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Class: CS3100

Assignment: Project 6

Professor: James Anderson

Due Date: 12/5/25

1.A - Introduction

My design will model the scene dataset for a given loading zone. The SceneTable datastructure will be based off of HashTable<>

1.B - Background

For some background, I am in the process of a Ghidra reverse engineering project for Ys VI: The Ark of Napishtim on Windows, and I got the idea from my work on that game.

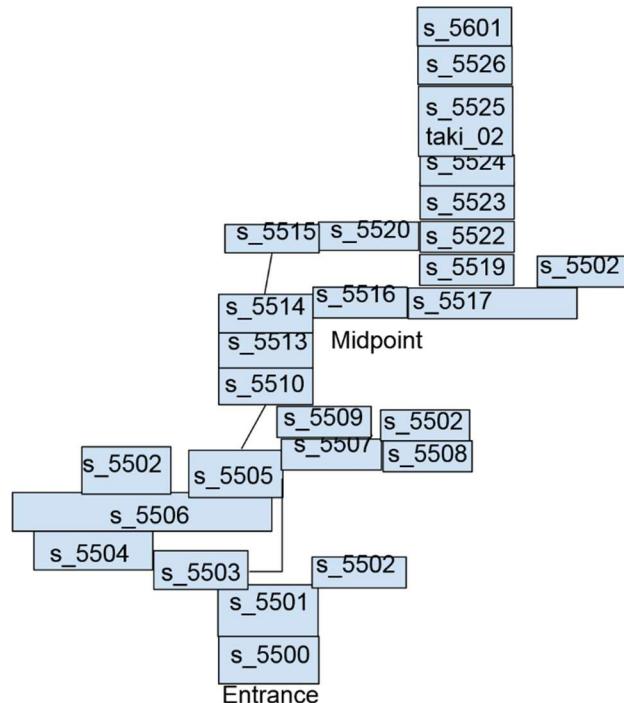
YS VI uses a system where game actors (characters, the player, enemies, etc.) have unique IDs that are used to fetch their dialogue, sprites, and other metadata. However, outside of the PS2 remake, the game lacks voice acting. As such, I have been trying to mod in voice acting to existing dialogue functions. In the original system, a pointer to the dialogue is found using the actor's ID and a scene identifier, and the lines are fed to the display in a single delimited string. Before going into function of the data structure:

1.C - Review of Terms

- Scene

Each level in Ys Vi is split into loading zones or **"scenes"** that have a unique ID (i.e. s_5502) that helps the game find the data for the scene (models, actors, etc.).

Ys Vi File Map:
Limewater Cave
By Grant Rynders
<https://github.com/GrantBenR/YsViDecomp>



- Actor

The term the game uses for its characters, whether the player, NPCs, or enemies. They have associated string IDs just like the scenes.

- Player (Adol)



- Moves around and talks (best of both worlds)

- Enemy



- Moves around but doesn't talk

- NPC



- Talks but doesn't move around

1.D - Design goal

The SceneTable structure should be able to efficiently fetch actor related data relevant to a scene when loading into that scene.

2 - Design Philosophy

The design of a system like this should be as simple and as lightweight as possible, but something that is flexible for the needs of a given game scene. The hashtable will take in a raw dataset from the game data when a scene is loaded in. This process should be streamlined for fast load times even in large or busy scenes with lots of actors.

3 - Core Operations

> As a hashtable derivative, SceneTable uses the primary CRUD actions of

- + insert(id): bool // O(1)-O(N)

> At the initialization of a scene, all the necessary actors are inserted.

- + update(id): bool // O(1)-O(N)

- + delete(id): bool // O(1)-O(N)

- + get(id): bucket* // O(1)

> The scene should be able to fetch the pointers to the actor's data

-- resize_table(): bool // O(N)

> As with other hashtables, the table size increases capacity when enough are added.

4 - Set Operations

- SceneTable(scene_data*): int

> Initialize the table on loading into scene from raw data.

> Returns success value

```
int SceneTable::SceneTable(SceneData* scene_data)
{
    scene_data ...;
    //
    // Extract actors from data
    //
    for (Actor* actor in actors)
    {
        // insert into table
        this->insert(actor->actor_id);
    }
    return 0;
}
```

5 - Extension Feature

- GetDialogue(actor_id): (audio*, text*)

> Returns pointers to raw audio and text values in the game data that are fed to the game front end functions. The text and audio have delimiters so that the display and audio functions know how many lines are in the dialogue and where the raw data begins and ends.

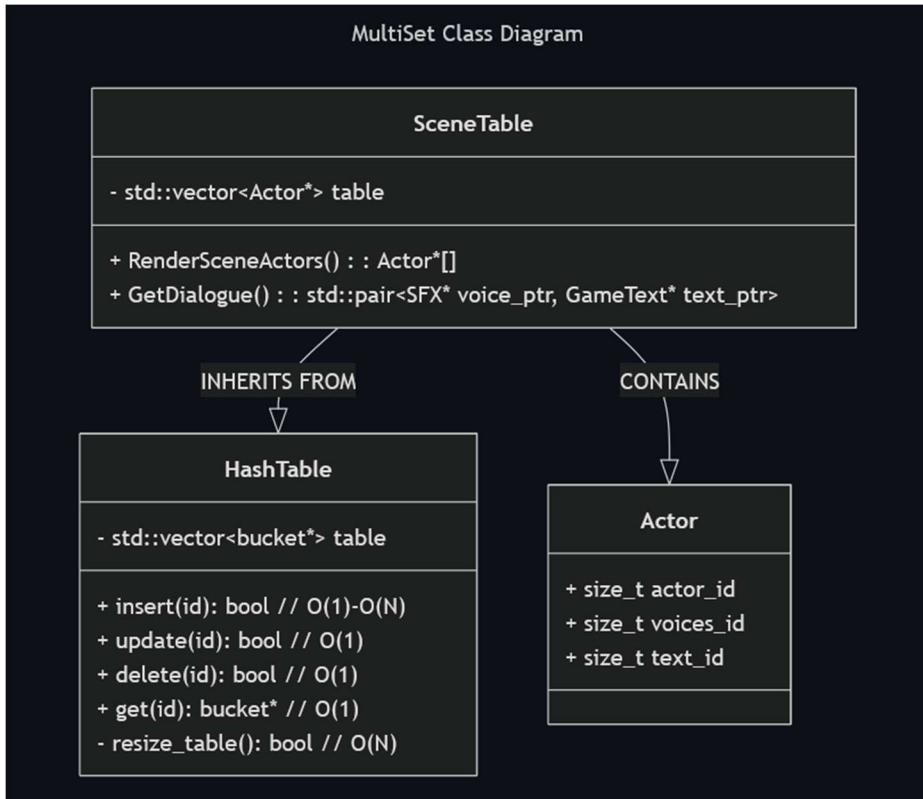
```
std::pair<SFX* voice_ptr, GameText* text_ptr> SceneTable::GetDialogue(size_t actor_id)
{
    try
    {
        Actor* actor = this->get(actor_id);
        SFX* voice = actor->get_voice(this->scene_id);
        GameText* game_text = actor->get_dialogue_text(this->scene_id);
        std::pair<SFX* voice_ptr, GameText* text_ptr> dialogue_pair(voice,game_text);
        return dialogue_pair;
    }
    catch (...)
    {
        ...
    }
}
```

- RenderSceneActors(): Actor*[]
- For each actor in the scene, load sprites and models from the pointers in the hashtable into memory and display.
- Returns pointers to actors' data in memory

```
Actor*[] SceneTable::RenderSceneActors()
{
    for (Actor* actor : this->table)
    {
        GameClass::Render(actor);
    }
    return this->table;
}
```

6 - UML Diagram / Abstraction Boundary

I have used mermaid class diagrams in [prior classes](#), so that is what I am using here



7.0 - Trade Off Analysis

Below are some comparisons to other possible data structures:

7.1 - Sequence Trade Offs

The main alternative for this multiset's underlying data-structure would be a Sequence (`std::vector`). A HashTable is already simply a more complex wrapper for a `std::vector`, so realistically, a Sequence would be more lightweight to implement.

However, assuming an actor has the same ID across scenes, then when making the vector you wouldn't have a clean [1,2,3,4,5...] of ID values (it would be something like [t_200,t_700,t_567,t_890...] which is not cleanly searchable).

As such, either you would hash the values (which is just a HashTable O(1) search) or insert them in the order of appearance making the search complexity O(N).

As such, given that the IDs are unique, a HashTable is the better option, especially in large scenes with lots of actors [1].

7.2 - AVLTree Trade Offs

HashTable has a worst case insertion complexity of O(N) and best case O(1) complexity. AVLTree consistently has complexity O(logN) for its insertion.

In this context, where the number of scene actors is not significant, these differences don't matter much. For getting a value though, HashTable has a typical complexity of O(1) whereas AVLTree consistently has a higher O(logN) complexity.

As such, for a series of values that is primarily only read from, HashTable is the better option

8 - Alternative Design Sketch

Alternative Designs Compared		
HashTable	Sequence	AVLTree
- std::vector<bucket*> table	+ SequenceNode* root + SequenceNode* tail + size_t size	+ AVLNode* root + size_t size
+ insert(id): bool // O(1)-O(N) + update(id): bool // O(1) + delete(id): bool // O(1) + get(id): bucket* // O(1)	+ insert(id): bool // O(N) + update(id): bool // O(N) + delete(id): bool // O(N) + get(id): bool // O(N)	+ insert(id): bool // O(logN) + update(id): bool // O(logN) + delete(id): bool // O(logN) + get(id): bool // O(logN)

9 - Evaluation Plan

Some basic tests to perform would be to compare load times between structures when entering a new scene. If performing get actions on the table is slowing down game-speed a solution with less functionality may need to be explored. How many times the HashTable needs to be resized would be the costliest operation performance-wise, so putting breakpoints on that happening is another good way to gauge performance. The underlying data structure of the SceneTable should not necessarily affect the outside program: the program should just be able to ask for data and get it.

10 - Conclusion/Reflection

This design for a scene data table works because once it is set up, most actions performed on it have $O(1)$ time complexity resulting in low performance overhead. While a Sequence or AVLTree are more performant to initialize at scale, because the other actions performed on the table significantly exceed that of the initialization, it performs better overall.

The SceneTable implementation has a low level of abstraction with the core functions of searching and indexing being made as fast and as simple as possible. It is a simple idea, and a simple execution that results in a low performance overhead. The encapsulation is very good since the load times will be non-substantial, and the end user will not realize the

extent of the behind-the-scenes work being done [2]. Other systems in the game will be able to easily access data for the actor_id resulting in a composition where the program does not have to know much of anything about how SceneTable even works. New methods can easily be implemented that take advantage of its simple HashTable derived CRUD operations.

Works Cited

Karimov, Elshad. *Data Structures and Algorithms in Swift*. Apress, Berkeley, CA. pp55-60. 2020. https://doi.org/10.1007/978-1-4842-5769-2_7.

[2] Cahill, Vinny. *Learning to Program the Object-Oriented Way with C#*. Springer-Verlag London. pp221-249. 2002. https://link.springer.com/chapter/10.1007/978-1-4471-0115-4_7.