+d) = m(k+l) \\ \Rightarrow m \mid (a+c) - (b+d) \Rightarrow a + c \equiv b + d \pmod{m}$$

~~#~~ $a \equiv b \& c \equiv d \Rightarrow a - c \equiv b - d \pmod{m}$

$a \equiv b \& c \equiv d \Rightarrow ac \equiv bd \pmod{m}$

$$a - b = mk, c - d = ml. \text{ Now, } ac \equiv bd \pmod{m} \Leftrightarrow m \mid ac - bd \\ \Leftrightarrow m \mid c(a - b) + b(c - d) \Leftrightarrow m \mid cmk + bml, \text{ which is true} \\ [\quad ac - bc + bc - bd \quad]$$

$$a \equiv b \pmod{m} \Rightarrow a^n \equiv b^n \pmod{m}$$

$$10 \equiv 6 \pmod{4} \quad 14 \equiv 8 \pmod{3}$$

$$\Rightarrow 5 \cdot 2 \equiv 3 \cdot 2 \pmod{4} \Rightarrow 7 \cdot 2 \equiv 4 \cdot 2 \pmod{3}$$

$$\Rightarrow 5 \equiv 3 \pmod{4} \times \Rightarrow 7 \equiv 4 \pmod{3} \checkmark$$

$$\Rightarrow \frac{5 \cdot 2 - 3 \cdot 2}{4} = 2 \cdot \frac{(5-3)}{4} \quad \frac{7 \cdot 2 - 4 \cdot 2}{3} = 2 \left(\frac{7-4}{3} \right)$$

$$\Rightarrow 2 \equiv -1 \pmod{3}$$

$$\Rightarrow 2^n \equiv (-1)^n \pmod{3}$$

if n even $3 \mid 2^n - 1$

n odd, $3 \mid 2^n + 1$

$$\Rightarrow 3 \mid 2^n + 1 - 3 = 2^n - 2$$

$$\# ad \equiv bd \pmod{m} \Leftrightarrow a \equiv b \pmod{m}, d \nmid m$$

$$\exists d', m' \text{ s.t. } dd' + mm' = \gcd(d, m) = 1$$

$$\Rightarrow dd' + mm' \equiv 1 \pmod{m} \Rightarrow dd' \equiv 1 \pmod{m}$$

$$\text{Now, } ad \equiv bd \pmod{m} \Rightarrow add' \equiv bdd' \pmod{m} \Rightarrow a \equiv b \pmod{m}$$

$d' \rightarrow 1/d$, multiplicative inverse of d modulo m .

$$\# ad \equiv bd \pmod{md} \Leftrightarrow a \equiv b \pmod{m} \quad [d \neq 0]$$

$$md \mid ad - bd \Rightarrow ad - bd = mdk \Rightarrow d(a-b) = mdk \Rightarrow a-b = mk$$

$$\Rightarrow m \mid a-b \Rightarrow a \equiv b \pmod{m}$$

$$\# ad \equiv bd \pmod{m} \Leftrightarrow a \equiv b \pmod{\frac{m}{\gcd(d, m)}} \quad \boxed{\frac{ad - bd}{m} = \frac{d(a-b)}{m}}$$

$$\exists d', m', \text{ s.t. } dd' + mm' = \gcd(d, m)$$

$$\Rightarrow dd' \equiv \gcd(d, m) \pmod{m}$$

$$\text{Now, } ad \equiv bd \pmod{m} \Rightarrow add' \equiv bdd' \pmod{m}$$

$$\Rightarrow a \frac{\gcd(d, m)}{m} \equiv b \frac{\gcd(d, m)}{m} \pmod{\frac{m}{\gcd(d, m)} * \cancel{\gcd(d, m)}}$$

$$\Rightarrow a \equiv b \pmod{\frac{m}{\gcd(d, m)}}$$

$$\# a \equiv b \pmod{md} \Rightarrow a \equiv b \pmod{m}$$

$$\# a \equiv b \pmod{m} \& a \equiv b \pmod{n} \Leftrightarrow a \equiv b \pmod{\text{lcm}(m, n)}$$

$$d \nmid m, d \nmid n \Rightarrow d \mid \gcd(m, n)$$

$$m \nmid d, n \nmid d \Rightarrow \text{lcm}(m, n) \mid d \quad [mn = \gcd(m, n) * \text{lcm}(m, n)]$$

$$\text{if } m \perp n, a \equiv b \pmod{m} \& a \equiv b \pmod{n} \Leftrightarrow a \equiv b \pmod{mn}$$

Chinese Remainder Theorem (Ex. 30)

Lect. 13: Congruence (Modular Arithmetic)

Sec. A

$$a = mq + r, 0 \leq r < m \Rightarrow a = m \lfloor a/m \rfloor + a \bmod m$$

Defn: $a \equiv b \pmod{m} \Leftrightarrow a \bmod m = b \bmod m$

$a \equiv b \pmod{m} \Leftrightarrow m \mid (a - b)$

1st part: $a = \lfloor a/m \rfloor m + a \bmod m, b = \lfloor b/m \rfloor m + b \bmod m$
 $\overset{r \rightarrow r}{\Rightarrow} a - b = (\lfloor a/m \rfloor - \lfloor b/m \rfloor)m + a \bmod m - b \bmod m \Rightarrow m \mid (a - b)$

2nd part: $m \mid (a - b)$

$\overset{r \rightarrow r}{\Rightarrow} a = \lfloor a/m \rfloor m + r_1, b = \lfloor b/m \rfloor m + r_2; 0 \leq r_1, r_2 < m$

$$a - b = (\lfloor a/m \rfloor - \lfloor b/m \rfloor)m + r_1 - r_2 \Rightarrow m \mid (r_1 - r_2)$$

$$\max(r_1, r_2) = m - 1, \min(r_1, r_2) = -(m - 1)$$

Multiples of m : $\dots - 3m, -2m, -m, 0, m, 2m, 3m, \dots$

$$r_1 - r_2 = 0 \Rightarrow r_1 = r_2 \Rightarrow a \bmod m = b \bmod m \Rightarrow a \equiv b \pmod{m}$$

• $a \equiv a \pmod{m}$ [reflexive law] $LHS = RHS \Rightarrow LHS \equiv RHS \pmod{m}$

• $a \equiv b \Leftrightarrow b \equiv a \pmod{m}$ [symmetric law]

• $a \equiv b \wedge b \equiv c \Rightarrow a \equiv c \pmod{m}$ [transitive law]

$a \equiv b \wedge c \equiv d \pmod{m} \Rightarrow a + c \equiv b + d \pmod{m}$

$$m \mid (a - b), m \mid c - d \Rightarrow m \mid (a - b) + (c - d) \Rightarrow m \mid (a + c) - (b + d) \Rightarrow (a + c) \equiv (b + d) \pmod{m}$$

$a \equiv b \wedge c \equiv d \pmod{m} \Rightarrow a - c \equiv b - d \pmod{m}$

$a \equiv b \wedge c \equiv d \Rightarrow ac \equiv bd \pmod{m}$ $[ac - bd = ac - bc + bc - bd]$

$$m \mid a - b, m \mid c - d; m \mid ac - bd \Leftrightarrow m \mid c(a - b) + b(c - d)$$

$\Leftrightarrow m \mid cmk + bml$, which is true

$a \equiv b \Rightarrow a^n \equiv b^n \pmod{m}$

$$3 \mid 2^n - 2$$

$$\begin{array}{ll} 10 \equiv 6 \pmod{4} & 14 \equiv 8 \pmod{3} \\ \Rightarrow 5, 2 \equiv 3, 2 \pmod{4} & \Rightarrow 7, 2 \equiv 4, 2 \pmod{3} \\ \Rightarrow 5 \equiv 3 \pmod{4} \times & \Rightarrow 7 \equiv 4 \pmod{3} \checkmark \\ \frac{10-6}{4} = \frac{2(5-3)}{4} & \frac{14-8}{3} = \frac{2(7-4)}{3} \end{array}$$

$$\begin{array}{l} 2 \equiv -1 \pmod{3} \\ \Rightarrow 2^n \equiv (-1)^n \pmod{3} \\ \text{n even } 3 \mid 2^n - 1 \\ \text{n odd } 3 \mid 2^n + 1 \Rightarrow 3 \mid 2^n + 1 - 3 \\ \Rightarrow 3 \mid 2^n - 2 \end{array}$$

$\# ad \equiv bd \Leftrightarrow a \equiv b \pmod{m}$ when $d \perp m$

$$\begin{aligned} \gcd(d, m) = 1; \exists d', m', \text{ s.t. } dd' + mm' = 1 \\ \Rightarrow dd' \equiv 1 \pmod{m} \end{aligned}$$

$$\text{Now, } ad \equiv bd \Rightarrow add' \equiv bdd' \Rightarrow a \equiv b \pmod{m}$$

$d' \rightarrow$ multiplicative inverse of d modulo m
 Cyclic group of order m . $(\mathbb{Z}/m\mathbb{Z})^\times$

$\# ad \equiv bd \pmod{md} \Leftrightarrow a \equiv b \pmod{m}; d \neq 0$

$ad \equiv bd \pmod{m} \Leftrightarrow a \equiv b \pmod{m/\gcd(d, m)}$

$$\begin{aligned} \exists d', m', \text{ s.t. } dd' + mm' = \gcd(d, m) \\ \Rightarrow dd' \equiv \gcd(d, m) \pmod{m} \end{aligned}$$

$$ad \equiv bd \pmod{m} \Leftrightarrow add' \equiv bdd' \pmod{m}$$

$$\begin{aligned} \Leftrightarrow a \frac{\gcd(d, m)}{\gcd(d, m)} \equiv b \frac{\gcd(d, m)}{\gcd(d, m)} \pmod{m} \\ \Leftrightarrow a \equiv b \pmod{m/\gcd(d, m)} \end{aligned}$$

$\# a \equiv b \pmod{md} \Rightarrow a \equiv b \pmod{m}$

$\# a \equiv b \pmod{m} \& a \equiv b \pmod{n}$

$$\Leftrightarrow a \equiv b \pmod{\text{lcm}(m, n)}$$

$$40 \equiv 4 \pmod{12}$$

$$40 \equiv 4 \pmod{18}$$

$$40 \equiv 4 \pmod{216} \times$$

$$40 \equiv 4 \pmod{36} \checkmark$$

$$mn = \gcd(m, n) * \text{lcm}(m, n)$$

$$d \mid m \& d \mid n \Rightarrow d \mid \gcd(m, n)$$

$$m \mid d \& n \mid d \Rightarrow \text{lcm}(m, n) \mid d$$

$\# m \perp n; a \equiv b \pmod{m} \& a \equiv b \pmod{n} \Leftrightarrow a \equiv b \pmod{mn}$

\hookrightarrow Ex. 30 (Chinese Remainder Theorem)

$$a \equiv b \pmod{m} \Leftrightarrow a \equiv b \pmod{p^{m_p}} \forall p$$

Lect. 14: Residue Classes

Chinese Remainder Thm: Let m_1, \dots, m_r be integers greater than 1 with $m_j \perp m_k$ for $1 \leq j < k \leq r$.

$$\text{Let } M = m_1 \dots m_r.$$

If $\alpha_1, \dots, \alpha_r$ are any integers then

$$x \equiv \alpha_1 \pmod{m_1}$$

$$x \equiv \alpha_2 \pmod{m_2}$$

 \vdots

$$x \equiv \alpha_r \pmod{m_r}$$

has a unique solution modulo M .

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

$$x \equiv ? \pmod{105}$$

$$\hookrightarrow 2^3$$

Uniqueness: Let x, y be two solutions to the system s.t. $x \not\equiv y \pmod{M}$.

$$\left. \begin{array}{l} x \equiv \alpha_1 \pmod{m_1} \& y \equiv \alpha_1 \pmod{m_1} \Rightarrow x \equiv y \pmod{m_1} \\ x \equiv \alpha_2 \pmod{m_2} \& y \equiv \alpha_2 \pmod{m_2} \Rightarrow x \equiv y \pmod{m_2} \\ \vdots & \vdots \\ x \equiv \alpha_r \pmod{m_r} \& y \equiv \alpha_r \pmod{m_r} \Rightarrow x \equiv y \pmod{m_r} \end{array} \right\} x \equiv y \pmod{M}, \text{ a contradiction.}$$

Existence (implicit construction):

Proof by induction on r . Base case: $r=2$,

$$x \equiv \alpha_1 \pmod{m_1} \quad m_1 \perp m_2. \quad \exists m'_1, m'_2, \text{s.t. } m_1 m'_1 + m_2 m'_2 = 1$$

$$x \equiv \alpha_2 \pmod{m_2}$$

$$\text{We claim } x = \alpha_1 m_2 m'_2 + \alpha_2 m_1 m'_1$$

$$\text{Indeed, } x = \alpha_1 (1 - m_1 m'_1) + \alpha_2 m_1 m'_1 = \alpha_1 + (\alpha_2 - \alpha_1) m_1 m'_1$$

$$x \equiv \alpha_1 \pmod{m_1}$$

$$\text{Analogously, } x \equiv \alpha_2 \pmod{m_2}$$

Hypo.: True upto ($r-1$). Let $\hat{M} = m_1 \dots m_{r-1}$. By induction hypo. $\exists x \equiv \hat{x} \pmod{\hat{M}}$

Induction step: $\gcd(\hat{M}, m_r) = 1$. Using the same argument of the base case, we can show that $\exists \hat{M}', m'_r$ s.t. $\hat{M} \hat{M}' + m_r m'_r = 1$

and $x = \hat{x} m_r m'_r + \alpha_r \hat{M} \hat{M}'$ is a solution. [Proved]

$x \bmod 15$	$x \bmod 3$	$x \bmod 5$	
0	0	0	$7 * 13 \pmod{15}$
1	$\equiv 1$	$\equiv 1$	$7 * 13 \equiv 1 \cdot 1 \pmod{3}$
2	2	2	$7 * 13 \equiv 2 \cdot 3 \equiv 1 \pmod{5}$
3	0	3	
4	$\equiv 1$	4	
5	2	0	
6	0	$\equiv 1$	
7	$\equiv 1$	2	
8	2	3	
9	0	4	
10	$\equiv 1$	0	
11	2	$\equiv 1$	
12	0	2	
13	$\equiv 1$	3	
14	2	4	

Self-study: $x^2 \equiv 1 \pmod{m}$

$$\text{Res}(x) = (x \bmod m_1, x \bmod m_2, \dots, x \bmod m_r), \quad m_1 \neq m_r$$

$$\gcd(8, 12) = 4$$

$$0 \times 8 \bmod 12, 1 \times 8 \bmod 12, 2 \times 8 \bmod 12, \dots, 11 \times 8 \bmod 12$$

$$0, 8, 4, 0, 8, 4, 0, 8, 4$$

$$0 \times n \bmod m, 1 \times n \bmod m, 2 \times n \bmod m, \dots, (m-1) \times n \bmod m \quad [\# m \text{ numbers}]$$

$$\text{Let } \gcd(m, n) = d$$

There are $(0, d, 2d, 3d, \dots, m-d)$ in d copies $\left[\frac{m}{m/d} = d\right]$

$$\underbrace{m/d}_{m/d} \quad \underbrace{(m/d-1)d}_{(m/d-1)d}$$

$$\#\{kn \equiv sd \pmod{m}\} \quad \text{Let } m = m'd, n = n'd \text{ where } \gcd(m', n') = 1$$

$$\Leftrightarrow kn'd \equiv sd \pmod{m'd}$$

$$\Leftrightarrow kn' \equiv s \pmod{m'}$$

$$[d = -t]$$

$$\Leftrightarrow m' \mid kn' - s \Leftrightarrow kn' - s = tm' \Leftrightarrow kn' - tm' = s \Leftrightarrow kn' + dm' = s$$

$$\gcd(m', n') = 1; \exists k', l' \text{ s.t. } kn' + l'm' = 1 \Rightarrow \underbrace{k's}_{k} n' + \underbrace{l's}_{l} m' = s$$

$$\#\{jn \equiv (j+m/d)n \pmod{m}\}$$

$$\Leftrightarrow jn \equiv jn + m/d \neq n \pmod{m}$$

$\Leftrightarrow n/d \neq m \equiv 0 \pmod{m}$, which is true.

What if period $< m/d$?

Let $\exists 0 \leq j, k < m/d \ni jn = kn$. Now,

$$jn \equiv kn \pmod{m}$$

$$\Leftrightarrow jn/d \equiv kn/d \pmod{m/d}$$

$$\Leftrightarrow jn' \equiv kn' \pmod{m'}$$

$\Leftrightarrow j \equiv k \pmod{m/d} \Leftrightarrow j = k$, a contradiction.

Lect. 14: Residue Classes

Chinese Remainder Th^m: Let m_1, \dots, m_r be integers greater than 1 with $m_i \perp m_k$

Let $M = m_1 \dots m_r$. If a_1, \dots, a_r are any integers, then for $1 \leq j < k \leq r$.

$$x \equiv a_1 \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2}$$

!

$$x \equiv a_r \pmod{m_r}$$

has a unique solution modulo M .

$$x \equiv 2 \pmod{3}$$

$$x \equiv 3 \pmod{5}$$

$$x \equiv 2 \pmod{7}$$

$$x \equiv ? \pmod{105}$$

↪ 23

Uniqueness: $\exists x, y$ s.t. $x \not\equiv y \pmod{M}$ but both satisfy the systems

$$x \equiv a_1 \pmod{m_1} \& y \equiv a_1 \pmod{m_1} \Rightarrow x \equiv y \pmod{m_1}$$

$$x \equiv a_2 \pmod{m_2} \& y \equiv a_2 \pmod{m_2} \Rightarrow x \equiv y \pmod{m_2}$$

!

!

!

$$x \equiv a_r \pmod{m_r} \& y \equiv a_r \pmod{m_r} \Rightarrow x \equiv y \pmod{m_r}$$

$$x \equiv y \pmod{M}$$

, a contradiction

Existence (implicit construction):

Proof by induction on r . Base case: $r=2$

$$x \equiv a_1 \pmod{m_1} \quad m_1 \perp m_2, \exists m'_1, m'_2 \text{ s.t. } m_1 m'_1 + m_2 m'_2 = 1$$

$$x \equiv a_2 \pmod{m_2}$$

We claim, $x = a_1 m_2 m'_2 + a_2 m_1 m'_1$ is a solution

$$\text{Indeed, } x = a_1(1 - m_2 m'_2) + a_2 m_1 m'_1 = a_1 + (a_2 - a_1)m_1 m'_1 \Rightarrow x \equiv a_1 \pmod{m_1}$$

$$\text{Analogously, } x = a_1 m_2 m'_2 + a_2(1 - m_1 m'_1) = (a_1 - a_2)m_2 m'_2 + a_2 \Rightarrow x \equiv a_2 \pmod{m_2}$$

Hypothesis: True upto $(r-1)$. Let $\hat{M} = m_1 \dots m_{r-1}$. By induction hypothesis,

$$\exists \hat{A} \text{ s.t. } x \equiv \hat{A} \pmod{\hat{M}}$$

Induction step: $\gcd(\hat{M}, m_r) = 1$. Using the same argument as the base case, we can show that, $\exists \hat{M}', m'_r$ s.t. $\hat{M}' \hat{M}' + m_r m'_r = 1$

and $x = \hat{A} m_r m'_r + a_r \hat{M}' \hat{M}'$ is a solution. [Proved]

$\text{Res}(x) = (x \bmod m_1, x \bmod m_2, \dots, x \bmod m_r); m_j \neq m_k; 1 \leq j < k \leq r.$

$x \bmod 15 \quad x \bmod 3 \quad x \bmod 5$

	0	0	0	$7*13 \bmod 15$
1	$\equiv 1$	$\equiv 1$	$\equiv 1$	$7*13 \equiv 1 \cdot 1 \pmod{3}$
2	2	2	2	$7*13 \equiv 2 \cdot 3 \equiv 6 \equiv 1 \pmod{5}$
3	0	0	3	
4	$\equiv 1$	1	4	
5	2	2	0	
6	0	$\equiv 1$	$\equiv 1$	
7	$\equiv 1$	1	2	
8	2	2	3	
9	0	0	4	
10	$\equiv 1$	$\equiv 1$	0	
11	2	$\equiv 1$	$\equiv 1$	
12	0	0	2	
13	$\equiv 1$	1	3	
14	2	2	4	

Self-study: $x^2 \equiv 1 \pmod{m}$

$$0 \times 8 \bmod 12, 1 \times 8 \bmod 12, 2 \times 8 \bmod 12, \dots, 11 \times 8 \bmod 12 \quad \gcd(12, 8) = 4$$

$$0 \qquad \qquad 8 \qquad \qquad 4 \qquad \qquad 0, 8, 4, \quad 0, 8, 4, \quad 0, 8, 4$$

$$0 \times n \bmod m, 1 \times n \bmod m, 2 \times n \bmod m, \dots, (m-1) \times n \bmod m \quad [\gcd(m, n) = d] \rightarrow \text{total numbers}$$

$$\underbrace{0, d, 2d, 3d, \dots, m-d}_{m/d \text{ numbers}} \quad \# \text{ copies: } \frac{m}{m/d} = d$$

$$\# kn \equiv sd \pmod{m} \quad [m = m'd, n = n'd, \gcd(m', n') = 1]$$

$$\Leftrightarrow kn'd \equiv sd \pmod{m'd}$$

$$\Leftrightarrow kn' \equiv s \pmod{m'}$$

$$\Leftrightarrow m' \mid kn' - s \Leftrightarrow kn' - s = tm' \Leftrightarrow kn' - tm' = s \Leftrightarrow kn' + lm' = s \quad [l = -t]$$

$$\exists k', l', \text{s.t. } kn' + lm' = 1 \Rightarrow \underbrace{k's}_{k}, \underbrace{n'+l'm'}_{l} = s$$

$$\# jn \equiv (j+m_d)n \pmod{m}$$

$$\Leftrightarrow jm \equiv jn + m_d * n \pmod{m}$$

$$\Leftrightarrow m_d * m \equiv 0 \pmod{m}, \text{ which is true}$$

What if period < m_d ?

Let $\exists 0 \leq j, k < m_d \ni jn = kn$. Now,

$$jn \equiv kn \pmod{m}$$

$$\Leftrightarrow jn'd \equiv kn'd \pmod{m'd}$$

$$\Leftrightarrow jn' \equiv kn' \pmod{m'}$$

$$\Leftrightarrow j \equiv k \pmod{m/d} \Leftrightarrow j = k, \text{ a contradiction.}$$

Lect. 15: Fermat's th^m, Wilson's th^m, Euler's th^m

$$x^2 \equiv 1 \pmod{m} \Rightarrow x^2 \equiv 1 \pmod{p}; p > 2$$

$$p \nmid x^2 - 1 \Rightarrow p \nmid (x+)(x-). \text{ If } p \nmid x+ \& p \nmid x-1, p \nmid (x+)-(x-1) \Rightarrow p \nmid 2 \Rightarrow p \nmid x$$

$$p \nmid x+ \text{ or } p \nmid x-1$$

$$\Rightarrow x \equiv -1 \pmod{p} \text{ or } x \equiv +1 \pmod{p} \Rightarrow x \equiv \pm 1 \pmod{p}$$

The m numbers $0 \pmod{m}, n \pmod{m}, 2n \pmod{m}, \dots, (m-1)n \pmod{m}$
 consist of precisely d copies of the m/d numbers $0, d, 2d, \dots, m-d$ in some order.
 where $d = \gcd(m, n)$

If $m \nmid n$, then there is exactly 1 copy of the m numbers

$$0, d, 2d, \dots, (m-1)d \equiv 0, 1, 2, \dots, m-1. \text{ (complete residue class)}$$

If m is a prime p , then

$$n \cdot 2n \cdot 3n \cdots (p-1)n \equiv 1 \cdot 2 \cdot 3 \cdots (p-1) \pmod{p}$$

$$\Rightarrow n^{p-1} \pmod{(p-1)!} \equiv (p-1)! \pmod{p}$$

$$\Rightarrow n^{p-1} \equiv 1 \pmod{p}, n \neq p \quad [\text{Fermat's little th}^m]$$

Wilson's th^m: $(n-1)! \equiv -1 \pmod{n} \Leftrightarrow n \text{ is prime; if } n > 1.$

Proof: Sufficiency: If $n > 1$ is not a prime, it has a prime divisor $p < n$
 $p \nmid (n-1)!$, a contradiction. Because if $(n-1)! \equiv -1 \pmod{n}$

$$\Rightarrow (n-1)! \equiv -1 \pmod{p}, \text{ but } p \nmid (n-1)! \Rightarrow (n-1)! \equiv 0 \pmod{p}.$$

Necessity: $(p-1)! \equiv -1 \pmod{p}$ for any prime p .

Case 1: $p = 2$. LHS = $(2-1)! = 1! = 1 \equiv -1 \pmod{2}$, which is true.

Case 2: $p > 2$. $\forall n < p$, and $n \neq p \Rightarrow \exists n'$, s.t. $n' \cdot n \equiv 1 \pmod{p}$

$$p=5$$

$$\begin{aligned} & \underbrace{1 \cdot 2 \cdot 3 \cdot 4 \cdots (p-2)}_{\equiv 1 \pmod{p}} \underbrace{(p-1)}_{(-1)} \\ & \equiv -1 \pmod{p} \end{aligned}$$

$$\begin{aligned} & 1' = 1, 2' = 3, 3' = 2, 4' = 4 \\ & n' \equiv n \Rightarrow n^2 \equiv 1 \pmod{p} \\ & \Rightarrow n \equiv \pm 1 \pmod{p} \\ & \Rightarrow n \neq 1, p-1 \end{aligned}$$

Euler's thm:

How many integers $\leq m$ i.e., $\{0, 1, 2, \dots, m-1\}$ are relatively prime to m ?

$\varphi(m) \Rightarrow$ Euler's totient function.

$$\varphi(1) = 1, \varphi(p) = p-1$$

Euler's thm: $n^{\varphi(m)} \equiv 1 \pmod{m}$, $n \perp m$.

If $m = p^k$, $n \perp m \Rightarrow n \perp p^k \Rightarrow n \perp p$. $\Rightarrow p \nmid n$

The multiples of p in $\{0, 1, \dots, p^{k-1}\}$ are $\{0, p, 2p, \dots, p^{k-1} - 1\}$

$$\varphi(p^k) = p^k - p^{k-1}$$

$m > 1$ and not a prime power, $m = m_1 m_2$, $m_1 > 1, m_2 > 1, m_1 \perp m_2$

$n \perp m \Leftrightarrow n \perp m_1 \& n \perp m_2 \Leftrightarrow \gcd(n, m_1) = 1 \& \gcd(n, m_2) = 1$

$\varphi(m) \Leftrightarrow \gcd(n \bmod m_1, m_1) = 1 \& \gcd(n \bmod m_2, m_2) = 1$

$\Leftrightarrow \underbrace{n \bmod m_1 \perp m_1}_{\varphi(m_1)} \& \underbrace{n \bmod m_2 \perp m_2}_{\varphi(m_2)}$

$$\varphi(m) = \varphi(m_1) \varphi(m_2), m_1 \perp m_2 \quad [\text{Multiplicative function}]$$

$$\varphi(m) = \varphi\left(\prod_{p|m} p^{m_p}\right) = \prod_{p|m} \varphi(p^{m_p}) = \prod_{p|m} p^{m_p} - p^{m_p-1} = \prod_{p|m} p^{m_p} \left(1 - \frac{1}{p}\right)$$

$$= p_1^{m_1} \left(1 - \frac{1}{p_1}\right) p_2^{m_2} \left(1 - \frac{1}{p_2}\right) \cdots p_k^{m_k} \left(1 - \frac{1}{p_k}\right)$$

$$= m \prod_{p|m} \left(1 - \frac{1}{p}\right)$$

Lect. 15: Fermat's th^m, Wilson's th^m, Euler's th^m

Sec. A

The m numbers $0 \pmod{m}, n \pmod{m}, 2n \pmod{m}, \dots, (m-1)n \pmod{m}$ consist of precisely d copies of the $\frac{m}{d}$ numbers $0, d, 2d, \dots, m-d$ in some order, where $d = \gcd(m, n)$.

If $m \nmid n$, there is exactly 1 copy of the m numbers $0, d, 2d, \dots, m-d$

If m is a prime p , then

$$\equiv 0, 1, 2, \dots, m-1$$

(complete residue class)

$$n \cdot 2n \cdot 3n \cdots (p-1)n \equiv 1 \cdot 2 \cdots (p-1) \pmod{p}$$

$$\Rightarrow n^{p-1} (p-1)! \equiv (p-1)! \pmod{p} \Rightarrow n^{p-1} \equiv 1 \pmod{p} \quad [\text{Fermat's little th}^m]$$

$$x^2 \equiv 1 \pmod{p}, p > 2. \Rightarrow p \nmid x^2 - 1 \Rightarrow p \nmid (x+1)(x-1)$$

$$\text{If } p \nmid x+1 \text{ & } p \nmid x-1 \Rightarrow p \mid (x+1) - (x-1) \Rightarrow p \mid 2 \Rightarrow p = 2, \text{ a contradiction.}$$

$$\text{So, } p \mid x+1 \text{ or, } p \mid x-1 \Rightarrow x \equiv -1 \pmod{p} \text{ or, } x \equiv +1 \pmod{p} \Rightarrow x \equiv \pm 1 \pmod{p}.$$

$$\Rightarrow x \equiv 1, p-1 \pmod{p}$$

Wilson's th^m: $(n-1)! \equiv -1 \pmod{n} \Leftrightarrow n \text{ is prime, if } n > 1$

Proof: Sufficiency: Given $(n-1)! \equiv -1 \pmod{n}$, we need to prove n is prime.

Assume $n > 1$ is not a prime. If has a prime divisor $p < n$.

$p \nmid (n-1)!$, this is in violation of our assumption. Because;

$(n-1)! \equiv -1 \pmod{n} \Rightarrow (n-1)! \equiv -1 \pmod{p}$. But $p \nmid (n-1)! \Rightarrow (n-1)! \equiv 0 \pmod{p}$.

Necessity: $(p-1)! \equiv -1 \pmod{p}$, p is prime.

Case 1 ($p=2$): If $p=2$, $(2-1)! = 1 \equiv -1 \pmod{2}$

Case 2 ($p > 2$): $\forall n < p$, $n \perp p$. And $n \perp p$ implies $\exists n'$, s.t. $n'n \equiv 1 \pmod{p}$

$$\begin{aligned} & 1, 2, 3, \dots, (p-2), (p-1) \pmod{p} \\ & \equiv \underbrace{1 \cdot 2 \cdot 3 \cdots}_{(1,1,1,\dots)} (p-1) \pmod{p} \\ & \equiv -1 \pmod{p} \end{aligned}$$

$$\begin{aligned} & p=5 \\ & 1' = 1, 2' = 3, 3' = 2, 4' = 4 \\ & n' \equiv n \Rightarrow n^2 \equiv 1 \pmod{p} \\ & \Rightarrow n \equiv \pm 1 \pmod{p} \\ & \equiv 1, (p-1) \pmod{p} \end{aligned}$$

Euler's thm:

How many integers $\{0, 1, \dots, m-1\}$ are relatively prime to m ?

$\varphi(m) \rightarrow$ Euler's totient function.

$\varphi(1) = 1$, $\varphi(p) = p-1$. Euler's thm: $n^{\varphi(m)} \equiv 1 \pmod{m}$, if $n \perp m$.

If $m = p^k$, $n \perp m \Rightarrow n \perp p^k \Rightarrow n \perp p \Rightarrow p \nmid n$

The multiples of p in $\{0, 1, 2, \dots, p^k-1\}$ are $\{0, p, 2p, \dots, p^k-p\}$
 $\underbrace{0}_{p^k}, \underbrace{p, 2p, \dots, (p^{k-1}-1)p}_{p^{k-1}}$

$$\varphi(p^k) = p^k - p^{k-1}$$

$m > 1$ is not a prime. Then, $m = m_1 m_2$; $m_1 > 1, m_2 > 1 \& m_1 \perp m_2$

$n \perp m \Leftrightarrow n \perp m_1 \& n \perp m_2 \Leftrightarrow \gcd(n, m_1) = 1 \& \gcd(n, m_2) = 1$

$\varphi(m) \Leftrightarrow \gcd(n \bmod m_1, m_1) = 1 \& \gcd(n \bmod m_2, m_2) = 1$

$\Leftrightarrow \underbrace{n \bmod m_1 \perp m_1}_{\varphi(m_1)} \& \underbrace{n \bmod m_2 \perp m_2}_{\varphi(m_2)}$

$$\varphi(m) = \varphi(m_1) \varphi(m_2), m_1 \perp m_2 \quad [\text{Multiplicative function}]$$

$$m = p_1^{m_1} p_2^{m_2} \cdots p_k^{m_k}$$

$$\begin{aligned}\varphi(m) &= \varphi(\prod_{p \mid m} p^{m_p}) = \prod_{p \mid m} \varphi(p^{m_p}) = \prod_{p \mid m} (p^{m_p} - p^{m_p-1}) = \prod_{p \mid m} p^{m_p} \left(1 - \frac{1}{p}\right) \\ &= p_1^{m_1} \left(1 - \frac{1}{p_1}\right) p_2^{m_2} \left(1 - \frac{1}{p_2}\right) \cdots p_k^{m_k} \left(1 - \frac{1}{p_k}\right) \\ &= m \prod_{p \mid m} \left(1 - \frac{1}{p}\right)\end{aligned}$$

Lect. 16: Generating functions (and Euler's th^m contd.)

"There's a time in your life before you understand generating functions and a time after, and I can't think of anything that connects them other than a leap of faith." - Grant Sanderson (3Blue1Brown)

Find the number of subsets of $\{1, \dots, 2000\}$ where the sum of the elements is divisible by 5.

$$\text{sum } \{3, 1, 4\} = 8 \times \text{sum } \{2, 3, 5\} = 10 \checkmark$$

Total subset: 2^{2000}

$$\text{Guess: } \frac{1}{5} \cdot 2^{2000}$$

Simpler example: $\{1, 2, 3, 4, 5\}$

$$\begin{array}{ccccc} \{ \} \rightarrow 0 & \{1\} \rightarrow 1 & \{2\} \rightarrow 2 & \{3\} \rightarrow 3 & \{4\} \rightarrow 4 \\ & & & \{1, 2\} \rightarrow 3 & \{1, 3\} \rightarrow 4 \end{array}$$

$$\begin{array}{ccccc} \{5\} \rightarrow 5 & \{1, 5\} \rightarrow 6 & \{2, 5\} \rightarrow 7 & \{3, 5\} \rightarrow 8 & \{4, 5\} \rightarrow 9 \\ \{1, 4\} \rightarrow 5 & \{2, 4\} \rightarrow 6 & \{3, 4\} \rightarrow 7 & \{1, 2, 5\} \rightarrow 8 & \{1, 3, 5\} \rightarrow 9 \\ \{2, 3\} \rightarrow 5 & \{1, 2, 3\} \rightarrow 6 & \{1, 2, 4\} \rightarrow 7 & \{1, 3, 4\} \rightarrow 8 & \{2, 3, 4\} \rightarrow 9 \end{array}$$

$$\begin{array}{ccccc} \{1, 4, 5\} \rightarrow 10 & \{2, 4, 5\} \rightarrow 11 & \{3, 4, 5\} \rightarrow 12 & \{1, 3, 4, 5\} \rightarrow 13 & \{2, 3, 4, 5\} \rightarrow 14 \\ \{2, 3, 5\} \rightarrow 10 & \{1, 2, 3, 5\} \rightarrow 11 & \{1, 2, 4, 5\} \rightarrow 12 & & \end{array}$$

$$\{1, 2, 3, 4, 5\} \rightarrow 10$$

$$\text{Ans: } 8 > \frac{1}{5} \cdot 32$$

$$\{1, 2, 3, 4, 5\} \rightarrow 15$$

$$P(x) = (1+x)(1+x^2)(1+x^3)(1+x^4)(1+x^5)$$

$x = ? \rightarrow$ just a symbol

$$\begin{aligned} \text{Trivial ex. of expansion: } (a+b)^2 \\ = (a+b)(a+b) &= a \cdot a + a \cdot b + b \cdot a + b \cdot b \\ &= a^2 + 2ab + b^2 \end{aligned}$$

expansion: 5 binary choices (which term from each parenthesis to choose from)

$$1 + x + x^2 + x^3 + x^4 + x^5 + x^6 + x^7 + \dots$$

$$= 1 + x + x^2 + 2x^3 + 2x^4 + 3x^5 + \dots \quad | \text{ Fibonacci numbers: } f_n = f_{n-1} + f_{n-2}$$

$\hookrightarrow P(x)$: generating function.

$$F(x) = 0 + 1x^1 + 1x^2 + 2x^3 + 3x^4 + 5x^5 + \dots$$

$$F(x) = xF(x) + x^2 F(x) + x \Rightarrow F(x) = \frac{x}{1-x-x^2}$$

Art: Deduce the co-efficients without doing the expansion.

(Next class)

Reduced residue class: a subset X of the integers is called a "reduced residue class" (mod m) if:

$$1. \gcd(x, m) = 1 \quad \forall x \in X.$$

2. X contains $\varphi(m)$ elements.

3. No two elements of X are congruent (mod m). i.e., $x_i \not\equiv x_j \pmod{m} \quad \forall 1 \leq i < j \leq \varphi(m)$

Now, let $X = \{x_1, x_2, \dots, x_{\varphi(m)}\}$ be a reduced residue class (mod m).

Claim: $X_a = \{\alpha x_1, \alpha x_2, \dots, \alpha x_{\varphi(m)}\}$ ($\alpha \perp m$) is also a reduced residue class (mod m).

$$1. \gcd(\alpha x_i, m) = 1 \quad \forall x_i \in X_a.$$

$$\text{pf: } \gcd(\alpha, m) = 1, \gcd(x_i, m) = 1; \alpha \perp m \& x_i \perp m \Rightarrow \alpha x_i \perp m \Rightarrow \gcd(\alpha x_i, m) = 1$$

2. X_a contains $\varphi(m)$ elements. ✓

3. $\alpha x_i \not\equiv \alpha x_j \pmod{m}$ where $1 \leq i < j \leq \varphi(m)$.

pf: Assume, $\exists i, j$, st. $1 \leq i < j \leq \varphi(m)$ and $\alpha x_i \equiv \alpha x_j \pmod{m}$

$\Rightarrow x_i \equiv x_j \pmod{m}$, but this violates the third property of X being a reduced residue class. so a contradiction.

Hence, X_a is a permutation of X . \Rightarrow They will yield the same product (mod m)

$$\Rightarrow \prod_{i=1}^{\varphi(m)} x_i \equiv \prod_{i=1}^{\varphi(m)} \alpha x_i \equiv \alpha x_1 x_2 \dots x_{\varphi(m)} \equiv \alpha^{\varphi(m)} \prod_{i=1}^{\varphi(m)} x_i \pmod{m}$$

$$\Rightarrow \alpha^{\varphi(m)} \equiv 1 \pmod{m}, \alpha \perp m. \text{ [Proved]}$$

$$\text{Self-study: } \sum_{d \mid m} \varphi(d) = m$$

Lect. 16: Generating functions (and Euler's thⁿ contd.)

"There's a time in your life before you understand generating functions and a time after, and I can't think of anything that connects them other than a leap of faith." - Grant Sanderson (3Blue1Brown)

Find the number of subsets of $\{1, 2, \dots, 2000\}$ where the sum of those subsets is divisible by 5. $\sum \{3, 1, 4\} = 8 \times$ $\sum \{2, 3, 5\} = 10 \checkmark$

Total subsets: 2^{2000}

Guess: $\frac{1}{5} \cdot 2^{2000}$ not an integer

Simpler example: $\{1, 2, 3, 4, 5\}$

$$\begin{array}{lllll} \{\} \rightarrow 0 & \{1\} \rightarrow 1 & \{2\} \rightarrow 2 & \{3\} \rightarrow 3 & \{4\} \rightarrow 4 \\ & & \{1, 2\} \rightarrow 3 & & \{1, 3\} \rightarrow 4 \end{array}$$

$$\begin{array}{ccccc} \{5\} \rightarrow 5 & \{1, 5\} \rightarrow 6 & \{2, 5\} \rightarrow 7 & \{3, 5\} \rightarrow 8 & \{4, 5\} \rightarrow 9 \\ \{1, 4\} \rightarrow 5 & \{2, 4\} \rightarrow 6 & \{3, 4\} \rightarrow 7 & \{1, 2, 5\} \rightarrow 8 & \{1, 2, 5\} \rightarrow 9 \\ \{2, 3\} \rightarrow 5 & \{1, 2, 3\} \rightarrow 6 & \{1, 2, 4\} \rightarrow 7 & \{1, 3, 4\} \rightarrow 8 & \{2, 3, 4\} \rightarrow 9 \end{array}$$

$$\begin{array}{ccccc} \{1, 4, 5\} \rightarrow 10 & \{2, 4, 5\} \rightarrow 11 & \{3, 4, 5\} \rightarrow 12 & \{1, 3, 4, 5\} \rightarrow 13 & \{2, 3, 4, 5\} \rightarrow 14 \\ \{2, 3, 5\} \rightarrow 10 & \{1, 2, 3, 5\} \rightarrow 11 & \{1, 2, 4, 5\} \rightarrow 12 & & \\ \{1, 2, 3, 4\} \rightarrow 10 & & & & \end{array}$$

$$\text{Ans: } 8 > \frac{1}{5} \cdot 32 \quad \rightarrow \text{Trivial ex. of expansion: } (a+b)^5$$

$$\{1, 2, 3, 4, 5\} \rightarrow 15 \quad \rightarrow P(x) = (1+x)(1+x^2) \cdots (1+x^{2000}) \quad = a^2 + 2ab + b^2$$

$$P(x) = (1+x)(1+x^2)(1+x^3)(1+x^4)(1+x^5) \quad x=? \Rightarrow \text{just a symbol}$$

Expansion - 5 binary choices (which term for each parenthesis to choose)

$$P(x) = 1 + x + x^2 + x^3 + x^4 + x^5 + x^6 + \dots = 1 + x + x^2 + 2x^3 + 2x^4 + 3x^5 + \dots$$

Generating function.

↳ encodes the number of subsets with a particular sum, namely, the corresponding exponent.

Fibonacci number: $F(x) = 0 + 1x^1 + 1x^2 + 2x^3 + 3x^4 + 5x^5 + 8x^6 \dots$ $f_m = f(m) = f(m-1) + f(m-2)$

$$F(x) = x F(x) + x^2 F(x) + x \Rightarrow F(x) = \frac{x}{1-x-x^2}$$

Art: Deduce the co-efficients without doing the expansion.

Reduced residue system: a subset X of the integers is called a "reduced residue class" $(\text{mod } m)$ if:

1. $\gcd(x_i, m) = 1 \forall x_i \in X$.
2. X contains $\varphi(m)$ elements.

3. No two elements of X are congruent $(\text{mod } m)$. i.e., $x_i \not\equiv x_j (\text{mod } m)$

$$\forall 1 \leq i \neq j \leq \varphi(m)$$

Now, let $X = \{x_1, x_2, \dots, x_{\varphi(m)}\}$ be a reduced residue class $(\text{mod } m)$.

Claim: $X_a = \{ax_1, ax_2, \dots, ax_{\varphi(m)}\} (a \perp m)$ is also a reduced residue class $(\text{mod } m)$.

1. $\gcd(ax_i, m) = 1 \forall ax_i \in X_a$.

Pf: $\gcd(x_i, m) = 1 \Rightarrow x_i \perp m, a \perp m \Rightarrow ax_i \perp m \Rightarrow \gcd(ax_i, m) = 1 \forall ax_i \in X_a$.

2. X_a contains $\varphi(m)$ elements. ✓

3. Assume $\exists i, j, 1 \leq i \neq j \leq \varphi(m)$ s.t. $ax_i \equiv ax_j (\text{mod } m) [1 \leq i \neq j \leq \varphi(m)]$

$\Rightarrow x_i \equiv x_j (\text{mod } m)$, but this violates the third property of X being a reduced residue class, so a contradiction.

Hence X_a is a permutation of $X \Rightarrow$ they will yield the same product $(\text{mod } m)$

$$\Rightarrow \prod_{i=1}^{\varphi(m)} x_i \equiv \prod_{i=1}^{\varphi(m)} ax_i = ax_1 \cdot ax_2 \cdots ax_{\varphi(m)} = a^{\varphi(m)} \prod_{i=1}^{\varphi(m)} x_i (\text{mod } m)$$

$$\Rightarrow a^{\varphi(m)} \equiv 1 (\text{mod } m) (a \perp m) [\text{Proved}]$$

Self-study: $\sum_{d \mid m} \varphi(d) = m$

Lect. 17: Generating functions (contd.)

$$G(z) = g_0 + g_1 z + g_2 z^2 + g_3 z^3 + \dots = \sum_{n \geq 0} g_n z^n; g_n = [z^n] G(z)$$

$$\#a_{n+1} = 2a_n + 1 \quad (n \geq 0, a_0 = 0) \quad [\text{Generating functionalogy}]$$

$$\text{Generating function: } A(x) = \sum_{n \geq 0} a_n x^n = a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots$$

Multiply both sides by x^n and sum over all values of $n \geq 0$

$$\text{LHS: } \sum_{n \geq 0} a_{n+1} x^n = a_1 + a_2 x + a_3 x^2 + \dots = \frac{(a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots) - a_0}{x} = A(x)/x$$

$$\text{RHS: } \sum_{n \geq 0} (2a_n + 1)x^n = 2 \sum_{n \geq 0} a_n x^n + \sum_{n \geq 0} x^n = 2A(x) + \frac{1}{1-x}; |x| < 1$$

$$A(x)/x = 2A(x) + \frac{1}{1-x} \Rightarrow A(x) = \frac{x}{(1-x)(1-2x)} = x \left(\frac{2}{1-2x} - \frac{1}{1-x} \right)$$

$$\begin{aligned} \Rightarrow A(x) &= 2x(1+2x+2^2x^2+\dots) - x(1+x+x^2+\dots) \\ &= (2x+2^2x^2+2^3x^3+\dots) - (x+x^2+x^3+\dots) \\ &= (2-1)x + (2^2-1)x^2 + (2^3-1)x^3 + \dots + (2^n-1)x^n + \dots \end{aligned}$$

$$\#a_{n+1} = 2a_n + n \quad (n \geq 0; a_0 = 1) \quad \text{Generating function: } A(x) = \sum_{n \geq 0} a_n x^n$$

$$\text{LHS: } \sum_{n \geq 0} a_{n+1} x^n = a_1 + a_2 x + a_3 x^2 + \dots = \frac{(a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots) - a_0}{x} = \frac{A(x) - 1}{x}$$

$$\text{RHS: } \sum_{n \geq 0} (2a_n + n)x^n = 2 \sum_{n \geq 0} a_n x^n + \sum_{n \geq 0} nx^n = 2A(x) + \frac{x}{(1-x)^2}; |x| < 1 \quad \begin{cases} \sum_{n \geq 0} nx^n = \frac{x}{(1-x)^2} \\ \sum_{n \geq 0} nx^{n-1} = \frac{1}{(1-x)^2} \end{cases}$$

$$\begin{aligned} \frac{A(x) - 1}{x} &= 2A(x) + \frac{x}{(1-x)^2} \Rightarrow A(x) = \frac{1-2x+2x^2}{(1-x)^2(1-2x)} = \frac{A}{(1-x)^2} + \frac{B}{1-x} + \frac{C}{1-2x} \\ \Rightarrow A(x) &= \frac{-1}{(1-x)^2} + \frac{2}{1-2x} \end{aligned}$$

$$[x^n] \frac{-1}{(1-x)^2} = -(n+1); [x^n] \frac{2}{1-2x} = [x^n] 2(1+2x+2^2x^2+\dots+2^n x^n+\dots)$$

$$= [x^n] (2+2^2x+2^3x^2+\dots+2^{n+1}x^n+\dots)$$

$$[x^n] A(x) = 2^{n+1} - (n+1)$$

$f_{n+1} = f_n + f_{n-1}$, ($n \geq 1$; $f_0 = 0$, $f_1 = 1$) Generating function: $F(x) = \sum_{n \geq 0} f_n x^n$
 Multiply both sides by x^n and sum over all values of $n \geq 1$

$$\text{LHS} = f_2 x + f_3 x^2 + f_4 x^3 + \dots = \frac{(f_0 + f_1 x + f_2 x^2 + f_3 x^3 + f_4 x^4 + \dots) - f_0 x - f_1}{x} = \frac{F(x) - x}{x}$$

$$\text{RHS} = (f_0 + f_1 x + f_2 x^2 + \dots) + (f_0 x + f_1 x^2 + f_2 x^3 + \dots) = F(x) + x F(x)$$

$$\frac{F(x) - x}{x} = F(x) + x F(x) \Rightarrow F(x) = \frac{x}{1 - x - x^2} = \frac{x}{(1 - x r_+)(1 - x r_-)}$$

$$= \frac{1}{(r_+ - r_-)} \left(\frac{1}{1 - x r_+} - \frac{1}{1 - x r_-} \right) \quad [r_{\pm} = (1 \pm \sqrt{5})/2]$$

$$= \frac{1}{\sqrt{5}} \left(\sum_{n \geq 0} r_+^n x^n - \sum_{n \geq 0} r_-^n x^n \right)$$

$$\Rightarrow x^n [F(x)] = f(n) = \frac{1}{\sqrt{5}} (r_+^n - r_-^n) = \frac{1}{\sqrt{5}} \left\{ \left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2}\right)^n \right\}$$

Lect. 17: Generating functions contd.

$$G(z) = g_0 + g_1 z + g_2 z^2 + g_3 z^3 + \dots = \sum_{n \geq 0} g_n z^n; [z^n]G(z) = g_n$$

$$a_{n+1} = 2a_n + 1 \quad (n \geq 0; a_0 = 0)$$

[Generating functionology]

$$\text{Generating function: } A(x) = \sum_{n \geq 0} a_n x^n$$

Multiply both sides of the recurrence by x^n and sum over $\forall n \geq 0$.

$$\begin{aligned} \text{LHS: } \sum_{n \geq 0} a_{n+1} x^n &= a_1 + a_2 x + a_3 x^2 + \dots = \frac{a_1 x + a_2 x^2 + a_3 x^3 + \dots}{x} \\ &= \frac{(a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots) - a_0}{x} = A(x)/x \end{aligned}$$

$$\text{RHS} = \sum_{n \geq 0} 2a_n x^n = 2A(x) + \sum_{n \geq 0} x^n = 2A(x) + \frac{1}{1-x}; |x| < 1$$

$$A(x)/x = 2A(x) + \frac{1}{1-x} \Rightarrow A(x) = \frac{x(1+x+x^2+x^3+\dots)}{(1-2x)(1-x)} = x \left\{ \frac{2}{1-2x} - \frac{1}{1-x} \right\}$$

$$= 2x(1 + 2x + 2^2x^2 + 2^3x^3 + \dots) - x(1 + x + x^2 + \dots)$$

$$= (2x + 2^2x^2 + 2^3x^3 + \dots) - (x + x^2 + x^3 + \dots)$$

$$= (2-1)x + (2^2-1)x^2 + (2^3-1)x^3 + \dots + (2^n-1)x^n + \dots$$

$$a_{n+1} = 2a_n + n \quad (n \geq 0, a_0 = 1)$$

$$\sum_{n \geq 0} nx^n = x + 2x^2 + 3x^3 + \dots$$

$$\text{LHS} = A(x) - 1/x$$

$$= \frac{x}{(1-x)^2}; |x| < 1$$

$$\text{RHS} = 2A(x) + \sum_{n \geq 0} nx^n$$

$$\Rightarrow \sum_{n \geq 0} nx^{n-1} = \frac{1}{(1-x)^2}$$

$$= 2A(x) + \frac{x}{(1-x)^2}$$

$$\Rightarrow A(x) - 1/x = 2A(x) + \frac{x}{(1-x)^2} \Rightarrow A(x) = \frac{1-2x+2x^2}{(1-x)^2(1-2x)} = \frac{A}{(1-x)^2} + \frac{B}{1-x} + \frac{C}{1-2x}$$

$$\Rightarrow A(x) = \frac{-1}{(1-x)^2} + \frac{2}{1-2x}$$

$$[x^n] \frac{1}{(1-x)^2} = -(n+1)$$

$$\frac{2}{1-2x} = 2(1 + 2x + 2^2x^2 + 2^3x^3 + \dots) = 2 + 2^2x + 2^3x^2 + \dots + 2^{n+1}x^{n+1}$$

$$[x^n] A(x) = 2^{n+1} - (n+1)$$

$$f_{n+1} = f_n + f_{n-1} \quad (n \geq 1, f_0 = 0, f_1 = 1)$$

$$F(x) = \sum_{n \geq 0} f_n x^n = f_0 + f_1 x + f_2 x^2 + \dots$$

Multiply both sides by x^n and sum over $\forall n \geq 1$

$$\begin{aligned} \text{LHS: } f_2 x + f_3 x^2 + f_4 x^3 + \dots &= \frac{f_2 x^2 + f_3 x^3 + f_4 x^4 + \dots}{x} \\ &= \frac{(f_0 + f_1 x + f_2 x^2 + f_3 x^3 + \dots) - f_1 x - f_0}{x} = \frac{F(x) - x}{x} \end{aligned}$$

$$\begin{aligned} \text{RHS} &= F(x) + (f_0 x + f_1 x^2 + f_2 x^3 + \dots) \\ &= F(x) + x (f_0 + f_1 x + f_2 x^2 + \dots) \\ &= F(x) + x F(x) \end{aligned}$$

$$\frac{F(x) - x}{x} = F(x) + x F(x) \Rightarrow F(x) = \frac{x}{1 - x - x^2}$$

$$1 - x - x^2 = (1 - x r_+) (1 - x r_-), \quad [r_{\pm} = (1 \pm \sqrt{5})/2]$$

$$\begin{aligned} F(x) &= \frac{x}{1 - x - x^2} = \frac{1}{r_+ - r_-} \left(\frac{1}{1 - x r_+} - \frac{1}{1 - x r_-} \right) \\ &= \frac{1}{\sqrt{5}} \left\{ \sum_{n \geq 0} r_+^n x^n - \sum_{n \geq 0} r_-^n x^n \right\} \end{aligned}$$

$$f_n = \frac{1}{\sqrt{5}} (r_+^n - r_-^n) = \frac{1}{\sqrt{5}} \left\{ \left(\frac{1 + \sqrt{5}}{2} \right)^n - \left(\frac{1 - \sqrt{5}}{2} \right)^n \right\}$$

Lect 18: Binomial coefficients and dominoes

GFs with two variables:

$$f(n, k) = f(n-1, k) + f(n-1, k-1) \quad (f(n, 0) = 1, k \leq n)$$

For each $n = 0, 1, 2, \dots$ let $B_n(x) = \sum_{k \geq 0} f(n, k) x^k = f(n, 0) + f(n, 1)x + f(n, 2)x^2 + \dots$

Multiply the recurrence by x^k and sum over all $k \geq 1$

$$\text{LHS: } \sum_{k \geq 1} f(n, k) x^k = (f(n, 0) + f(n, 1)x + f(n, 2)x^2 + \dots) - f(n, 0) = B_n(x) - 1$$

$$\begin{aligned} \text{RHS: } & \sum_{k \geq 1} f(n-1, k) x^k + \sum_{k \geq 1} f(n-1, k-1) x^k \\ &= \{f(n-1, 0) + f(n-1, 1)x + f(n-1, 2)x^2 + \dots\} - f(n-1, 0) \\ &\quad + f(n-1, 0)x + f(n-1, 1)x^2 + \dots \\ &= B_{n-1}(x) - 1 + x B_{n-1}(x) \end{aligned}$$

$$\Rightarrow B_n(x) - 1 = B_{n-1}(x) - 1 + x B_{n-1}(x) \Rightarrow B_n(x) = (1+x) B_{n-1}(x)$$

$$\Rightarrow B_n(x) = (1+x)^n B_0(x) \quad [B_0(x) = \sum_{k \geq 0} f(0, k) x^k = f(0, 0) = 1]$$

$$\Rightarrow B_n(x) = (1+x)^n \quad [\text{Binomial function}]$$

$$[x^k] B_n(x) = \binom{n}{k} = \frac{n!}{(n-k)!k!} \quad ; \quad \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$$

How many ways T_n to cover a $2 \times n$ rectangle with 2×1 dominos?



$$n=0 \quad |, \quad n=1 \quad \square, \quad n=2 \quad \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array}, \quad \begin{array}{|c|c|} \hline \square & \\ \hline \end{array}; \quad n=3 \quad \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \end{array}, \quad \begin{array}{|c|c|c|} \hline \square & \square & \\ \hline \end{array}, \quad \begin{array}{|c|c|} \hline \square & \\ \hline \square & \square \\ \hline \end{array} \quad T_n = f_{n+1}$$

$$T_0 = 1 \quad T_1 = 1 \quad T_2 = 2 \quad T_3 = 3 \quad T_n = T_{n-1} + T_{n-2}$$

Generating function: "sum" of all possible $2 \times n$ tilings

$$T = 1 + \square + \square\square + \square\square\square + \square\square\square\square + \dots$$

$\hookrightarrow 1 \times \square = \square = \square \times 1 \rightarrow$ multiplicative identity \hookrightarrow combinatorial objects

$$= 1 + \square(1 + \square + \square\square + \dots) + \square\square(1 + \square + \dots)$$

$$= 1 + \square T + \square\square T \Rightarrow T = 1 + \square T + \square\square T \Rightarrow T - \square T - \square\square T = 1$$

$$\Rightarrow T(1 - \square - \square\square) = 1 \Rightarrow T = \frac{1}{1 - \square - \square\square} = 1 + (\square + \square\square) + (\square + \square\square)^2 + \dots$$

$$= 1 + (\square + \square\square) + (\square\square\square + \square\square\square\square + \square\square\square\square\square + \dots)$$

Compress:



$$T = 1 + \square + \square^2 + \square^3 + \square^4 + 2\square^2\square^2 + \square^4 + 3\square^2\square^2 + \square^4 + \dots$$

$$T = \frac{1}{1 - (\square + \square^2)} = 1 + (\square + \square^2) + (\square + \square^2)^2 + \dots$$

$$= \sum_{k \geq 0} (\square + \square^2)^k$$

$$= \sum_{k \geq j \geq 0} \binom{k}{j} \square^j \square^{2k-2j} \quad [m = k-j \Rightarrow k = j+m]$$

$$= \sum_{j,m \geq 0} \binom{j+m}{j} \square^j \square^{2m}$$

Hence, $\binom{j+m}{j}$ is the number of ways to tile a $2 \times (j+2m)$ rectangle with j vertical & $2m$ horizontal dominos.

$n = j+2m \Rightarrow j+m = n-m$; Now $\binom{j+m}{j} = \binom{n-m}{m}$, which is the number of tilings with exactly m pairs of horizontal dominos.

$$T_n = \sum_{0 \leq m \leq \lfloor \frac{n}{2} \rfloor} \binom{n-m}{m}$$

$$0 \leq 2m \leq n$$

$$\Rightarrow 0 \leq m \leq \lfloor \frac{n}{2} \rfloor$$

Self-study: $3 \times n$ rectangle.

Lect. 18: Binomial coefficients and dominoes

GFs with two variables: $f(n, k) = f(n-1, k) + f(n-1, k-1)$ [$f(n, 0) = 1; k \leq n$]

For each $n=0, 1, 2, \dots$, let $B_n(x) = \sum_{k \geq 0} f(n, k) x^k = f(n, 0) + f(n, 1)x + f(n, 2)x^2 + \dots$
 $B_{n-1}(x) = f(n-1, 0) + f(n-1, 1)x + f(n-1, 2)x^2 + \dots$

Multiply the recurrence by x^k & sum over all $k \geq 1$

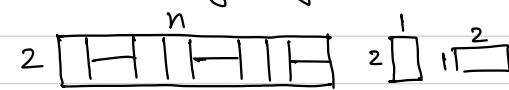
$$\text{LHS: } \sum_{k \geq 1} f(n, k) x^k = \{f(n, 0) + f(n, 1)x + f(n, 2)x^2 + \dots\} - f(n, 0) = B_n(x) - 1$$

$$\begin{aligned} \text{RHS: } & \sum_{k \geq 1} f(n-1, k) x^k + \sum_{k \geq 1} f(n-1, k-1) x^k \\ &= \{f(n-1, 0) + f(n-1, 1)x + f(n-1, 2)x^2 + \dots\} - f(n-1, 0) \\ &+ f(n-1, 0)x + f(n-1, 1)x^2 + \dots \\ &= B_{n-1}(x) - 1 + x B_{n-1}(x) \end{aligned}$$

$$\begin{aligned} \Rightarrow B_n(x) - 1 &= B_{n-1}(x) - 1 + x B_{n-1}(x) \Rightarrow B_n(x) = (1+x) B_{n-1}(x) \\ \Rightarrow B_n(x) &= (1+x)^n B_0(x) \quad [B_0(x) = \sum_{k \geq 0} f(0, k) x^k] \\ &\quad \left. \begin{array}{l} \stackrel{k \geq 0}{=} f(0, 0) = 1 \\ \stackrel{k \geq 1}{=} (1+x)(1+x) B_{n-1}(x) \\ \vdots \\ \stackrel{n}{=} (1+x)^n B_0(x) \end{array} \right\} \\ &= (1+x)^n \quad [\text{Binomial function}] \end{aligned}$$

$$[x^k] B_n(x) = \binom{n}{k} = \frac{n!}{(n-k)! k!} ; \binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}$$

How many ways T_n to cover a $2 \times n$ rectangle with 2×1 dominoes?



$$n=0, n=1, n=2, n=3, T_0=1, T_1=1, T_2=2, T_3=3, T_n = T_{n-1} + T_{n-2}, T_n = f_{n+1}$$

Generating function: "sum" of all possible $2 \times n$ tilings.

$(1 \times \square = \square = \square \times 1) \rightarrow$ multiplicative identity \rightarrow combinatorial object

$$\begin{aligned} T &= 1 + \square + \square\square + \square\square + \square\square\square + \square\square\square + \square\square\square\square + \dots \\ &= 1 + \square(1 + \square + \square\square + \square\square + \dots) + \square\square(1 + \square + \dots) = 1 + \square T + \square\square T \\ \Rightarrow T - \square T - \square\square T &= 1 \Rightarrow T(1 - \square - \square\square) = 1 \Rightarrow T = \frac{1}{1 - \square - \square\square} \end{aligned}$$

$$\Rightarrow T = 1 + (\square + \square\square) + (\square + \square\square)^2 + \dots = 1 + (\square + \square\square) + (\square\square\square + \square\square\square + \square\square\square\square + \square\square\square\square) + \dots$$

Compress:

$$T = 1 + \square + \square^2 + \square^3 + \dots + 2\square\Box^2 + \square^4 + 3\square^2\Box^2 + \square^6 + \dots$$

$$\begin{aligned} T &= \frac{1}{1 - (\square + \square\square)} = 1 + (\square + \square\square) + (\square + \square\square)^2 + (\square + \square\square)^3 + \dots \\ &= \sum_{k \geq 0} (\square + \square\square)^k & k-j = m \Rightarrow k = j+m \\ &= \sum_{k \geq j, j \geq 0} \binom{k}{j} \square^j \square^{2k-2j} & n = j+2m \\ &= \sum_{j, m \geq 0} \binom{j+m}{j} \square^j \square^{2m} \end{aligned}$$

Hence, $\binom{j+m}{j}$ is the number of ways to tile a $2 \times (j+2m)$ rectangle with j vertical dominoes and $2m$ horizontal dominoes.
 $n = j+2m \Rightarrow j+m = n-m$; Now, $\binom{j+m}{j} = \binom{n-m}{m}$, which is the number of tilings with exactly m pairs of horizontal dominoes.

$$T_n = \sum_{0 \leq m \leq \lfloor \frac{n}{2} \rfloor} \binom{n-m}{m}$$

$$\begin{aligned} 0 &\leq 2m \leq n \\ \Rightarrow 0 &\leq m \leq \lfloor \frac{n}{2} \rfloor \end{aligned}$$

Self-study: $3 \times n$ rectangle

Lect 19: Coin change and integer partitioning

How many ways are there to pay 50 cents?

① pennies, ② nickels, ③ dimes, ④ quarters, ⑤ half.

Only pennies: $P = 1 + 1 + 1 + \dots = 1 + 1 + 1^2 + 1^3 + \dots$

+ Nickels: $N = P + ②P + ②^2P + \dots = (1 + ② + ②^2 + \dots)P$

+ Dimes: $D = (10 + ③ + ③^2 + \dots)N$

+ Quarter: $Q = (25 + ④ + ④^2 + \dots)D$

+ Half: $H = (50 + ⑤ + ⑤^2 + \dots)Q$

$$P = 1 + z + z^2 + z^3 + \dots \Rightarrow P = \frac{1}{1-z} \Rightarrow P(1-z) = 1 \quad P_0 + P_1 z + P_2 z^2 + \dots + P_m z^m + P_{m+1} z^{m+1} + \dots$$

$$N = (1 + z^5 + z^{10} + z^{15} + \dots)P \Rightarrow N = \frac{P}{1-z^5} \Rightarrow N(1-z^5) = P$$

$$D = (1 + z^{10} + z^{20} + z^{30} + \dots)N \Rightarrow D = \frac{N}{1-z^{10}} \Rightarrow D(1-z^{10}) = N$$

$$Q = (1 + z^{25} + z^{50} + z^{75} + \dots)D \Rightarrow Q = \frac{D}{1-z^{25}} \Rightarrow Q(1-z^{25}) = D$$

$$H = (1 + z^{50} + z^{100} + z^{150} + \dots)Q \Rightarrow H = \frac{Q}{1-z^{50}} \Rightarrow H(1-z^{50}) = Q$$

Let P_n, N_n, D_n, Q_n, H_n be the number of ways to pay n cents when max coin allowed are 1, 5, 10, 25, 50, respectively. [coefficients of z^n in the series above]

$$P_n - P_{n-1} = [n=0] \Rightarrow P_n = P_{n-1} + [n=0]$$

$$N_n - N_{n-5} = P_n \Rightarrow N_n = N_{n-5} + P_n$$

$$D_n - D_{n-10} = N_n \Rightarrow D_n = D_{n-10} + N_n$$

$$Q_n - Q_{n-25} = D_n \Rightarrow Q_n = Q_{n-25} + D_n$$

$$H_n - H_{n-50} = Q_n \Rightarrow H_n = H_{n-50} + Q_n$$

$$H = \frac{Q}{1-z^{50}} = \frac{D}{1-z^{25}} \cdot \frac{1}{1-z^{50}} = \frac{N}{1-z^{10}} \cdot \frac{1}{1-z^{25}} \cdot \frac{1}{1-z^{50}} = \frac{P}{1-z^5} \cdot \frac{1}{1-z^{10}} \cdot \frac{1}{1-z^{25}} \cdot \frac{1}{1-z^{50}}$$

$$= \frac{1}{1-z} \cdot \frac{1}{1-z^5} \cdot \frac{1}{1-z^{10}} \cdot \frac{1}{1-z^{25}} \cdot \frac{1}{1-z^{50}}$$

A partition of a positive integer n is a way of writing n as a sum of positive integers (each integer in the sum is called a part). For example, 4 can be partitioned into five distinct ways: $4, 3+1, 2+2, 2+1+1, 1+1+1+1$.

Show that for each positive integer n , the number of partitions of n into unequal parts is equal to the number of partitions of n into odd parts. For instance, if $n=6$, there are 4 partitions into unequal parts: $6, 5+1, 4+2, 3+2+1$. And there are four partitions into odd parts: $5+1, 3+3, 3+1+1+1, 1+1+1+1+1$.

GF for partitioning into odd parts: $\frac{1}{1-x} \cdot \frac{1}{1-x^3} \cdot \frac{1}{1-x^5} \cdot \frac{1}{1-x^7} \cdots$

GF for partitioning into unequal parts: $(1+x)(1+x^2)(1+x^3)\cdots$
 $= \frac{1-x^2}{1-x} \cdot \frac{1-x^4}{1-x^2} \cdot \frac{1-x^6}{1-x^3} \cdot \frac{1-x^8}{1-x^4} \cdot \frac{1-x^{10}}{1-x^5} \cdots$

$$\# a_n = \begin{cases} 1, & \text{if } n=0 \\ 4, & \text{if } n=1 \\ a_{n-1} + a_{n-2} \\ 2a_{n-2} - a_{n-1} \end{cases}$$

$$a_n = \frac{4}{7n+4} \quad \text{Top-down, bottom-up both work!}$$

$$[n-2, n-3, n-4]$$

$$\# \sum_{k=1}^n (-1)^k (2k-1) \rightarrow \begin{array}{l} \text{Calculate first few terms, generalize, prove by induction} \\ \text{Convert to recurrence, solve using repertoire} \\ \text{perturbation} \\ \text{Convert to multiple sums} \end{array}$$

Lect. 19: Coin change and integer partitioning

How many ways are there to pay 50 cents? (denominations)

① pennies, ② nickels ③ dimes, ④ quarters ⑤ half.

Only pennies: $P = 1 + 1 + 1 \cdot 1 + 1 \cdot 1 \cdot 1 + \dots = 1 + 1 + 1^2 + 1^3 + \dots$ + Nickels: $N = P + ②P + ②^2P + \dots = (1 + ② + ②^2 + \dots)P$ + Dimes: $D = (10 + ⑩ + ⑩^2 + \dots)N$ + Quarter: $Q = (25 + ④ + ④^2 + \dots)D$ + Half: $H = (50 + ⑤ + ⑤^2 + \dots)Q$

$$P = P_0 + P_1 z + P_2 z^2 + P_3 z^3 + \dots \Rightarrow P_2 = P_0 z + P_1 z^2 + \dots \\ P_3 = P_0 z^3 + \dots$$

$$P = 1 + z + z^2 + z^3 + \dots \Rightarrow P = \frac{1}{1-z} \Rightarrow P(1-z) = 1$$

$$N = (1 + z^5 + z^{10} + \dots)P \Rightarrow N = \frac{P}{1-z^5} \Rightarrow N(1-z^5) = P$$

$$D = (1 + z^{10} + z^{20} + \dots)N \Rightarrow D = \frac{N}{1-z^{10}} \Rightarrow D(1-z^{10}) = N$$

$$Q = (1 + z^{25} + z^{50} + \dots)D \Rightarrow Q = \frac{D}{1-z^{25}} \Rightarrow Q(1-z^{25}) = D$$

$$H = (1 + z^{50} + z^{100} + \dots)Q \Rightarrow H = \frac{Q}{1-z^{50}} \Rightarrow H(1-z^{50}) = Q$$

Let P_n, N_n, D_n, Q_n, H_n be the number of ways to pay n cents when max coin allowed are 1, 5, 10, 25, 50, respectively. (by defn, coefficients of z^n in the series above)

Equating the coefficients of z^n in both sides:

$$P_n - P_{n-1} = [n=0] \Rightarrow P_n = P_{n-1} + [n=0]$$

$$N_n - N_{n-5} = P_n \Rightarrow N_n = N_{n-5} + P_n$$

$$D_n - D_{n-10} = N_n \Rightarrow D_n = D_{n-10} + N_n$$

$$Q_n - Q_{n-25} = D_n \Rightarrow Q_n = Q_{n-25} + D_n$$

$$H_n - H_{n-50} = Q_n \Rightarrow H_n = H_{n-50} + Q_n$$

$$H = \frac{Q}{1-z^{50}} = \frac{D}{1-z^{25}} \cdot \frac{1}{1-z^{50}} = \frac{N}{1-z^{10}} \cdot \frac{1}{1-z^{25}} \cdot \frac{1}{1-z^{50}} = \frac{P}{1-z^5} \cdot \frac{1}{1-z^{10}} \cdot \frac{1}{1-z^{25}} \cdot \frac{1}{1-z^{50}}$$

$$= \frac{1}{1-z} \cdot \frac{1}{1-z^5} \cdot \frac{1}{1-z^{10}} \cdot \frac{1}{1-z^{25}} \cdot \frac{1}{1-z^{50}}$$

A partition of a positive integer n is a way of writing n as a sum of positive integers (each integer in the sum is called a part). For example, 4 can be partitioned into five distinct ways: $4, 3+1, 2+2, 2+1+1, 1+1+1+1$.

Show that for each positive integer n , the number of partitions of n into unequal parts is equal to the number of partitions of n into odd parts. For instance, if $n=6$, there are 4 partitions into unequal parts: $6, 5+1, 4+2, 3+2+1$. And there are four partitions into odd parts: $5+1, 3+3, 3+1+1+1, 1+1+1+1+1$.

GF for partitioning into odd parts: $\frac{1}{1-z} \cdot \frac{1}{1-z^3} \cdot \frac{1}{1-z^5} \cdot \frac{1}{1-z^7} \cdots$

GF for partitioning into unequal parts: $(1+z)(1+z^2)(1+z^3) \cdots$

$$= \frac{1+z^2}{1-z} \cdot \frac{1+z^4}{1-z^2} \cdot \frac{1+z^6}{1-z^3} \cdot \frac{1+z^{10}}{1-z^5} \cdots = \frac{1}{1-z} \cdot \frac{1}{1-z^2} \cdot \frac{1}{1-z^5} \cdots$$

$$\# a_n = \begin{cases} 1, & \text{if } n=0 \\ 4/11, & \text{if } n=1 \\ \frac{a_{n-1}a_{n-2}}{2a_{n-2}-a_{n-1}}, & [n-2, n-3, n-4] \end{cases}$$

Top-down, bottom-up both work!

$$\# \sum_{k=1}^n (-1)^{k+1} (2k-1)$$

- calculate first few terms, generalize, prove by induction
- convert to recurrence, solve using repertoire
- perturbation
- convert to multiple sum