





MOVERY: A Precise Approach for Modified Vulnerable Code Clone Discovery from Modified Open-Source Software Components

Seunghoon Woo, Hyunji Hong, Eunjin Choi, Heejo Lee

Korea University

USENIX security 2022

Modified open-source software reuse is prevalent

- Reuse of open-source software (OSS) becomes a trend in software development
- Unmanaged OSS reuse can pose security threats (e.g., vulnerability propagation)
- Most developers reuse OSS projects with code/structural modifications*
 - Hard to discover propagated vulnerable codes with code changes

How can we precisely discover propagated vulnerable codes with various syntaxes?

^{* [}CCS 2017] "Identifying Open-Source License Violation and 1-day Security Risk at Large Scale", Ruian Duan, Ashish Bijlani, Meng Xu, Taesoo Kim, and Wenke Lee
* [ICSE 2021] "CENTRIS: A Precise and Scalable Approach for Identifying Modified Open-Source Software Reuse", Seunghoon Woo, Sunghan Park, Seulbae Kim, Heejo Lee, and Hakjoo Oh

Addressing syntax diversity of vulnerable code

- Syntax diversity of vulnerable code
 - Internal modification of OSS
 - OSS source code frequently changes during OSS updates
 - ❖ Vulnerable code may exist in various syntax depending on the reused OSS version
 - External modification of OSS
 - ❖ Vulnerable code can be modified during the OSS reuse process

Example: Syntax diversity caused by the internal OSS modification

- CVE-2014-5461 vulnerability in Lua (DoS vulnerability)
 - This vulnerable code existed in Redis (using Lua)
 - The syntax of the two vulnerable functions is quite different

```
int luaD_precall (lua_State *L, StkId func, int nresults) {
    lua_CFunction f;
    CallInfo *ci;
    int n; /* number of arguments (Lua) or returns (C) */
    ptrdiff_t funcr = savestack(L, func);
    switch (ttype(func)) {
        ...
        case LUA_TLCL: { /* Lua function: prepare its call */
            StkId base;
        Proto *p = clLvalue(func)->p;
        - luaD_checkstack(L, p->maxstacksize);
        - func = restorestack(L, funcr);
        n = cast_int(L->top - func) - 1;
        + luaD_checkstack(L, p->maxstacksize);
```

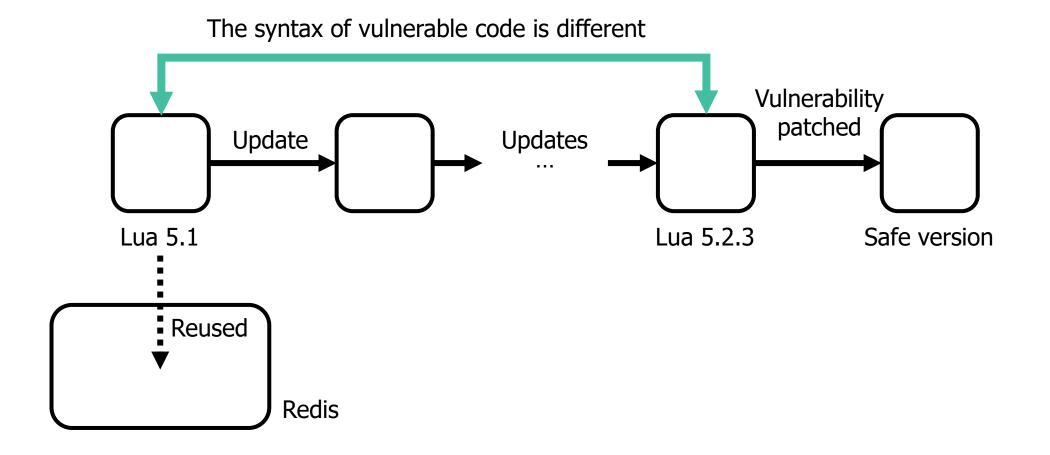
```
int luaD_precall (lua_State *L, StkId func, int nresults) {
    LClosure *cl;
    ptrdiff_t funcr;
    if (!ttisfunction(func)) /* `func' is not a function? */
        func = tryfuncTM(L, func); /* check the `function' tag method */
    funcr = savestack(L, func);
    cl = &clvalue(func)->l;
    L->ci->savedpc = L->savedpc;
    if (!cl->isC) {        /* Lua function? prepare its call */
        CallInfo *ci;
        StkId st, base;
        Proto *p = cl->p;
        - luaD_checkstack(L, p->maxstacksize);
        + luaD_checkstack(L, p->maxstacksize + p->numparams);
```

A patch snippet for CVE-2014-5461 in Lua 5.2.3

A patch snippet for CVE-2014-5461 in Redis (using Lua 5.1)

Example: Syntax diversity caused by the internal OSS modification

• CVE-2014-5461 vulnerability in Lua (DoS vulnerability)



Example: Syntax diversity caused by the internal OSS modification

- Existing approaches fail to detect this propagated vulnerable code
 - ReDeBug [S&P 2012]
 - Considering nearby three (by default) lines of deleted and added code lines from the patch

```
int luaD precall (lua State *L, StkId func, int nresults) {
                                                                int luaD precall (lua State *L, StkId func, int nresults) {
 lua CFunction f;
                                                                  LClosure *cl;
 CallInfo *ci;
                                                                  ptrdiff t funcr;
 int n; /* number of arguments (Lua) or returns (C) */
                                                                  if (!ttisfunction(func)) /* `func' is not a function? */
 ptrdiff t funcr = savestack(L, func);
                                                                    func = tryfuncTM(L, func); /* check the `function' tag method */
 switch (ttype(func)) {
                                                                  funcr = savestack(L, func);
                                                                  cl = &clvalue(func)->l;
   case LUA TLCL: { /* Lua function: prepare its call */
                                                                  L->ci->savedpc = L->savedpc;
      StkId base;
                                                                  if (!cl->isC) { /* Lua function? prepare its call */
                                                                    CallInfo *ci;
      Proto *p = clLvalue(func)->p;
    luaD checkstack(L, p->maxstacksize);
                                                                    StkId st, base;
    - func = restorestack(L, funcr);
                                                                    Proto *p = cl->p;
      n = cast_int(L->top - func) - 1;
                                                                   - luaD_checkstack(L, p->maxstacksize);
                                                                    luaD checkstack(L, p->maxstacksize + p->numparams);
    + luaD checkstack(L, p->maxstacksize);
```

Example: Syntax diversity caused by the internal OSS modification

- Existing approaches fail to detect this propagated vulnerable code
 - VUDDY [S&P 2017]
 - Considering a whole vulnerable function

```
int luaD_precall (lua_State *L, StkId func, int nresults) {
    lua_CFunction f;
    CallInfo *ci;
    int n; /* number of arguments (Lua) or returns (C) */
    ptrdiff_t funcr = savestack(L, func);
    switch (ttype(func)) {
        ...
        case LUA_TLCL: { /* Lua function: prepare its call */
            StkId base;
            Proto *p = clLvalue(func)->p;
        - luaD_checkstack(L, p->maxstacksize);
        - func = restorestack(L, funcr);
        n = cast_int(L->top - func) - 1;
        + luaD_checkstack(L, p->maxstacksize);
```

```
int luaD_precall (lua_State *L, StkId func, int nresults) {
   LClosure *cl;
   ptrdiff_t funcr;
   if (!ttisfunction(func)) /* `func' is not a function? */
      func = tryfuncTM(L, func); /* check the `function' tag method */
   funcr = savestack(L, func);
   cl = &clvalue(func)->l;
   L->ci->savedpc = L->savedpc;
   if (!cl->isC) {      /* Lua function? prepare its call */
        CallInfo *ci;
      StkId st, base;
      Proto *p = cl->p;
      - luaD_checkstack(L, p->maxstacksize);
      + luaD_checkstack(L, p->maxstacksize + p->numparams);
```

MOVERY: A Precise Approach for Modified Vulnerable Code Clone Discovery from Modified Open-Source Software Components

Design of MOVERY

MOdified Vulnerable code clone discovERY

- A novel approach to precisely detect modified vulnerable code clones
- Key techniques
 - (1) Function collation
 - (2) Core line extraction
 - For addressing internal/external modifications of OSS

Notations



Working example: Heap-buffer overflow vulnerability (CVE-2016-8654) in Jasper

```
void jpc_qmfb_split_col (...) {
if (bufsize > QMFB SPLITBUFSIZE) {
   if (!(buf = jas alloc2(bufsize, sizeof(jpc fix t)))) {
      abort();
if (numrows >= 2) {
  - hstartcol = (numrows + 1 - parity) > 1;
  - // ORIGINAL (WRONG): m = (parity) ? hstartcol : (numrows - hstartcol);
  - m = numrows - hstartcol;
 + hstartrow = (numrows + 1 - parity) > 1;
  + // ORIGINAL (WRONG): m = (parity) ? hstartrow : (numrows - hstartrow);
 + m = numrows - hstartrow;
   n = m;
   dstptr = buf;
   srcptr = &a[(1 - parity) * stride]
```

A patch snippet for CVE-2016-8654 in Jasper

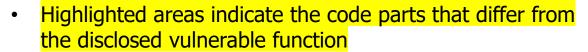
Working example

Function collation



Oldest vulnerable function

```
1 void jpc qmfb split col (...) {
2
  . . .
   if (bufsize > QMFB SPLITBUFSIZE) {
      if (!(buf = jas alloc(bufsize * sizeof(jpc fix t)))) {
         abort();
6
   if (numrows >= 2) {
      hstartcol = (numrows + 1 - parity) > 1;
10
      m = (parity) ? hstartcol : (numrows - hstartcol);
11
12
      n = m;
13
      dstptr = buf;
      srcptr = &a[(1 - parity) * stride]
14
15
      . . .
```





Disclosed vulnerable function

```
1 void jpc qmfb split col (...) {
   if (bufsize > QMFB SPLITBUFSIZE) {
      if (!(buf = jas alloc2(bufsize, sizeof(jpc fix t)))) {
         abort();
6
   if (numrows >= 2) {
    - hstartcol = (numrows + 1 - parity) > 1;
10
    - // ORIGINAL (WRONG): m = (parity) ?
                              hstartcol : (numrows - hstartcol);
    - m = numrows - hstartcol;
11
12
      n = m;
13
      dstptr = buf;
      srcptr = &a[(1 - parity) * stride]
14
15
      . . .
```

Definition of core lines

Core lines in vulnerability signature generation

1. Essential code lines

Code lines that were deleted from the patch and exist in both fo and fd

2. Dependent code lines

❖ Code lines that have control/data dependencies with the essential code lines

3. Control flow code lines

Control statements that exist in both fo and fd

Working example

1) Extracting essential code lines (Ev)

- Code lines that were deleted from the patch (existing in both fo and fd)
- Essential code lines are closely related to the vulnerability manifestation

```
1 void jpc qmfb split col (...) {
2
   if (bufsize > QMFB SPLITBUFSIZE) {
      if (!(buf = jas alloc(bufsize * sizeof(jpc fix t)))) {
         abort();
6
  if (numrows >= 2) {
      hstartcol = (numrows + 1 - parity) > 1;
10
      m = (parity) ? hstartcol : (numrows - hstartcol);
11
12
      n = m;
13
      dstptr = buf;
      srcptr = &a[(1 - parity) * stride]
14
15
      . . .
```

```
1 void jpc qmfb split col (...) {
  if (bufsize > QMFB SPLITBUFSIZE) {
      if (!(buf = jas alloc2(bufsize, sizeof(jpc fix t)))) {
         abort();
6
8 if (numrows >= 2) {
    - hstartcol = (numrows + 1 - parity) > 1;
10
    - // ORIGINAL (WRONG): m = (parity) ?
                             hstartcol : (numrows - hstartcol);
    - m = numrows - hstartcol;
12
      n = m;
13
      dstptr = buf;
14
      srcptr = &a[(1 - parity) * stride]
15
```

Working example

2) Extracting dependent code lines (Dv):

- Code lines that have control/data dependency with the essential code lines
- To determine whether the vulnerability trigger environment has propagated

```
1 void jpc qmfb split col (...) {
2
   if (bufsize > QMFB SPLITBUFSIZE) {
      if (!(buf = jas alloc(bufsize * sizeof(jpc fix t)))) {
         abort();
6
   if (numrows >= 2) {
      hstartcol = (numrows + 1 - parity) > 1;
10
      m = (parity) ? hstartcol : (numrows - hstartcol);
11
12
      n = m;
13
      dstptr = buf;
      srcptr = &a[(1 - parity) * stride]
15
      . . .
```

```
1 void jpc qmfb split col (...) {
   if (bufsize > QMFB SPLITBUFSIZE) {
      if (!(buf = jas alloc2(bufsize, sizeof(jpc fix t)))) {
         abort();
   if (numrows >= 2) {
      hstartcol = (numrows + 1 - parity) > 1;
10
      // ORIGINAL (WRONG): m = (parity) ?
                             hstartcol : (numrows - hstartcol);
11
      m = numrows - hstartcol;
12
      n = m;
13
      dstptr = buf;
14
      srcptr = &a[(1 - parity) * stride]
15
```

Working example

- 3) Extracting control flow code lines (Fv)
 - To determine whether the essential code line has still reachable with the same conditions.

```
1 void jpc qmfb split col (...) {
2
  . . .
   if (bufsize > QMFB SPLITBUFSIZE) {
      if (!(buf = jas alloc(bufsize * sizeof(jpc fix t)))) {
         abort();
6
   if (numrows >= 2) {
      hstartcol = (numrows + 1 - parity) > 1;
10
      m = (parity) ? hstartcol : (numrows - hstartcol);
11
12
      n = m;
13
      dstptr = buf;
      srcptr = &a[(1 - parity) * stride]
15
      . . .
```

```
1 void jpc qmfb split col (...) {
  if (bufsize > QMFB SPLITBUFSIZE) {
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         abort();
   if (numrows >= 2) {
      hstartcol = (numrows + 1 - parity) > 1;
10
      // ORIGINAL (WRONG): m = (parity) ?
                             hstartcol : (numrows - hstartcol);
11
      m = numrows - hstartcol;
12
      n = m;
13
      dstptr = buf;
14
      srcptr = &a[(1 - parity) * stride]
15
```

Gathering code lines and generating signatures

4) Generating signatures

- Vulnerability signature (Sv)
- Patch signature (Sp)
 - ❖ An approach similar to generating a vulnerability signature is performed (deleted -> added)
 - ❖ Control flow lines (Fp) that exist only in the patch function are already included in Ep

$$S_{\nu}=(E_{\nu},\,D_{\nu},\,F_{\nu})$$

$$S_p = (E_p, D_p)$$

Phase (2) Vulnerable code clone discovery

Detecting vulnerable code clones in the target program (T)

A function f in T is a vulnerable code clone if it satisfies:

• Cond 1) f should contain all code lines in S_v .

$$\forall_{l \in S_{v}}. (l \in f)$$

• Cond 2) f should not contain any code lines in S_p .

$$\forall_{l \in S_p} . (l \notin f)$$

• Cond 3) The syntax of f should be similar to f_o or f_d .

$$(\mathtt{Sim}(f,f_o) \geq \mathbf{\theta}) \vee (\mathtt{Sim}(f,f_d) \geq \mathbf{\theta})$$

Phase (2) Vulnerable code clone discovery

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$$\forall_{l \in S_p} . (l \notin f)$$

• Cond 3) The syntax of f should be similar to f_o or f_d .

$$(\mathtt{Sim}(f,f_o) \geq \theta) \vee (\mathtt{Sim}(f,f_d) \geq \theta)$$



Dataset and parameter setting

CVE dataset

- 4,219 C/C++ CVE vulnerabilities (patches)
 - ❖ Collected from NVD
 - ❖ 7,762 vulnerable/patched function pairs
 - ❖ 5,936 oldest vulnerable functions

Target programs

10 software programs that are popular (based on GitHub stars) and contain a sufficient number of OSS components

Parameter

Target program overview

IDX	Name	Version	#Line*	#Comp [†]	Domain
T1	FreeBSD	v12.2.0	14,489,534	47	Operating system
T2	ReactOS	v0.4.13	6,419,855	23	Operating system
T3	ArangoDB	v3.7.9	3,064,973	22	Database
T4	FFmpeg	n4.3.2	1,230,520	4	Multimedia processing
T5	OpenCV	v4.5.1	1,092,317	15	Computer vision
T6	Emscripten	v2.0.15	759,020	11	Compiler
T7	Crown	v0.42.0	723,372	20	Game engine
T8	Git	v2.31.0	293,467	5	Version control system
T9	OpenMVG	v1.6	262,610	8	Image processing
T10	Redis	v5.0.12	212,672	8	Database
Total	-	-	28,548,340	190	-

^{*:} Counting only C/C++ code lines, †: The number of modified OSS components.

[•] $\theta = 0.5$

Accuracy measurement

- Comparison targets
 - Two existing vulnerable code clone detection tools: <u>VUDDY</u> [S&P 2017] and <u>ReDeBug</u> [S&P 2012]
 - MOVERY significantly outperformed existing approaches

Target software	#Discovered VCCs*	ReDeBug						VUDDY					Movery				
Target software		#TP	#FP	#FN	Precision	Recall	#TP	#FP	#FN	Precision	Recall	#TP	#FP	#FN	Precision	Recall	
ReactOS	210	31	9	179	0.78	0.15	8	0	202	1.00	0.04	207	3	3	0.99	0.99	
OpenCV	72	38	15	34	0.72	0.53	26	2	46	0.93	0.36	72	3	0	0.96	1.00	
Emscripten	56	22	8	34	0.73	0.39	9	1	47	0.90	0.16	50	4	6	0.93	0.89	
FreeBSD	33	25	44	8	0.36	0.76	6	16	27	0.27	0.18	27	4	6	0.87	0.82	
Crown	23	22	2	1	0.92	0.96	14	2	9	0.88	0.61	23	2	0	0.92	1.00	
OpenMVG	23	15	5	8	0.75	0.65	4	0	19	1.00	0.17	19	0	4	1.00	0.83	
ArangoDB	6	4	1	2	0.80	0.67	2	0	4	1.00	0.33	6	2	0	0.75	1.00	
FFmpeg	5	2	2	3	0.50	0.40	0	1	5	0.00	0.00	5	1	0	0.83	1.00	
Redis	5	3	0	2	1.00	0.60	3	0	2	1.00	0.60	5	0	0	1.00	1.00	
Git	1	1	1	0	0.50	1.00	0	0	1	N/A	0.00	1	0	0	1.00	1.00	
Total	434	163	87	271	0.65	0.38	72	22	362	0.77	0.17	415	19	19	0.96	0.96	

*VCCs: Vulnerable Code Clones

Accuracy measurement

- Comparison targets
 - Two existing vulnerable code clone detection tools: <u>VUDDY</u> [S&P 2017] and <u>ReDeBug</u> [S&P 2012]
 - MOVERY significantly outperformed existing approaches

Target software	#Discovered VCCs*	ReDeBug						VUDDY					MOVERY				
Target software		#TP	#FP	#FN	Precision	Recall	#TP	#FP	#FN	Precision	Recall	#TP	#FP	#FN	Precision	Recall	
ReactOS	210	31	9	179	0.78	0.15	8	0	202	1.00	0.04	207	3	3	0.99	0.99	
OpenCV	72	38	15	34	0.72	0.53	26	2	46	0.93	0.36	72	3	0	0.96	1.00	
Emscripten	56	22	8	34	0.73	0.39	9	1	47	0.90	0.16	50	4	6	0.93	0.89	
FreeBSD	33	25	44	8	0.36	0.76	6	16	27	0.27	0.18	27	4	6	0.87	0.82	
Crown	23	22	2	1	0.92	0.96	14	2	9	0.88	0.61	23	2	0	0.92	1.00	
OpenMVG	23	15	5	8	0.75	0.65	4	0	19	1.00	0.17	19	0	4	1.00	0.83	
ArangoDB	6	4	1	2	0.80	0.67	2	0	4	1.00	0.33	6	2	0	0.75	1.00	
FFmpeg	5	2	2	3	0.50	0.40	0	1	5	0.00	0.00	5	1	0	0.83	1.00	
Redis	5	3	0	2	1.00	0.60	3	0	2	1.00	0.60	5	0	0	1.00	1.00	
Git	1	1	1	0	0.50	1.00	0	0	1	N/A	0.00	1	0	0	1.00	1.00	
Total	434	163	87	271	0.65	0.38	72	22	362	0.77	0.17	415	19	19	0.96	0.96	

MOVERY could discover 2.5x and 5.8x more vulnerable codes than ReDeBug and VUDDY

Accuracy measurement

MOVERY could discover more VCCs than VUDDY and ReDeBug

VCCs that are hardly discovered by existing techniques

Types	Description	#Discovered VCCs
T1	VCCs without code lines deleted in security patches.	32
T2	VCCs with various syntaxes derived from the oldest vulnerable function (fo).	221 (221 VCCs closer to fo than fd)
Т3	VCCs with heavy syntax change.	166 (166 VCCS: $Sim(f, f_d) \le 0.5$)

Conclusion

Conclusion

- Many vulnerable codes are propagated with <u>syntax modifications</u>
 - 396 (91%) out of 434 VCCs existed in a different syntax to the disclosed vulnerable function
- MOVERY
 - A precise approach for discovering modified VCCs from modified components
 - MOVERY significantly outperformed existing approaches in vulnerable code clone discovery
 - ❖ High vulnerability discovery accuracy: 96% precision and 96% recall
- Equipped with VCC discovery results from MOVERY,
 - Developers can address threats caused by propagated vulnerabilities in modified components

Thank you for your attention!

MOVERY repository (https://github.com/wooseunghoon/MOVERY-public)

CONTACT

- Seunghoon Woo (<u>seunghoonwoo@korea.ac.kr</u>, <u>https://wooseunghoon.github.io</u>)
- Computer & Communication Security Lab (https://ccs.korea.ac.kr)
- Center for Software Security and Assurance (https://cssa.korea.ac.kr)

APPENDIX I

Preprocessing

- Abstraction
 - Replacing every occurrence of parameters, variable names, variable types, and callee function names in each function with symbols PARAM, DNAME, DTYPE, and FCALL

```
3 if (bufsize > QMFB_SPLITBUFSIZE) {
8 if (numrows >= 2) {
9  hstartcol = (numrows + 1 - parity) > 1;
14 srcptr = &a[(1 - parity) * stride]

3 if (DVAL > QMFB_SPLITBUFSIZE) {
8 if (PARAM >= 2) {
9  DVAL = (PARAM + 1 - PARAM) > 1;
14 DVAL = &PARAM[(1 - PARAM) * PARAM]
```

- Selective abstraction
 - Abstraction is applied only when the abstraction code before and after the patch is <u>different</u>

APPENDIX II

Speed and scalability measurement

- MOVERY requires the least amount of time in the vulnerability discovery
- MOVERY discovers VCCs from the target programs varied from 213 K to 14.5 M LoC
 - The required time is not significantly increased

