# Collecting a Swarm in a 2D Environment Using Shared, Global Inputs

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Abstract—Consider a swarm of robots initialized in a 2D planar grid world where each position is either free-space or obstacle. Assume all robots respond equally to the same global input, such as an applied fluidic flow or electric field. This paper investigates what input sequence should be applied to collect all the robots to one position.

#### I. INTRODUCTION

This paper builds on the techniques for controlling many simple robots with uniform control inputs presented in [1]-[3], and outlines new research problems.

a) Definitions: The 'robots' in this paper are simple particles without autonomy. A planar grid workspace W is filled with a number of unit-square robots (each occupying one cell of the grid) and some fixed unit-square blocks. Each unit square in the workspace is either free, which a robot may occupy or obstacle which a robot may not occupy. Each square in the grid can be referenced by its Cartesian coordinates x = (x, y). All robots are commanded in unison: the valid commands are "Go Up" (u), "Go Right" (r), "Go Down" (d), or "Go Left" (l).

We consider two classes of commands, discrete and maximal moves. Discrete moves: robots all move in the commanded direction one unit unless they are prevented from moving by an obstacle or a stationary robot. Maximal moves: robots all move in the commanded direction until they hit an obstacle or a stationary robot. For maximal moves, we assume the area of W is finite and issue each command long enough for the robots to reach their maximum extent. A command sequence m consists of an ordered sequence of moves  $m_k$ , where each  $m_k \in \{u, d, r, l\}$  A representative command sequence is  $\langle u, r, d, l, d, r, u, \ldots \rangle$ .

We consider two notions of a collected swarm: if robots are allowed to overlap, the swarm is connected when all robots share the same (x, y) coordinates. If the robots may not overlap, the swarm is collected when it forms one 4connected component.

Research questions include

- 1) What connected topologies allow collecting a swarm?
- 2) What connected topologies do not allow collecting a swarm?
- 3) What connected topologies can be collected with a given command sequence?

#### II. RELATED WORK

#### A. Localization

The strongest parallel in the literature is the field of *Almost* sensor-less localization or "localizing a blind robot in a known map", where a mobile robot using a compass, a bumpsensor that detects when the robot contacts a wall, and a map of the workspace must localize itself in the map [4], [5]. This has been recently extended to robots with bounded uncertainty in their inputs [6]. The basic methodology is to construct motion plans that: "First, actions are selected which reduce the uncertainty in the robots position to a finite set of possibilities. Second, actions are selected to reduce the uncertainty from this finite set to a single point. "[4].

"The problem of finding a localizing sequence for a given environment can be seen as a planning problem in which the initial state is unknown and the current state is unobservable. To manage this uncertainty, we transform the problem from an unobservable planning problem in state space to an observable problem in a more complex space called the robots information space, which we now define."

"it remains an open problem to generate localizing sequences that are optimal." [4].

This work assumes there is only one robot, and works in a simply connected polygonal environment. We want to relax this to handle unit robots.

Also related is work on sensorless part orientation, where flat tray is tilted in a series of directions to bring a polygonal part, initially placed at random orientation and position in the tray, to a known position and orientation [7]. This may be similar to localizing a robot with minimum travel, but this moves the robot and has it take additional measurements [8]

#### B. Robot Rendezvous

Rendezvous usually concerns two independent, intelligent agents that must meet at a predetermined time.

#### C. Manufacturing

This is similar to work on draining a polygon [9], but with discrete inputs.

Or placing ports for injection-molding

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#### E. Topologies that do not allow collecting a swarm

Robots initialized in two unconnected components i and j of a freespace cannot be collected. The proof is trivial, since robot in freespace i can not reach freespace j.

Under maximal inputs, the freespace can be constructed with spaces resembling bottles or fish weirs from which a single robot cannot escape, as shown in Fig. [TODO]. If the freespace contains at least two such bottles with at least one robot in each, the swarm cannot be collected.

The world must be bounded. Two initially-separated robots in an unbounded world without obstacles cannot be collected, however with discrete inputs, one obstacle is sufficient [10].

#### F. connected topologies allow collecting a swarm

With discrete inputs and robots that are allowed to overlap, the problem can be reduced to localizing a sensorless robot in a known workspace. This is similar to work on draining a polygon [9], but with discrete inputs.

Algorithm 1, brute-force-search: initializes a tree with the root node  $\{C_0,\}$ , where  $C_0$  is the current configuration of possible robot locations, where  $C_0$  is an ordered vector containing the locations of every possible location for a robot in the workspace and  $\{\}$  is the shortest corresponding movement sequence to achieve this configuration. We then construct a breadth-first tree of possible configurations  $\{u,r,d,l\}$ , pruning leaves that already exist in the tree. We stop with the cardinality at a leaf  $C_i=1$ , which indicates that the swarm has been collected (equivalently, the robot has been localized). This algorithm produces the optimal path, but requires O(XX) time to learn and O(YY) memory.

## **Algorithm 1** BlindLocalizationBruteForce( $\overline{W}$ , $\overline{C_0}$

#### Require:

Ensure:  $C_{\boldsymbol{m}} \leftarrow \text{ApplyMoves}[\boldsymbol{m}, C_0, W]$ ,  $|C_{\boldsymbol{m}}| = 1$ 

1:  $T \leftarrow C_0$ ,

2: **return**  $(r_1, r_2)$ 

### Algorithm 2, greedy-approach:

- 1) possible configurations = current config
- 2) initial cardinality = current config
- 3) construct a breadth-first tree of possible configurations  $\{u, r, d, l\}$  prune leaves that already exist.
- 4) stop when the cardinality of possible positions in any leaf is less than initial cardinality, apply this move, and restart the algorithm until cardinality is 1

Given two robots in a bounded  $L \times L$  workspace, the worst case to bring them to the same position is  $< L^2$ . For n robots,  $O(nL^2)$ .

# III. CONNECTED TOPOLOGIES THAT ALLOW COLLECTING A SWARM

Assume that passageways are large compared to the maximum diameter of a swarm. We can accomplish with discrete inputs, we can move a

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