

Forkable-Strings

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Contents

theory *Forkable-Strings* **imports** *Main* $\sim\sim$ */src/HOL/List* **begin**

We will use True as 1 and False as 0 in characteristic strings; we might think about each bool value answers to the question 'is this slot controlled by an adversarial player?'.
datatype *nattree* = *Empty* | *Node* *nat* *nattree* *list*

fun *lt-nat-tree* :: *nat* \Rightarrow *nattree* \Rightarrow *bool* **where**
 lt-nat-tree *n* *Empty* = *True*
 | *lt-nat-tree* *n* (*Node* *m* *-*) = (*n* < *m*)

lemma *lt-nat-tree-lt* [*simp*]: (*n* < *m*) \longleftrightarrow *lt-nat-tree* *n* (*Node* *m* *l*)
 by *simp*

lemma *lt-nat-tree-ge* [*simp*]: (*n* \geq *m*) \longleftrightarrow \neg *lt-nat-tree* *n* (*Node* *m* *l*)
 by *auto*

fun *increasing-tree* :: *nattree* \Rightarrow *bool* **where**
 increasing-tree *Empty* = *True*
 | *increasing-tree* (*Node* *-* []) = *True*
 | *increasing-tree* (*Node* *n* *l*) = ($\forall x \in \text{set } l. \text{increasing-tree } x \wedge \text{lt-nat-tree } n \ x$)

lemma *increasing-tree-empty-branch-list* [*simp*]: *increasing-tree* (*Node* *n* [])
 by *simp*

lemma *increasing-tree-ind* [*simp*] : ($\forall x \in \text{set } l. \text{increasing-tree } x \wedge \text{lt-nat-tree } n \ x$) \longleftrightarrow *increasing-tree* (*Node* *n* *l*)
 by (*smt* *increasing-tree.elims*(2) *increasing-tree.elims*(3) *length-greater-0-conv* *length-pos-if-in-set* *nattree.distinct*(1) *nattree.inject*)

definition *ListMax* :: *nat* *list* \Rightarrow *nat* **where**
 ListMax *l* = *foldr* *max* *l* 0

lemma *max-of-the-list-0* [*simp*]: *ListMax* [] = 0

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by (simp add: ListMax-def)

lemma max-of-the-list [simp]:  $\forall x \in \text{set } l. x \leq \text{ListMax } l$ 
proof (induction l)
  case Nil
  then show ?case
    by auto
next
  case (Cons a l)
  have ListMax (Cons a l) = max a (ListMax l)
    using ListMax-def by auto
  have ListMax l  $\leq$  ListMax (Cons a l)  $\wedge$   $a \leq$  ListMax (Cons a l)
    by (simp add:  $\langle \text{ListMax } (a \# l) = \max a (\text{ListMax } l) \rangle$ )
  then show ?case
    using Cons.IH by auto
qed

fun height-branch :: nattree  $\Rightarrow$  nat where
  height-branch Empty = 0
| height-branch (Node n l) = 1 + ListMax (map height-branch l)

lemma height-branch-empty-list [simp]: height-branch (Node n []) = 1
  by simp

lemma height-branch-lt [simp]:  $\forall x \in \text{set } l. \text{height-branch } x < \text{height-branch } (\text{Node } n \ l)$ 
  by (simp add: le-imp-less-Suc)

fun height :: nattree  $\Rightarrow$  nat where
  height Empty = 0
| height (Node n bl) = ListMax (map height-branch bl)

lemma height-empty-list [simp]: height (Node n []) = 0
  by simp

lemma height-ge [simp]:  $\forall x \in \text{set } l. \text{height-branch } x \leq \text{height } (\text{Node } n \ l)$ 
  by auto

fun depth-branch :: nat  $\Rightarrow$  nattree  $\Rightarrow$  nat where
  depth-branch m Empty = 0
| depth-branch m (Node n l) = (if m = n then 1 else
                                (case ListMax (map (depth-branch m) l) of
                                 0  $\Rightarrow$  0 |
                                 Suc n  $\Rightarrow$  Suc (Suc n)))

lemma depth-branch-eq [simp]: depth-branch m (Node m l) = 1
  by simp

lemma listmax-0 [simp]:  $(\forall x \in \text{set } l. f x = 0) \longrightarrow \text{ListMax } (\text{map } f \ l) = 0$ 

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proof (induction l)
  case Nil
  then show ?case by simp
next
  case (Cons a l)
  have ListMax (map f (Cons a l)) = max (f a) (ListMax (map f l))
    using ListMax-def by auto
  then have (f a = 0)  $\wedge$  ( $\forall x \in \text{set } l. f x = 0$ )  $\longrightarrow$  ListMax (map f (Cons a l))
    = 0
    using Cons.IH by linarith
  then show ?case
    by simp
qed

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lemma depth-branch-empty-branch-list [simp]: depth-branch m (Node n []) = of-bool
(m = n)
proof (cases m = n)
  case True
  then show ?thesis by simp
next
  case False
  then show ?thesis by auto
qed

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lemma depth-branch-ne-nf [simp]:
( $\forall x \in \text{set } l. \text{depth-branch } m x = 0$ )  $\wedge m \neq n \longrightarrow \text{depth-branch } m (\text{Node } n l) = 0$ 
by simp

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fun depth :: nattree  $\Rightarrow$  nat  $\Rightarrow$  nat where
  depth - 0 = 0
| depth Empty n = 0
| depth (Node n bl) m = ListMax (map (depth-branch m) bl)

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lemma depth-root-ge [simp] :  $m > 0 \longrightarrow (\forall x \in \text{set } l. \text{depth-branch } m x \leq \text{depth} (\text{Node } n l) m)$ 
using gr0-conv-Suc image-eqI by auto

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fun getFrom' :: bool list  $\Rightarrow$  nat  $\Rightarrow$  bool where
  getFrom' [] n = True
| getFrom' (x#xs) 0 = x
| getFrom' (x#xs) (Suc n) = getFrom' xs n

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definition H :: bool list  $\Rightarrow$  nat set where
  H l = {x.  $\neg \text{getFrom}' (\text{False}\#l) x$ }

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lemma *H-0* [*simp*]: $0 \in H\ l$
by (*simp add: H-def*)

lemma *getFrom-suc-eq-H* [*simp*]: $\neg \text{getFrom}'\ l\ x \longleftrightarrow \text{Suc}\ x \in H\ l$
using *H-def* **by** *auto*

fun *ListSum* :: $\text{nat list} \Rightarrow \text{nat}$ **where**
ListSum *l* = *foldr plus l 0*

fun *count-node* :: $\text{nat} \Rightarrow \text{nattree} \Rightarrow \text{nat}$ **where**
count-node - *Empty* = 0
| *count-node* *m* (*Node n bl*) = (*of-bool* ($m = n$)) + *ListSum* (*map* (*count-node* *m*) *bl*)

lemma *count-node-empty-branch-list* [*simp*]: *count-node* *m* (*Node n []*) = *of-bool* ($m = n$) **by** *simp*

definition *unique-node* :: $\text{nattree} \Rightarrow \text{nat} \Rightarrow \text{bool}$ **where**
unique-node *t n* = (*count-node* *n t* = 1)

fun *unique-nodes-by-nat-set* :: $\text{nattree} \Rightarrow \text{nat set} \Rightarrow \text{bool}$ **where**
unique-nodes-by-nat-set *t s* = ($\forall x \in s. \text{unique-node}\ t\ x$)

definition *uniqueH-tree* :: $\text{nattree} \Rightarrow \text{bool list} \Rightarrow \text{bool}$ **where**
uniqueH-tree *t l* = *unique-nodes-by-nat-set* *t* (*H l*)

lemma *uniqueH-tree-in-imp-l* [*simp*]: $\forall x \in H\ l. \text{uniqueH-tree}\ t\ l \longrightarrow \text{unique-node}\ t\ x$
using *uniqueH-tree-def* **by** *auto*

lemma *uniqueH-tree-in-imp-r* [*simp*]: $(\forall x \in H\ l. \text{unique-node}\ t\ x) \longrightarrow \text{uniqueH-tree}\ t\ l$
using *uniqueH-tree-def unique-nodes-by-nat-set.simps* **by** *blast*

fun *max-node* :: $\text{nattree} \Rightarrow \text{nat}$ **where**
max-node *Empty* = 0
| *max-node* (*Node n bl*) = *ListMax* (*n* # (*map* *max-node* *bl*))

lemma *ListSum-0* [*simp*]: $(\forall x \in \text{set}\ l. x = 0) \longrightarrow \text{ListSum}\ l = 0$
proof (*induction l*)
case *Nil*
then show ?*case* **by** *simp*
next
case (*Cons a l*)
then show ?*case*
by *simp*
qed

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lemma max-node-max [simp]:  $\forall m. \text{max-node } t < m \longrightarrow \text{count-node } m \ t = 0$ 
proof (induction t)
  case Empty
  then show ?case
  by simp
next
  case (Node x1 x2)
  have max-node (Node x1 x2) = ListMax (x1 # (map max-node x2))
  by simp
  then have  $\forall x \in \text{set } x2. \text{max-node } x \leq \text{max-node } (\text{Node } x1 \ x2) \wedge x1 \leq \text{max-node } (\text{Node } x1 \ x2)$ 
  by simp
  then have  $\forall x. \forall y \in \text{set } x2. \text{max-node } (\text{Node } x1 \ x2) < x \longrightarrow \text{max-node } y < x$ 
 $\wedge x1 < x$ 
  using le-less-trans by blast
  then have  $\forall x. \forall y \in \text{set } x2. \text{max-node } (\text{Node } x1 \ x2) < x \longrightarrow \text{count-node } x \ y = 0$ 
  by (simp add: Node.IH)
  then have  $\forall x. \text{max-node } (\text{Node } x1 \ x2) < x \longrightarrow \text{count-node } x \ (\text{Node } x1 \ x2) = \text{ListSum } (\text{map } (\text{count-node } x) \ x2)$ 
  by (smt (max-node (Node x1 x2) = ListMax (x1 # map max-node x2))
add.commute add-cancel-left-right count-node.simps(2) less-irrefl-nat less-le-trans list.set-intros(1) max-of-the-list of-bool-def)
  then show ?case
  using ListSum-0 ( $\forall x. \forall y \in \text{set } x2. \text{max-node } (\text{Node } x1 \ x2) < x \longrightarrow \text{count-node } x \ y = 0$ ) by auto
qed

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fun increasing-depth-H :: nattree  $\Rightarrow$  bool list  $\Rightarrow$  bool where
  increasing-depth-H t l = ( $\forall x \in H \ l. \forall y \in H \ l. x < y \longrightarrow \text{depth } t \ x < \text{depth } t \ y$ )

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fun root-label0 :: nattree  $\Rightarrow$  bool where
  root-label0 Empty = False
| root-label0 (Node 0 -) = True
| root-label0 (Node (Suc n) -) = False

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F — w

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fun isFork :: bool list  $\Rightarrow$  nattree  $\Rightarrow$  bool where
  isFork w F = ((length w  $\geq$  max-node F)
     $\wedge$  (increasing-tree F)
     $\wedge$  (uniqueH-tree F w)
     $\wedge$  (increasing-depth-H F w)
     $\wedge$  root-label0 F)

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lemma isFork-max-not-exceed [simp] : isFork w F  $\longrightarrow$  length w  $\geq$  max-node F by
simp

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lemma isFork-root-0 [simp] : isFork w F  $\longrightarrow$  root-label0 F by simp

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lemma *isFork-increasing-tree* [simp] : *isFork* *w F* \longrightarrow *increasing-tree* *F*
using *isFork.simps* **by** *blast*

lemma *isFork-uniqueH-tree* [simp] : *isFork* *w F* \longrightarrow $(\forall x \in H \ w. \text{unique-node } F \ x)$
by (*meson isFork.elims*(2) *uniqueH-tree-in-imp-l*)

lemma *isFork-increasing-depth-H* [simp] : *isFork* *w F* \longrightarrow $(\forall x \in H \ w. \forall y \in H \ w. \ x < y \longrightarrow \text{depth } F \ x < \text{depth } F \ y)$
by (*meson increasing-depth-H.elims*(2) *isFork.elims*(2))

fun *flatTree* :: *nattree* \Rightarrow *bool* **where**
flatTree Empty = *False*
| *flatTree (Node - [])* = *True*
| *flatTree (Node 0 l)* = (*Suc (Suc 0)* \leq
foldr (($\lambda x. (\lambda y. (\text{if } x = (\text{height } (\text{Node } 0 \ l)) \text{ then } \text{Suc } y \text{ else } y))$)) (*map height-branch*
l) 0)
| *flatTree (Node (Suc n) l)* = *False*

definition *isForkable* :: *bool list* \Rightarrow *bool* **where**
isForkable w = $(\exists F. (\text{isFork } w \ F) \wedge \text{flatTree } F)$

fun *order-map* :: $(\text{nat} \Rightarrow 'b \Rightarrow 'c \text{ list}) \Rightarrow 'b \text{ list} \Rightarrow \text{nat} \Rightarrow 'c \text{ list}$ **where**
order-map f [] = []
| *order-map f (x#xs) n* = *f n x @ (order-map f xs (Suc n))*

fun *order-map-disjoint* :: $(\text{nat} \Rightarrow 'b \Rightarrow 'c \text{ list}) \Rightarrow \text{nat} \Rightarrow 'b \text{ list} \Rightarrow ('c \text{ list}) \text{ list}$
where
order-map-disjoint f - [] = []
| *order-map-disjoint f n (x#xs)* = *f n x # (order-map-disjoint f (Suc n) xs)*

function *firstN-R* :: $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat list}$ **where**
firstN-R i n = (*if i* \geq *n* *then []* *else i#(firstN-R (Suc i) n)*)
apply *auto[1]* **by** *blast*

termination *firstN-R*
apply (*relation measure* ($\lambda(i,n). \ n - i$))
apply *simp*
by *simp*

definition *firstN* :: $\text{nat} \Rightarrow \text{nat list}$ **where**
firstN n = *firstN-R 0 n*

function *tinelist-nat* :: $\text{nat} \Rightarrow \text{nattree} \Rightarrow (\text{nat list}) \text{ list}$ **where**
tinelist-nat - Empty = [[]]
| *tinelist-nat m (Node - [])* = [[*m*]]
| *tinelist-nat m (Node - (x#xs))* = *map (Cons m)*

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(foldr append (map ( $\lambda(n,t).$  tinelist-nat n t) (zip (firstN (length (x#xs))) (x#xs)))
[])
  apply (metis depth-branch.cases neq-Nil-conv)
  by auto

fun total-node :: nattree  $\Rightarrow$  nat where
  total-node Empty = 0
| total-node (Node n l) = Suc (ListSum (map total-node l))

lemma total-nod-dec [simp] :  $\forall x \in \text{set } l. \text{total-node } x < \text{total-node } (\text{Node } n \ l)$ 
proof (induction l)
  case Nil
  then show ?case by auto
next
  case (Cons a l)
  then show ?case by auto
qed

lemma [simp] :  $(a, b) \in \text{set } (\text{zip } (\text{firstN } (\text{Suc } (\text{length } xs))) (x \# xs)) \Rightarrow$ 
  total-node b < Suc (total-node x + foldr op + (map total-node xs) 0)
using total-nod-dec set-zip-rightD by force

termination tinelist-nat
  apply (relation measure ( $\lambda(n,nt).$  total-node nt))
  apply auto
  done

fun tinelist :: nattree  $\Rightarrow$  (nat list) list where
  tinelist Empty = []
| tinelist (Node 0 l) = (foldr append (map ( $\lambda(n,t).$  tinelist-nat n t) (zip (firstN
(length l)) l)) [])
| tinelist (Node (Suc n) l) = []

fun getLabelFromTine :: nattree  $\Rightarrow$  nat list  $\Rightarrow$  nat list where
  getLabelFromTine Empty l = []
| getLabelFromTine - [] = []
| getLabelFromTine (Node - l) (x#xs) = (case hd (drop x l) of
  Empty  $\Rightarrow$  [] | (*it runs out of nodes before we
can trace down all paths*)
  Node n -  $\Rightarrow$  n # getLabelFromTine (hd (drop
x l)) xs)

fun isPrefix-lists :: 'a list  $\Rightarrow$  'a list  $\Rightarrow$  bool where
  isPrefix-lists [] - = True
| isPrefix-lists (l#ls) [] = False
| isPrefix-lists (l#ls) (r#rs) = ((l=r)  $\wedge$  isPrefix-lists ls rs)

definition isPrefix-tines :: nattree  $\Rightarrow$  nattree  $\Rightarrow$  nat list  $\Rightarrow$  nat list  $\Rightarrow$  bool where
  isPrefix-tines nt1 nt2 t1 t2 =

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$(isPrefix-lists\ t1\ t2 \wedge isPrefix-lists\ (getLabelFromTine\ nt1\ t1)\ (getLabelFromTine\ nt2\ t2))$

definition $isPrefix-trees :: nattree \Rightarrow nattree \Rightarrow bool$ **where**
 $isPrefix-trees\ nt1\ nt2 =$
 $(\forall\ t1. (\forall\ t2. ListMem\ t1\ (tinelist\ nt1) \wedge ListMem\ t2\ (tinelist\ nt2)) \longrightarrow isPrefix-tines\ nt1\ nt2\ t1\ t2))$

definition $isPrefix-forks :: bool\ list \Rightarrow bool\ list \Rightarrow nattree \Rightarrow nattree \Rightarrow bool$ **where**
 $isPrefix-forks\ w1\ w2\ nt1\ nt2 =$
 $(isFork\ w1\ nt1 \wedge isFork\ w2\ nt2 \wedge isPrefix-lists\ w1\ w2 \wedge isPrefix-trees\ nt1\ nt2)$

inductive $ListOfEmpty :: nattree\ list \Rightarrow bool$ **where**
 $Nil : ListOfEmpty\ []$
 $| Cons : ListOfEmpty\ l \Longrightarrow ListOfEmpty\ (Empty\#l)$

fun $closedFork-Hgiven :: nattree \Rightarrow nat\ set \Rightarrow bool$ **where**
 $closedFork-Hgiven\ Empty = True$
 $| closedFork-Hgiven\ (Node\ n\ l)\ h = (if\ ListOfEmpty\ l$
 $\quad then\ (n \in h)$
 $\quad else\ foldr\ conj\ (map\ (\lambda x. closedFork-Hgiven\ x\ h)\ l)$
 $\quad True)$

definition $closedFork :: nattree \Rightarrow bool\ list \Rightarrow bool$ **where**
 $closedFork\ t\ w = closedFork-Hgiven\ t\ (H\ w)$

definition $gap :: nattree \Rightarrow nat\ list \Rightarrow nat$ **where**
 $gap\ nt\ tine = height\ nt - (length\ tine)$

definition $reserve :: bool\ list \Rightarrow nat\ list \Rightarrow nat$ **where**
 $reserve\ w\ labeledTine = foldr\ (\lambda x. (plus\ (of-bool\ x)))\ (drop\ (ListMax\ labeledTine)\ w)\ 0$

definition $reach :: nattree \Rightarrow bool\ list \Rightarrow nat\ list \Rightarrow nat$ **where**
 $reach\ nt\ w\ tine = reserve\ w\ tine - gap\ nt\ tine$

definition $lambda :: nattree \Rightarrow bool\ list \Rightarrow nat$ **where**
 $lambda\ t\ w = ListMax\ (map\ (reach\ t\ w)\ (tinelist\ t))$

fun $crosslist :: 'a\ list \Rightarrow 'a\ list \Rightarrow (('a, 'a)\ prod)\ list$ **where**
 $crosslist\ [] = []$
 $| crosslist\ (x\#xs)\ ys = (map\ (Pair\ x)\ ys) @ (crosslist\ xs\ ys)$

fun *cross-all-pair'* :: ('a list) list \Rightarrow ('a list, ('a, 'a) prod) list prod **where**
cross-all-pair' [] = ([], [])
| *cross-all-pair'* (x#xs) = (x @ fst (*cross-all-pair'* xs),
(*crosslist* x (fst (*cross-all-pair'* xs))) @ snd (*cross-all-pair'* xs))

definition *cross-all-pair* :: ('a list) list \Rightarrow (('a, 'a) prod) list **where**
cross-all-pair l = snd (*cross-all-pair'* l)

fun *list-of-disjoint-edged-tines* :: nattree \Rightarrow ((nat list, nat list) prod) list **where**
list-of-disjoint-edged-tines Empty = []
| *list-of-disjoint-edged-tines* (Node n l)
= *cross-all-pair* (order-map-disjoint ($\lambda x.$ map (Cons x)) 0 (map tinelist l))

definition *margin* :: nattree \Rightarrow bool list \Rightarrow nat **where**
margin t w = foldr ($\lambda(a,b).$ max (min (reach t w a) (reach t w b))) (*list-of-disjoint-edged-tines* t) (0 - (height t))

proposition *proposition-4-17* : *isForkable* w \longleftrightarrow ($\exists F.$ (*isFork* w F \wedge *margin* F w \geq 0))
sorry

definition *lambda-of-string* :: bool list \Rightarrow nat **where**
lambda-of-string w = (GREATEST t. ($\exists F.$ (*isFork* w F \wedge *closedFork* F w \wedge t = *lambda* F w)))

definition *margin-of-string* :: bool list \Rightarrow nat **where**
margin-of-string w = (GREATEST t. ($\exists F.$ (*isFork* w F \wedge *closedFork* F w \wedge t = *margin* F w)))

definition *m* :: bool list \Rightarrow (nat, nat) prod **where**
m w = (*lambda-of-string* w, *margin-of-string* w)

lemma *lemma-4-18* : (*m* [] = (0,0)) \wedge
($\forall w.$ ((length w > 0) \longrightarrow (
(*m* (w @ [True]) = (*lambda-of-string* w + 1, *margin-of-string* w + 1))
 \wedge ((*lambda-of-string* w > *margin-of-string* w) \wedge (*margin-of-string* w = 0)
 \longrightarrow (*m* (w @ [False]) = (*lambda-of-string* w - 1, 0)))
 \wedge (*lambda-of-string* w = 0 \longrightarrow (*m* (w @ [False]) = (0, *margin-of-string* w - 1))))
 \wedge (*lambda-of-string* w > 0 \wedge *margin-of-string* w \neq 0 \longrightarrow (*m* (w @ [False])
= (*lambda-of-string* w - 1, *margin-of-string* w - 1))))))
 \wedge ($\exists F.$ (*isFork* w F \wedge *closedFork* F w \wedge (*m* w = (*lambda* F w, *margin* F w))))))
sorry

end