An FPT algorithm for counting subgraphs

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December 30, 2018

What is the subgraph counting problem?

Problem Statement

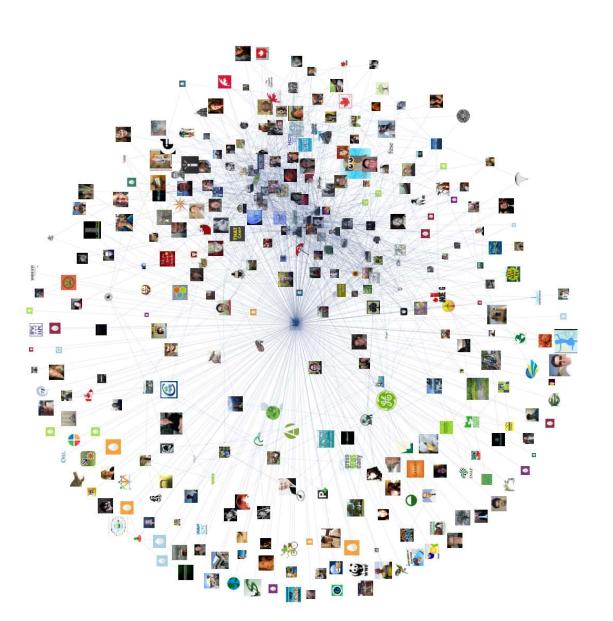
How many (unlabelled) copies of the graph H are contained in the graph ĊЭ

We call the graph G the host graph and H the pattern graph.

Why do we care?

subgraph isomorphism Generalisation of

Network analysis



Algorithmic complexity

NP-hard

NP-hard problems cannot be solved in polynomial time unless P=NP, which we do not expect is true.

Fixed-parameter tractable (FPT)

A problem is FPT if it can be solved in time depending polynomially on the input size and exponentially on some other parameter.

Subgraph counting is hard!

ullet Subgraph isomorphism is NP-hard o subgraph counting is NP-hard

 Assuming Exponential Time Hypothesis → subgraph counting is not in FPT in general

• (Enright and Meeks) Subgraph counting is in FPT when the host graph has almost bounded degree

My PhD Mini-Project

Bounded degree graph

A graph G has bounded degree k if each vertex in G has degree at most k.

Almost bounded degree graph

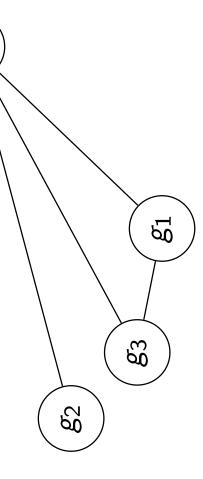
A graph G has almost bounded degree k if G contains at most k vertices with degree greater than *k*.

Project objective:

Design and implement an FPT algorithm for subgraph counting in host graphs with almost bounded degree.

Algorithm: General Idea

- Consider each way to assign part of H to high degree vertices in G
- For each feasible assignment, count ways to assign the rest of H to the bounded degree part of G
- ullet Sum up the counts to obtain number of labelled copies of H in G
- Divide by the number of copies of H in H to obtain number of unlabelled copies of H in G



h₃

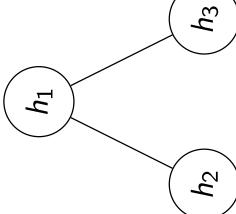
 h_2

 $h_1\colon g_1,g_2,g_3,g_4$

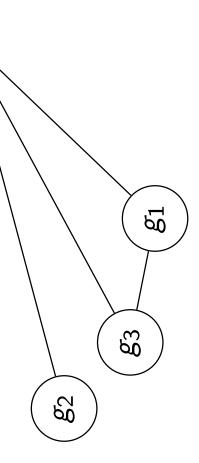
*h*2: *g*1, *g*2, *g*3, *g*4 *h*3: *g*1, *g*2, *g*3, *g*4

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 $h_1\colon g_1,g_2,g_3,g_4$

*h*2: *g*1, *g*2, *g*3, *g*4 *h*3: *g*1, *g*2, *g*3, *g*4

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 h_2

Subset of V(H): \emptyset Count = 0

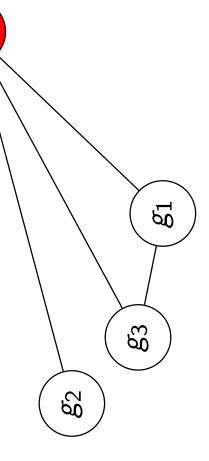
 $h_1: g_1, g_2, g_3, g_4$

 h_2 : g_1, g_2, g_3, g_4

h3: g1, g2, g3, g4

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*g*4

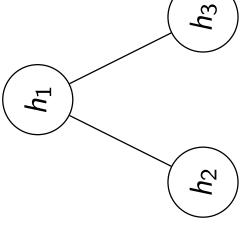


Subset of V(H): \emptyset Count = 0

 $h_1 o g_1$ $h_2 : g_1, g_2, g_3$ $h_3 : g_1, g_2, g_3$

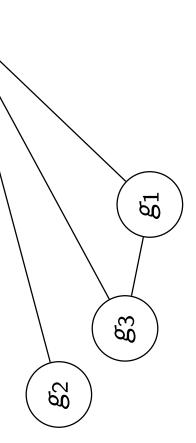
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 g_4



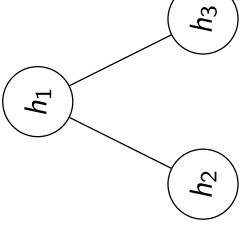
Subset of V(H): \emptyset Count = 0

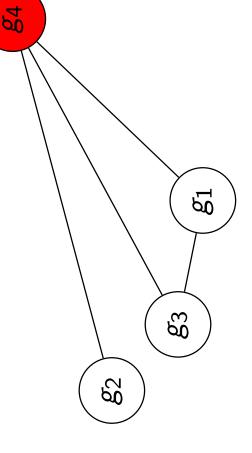
 $h_1 o g_1$

 $h_2 o g_3$

h3: 83

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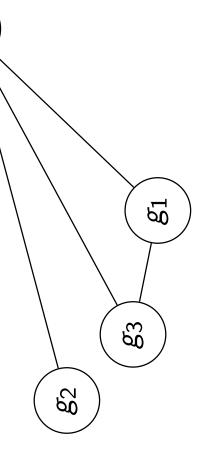


Subset of V(H): \emptyset Count = 0

 $h_1 o g_2$ $h_2 : g_1, g_2, g_3$ $h_3 : g_1, g_2, g_3$

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*g*4



Subset of V(H): \emptyset Count = 0

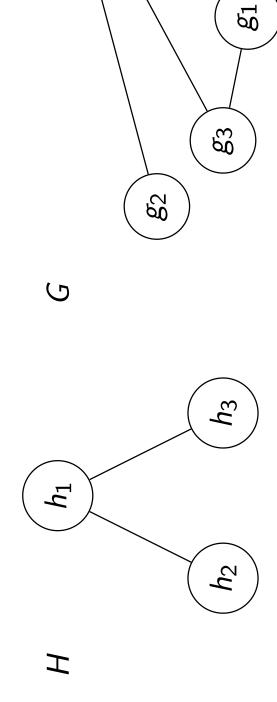
 $h_1 o g_3$ $h_2 : g_1, g_2, g_3$

h3: g1, g2, g3

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Subset of V(H): \emptyset $\mathsf{Count} = 0$

 $h_1 o g_3$ $h_2 o g_1$ 48

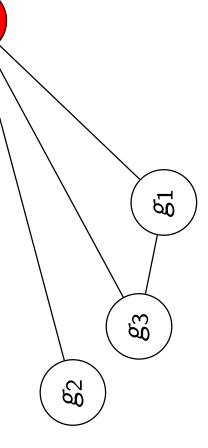
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 h_1

*g*4



h3

 h_2

Subset of V(H): h_1 Count = 0

h₁: g₁, g₂, g₃, g₄ h₂: g₁, g₂, g₃, g₄

h3: g1, g2, g3, g4

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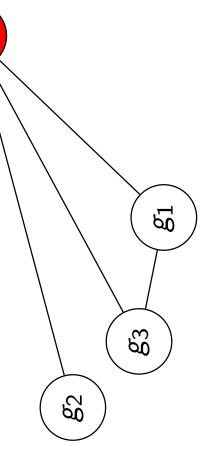
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 h_2

*g*4



Subset of V(H): h_1 Count = 0

 $h_1 o g_4$

 h_2 : g_1, g_2, g_3

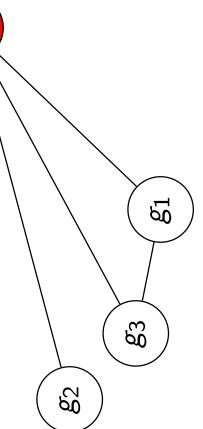
 $h_3:g_1,g_2,g_3$

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Connected components of $H \setminus h_1$:

$$C_1 = h_2$$

$$C_1 = h_2$$
$$C_2 = h_3$$

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copies of C_1 and C_2 in $G\setminus g_4=$ (copies of C_1 in $G\setminus g_4$

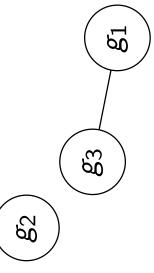
 \times copies of C_2 in $G \setminus g_4$)

overlapping copies of C_1 and C_2 in $G \setminus g_4$



$$G \setminus g_4$$





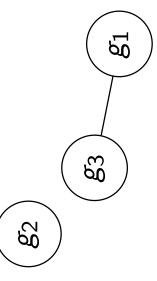
Count =0

 $h_2: g_1, g_2, g_3$

 \mathcal{C}_1

 $G \setminus g_4$





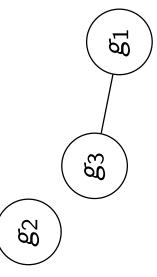
Count =0+1

 $h_2 o g_1$

$$\mathcal{C}_1$$

$$G \setminus g_4$$





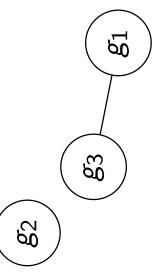
Count =1+1

$$h_2 \rightarrow g_2$$

$$\mathcal{C}_1$$

$$G \setminus g_4$$





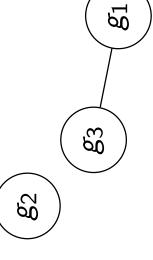
Count
$$=2+1=3$$

$$h_2 \rightarrow g_3$$

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 C_2

 $G \setminus g_4$

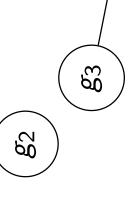


Count =0

 $h_3:g_1,g_2,g_3$

 C_2

$$G \setminus g_4$$



 g_1

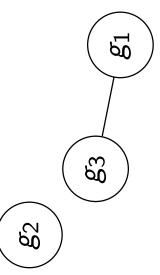
Count =0+1

$$h_3 o g_1$$

 \mathcal{C}^{2}

$$G \setminus g_4$$





Count =1+1

$$h_3 o g_2$$

$$G \setminus g_4$$



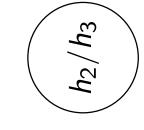
$$\begin{pmatrix} g_2 \\ g_3 \end{pmatrix}$$

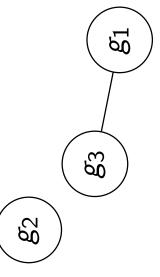
Count =2+1=3

$$h_3 o g_3$$

$$C_1 \cap C_2$$

$$G \setminus g_4$$



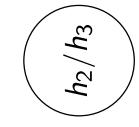


Count =0

 h_2/h_3 : g_1, g_2, g_3

$$C_1\cap C_2$$

$$G \setminus g_4$$



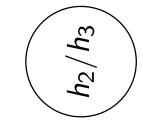
$$\begin{pmatrix} g_2 \\ g_3 \end{pmatrix}$$

Count =0+1

$$h_2/h_3 \rightarrow g_1$$

$$C_1 \cap C_2$$

$$G \setminus \mathcal{B}_4$$



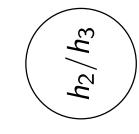
$$\begin{pmatrix} g_2 \\ g_3 \end{pmatrix}$$

Count =1+1

$$h_2/h_3 \rightarrow g_2$$

$$C_1 \cap C_2$$

$$G \setminus g_4$$



$$\begin{pmatrix} g_2 \\ g_3 \end{pmatrix}$$

Count =2+1=3

$$h_2/h_3 \rightarrow g_3$$

copies of C_1 and C_2 in $G\setminus g_4=$ (copies of C_1 in $G\setminus g_4$

 \times copies of C_1 in $G \setminus g_4$)

overlapping copies of C_1 and C_2 in $G \setminus g_4$

copies of
$$C_1$$
 and C_2 in $G\setminus g_4=(3 imes3)-3$

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 h_2

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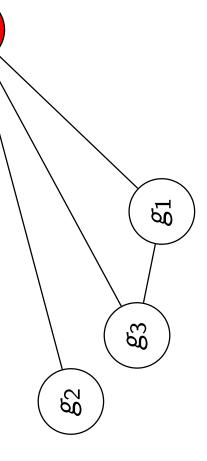
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h3

*g*4



Subset of V(H): h_2 Count = 6

 $h_1: g_1, g_2, g_3, g_4$

 h_2 : g_1, g_2, g_3, g_4

h3: g1, g2, g3, g4

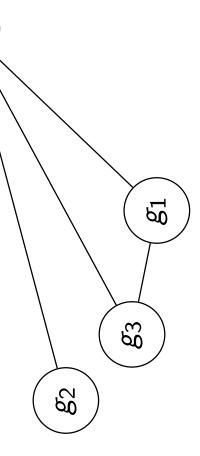
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h3

 h_2

9

*g*4



Subset of V(H): h_2 Count = 6

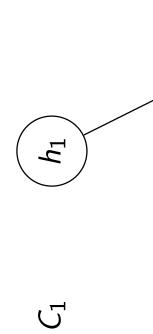
 $h_1\colon g_1,g_2,g_3$

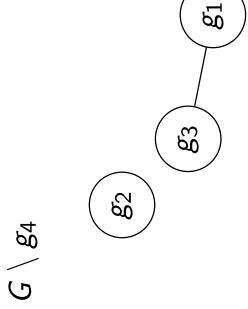
 $h_2 o g_4$

 $h_3:g_1,g_2,g_3$

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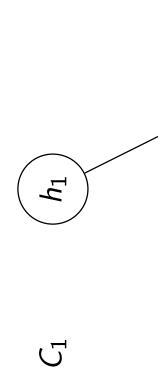
Connected components of $H \setminus h_2$: $C_1 = h_1, h_3$

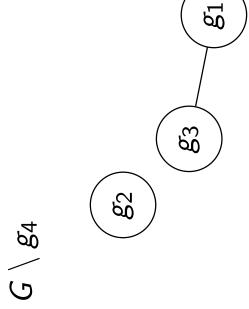




Count =0

*h*₁: *g*₁, *g*₂, *g*₃ *h*₃: *g*₁, *g*₂, *g*₃

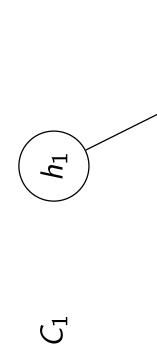


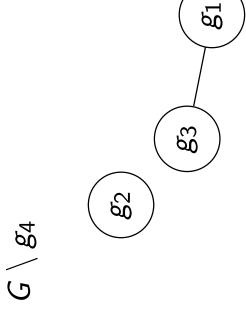


Count =0

 $h_1 o g_1$

 $h_3: g_1, g_2, g_3$



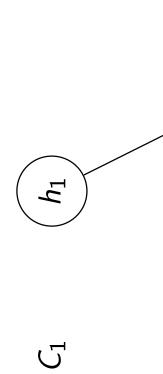


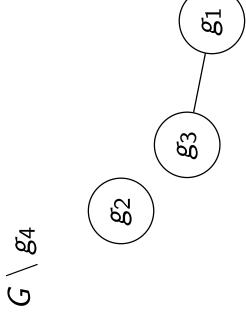
Count =0+1

$$h_1 \rightarrow g_1$$
 $h_3 \rightarrow g_3$

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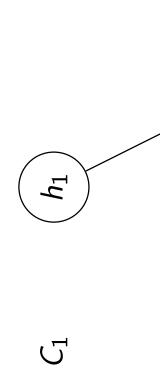


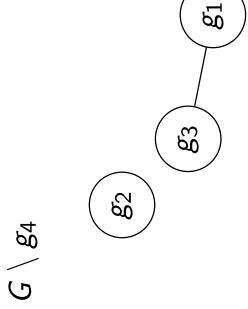


Count =1

 $h_1 o g_2$

h3: g1, g2, g3

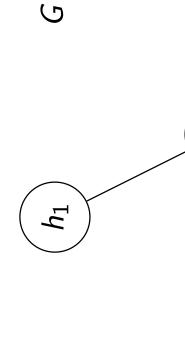




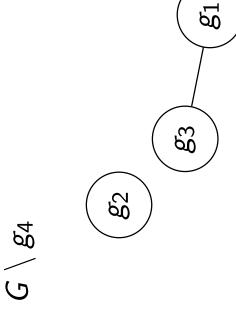
Count =1

 $h_1 o g_3$

 $h_3: g_1, g_2, g_3$



 C_1

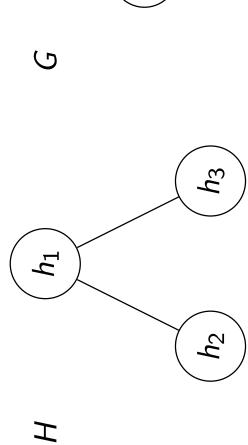


Count =1+1

$$h_1 \rightarrow g_3$$

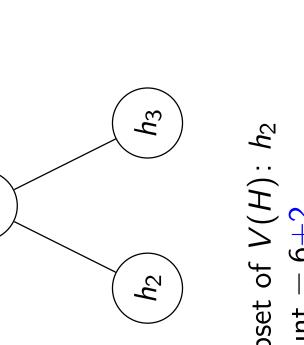
 $h_3 \rightarrow g_1$

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 g_1

 g_3

Subset of V(H): h_2 Count = 6+2

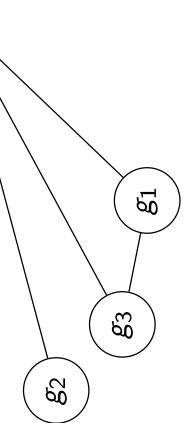
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Subset of V(H): h_3 Count = 8+2

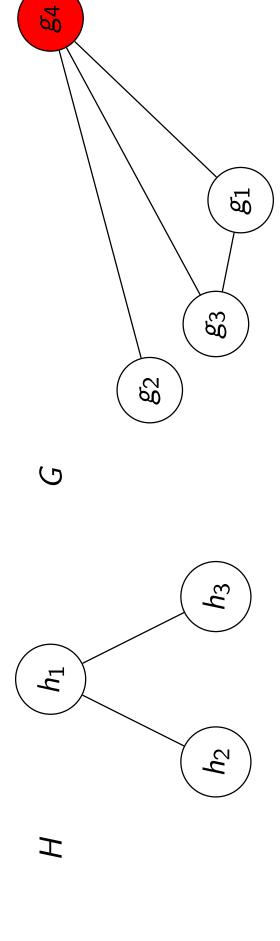
 g_4



Count = 10

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Count = 10

Number of copies of H in H=2

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 $\mathsf{Count} = 10$

Number of copies of H in H=2

ightarrow Number of unlabelled copies of H in G=10/2=5