

template

WUST

So Like Coding? You Baldy

Contents

0	Hea			1
	0.1	pbds		. 1
		0.1.1	head	. 1
	0.2	FastIC)	. 1
		0.2.1	FastScanner	. 1
		0.2.2	FastPowAndAdd	. 1
		0.2.3	PythonInput	
		0.2.4	SpecialInput	
	0.3		r	
	0.0		comp	
		0.0.1		
1	Mat	th		5
	1.1	素数		. 5
		1.1.1	Eratosthenes 筛法	
		1.1.2	Euler 筛	
		1.1.3	MillerRabin 素性测试	
		1.1.4	PollardRho 快速因数分解	_
			求素因子	_
	1.0	1.1.5	* ****	
	1.2	. 4,,,,	D 1 DI:	
		1.2.1	EulerPhi	
		1.2.2	Sieve	
		1.2.3	gcd	
		1.2.4	解乘法逆元	
	1.3	1 3/3.		_
		1.3.1	扩展欧几里得算法	
		1.3.2	中国剩余定理	
		1.3.3	扩展中国剩余定理	. 9
		1.3.4	BSGS	. 10
		1.3.5	exBSGS	. 10
		1.3.6	逆元	. 11
		1.3.7	模素数二次同余方程	. 12
	1.4	矩阵与	5线性方程组	
		1.4.1		
		1.4.2	高斯消元....................................	
		1.4.3	线性基	
	1.5		ஜ	
	1.0	1.5.1	Lucas	
		1.5.1 $1.5.2$	exLucas	
		1.5.2 $1.5.3$	extituteds · · · · · · · · · · · · · · · · · · ·	
		1.5.3 $1.5.4$	小模数组合数	
		-		
	1.0	1.5.5	大模数组合数	
	1.6	J , .	TDDD	
		1.6.1	FFT	
		1.6.2	NTT	
		1.6.3	原根	
		1.6.4	FWT	
	1.7	多项式	<u></u>	. 22
		1.7.1	拉格朗日插值	. 22
		1.7.2	拉格朗日插值 (连续取值)	. 22
	1.8	Others	S	. 23
		1.8.1	BM	. 25
		1.8.2	exBM	. 27
		1.8.3	 杜教筛	
			财拉降 复	. 33

2	Gra	ph Th	neory 35
		2.1.1	Dijkstra
		2.1.2	Euler Path
		2.1.3	K shortest Path(Astar)
		2.1.4	K shortest Path(可持久化可并堆)
	2.2	生成林	
	2.2	主.风怀 2.2.1	Kruskal
		2.2.2	Prim
		2.2.3	最小树形图
		2.2.4	Matrix Tree
	2.3	连通性	
		2.3.1	割点
		2.3.2	桥 45
		2.3.3	强连通分量 44
		2.3.4	点双联通分量 45
		2.3.5	边双联通分量 46
	2.4	二分图]匹配
		2.4.1	Hungary Algorithm
		2.4.2	Hopcroft-karp Algorithm
		2.4.3	二分图多重匹配
		2.4.4	二分图最大权匹配 (KM 算法) 50
		2.4.5	一般图匹配带花树
	2.5	网络流	[
		2.5.1	Dinic
		2.5.2	ISAP
		2.5.3	MCMF
		2.5.4	Trick
	2.6		s
	2.0	2.6.1	
		2.6.2	2-SAT
		2.6.3	差分约束系统
		2.6.4	支配树
		2.0.4	又癿例
3	Dat	aStru	ctrue 63
-	3.1		entTreeDS
	0.1	3.1.1	SegmentTree
		3.1.2	离散化区间
		3.1.3	动态区间最大子段和
		3.1.4	动态开点权值线段树
		3.1.5	扫描线
	3.2		1
	ე.∠	3.2.1	HLD
	9 9	-	
	3.3	•	
		3.3.1	RMQ
		3.3.2	RMQbyIndex
	0.4	3.3.3	RMQinNM
	3.4	MO.	76
		3.4.1	MO
		3.4.2	MObyModify
		3.4.3	分块
		3.4.4	弾飞绵羊
		3.4.5	树莫队
	3.5	Virtua	
		3.5.1	VirtualTree
	3.6	Persis	tentDS
		3.6.1	主席树区间 k 大
		3.6.2	可持久化数组
	3.7	Tree	87

template, So Like Coding? You Baldy

		3.7.1 LCA	88
	3.8	3.7.3 点分治	
		3.8.1 BITinNM	
		3.8.2 静态区间 k 大划分树	91
4	Geo	netry	93
	4.1	Class	
		4.1.1 geo	
		4.1.2 3D 凸包	102
5	Stri		108
	5.1	KMP	
		5.1.1 KMP	
	5.2	Trie	
		5.2.1 Trie	109
		5.2.2 Persistence Trie	
	5.3	5.2.3 01Trie	
	0.0	5.3.1 Manacher	
	5.4	Aho-Corasick Automation	112
		5.4.1 AC Automation	
	5.5	Suffix Array	
	5.6	5.5.1 Suffix Array	
	0.0	5.6.1 PalindromicTree	
	5.7	Hash	116
		5.7.1 hash	
		5.7.2 doubleHash	
		5.7.4 树 hash 同构	
	5.8	Suffix Automation	
		5.8.1 SAM	
	5.9	Others	
		5.9.1 最小表示法	123
6	$\mathbf{d}\mathbf{p}$		124
	6.1	BitDP	
		6.1.1 数位 dp 计和 . .	
	6.2	Subsequence	
		6.2.1 MaxSum	
		5.2.2 LIS	_
		6.2.3 LongestCommonIncrease	
	6.3	6.2.4 LCS	
	0.0	5.3.1 矩阵快速幂	
_	0.1		100
7	Oth 7.1	nrs mint 类	130 130
	7.1	************************************	
	7.3	BigInt 类	
	7.4	date	
	7.5	Frac 类	
	$7.6 \\ 7.7$	模拟退火 (最小圆覆盖)	
	7.8	前缀异或和	
	7.9	め 蒙夫 环 第 k 个	139

template, So Like Coding? You Baldy

7.10	:分	39
7.11	i男 IO 挂	10
7.12	[心结论	11
7.13	uiltin	11
7.14	以内 k 因子的个数 14	11

0 Header

0.1 pbds

0.1.1 head

```
1 #include <bits/extc++.h>
2 #pragma comment(linker, "/STACK:102400000,102400000")
3 using namespace __gnu_pbds; // tree, gp_hash_table, trie
4 using namespace __gnu_cxx; // rope
5 tree<TYPE, null_type, less<>, rb_tree_tag, tree_order_statistics_node_update> tr;
6 // 可并堆
7 #include <ext/pb_ds/priority_queue.hpp>
8 using namespace __gnu_pbds;
9 __gnu_pbds::priority_queue<int,greater<int>,pairing_heap_tag> q[maxn];
10 //q[i].join(q[j]) 将j堆并入i
```

0.2 FastIO

0.2.1 FastScanner

```
1 // 适用于正负整数
 2 template <class T>
3 inline bool scan(T &ret){
        char c;
4
5
        int sgn;
        if (c = getchar(), c == EOF) return 0; //EOF
6
7
        while (c != '-' \&\& (c < '0' || c > '9')) c = getchar();
8
        sgn = (c == '-') ? -1 : 1;
9
        ret = (c == '-') ? 0 : (c - '0');
10
        while (c = getchar(), c >= '0' && c <= '9') ret = ret * 10 + (c - '0');
11
        ret *= sgn;
12
        return 1;
13 }
14
15 template <class T>
16 inline void out(T x) {
        if (x > 9) out(x / 10);
17
18
        putchar(x % 10 + '0');
19 }
20
21 inline int read() {
22
        int x = 0;
23
        char ch = getchar();
24
        while (ch > '9' || ch < '0')ch = getchar();</pre>
25
        while (ch >= '0' && ch <= '9') {</pre>
26
            x = x * 10 + ch - '0';
27
            ch = getchar();
28
        }
29
        return x;
30 }
```

0.2.2 FastPowAndAdd

```
1 // 精确快速乘
2 ll qpmul(ll a, ll b) {
3 a %= mod; b %= mod;
```

```
ll res = 0;
4
5
       while (b > 0) {
6
            if (b & 1) {
7
                res = (res + a);
8
                if (res >= mod) res -= mod;
9
            }
10
            a = (a + a);
            if (a >= mod) a -= mod;
11
12
            b >>= 1;
13
       }
14
       return res;
15 }
16
17 // 0(1)快速乘
18
   11 mul2(11 x,11 y,11 p) {
        11 res=(x*y-(11)((long double)x/p*y+1.0e-8)*p);
19
20
        return res<0?res+p:res;</pre>
21 }
22
23 //int128
24 ll ans = ((__int128) a * b) % p;
26 // 10进制快速幂,直接读入%s,c 预处理字符串len
27 char c[1000005], len;
28 11 qp(11 a) {
29
       len --;
30
        a %= mod;
31
       11 s = a;
       ll res = 1;
32
33
        while (len \geq = 0) {
34
            11 cur = s;
35
            for (int i = 1; i <= c[len] - '0'; ++i) {</pre>
36
                res = res * s % mod;
37
            }
38
            for (int i = 1; i < 10; ++i) {</pre>
39
                cur = cur * s % mod;
40
            }
41
            s = cur;
42
            len --;
43
        }
44
        return res;
45 }
   0.2.3 PythonInput
1 // python一行读入
2 a,b = map(int, input().split())
3
4 a = []
5 for i in input().split():
6
       a.append(int(i))
    0.2.4 SpecialInput
1 // 代替gets
   scanf("%[^\n]%*c", ss)
2
4 void out2(int x, int flag = 1) {
```

```
if (x == 0) {
5
6
            if (flag) putchar('0');
7
            return;
        }
8
        out2(x >> 1, 0);
9
10
        putchar('0' + x % 2);
11 }
    0.3 header
1 // Editor -> Live Templates
2 // add template group acm
3 // add template main
4 // C++ Declaration
6 #include <bits/stdc++.h>
7
    #define 11 long long
8 using namespace std;
9
10 template <class T>
11 inline bool scan(T &ret){
12
        char c;
13
        int sgn;
14
        if (c = getchar(), c == EOF) return 0; //EOF
        while (c != '-' && (c < '0' || c > '9')) c = getchar();
15
16
        sgn = (c == '-') ? -1 : 1;
        ret = (c == '-') ? 0 : (c - '0');
17
18
        while (c = getchar(), c >= '0' && c <= '9') ret = ret * 10 + (c - '0');
19
        ret *= sgn;
20
        return 1;
21 }
22
23 const 11 mod = 1e9+7;
24 const int maxn = $MAXN$;
25
26 ll qp(ll x, ll n) {
27
       11 res = 1; x %= mod;
28
        while (n > 0) {
29
            if (n & 1) res = res * x % mod;
30
            x = x * x \% mod;
31
            n >>= 1;
32
        }
33
        return res;
34 }
35
36 int main(int argc, char* argv[]) {
        $CONTENT$
37
38
        return 0;
39 }
40
41 // C++ Expression debug
42 freopen("data.in", "r", stdin);
43 freopen("data.out", "w", stdout);
44 clock_t ST = clock();
45 cerr << "time: " << ((clock()-ST)*1000.0 / CLOCKS_PER_SEC) << "ms" << endl;
46
47
```

```
49 int T;
50 scanf("%d", &T);
51 for (int kase = 1; kase <= T; ++kase) { $CONTENT$;</pre>
52 }
   0.3.1 comp
1 // 1 create directory comp
3 // 3 create duipai.cpp
4 // add text
6 #include<bits/stdc++.h>
7 using namespace std;
8 int main(){
9
       int i;
10
     for (i=1;;i++){
           printf("The result of No. %d Case is: ",i);
11
12
           system("python3 rand.py");
           system("./std < test/data.in > test/std.out");
13
14
           system("./my < test/data.in > test/my.out");
           if (system("diff test/std.out test/my.out")){
15
16
               printf("Wrong Answer\n");
17
               return 0;
           }
18
19
           else printf("Accepted\n");
       }
20
21
       return 0;
22 }
23
24
25 // 4 create duipai.sh
26 #!/bin/bash
27 g++ std.cpp -o std
28 g++ my.cpp -o my
29 python3 rand.py
30 ./a.out
31
32 // 5 create rand.py
33 # coding=utf-8
34 from random import randint, choice, shuffle
35\, # with open("../cmake-build-debug/data.in", "w") as f:
36 with open("test/data.in", "w") as f:
37
       n = randint(1, 10)
38
       m = randint(1, 10)
39
       f.write(f"{n} {m}")
41 // 6 terminal: g++ duipai.cpp
42 // 7 terminal: sudo chmod 777 duipai.sh
43 // 8 add my.cpp and std.cpp
44 // 9 ./duipai.sh
```

1 Math

1.1 素数

7 8

9

10 11

12

13

14

15 }

1.1.1 Eratosthenes 筛法

```
1 bool vis[(int)1e6+5];
  int prim[(int)1e5], tot; // tot = 78499
3
   void init() {
4
       for (int i = 2; i < 1e6+5; ++i) {</pre>
           if (vis[i]) continue;
5
6
           prim[++tot] = i;
7
           for (int j = i + i; j < 1e6+5; j += i) vis[j] = 1;</pre>
8
       }
9 }
   1.1.2 Euler 筛
1 const int maxn = 1e6 + 10;
2 int prime[maxn], v[maxn], n, cnt; //每个合数只会被它的最小质因子p筛一次
4 void Euler_Sieve()
5 {
6
       for(int i = 2; i <= n; i ++)</pre>
```

if(!v[i]) v[i] = i, prime[++cnt] = i;

v[i * prime[j]] = prime[j];

if(i % prime[j] == 0) break;

for(int j = 1; j <= cnt && i * prime[j] <= n; j ++)</pre>

1.1.3 MillerRabin 素性测试

}

}

```
typedef long long 11;
2
3 bool check(ll a, ll n)
4
5
        if(n == 2 || a >= n) return true;
6
        if(n == 1 || !(n & 1)) return false;
7
        11 d = n - 1;
        while(!(d & 1)) d >>= 1;
8
9
        11 t = qp(a, d, n);
10
        while(d != n - 1 && t != 1 && t != n - 1)
11
        {
12
            t = mul(t, t, n);
13
            d <<= 1;
14
        }
15
        return t == n - 1 || d & 1;
16
   }
17
18 bool Miller_Rabin(ll n)
19
20
        static vector<11> t = {2, 325, 9375, 28178, 450775, 9780504, 1795265022};
21
        if (n <= 1) return false;</pre>
22
        for (ll k: t) if (!check(k, n)) return false;
```

```
23
        return true;
24 }
   1.1.4 PollardRho 快速因数分解
   mt19937 mt(time(0));
   ll pollard_rho(ll n, ll c) {
3
        ll x = uniform_int_distribution<ll>(1, n - 1)(mt), y = x;
        auto f = [\&](11 \ v) \{ 11 \ t = mul(v, v, n) + c; return \ t < n ? t : t - n; \};
4
5
       while (1) {
6
           x = f(x); y = f(f(y));
7
           if (x == y) return n;
8
           11 d = \_gcd(abs(x - y), n);
9
           if (d != 1) return d;
10
       }
   }
11
12
13 ll fac[100], fcnt;
   void get_fac(ll n, ll cc = 19260817) {
15
        if (n == 4) { fac[fcnt++] = 2; fac[fcnt++] = 2; return; }
16
        if (Miller_Rabin(n)) { fac[fcnt++] = n; return; }
17
       11 p = n;
18
       while (p == n) p = pollard_rho(n, --cc);
19
        get_fac(p); get_fac(n / p);
20 }
   1
   vector<pair<11, int>> getFactors(11 x)
2
   {
3
       vector<pair<ll, int>> fact;
        for (int i = 0; prime[i] <= x / prime[i]; i++)</pre>
4
5
           if (x % prime[i] == 0)
6
7
           {
8
               fact.emplace_back(prime[i], 0);
9
               while (x % prime[i] == 0) fact.back().second++, x /= prime[i];
10
11
        }
12
        if (x != 1) fact.emplace_back(x, 1);
13
       return fact;
14 }
   1.2 约数
   1.2.1 EulerPhi
1 //计算欧拉phi函数, phi(n)且与n互素的正整数个数
2
3
   int Euler(int n){
4
       int rea=n:
       for(int i=2; i*i<=n; i++)</pre>
5
6
           if(n%i==0)//第一次找到的必为素因子
7
8
               rea=rea-rea/i;
9
10
                   n/=i;//把该素因子全部约掉
```

```
11
                while(n%i==0);
12
            }
13
        if(n>1)
14
            rea=rea-rea/n;
15
        return rea;
16 } //单点欧拉 O(sqrt(n))
17
18 bool boo[50000];
   int p[20000];
19
   void prim(){
21
        memset(boo,0,sizeof(boo));
22
        boo[0]=boo[1]=1;
23
        int k=0;
24
        for(int i=2; i<50000; i++)</pre>
25
26
            if(!boo[i])
27
                p[k++]=i;
28
            for(int j=0; j<k&&i*p[j]<50000; j++)</pre>
29
30
                boo[i*p[j]=1;
31
                if(!(i%p[j]))
32
                    break;
33
            }
34
    }//筛选法打表
36
   int phi(int n)
37
    {
38
        int rea=n;
        for(int i=0; p[i]*p[i]<=n; i++)//对于一些不是素数的可不遍历
39
40
            if(n%p[i]==0)
41
            {
42
                rea=rea-rea/n;
43
                do
44
                     n/=p[i];
                while(n%p[i]==0);
45
            }
46
47
        if(n>1)
48
            rea=rea-rea/n;
49
        return rea;
   } //素数+欧拉
50
51
   int euler[maxn];
52
53
    void init() {
        int i, j;
54
        for(i=1; i<maxn; i++)</pre>
55
56
            euler[i]=i;
57
        for(i=2; i<maxn; i+=2)</pre>
            euler[i]/=2;
58
59
        for(i=3; i<maxn; i+=2)</pre>
60
            if(euler[i]==i) {
61
                for(j=i; j<=maxn; j+=i)</pre>
62
                     euler[j]=euler[j]/i*(i-1);
63
64 } //递推欧拉表
```

1.2.2 Sieve

1 //用类似筛法的方法计算phi(1),phi(2),...,phi(n)

```
2 int phi[maxn];
3
    void phi_table(int n)
4
5
6
        for (int i = 2; i <= n; i++) phi[i] = 0;</pre>
7
        phi[1] = 1;
8
        for (int i = 2; i <= n; i++) if (!phi[i])</pre>
9
            for (int j = i; j \le n; j += i)
10
11
                if (!phi[j]) phi[j] = j;
12
                phi[j] = phi[j] / i * (i - 1);
13
            }
14 }
    1.2.3 gcd
1 11 gcd(11 a,11 b) {while(b^=a^=b^=a%=b);return a;}
3 void exgcd(ll a, ll b, ll& x, ll& y, ll& c) {
        if(!b) {y = 0; x = 1; c = a; return;}
5
        exgcd(b, a \% b, y, x); y == a / b * x;
6 }
    1.2.4 解乘法逆元
1 void exgcd(ll a, ll b, ll c, ll d, ll &x, ll &y) {
2
        11 z = (a + b - 1) / b;
3
        if (z <= c / d) {
4
            x = z;
5
            y = 1;
6
            return;
7
        a = (z - 1) * b; c = (z - 1) * d;
8
9
        exgcd(d, c, b, a, y, x);
10
        x += (z - 1) * y;
11 }
12
13
   int main(int argc, char* argv[]) {
14
        int T;
        scanf("%d", &T);
15
16
        11 p, x;
17
        for (int kase = 1; kase <= T; ++kase) {</pre>
            scanf("%lld%lld", &p, &x);
18
            11 b, y;
19
            exgcd(p, x, p, x - 1, b, y);
20
21
            printf("%lld/%lld\n", b * x - p * y, b);
22
        }
23
        return 0;
24 }
    1.3 同余
    1.3.1 扩展欧几里得算法
1 void exgcd(int a, int b, int &x, int &y)
3
        if(b == 0) { x = 1; y = 0; return; }
```

```
exgcd(b, a % b, x, y);
5
        int t = x; x = y, y = t - a / b * y;
6 }
    1.3.2 中国剩余定理
1 typedef long long 11;
2
3
   void exgcd(ll a, ll b, ll &x, ll &y)
4
5
        if(b == 0) { x = 1; y = 0; return; }
6
        exgcd(b, a % b, x, y);
7
        11 t = x; x = y, y = t - a / b * y;
8
    }
9
10 ll crt(ll *a, ll *m, int n)
12
        11 M = 1, ans = 0;
13
        for(int i = 1; i <= n; i ++) M *= m[i];</pre>
        for(int i = 1; i <= n; i ++)</pre>
14
15
            11 x = 0, y = 0;
16
            11 Mi = M / m[i];
17
            exgcd(Mi, m[i], x, y);
18
19
            ans = (ans + Mi \% M * x \% M * a[i] \% M + M) \% M;
20
        }
21
        if(ans < 0) ans += M;
22
        return ans;
23 }
    1.3.3 扩展中国剩余定理
1 typedef long long 11;
2
3 const int N = 1e5 + 10;
4
5 int n;
6 ll a[N], r[N];
8 ll exgcd(ll a, ll b, ll& x, ll& y)
9 {
        if(b == 0) { x = 1, y = 0; return a; }
10
11
        11 ret = exgcd(b, a % b, y, x); y -= a / b * x;
12
        return ret;
13 }
14
15 ll excrt()
16 {
17
        11 M = a[1], R = r[1], x, y, d;
18
        for(int i = 2; i <= n; i ++)
19
20
            d = exgcd(M, a[i], x, y);
21
            if((R - r[i]) % d) return -1;
22
            x = (R - r[i]) / d * x % a[i];
23
            R -= M * x;
24
            M = M / d * a[i];
25
            R \%= M;
26
        }
```

```
27
        return (R % M + M) % M;
28 }
    1.3.4 BSGS
   int qp(int a, int n, int mod)
2
3
        long long ans = 1, base = a;
        while(n)
4
5
6
            if(n & 1) (ans *= base) %= mod;
7
            (base *= base) %= mod;
8
            n >>= 1;
9
        }
10
        return ans;
11 }
12
13 int BSGS(int a, int b, int p)
14 {
        map<int, int> hash;
15
16
        b %= p;
        int t = (int)sqrt(p) + 1;
17
18
        for(int j =0; j < t; j ++)</pre>
19
20
            int val = 111 * b * qp(a, j, p) % p;
21
            hash[val] = j;
22
        }
23
        a = qp(a, t, p);
24
        if(a == 0) return b == 0 ? 1 : -1;
25
        for(int i = 0; i <= t; i ++)</pre>
26
27
            int val = qp(a, i, p);
28
            int j = hash.find(val) == hash.end() ? -1 : hash[val];
29
            if(j \ge 0 \&\& i * t - j \ge 0) return i * t - j;
30
        }
31
        return -1;
32 }
    1.3.5 exBSGS
1 unordered_map<int, int> Hash;
2
3
    int exBSGS(int a, int b, int p)
4
    {
5
        a %= p, b %= p;
6
        if(b == 1) return 0;
7
        if(!b && !a) return 1;
8
        if(!a) return -1;
9
        if(!b)
10
11
            int ret = 0, d;
12
            while((d = \_gcd(a, p)) != 1)
13
14
                ++ ret, p /= d;
15
                if(p == 1) return ret;
16
17
            return -1;
18
        }
```

```
19
        int ret = 0, A = a, B = b, P = p, C = 1, d;
20
        while((d = __gcd(A, P)) != 1)
21
            if(B % d) return -1;
22
23
            P /= d, B /= d;
24
            C = 111 * C * (A / d) % P;
25
            ++ ret;
26
            if(C == B) return ret;
27
        }
28
       Hash.clear();
29
        int f = 1, t = sqrt(P) + 1;
30
       for(int i = 0; i < t; i ++)</pre>
31
32
            Hash[1ll * f * B % P] = i;
33
            f = 111 * f * A % P;
       }
34
35
       int tf = f;
36
       f = 111 * f * C % P;
37
       for(int i = 1; i <= t; i ++)</pre>
38
39
            if(Hash.find(f) != Hash.end()) return ret + i * t - Hash[f];
            f = 111 * f * tf % P;
40
       }
41
42
       return -1;
43 }
    1.3.6 逆元
1 /*
2
   1. 费马小定理
3 条件:mod为素数
4 */
5 11 inv(11 x){return qp(x,mod-2);}
6
7 /*
8 2.扩展欧几里得
9 条件:gcd(a,mod)==1
10 如果gcd(a,mod)!=1 返回-1
11 */
12 11 inv(11 a,11 p)
13 {
14
        11 g,x,y;
15
        g=exgcd(a,p,x,y);
16
       return g==1?(x+p)%p:-1;
17 }
18
19 /*
20 3.公式
21 \text{ a/b}/\text{mod=c}
22 \rightarrow a\%(b*mod)/b=c
23 */
24
25 /*
26 4.逆元打表
27 p是模
28 p要求是奇素数
29 */
30 ll inv[MAX];
```

```
31 void getinv(int n,ll p)
32 {
33
        11 i;
        inv[1]=1;
34
35
        for(i=2;i<=n;i++) inv[i]=(p-p/i)*inv[p%i]%p;</pre>
36 }
37
38 // log逆元
39 11 dlog(11 g, 11 b, 11 p) {
       11 m = sqrt(p - 1);
41
       map<11, 11> powers;
42
       for (long j = 0; j < m; j++) powers[qp(g, j, p)] = j;
43
       long gm = qp(g, -m + 2 * (p - 1), p);
       for (int i = 0; i < m; i++) {</pre>
44
45
           if (powers[b]) return i * m + powers[b];
46
           b = b * gm \% p;
        }
47
48
       return -1;
49 }
    1.3.7 模素数二次同余方程
   // 要求模为素数,输入n, mod,返回 x^2 % mod = n,可解任意一次二元方程
1
2
3
   bool Legendre(ll a,ll p) {
4
        return qp(a,p-1>>1,p)==1;
5
   }
6
7
   11 modsqr(ll a,ll p) {
8
       11 x;
9
        ll i,k,b;
10
        if(p==2) x=a\%p;
        else if(p\%4==3) x=qp(a,p+1>>2,p);
11
12
13
           for(b=1;Legendre(b,p);++b);
14
           i=p-1>>1;
15
           k=0;
16
           do
17
           {
18
               i>>=1;
19
20
               if(!((1LL*qp(a,i,p)*qp(b,k,p)+1)%p)) k+=p-1>>1;
21
           }while(!(i&1));
           x=111*qp(a,i+1>>1,p)*qp(b,k>>1,p)%p;
22
23
       }
24
       return min(x, p - x);
25 //
         if (p-x<x) x=p-x;
26
   //
          if(x==p-x) printf("%d\n",x);
27
   11
         else printf("%d %d\n",x,p-x);
28
    1.4 矩阵与线性方程组
    1.4.1 矩阵快速幂
 1 const int mod = 1e9 + 7;
   typedef long long 11;
3
```

```
4 int cur;
   struct Matrix {ll a[105][105]; };
7
   Matrix mul(Matrix a, Matrix b)
8
   {
9
        Matrix res;
10
        memset(res.a, 0, sizeof res.a);
11
        for(int i = 0; i < cur; i++)</pre>
12
            for(int j = 0; j < cur; j++)</pre>
13
                for(int k = 0; k < cur; k++)</pre>
                     (res.a[i][j] += a.a[i][k] * b.a[k][j] % mod) %= mod;
14
15
        return res;
16 }
17
18
   Matrix pow(Matrix a, ll n)
19
20
        Matrix ans, base = a;
21
        for(int i = 0; i < cur; i++) ans.a[i][i] = 1;</pre>
        while(n)
22
23
24
            if(n & 1) ans = mul(ans, base);
25
            base = mul(base, base);
26
            n >>= 1;
27
28
        return ans;
29 }
   1.4.2 高斯消元
   const int N = 20 + 10;
2
3 int n;
4 double b[N], c[N][N];
5 //c: 系数矩阵, b: 常数; 二者一起构成增广矩阵
6
7
   void Gaussian_Elimination()
8
9
        for(int i = 1; i <= n; i ++)</pre>
10
            //找到x[i]的系数不为0的一个方程
11
12
            for(int j = i; j \le n; j \leftrightarrow if(fabs(c[j][i]) > 1e-8)
13
14
                for(int k = 1; k <= n; k ++) swap(c[i][k], c[j][k]);</pre>
15
                swap(b[i], b[j]);
16
            //消去其他方程的x[i]的系数
17
18
            for(int j = 1; j <= n; j ++)</pre>
19
20
                if(i == j) continue;
21
                double rate = c[j][i] / c[i][i];
22
                for(int k = i; k <= n; k ++) c[j][k] -= c[i][k] * rate;</pre>
23
                b[j] -= b[i] * rate;
24
            }
25
        }
26 }
```

1.4.3 线性基

```
1 typedef long long ll;
2
3
    struct LinearBasis
4
    {
5
        11 d[64], tot;
6
7
        void ins(ll x) //插入线性基
8
9
            for(int i = 63; i >= 0; i --)
10
                if((x >> i) & 1)
11
12
                {
13
                    if(!d[i]) return void(d[i] = x);
14
                    x ^= d[i];
15
                }
16
            }
17
18
19
                        //在一个序列中取若干个数,使其异或和最大
        ll max_xor()
20
21
            11 \text{ ans} = 0;
22
            for(int i = 63; i >= 0; i --)
23
                if((ans ^ d[i]) > ans) ans ^= d[i];
24
            return ans;
25
        }
26
27
        void init()
28
29
            for(int i = 0; i < 64; i ++) if(d[i])</pre>
30
                for(int j = 0; j < i; j ++)
31
                    if(d[i] & (111 << j)) d[i] ^= d[j];</pre>
32
            for(int i = 0; i < 64; i ++) if(d[i]) d[tot ++] = d[i];</pre>
33
        }
34
        11 k_th(11 k) //取任意个元素进行异或的第k小个数
35
36
37
            //考虑能异或出0的情况, tot表示线性基中的元素个数
38
            k -= (n != tot);
            if(k > (1ll << tot)) return -1;</pre>
39
40
            11 \text{ ans} = 0;
41
            for(int i = 0; i < tot; i ++) if(k & (1ll << i)) ans ^= d[i];</pre>
42
            return ans;
43
        }
44 };
    1.5 组合数学
    1.5.1 Lucas
1 const int maxn = 1e6 + 10;
3 11 fac[maxn], inv[maxn], facinv[maxn];
4
5 void init()
6
7
        fac[0] = inv[0] = facinv[0] = 1;
8
        fac[1] = inv[1] = facinv[1] = 1;
9
        for(int i = 2; i < maxn; i++)</pre>
10
        {
```

```
11
            fac[i] = fac[i - 1] * i % mod;
12
            inv[i] = mod - mod / i * inv[mod % i] % mod;
13
            facinv[i] = facinv[i - 1] * inv[i] % mod;
14
15 }
16
17 11 C(int n, int k)
18 {
19
        if(k > n \mid \mid k < 0) return 0;
20
        return fac[n] * facinv[k] % mod * facinv[n - k] % mod;
21 }
22
23 11 lucas(11 n, 11 m)
24 {
25
        11 res = 1;
26
        while(n && m)
27
28
            res = res * C(n % mod, m % mod) % mod;
29
            n /= mod;
30
            m /= mod;
        }
31
32
        return res;
33 }
    1.5.2 exLucas
 1 typedef long long 11;
 2
 3 11 p, n, m;
 4
 5 ll exgcd(ll a, ll b, ll &x, ll &y)
 6
 7
        if(!b) { x = 1; y = 0; return a; }
 8
        ll res = exgcd(b, a \% b, x, y), t;
 9
        t = x, x = y, y = t - a / b * y;
10
        return res;
11 }
12
13 ll qp(ll a, ll n, ll mod)
14 {
15
        11 \text{ ans} = 1, \text{ base} = a;
16
        for(; n; n >>= 1, (base *= base) %= mod) if(n & 1) (ans *= base) %= mod;
17
        return ans;
18 }
19
20 ll fac(ll n, ll a, ll b)
21 {
22
        if(!n) return 1;
23
        11 \text{ res} = 1;
24
        for(ll i = 2; i <= b; i ++)
25
            if(i % a) (res *= i) %= b;
26
        res = qp(res, n / b, b);
27
        for(l1 i = 2; i <= n % b; i ++)</pre>
            if(i % a) (res *= i) %= b;
28
        return res * fac(n / a, a, b) % b;
29
30 }
32 ll inv(ll n, ll mod)
```

```
33 {
34
        11 x, y;
35
        exgcd(n, mod, x, y);
        return (x += mod) > mod ? x - mod : x;
36
37
38
39
   ll CRT(ll b, ll mod) { return b * inv(p / mod, mod) % p * (p / mod) % p; }
40
41
   ll C(ll n, ll m, ll a, ll b)
42 {
        ll up = fac(n, a, b), d1 = fac(m, a, b), d2 = fac(n - m, a, b);
43
44
        11 k = 0;
45
        for(ll i = n; i; i /= a) k += i / a;
46
        for(ll i = m; i; i /= a) k -= i / a;
47
        for(ll i = n - m; i; i /= a) k -= i / a;
48
        return up * inv(d1, b) % b * inv(d2, b) % b * qp(a, k, b) % b;
49
50
51 ll exlucas(ll n, ll m)
52 {
53
        ll res = 0, tmp = p, b;
54
        int lim = sqrt(p) + 5;
55
        for(int i = 2; i <= lim; i ++) if(tmp % i == 0)</pre>
56
57
            b = 1;
58
            while(tmp % i == 0) tmp /= i, b *= i;
59
            (res += CRT(C(n, m, i, b), b)) \%= p;
60
61
        if(tmp > 1) (res += CRT(C(n, m, tmp, tmp), tmp)) %= p;
62
        return res;
63 }
   1.5.3 递推组合数
   0 \le m \le n \le 1000
1 const int maxn = 1010;
   11 C[maxn] [maxn];
3
   void init() {
        C[0][0] = 1;
4
5
        for (int i = 1; i < maxn; i++)</pre>
6
        {
7
            C[i][0] = 1;
8
            for (int j = 1; j \le i; j++) C[i][j] = (C[i-1][j-1] + C[i-1][j]) % mod;
9
10 }
   0 \le m \le n \le 10^5, 模 p 为素数
1 const int maxn = 100010;
2 11 f[maxn];
   ll inv[maxn]; // 阶乘的逆元
4
   void CalFact() {
5
        f[0] = 1;
6
        for (int i = 1; i < maxn; i++) f[i] = (f[i - 1] * i) % p;</pre>
7
        inv[maxn - 1] = qp(f[maxn - 1], p - 2);
8
        for (int i = maxn - 2; ~i; i--) inv[i] = inv[i + 1] * (i + 1) % p;
9 }
10 ll C(int n, int m) { return f[n] * inv[m] % p * inv[n - m] % p; }
```

1.5.4 小模数组合数

p 小 n,m 大

```
1
2
   const int NICO = 100000+10;
3
   const int MOD = 99991;
4 11 f[NICO];
6 ll Lucas(ll a, ll k)
7
8
        11 res = 1;
9
        while(a && k)
10
            11 a1 = a % MOD;
11
12
            ll b1 = k \% MOD;
13
            if(a1 < b1) return 0;</pre>
            res = res*f[a1]*qp(f[b1]*f[a1-b1]%MOD,MOD-2)%MOD;
14
            a /= MOD;
15
16
            k /= MOD;
        }
17
18
        return res;
19
   }
20
   void init()
21
22 {
23
        f[0] = 1;
24
        for(int i=1;i<=MOD;i++)</pre>
25
26
            f[i] = f[i-1]*i%MOD;
27
        }
28 }
29
30 int main()
31 {
32
        init();
33
        cout << Lucas(5,2) << endl;</pre>
34 }
    1.5.5 大模数组合数
    n,m 小 p 大
1 map<int, 11> m;
    const int MOD = 1e9+7;
3
    void fun(int n, int k) {
4
        for (int i = 2; i <= sqrt(n * 1.0); i++) {</pre>
5
            while (n % i == 0) {
6
7
                n /= i;
8
                m[i] += k;
9
            }
10
        }
11
        if (n > 1) {
12
            m[n] += k;
13
        }
14 }
15
```

```
16
   11 C(11 a, 11 b) {
17
        if (a < b || a < 0 || b < 0)</pre>
18
            return 0;
19
        m.clear();
20
        ll ret = 1;
21
        b = min(a - b, b);
22
        for (int i = 0; i < b; i++) {</pre>
23
            fun(a - i, 1);
24
        }
25
        for (int i = b; i >= 1; i--) {
26
            fun(i, -1);
27
        }
28
        for (__typeof(m.begin()) it = m.begin(); it != m.end(); it++) {
29
            if ((*it).second != 0) {
30
                ret *= qp((*it).first, (*it).second);
                ret %= MOD;
31
32
33
        }
34
        return ret;
35 }
36
37 int main(int argc,char *argv[])
38 {
39
        11 a, b;
40
        while (scanf("%11d%11d", &a, &b) != EOF) {
41
            printf("%lld\n", C(a, b));
42
        }
43
        return 0;
44 }
          卷积
    1.6
    1.6.1 FFT
1 const int maxn = 1e7 + 10;
2
   const double Pi = acos(-1.0);
3
4 struct complex
5
   {
6
        double x, y;
7
        complex (double xx = 0, double yy = 0) { x = xx, y = yy; }
8
   }a[maxn], b[maxn];
9
10
  complex operator + (complex a, complex b) { return complex(a.x + b.x, a.y + b.y); }
   complex operator - (complex a, complex b) { return complex(a.x - b.x, a.y - b.y); }
   complex operator * (complex a, complex b) { return complex(a.x * b.x - a.y * b.y, a.x * b.y + a.y *
         b.x); }
13
14 int n, m;
15 int 1, r[maxn];
16 int limit;
17
18 void FFT(complex *A, int type)
19
20
        for(int i = 0; i < limit; i ++)</pre>
21
            if(i < r[i]) swap(A[i], A[r[i]]);</pre>
22
        for(int mid = 1; mid < limit; mid <<= 1)</pre>
23
            complex Wn(cos(Pi / mid), type * sin(Pi / mid));
24
```

```
25
            for(int R = mid << 1, j = 0; j < limit; j += R)</pre>
26
27
                 complex w(1, 0);
28
                 for(int k = 0; k < mid; k ++, w = w * Wn)</pre>
29
30
                     complex x = A[j + k], y = w * A[j + mid + k];
31
                     A[j + k] = x + y;
32
                     A[j + mid + k] = x - y;
33
                 }
34
            }
35
        }
36
    }
37
38 void mul()
39
        1 = 0, limit = 1;
40
41
        while(limit <= n + m) limit <<= 1, 1 ++;</pre>
42
        for(int i = 0; i < limit; i ++)</pre>
43
            r[i] = (r[i >> 1] >> 1) | ((i & 1) << (1 - 1));
44
        FFT(a, 1);
        FFT(b, 1);
45
        for(int i = 0; i <= limit; i ++) a[i] = a[i] * b[i];</pre>
46
47
        FFT(a, -1);
48
        for(int i = 0; i <= n + m; i ++)</pre>
49
            printf("%d ", (int)(a[i].x / limit + 0.5));
50 }
    1.6.2 NTT
 1 const int maxn = 2097152;
 2 const int mod = 998244353;
 3 const int root = 3;
 4 // 998244353 -> 3, 1e9+7 -> 5,
 5
 6 template<long long mod, long long root>
 7
    struct NTT {
 8
        vector<long long> omega;
 9
10
        NTT() {
11
            omega.resize(maxn + 1);
12
            long long x = fpow(root, (mod - 1) / maxn);
13
            omega[0] = 111;
14
            for (int i = 1; i <= maxn; ++i)</pre>
15
                 omega[i] = omega[i - 1] * x % mod;
16
        }
17
18
        long long fpow(long long a, long long n) {
19
             (n += mod - 1) \% = mod - 1;
20
            long long r = 1;
21
            for (; n; n >>= 1) {
22
                 if (n & 1) (r *= a) %= mod;
23
                 (a *= a) \% = mod;
24
            }
25
            return r;
26
        }
27
28
        void bitrev(vector<long long> &v, int n) {
            int z = __builtin_ctz(n) - 1;
29
```

```
30
             for (int i = 0; i < n; ++i) {</pre>
31
                 int x = 0;
32
                 for (int j = 0; j \le z; ++j) x = (i >> j & 1) << (z - j);
33
                 if (x > i) swap(v[x], v[i]);
34
             }
35
        }
36
        void ntt(vector<long long> &v, int n) {
37
38
             bitrev(v, n);
39
             for (int s = 2; s <= n; s <<= 1) {
                 int z = s >> 1;
40
41
                 for (int i = 0; i < n; i += s) {</pre>
42
                     for (int k = 0; k < z; ++k) {
43
                          long long x = v[i + k + z] * omega[maxn / s * k] % mod;
44
                          v[i + k + z] = (v[i + k] + mod - x) \% mod;
45
                          (v[i + k] += x) \% = mod;
46
47
                 }
48
             }
        }
49
50
        void intt(vector<long long> &v, int n) {
51
52
             ntt(v, n);
             for (int i = 1; i < n / 2; ++i) swap(v[i], v[n - i]);</pre>
53
54
             long long inv = fpow(n, -1);
55
             for (int i = 0; i < n; ++i) (v[i] *= inv) %= mod;</pre>
        }
56
57
58
        vector<long long> operator()(vector<long long> a, vector<long long> b) {
59
             int sz = 1;
             while (sz < a.size() + b.size() - 1) sz <<= 1;</pre>
60
61
             while (a.size() < sz) a.push_back(0);</pre>
62
             while (b.size() < sz) b.push_back(0);</pre>
63
             ntt(a, sz), ntt(b, sz);
64
             vector<long long> c(sz);
65
             for (int i = 0; i < sz; ++i) c[i] = a[i] * b[i] % mod;</pre>
66
             intt(c, sz);
67
             while (c.size() && c.back() == 0) c.pop_back();
68
             return c;
        }
69
70
71
        vector<long long> operator()(vector<long long> a, int n) {
72
             int sz = 1;
73
             while (sz < n * a.size()) sz <<= 1;</pre>
74
             while (a.size() < sz) a.push_back(0);</pre>
75
             ntt(a, sz);
76
             for (int i = 0; i < sz; ++i) a[i] = fpow(a[i], n);</pre>
77
             intt(a, sz);
78
             while (a.size() && a.back() == 0) a.pop_back();
79
             return a;
80
        }
81 };
82
83 NTT<mod, root> conv;
```

1.6.3 原根

1 #include<bits/stdc++.h>

```
2 #define 11 long long
3 #define IL inline
4 #define RG register
5 using namespace std;
7
   11 prm[1000],tot,N,root;
8
9
    11 Power(11 bs,11 js,11 MOD){
10
        11 S = 1,T = bs;
11
        while(js){
12
            if(js&1)S = S*T%MOD;
13
            T = T*T\%MOD;
            js >>= 1;
14
15
        } return S;
   }
16
17
18
    IL 11 GetRoot(RG 11 n){
19
        RG ll tmp = n - 1 , tot = 0;
        for(RG 11 i = 2; i <= sqrt(tmp); i ++){</pre>
20
21
            if(tmp%i==0){
22
                prm[++tot] = i;
23
                while(tmp%i==0)tmp /= i;
24
            }
25
26
        if(tmp != 1)prm[++tot] = tmp;
                                                 //质因数分解
27
        for(RG 11 g = 2; g <= n-1; g ++){</pre>
28
            bool flag = 1;
29
            for(RG int i = 1; i <= tot; i ++){</pre>
                                                   //检测是否符合条件
                if(Power(g,(n-1)/prm[i],n) == 1)
30
31
                    { flag = 0; break; }
32
            }
33
            if(flag)return g;
34
                                          //无解
        }return 0;
35 }
36
37
   int main(){
38
        cin >> N;
39
        root = GetRoot(N);
40
        cout<<root<<endl;</pre>
41
        return 0;
42 }
    1.6.4 FWT
1 //C_k=\sum_{i \in A_i B_j}
2 //FWT 完后需要先模一遍
3 template<typename T>
4 void fwt(ll a[], int n, T f) {
5
        for (int d = 1; d < n; d *= 2)
6
            for (int i = 0, t = d * 2; i < n; i += t)
7
                for(int j = 0; j < d; j ++)
8
                    f(a[i + j], a[i + j + d]);
9
   }
10
11 void AND(11& a, 11& b) { a += b; }
12 void OR(11& a, 11& b) { b += a; }
13 void XOR (11& a, 11& b) {
        11 x = a, y = b;
14
```

```
15
       a = (x + y) \% mod;
16
       b = (x - y + mod) \% mod;
17 }
18 void rAND(ll& a, ll& b) { a -= b; }
19 void rOR(11& a, 11& b) { b -= a; }
20 void rXOR(11& a, 11& b) {
21
       static 11 INV2 = (mod + 1) / 2;
22
       11 x = a, y = b;
23
       a = (x + y) * INV2 % mod;
24
       b = (x - y + mod) * INV2 % mod;
25 }
26
27 //FWT 子集卷积
28 a[popcount(x)][x] = A[x]
29 b[popcount(x)][x] = B[x]
30 fwt(a[i]) fwt(b[i])
31 c[i + j][x] += a[i][x] * b[j][x]
32 rfwt(c[i])
33 ans[x] = c[popcount(x)][x]
```

1.7 多项式

1.7.1 拉格朗日插值

```
1 typedef long long 11;
 2
 3 const int mod = 998244353;
 4 const int maxn = 1e5 + 10;
 5
 6 int x[maxn], y[maxn];
 7
 8 int qp(int a, int n)
 9 {
10
        11 \text{ ans} = 1, \text{ base} = a;
11
        for(; n; (base *= base) %= mod, n >>= 1) if(n & 1) (ans *= base) %= mod;
        return ans;
12
13 }
14
    int lagrange(int n, int *x, int *y, int xi)
15
16
    {
17
        int ans = 0;
18
        for(int i = 0; i <= n; i ++)</pre>
19
20
            int s1 = 1, s2 = 1;
21
            for(int j = 0; j \le n; j ++) if(i != j)
22
23
                s1 = 111 * s1 * (xi - x[j]) % mod;
24
                s2 = 111 * s2 * (x[i] - x[j]) % mod;
25
            }
26
            ans = (111 * ans + 111 * y[i] * s1 % mod * qp(s2, mod - 2) % mod) % mod;
27
        }
28
        return (ans + mod) % mod;
29 }
```

1.7.2 拉格朗日插值(连续取值)

```
1 const int mod = 'edit';
2 const int maxn = 'edit';
```

```
3
   int x[maxn], y[maxn];
   int s1[maxn], s2[maxn], ifac[maxn];
   //如果x的取值是连续一段,可以做到0(n)求解
7
   int lagrange(int n, int *x, int *y, int xi)
8
9
10
        int ans = 0;
        s1[0] = (xi - x[0]) \% mod, s2[n + 1] = 1;
11
12
        for(int i = 1; i <= n; i ++) s1[i] = 111 * s1[i - 1] * (xi - x[i]) % mod;
        for(int i = n; i >= 0; i --) s2[i] = 111 * s2[i + 1] * (xi - x[i]) % mod;
13
        ifac[0] = ifac[1] = 1;
14
15
        for(int i = 2; i <= n; i ++) ifac[i] = -111 * mod / i * ifac[mod % i] % mod;</pre>
        for(int i = 2; i <= n; i ++) ifac[i] = 111 * ifac[i] * ifac[i - 1] % mod;</pre>
16
17
        for(int i = 0; i <= n; i ++)</pre>
            (ans += 1ll * y[i] * (i == 0 ? 1 : s1[i - 1]) % mod * s2[i + 1] % mod * ifac[i] % mod * (((
18
        n - i) & 1) ? -1 : 1) * ifac[n - i] % mod) %= mod;
19
        return (ans + mod) % mod;
20
   }
```

1.8 Others

- 1. 约数定理: 若 $n = \prod_{i=1}^{k} p_i^{a_i}$,则
 - (a) 约数个数 $f(n) = \prod_{i=1}^{k} (a_i + 1)$
 - (b) 约数和 $g(n) = \prod_{i=1}^{k} (\sum_{j=0}^{a_i} p_i^j)$
- 2. 小于 n 且互素的数之和为 $n\varphi(n)/2$
- 3. 若 gcd(n, i) = 1, 则 $gcd(n, n i) = 1(1 \le i \le n)$
- 4. 错排公式: $D(n) = (n-1)(D(n-2) + D(n-1)) = \sum_{i=2}^{n} \frac{(-1)^{k} n!}{k!} = \left[\frac{n!}{e} + 0.5\right]$
- 5. 威尔逊定理: $p \text{ is } prime \Rightarrow (p-1)! \equiv -1 \pmod{p}$
- 6. 欧拉定理: $gcd(a,n) = 1 \Rightarrow a^{\varphi(n)} \equiv 1 \pmod{n}$
- 7. 欧拉定理推广: $\gcd(n,p) = 1 \Rightarrow a^n \equiv a^{n\%\varphi(p)} \pmod{p}$
- 8. 模的幂公式: $a^n \pmod m = \begin{cases} a^n \mod m & n < \varphi(m) \\ a^{n\%\varphi(m)+\varphi(m)} \mod m & n \ge \varphi(m) \end{cases}$
- 9. 素数定理: 对于不大于 n 的素数个数 $\pi(n)$, $\lim_{n\to\infty}\pi(n)=\frac{n}{\ln n}$
- 10. 位数公式: 正整数 x 的位数 $N = \log_{10}(n) + 1$
- 11. 斯特灵公式 $n! \approx \sqrt{2\pi n} \left(\frac{n}{r}\right)^n$
- 13. 设 a > b, gcd(a, b) = 1, 则 $gcd(a^m b^m, a^n b^n) = a^{gcd(m, n)} b^{gcd(m, n)}$

$$G=\gcd(C_n^1,C_n^2,...,C_n^{n-1})=\begin{cases} n, & n \text{ is prime}\\ 1, & n \text{ has multy prime factors}\\ p, & n \text{ has single prime factor } p \end{cases}$$

gcd(Fib(m), Fib(n)) = Fib(gcd(m, n))

- 14. 若 gcd(m, n) = 1, 则:
 - (a) 最大不能组合的数为 m*n-m-n
 - (b) 不能组合数个数 $N = \frac{(m-1)(n-1)}{2}$
- 15. $(n+1)lcm(C_n^0, C_n^1, ..., C_n^{n-1}, C_n^n) = lcm(1, 2, ..., n+1)$
- 16. 若 p 为素数,则 $(x+y+...+w)^p \equiv x^p + y^p + ... + w^p \pmod{p}$

17. 卡特兰数: 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012

$$h(0) = h(1) = 1, h(n) = \frac{(4n-2)h(n-1)}{n+1} = \frac{C_{2n}^n}{n+1} = C_{2n}^n - C_{2n}^{n-1}$$

18. 伯努利数: $B_n = -\frac{1}{n+1} \sum_{i=0}^{n-1} C_{n+1}^i B_i$

$$\sum_{i=1}^{n} i^{k} = \frac{1}{k+1} \sum_{i=1}^{k+1} C_{k+1}^{i} B_{k+1-i} (n+1)^{i}$$

19. 二项式反演:

$$f_n = \sum_{i=0}^n (-1)^i \binom{n}{i} g_i \Leftrightarrow g_n = \sum_{i=0}^n (-1)^i \binom{n}{i} f_i$$

$$f_n = \sum_{i=0}^n \binom{n}{i} g_i \Leftrightarrow g_n = \sum_{i=0}^n (-1)^{n-i} \binom{n}{i} f_i$$

20. 莫比乌斯反演:

(a)
$$\Rightarrow f(d) = \sum_{i=1}^{n} \sum_{j=1}^{m} [gcd(i,j) = d]$$

(b)
$$F(n) = \sum_{n|d} f(d) = \lfloor \frac{N}{n} \rfloor \lfloor \frac{M}{n} \rfloor$$

(c) 有
$$f(n) = \sum_{n|d} \mu(\lfloor \frac{d}{n} \rfloor) F(d)$$

(d)
$$\phi(n) = \sum_{d|n} d * \mu(n/d)$$

- 21. 2 的 n 次方, 在 pow 时可以精确输出最大 2¹023, pow(2,1023)
- 22. FFT 常用素数

$r 2^k + 1$	r	k	g
3	1	1	2
5	1	2	2
17	1	4	3
97	3	5	5
193	3	6	5
257	1	8	3
7681	15	9	17
12289	3	12	11
40961	5	13	3
65537	1	16	3
786433	3	18	10
5767169	11	19	3
7340033	7	20	3
23068673	11	21	3
104857601	25	22	3
167772161	5	25	3
469762049	7	26	3
998244353	119	23	3
1004535809	479	21	3
2013265921	15	27	31
2281701377	17	27	3
3221225473	3	30	5
75161927681	35	31	3
77309411329	9	33	7
206158430209	3	36	22
2061584302081	15	37	7
2748779069441	5	39	3
6597069766657	3	41	5
39582418599937	9	42	5
79164837199873	9	43	5
263882790666241	15	44	7
1231453023109121	35	45	3
1337006139375617	19	46	3
3799912185593857	27	47	5
4222124650659841	15	48	19
7881299347898369	7	50	6
31525197391593473	7	52	3
180143985094819841	5	55	6
1945555039024054273	27	56	5
4179340454199820289	29	57	3

1.8.1 BM

```
1 //Berlekamp-Massey
2 typedef vector<int> VI;
3 namespace linear_seq
4 {
5 #define rep(i,a,n) for (int i=a;i<n;i++)</pre>
6 #define SZ(x) ((int)(x).size())
7 #define pb(x) push_back(x)
        const ll mod=1e9+7;
8
9
        ll powmod(ll a,ll b){ll res=1;a%=mod; assert(b>=0); for(;b;b>>=1){if(b&1)res=res*a%mod;a=a*a%
        mod;}return res;}
10
        const int N=10010;
11
        11 res[N],base[N],_c[N],_md[N];
12
        vector<int> Md;
13
        void mul(ll *a,ll *b,int k)
14
        {
```

```
15
                                        rep(i,0,k+k) _c[i]=0;
                                         rep(i,0,k) if (a[i]) rep(j,0,k) _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;
16
17
                                         for (int i=k+k-1;i>=k;i--) if (_c[i])
18
                                                                    rep(j,0,SZ(Md)) _c[i-k+Md[j]]=(_c[i-k+Md[j]]-_c[i]*_md[Md[j]])%mod;
19
                                        rep(i,0,k) a[i]=_c[i];
20
                           }
21
                           int solve(ll n,VI a,VI b){
22
                                        11 ans=0,pnt=0;
23
                                         int k=SZ(a);
24
                                         assert(SZ(a)==SZ(b));
25
                                        rep(i,0,k) _md[k-1-i]=-a[i];_md[k]=1;
26
                                         Md.clear();
27
                                        rep(i,0,k) if (_md[i]!=0) Md.push_back(i);
28
                                         rep(i,0,k) res[i]=base[i]=0;
29
                                         res[0]=1;
                                         while ((111<<pnt)<=n) pnt++;</pre>
30
31
                                         for (int p=pnt;p>=0;p--) {
32
                                                      mul(res,res,k);
33
                                                      if ((n>>p)&1) {
                                                                   for (int i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
34
35
                                                                   \label{eq:condition} $\operatorname{rep}(j,0,SZ(Md))$ $\operatorname{res}[Md[j]] = (\operatorname{res}[Md[j]] - \operatorname{res}[k] *_m d[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[k] *_m d[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[k] *_m d[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[k] *_m d[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[k] *_m d[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[k] *_m d[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[Md[j]] - \operatorname{res}[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[Md[j]] - \operatorname{res}[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{res}[Md[j]] - \operatorname{res}[Md[j]]) \% $\operatorname{mod};$ $\operatorname{rep}(j,0,SZ(Md)) = (\operatorname{rep}(j,0,SZ(Md)) - \operatorname{rep}(j,SZ(Md))) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md))) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md))) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md))) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) = (\operatorname{rep}(j,SZ(Md)) - \operatorname{rep}(j,SZ(Md)) = (\operatorname{rep}(j,SZ(M
                                                      }
36
                                        }
37
38
                                        rep(i,0,k) ans=(ans+res[i]*b[i])%mod;
39
                                         if (ans<0) ans+=mod;</pre>
40
                                        return ans;
                           }
41
                           VI BM(VI s){
42
                                        VI C(1,1),B(1,1);
43
44
                                         int L=0,m=1,b=1;
45
                                        rep(n,0,SZ(s)){
46
                                                      11 d=0;
                                                      rep(i,0,L+1) d=(d+(ll)C[i]*s[n-i])%mod;
47
                                                      if(d==0) ++m;
48
                                                      else if(2*L \le n){
49
50
                                                                   VI T=C;
51
                                                                    11 c=mod-d*powmod(b,mod-2)%mod;
52
                                                                    while (SZ(C) < SZ(B) + m) C.pb(0);
53
                                                                   rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
                                                                   L=n+1-L; B=T; b=d; m=1;
54
55
                                                      } else {
                                                                    11 c=mod-d*powmod(b,mod-2)%mod;
56
57
                                                                    while (SZ(C) < SZ(B) + m) C.pb(0);
58
                                                                   rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
59
                                                                    ++m;
60
                                                      }
61
                                        }
62
                                         return C;
                           }
63
64
                           int gao(VI a,ll n)
65
                           {
66
                                         VI c=BM(a);
67
                                         c.erase(c.begin());
68
                                         rep(i,0,SZ(c)) c[i]=(mod-c[i])%mod;
69
                                         return solve(n,c,VI(a.begin(),a.begin()+SZ(c)));
70
71 };//linear_seq::gao(VI{},n-1)
```

1.8.2 exBM

```
2 // given first 2 *m items init[0..2m-1], it will compute trans[0..m-1]
3 // for you. trans[0..m] should be given as that
           init[m] = sum_{i=0}^{m-1} init[i] * trans[i]
5
   struct LinearRecurrence
6
   {
7
        using int64 = long long;
8
        using vec = std::vector<int64>;
9
10
        static void extand(vec& a, size_t d, int64 value = 0)
11
12
            if (d <= a.size()) return;</pre>
13
            a.resize(d, value);
14
        }
15
        static vec BerlekampMassey(const vec& s, int64 mod)
16
17
            std::function<int64(int64)> inverse = [&](int64 a) {
                return a == 1 ? 1 : (int64)(mod - mod / a) * inverse(mod % a) % mod;
18
19
            };
20
            vec A = \{1\}, B = \{1\};
21
            int64 b = s[0];
22
            for (size_t i = 1, m = 1; i < s.size(); ++i, m++)</pre>
23
24
                int64 d = 0;
25
                for (size_t j = 0; j < A.size(); ++j)</pre>
26
27
                    d += A[j] * s[i - j] % mod;
28
29
                if (!(d %= mod)) continue;
30
                if (2 * (A.size() - 1) <= i)</pre>
31
32
                    auto temp = A;
33
                    extand(A, B.size() + m);
                    int64 coef = d * inverse(b) % mod;
34
35
                    for (size_t j = 0; j < B.size(); ++j)</pre>
36
                        A[i + m] -= coef * B[j] % mod;
37
38
                        if (A[j + m] < 0) A[j + m] += mod;
39
40
                    B = temp, b = d, m = 0;
41
                }
42
                else
43
44
                    extand(A, B.size() + m);
45
                    int64 coef = d * inverse(b) % mod;
46
                    for (size_t j = 0; j < B.size(); ++j)</pre>
47
48
                        A[j + m] -= coef * B[j] % mod;
49
                        if (A[j + m] < 0) A[j + m] += mod;
50
51
                }
52
            }
53
            return A;
54
        }
55
        static void exgcd(int64 a, int64 b, int64& g, int64& x, int64& y)
56
            if (!b)
57
```

```
58
                 x = 1, y = 0, g = a;
 59
             else
 60
             {
 61
                 exgcd(b, a % b, g, y, x);
 62
                 y = x * (a / b);
 63
             }
 64
         }
 65
         static int64 crt(const vec& c, const vec& m)
 66
 67
             int n = c.size();
 68
             int64 M = 1, ans = 0;
 69
             for (int i = 0; i < n; ++i) M *= m[i];</pre>
 70
             for (int i = 0; i < n; ++i)</pre>
 71
             {
 72
                  int64 x, y, g, tm = M / m[i];
                  exgcd(tm, m[i], g, x, y);
 73
 74
                 ans = (ans + tm * x * c[i] \% M) \% M;
 75
 76
             return (ans + M) % M;
         }
 77
 78
         static vec ReedsSloane(const vec& s, int64 mod)
 79
 80
             auto inverse = [](int64 a, int64 m) {
 81
                 int64 d, x, y;
 82
                 exgcd(a, m, d, x, y);
 83
                 return d == 1 ? (x % m + m) % m : -1;
 84
             };
 85
             auto L = [](const vec& a, const vec& b) {
 86
                  int da = (a.size() > 1 || (a.size() == 1 && a[0])) ? a.size() - 1 : -1000;
 87
                  int db = (b.size() > 1 || (b.size() == 1 && b[0])) ? b.size() - 1 : -1000;
 88
                 return std::max(da, db + 1);
 89
             };
 90
             auto prime_power = [&](const vec& s, int64 mod, int64 p, int64 e) {
 91
                 // linear feedback shift register mod p^e, p is prime
 92
                 std::vector<vec> a(e), b(e), an(e), bn(e), ao(e), bo(e);
 93
                 vec t(e), u(e), r(e), to(e, 1), uo(e), pw(e + 1);
 94
 95
                 pw[0] = 1;
 96
                 for (int i = pw[0] = 1; i <= e; ++i) pw[i] = pw[i - 1] * p;</pre>
                 for (int64 i = 0; i < e; ++i)</pre>
 97
98
99
                      a[i] = {pw[i]}, an[i] = {pw[i]};
100
                      b[i] = \{0\}, bn[i] = \{s[0] * pw[i] % mod\};
101
                      t[i] = s[0] * pw[i] % mod;
102
                      if (t[i] == 0)
103
                      {
104
                          t[i] = 1, u[i] = e;
                      }
105
106
                      else
                      {
107
108
                          for (u[i] = 0; t[i] % p == 0; t[i] /= p, ++u[i])
109
110
111
112
                 for (size_t k = 1; k < s.size(); ++k)</pre>
113
114
                      for (int g = 0; g < e; ++g)
115
116
                          if (L(an[g], bn[g]) > L(a[g], b[g]))
```

```
{
117
118
                               ao[g] = a[e - 1 - u[g]];
119
                               bo[g] = b[e - 1 - u[g]];
120
                               to[g] = t[e - 1 - u[g]];
121
                               uo[g] = u[e - 1 - u[g]];
122
                               r[g] = k - 1;
                          }
123
124
                      }
125
                      a = an, b = bn;
126
                      for (int o = 0; o < e; ++o)</pre>
127
128
                           int64 d = 0;
129
                          for (size_t i = 0; i < a[o].size() && i <= k; ++i)</pre>
130
131
                               d = (d + a[o][i] * s[k - i]) % mod;
132
                           }
133
                           if (d == 0)
134
                           {
135
                               t[o] = 1, u[o] = e;
                          }
136
137
                           else
                           {
138
139
                               for (u[o] = 0, t[o] = d; t[o] % p == 0; t[o] /= p, ++u[o])
140
141
                               int g = e - 1 - u[o];
142
                               if (L(a[g], b[g]) == 0)
143
144
                                   extand(bn[o], k + 1);
145
                                   bn[o][k] = (bn[o][k] + d) \% mod;
                               }
146
147
                               else
148
                               {
                                   int64 coef = t[o] * inverse(to[g], mod) % mod * pw[u[o] - uo[g]] % mod;
149
150
                                   int m = k - r[g];
                                   extand(an[o], ao[g].size() + m);
151
152
                                   extand(bn[o], bo[g].size() + m);
153
                                   for (size_t i = 0; i < ao[g].size(); ++i)</pre>
154
                                   {
155
                                       an[o][i + m] -= coef * ao[g][i] % mod;
                                       if (an[o][i + m] < 0) an[o][i + m] += mod;</pre>
156
157
158
                                   while (an[o].size() && an[o].back() == 0) an[o].pop_back();
159
                                   for (size_t i = 0; i < bo[g].size(); ++i)</pre>
160
                                   {
161
                                       bn[o][i + m] -= coef * bo[g][i] % mod;
                                       if (bn[o][i + m] < 0) bn[o][i + m] -= mod;</pre>
162
163
                                   while (bn[o].size() && bn[o].back() == 0) bn[o].pop_back();
164
165
                               }
                          }
166
167
                      }
168
169
                  return std::make_pair(an[0], bn[0]);
170
             };
171
172
              std::vector<std::tuple<int64, int64, int>> fac;
173
             for (int64 i = 2; i * i <= mod; ++i)</pre>
174
             {
175
                  if (mod % i == 0)
```

```
176
                 {
177
                      int64 cnt = 0, pw = 1;
178
                      while (mod % i == 0) mod /= i, ++cnt, pw *= i;
179
                      fac.emplace_back(pw, i, cnt);
180
181
             }
182
             if (mod > 1) fac.emplace_back(mod, mod, 1);
183
             std::vector<vec> as;
184
             size_t n = 0;
185
             for (auto&& x : fac)
186
             {
187
                 int64 mod, p, e;
188
                 vec a, b;
189
                 std::tie(mod, p, e) = x;
190
                 auto ss = s;
191
                 for (auto&& x : ss) x %= mod;
192
                 std::tie(a, b) = prime_power(ss, mod, p, e);
193
                 as.emplace_back(a);
194
                 n = std::max(n, a.size());
             }
195
196
             vec a(n), c(as.size()), m(as.size());
197
             for (size_t i = 0; i < n; ++i)</pre>
198
                 for (size_t j = 0; j < as.size(); ++j)</pre>
199
200
                 {
201
                      m[j] = std::get<0>(fac[j]);
202
                      c[j] = i < as[j].size() ? as[j][i] : 0;
203
204
                 a[i] = crt(c, m);
205
206
             return a;
207
         }
208
209
         LinearRecurrence(const vec& s, const vec& c, int64 mod) : init(s), trans(c), mod(mod), m(s.size
210
         LinearRecurrence(const vec& s, int64 mod, bool is_prime = true) : mod(mod)
211
212
             vec A;
213
             if (is_prime)
214
                 A = BerlekampMassey(s, mod);
215
             else
216
                 A = ReedsSloane(s, mod);
217
             if (A.empty()) A = {0};
218
             m = A.size() - 1;
219
             trans.resize(m);
220
             for (int i = 0; i < m; ++i)</pre>
221
             {
222
                 trans[i] = (mod - A[i + 1]) \% mod;
             }
223
224
             std::reverse(trans.begin(), trans.end());
225
             init = {s.begin(), s.begin() + m};
226
         }
227
         int64 calc(int64 n)
228
229
             if (mod == 1) return 0;
230
             if (n < m) return init[n];</pre>
231
             vec v(m), u(m << 1);
232
             int msk = !!n;
233
             for (int64 m = n; m > 1; m >>= 1) msk <<= 1;</pre>
```

```
234
             v[0] = 1 \% mod;
235
             for (int x = 0; msk; msk >>= 1, x <<= 1)</pre>
236
237
                 std::fill_n(u.begin(), m * 2, 0);
238
                 x |= !!(n & msk);
239
                 if (x < m)
240
                      u[x] = 1 \% mod;
241
                 else
242
                 { // can be optimized by fft/ntt
243
                      for (int i = 0; i < m; ++i)</pre>
244
245
                          for (int j = 0, t = i + (x & 1); j < m; ++j, ++t)
246
247
                              u[t] = (u[t] + v[i] * v[j]) % mod;
248
                      }
249
250
                      for (int i = m * 2 - 1; i >= m; --i)
251
252
                          for (int j = 0, t = i - m; j < m; ++j, ++t)
253
254
                              u[t] = (u[t] + trans[j] * u[i]) % mod;
255
                          }
256
                      }
257
258
                 v = \{u.begin(), u.begin() + m\};
             }
259
260
             int64 ret = 0;
261
             for (int i = 0; i < m; ++i)</pre>
262
             {
263
                 ret = (ret + v[i] * init[i]) % mod;
264
265
             return ret;
266
         }
267
268
         vec init, trans;
269
         int64 mod;
270
         int m;
271 };
     1.8.3 杜教筛
    #include <bits/stdc++.h>
     #include <tr1/unordered_map>
 3
 4
    using namespace std;
 5
     typedef long long 11;
 6
 7
     const int N = 5e6;
 8
 9
    bool vis[N + 1];
 10 int mu[N + 1], sumu[N + 1], prim[N + 1], cnt;
 11
 12 tr1::unordered_map<int, int> Smu;
13
 14 void get_mu(int n)
 15
    {
 16
         mu[1] = 1;
 17
         for(int i = 2; i <= n; i ++)</pre>
```

```
{
18
19
             if(!vis[i]) { prim[++ cnt] = i; mu[i] = -1; }
20
            for(int j = 1; j <= cnt && prim[j] * i <= n; j ++)</pre>
21
            {
22
                 vis[prim[j] * i] = 1;
23
                 if(i % prim[j] == 0) break;
24
                 else mu[i * prim[j]] = -mu[i];
25
            }
26
        }
27
        for(int i = 1; i <= n; i ++) sumu[i] = sumu[i - 1] + mu[i];</pre>
28 }
29
30 int phi[N + 1]; ll sumphi[N + 1];
31 tr1::unordered_map<11, 11> Sphi;
32
33 void get(int n)
34
35
        phi[1] = mu[1] = 1;
36
        for(int i = 2; i <= n; i ++)</pre>
37
38
            if(!vis[i])
39
            {
40
                 prim[++ cnt] = i;
41
                 mu[i] = -1; phi[i] = i - 1;
42
43
            for(int j = 1; j <= cnt && prim[j] * i <= n; j ++)</pre>
44
45
                 vis[i * prim[j]] = 1;
46
                 if(i % prim[j] == 0)
47
48
                     phi[i * prim[j]] = phi[i] * prim[j];
49
                     break;
50
                 }
                 else mu[i * prim[j]] = -mu[i], phi[i * prim[j]] = phi[i] * (prim[j] - 1);
51
            }
52
53
54
        }
55
        for(int i = 1; i <= n; i ++) sumu[i] = sumu[i - 1] + mu[i], sumphi[i] = sumphi[i - 1] + phi[i];</pre>
56
57
58 ll getSum_mu(int x)
59
60
        if(x <= N) return sumu[x];</pre>
        if(Smu[x]) return Smu[x];
61
62
        int ans = 1;
63
        for(int 1 = 2, r; 1 >= 0 && 1 <= x && r < 2147483647; 1 = r + 1)
64
65
             r = x / (x / 1);
66
             ans -= (r - 1 + 1) * getSum_mu(x / 1);
67
68
        return Smu[x] = ans;
69 }
70
71 ll getSum_phi(ll x)
72
73
        if(x <= N) return sumphi[x];</pre>
74
        if(Sphi[x]) return Sphi[x];
75
        11 ans = x * (x + 1) / 2;
76
        for(11 1 = 2, r; 1 \le x; 1 = r + 1)
```

```
77
        {
            r = x / (x / 1);
78
79
            ans -= (r - 1 + 1) * getSum_phi(x / 1);
80
        return Sphi[x] = ans;
81
82 }
    1.8.4 欧拉降幂
1 int prim[maxn], vis[maxn];
2 int tot, phi[maxn];
3
    struct node {
4
        ll res;
5
        bool v;
6
   };
7
8
    node qpow(11 A, 11 B, 11 C) {
9
        ll re = 1;
10
        bool flag = true;
        while (B) {
11
12
            if (B & 1) {
                if ((re *= A) >= C) flag = 0;
13
14
                re = re % C;
            }
15
16
            B = B >> 1;
17
            if (B) {
18
                if (A >= C) flag = 0;
19
                A %= C;
20
                if ((A *= A) >= C) flag = 0;
21
                A \%= C;
22
            }
23
        }
24
        return node{re, flag};
25 }
26
27 void init(int n) {
28
        phi[1] = 1;
        for (int i = 2; i <= n; i++) {</pre>
29
30
            if (!vis[i]) {
31
                prim[++tot] = i;
32
                phi[i] = i - 1;
            }
33
34
            for (int j = 1; j <= tot && prim[j] * i <= n; j++) {</pre>
35
                vis[i * prim[j]] = 1;
36
                if (i % prim[j] == 0) {
37
                    phi[i * prim[j]] = phi[i] * prim[j];
38
39
                } else phi[i * prim[j]] = phi[i] * (prim[j] - 1);
            }
40
41
42
        }
   }
43
44
45
46
   inline ll Euler(ll x) {
47
        if (x <= maxn) return phi[x];</pre>
48
        return 0;
49 }
```

```
50
    node f(ll a, ll k, ll p) {
51
52
        if (p == 1) return node{0, 0};
53
        if (k == 0) return node{a % p, a < p};</pre>
54
        11 ep = Euler(p);
55
        node tmp = f(a, k - 1, ep);
        if (_gcd(a, p) == 1)return qpow(a, tmp.res, p);
56
57
        if (!tmp.v) {
            tmp.res += ep;
58
        }
59
60
        return qpow(a, tmp.res, p);
61
    }
62
63
    int main() {
64
        ll a, k, p;
        init(1e7+2);
65
66
        int T;
67
        scanf("%d", &T);
        for (int kase = 1; kase <= T; ++kase) {</pre>
68
            scanf("%lld%lld%lld", &a, &k, &p);
69
70
            else printf("%lld\n", f(a, k, p).res);
        }
71
72
        return 0;
73 }
```

2 Graph Theory

2.1 路径

2.1.1 Dijkstra

```
const int maxn = 1e5 + 10;
   const int inf = 0x3f3f3f3f;
3
   int head[maxn], dis[maxn], cnt, n;
4
6
   struct Edge { int nex,to,w; }edge[20*maxn];
8
   void add(int u,int v,int w)
9
10
        edge[++cnt].nex=head[u];
11
        edge[cnt].w=w;
12
        edge[cnt].to=v;
13
        head[u]=cnt;
14
   }
15
16
   void dijkstra(int s)
17
18
        priority_queue<pair<int, int>, vector<pair<int, int> >, greater<pair<int, int> > > que;
19
        memset(dis, 0x3f, sizeof dis);
20
        que.push(\{0, s\}); dis[s] = 0;
21
        while(!que.empty())
22
23
            auto f = que.top(); que.pop();
24
            int u = f.second, d = f.first;
25
            if(d != dis[u]) continue;
            for(int i = head[u]; ~i; i = edge[i].nex)
27
28
                int v = edge[i].to, w = edge[i].w;
                if(dis[u] + w < dis[v])</pre>
29
30
31
                    dis[v] = dis[u] + w;
32
                    que.push({dis[v], v});
33
34
            }
35
        }
36 }
```

2.1.2 Euler Path

```
1 int S[N << 1], top;</pre>
    Edge edges[N << 1];</pre>
    set<int> G[N];
3
4
5
    void DFS(int u) {
6
        S[top++] = u;
7
        for (int eid: G[u])
8
9
             int v = edges[eid].get_other(u);
             G[u].erase(eid);
10
             G[v].erase(eid);
11
12
             DFS(v);
13
             return;
14
        }
```

```
15 }
16
   void fleury(int start)
17
18
    {
19
        int u = start;
20
        top = 0; path.clear();
21
        S[top++] = u;
22
        while (top)
23
24
            u = S[--top];
25
            if (!G[u].empty())
26
                DFS(u);
27
            else path.push_back(u);
28
        }
29 }
    2.1.3 K shortest Path(Astar)
1 const int inf = 0x3f3f3f3f;
2 const int maxn = 1000 + 10;
3 const int maxm = 100000 + 10;
4
5 int n, k, cnt, head[maxn], revhead[maxn], dis[maxn];
6 bool vis[maxn];
7
8
   struct node { int v, w, nex; } edge[maxm], revedge[maxm];
9
10 void init()
11
12
        cnt = 0;
13
        memset(head, Oxff, sizeof head);
14
        memset(revhead, 0xff, sizeof revhead);
15
   }
16
17 void add(int u, int v, int w)
18 {
19
        edge[cnt].v = v, revedge[cnt].v = u;
20
        edge[cnt].w = revedge[cnt].w = w;
21
        edge[cnt].nex = head[u];
22
        revedge[cnt].nex = revhead[v];
23
        head[u] = revhead[v] = cnt;
24
        cnt++;
25 }
26
27
   void spfa(int src)
                            //建立反向图, 求图中所有点到终点的最短路径
28
    {
29
        for (int i = 1; i <= n; i++) dis[i] = inf;</pre>
30
        memset(vis, false, sizeof vis);
31
        vis[src] = 0;
32
        queue<int> que;
33
        que.push(src);
34
        dis[src] = 0;
35
        while (!que.empty())
36
37
            int u = que.front();
38
            que.pop();
39
            vis[u] = false;
            for (int i = revhead[u]; ~i; i = revedge[i].nex)
40
```

```
{
41
42
               int v = revedge[i].v, w = revedge[i].w;
               if (dis[v] > dis[u] + w)
43
44
               {
                   dis[v] = dis[u] + w;
45
                   if (!vis[v])
46
47
                   {
48
                       que.push(v);
                       vis[v] = true;
49
50
                   }
51
               }
52
           }
       }
53
   }
54
55
56
   struct A
57
58
        int f, g, h;
                     //f(n),g(n),h(n)函数
59
        int id;
                      //当前点的编号
60
       bool operator<(const A a) const</pre>
61
                //定义比较函数
62
           if (a.f == f) return a.g < g;</pre>
63
           return a.f < f;</pre>
64
       }
65 };
66
   int Astar(int src, int des)
67
68
   {
        int cnt = 0;
69
70
       priority_queue <A> Q;
71
       if (src == des) k++; //如果起点即为终点
72
       if (dis[src] == inf) return -1; //如果起点不能到达终点
73
       A st, now, tmp;
74
       st.id = src, st.g = 0, st.f = st.g + dis[src]; //定义起始节点
75
       Q.push(st);
76
       while (!Q.empty())
77
        {
78
           now = Q.top();
79
           Q.pop();
80
           if (now.id == des)
                                 //如果当前节点为终点
81
           {
82
83
               if (cnt == k) return now.g;
                                             //找到第k短路
84
85
           for (int i = head[now.id]; ~i; i = edge[i].nex)
86
           {
87
               tmp.id = edge[i].v;
                                             //到该点的实际花费
88
               tmp.g = now.g + edge[i].w;
89
               tmp.f = tmp.g + dis[tmp.id];
                                            //到最终状态的估计花费
90
               Q.push(tmp);
91
           }
92
       }
93
       return -1; //路径总数小于k
94 }
95
96
   int main()
97
   {
98
        int m, s, t, u, v, w;
99
       while (scanf("%d%d", &n, &m) != EOF)
```

```
100
         {
101
             init();
102
             while (m--)
103
             {
104
                 scanf("%d%d%d", &u, &v, &w);
                 add(u, v, w);
105
106
             }
107
             scanf("%d%d%d", &s, &t, &k);
108
                        //求所有点到终点的最短路
             spfa(t);
109
             printf("%d\n", Astar(s, t));
110
111
        return 0;
112 }
    2.1.4 K shortest Path(可持久化可并堆)
    #include <bits/stdc++.h>
    #include<ext/pb_ds/priority_queue.hpp>
 3
 4 using namespace std;
 5
 6 const int N = '';
 7 const int M = '';
 8 const int logM = 20;
 9 const int inf = 0x3f3f3f3f;
10
11 int n, m, k, S, T;
12
13 struct Edge{ int nex, to, w; };
14
15 struct Graph
16 {
17
         int head[N], cnt;
18
         Edge edge[M];
19
         void init(int n) { for(int i = 0; i <= n; i ++) head[i] = 0; cnt = 0; }</pre>
         void addedge(int u, int v, int val) { edge[++ cnt].nex = head[u], edge[cnt].to = v, edge[cnt].w
20
          = val, head[u] = cnt; }
21 }g, rg;
22
23 int dis[N];
24
25
    void dijkstra()
26
    {
27
         priority_queue<pair<int, int>, vector<pair<int, int> >, greater<pair<int, int> > > que;
28
         memset(dis, inf, sizeof dis);
29
         que.push(\{0, T\}); dis[T] = 0;
30
         const int *head = rg.head; const Edge *edge = rg.edge;
31
         while(!que.empty())
32
         {
33
             auto f = que.top(); que.pop();
34
             int u = f.second, d = f.first;
35
             if(d != dis[u]) continue;
```

 $if(dis[u] + w < dis[v]) { dis[v] = dis[u] + w; que.push({dis[v], v}); }$

for(int i = head[u]; i; i = edge[i].nex)

int v = edge[i].to, w = edge[i].w;

36

37 38

39

40

41

}

}

```
42 }
43
44
    bool tree_edge[M], vis[N];
45
    int fa[N], st[N], top;
46
47
    void dfs(int u)
48
    {
49
         vis[u] = true;
50
         st[++ top] = u;
51
         for(int i = rg.head[u]; i; i = rg.edge[i].nex)
52
53
             int v = rg.edge[i].to;
54
             if(!vis[v] && dis[v] == dis[u] + rg.edge[i].w)
55
             {
56
                 fa[v] = u;
57
                 tree_edge[i] = true;
58
                 dfs(v);
59
             }
60
         }
61 }
62
63 namespace LT
64 {
         int son[M * logM][2];
65
66
         int ht[M * logM], val[M * logM], id[M * logM];
67
         int tot;
68
69
         int newnode(int _val, int _id, int _dis = 0)
70
71
             int now = ++ tot;
72
             val[now] = _val, id[now] = _id;
73
             ht[now] = _dis, son[now][0] = son[now][1] = 0;
74
             return now;
75
         }
76
77
         int _copy(int ori)
78
79
             int now = ++tot;
80
             val[now] = val[ori], id[now] = id[ori];
             ht[now] = ht[ori], son[now][0] = son[ori][0], son[now][1] = son[ori][1];
81
82
             return now;
         }
83
84
85
         int merge(int a, int b)
86
87
             if(!a || !b) return a | b;
             if(val[a] > val[b]) swap(a, b);
88
89
             int now = _copy(a);
             son[now][1] = merge(son[now][1], b);
90
91
             if(ht[son[now][0]] < ht[son[now][1]]) swap(son[now][0], son[now][1]);</pre>
92
             ht[now] = ht[son[now][1]] + 1;
93
             return now;
94
         }
95
         void insert(int &rt, int val, int id) { rt = merge(newnode(val, id), rt); }
96
97
    }
98
99
    int rt[M];
100
```

```
101 void build_heap()
102
    {
103
         for(int i = 1; i <= top; i ++)</pre>
104
         {
105
             int u = st[i];
106
             rt[u] = rt[fa[u]];
107
             for(int i = g.head[u]; i; i = g.edge[i].nex)
108
             {
109
                 int v = g.edge[i].to;
110
                 if(!tree_edge[i] && dis[v] != inf) LT::insert(rt[u], dis[v] - dis[u] + g.edge[i].w, v);
             }
111
112
         }
113
    }
114
115
    int solve(int k)
116
117
         if(k == 1) return dis[S];
         __gnu_pbds::priority_queue<pair<int, int>, greater<pair<int, int> > que;
118
119
         que.push({dis[S] + LT::val[rt[S]], rt[S]});
120
         while(!que.empty())
121
         {
             pair<int, int> f = que.top(); que.pop();
122
123
             if((--k) == 1) return f.first;
             int v = f.first, u = f.second;
124
125
             int lc = LT::son[u][0], rc = LT::son[u][1], o = LT::id[u];
126
             if(rt[o]) que.push({v + LT::val[rt[o]], rt[o]});
127
             if(lc) que.push({v + LT::val[lc] - LT::val[u], lc});
128
             if(rc) que.push({v + LT::val[rc] - LT::val[u], rc});
         }
129
130
         return -1;
131
132
133
    void init()
134
    {
135
         g.init(n), rg.init(n);
136
         memset(rt, 0, sizeof rt);
137
         memset(tree_edge, 0, sizeof tree_edge);
138
         top = LT::tot = 0;
139
    }
140
141
    void getans()
142
143
         //input S-T
144
         init();
145
         dijkstra();
146
         dfs(T);
147
         build_heap();
148
         cout << solve(k);</pre>
149 }
     2.2
           生成树
     2.2.1 Kruskal
 1
    const int maxn = 1e5 + 10;
 2
 3 int n, m, pre[maxn];
 4 struct edge {int u, v, w; } es[maxn];
 5 int Find(int x) { return x == pre[x] ? x : pre[x] = Find(pre[x]); }
```

```
bool cmp(const edge &x, const edge &y) { return x.cost < y.cost; }</pre>
7
8
    int kruskal()
9
    {
10
        sort(es, es + m, cmp);
        int res = 0;
11
12
        for(int i = 0; i < m; i ++)</pre>
13
14
            int fx = Find(es[i].u), fy = Find(es[i].v);
15
            if(fx != fy) pre[fx] = fy, res += es[i].cost;
16
17
        return res;
18 }
    2.2.2 Prim
1 const int maxn = 1000 + 10;
   const int inf = 0x3f3f3f3f;
3
4 int n, mp[maxn][maxn], cost[maxn];
5 bool vis[maxn];
6
7
   int prim()
8
9
        for(int i = 0; i < n; i ++) cost[u] = inf, vis[u] = false;</pre>
10
        int res = 0; cost[0] = 0;
11
        for(;;)
12
        {
13
            int v = -1;
            for(int u = 0; u < n; u ++)
                if(!vis[u] && (v == -1 || cost[u] < cost[v])) v = u;</pre>
15
16
            if(v == -1) break;
            res += cost[v];
17
            vis[v] = true;
18
19
            for(int u = 0; u < n; u ++) cost[u] = min(cost[u], mp[v][u]);</pre>
20
        }
21
        return res;
22 }
    2.2.3 最小树形图
   const int INF = 0x3f3f3f3f;
2
    const int maxn = 10000;
3
    const int maxm = 10000;
4
5
   struct Edge{int u,v,cost; } edge[maxm];
6
7
   int pre[maxn], id[maxn], vis[maxn], in[maxn];
8
9
   int zhuliu(int root, int n, int m)
10 {
11
        int res=0, u, v;
12
        for(;;)
13
14
            for(int i=0; i<n; i++) in[i] = INF;</pre>
15
            for(int i=0; i<m; i++) if(edge[i].u != edge[i].v && edge[i].cost < in[edge[i].v])</pre>
16
                pre[edge[i].v] = edge[i].u;
17
```

```
18
                in[edge[i].v] = edge[i].cost;
            }
19
            for(int i=0; i<n; i++) if(i != root && in[i] ==INF) return -1;</pre>
20
21
            int tn=0;
22
            memset(id, 0xff, sizeof id);
23
            memset(vis, 0xff, sizeof vis);
24
            in[root] = 0;
25
            for(int i=0; i<n;i++)</pre>
26
            {
27
                res += in[i];
28
                v = i;
29
                while( vis[v] != i && id[v] == -1 && v!= root) vis[v] = i, v = pre[v];
30
                if(v != root && id[v] == -1)
31
32
                     for(int u = pre[v]; u != v; u = pre[u]) id[u] = tn;
33
                     id[v] = tn++;
                }
34
35
            }
36
            if(tn == 0) break;
37
            for(int i=0; i<n; i++) if(id[i] == -1) id[i] = tn++;</pre>
38
            for(int i=0; i<m; )</pre>
39
            {
40
                v = edge[i].v;
41
                edge[i].u = id[edge[i].u];
42
                edge[i].v = id[edge[i].v];
43
                if(edge[i].u != edge[i].v) edge[i++].cost -= in[v];
44
                else swap(edge[i], edge[--m]);
            }
45
46
            n = tn;
47
            root = id[root];
48
        }
49
        return res;
50 }
    2.2.4 Matrix Tree
1 const int N = 305;
    const int mod = 1e9 + 7;
3
4
    int n, m, a[N][N];
5
6
    int Gauss(int n) {
7
        int ans = 1;
8
        for (int i = 1; i <= n; i++) {</pre>
9
            for (int k = i + 1; k <= n; k++) {</pre>
10
                while (a[k][i]) {
11
                     int d = a[i][i] / a[k][i];
12
                     for (int j = i; j <= n; j++) {</pre>
13
                         a[i][j] = (a[i][j] - 1LL * d * a[k][j] % mod + mod) % mod;
14
15
                     std::swap(a[i], a[k]);
                     ans = -ans;
16
17
                }
18
            }
19
            ans = 1LL * ans * a[i][i] % mod;
20
21
        return (ans % mod + mod) % mod;
22 }
```

```
23
   int main() {
24
        scanf("%d%d", &n, &m);
25
        for (int i = 1; i <= m; i++) {</pre>
26
            int u, v;
27
           scanf("%d%d", &u, &v);
28
           a[u][v]--, a[v][u]--;
29
           a[u][u]++, a[v][v]++;
30
        }
31
       printf("%d\n", Gauss(n - 1));
32
       return 0;
33 }
         连通性
   2.3
    2.3.1 割点
1 const int maxn = 1e4 + 10;
3 vector<int> edge[maxn];
4 int n, dfn[maxn], low[maxn], cnt = 0;
5 bool vis[maxn], cut[maxn];
6
7
   void Tarjan(int u, int fa)
8
9
        dfn[u] = low[u] = ++cnt;
10
        vis[u] = true;
11
        int children = 0;
       for (int i = 0; i < edge[u].size(); i++)</pre>
12
13
14
           int v = edge[u][i];
           if (v != fa && vis[v])
15
               low[u] = min(low[u], dfn[v]);
16
17
           else if (!vis[v])
18
            {
19
               Tarjan(v, u);
20
               children++;
21
               low[u] = min(low[u], low[v]);
               if (fa == -1 && children > 1) //若u是根节点且子节点数大于1
22
23
                    cut[u] = true;
                                    //u是割点
               else if (fa != -1 && low[v] >= dfn[u])
                                                         //若u不是根节点且v不能访问到u的父节点
24
25
                   cut[u] = true;
                                     //u是割点
26
           }
27
        }
28 }
   2.3.2 桥
1 const int maxn = 1e4 + 10;
3 vector<int> edge[maxn];
4 int n, dfn[maxn], low[maxn], father[maxn], cnt = 0;
5 bool bridge[maxn][maxn];
6
7
   void Tarjan(int u, int fa)
8
9
        dfn[u] = low[u] = ++cnt;
10
        for (int i = 0; i < edge[u].size(); i++)</pre>
11
        {
```

```
12
           int v = edge[u][i];
                          //未访问节点v
13
           if (!dfn[v])
14
15
               Tarjan(v, u);
16
               low[u] = min(low[u], low[v]);
17
               if (low[v] > dfn[u]) //节点v到达祖先必须经过(u,v)
                  bridge[u][v] = bridge[v][u] = true;
18
                                                       //(u,v)是桥
19
           }
20
                               //u的父节点不是v, (u,v)不存在重边
           else if (fa != v)
               low[u] = min(low[u], dfn[v]);
21
22
       }
23
   }
```

2.3.3 强连通分量

```
const int maxn=1000+10;
 2
 3
    vector<int> edge[maxn];
 4
 5 int dfn[maxn], low[maxn];
   int stack[maxn], index, tot;
 7
    int belong[maxn], inde[maxn], outde[maxn], scc;
    bool vis[maxn];
 9
10
   void add(int u, int v)
11
12
        edge[u].push_back(v);
13
        edge[v].push_back(u);
14
    }
15
16 void Tarjan(int u)
17
    {
18
        dfn[u] = low[u] = ++tot;
19
        stack[++index] = u;
20
        vis[u] = true;
21
        int v;
22
        for(int i = 0;i < edge[u].size(); i++)</pre>
23
24
            v=edge[u][i];
25
            if(!dfn[v])
26
            {
27
                Tarjan(v);
28
                low[u] = min(low[v], low[u]);
29
30
            else if(vis[v]) low[u] = min(low[v], dfn[u]);
31
        }
32
        if(dfn[u] == low[u])
33
34
            scc++;
35
            do
36
            {
37
                v = stack[index--];
38
                vis[v] = false;
39
                belong[v] = scc;
40
            }while(v != u);
41
        }
42 }
```

2.3.4 点双联通分量

```
1
    const int maxn = 10000 + 10;
 2
 3
    struct Edge{ int u, v; };
    vector<int> G[maxn], bcc[maxn];
 4
 5
   int dfn[maxn], low[maxn], bccno[maxn], idx, bcc_cnt, bridge;
 6
 7
    bool iscut[maxn];
 8
 9
    stack<Edge> st;
10
11 void dfs(int u, int pre)
12 {
13
        dfn[u] = low[u] = ++idx;
14
        int child = 0;
15
        for(auto v : G[u])
16
17
            if(v == pre) continue;
            if(!dfn[v])
18
19
            {
20
                child ++;
21
                st.push({u, v});
22
                dfs(v, u);
23
                low[u] = min(low[u], low[v]);
                if(low[v] >= dfn[u])
24
25
26
                     iscut[u] = true;
27
                     bcc[++bcc_cnt].clear();
28
                     Edge x;
29
                     do
30
                     {
31
                         x = st.top(); st.pop();
32
                         if(bccno[x.u] != bcc_cnt) { bcc[bcc_cnt].push_back(x.u); bccno[x.u] = bcc_cnt;
        }
33
                         if(bccno[x.v] != bcc_cnt) { bcc[bcc_cnt].push_back(x.v); bccno[x.v] = bcc_cnt;
        }
34
                     } while(x.u != u || x.v != v);
35
36
                if(low[v] > dfn[u]) ++ bridge;
37
            }
38
            else if(dfn[v] < dfn[u])</pre>
39
40
                st.push({u, v});
41
                low[u] = min(low[u], dfn[v]);
42
43
44
        if(pre < 0 && child == 1) iscut[u] = 0;</pre>
    }
45
46
47
   void find_bcc(int n)
48 {
49
        memset(dfn, 0, sizeof dfn);
50
        memset(iscut, 0, sizeof iscut);
51
        memset(bccno, 0, sizeof bccno);
52
        for(int i = 1; i <= bcc_cnt; i ++) bcc[i].clear();</pre>
53
        idx = bcc_cnt = bridge = 0;
54
        for(int i = 0; i < n; i ++) if(!dfn[i]) dfs(i, -1);</pre>
55 }
```

2.3.5 边双联通分量

```
1 const int maxn = 10000 + 10;
2
3 int low[maxn], dfn[maxn], head[maxn], cnt, idx;
4 int cutEdge[maxn << 2];</pre>
5 struct Edge { int nex, v; }edge[maxn << 2];</pre>
6
7
   void add(int u, int v) { edge[cnt].nex = head[u], edge[cnt].v = v, head[u] = cnt ++; }
8
9
   void dfs(int u, int pre)
10
11
        low[u] = dfn[u] = ++idx;
12
        for(int i = head[u]; ~i; i = edge[i].nex)
13
            int v = edge[i].v;
14
            if(v == pre) continue;
15
            if(!dfn[v])
16
17
18
                dfs(v, u);
                low[u] = min(low[u], low[v]);
19
20
                if(low[v] > dfn[u]) cutEdge[i] = cutEdge[i ^ 1] = 1;
21
            else if(dfn[v] < dfn[u]) low[u] = min(low[u], dfn[v]);</pre>
22
23
        }
24 }
```

2.4 二分图匹配

- 1. 二分图中的最大匹配数 = 最小点覆盖数
- 2. 最小路径覆盖 = 最小路径覆盖 = | G | 最大匹配数
- 3. 二分图最大独立集 = 顶点数-最小点覆盖
- 4. 二分图的最大团 = 补图的最大独立集

2.4.1 Hungary Algorithm

```
const int maxn = 150;
2
3 int n;
4 int edge[maxn][maxn];
   int linker[maxn];
    bool vis[maxn];
8
    bool path(int u)
9
    {
10
        for (int v = 1; v <= n; v++)</pre>
11
12
            if (edge[u][v] && !vis[v])
13
14
                vis[v] = true;
                if (linker[v] == -1 || path(linker[v]))
15
16
17
                     linker[v] = u;
18
                     return true;
19
                }
20
            }
21
        }
```

```
22
        return false;
23 }
24
25 int hungary()
26 {
27
        int res = 0;
28
        memset(linker, 0xff, sizeof(linker));
29
        for (int i = 1; i <= n; i++)</pre>
30
31
            memset(vis, false, sizeof(vis));
32
            res += path(i);
33
        }
34
        return res;
35 }
```

2.4.2 Hopcroft-karp Algorithm

```
1 //复杂度O(n^0.5*m),注意这个板子的下标是从O开始的
3 const int MAXN = 3010;//左边节点数量、右边节点数量
4 const int MAXM = 3010 * 3010;//边的数量
5 const int INF = 0x3f3f3f3f;
6
7 struct Edge
8
9
       int v;
10
       int next;
11 } edge[MAXM];
12
13 int nx, ny;
14 int cnt;
15 int dis;
16
17 int first[MAXN];
18 int xlink[MAXN], ylink[MAXN];
19 /*xlink[i]表示左集合顶点所匹配的右集合顶点序号, ylink[i]表示右集合i顶点匹配到的左集合顶点序号。*/
20 int dx[MAXN], dy[MAXN];
21 /*dx[i]表示左集合i顶点的距离编号, dy[i]表示右集合i顶点的距离编号*/
22 int vis[MAXN]; //寻找增广路的标记数组
23
24 void init()
25 {
26
       cnt = 0;
27
       memset(first, -1, sizeof(first));
28
       memset(xlink, -1, sizeof(xlink));
29
       memset(ylink, -1, sizeof(ylink));
30 }
31
32 void read_graph(int u, int v)
33 {
34
       edge[cnt].v = v;
35
       edge[cnt].next = first[u], first[u] = cnt++;
36 }
37
38 int bfs()
39
40
       queue<int> q;
       dis = INF;
41
```

```
42
         memset(dx, -1, sizeof(dx));
 43
         memset(dy, -1, sizeof(dy));
         for (int i = 0; i < nx; i++)</pre>
44
 45
         {
 46
             if (xlink[i] == -1)
47
             {
 48
                 q.push(i);
 49
                 dx[i] = 0;
50
51
         }
52
         while (!q.empty())
53
54
             int u = q.front();
55
             q.pop();
56
             if (dx[u] > dis) break;
             for (int e = first[u]; e != -1; e = edge[e].next)
57
 58
59
                 int v = edge[e].v;
60
                 if (dy[v] == -1)
61
62
                     dy[v] = dx[u] + 1;
63
                      if (ylink[v] == -1) dis = dy[v];
64
                      else
65
                      {
66
                          dx[ylink[v]] = dy[v] + 1;
67
                          q.push(ylink[v]);
68
                     }
 69
                 }
             }
 70
 71
 72
         return dis != INF;
73 }
74
75
    int find(int u)
76
77
         for (int e = first[u]; e != -1; e = edge[e].next)
78
79
             int v = edge[e].v;
             if (!vis[v] && dy[v] == dx[u] + 1)
80
81
82
                 vis[v] = 1;
                 if (ylink[v] != -1 && dy[v] == dis) continue;
 83
 84
                 if (ylink[v] == -1 || find(ylink[v]))
 85
 86
                     xlink[u] = v, ylink[v] = u;
 87
                     return 1;
 88
                 }
 89
             }
         }
90
 91
         return 0;
92 }
93
    int MaxMatch()
94
     {
95
         int ans = 0;
96
97
         while (bfs())
98
99
             memset(vis, 0, sizeof(vis));
100
             for (int i = 0; i < nx; i++)</pre>
```

```
101
               if (xlink[i] == -1)
102
                   ans += find(i);
        }
103
104
        return ans;
105
    2.4.3 二分图多重匹配
 1 const int maxn = 1e2 + 5;//左边最大点数
 2 const int maxm = 1e2 + 5;//右边最大点数
 3 int graph[maxn][maxm], vis[maxm];//图G和增广路访问标记
 4 int match[maxm] [maxn];//左边元素与右边元素第n次匹配
 5 int nx, ny, m;//左边点数, 右边点数,边数
   int vol[maxm];//右边点多重匹配可容纳值
    int cnt[maxm];//右边点已匹配值
 7
 9
    bool find_path(int u)//找增广路
 10
    {
        for (int i = 0; i < ny; i++)//注意, 这里节点是从0开始编号, 题目有时是从1开始编号
11
12
13
           if (graph[u][i] && !vis[i])//不在增广路
           {
14
               vis[i] = 1;//放进增广路
15
16
               if (cnt[i] < vol[i])//如果当前已匹配数量小于可容纳量,则直接匹配
17
               {
18
                   match[i][cnt[i]++] = u;
19
                   return true;
20
               }
21
               for (int j = 0; j < cnt[i]; j++)</pre>
22
23
                   if (find_path(match[i][j]))//如果先前已匹配右边的点能另外找到增广路,则此点仍可匹配
24
25
                      match[i][j] = u;
26
                      return true;
27
                   }
28
               }
29
           }
30
31
        return false;
32 }
33
34
    int max_match()//计算多重匹配的最大匹配数
35
    {
36
        int res = 0;
37
        memset(match, -1, sizeof(match));
38
        memset(cnt, 0, sizeof(cnt));
        for (int i = 0; i < nx; i++)</pre>
39
40
41
           memset(vis, 0, sizeof(vis));
42
           if (find_path(i)) res++;
43
        }
        return res;
44
    }
45
46
47
    bool all_match()//判断左边的点是否都与右边的点匹配了
48
49
        memset(cnt, 0, sizeof(cnt));
        for (int i = 0; i < nx; i++)</pre>
50
```

```
{
51
            memset(vis, 0, sizeof(vis));
52
            if (!find_path(i)) return false;
53
        }
54
55
        return true;
56
   }
    2.4.4 二分图最大权匹配 (KM 算法)
   const int maxn=1000+10;
    const int inf=0x3f3f3f3f;
3
4 int n;
5 int lx[maxn],ly[maxn],edge[maxn][maxn];
   int match[maxn],delta;
6
7
    bool vx[maxn],vy[maxn];
8
   bool dfs(int x) //DFS增广, 寻找相等子图的完备匹配
9
10
    {
11
        vx[x]=true;
12
        for(int y=1;y<=n;y++)</pre>
13
        {
            if(!vy[y])
14
15
            {
16
                int tmp=lx[x]+ly[y]-edge[x][y];
17
                if(!tmp)
                            //edge(x,y)为可行边
18
                {
19
                    vy[y]=true;
20
                    if(!match[y]||dfs(match[y]))
21
22
                        match[y]=x;
23
                        return true;
24
25
                }
26
                else delta=min(delta,tmp);
27
            }
28
        }
29
        return false;
30 }
31
32
   void KM()
33
    {
34
        for(int i=1;i<=n;i++)</pre>
                               //初始化可行顶标的值
35
36
            lx[i]=-inf;
37
            ly[i]=0;
38
            for(int j=1;j<=n;j++)</pre>
39
                lx[i]=max(lx[i],edge[i][j]);
        }
40
        memset(match,0,sizeof(match));
41
42
        for(int x=1;x<=n;x++)</pre>
43
        {
44
            for(;;)
45
            {
46
                delta=inf;
47
                memset(vx,0,sizeof(vx));
48
                memset(vy,0,sizeof(vy));
49
                if(dfs(x)) break;
```

```
50
               for(int i=1;i<=n;i++) //修改顶标
51
52
                   if(vx[i]) lx[i]-=delta;
53
                   if(vy[i]) ly[i]+=delta;
54
55
           }
56
       }
57 }
    2.4.5 一般图匹配带花树
1 //一般图匹配, 带花树算法
   const int maxn = 1000 + 10;
3
4
   vector<int> edge[maxn];
5
   queue<int> que;
   int n, pre[maxn], type[maxn], link[maxn], nex[maxn], vis[maxn];
8
9 void add(int u, int v)
10 {
        edge[u].push_back(v);
11
        edge[v].push_back(u);
12
13
14
15 int Find(int x)
16
17
       return x == pre[x] ? x : pre[x] = Find(pre[x]);
18
19
20 void combine(int x, int lca)
                                  //如果找到奇环,对当前点x和找到的
21 {
22
       while (x != lca)
23
24
           int u = link[x], v = nex[u];
           if (Find(v) != lca) nex[v] = u;
25
26
           if (type[u] == 1) type[u] = 2, que.push(u);
27
           pre[Find(x)] = Find(u);
28
           pre[Find(u)] = Find(v);
29
           x = v;
30
       }
31 }
32
33
   void contrack(int x, int y)
34
   {
35
       int lca = x;
36
       memset(vis, 0, sizeof(vis));
37
       for (int i = x; i; i = nex[link[i]])
38
        {
39
            i = Find(i);
40
           vis[i] = 1;
41
       }
42
       for (int i = y; i; i = nex[link[i]])
43
44
           i = Find(i);
45
           if (vis[i])
46
           {
47
               lca = i;
```

```
48
                  break;
             }
 49
         }
 50
 51
         if (lca != Find(x)) nex[x] = y;
 52
         if (lca != Find(y)) nex[y] = x;
53
         combine(x, lca);
54
         combine(y, lca);
    }
55
56
 57
    void bfs(int s)
58
59
         memset(type, 0, sizeof(type));
60
         memset(nex, 0, sizeof(nex));
         for (int i = 1; i <= n; i++) pre[i] = i;</pre>
61
62
         while (!que.empty()) que.pop();
63
         que.push(s);
 64
         type[s] = 2;
 65
         while (!que.empty())
 66
 67
             int x = que.front();
 68
             que.pop();
 69
             for (int i = 0; i < edge[x].size(); i++)</pre>
 70
 71
                  int y = edge[x][i];
 72
                  if (Find(x) == Find(y) || link[x] == y || type[y] == 1) continue;
                  if (type[y] == 2) contrack(x, y);
 73
 74
                  else if (link[y])
 75
                      nex[y] = x;
 76
 77
                      type[y] = 1;
 78
                      type[link[y]] = 2;
 79
                      que.push(link[y]);
 80
                  } else
81
 82
                      nex[y] = x;
 83
                      int pos = y, u = nex[pos], v = link[u];
 84
                      while (pos)
 85
                      {
 86
                          link[pos] = u;
 87
                          link[u] = pos;
 88
                          pos = v;
 89
                          u = nex[pos];
 90
                          v = link[u];
 91
                      }
 92
                      return;
 93
                  }
 94
             }
95
         }
96 }
97
98 int maxmatch()
99
    {
100
         for (int i = 1; i <= n; i++) if (!link[i]) bfs(i);</pre>
101
         int ans = 0;
102
         for (int i = 1; i <= n; i++) if (link[i]) ans++;</pre>
103
         return ans / 2;
104
105
106 void init()
```

```
107 {
         for (int i = 1; i <= n; i++) edge[i].clear();</pre>
108
109
        memset(link, 0, sizeof(link));
110 }
           网络流
    2.5
    2.5.1 Dinic
    const int MAX_V = 1000 + 10;
    const int INF = 0x3f3f3f3f;
 3
 4 //用于表示边的结构体 (终点,流量,反向边)
 5 struct edge{int to, cap, rev;};
 6
 7
    vector<edge> G[MAX_V]; //图的邻接表表示
    int level[MAX_V]; //顶点到源点的距离标号
 9
    int iter[MAX_V];
                        //当前弧
 10
    void add(int from, int to, int cap)
11
12
    {
13
         G[from].push_back((edge){to, cap, (int)G[to].size()});
         G[to].push_back((edge){from, 0, (int)G[from].size() - 1});
14
    }
15
16
   //计算从源点出发的距离标号
17
   void bfs(int s)
18
19
20
        memset(level, -1, sizeof(level));
21
         queue<int> que;
22
         level[s] = 0;
23
         que.push(s);
24
        while(!que.empty())
25
26
            int v = que.front(); que.pop();
27
            for(int i = 0; i < G[v].size(); i++)</pre>
28
29
                edge &e = G[v][i];
30
                if(e.cap > 0 && level[e.to] < 0)</pre>
31
                {
32
                    level[e.to] = level[v] + 1;
33
                    que.push(e.to);
34
35
            }
36
         }
37
    }
38
    //通过DFS寻找增广路
39
40
    int dfs(int v, int t, int f)
41
42
         if(v == t) return f;
43
         for(int &i = iter[v]; i<G[v].size(); i++)</pre>
44
45
             edge &e = G[v][i];
46
            if(e.cap > 0 && level[v] < level[e.to])</pre>
47
48
                int d = dfs(e.to, t, min(f, e.cap));
49
                if(d > 0)
50
                {
```

```
51
                     e.cap -= d;
                    G[e.to][e.rev].cap += d;
52
53
                     return d;
54
                }
55
            }
56
        }
57
        return 0;
    }
58
59
   //求解从s到t的最大流
   int max_flow(int s, int t)
61
62
        int flow = 0;
63
64
        for(;;)
65
66
            bfs(s);
67
            if(level[t] < 0) return flow;</pre>
68
            memset(iter, 0, sizeof(iter));
69
            int f;
70
            while((f = dfs(s,t,INF)) > 0) flow += f;
71
        }
72 }
    2.5.2 ISAP
1 struct Edge {
2
      int from, to, cap, flow;
3
      Edge(int u, int v, int c, int f) : from(u), to(v), cap(c), flow(f) {}
4 };
5
6
   bool operator<(const Edge& a, const Edge& b) {</pre>
7
     return a.from < b.from || (a.from == b.from && a.to < b.to);</pre>
8
9
10 struct ISAP {
11
      int n, m, s, t;
12
      vector<Edge> edges;
13
      vector<int> G[maxn];
14
      bool vis[maxn];
15
      int d[maxn];
16
      int cur[maxn];
17
      int p[maxn];
18
      int num[maxn];
19
20
      void AddEdge(int from, int to, int cap) {
21
        edges.push_back(Edge(from, to, cap, 0));
22
        edges.push_back(Edge(to, from, 0, 0));
23
        m = edges.size();
24
        G[from].push_back(m - 2);
25
        G[to].push_back(m - 1);
26
27
28
      bool BFS() {
29
        memset(vis, 0, sizeof(vis));
30
        queue<int> Q;
31
        Q.push(t);
32
        vis[t] = 1;
33
        d[t] = 0;
```

```
34
        while (!Q.empty()) {
35
          int x = Q.front();
36
          Q.pop();
37
          for (int i = 0; i < G[x].size(); i++) {</pre>
             Edge& e = edges[G[x][i] ^ 1];
38
39
             if (!vis[e.from] && e.cap > e.flow) {
40
               vis[e.from] = 1;
41
               d[e.from] = d[x] + 1;
               Q.push(e.from);
42
43
44
          }
45
        }
46
        return vis[s];
      }
47
48
49
      void init(int n) {
50
        this -> n = n;
        for (int i = 0; i < n; i++) G[i].clear();</pre>
51
52
        edges.clear();
53
      }
54
      int Augment() {
55
56
        int x = t, a = INF;
57
        while (x != s) {
58
          Edge& e = edges[p[x]];
59
          a = min(a, e.cap - e.flow);
60
          x = edges[p[x]].from;
        }
61
62
        x = t;
63
        while (x != s) {
64
          edges[p[x]].flow += a;
          edges[p[x] ^ 1].flow -= a;
65
66
          x = edges[p[x]].from;
        }
67
68
        return a;
69
70
71
      int Maxflow(int s, int t) {
72
        this \rightarrow s = s;
73
        this->t = t;
74
        int flow = 0;
75
        BFS();
76
        memset(num, 0, sizeof(num));
77
        for (int i = 0; i < n; i++) num[d[i]]++;</pre>
78
        int x = s;
        memset(cur, 0, sizeof(cur));
79
        while (d[s] < n) {
80
          if (x == t) {
81
82
             flow += Augment();
83
             x = s;
84
85
          int ok = 0;
86
          for (int i = cur[x]; i < G[x].size(); i++) {</pre>
87
             Edge& e = edges[G[x][i]];
             if (e.cap > e.flow && d[x] == d[e.to] + 1) {
88
89
               ok = 1;
90
               p[e.to] = G[x][i];
91
               cur[x] = i;
92
               x = e.to;
```

```
93
               break;
94
             }
           }
95
 96
           if (!ok) {
97
             int m = n - 1;
98
             for (int i = 0; i < G[x].size(); i++) {</pre>
99
               Edge& e = edges[G[x][i]];
100
               if (e.cap > e.flow) m = min(m, d[e.to]);
101
102
             if (--num[d[x]] == 0) break;
103
             num[d[x] = m + 1]++;
             cur[x] = 0;
104
105
             if (x != s) x = edges[p[x]].from;
106
         }
107
108
         return flow;
109
110 };
     2.5.3 MCMF
 1 const int maxn = 10000 + 10;
     const int inf = 0x3f3f3f3f;
 3
 4 struct Edge { int from, to, cap, flow, cost; };
 5
 6
    struct MCMF
 7
    {
 8
         int n, m;
 9
         vector<Edge> edges;
 10
         vector<int> G[maxn];
 11
         bool inq[maxn];
 12
         int dis[maxn], path[maxn], a[maxn];
 13
         void init(int n)
 14
 15
 16
             this->n = n;
 17
             for(int i = 0;i <= n;i ++)</pre>
 18
                 G[i].clear();
 19
             edges.clear();
 20
         }
 21
 22
         void addEdge(int from, int to, int cap, int cost)
 23
 24
             edges.push_back(Edge{from, to, cap, 0, cost});
 25
             edges.push_back(Edge{to, from, 0, 0, -cost});
 26
             m = edges.size();
 27
             G[from].push_back(m - 2);
 28
             G[to].push_back(m - 1);
 29
         }
 30
 31
         bool Bellman_Ford(int s, int t, int& flow, int& cost)
 32
 33
             for(int i = 0; i <= n; i++) dis[i] = inf;</pre>
 34
             memset(inq, 0, sizeof inq);
 35
             dis[s]=0, inq[s]=true, path[s]=0, a[s]=inf;
 36
             queue<int> Q;
 37
             Q.push(s);
```

```
38
            while(!Q.empty())
39
40
                int u = Q.front(); Q.pop();
                inq[u] = false;
41
                for(int i = 0; i < G[u].size(); i++)</pre>
42
43
44
                    Edge& e = edges[G[u][i]];
                    if(e.cap > e.flow && dis[e.to] > dis[u] + e.cost)
45
46
47
                        dis[e.to] = dis[u] + e.cost;
                        path[e.to] = G[u][i];
48
49
                        a[e.to] = min(a[u], e.cap - e.flow);
50
                        if(!inq[e.to])
51
                        {
52
                            Q.push(e.to);
53
                            inq[e.to] = true;
54
                    }
55
                }
56
            }
57
            if(dis[t] == inf) return false;
                                                //求最小费用最大流
58
            //if(111 * dis[t] * a[t] > 0) return false; 求可行流最小费用, 因此当费用增量大于0时不继续增加
59
60
            flow += a[t];
61
            cost += dis[t] * a[t];
62
            for(int u = t; u != s; u = edges[path[u]].from)
63
            {
64
                edges[path[u]].flow += a[t];
65
                edges[path[u] ^ 1].flow -= a[t];
66
67
            return true;
68
        }
69
70
        int mincostMaxFlow(int s, int t)
71
72
            int flow = 0, cost = 0;
73
            while(Bellman_Ford(s, t, flow, cost));
74
            return cost;
75
        }
76 };
```

2.5.4 Trick

建模技巧

二**分图带权最大独立集**。给出一个二分图,每个结点上有一个正权值。要求选出一些点,使得这些点之间没有边相连,且权值和最大。

解: 在二分图的基础上添加源点 S 和汇点 T,然后从 S 向所有 X 集合中的点连一条边,所有 Y 集合中的点向 T 连一条边,容量均为该点的权值。X 结点与 Y 结点之间的边的容量均为无穷大。这样,对于图中的任意一个割,将割中的边对应的结点删掉就是一个符合要求的解,权和为所有权减去割的容量。因此,只需要求出最小割,就能求出最大权和。

公平分配问题。把 m 个任务分配给 n 个处理器。其中每个任务有两个候选处理器,可以任选一个分配。要求所有处理器中,任务数最多的那个处理器所分配的任务数尽量少。不同任务的候选处理器集 $\{p_1, p_2\}$ 保证不同。

解:本题有一个比较明显的二分图模型,即 X 结点是任务,Y 结点是处理器。二分答案 x,然后构图,首先从源点 S 出发向所有的任务结点引一条边,容量等于 1,然后从每个任务结点出发引两条边,分别到达它所能分配到的两个处理器结点,容量为 1,最后从每个处理器结点出发引一条边到汇点 T,容量为 x,表示选择该处理器的任务不能超过 x。

这样网络中的每个单位流量都是从 S 流到一个任务结点,再到处理器结点,最后到汇点 T。只有当网络中的总流量等于m 时才意味着所有任务都选择了一个处理器。这样,我们通过 $O(\log m)$ 次最大流便算出了答案。

区间 k **覆盖问题**。数轴上有一些带权值的左闭右开区间。选出权和尽量大的一些区间,使得任意一个数最多被 k 个区间覆盖。

解: 本题可以用最小费用流解决,构图方法是把每个数作为一个结点,然后对于权值为 w 的区间 [u,v) 加边 $u \rightarrow v$,容量为 1,费用为 -w。再对所有相邻的点加边 $i \rightarrow i + 1$,容量为 k,费用为 0。最后,求最左点到最右点的最小费用最大流即可,其中每个流量对应一组互不相交的区间。如果数值范围太大,可以先进行离散化。

最大闭合子图。给定带权图 G(权值可正可负),求一个权和最大的点集,使得起点在该点集中的任意弧,终点也在该点集中。

解: 新增附加源 s 和附加汇 t, 从 s 向所有正权点引一条边,容量为权值;从所有负权点向汇点引一条边,容量为权值的相反数。求出最小割以后, $S-\{s\}$ 就是最大闭合子图。

最大密度子图。给出一个无向图,找一个点集,使得这些点之间的边数除以点数的值(称为子图的密度)最大。

解:如果两个端点都选了,就必然要选边,这就是一种推导。如果把每个点和每条边都看成新图中的结点,可以把问题转化为最大闭合子图。

无源汇有上下界可行流: 附加源 S 和汇 T; 对于边 (u,v,min,max),记 d[u]-=min,d[v]+=max,并添加弧 (u,v,max-min);对于流量不平衡的点 u,设多余流量为 W,如果 W>0,添加弧 S->u:W,否则若 W<0,添加弧 u->T:-W,求改造后的网络 S-T 最大流即可,当且仅当所有附加弧满载时原图有可行流。

有源汇有上下界可行流: 建 t->s, 容量为 inf, 然后和无源汇相同。

有源汇有上下界最大/最小流:与上面相同,跑完可行流 S->T 后去掉边 t->s,最大流为加 s->t,最小流为 $G[s][t].cap-max_flow(t,s)$ 。

2.6 Others

2.6.1 拓扑排序

```
const int maxn = 1e5 + 10;
1
2
   vector<int> edge[maxn];
3
4
   int indegree[maxn];
5
6
   void add(int u, int v)
7
   {
8
        edge[u].push_back(v);
9
        indegree[v]++;
10
   }
11
12
   void Toposort(int n)
13
   {
14
        queue<int> que;
        for (int i = 1; i <= n; i++)</pre>
15
16
           if (!indegree[i]) que.push(i);
                                           //将图中没有前驱,即入度为0的点加入队列
17
        while (!que.empty())
18
           int u = que.front();
19
20
           que.pop();
21
           indegree[u] = -1;
                               //从图中删去此顶点
           for (int i = 0; i < edge[u].size(); i++)</pre>
22
23
           {
24
               int v = edge[u][i];
                                //删去图中以u为尾的弧
               indegree[v]--;
25
```

//将新增的当前入度为0的点压入队列中

if (!indegree[v]) que.push(v);

26

```
27
           }
       }
28
29 }
    2.6.2 2-SAT
   /*2-SAT连边含义:选A必选B
1
2
        点$x_i$表示选, $x_i'$表示不选
        1.必选$x_i$, 等价于$x_i=1$: $x_i'→x_i$
3
       2. 必不选$x_i$, 等价于$x_i=0$, $x_i→x_i'$
 4
       3.$x_i$与$x_j$中至少选择一个,等价于$x_iORx_j=1$, 连边$x_i'→x_j$,$x_j'→x_i$
5
6
       4.$x_i$与$x_j$不都选,等价于$x_iANDx_j=0$, 连边$x_i→x_j',x_j→x_i'$
       5.$x_i$与$x_j$情况相同,等价于$x_iXORx_j=O$,连边$x_i→x_j$,$x_i'→x_j'$,$x_j→x_i$,$x_j'→x_i'$
7
8
        6.$x_i$与$x_j$情况相反,等价于$x_iX0Rx_j=1$,连边$x_i→x_j'$,$x_i'→x_j$,$x_j→x_i'$,$x_j'→x_i$
9
   */
10
11
   const int maxn = 2e6 + 10;
12
13 int n, m, a, va, b, vb;
   int low[maxn], dfn[maxn], color[maxn], cnt, scc_cnt;
14
15 bool instack[maxn];
16
17 vector<int> g[maxn];
18 stack<int> st;
19
20 void Tarjan(int u)
21
22
        low[u] = dfn[u] = ++cnt;
23
        st.push(u);
24
        instack[u] = true;
25
       for(const auto &v : g[u])
26
27
           if(!dfn[v]) Tarjan(v), low[u] = min(low[u], low[v]);
28
           else if(instack[v]) low[u] = min(low[u], dfn[v]);
29
       }
30
       if(low[u] == dfn[u])
31
32
           ++scc_cnt;
33
           do {
34
               color[u] = scc_cnt;
35
               u = st.top(); st.pop();
36
               instack[u] = false;
37
           } while(low[u] != dfn[u]);
38
       }
39
   }
40
   inline void add(int a, int b) { g[a].push_back(b); }
41
42
43 inline void AND(int a, int b, int c)
44
45
        if(c == 1) add(a, a + n), add(b, b + n);
46
        else add(a + n, b), add(b + n, a);
47
   }
48
49
   inline void OR(int a, int b, int c)
50
       if(c == 0) add(a + n, a), add(b + n, b);
51
```

```
52
        else add(a, b + n), add(b, a + n);
53 }
54
55
   inline void XOR(int a, int b, int c)
56
    {
57
        if(c == 0) add(a, b), add(a + n, b + n), add(b, a), add(b + n, a + n);
        else add(a, b + n), add(a + n, b), add(b, a + n), add(b + n, a);
58
59 }
60
   bool TWO_SAT()
61
62 {
63
        input();
64
        for(int i = 1; i <= (n << 1); i ++) if(!dfn[i]) Tarjan(i);</pre>
65
        for(int i = 1; i <= n; i ++)</pre>
66
            if(color[i] == color[i + n]) return false;
67
        for(int i = 1; i <= n; i ++)</pre>
68
            printf("%d ", color[i] > color[i + n]);
69
        return true;
70 }
```

2.6.3 差分约束系统

```
1 //以$x_i-x_j y$为约束条件, 建图求最短路后得到的是最大解。所有的解都不大于且尽可能逼近$dis[x0]$
2 //最短路对应最大解,最长路对应最小解
3
4 \quad const int maxn = 1000 + 10;
   const int inf = 0x3f3f3f3f;
5
6
7
   struct Edge
8
9
       int nex, to, w;
10 } edge[10 * maxn];
11
12 int head[maxn], cnt, dis[maxn], n;
13 bool vis[maxn];
14
15 void init()
16 {
17
       cnt = 0;
18
       memset(head, Oxff, sizeof head);
19 }
20
21
   void add(int u, int v, int w)
22
   {
23
       edge[cnt].nex = head[u];
24
       edge[cnt].to = v;
25
       edge[cnt].w = w;
26
       head[u] = ++cnt;
27 }
28
29 void spfa(int u)
30 {
31
       int u, v, w;
32
       for (int i = 1; i <= n; i++) dis[i] = inf, vis[i] = false;</pre>
33
       dis[u] = 0;
34
       queue<int> que;
35
       que.push(u);
36
       vis[u] = true;
```

```
37
        while (!que.empty())
38
39
            u = que.front();
40
            que.pop();
41
            vis[u] = false;
42
            for (int i = head[u]; ~i; i = edge[i].nex)
43
44
                v = edge[i].v, w = edge[i].w;
                if (dis[u] + w < dis[v])</pre>
45
46
                {
47
                    dis[v] = dis[u] + w;
                    if (!vis[v])
48
49
50
                         que.push(v);
51
                         vis[v] = true;
52
53
                }
54
            }
55
        }
56 }
    2.6.4 支配树
    const int N = 2e5 + 10;
2
3 int n, m;
4
5
   struct G
6
7
        vector<int> edge[N];
8
        inline void add(int u, int v) { edge[u].push_back(v); }
9
   }a, b, c, d;
10
   int dfn[N], id[N], fa[N], cnt;
11
12
13 void dfs(int u)
14 {
15
        dfn[u] = ++ cnt; id[cnt] = u;
16
        int len = a.edge[u].size();
17
        for(auto v : a.edge[u]) if(!dfn[v]) { fa[v] = u; dfs(v); }
18 }
19
20 int semi[N], idom[N], belong[N], val[N];
21
22
   int find(int x)
23 {
24
        if(x == belong[x]) return x;
25
        int tmp = find(belong[x]);
26
        if(dfn[semi[val[belong[x]]]] < dfn[semi[val[x]]]) val[x] = val[belong[x]];</pre>
27
        return belong[x] = tmp;
28 }
29
30 void tarjan()
31
32
        for(int i = cnt; i > 1; i --)
33
34
            int u = id[i];
35
            for(auto v : b.edge[u])
```

```
{
36
                 if(!dfn[v]) continue;
37
38
                 find(v);
                 if(dfn[semi[val[v]]] < dfn[semi[u]]) semi[u] = semi[val[v]];</pre>
39
40
41
            c.add(semi[u], u);
42
            belong[u] = fa[u];
            u = fa[u];
43
            for(auto v : c.edge[u])
44
45
            {
46
                 find(v);
47
                 if(semi[val[v]] == u) idom[v] = u;
48
                 else idom[v] = val[v];
            }
49
        }
50
        for(int i = 2; i <= cnt; i ++)</pre>
51
52
53
            int u = id[i];
            if(idom[u] != semi[u]) idom[u] = idom[idom[u]];
54
        }
55
56 }
57
58 int ans[N];
60 void dfs_ans(int u)
61 {
62
        ans[u] = 1;
        for(auto v : d.edge[u]) dfs_ans(v), ans[u] += ans[v];
63
64 }
65
66
   void solve()
67
   {
68
        int u, v;
        scanf("%d%d", &n, &m);
69
        while(m --)
70
71
72
            scanf("%d%d", &u, &v);
73
            a.add(u, v);
            b.add(v, u);
74
75
76
        for(int i = 1; i <= n; i ++) semi[i] = belong[i] = val[i] = i;</pre>
77
        dfs(1);
78
        tarjan();
79
        for(int i = 2; i <= n; i ++) d.add(idom[i], i);</pre>
80
        dfs_ans(1);
81
        for(int i = 1; i <= n; i ++) printf("%d ", ans[i]);</pre>
82 }
```

3 DataStructrue

3.1 SegmentTreeDS

3.1.1 SegmentTree

```
struct SegmentTree {
 1
    #define TYPE int
 2
 3
        TYPE val[maxn << 2];</pre>
 4
        int sz;
 5
        // check this type
 6
        int lazy[maxn << 2];</pre>
 7
 8
        inline TYPE comb(const TYPE& a, const TYPE& b) {
 9
             TYPE res;
10
             res = a + b;
11
             return res;
12
13
14
        int le, re, k;
15
16
        inline void init(int sz_) {
17
             sz = sz_;
18
             memset(lazy, 0, sizeof lazy);
19
             memset(val, 0, sizeof val);
20
21
        inline void pushup(int rt) {val[rt] = comb(val[rt << 1], val[rt << 1 | 1]);}</pre>
22
        inline void deal(int rt, int kt) {
23
24
25
        inline void pushdown(int rt, int len) {
26
             if (lazy[rt]) {
                 // check the lazy change
27
28
                 lazy[rt << 1] += lazy[rt];</pre>
                 lazy[rt << 1 | 1] += lazy[rt];</pre>
29
30
                 deal(rt << 1, lazy[rt]);</pre>
                 deal(rt << 1 | 1, lazy[rt]);</pre>
31
32
                 lazy[rt] = 0;
33
             }
34
        }
35
36
         inline void build(int rt, int 1, int r) {
37
             if (1 == r) {
38
                 val[rt] = a[1];
39
                 return;
40
41
             int mid = 1 + r >> 1;
42
             build(rt << 1, 1, mid);</pre>
43
             build(rt << 1 | 1, mid + 1, r);
44
             pushup(rt);
45
46
        inline void build() {build(1, 1, sz);}
47
48
        inline void update(int rt, int 1, int r) {
49
             if (le <= 1 && r <= re) {</pre>
50
                 deal(rt, k);
51
                 return;
52
53
             pushdown(rt, r - 1 + 1);
```

```
54
            int mid = 1 + r >> 1;
55
            if (le <= mid) update(rt << 1, 1, mid);</pre>
56
            if (re > mid) update(rt << 1 | 1, mid + 1, r);</pre>
57
            pushup(rt);
58
59
60
        inline TYPE query(int rt, int 1, int r) {
            if (le <= 1 && r <= re) {</pre>
61
                return val[rt];
62
63
            }
            pushdown(rt, r - l + 1);
64
65
            // check the zero type
            TYPE res;
66
            int mid = 1 + r >> 1;
67
68
            if (le <= mid) res = comb(res, query(rt << 1, 1, mid));</pre>
69
            if (re > mid) res = comb(res, query(rt << 1 | 1, mid + 1, r));
70
            return res;
71
        }
72
73
        // check return type
74
        inline int query(int 1, int r) {
75
            le = 1, re = r;
76
            return query(1, 1, sz);
77
78
        inline void modify(int 1, int r, int kt) {
79
            le = 1, re = r, k = kt;
80
            update(1, 1, sz);
        }
81
82
83
        inline void pt(int rt, int 1, int r) {
84
            if (1 == r) {
85
                printf("%d ", val[1]);
86
                return;
            }
87
88
            pushdown(rt, r - l + 1);
89
            int mid = 1 + r >> 1;
90
            if (le <= mid) pt(rt << 1, 1, mid);</pre>
91
            if (re > mid) pt(rt << 1 | 1, mid + 1, r);</pre>
92
        }
93
94 #undef TYPE
95 };
    3.1.2 离散化区间
1 // 原题1e5个个区间有2e5个端点,离散化出来4e5个区间
2 // 然后线段树需要4e5*4=16e5的大小
3 // 注意三个数组要开离散化数量的四倍,如果不需要sz可以不用这个数组。
4 \quad int \quad val[maxn << 4];
5 int lpos[maxn << 2], rpos[maxn << 2], tot, sz[maxn << 2];</pre>
6 vector<int> xpos;
7 sort(xpos.begin(), xpos.end());
8 xpos.erase(unique(xpos.begin(), xpos.end()), xpos.end());
9 \text{ tot = 1;}
10 \ lpos[1] = rpos[1] = xpos[0];
11 \text{ sz}[1] = 1;
12 for (int i = 1; i < xpos.size(); ++i) {
        if (xpos[i] - xpos[i - 1] != 1) {
13
```

```
14
            lpos[++tot] = xpos[i - 1] + 1;
15
            rpos[tot] = xpos[i] - 1;
16
            sz[tot] = rpos[tot] - lpos[tot] + 1;
17
        }
18
        ++tot;
19
        lpos[tot] = rpos[tot] = xpos[i];
20
        sz[tot] = 1;
21
   }
22 le = lower_bound(lpos + 1, lpos + 1 + tot, p[i].x) - lpos;
23 re = upper_bound(rpos + 1, rpos + 1 + tot, p[i].y) - rpos - 1;
    3.1.3 动态区间最大子段和
    namespace ST {
2
        struct node{
3
            11 ans,ls,rs,sum;
 4
        xx[maxn << 2];
5
        inline void pushdown(int x){
            xx[x].sum=xx[x<<1].sum+xx[x<<1|1].sum;
6
7
            xx[x].ls=max(xx[x<<1].ls,xx[x<<1].sum+xx[x<<1|1].ls);</pre>
8
            xx[x].rs=max(xx[x<<1|1].rs,xx[x<<1|1].sum+xx[x<<1].rs);
9
            xx[x].ans=max(xx[x<<1].ans,max(xx[x<<1|1].ans,xx[x<<1].rs+xx[x<<1|1].ls));
10
            return;
11
        }
        inline void build(int k,int l,int r){
12
13
            if(l==r){
14
                xx[k].ls=xx[k].rs=xx[k].ans=xx[k].sum=0;
15
                return;
16
            }
17
            int mid=l+r>>1;
            build(k<<1,1,mid),build(k<<1|1,mid+1,r);
18
19
            pushdown(k);
20
            return;
21
        }
22
        inline void change(int k,int l,int r,int x,int y,int w){ // 1, 1, n
23
            if(x<=1&&r<=y){</pre>
24
                xx[k].ls += w;
25
                xx[k].rs += w;
26
                xx[k].ans += w;
27
                xx[k].sum += w;
28
   //
                  xx[k].ls=xx[k].rs=xx[k].ans=xx[k].sum=w;
29
                return;
30
            }
31
            int mid=l+r>>1;
32
            if(x<=mid) change(k<<1,1,mid,x,y,w);</pre>
33
            if(mid<y) change(k<<1|1,mid+1,r,x,y,w);</pre>
34
            pushdown(k);
35
            return;
36
        }
37
        inline node query(int k,int l,int r,int x,int y){
38
            if(x<=1&&r<=y) {</pre>
39
                return xx[k];
40
            }
41
            int mid=l+r>>1;
42
            if(x<=mid&&!(mid<y)) return query(k<<1,1,mid,x,y);</pre>
43
            else if(!(x<=mid)&&mid<y) return query(k<<1|1,mid+1,r,x,y);
44
                node st,t1=query(k <<1,1,mid,x,y),t2=query(k <<1|1,mid+1,r,x,y);
45
```

```
46
                st.sum=t1.sum+t2.sum;
47
                st.ls=max(t1.ls,t1.sum+t2.ls);
48
                st.rs=max(t2.rs,t2.sum+t1.rs);
49
                st.ans=max(t1.ans,max(t2.ans,t1.rs+t2.ls));
50
                return st;
51
            }
52
        }
53
  }
    3.1.4 动态开点权值线段树
1 int root[100005];
   int ls[1800000], rs[1800000], sum[1800000];
3
   int sz = 0;
4
5
   void insert(int &k, int 1, int r, int val){
6
        if (!k) k = ++sz;
7
        if (1 == r) {
8
            sum[k] = 1;
9
            return;
10
       }
11
        int mid = (1 + r) >> 1;
12
        if (val <= mid) insert(ls[k], l, mid, val);</pre>
13
        else insert(rs[k], mid + 1, r, val);
        sum[k] = sum[ls[k]] + sum[rs[k]];
14
15
   }
16
17
   int query(int k, int l, int r, int rank) {
18
        if (1 == r) return 1;
19
        int mid = (1 + r) >> 1;
20
        if (sum[ls[k]] >= rank) return query(ls[k], 1, mid, rank);
21
        else return query(rs[k], mid + 1, r, rank - sum[ls[k]]);
22 }
23 int merge(int x, int y)
24 {
25
        if (!x) return y;
26
        if (!y) return x;
27
       ls[x] = merge(ls[x], ls[y]);
28
       rs[x] = merge(rs[x], rs[y]);
29
        sum[x] = sum[ls[x]] + sum[rs[x]];
30
        return x;
31 }
32 insert(root[i], 1, n, a[i]);
33 query(root[p], 1, n, x);
    3.1.5 扫描线
1 // 范用型扫描线, del储存上界+1, add储存下届, 先del后add即可
   struct node {
3
        int lpos, rpos, linepos;
4
        bool operator < (const node& oth) const {</pre>
5
            return linepos < oth.linepos;</pre>
6
       }
7
   };
8
   vector<node> add, del;
9 int delpos = 0;
10 int res = 0;
11 for (int addpos = 0; addpos < add.size(); ++addpos) {</pre>
```

```
12
        while (delpos < del.size() && del[delpos].linepos <= add[addpos].linepos) {</pre>
13
            up(del[delpos].lpos, del[delpos].rpos, -1);
14
            delpos ++;
15
16
        up(add[addpos].lpos, add[addpos].rpos, 1);
17
        res = max(res, val[1]);
18 }
19
20 // 求面积并
21 #define maxn 222
22 #define tmp (st<<1)
23 #define mid ((1+r)>>1)
24 #define lson l,mid,tmp
25 #define rson mid+1,r,tmp|1
26 using namespace std;
27 int cnt[maxn<<2];</pre>
28 double sum[maxn<<2];
29 double x[maxn];
30 struct Seg{
        double h,1,r;
31
32
        int s;
33
        Seg(){}
        Seg(double a, double b, double c, int d):1(a),r(b),h(c),s(d){}
34
35
        bool operator<(const Seg &cmp)const{</pre>
36
            return h<cmp.h;</pre>
        }
37
    }ss[maxn];
38
39
    void push_up(int st,int l,int r){
        if(cnt[st])sum[st]=x[r+1]-x[1];
40
41
        else if(l==r)sum[st]=0;
42
        else sum[st] = sum[tmp] + sum[tmp|1];
43
    }
    void update(int L,int R,int c,int 1,int r,int st){
44
45
        if(L<=1&&r<=R){</pre>
46
            cnt[st]+=c;
47
            push_up(st,1,r);
48
            return ;
49
50
        if(L<=mid)update(L,R,c,lson);</pre>
51
        if(R>mid)update(L,R,c,rson);
52
        push_up(st,1,r);
53
   }
54
    int main(){
55
        int n,tot=1,m;
        while(scanf("%d",&n)&&n){
56
57
            double a,b,c,d;
            m=0;
58
59
            while(n--){
                scanf("%lf%lf%lf",&a,&b,&c,&d);
60
61
                ss[m++]=Seg(a,c,b,1);
62
63
                x[m]=c;
64
                ss[m++]=Seg(a,c,d,-1);
65
            }
            sort(x,x+m);
66
67
            sort(ss,ss+m);
68
            double ans=0;
69
            for(int i=0;i<m;++i){</pre>
                int l=lower_bound(x,x+m,ss[i].1)-x;
70
```

```
71
                 int r=lower_bound(x,x+m,ss[i].r)-x-1;
 72
                 update(l,r,ss[i].s,0,m-1,1);
 73
                 ans+=sum[1]*(ss[i+1].h-ss[i].h);
 74
 75
             printf("Test case #%dnTotal explored area: %.21fnn",tot++,ans);
 76
         }
 77
         return 0;
     }
 78
 79
 80 // 面积交
 81 #include<br/>bits/stdc++.h>
 82 #define maxn 100005
 83 #define lson l,mid,rt<<1
 84 #define rson mid+1,r,rt<<1|1
 85 #define pb push_back
 86 using namespace std;
 87
 88 double tree[maxn<<2],tree2[maxn<<2];
 89 int lazy[maxn<<2];
 90 vector<double>ve;
 91
 92
     struct seg{
 93
         double 1,r,h;
 94
         int flag;
 95
         seg(){}
 96
         seg(double _l,double _r,double _h,int _flag){l=_l,r=_r,h=_h,flag=_flag;}
         bool operator<(const seg &b)const{return h<b.h;}</pre>
 97
 98
     }s[maxn];
99
100
     void push_up(int 1,int r,int rt){
101
         if(lazy[rt]) tree[rt]=ve[r]-ve[l-1];
102
         else if(l==r) tree[rt]=0;
103
         else tree[rt]=tree[rt<<1]+tree[rt<<1|1];</pre>
104
    }
105
     void push_up2(int 1,int r,int rt){
106
107
         if(lazy[rt]>1) tree2[rt]=ve[r]-ve[l-1];
108
         else if(l==r) tree2[rt]=0;
         else if(lazy[rt]==1)tree2[rt]=tree[rt<<1]+tree[rt<<1|1];</pre>
109
110
         else tree2[rt]=tree2[rt<<1]+tree2[rt<<1|1];</pre>
111 }
112
113 void build(int l,int r,int rt){
114
         tree[rt]=0,lazy[rt]=0;
115
         if(l==r) return;
         int mid=l+r>>1;
116
117
         build(lson);
         build(rson);
118
119 }
120
121
     void add(int L,int R,int v,int l,int r,int rt){
122
         if(L<=1&&R>=r){
123
             lazy[rt]+=v;
124
             push_up(1,r,rt);
125
             push_up2(1,r,rt);
126
             return;
127
128
         int mid=l+r>>1;
129
         if(L<=mid) add(L,R,v,lson);</pre>
```

```
130
         if(R>mid) add(L,R,v,rson);
131
         push_up(1,r,rt);
132
         push_up2(1,r,rt);
133
     }
134
135
     int getid(double x){ return lower_bound(ve.begin(), ve.end(),x)-ve.begin()+1;}
136
137
     int main(){
138
         int n;
139
         int Case=1;
140
         int T;
141
         scanf("%d",&T);
142
         while(T--){
             scanf("%d",&n);
143
144
             ve.clear();
145
             int tot=0;
146
             double x1,y1,x2,y2;
             for(int i=1;i<=n;i++){</pre>
147
148
                  scanf("%lf %lf %lf",&x1,&y1,&x2,&y2);
                  ve.pb(x1), ve.pb(x2);
149
150
                  s[++tot] = seg(x1,x2,y1,1);
                  s[++tot] = seg(x1,x2,y2,-1);
151
             }
152
153
             sort(ve.begin(),ve.end());
154
             ve.erase(unique(ve.begin(),ve.end()),ve.end());
155
             sort(s+1,s+tot+1);
156
             int N=ve.size();
157
             build(1,N,1);
158
             double ans=0;
159
             for(int i=1;i<tot;i++){</pre>
160
                  int L=getid(s[i].1);
161
                  int R=getid(s[i].r)-1;
162
                  add(L,R,s[i].flag,1,N,1);
163
                  ans+=tree2[1]*(s[i+1].h-s[i].h);
164
             }
165
             printf("%.2f\n",ans);
166
         }
167
     }
168
    // 求周长并
169
170 #include<bits/stdc++.h>
171 #define maxn 100005
     #define lson l,mid,rt<<1</pre>
    #define rson mid+1,r,rt<<1|1</pre>
174 #define pb push_back
175
     using namespace std;
176
177   int tree[maxn<<2];</pre>
178 int lazy[maxn<<2];
179 vector<int>ve[2];
180 int k;
181
182
     struct seg{
183
         int 1,r,h;
184
         int flag;
185
         seg(){}
186
         seg(int _l,int _r,int _h,int _flag){l=_l,r=_r,h=_h,flag=_flag;}
187
         bool operator<(const seg &b)const{return h<b.h;}</pre>
188 }s[maxn];
```

```
189
190
     void push_up(int 1,int r,int rt){
         if(lazy[rt]) tree[rt]=ve[k][r]-ve[k][1-1];
191
192
         else if(l==r) tree[rt]=0;
193
         else tree[rt] = tree[rt << 1] + tree[rt << 1 | 1];</pre>
194
    }
195
     void build(int l,int r,int rt){
196
         tree[rt]=0,lazy[rt]=0;
197
198
         if(l==r) return;
199
         int mid=l+r>>1;
200
         build(lson);
201
         build(rson);
202
     }
203
204
     void add(int L,int R,int v,int l,int r,int rt){
205
         if(L<=1&&R>=r){
206
              lazy[rt]+=v;
207
              push_up(1,r,rt);
208
              return;
         }
209
210
         int mid=l+r>>1;
211
         if(L<=mid) add(L,R,v,lson);</pre>
212
         if(R>mid) add(L,R,v,rson);
213
         push_up(1,r,rt);
214 }
215
216
     int getid(int x){return lower_bound(ve[k].begin(),ve[k].end(),x)-ve[k].begin()+1;}
217
218
     int main(){
219
         int n;
220
         while(~scanf("%d",&n)){
221
              ve[0].clear();
222
              ve[1].clear();
223
              int x1,y1,x2,y2;
              for(int i=1;i<=n;i++){</pre>
224
225
                  scanf("%d %d %d",&x1,&y1,&x2,&y2);
                  ve[0].pb(x1),ve[0].pb(x2);
226
227
                  ve[1].pb(y1),ve[1].pb(y2);
228
                  s[i] = seg(x1,x2,y1,1);
229
                  s[i+n]=seg(x1,x2,y2,-1);
230
                  s[i+n+n]=seg(y1,y2,x1,1);
231
                  s[i+n+n+n] = seg(y1,y2,x2,-1);
232
             }
233
             int ans=0;
234
              int pos=1;
235
              for(k=0;k<2;k++){</pre>
236
                  sort(ve[k].begin(),ve[k].end());
                  ve[k].erase(unique(ve[k].begin(),ve[k].end()),ve[k].end());
237
238
                  sort(s+pos,s+pos+n+n);
239
                  int N=ve[k].size();
                  build(1,N,1);
240
241
                  int pre=0;
242
                  for(int i=pos;i<pos+n+n;i++){</pre>
243
                      int L=getid(s[i].1);
244
                      int R=getid(s[i].r)-1;
245
                      add(L,R,s[i].flag,1,N,1);
246
                      ans+=abs(tree[1]-pre);
                      pre=tree[1];
247
```

```
}
248
249
               pos+=n+n;
250
           }
251
           printf("%d\n",ans);
252
        }
253 }
    3.2 HLD
    3.2.1 HLD
 1 #include <bits/stdc++.h>
 2 #define 11 long long
 3 using namespace std;
 4 /*
 5 node 计算点权, path 下放后计算边权, edge 根据边的编号计算边权
 6 work 中没有build需手动写
    sz[]数组,以x为根的子树节点个数
 8 top[]数组, 当前节点的所在链的顶端节点
 9 son[]数组, 重儿子
10 deep[]数组, 当前节点的深度
11 fa[]数组,当前节点的父亲
12 idx[]数组,树中每个节点剖分后的新编号
13 rnk[]数组, idx的逆, 表示线段上中当前位置表示哪个节点
14 */
15
16 const int maxn = 1e5+5;
17
18 int sz[maxn], top[maxn], son[maxn], deep[maxn], fa[maxn], idx[maxn], rnk[maxn];
19 int tot;
20
    int n, le, re;
21
    11 k;
22
23 struct HLD {
24
    #define type int
25
26
        struct edge {
27
           int a, b;
28
           type v;
29
30
           edge(int _a, int _b, type _v = 0) : a(_a), b(_b), v(_v) {}
31
        };
32
33
        struct node {
34
           int to;
35
           type w;
36
37
           node() {}
38
39
           node(int _to, type _w) : to(_to), w(_w) {}
40
        };
41
42
        vector<int> mp[maxn];
43
        vector<edge> e;
44
45
        void init(int _n) {
46
47
           for (int i = 0; i <= n; i++) mp[i].clear();</pre>
48
           e.clear();
```

```
49
             e.push_back(edge(0, 0));
         }
 50
51
 52
         void add_edge(int a, int b, type v = 0) {
 53
     //
                e.push_back(edge(a,b,v));
 54
             mp[a].push_back(b);
 55
             mp[b].push_back(a);
         }
56
 57
 58
         void dfs1(int x, int pre, int h) {
 59
             int i, to;
60
             deep[x] = h;
61
             fa[x] = pre;
62
             sz[x] = 1;
63
             for (i = 0; i < (int) (mp[x].size()); i++) {</pre>
 64
                  to = mp[x][i];
 65
                  if (to == pre) continue;
 66
                  dfs1(to, x, h + 1);
 67
                  sz[x] += sz[to];
                  if (son[x] == -1 \mid \mid sz[to] > sz[son[x]]) son[x] = to;
 68
 69
             }
         }
 70
 71
 72
         void dfs2(int x, int tp) {
 73
             int i, to;
 74
             top[x] = tp;
 75
             idx[x] = ++tot;
 76
             rnk[idx[x]] = x;
 77
             if (son[x] == -1) return;
 78
             dfs2(son[x], tp);
 79
             for (i = 0; i < (int) (mp[x].size()); i++) {</pre>
 80
                  to = mp[x][i];
 81
                  if (to != son[x] && to != fa[x]) dfs2(to, to);
             }
 82
         }
 83
 84
 85
         void work(int _rt = 1) {
 86
             memset(son, -1, sizeof son);
 87
             tot = 0;
 88
             dfs1(_rt, 0, 0);
 89
             dfs2(_rt, _rt);
         }
 90
 91
 92
         int LCA(int x, int y) {
             while (top[x] != top[y]) {
 93
                  if (deep[top[x]] < deep[top[y]]) swap(x, y);</pre>
 94
 95
                  x = fa[top[x]];
             }
 96
             if (deep[x] > deep[y]) swap(x, y);
 97
 98
             return x;
99
100
101
         void modify_node(int x, int y, type val) {
             while (top[x] != top[y]) {
102
                  if (deep[top[x]] < deep[top[y]]) swap(x, y);</pre>
103
104
                  le = idx[top[x]], re = idx[x];
105
                  k = val;
106
                  update(1, 1, n);
107
                  x = fa[top[x]];
```

```
108
             }
109
             if (deep[x] > deep[y]) swap(x, y);
             le = idx[x], re = idx[y];
110
111
             k = val;
112
             update(1, 1, n);
113
         }
114
         type query_node(int x, int y) {
115
             type res = 0;
116
117
             while (top[x] != top[y]) {
118
                 if (deep[top[x]] < deep[top[y]]) swap(x, y);</pre>
119
                 le = idx[top[x]], re = idx[x];
120
                 res += query(1, 1, n);
121
                 x = fa[top[x]];
             }
122
             if (deep[x] > deep[y]) swap(x, y);
123
124
             le = idx[x], re = idx[y];
125
             res += query(1, 1, n);
126
             return res;
127
         }
128
129
         //path
130 //
           void init_path()
131 //
           {
132 //
               v[idx[rt]]=0;
133 //
               for(int i=1;i<n;i++)</pre>
134 //
135
    //
                    if(deep[e[i].a] < deep[e[i].b]) swap(e[i].a,e[i].b);
136
    //
                   a[idx[e[i].a]]=e[i].v;
               }
137
    //
138
    //
               build(n);
139
    //
140
         void modify_edge(int id, type val) {
141
             if (deep[e[id].a] > deep[e[id].b]) {
142
                 le = idx[e[id].a], re = idx[e[id].a];
143
                 k = val;
144
                 update(1, 1, n);
145
             } else {
146
                 le = idx[e[id].b], re = idx[e[id].b];
147
                 k = val;
148
                 update(1, 1, n);
             }
149
         }
150
151
         void modify_path(int x, int y, type val) {
152
             while (top[x] != top[y]) {
153
154
                 if (deep[top[x]] < deep[top[y]]) swap(x, y);</pre>
                 le = idx[top[x]], re = idx[x];
155
156
                 k = val;
157
                 update(1, 1, n);
158
                 x = fa[top[x]];
159
             }
160
             if (deep[x] > deep[y]) swap(x, y);
161
             if (x != y) {
                 le = idx[x] + 1, re = idx[y];
162
163
                 k = val;
164
                 update(1, 1, n);
165
             }
166
         }
```

```
167
168
         type query_path(int x, int y) {
169
             type res = 0;
170
             while (top[x] != top[y]) {
171
                 if (deep[top[x]] < deep[top[y]]) swap(x, y);</pre>
172
                 le = idx[top[x]], re = idx[x];
173
                 res += query(1, 1, n);
174
                 x = fa[top[x]];
             }
175
176
             if (deep[x] > deep[y]) swap(x, y);
177
             if (x != y) {
178
                 le = idx[x] + 1, re = idx[y];
179
                 res += query(1, 1, n);
180
             }
181
             return res;
         }
182
183
184 #undef type
185 } hld;
     3.3 RMQ
     3.3.1 RMQ
 1 int A[maxn];
   int maxx[maxn][22];
    void RMQ(int n) {
         for (int i = 1; i <= n; i++)</pre>
 4
 5
             maxx[i][0] = A[i];
 6
         for (int j = 1; (1 << j) <= n; j++) {
 7
             for (int i = 1; i + (1 << j) - 1 <= n; i++) {
 8
                 \label{eq:maxx} \max[i][j] = \max(\max[i][j-1], \ \max[i+(1<<(j-1))][j-1]);
 9
             }
10
         }
11 }
12 int query(int 1, int r) {
13
         int k = 0;
14
         while ((1 << (k + 1)) <= r - 1 + 1) k++;
         return max(maxx[1][k], maxx[r - (1 << k) + 1][k]);</pre>
15
16 }
17
18 template <typename T, class F = function<T(const T&, const T&)>>
19
    class SparseTable {
20
     public:
21
      int n;
22
       vector<vector<T>> mat;
23
       F func;
24
25
       SparseTable(const vector<T>& a, const F& f) : func(f) {
26
         n = static_cast<int>(a.size());
27
         int max_log = 32 - __builtin_clz(n);
28
         mat.resize(max_log);
29
         mat[0] = a;
30
         for (int j = 1; j < max_log; j++) {</pre>
31
           mat[j].resize(n - (1 << j) + 1);
32
           for (int i = 0; i <= n - (1 << j); i++) {</pre>
33
             mat[j][i] = func(mat[j-1][i], mat[j-1][i+(1 << (j-1))]);
34
35
         }
```

```
36
      }
37
38
      T get(int from, int to) const {
39
        assert(0 <= from && from <= to && to <= n - 1);
40
        int lg = 32 - __builtin_clz(to - from + 1) - 1;
41
        return func(mat[lg][from], mat[lg][to - (1 << lg) + 1]);</pre>
42
      }
43 };
   //静态区间最大值
44
    SparseTable<int> solve(v, [&](int i, int j) {return max(i, j);});
    3.3.2 RMQbyIndex
1 //下标RMQ
2 int v[MAX], maxx[MAX][22], minn[MAX][22];
3 int pmax(int a,int b){return v[a]>v[b]?a:b;}
 4 int pmin(int a,int b){return v[a]<v[b]?a:b;}</pre>
 5 void RMQ(int n) {
6
        int i,j;
7
        for(i=1;i<=n;i++) {</pre>
            maxx[i][0]=minn[i][0]=i;
8
9
        for(j=1;1<<(j-1)<=n;j++) {
10
            for(i=1;i+(1<<j)-1<=n;i++) {</pre>
11
                int t=1<<(j-1);</pre>
12
13
                maxx[i][j]=pmax(maxx[i][j-1],maxx[i+t][j-1]);
14
                minn[i][j]=pmin(minn[i][j-1],minn[i+t][j-1]);
15
            }
16
        }
17 }
18
   int query(int 1,int r) {
        int j=(int)(log10(r-l+1)/log10(2))+1;
19
20
        int i=r-(1<<(j-1))+1;</pre>
        return pmax(maxx[1][j-1],maxx[i][j-1]);
22 // return pmin(minn[l][j-1],minn[i][j-1]);
23 }
    3.3.3 RMQinNM
1 //二维RMQ
   int v[302][302];
    int maxx[302][302][9][9],minn[302][302][9][9];
3
 4
    void RMQ(int n,int m)
5
    {
6
        int i,j,ii,jj;
7
        for(i=1;i<=n;i++)</pre>
8
        {
9
            for(j=1;j<=m;j++)</pre>
10
            {
11
                maxx[i][j][0][0]=minn[i][j][0][0]=v[i][j];
12
13
        }
14
        for(ii=0;(1<<ii)<=n;ii++)</pre>
15
16
            for(jj=0;(1<<jj)<=m;jj++)</pre>
17
                if(ii+jj)
18
19
                {
```

```
20
                    for(i=1;i+(1<<ii)-1<=n;i++)</pre>
21
22
                        for(j=1;j+(1<<jj)-1<=m;j++)</pre>
23
                        {
24
                             if(ii)
25
                             {
26
                                 minn[i][j][ii][jj]=min(minn[i][j][ii-1][jj],minn[i+(1<<(ii-1))][j][ii
        -1][jj]);
27
                                 maxx[i][j][ii][jj]=max(maxx[i][j][ii-1][jj],maxx[i+(1<<(ii-1))][j][ii
        -1][jj]);
28
                            }
29
                             else
30
                             {
                                 minn[i][j][ii][jj]=min(minn[i][j][ii][jj-1],minn[i][j+(1<<(jj-1))][ii][
31
        jj-1]);
                                 \max[i][j][ii][jj] = \max(\max[i][j][ii][jj-1], \max[i][j+(1<<(jj-1))][ii][
32
        jj-1]);
33
                            }
34
                        }
                    }
35
36
                }
37
            }
        }
38
39 }
40 int query(int x1,int y1,int x2,int y2)
41 {
42
        int k1=0;
43
        while((1<<(k1+1))<=x2-x1+1) k1++;</pre>
44
        int k2=0;
45
        while((1<<(k2+1))<=y2-y1+1) k2++;</pre>
46
        x2=x2-(1<< k1)+1;
47
        y2=y2-(1<< k2)+1;
        return max(maxx[x1][y1][k1][k2], maxx[x1][y2][k1][k2]), max(maxx[x2][y1][k1][k2], maxx[x2][y2
48
        ][k1][k2]))
   // return min(min(minn[x1][y1][k1][k2], minn[x1][y2][k1][k2]), min(minn[x2][y1][k1][k2], minn[x2][y2
        ][k1][k2]));
50 }
    3.4 MO
   3.4.1 MO
   // const int maxn = 50005;
2
3
   struct MO {
4
        int 1, r, id;
5
   }q[maxn];
6
7 int n, m, col[maxn], block, belong[maxn];
8 int vis[maxn * 10];
9 11 res[maxn], ans;
10 bool cmp(const MO& a, const MO& b) { return belong[a.1] == belong[b.1] ? a.r < b.r : a.1 < b.1; }
11 void add(int x) {
12
        vis[x] ++;
13
        ans += 111 * x * (vis[x] * vis[x] - (vis[x] - 1) * (vis[x] - 1));
14 }
15
16 void del(int x) {
        vis[x] --;
17
```

```
ans -= 111 * x * ((vis[x] + 1) * (vis[x] + 1) - vis[x] * vis[x]);
18
   }
19
20
21
    int main() {
22
        scanf("%d%d", &n, &m);
23
        block = sqrt(n);
24
        for (int i = 1; i <= n; ++i) {</pre>
25
            scanf("%d", &col[i]);
26
            belong[i] = i / block + 1;
27
        }
28
        for (int i = 1; i <= m; ++i) {</pre>
29
            scanf("%d%d", &q[i].1, &q[i].r);
30
            q[i].id = i;
        }
31
32
        sort(q + 1, q + 1 + m, cmp);
33
        int 1 = 1, r = 0;
34
        for (int i = 1; i <= m; ++i) {</pre>
            while(r < q[i].r) add(col[++r]);</pre>
35
36
            while(r > q[i].r) del(col[r--]);
            while(1 < q[i].1) del(col[1++]);</pre>
37
38
            while(1 > q[i].1) add(col[--1]);
39
            res[q[i].id] = ans;
        }
40
41
        for (int i = 1; i <= m; ++i) printf("%lld\n", res[i]);</pre>
42
        return 0;
43 }
    3.4.2 MObyModify
 1 #include <bits/stdc++.h>
 2 #define 11 long long
3 using namespace std;
4 const int maxn = 50005;
5
6 struct MO {
7
        int 1, r, id, oppre;
8 }q[maxn];
10 int n, m, col[maxn], block, belong[maxn], colpre[maxn];
int changepos[maxn], changepre[maxn], changenow[maxn];
12 int vis[maxn * 20];
13 int ans;
    int res[maxn];
15
    bool cmp(const MO& a, const MO& b) {
16
        if (belong[a.1] != belong[b.1]) return a.1 < b.1;</pre>
17
        if (belong[a.r] != belong[b.r]) return a.r < b.r;</pre>
18
        return a.oppre < b.oppre;</pre>
19 }
20 void add(int x) {}
21
22 void del(int x) {}
23
   void unmodify(int pos, int now) {
24
25
        if (q[pos].1 <= changepos[now] && changepos[now] <= q[pos].r) {</pre>
26
            del(changenow[now]);
27
            add(changepre[now]);
28
        col[changepos[now]] = changepre[now];
29
```

```
30 }
31
    void modify(int pos, int now) {
32
33
        if (q[pos].l <= changepos[now] && changepos[now] <= q[pos].r) {</pre>
34
            del(changepre[now]);
35
            add(changenow[now]);
36
        }
37
        col[changepos[now]] = changenow[now];
    }
38
39
40
    int main() {
41
        scanf("%d%d", &n, &m);
42
        block = pow(n, 0.66666);
        for (int i = 1; i <= n; ++i) {</pre>
43
44
            scanf("%d", &col[i]);
45
            colpre[i] = col[i];
46
            belong[i] = i / block + 1;
47
        }
48
        char s[2];
49
        int t = 0, t2 = 0;
        for (int i = 1; i <= m; ++i) {</pre>
50
            scanf("%s", s);
51
52
            if (s[0] == 'Q') {
53
                ++t;
54
                scanf("%d%d", &q[t].1, &q[t].r);
55
                q[t].oppre = t2;
56
                q[t].id = t;
57
            } else {
58
                ++t2;
59
                scanf("%d%d", &changepos[t2], &changenow[t2]);
60
                changepre[t2] = colpre[changepos[t2]];
61
                colpre[changepos[t2]] = changenow[t2];
62
            }
        }
63
64
        sort(q + 1, q + 1 + t, cmp);
65
        int 1 = 1, r = 0, now = 0;
        for (int i = 1; i <= t; ++i) {</pre>
66
67
            while(r < q[i].r) add(col[++r]);</pre>
68
            while(r > q[i].r) del(col[r--]);
            while(1 < q[i].1) del(col[1++]);</pre>
69
70
            while(1 > q[i].1) add(col[--1]);
71
            while (now < q[i].oppre) modify(i, ++now);</pre>
72
            while (now > q[i].oppre) unmodify(i, now--);
73
            res[q[i].id] = ans;
        }
74
75
        for (int i = 1; i <= t; ++i) printf("%d\n", res[i]);</pre>
76
        return 0;
77 }
    3.4.3 分块
1 // 非预处理数组版
2 inline int belong(int x) { return (x - 1) / block + 1; }
3 inline int lpos(int x) { return 1 + (x - 1) * block; }
   inline int rpos(int x) { return min(n, x * block); }
5
   int sz = (n - 1) / block + 1;
7 // 预处理版, maxn大于1e6已经不可能处理了
```

```
8 const int maxb = 1005;
9
   int n, m;
10 int belong[maxn], lpos[maxb], rpos[maxb];
   int val[maxn], lazy[maxb];
12 int block;
13
14 scanf("%d", &n);
15 block = sqrt(n);
16 for (int i = 1; i <= n; ++i) {
17
        scanf("%d", &val[i]);
18
        belong[i] = (i - 1) / block + 1;
19 }
20 int sz = (n - 1) / block + 1;
21 for (int i = 1; i <= sz; ++i) {
        lpos[i] = 1 + (i - 1) * block;
22
23
        rpos[i] = i * block;
24 }
25 \quad rpos[sz] = n;
    3.4.4 弹飞绵羊
1 int n, m;
2 int belong[maxn], lpos[maxn], rpos[maxn];
3 int val[maxn], nxt[maxn], k[maxn], lst[maxn];
4 int block;
5
6
    void update(int pos) {
7
        int llim = lpos[belong[pos]], rlim = rpos[belong[pos]];
8
        for (int i = pos; i >= llim; --i) {
9
            if (val[i] + i > rlim) {
10
                k[i] = 1;
11
                nxt[i] = val[i] + i;
12
                if (val[i] + i > n) lst[i] = i;
13
                else lst[i] = lst[nxt[i]];
            } else {
14
15
                k[i] = 1 + k[val[i] + i];
16
                nxt[i] = nxt[val[i] + i];
17
                lst[i] = lst[val[i] + i];
18
            }
19
        }
20 }
21
22
    void init() {
        for (int i = n; i >= 1; --i) {
23
24
            int rlim = rpos[belong[i]];
            if (val[i] + i > rlim) {
25
26
                k[i] = 1;
27
                nxt[i] = val[i] + i;
28
                if (val[i] + i > n) lst[i] = i;
29
                else lst[i] = lst[nxt[i]];
30
            } else {
                k[i] = 1 + k[val[i] + i];
31
32
                nxt[i] = nxt[val[i] + i];
33
                lst[i] = lst[val[i] + i];
34
            }
35
        }
36 }
37
```

```
38
    int query(int pos) {
39
        int res = 0;
40
        while (pos <= n) {</pre>
41
            res += k[pos];
42
            if (nxt[pos] > n) printf("%d ", lst[pos]);
            pos = nxt[pos];
43
        }
44
45
        return res;
    }
46
47
    int main(int argc, char* argv[]) {
48
49
        scanf("%d%d", &n, &m);
50
        block = sqrt(n) * 1.6 + 1;
        for (int i = 1; i <= n; ++i) {</pre>
51
            scanf("%d", &val[i]);
52
53
            belong[i] = (i - 1) / block + 1;
        }
54
55
        int sz = (n - 1) / block + 1;
56
        for (int i = 1; i <= sz; ++i) {</pre>
            lpos[i] = 1 + (i - 1) * block;
57
            rpos[i] = i * block;
58
        }
59
60
        rpos[sz] = n;
61
        init();
62
        while (m--) {
63
            int op;
            scanf("%d", &op);
64
65
            if (op == 1) {
66
                int pos;
67
                scanf("%d", &pos);
68
                printf("%d\n", query(pos));
69
            } else {
70
                int pos, kl;
71
                scanf("%d%d", &pos, &kl);
72
                val[pos] = kl;
73
                update(pos);
74
            }
75
        }
76
        return 0;
77 }
    3.4.5 树莫队
1 // rnk保存欧拉序
   int sz[maxn], top[maxn], son[maxn], deep[maxn], fa[maxn], idx[maxn], ed[maxn], rnk[maxn*2];
3 int tot, n, m;
4 vector<int> edge[maxn];
5 int val[maxn];
6
   vector<int> xpos;
7
8
    inline void dfs1(int u, int pre, int h) {
9
        deep[u] = h;
10
        fa[u] = pre;
11
        sz[u] = 1;
12
        for (auto to : edge[u]) {
13
            if (to == pre) continue;
14
            dfs1(to, u, h + 1);
15
            sz[u] += sz[to];
```

```
16
            if (son[u] == 0 \mid \mid sz[to] > sz[son[u]]) son[u] = to;
        }
17
    }
18
19
20
    inline void dfs2(int u, int tp) {
21
        top[u] = tp;
22
        idx[u] = ++tot, rnk[tot] = u;
23
        if (son[u] == 0) {
24
            ed[u] = ++tot, rnk[tot] = u;
25
            return;
26
        }
27
        dfs2(son[u], tp);
28
        for (auto to : edge[u]) {
29
            if (to != son[u] && to != fa[u]) dfs2(to, to);
30
31
        ed[u] = ++tot, rnk[tot] = u;
32 }
33
34 inline int LCA(int x, int y) {
35
        while (top[x] != top[y]) {
36
            if (deep[top[x]] < deep[top[y]]) swap(x, y);</pre>
37
            x = fa[top[x]];
38
39
        if (deep[x] > deep[y]) swap(x, y);
40
        return x;
41 }
42
43 int belong[maxn*2], block;
    int res[maxn], ans;
   // 每个点是否访问(欧拉序去重)
46
   int vis[maxn];
47
   // 标记数组
48 int pre[maxn];
49 struct MO {
50
        int 1, r, id, lca;
51
        bool operator < (const MO& oth) const {</pre>
52
            return belong[1] == belong[oth.1] ? r < oth.r : belong[1] < belong[oth.1];</pre>
53
54 }q[maxm];
55
    inline void add(int x) {
56
57
        pre[x] ++;
58
        if (pre[x] == 1) ans ++;
59
60
    inline void del(int x) {
61
62
        pre[x] --;
63
        if (pre[x] == 0) ans --;
64 }
65
66
    inline void deal(int x) {
67
        vis[x] ? del(val[x]) : add(val[x]);
68
        vis[x] = !vis[x];
69 }
70
71
   int main(int argc, char* argv[]) {
72
        scanf("%d%d", &n, &m);
73
        block = sqrt(n);
74
        xpos.resize(n + 1);
```

```
75
         for (int i = 1; i <= n; ++i) {</pre>
 76
             scanf("%d", &val[i]);
 77
             xpos[i] = val[i];
 78
 79
         sort(xpos.begin(), xpos.end());
 80
         xpos.erase(unique(xpos.begin(), xpos.end()), xpos.end());
 81
         for (int i = 1; i <= n; ++i) val[i] = lower_bound(xpos.begin(), xpos.end(), val[i]) - xpos.</pre>
         begin();
 82
         // 欧拉序长度为n两倍所以分块要分两倍大小
 83
         for (int i = 1; i <= n * 2; ++i) {
             belong[i] = (i - 1) / block + 1;
 84
 85
         }
 86
         for (int i = 1, u, v; i < n; ++i) {</pre>
 87
             scanf("%d%d", &u, &v);
 88
             edge[u].push_back(v);
 89
             edge[v].push_back(u);
 90
         // 树剖预处理1ca
 91
 92
         dfs1(1, 0, 0);
 93
         dfs2(1, 1);
 94
         for (int i = 1, x, y; i <= m; ++i) {</pre>
 95
             scanf("%d%d", &x, &y);
 96
             if (idx[x] > idx[y]) swap(x, y);
 97
             int _lca = LCA(x, y);
 98
             q[i].id = i;
99
             if (_lca == x) q[i].l = idx[x], q[i].r = idx[y], q[i].lca = 0;
100
             else q[i].l = ed[x], q[i].r = idx[y], q[i].lca = _lca;
               \texttt{cerr} << q[i].1 << " " << q[i].r << " " << q[i].id << " " << q[i].lca << endl;
101
    //
         }
102
103
         sort(q + 1, q + 1 + m);
104
         int 1 = 1, r = 0;
105
         for (int i = 1; i <= m; ++i) {</pre>
             while(r < q[i].r) deal(rnk[++r]);</pre>
106
             while(r > q[i].r) deal(rnk[r--]);
107
             while(l < q[i].l) deal(rnk[l++]);</pre>
108
109
             while(1 > q[i].1) deal(rnk[--1]);
             if (q[i].lca) deal(q[i].lca);
110
111
             res[q[i].id] = ans;
112
             if (q[i].lca) deal(q[i].lca);
         }
113
         for (int i = 1; i <= m; ++i) {</pre>
114
             printf("%d\n", res[i]);
115
116
117
         return 0;
118 }
          VirtualTree
     3.5
     3.5.1 VirtualTree
  1 const int pow2 = 19;
  2 const int maxn = 1 << pow2;</pre>
  3 vector<int> adj0[maxn], adj1[maxn];
  4 int st[maxn << 1][pow2 + 1], dep[maxn], euler[maxn], euler_clock;</pre>
 5 // fa0 是原树的父节点
 6 // fa1 是虚树的父节点
```

7 // len 是虚树每个节点的权重,每个节点代表原树的几个节点,也是虚树到它父节点的链的长度

8 int stk[maxn], fa0[maxn], fa1[maxn], len[maxn];

9 11 val[maxn];

```
10
11
   void link0(int u, int v) { adj0[u].emplace_back(v); adj0[v].emplace_back(u); }
   void link1(int u, int v) { adj1[u].emplace_back(v); adj1[v].emplace_back(u); }
12
13
   void dfs0(int u, int p) {
        fa0[u] = p;
14
15
        dep[u] = dep[p] + 1;
16
        st[++euler_clock][0] = u;
17
        euler[u] = euler_clock;
        for (const auto& v : adj0[u]) if (v != p) {
18
19
                dfs0(v, u);
20
                st[++euler_clock][0] = u;
21
            }
22 }
23 inline bool cmp(int u, int v) {return dep[u] < dep[v];}
24
   inline int upper(int u, int v) {return cmp(u, v) ? u : v;}
   void lca_init() {
25
        for (int i = 0; i != 31 - __builtin_clz(euler_clock); ++i)
26
27
            for (int j = 1; j + (1 << (i + 1)) <= euler_clock; ++j)
28
                st[j][i + 1] = upper(st[j][i], st[j + (1 << i)][i]);
29 }
30 inline int lca(int u, int v) {
        if (u == v) return u;
31
32
        u = euler[u];
33
        v = euler[v];
34
        if (u > v) swap(u, v);
        int temp = 31 - __builtin_clz(++v - u);
35
36
        return upper(st[u][temp], st[v - (1 << temp)][temp]);</pre>
37 }
   void build(vector<int>& key) {
38
        sort(key.begin(), key.end(), [&] (int u, int v) { return euler[u] < euler[v]; });</pre>
39
40
        key.resize(unique(key.begin(), key.end()) - key.begin());
41
        int top = 0;
42
        for (const auto& u : key) {
43
            if (!top) {
                stk[++top] = u;
44
45
                continue;
46
            }
47
            int p = lca(u, stk[top]);
            while (euler[p] < euler[stk[top]]) {</pre>
48
                if (euler[p] >= euler[stk[top - 1]]) {
49
                    link1(p, stk[top]);
50
51
                    if (stk[--top] != p) stk[++top] = p;
52
53
                link1(stk[top - 1], stk[top]);
54
55
                --top;
            }
56
            stk[++top] = u;
57
        }
58
59
        while (top > 1) {
60
            link1(stk[top - 1], stk[top]);
61
            --top;
62
        }
63 }
64
65
   void dfs1(int u, int p) {
66
        fa1[u] = p;
67
        val[u] = 0;
68
        len[u] = dep[u] - dep[p];
```

```
69
        for (const auto& v : adj1[u]) if (v != p) dfs1(v, u);
70 }
71
   int main() {
72
73
        // 多组清空操作
74
        for (int i = 1; i <= n; ++i) {</pre>
75
            adj0[i].clear();
76
            adj1[i].clear();
       }
77
78
        euler_clock = 0;
79
80
       // 读入原树 link0 加边
81
        // 读入处理关键节点存入vector key, 包含1和链的端点和他们的lca的父节点(lca如果为1就不加)。
82
        dfs0(1, 0);
83
       lca_init();
84
85
       vector<int> key(1, 1);
86
        for (auto& q : query) {
87
           cin >> q.u >> q.v;
88
           key.emplace_back(q.u);
89
           key.emplace_back(q.v);
90
           int p = lca(q.u, q.v);
91
           if (p != 1) key.emplace_back(fa0[p]);
92
93
94
        build(key);
95
        dfs1(1, 0);
        return 0;
96
97 }
    3.6 PersistentDS
    3.6.1 主席树区间 k 大
1 // const int maxn = 100005;
2 int n, m;
3 int a[maxn];
 4 int root[maxn];
5 int cnt = 0;
6 vector<int> b;
7
   struct node {
8
        int 1, r, val;
9
   p[maxn * 40];
10
11
   void update(int 1, int r, int pre, int &now, int pos) {
12
        now = ++cnt;
13
       p[now] = p[pre];
14
       p[now].val++;
15
        if (1 == r) {
16
           return;
        }
17
18
        int mid = 1 + r >> 1;
19
        if (pos <= mid) update(1, mid, p[pre].1, p[now].1, pos);</pre>
20
        else update(mid + 1, r, p[pre].r, p[now].r, pos);
21 }
22
23 int query(int 1, int r, int x, int y, int k) {
```

24

25

if (1 == r) return b[1 - 1];

int mid = 1 + r >> 1;

```
26
        int temp = p[p[y].1].val - p[p[x].1].val;
27
        if (k <= temp) return query(1, mid, p[x].1, p[y].1, k);</pre>
28
        return query(mid + 1, r, p[x].r, p[y].r, k - temp);
29 }
30
31
   int main(int argc,char *argv[])
32
    {
33
        while (scanf("%d%d", &n, &m) != EOF) {
34
            b.clear();
35
            cnt = 0;
            for (int i = 1; i <= n; ++i) scanf("%d", &a[i]), b.push_back(a[i]);</pre>
36
37
            sort(b.begin(), b.end());
38
            b.erase(unique(b.begin(), b.end()), b.end());
39
            for (int i = 1; i <= n; ++i) {</pre>
40
                update(1, b.size(), root[i - 1], root[i], lower_bound(b.begin(), b.end(), a[i]) - b.
        begin() + 1);
41
            }
42
            int L, R, k;
            while (m--) {
43
                scanf("%d%d%d", &L, &R, &k);
44
45
                printf("%d\n", query(1, b.size(), root[L - 1], root[R], k));
            }
46
        }
47
48
        return 0;
49 }
    3.6.2 可持久化数组
 1 /*1、操作将u, v合并 2、操作回退 */
    const int maxn = 2e5+5;
    int n, m, sz;
4 int root [maxn], ls [maxn*40], rs [maxn*40], v [maxn*40], deep [maxn*40];
5 int has[maxn];
6
7
    void build(int &k, int 1, int r) {
8
        if (!k)k = ++sz;
9
        if (1 == r) {
10
            v[k] = 1;
11
            return;
12
        }
13
        int mid = (1 + r) >> 1;
14
        build(ls[k], 1, mid);
15
        build(rs[k], mid + 1, r);
16
   }
17
18
    void modify(int 1, int r, int x, int &y, int pos, int val) {
19
        y = ++sz;
20
        if (1 == r) {
21
            v[y] = val;
22
            deep[y] = deep[x];
23
            return;
24
        }
25
        ls[y] = ls[x];
26
        rs[y] = rs[x];
27
        int mid = (1 + r) >> 1;
28
        if (pos <= mid)</pre>
29
            modify(1, mid, ls[x], ls[y], pos, val);
30
        else modify(mid + 1, r, rs[x], rs[y], pos, val);
```

```
31 }
32
33
    int query(int k, int l, int r, int pos) {
34
        if (1 == r)return k;
35
        int mid = (1 + r) >> 1;
36
        if (pos <= mid)return query(ls[k], 1, mid, pos);</pre>
37
        else return query(rs[k], mid + 1, r, pos);
    }
38
39
40
    void add(int k, int l, int r, int pos) {
        if (1 == r) {
41
42
            deep[k]++;
43
            return;
        }
44
45
        int mid = (1 + r) >> 1;
46
        if (pos <= mid)add(ls[k], l, mid, pos);</pre>
47
        else add(rs[k], mid + 1, r, pos);
48
   }
49
50
    int find(int k, int x) {
51
        int p = query(k, 1, n, x);
52
        if (x == v[p])return p;
53
        return find(k, v[p]);
54 }
55
56
    int main() {
57
        int T = read();
58
        while (T--) {
59
            sz = 0;
60
            memset(root, 0, sizeof root);
61
            memset(ls, 0, sizeof ls);
62
            memset(rs, 0, sizeof rs);
63
            n = read();
64
            has[0] = n;
65
            m = read();
66
            build(root[0], 1, n);
67
            int f, k, a, b;
68
            for (int i = 1; i <= m; i++) {</pre>
                 f = read();
69
70
                 if (f == 1) {
71
                     root[i] = root[i - 1];
                     has[i] = has[i - 1];
72
73
                     a = read();
74
                     b = read();
                     int p = find(root[i], a), q = find(root[i], b);
75
76
                     if (v[p] == v[q])continue;
77
                     has[i]--;
                     if (deep[p] > deep[q])swap(p, q);
78
79
                     modify(1, n, root[i - 1], root[i], v[p], v[q]);
80
                     if (deep[p] == deep[q])add(root[i], 1, n, v[q]);
81
                 } else if (f == 2) {
82
                     k = read();
83
                     root[i] = root[k];
84
                     has[i] = has[k];
85
86
                 printf("%d\n", has[i]);
87
            }
88
        }
89
        return 0;
```

90 }

3.7 Tree

```
3.7.1 LCA
```

```
1
   // const int maxn = 1e5 + 10;
2
3 // 普通倍增1ca
4 int n, dep[maxn], fa[maxn][30];
5 vector<int> edge[maxn];
7
   void dfs(int u, int pre) {
        dep[u] = dep[pre] + 1, fa[u][0] = pre;
8
9
        for(int i = 1; (1 << i) <= n; i ++)</pre>
10
            fa[u][i] = fa[fa[u][i - 1]][i - 1];
11
        for(auto v : edge[u]) if(v != pre) dfs(v, u);
   }
12
13
14
   int LCA(int u, int v) {
        if(dep[u] < dep[v]) swap(u, v);</pre>
15
16
        int d = dep[u] - dep[v];
        for(int i = 0; (1 << i) <= d; i ++)</pre>
17
18
            if((1 << i) & d) u = fa[u][i];
19
        if(u == v) return u;
        for(int i = 20; i >= 0; i --)
20
21
            if(fa[u][i] != fa[v][i])
22
                u = fa[u][i], v = fa[v][i];
23
        return fa[u][0];
24 }
25
26
27 // 欧拉序1ca
28 // pow2 = 19
29 // maxn = 1 << pow2
30 int st[maxn << 1][pow2 + 1], dep[maxn], euler[maxn], euler_clock, fa[maxn];</pre>
31 void dfs(int u, int p) {
32
        fa[u] = p;
        dep[u] = dep[p] + 1;
33
34
        st[++euler_clock][0] = u;
35
        euler[u] = euler_clock;
36
        for (const auto& v : adj0[u]) if (v != p) {
37
                dfs(v, u);
38
                st[++euler_clock][0] = u;
            }
39
40 }
41
   void lca_init() {
42
        for (int i = 0; i != 31 - __builtin_clz(euler_clock); ++i)
43
            for (int j = 1; j + (1 << (i + 1)) <= euler_clock; ++j)</pre>
44
                st[j][i + 1] = upper(st[j][i], st[j + (1 << i)][i]);
45
  }
46
   inline int lca(int u, int v) {
47
        if (u == v) return u;
        u = euler[u];
48
49
        v = euler[v];
50
        if (u > v) swap(u, v);
        int temp = 31 - __builtin_clz(++v - u);
52
        return upper(st[u][temp], st[v - (1 << temp)][temp]);</pre>
53 }
```

```
54
55 // dfs(1, 0);
   // lca_init();
57
58 // 另有树剖lca详见hld模板
    3.7.2 前向星
1 // 清零 head 和 tot
    const int maxm = 4e5+5;
    int ver[maxm], Next[maxm], head[maxn], edge[maxm];
    void addEdge(int u, int v, int w){
4
5
        ver[++tot]=v;
6
        Next[tot] = head[u];
7
        head[u]=tot;
8
        edge[tot]=w;
9
   }
10
11 for(int i = head[u]; i; i=Next[i])
    3.7.3 点分治
1 int n, k;
2
3 // 清零 head 和 tot
4 const int maxm = 2e4+5;
5 int ver[maxm], Next[maxm], head[maxn], edge[maxm];
6
   int tot;
7
    void addEdge(int u, int v, int w){
8
        ver[++tot]=v;
9
        Next[tot] = head[u];
10
        head[u]=tot;
11
        edge[tot]=w;
12 }
13
14 int sz[maxn], vis[maxn];
15 int rt, mxsz, has;
16
17
    void getrt(int u, int pre) {
18
        sz[u] = 1;
19
        int mxnow = 0;
20
        for (int i = head[u]; i; i = Next[i]) {
21
            int v = ver[i];
22
            if (v == pre || vis[v]) continue;
23
            getrt(v, u);
24
            sz[u] += sz[v];
25
            mxnow = max(mxnow, sz[v]);
26
        }
27
        mxnow = max(mxnow, has - sz[u]);
28
        if (mxnow < mxsz) {</pre>
29
            mxsz = mxnow, rt = u;
30
        }
31 }
32
33 int dl[maxn], r;
34
   int val[maxn];
35
36 void getdis(int u, int pre) {
```

```
37
        dl[r++] = val[u];
        for (int i = head[u]; i; i = Next[i]) {
38
39
            int v = ver[i];
            if (v == pre || vis[v]) continue;
40
41
            val[v] = val[u] + edge[i];
42
            getdis(v, u);
43
        }
   }
44
45
    11 cal(int u, int pre) {
47
        r = 0;
48
        val[u] = pre;
49
        getdis(u, 0);
        11 sum = 0;
50
        sort(dl, dl + r);
51
52
        r --;
53
        int 1 = 0;
54
        while (1 < r) {
55
            if (dl[1] + dl[r] > k) r --;
56
            else sum += r - 1, 1 ++;
57
        }
58
        return sum;
59 }
60
61 11 res = 0;
62
   void dfs(int u) {
63
        res += cal(u, 0);
        vis[u] = 1;
64
        for (int i = head[u]; i; i = Next[i]) {
65
66
            int v = ver[i];
67
            if (vis[v]) continue;
            res -= cal(v, edge[i]);
68
69
            has = sz[v];
70
            mxsz = 0x3f3f3f3f;
71
            getrt(v, 0);
72
            dfs(rt);
73
        }
74 }
75
76
    int main(int argc, char* argv[]) {
77
        while (scanf("%d%d", &n, &k) != EOF && (n || k)) {
            tot = 0; memset(head, 0, sizeof head);
78
79
            memset(vis, 0, sizeof vis);
80
            res = 0;
            for (int i = 1, u, v, w; i < n; ++i) {</pre>
81
82
                scanf("%d%d%d", &u, &v, &w);
83
                addEdge(u, v, w);
84
                addEdge(v, u, w);
            }
85
86
            mxsz = 0x3f3f3f3f;
87
            has = n;
88
            getrt(1, 0);
89
            dfs(rt);
90
            printf("%lld\n", res);
91
        }
92
        return 0;
93 }
```

3.8 Others

3.8.1 BITinNM

```
1
   struct Fenwick_Tree {
    #define type int
3
        type bit[maxn][maxn];
4
        int n, m;
5
        void init(int _n, int _m) {
6
            n = _n;
7
            m = _m;
8
            mem(bit, 0);
9
10
       int lowbit(int x) { return x & (-x); }
        void update(int x, int y, type v) {
11
12
            int i, j;
13
            for (i = x; i <= n; i += lowbit(i)) {</pre>
14
                for (j = y; j <= m; j += lowbit(j)) {</pre>
15
                    bit[i][j] += v;
16
            }
17
18
        }
        type get(int x, int y) {
19
20
            type i, j, res = 0;
21
            for (i = x; i > 0; i -= lowbit(i)) {
                for (j = y; j > 0; j -= lowbit(j)) {
22
23
                    res += bit[i][j];
24
25
            }
26
            return res;
27
28
        type query(int x1, int x2, int y1, int y2) {
29
            x1--;
30
            y1--;
31
            return get(x2, y2) - get(x1, y2) - get(x2, y1) + get(x1, y1);
32
        }
33 #undef type
34 } tr;
36 // 二维区间前缀和写法(非树状数组)
   inline void range_add(int xa, int ya, int xb, int yb) { add(xa, ya, 1), add(xa, yb + 1, -1), add(xb
         + 1, ya, -1), add(xb + 1, yb + 1, 1); }
38
   inline 11 range_ask(int xa, int ya, int xb, int yb){ return ask(xb, yb) - ask(xb, ya - 1) - ask(xa
        - 1, yb) + ask(xa - 1, ya - 1); }
39
   inline void build() {
40
        for (int i = 1; i < n + 5; ++i) {</pre>
41
            for (int j = 1; j < m + 5; ++j) {
                if (st[i][j] > 1) st[i][j] = 1;
42
43
                st[i][j] += st[i - 1][j] + st[i][j - 1] - st[i - 1][j - 1];
44
            }
45
        }
46 }
47
48 // 二维树状数组区间加与求和
  11 t1[maxn] [maxn], t2[maxn] [maxn], t3[maxn] [maxn], t4[maxn] [maxn];
50
   void add(ll x, ll y, ll z){
51
        for(int X = x; X <= n; X += X & -X)</pre>
52
            for(int Y = y; Y <= m; Y += Y & -Y){
                t1[X][Y] += z;
53
```

```
54
               t2[X][Y] += z * x;
55
               t3[X][Y] += z * y;
56
               t4[X][Y] += z * x * y;
           }
57
58
   }
59
   11 ask(ll x, ll y){
60
       11 \text{ res} = 0;
61
        for(int i = x; i; i -= i & -i)
62
           for(int j = y; j; j -= j & -j)
63
               res += (x + 1) * (y + 1) * t1[i][j]
                    - (y + 1) * t2[i][j]
64
65
                    -(x + 1) * t3[i][j]
66
                    + t4[i][j];
67
        return res;
68
   }
69
70 // 区间加, 询问单点: 直接维护前缀差分数组, 求单点=普通求前缀和
```

3.8.2 静态区间 k 大划分树

```
1 // const int maxn = 100010;
2 int tree[20][maxn];
3 // 读入sorted并排序, 赋值给tree的第0层
4 int sorted[maxn];
5 int toleft[20][maxn];
6 // 保存左子树的和
7
   // 11 sum[20][maxn];
8
9
  // 1, n, 0
10
   void build(int 1, int r, int dep) {
        if (1 == r) return;
11
12
        // sum[dep][0] = 0;
13
        toleft[dep][0] = 0;
        int mid = 1 + r >> 1;
14
        int same = mid - 1 + 1;
15
16
        for (int i = 1; i <= r; ++i) {</pre>
17
            if (tree[dep][i] < sorted[mid]) same--;</pre>
18
19
        int lpos = 1, rpos = mid + 1;
20
        for (int i = 1; i <= r; ++i) {</pre>
21
            // sum[dep][i] = sum[dep][i - 1];
22
            if (tree[dep][i] < sorted[mid]) {</pre>
23
                // sum[dep][i] += tree[dep][i];
24
                tree[dep + 1][lpos++] = tree[dep][i];
25
            }
26
            else if (tree[dep][i] == sorted[mid] && same > 0) {
27
                // sum[dep][i] += tree[dep][i];
28
                tree[dep + 1][lpos++] = tree[dep][i];
29
                same --;
30
            } else tree[dep + 1][rpos ++] = tree[dep][i];
31
            toleft[dep][i] = toleft[dep][l - 1] + lpos - 1;
32
33
        build(1, mid, dep + 1);
34
        build(mid + 1, r, dep + 1);
35
36
37 //(1~k-1)的数的和,注意每次查询前初始化
38 // 11 \text{ ress} = 0;
```

```
39
   // L = 1, R = n, dep = 0, 1,r是查询区间
    int query(int L, int R, int 1, int r, int dep, int k) {
42
        if (1 == r) return tree[dep][1];
43
        int mid = (L + R) >> 1;
44
        int cnt = toleft[dep][r] - toleft[dep][l - 1];
45
        if (cnt >= k) {
            int newl = L + toleft[dep][l - 1] - toleft[dep][L - 1];
46
47
            int newr = newl + cnt - 1;
48
            return query(L, mid, newl, newr, dep + 1, k);
49
        } else {
50
            int newr = r + toleft[dep][R] - toleft[dep][r];
51
            int newl = newr - (r - 1 - cnt);
52
            // ress += sum[dep][r] - sum[dep][l - 1];
53
            return query(mid + 1, R, newl, newr, dep + 1, k - cnt);
        }
54
55
   }
56
57
58 \quad \text{scan(n), scan(m);}
59 for (int i = 1; i <= n; ++i) {
        scan(sorted[i]);
60
        tree[0][i] = sorted[i];
61
63 sort(sorted + 1, sorted + 1 + n);
64 build(1, n, 0);
65 int 1, r, k;
66 while (m--) {
        scan(1), scan(r), scan(k);
67
68
        printf("%d\n", query(1, n, 1, r, 0, k));
69 }
```

4 Geometry

4.1 Class

4.1.1 geo

```
1 #define mp make_pair
2 #define fi first
3 #define se second
4 #define pb push_back
5 typedef double db;
6 const db eps=1e-6;
7 const db pi=acos(-1);
8 int sign(db k){
9
        if (k>eps) return 1; else if (k<-eps) return -1; return 0;
10 }
int cmp(db k1,db k2){return sign(k1-k2);}
12 int inmid(db k1,db k2,db k3){return sign(k1-k3)*sign(k2-k3)<=0;}// k3 在 [k1,k2] 内
13 struct point{
14
        db x,y;
15
        point operator + (const point &k1) const{return (point){k1.x+x,k1.y+y};}
16
       point operator - (const point &k1) const{return (point){x-k1.x,y-k1.y};}
       point operator * (db k1) const{return (point){x*k1,y*k1};}
17
18
       point operator / (db k1) const{return (point){x/k1,y/k1};}
19
       int operator == (const point &k1) const{return cmp(x,k1.x)==0\&&cmp(y,k1.y)==0;}
20
        // 逆时针旋转
21
       point turn(db k1){return (point){x*cos(k1)-y*sin(k1),x*sin(k1)+y*cos(k1)};}
22
       point turn90(){return (point){-y,x};}
23
       bool operator < (const point k1) const{</pre>
24
            int a=cmp(x,k1.x);
25
            if (a==-1) return 1; else if (a==1) return 0; else return cmp(y,k1.y)==-1;
26
27
        db abs(){return sqrt(x*x+y*y);}
28
        db abs2(){return x*x+y*y;}
29
        db dis(point k1){return ((*this)-k1).abs();}
30
       point unit(){db w=abs(); return (point){x/w,y/w};}
31
       void scan(){double k1,k2; scanf("%lf%lf",&k1,&k2); x=k1; y=k2;}
32
        void print(){printf("%.11lf %.11lf\n",x,y);}
33
        db getw(){return atan2(y,x);}
34
       point getdel(){if (sign(x)=-1||(sign(x)==0\&\&sign(y)==-1)) return (*this)*(-1); else return (*
        this);}
35
        int getP() const{return sign(y)==1||(sign(y)==0&&sign(x)==-1);}
36 };
   int inmid(point k1,point k2,point k3){return inmid(k1.x,k2.x,k3.x)&&inmid(k1.y,k2.y,k3.y);}
38 db cross(point k1,point k2){return k1.x*k2.y-k1.y*k2.x;}
39 db dot(point k1,point k2){return k1.x*k2.x+k1.y*k2.y;}
40 db rad(point k1,point k2){return atan2(cross(k1,k2),dot(k1,k2));}
41 // -pi -> pi
42 int compareangle (point k1, point k2){
43
        return k1.getP()<k2.getP()||(k1.getP()==k2.getP()&&sign(cross(k1,k2))>0);
44 }
45 point proj(point k1, point k2, point q){ // q 到直线 k1, k2 的投影
46
       point k=k2-k1; return k1+k*(dot(q-k1,k)/k.abs2());
47 }
   point reflect(point k1,point k2,point q){return proj(k1,k2,q)*2-q;}
49 int clockwise(point k1,point k2,point k3){// k1 k2 k3 逆时针 1 顺时针 -1 否则 0
50
        return sign(cross(k2-k1,k3-k1));
52 int checkLL(point k1, point k2, point k3, point k4) {// 求直线 (L) 线段 (S)k1,k2 和 k3,k4 的交点
```

```
53
         return cmp(cross(k3-k1,k4-k1),cross(k3-k2,k4-k2))!=0;
54 }
55
    point getLL(point k1,point k2,point k3,point k4){
56
         db w1=cross(k1-k3,k4-k3),w2=cross(k4-k3,k2-k3); return (k1*w2+k2*w1)/(w1+w2);
57
58 int intersect(db 11,db r1,db 12,db r2){
59
         if (11>r1) swap(11,r1); if (12>r2) swap(12,r2); return cmp(r1,12)!=-1&&cmp(r2,11)!=-1;
60 }
61
    int checkSS(point k1,point k2,point k3,point k4){
62
         return intersect(k1.x,k2.x,k3.x,k4.x)&&intersect(k1.y,k2.y,k3.y,k4.y)&&
         sign(cross(k3-k1,k4-k1))*sign(cross(k3-k2,k4-k2))<=0&&
63
64
         sign(cross(k1-k3,k2-k3))*sign(cross(k1-k4,k2-k4))<=0;
65 }
66
    db disSP(point k1,point k2,point q){
67
         point k3=proj(k1,k2,q);
68
         if (inmid(k1,k2,k3)) return q.dis(k3); else return min(q.dis(k1),q.dis(k2));
69 }
70 db disSS(point k1,point k2,point k3,point k4){
71
         if (checkSS(k1,k2,k3,k4)) return 0;
72
         else return min(min(disSP(k1,k2,k3),disSP(k1,k2,k4)),min(disSP(k3,k4,k1),disSP(k3,k4,k2)));
73 }
74 int onS(point k1,point k2,point q){return inmid(k1,k2,q)&&sign(cross(k1-q,k2-k1))==0;}
75 struct circle{
         point o; db r;
76
77
         void scan(){o.scan(); scanf("%lf",&r);}
78
         int inside(point k){return cmp(r,o.dis(k));}
79 };
80
    struct line{
81
         // p[0]->p[1]
82
         point p[2];
         line(point k1,point k2){p[0]=k1; p[1]=k2;}
83
84
         point& operator [] (int k){return p[k];}
85
         int include(point k){return sign(cross(p[1]-p[0],k-p[0]))>0;}
86
         point dir(){return p[1]-p[0];}
87
         line push(){ // 向外 ( 左手边 ) 平移 eps
88
             const db eps = 1e-6;
89
             point delta=(p[1]-p[0]).turn90().unit()*eps;
90
             return {p[0]-delta,p[1]-delta};
91
         }
92 };
93
    point getLL(line k1,line k2){return getLL(k1[0],k1[1],k2[0],k2[1]);}
    int parallel(line k1,line k2){return sign(cross(k1.dir(),k2.dir()))==0;}
    int sameDir(line k1,line k2){return parallel(k1,k2)&&sign(dot(k1.dir(),k2.dir()))==1;}
    int operator < (line k1,line k2){</pre>
97
         if (sameDir(k1,k2)) return k2.include(k1[0]);
98
         return compareangle(k1.dir(),k2.dir());
99 }
int checkpos(line k1,line k2,line k3){return k3.include(getLL(k1,k2));}
    vector<line> getHL(vector<line> &L){ // 求半平面交 , 半平面是逆时针方向 , 输出按照逆时针
102
         sort(L.begin(),L.end()); deque<line> q;
103
         for (int i=0;i<(int)L.size();i++){</pre>
104
             if (i&&sameDir(L[i],L[i-1])) continue;
105
             while (q.size()>1&&!checkpos(q[q.size()-2],q[q.size()-1],L[i])) q.pop_back();
106
             while (q.size()>1&&!checkpos(q[1],q[0],L[i])) q.pop_front();
107
             q.push_back(L[i]);
108
         }
109
         while (q.size()>2\&\&!checkpos(q[q.size()-2],q[q.size()-1],q[0])) q.pop_back();
110
         while (q.size()>2\&\&!checkpos(q[1],q[0],q[q.size()-1])) q.pop_front();
111
         vector<line>ans; for (int i=0;i<q.size();i++) ans.push_back(q[i]);</pre>
```

```
112
         return ans;
113 }
114
    db closepoint(vector<point>&A,int l,int r){ // 最近点对 , 先要按照 x 坐标排序
115
         if (r-1<=5){</pre>
116
             db ans=1e20;
117
             for (int i=1;i<=r;i++) for (int j=i+1;j<=r;j++) ans=min(ans,A[i].dis(A[j]));</pre>
118
             return ans;
         }
119
         int mid=l+r>>1; db ans=min(closepoint(A,1,mid),closepoint(A,mid+1,r));
120
121
         vector<point>B; for (int i=1;i<=r;i++) if (abs(A[i].x-A[mid].x)<=ans) B.push_back(A[i]);</pre>
122
         sort(B.begin(),B.end(),[](point k1,point k2){return k1.y<k2.y;});</pre>
123
         for (int i=0;i<B.size();i++) for (int j=i+1;j<B.size()&&B[j].y-B[i].y<ans;j++) ans=min(ans,B[i</pre>
         ].dis(B[j]));
124
         return ans;
125 }
126 int checkposCC(circle k1, circle k2){// 返回两个圆的公切线数量
127
         if (cmp(k1.r,k2.r)=-1) swap(k1,k2);
         db dis=k1.o.dis(k2.o); int w1=cmp(dis,k1.r+k2.r),w2=cmp(dis,k1.r-k2.r);
128
129
         if (w1>0) return 4; else if (w1==0) return 3; else if (w2>0) return 2;
         else if (w2==0) return 1; else return 0;
130
131 }
132
    vector<point> getCL(circle k1,point k2,point k3){ // 沿着 k2->k3 方向给出 , 相切给出两个
133
         point k=proj(k2,k3,k1.o); db d=k1.r*k1.r-(k-k1.o).abs2();
134
         if (sign(d)==-1) return {};
135
         point del=(k3-k2).unit()*sqrt(max((db)0.0,d)); return {k-del,k+del};
136 }
137
    vector<point> getCC(circle k1,circle k2){// 沿圆 k1 逆时针给出 , 相切给出两个
138
         int pd=checkposCC(k1,k2); if (pd==0||pd==4) return {};
139
         db a=(k2.o-k1.o).abs2(),cosA=(k1.r*k1.r+a-k2.r*k2.r)/(2*k1.r*sqrt(max(a,(db)0.0)));
140
         db b=k1.r*cosA,c=sqrt(max((db)0.0,k1.r*k1.r-b*b));
141
         point k=(k2.o-k1.o).unit(),m=k1.o+k*b,del=k.turn90()*c;
142
         return {m-del,m+del};
143 }
    vector<point> TangentCP(circle k1, point k2){// 沿圆 k1 逆时针给出
144
         db a=(k2-k1.o).abs(),b=k1.r*k1.r/a,c=sqrt(max((db)0.0,k1.r*k1.r-b*b));
145
146
         point k=(k2-k1.o).unit(),m=k1.o+k*b,del=k.turn90()*c;
147
         return {m-del,m+del};
148 }
149 vector<line> TangentoutCC(circle k1,circle k2){
         int pd=checkposCC(k1,k2); if (pd==0) return {};
150
151
         if (pd==1){point k=getCC(k1,k2)[0]; return {(line){k,k}};}
152
         if (cmp(k1.r,k2.r)==0){
153
             point del=(k2.o-k1.o).unit().turn90().getdel();
154
             return {(line){k1.o-del*k1.r,k2.o-del*k2.r},(line){k1.o+del*k1.r,k2.o+del*k2.r}};
155
156
             point p=(k2.0*k1.r-k1.0*k2.r)/(k1.r-k2.r);
157
             vector<point>A=TangentCP(k1,p),B=TangentCP(k2,p);
             vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});</pre>
158
159
             return ans;
160
         }
161 }
162
    vector<line> TangentinCC(circle k1,circle k2){
163
         int pd=checkposCC(k1,k2); if (pd<=2) return {};</pre>
164
         if (pd==3){point k=getCC(k1,k2)[0]; return {(line){k,k}};}
165
         point p=(k2.o*k1.r+k1.o*k2.r)/(k1.r+k2.r);
166
         vector<point>A=TangentCP(k1,p),B=TangentCP(k2,p);
167
         vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});</pre>
168
         return ans;
169 }
```

```
170 vector<line> TangentCC(circle k1,circle k2){
171
         int flag=0; if (k1.r<k2.r) swap(k1,k2),flag=1;</pre>
172
         vector<line>A=TangentoutCC(k1,k2),B=TangentinCC(k1,k2);
173
         for (line k:B) A.push_back(k);
174
         if (flag) for (line &k:A) swap(k[0],k[1]);
175
         return A;
176 }
177
    db getarea(circle k1,point k2,point k3){
178
         // 圆 k1 与三角形 k2 k3 k1.o 的有向面积交
179
         point k=k1.o; k1.o=k1.o-k; k2=k2-k; k3=k3-k;
         int pd1=k1.inside(k2),pd2=k1.inside(k3);
180
181
         vector<point>A=getCL(k1,k2,k3);
182
         if (pd1>=0){
183
             if (pd2>=0) return cross(k2,k3)/2;
184
             return k1.r*k1.r*rad(A[1],k3)/2+cross(k2,A[1])/2;
185
         } else if (pd2>=0){
186
             return k1.r*k1.r*rad(k2,A[0])/2+cross(A[0],k3)/2;
187
         }else {
188
             int pd=cmp(k1.r,disSP(k2,k3,k1.o));
189
             if (pd<=0) return k1.r*k1.r*rad(k2,k3)/2;</pre>
190
             return cross(A[0],A[1])/2+k1.r*k1.r*(rad(k2,A[0])+rad(A[1],k3))/2;
         }
191
192 }
193
    circle getcircle(point k1,point k2,point k3){
194
         db a1=k2.x-k1.x,b1=k2.y-k1.y,c1=(a1*a1+b1*b1)/2;
195
         db a2=k3.x-k1.x,b2=k3.y-k1.y,c2=(a2*a2+b2*b2)/2;
196
         db d=a1*b2-a2*b1;
197
         point o=(point)\{k1.x+(c1*b2-c2*b1)/d,k1.y+(a1*c2-a2*c1)/d\};
198
         return (circle){o,k1.dis(o)};
199
     circle getScircle(vector<point> A){
200
201
         random_shuffle(A.begin(), A.end());
202
         circle ans=(circle){A[0],0};
203
         for (int i=1;i<A.size();i++)</pre>
204
             if (ans.inside(A[i])==-1){
205
                 ans=(circle){A[i],0};
206
                 for (int j=0;j<i;j++)</pre>
207
                      if (ans.inside(A[j])==-1){
208
                          ans.o=(A[i]+A[j])/2; \ ans.r=ans.o.dis(A[i]);
209
                          for (int k=0;k<j;k++)</pre>
210
                              if (ans.inside(A[k])==-1)
211
                                  ans=getcircle(A[i],A[j],A[k]);
212
                     }
213
             }
214
         return ans;
215 }
216 db area(vector<point> A){ // 多边形用 vector<point> 表示 , 逆时针
217
         db ans=0;
218
         for (int i=0;i<A.size();i++) ans+=cross(A[i],A[(i+1)%A.size()]);</pre>
219
         return ans/2;
220 }
221
    int checkconvex(vector<point>A){
222
         int n=A.size(); A.push_back(A[0]); A.push_back(A[1]);
223
         for (int i=0;i<n;i++) if (sign(cross(A[i+1]-A[i],A[i+2]-A[i]))==-1) return 0;</pre>
224
         return 1;
225
226
    int contain(vector<point>A,point q){ // 2 内部 1 边界 0 外部
227
         int pd=0; A.push_back(A[0]);
228
         for (int i=1;i<A.size();i++){</pre>
```

```
229
             point u=A[i-1],v=A[i];
230
             if (onS(u,v,q)) return 1; if (cmp(u.y,v.y)>0) swap(u,v);
231
             if (cmp(u.y,q.y) \ge 0 | | cmp(v.y,q.y) < 0) continue;
232
             if (sign(cross(u-v,q-v))<0) pd^=1;</pre>
233
234
         return pd<<1;</pre>
235 }
236
     vector<point> ConvexHull(vector<point>A,int flag=1){ // flag=0 不严格 flag=1 严格
237
         int n=A.size(); vector<point>ans(n*2);
238
         sort(A.begin(), A.end()); int now=-1;
239
         for (int i=0;i<A.size();i++){</pre>
240
             while (now>0&&sign(cross(ans[now]-ans[now-1],A[i]-ans[now-1]))<flag) now--;</pre>
241
             ans[++now]=A[i];
242
         } int pre=now;
243
         for (int i=n-2;i>=0;i--){
244
             while (now>pre&&sign(cross(ans[now]-ans[now-1],A[i]-ans[now-1]))<flag) now--;</pre>
245
             ans[++now]=A[i];
         } ans.resize(now); return ans;
246
247 }
248 db convexDiameter(vector<point>A){
249
         int now=0,n=A.size(); db ans=0;
250
         for (int i=0;i<A.size();i++){</pre>
251
             now=max(now,i);
252
             while (1){
253
                 db k1=A[i].dis(A[now%n]),k2=A[i].dis(A[(now+1)%n]);
254
                 ans=max(ans,max(k1,k2)); if (k2>k1) now++; else break;
255
             }
256
         }
257
         return ans;
258
     vector<point> convexcut(vector<point>A,point k1,point k2){
259
260
         // 保留 k1,k2,p 逆时针的所有点
261
         int n=A.size(); A.push_back(A[0]); vector<point>ans;
262
         for (int i=0;i<n;i++){</pre>
263
             int w1=clockwise(k1,k2,A[i]),w2=clockwise(k1,k2,A[i+1]);
264
             if (w1>=0) ans.push_back(A[i]);
265
             if (w1*w2<0) ans.push_back(getLL(k1,k2,A[i],A[i+1]));</pre>
266
         }
267
         return ans;
268
269
     int checkPoS(vector<point>A,point k1,point k2){
270
         // 多边形 A 和直线 ( 线段 )k1->k2 严格相交 , 注释部分为线段
271
         struct ins{
272
             point m,u,v;
273
             int operator < (const ins& k) const {return m<k.m;}</pre>
274
         }; vector<ins>B;
275
         //if (contain(A,k1)==2||contain(A,k2)==2) return 1;
276
         vector<point>poly=A; A.push_back(A[0]);
         for (int i=1;i<A.size();i++) if (checkLL(A[i-1],A[i],k1,k2)){</pre>
277
278
             point m=getLL(A[i-1],A[i],k1,k2);
279
             if (inmid(A[i-1],A[i],m)/*&&inmid(k1,k2,m)*/) B.push_back((ins){m,A[i-1],A[i]});
280
         }
281
         if (B.size()==0) return 0; sort(B.begin(),B.end());
282
         int now=1; while (now<B.size()&&B[now].m==B[0].m) now++;</pre>
283
         if (now==B.size()) return 0;
284
         int flag=contain(poly,(B[0].m+B[now].m)/2);
285
         if (flag==2) return 1;
286
         point d=B[now].m-B[0].m;
287
         for (int i=now;i<B.size();i++){</pre>
```

```
288
             if (!(B[i].m==B[i-1].m)&&flag==2) return 1;
289
             int tag=sign(cross(B[i].v-B[i].u,B[i].m+d-B[i].u));
290
             if (B[i].m==B[i].u||B[i].m==B[i].v) flag+=tag; else flag+=tag*2;
291
         }
292
         //return 0;
293
         return flag==2;
294
    }
295
    int checkinp(point r,point l,point m){
296
         if (compareangle(1,r)){return compareangle(1,m)&&compareangle(m,r);}
297
         return compareangle(1,m)||compareangle(m,r);
298
299
    int checkPosFast(vector<point>A,point k1,point k2){ // 快速检查线段是否和多边形严格相交
300
         if (contain(A,k1)==2||contain(A,k2)==2) return 1; if (k1==k2) return 0;
301
         A.push_back(A[0]); A.push_back(A[1]);
302
         for (int i=1;i+1<A.size();i++)</pre>
303
             if (checkLL(A[i-1],A[i],k1,k2)){
304
                point now=getLL(A[i-1],A[i],k1,k2);
                 if (inmid(A[i-1],A[i],now)==0||inmid(k1,k2,now)==0) continue;
305
306
                if (now==A[i]){
307
                     if (A[i]==k2) continue;
308
                     point pre=A[i-1],ne=A[i+1];
                     if (checkinp(pre-now,ne-now,k2-now)) return 1;
309
310
                } else if (now==k1){
                     if (k1==A[i-1]||k1==A[i]) continue;
311
312
                     if (checkinp(A[i-1]-k1,A[i]-k1,k2-k1)) return 1;
313
                } else if (now==k2||now==A[i-1]) continue;
314
                else return 1;
315
            }
316
         return 0;
317
    // 拆分凸包成上下凸壳 凸包尽量都随机旋转一个角度来避免出现相同横坐标
318
319
    // 尽量特判只有一个点的情况 凸包逆时针
320
    void getUDP(vector<point>A,vector<point>&U,vector<point>&D){
321
         db l=1e100,r=-1e100;
322
         for (int i=0;i<A.size();i++) l=min(l,A[i].x),r=max(r,A[i].x);</pre>
323
         int wherel, wherer;
324
         for (int i=0;i<A.size();i++) if (cmp(A[i].x,1)==0) wherel=i;</pre>
325
         for (int i=A.size();i;i--) if (cmp(A[i-1].x,r)==0) where r=i-1;
326
         U.clear(); D.clear(); int now=wherel;
327
         while (1){D.push_back(A[now]); if (now==wherer) break; now++; if (now>=A.size()) now=0;}
328
        now=where1:
329
         while (1) {U.push_back(A[now]); if (now==wherer) break; now--; if (now<0) now=A.size()-1;}
330 }
    // 需要保证凸包点数大于等于 3,2 内部 ,1 边界 ,0 外部
331
332
    int containCoP(const vector<point>&U,const vector<point>&D,point k){
333
         db lx=U[0].x,rx=U[U.size()-1].x;
334
         if (k==U[0]||k==U[U.size()-1]) return 1;
335
         if (cmp(k.x,lx) == -1 | cmp(k.x,rx) == 1) return 0;
336
         int where1=lower_bound(U.begin(), U.end(), (point) {k.x,-1e100}) - U.begin();
337
         int where2=lower bound(D.begin(),D.end(),(point){k.x,-1e100})-D.begin();
338
         int w1=clockwise(U[where1-1],U[where1],k),w2=clockwise(D[where2-1],D[where2],k);
339
         if (w1==1||w2==-1) return 0; else if (w1==0||w2==0) return 1; return 2;
340 }
    // d 是方向,输出上方切点和下方切点
341
    pair<point,point> getTangentCow(const vector<point> &U,const vector<point> &D,point d){
342
343
         if (sign(d.x)<0||(sign(d.x)==0\&\&sign(d.y)<0)) d=d*(-1);
344
         point whereU, whereD;
345
         if (sign(d.x)==0) return mp(U[0],U[U.size()-1]);
346
         int l=0,r=U.size()-1,ans=0;
```

```
347
         while (l<r){int mid=l+r>>1; if (sign(cross(U[mid+1]-U[mid],d))<=0) l=mid+1,ans=mid+1; else r=
         mid;}
348
         whereU=U[ans]; l=0,r=D.size()-1,ans=0;
349
         while (1<r){int mid=1+r>>1; if (sign(cross(D[mid+1]-D[mid],d))>=0) l=mid+1,ans=mid+1; else r=
350
         whereD=D[ans]; return mp(whereU, whereD);
351 }
352 // 先检查 contain, 逆时针给出
    pair<point,point> getTangentCoP(const vector<point>&U,const vector<point>&D,point k){
353
354
         db lx=U[0].x,rx=U[U.size()-1].x;
355
         if (k.x<lx){</pre>
356
             int l=0,r=U.size()-1,ans=U.size()-1;
357
             while (l<r){int mid=l+r>>1; if (clockwise(k,U[mid],U[mid+1])==1) l=mid+1; else ans=mid,r=
         mid;}
358
             point w1=U[ans]; l=0,r=D.size()-1,ans=D.size()-1;
359
             while (l<r){int mid=l+r>>1; if (clockwise(k,D[mid],D[mid+1])==-1) l=mid+1; else ans=mid,r=
360
             point w2=D[ans]; return mp(w1,w2);
361
         } else if (k.x>rx){
362
             int l=1,r=U.size(),ans=0;
363
             while (l<r){int mid=l+r>>1; if (clockwise(k,U[mid],U[mid-1])==-1) r=mid; else ans=mid,l=mid
         +1:}
364
             point w1=U[ans]; l=1,r=D.size(),ans=0;
             while (l<r){int mid=l+r>>1; if (clockwise(k,D[mid],D[mid-1])==1) r=mid; else ans=mid,l=mid
365
366
             point w2=D[ans]; return mp(w2,w1);
367
         } else {
368
             int where1=lower_bound(U.begin(), U.end(), (point) {k.x,-1e100}) - U.begin();
369
             int where2=lower_bound(D.begin(),D.end(),(point){k.x,-1e100})-D.begin();
370
             if ((k.x==lx\&k.y>U[0].y)||(where1\&clockwise(U[where1-1],U[where1],k)==1)){
371
                 int l=1,r=where1+1,ans=0;
372
                 while (l<r){int mid=l+r>>1; if (clockwise(k,U[mid],U[mid-1])==1) ans=mid,l=mid+1; else
         r=mid;}
373
                 point w1=U[ans]; l=where1,r=U.size()-1,ans=U.size()-1;
                 while (1<r){int mid=1+r>>1; if (clockwise(k,U[mid],U[mid+1])==1) 1=mid+1; else ans=mid,
374
         r=mid;}
375
                 point w2=U[ans]; return mp(w2,w1);
376
             } else {
377
                 int l=1,r=where2+1,ans=0;
378
                 while (l<r){int mid=l+r>>1; if (clockwise(k,D[mid],D[mid-1])==-1) ans=mid,l=mid+1; else
          r=mid;}
379
                 point w1=D[ans]; l=where2,r=D.size()-1,ans=D.size()-1;
380
                 while (1<r){int mid=1+r>>1; if (clockwise(k,D[mid],D[mid+1])==-1) l=mid+1; else ans=mid
         ,r=mid;}
                 point w2=D[ans]; return mp(w1,w2);
381
382
             }
383
         }
384
    }
385
    struct P3{
386
         db x.v.z:
387
         P3 operator + (P3 k1){return (P3){x+k1.x,y+k1.y,z+k1.z};}
388
         P3 operator - (P3 k1){return (P3){x-k1.x,y-k1.y,z-k1.z};}
389
         P3 operator * (db k1){return (P3){x*k1,y*k1,z*k1};}
         P3 operator / (db k1){return (P3){x/k1,y/k1,z/k1};}
390
391
         db abs2(){return x*x+y*y+z*z;}
392
         db abs(){return sqrt(x*x+y*y+z*z);}
393
         P3 unit(){return (*this)/abs();}
394
         int operator < (const P3 k1) const{</pre>
395
             if (cmp(x,k1.x)!=0) return x<k1.x;</pre>
```

```
396
              if (cmp(y,k1.y)!=0) return y<k1.y;</pre>
397
              return cmp(z,k1.z) == -1;
         }
398
399
          int operator == (const P3 k1){
400
              return cmp(x,k1.x) == 0 \& cmp(y,k1.y) == 0 \& cmp(z,k1.z) == 0;
401
         }
402
         void scan(){
403
              double k1,k2,k3; scanf("%lf%lf%lf",&k1,&k2,&k3);
404
              x=k1; y=k2; z=k3;
405
         }
406 };
     P3 cross(P3 k1,P3 k2){return (P3){k1.y*k2.z-k1.z*k2.y,k1.z*k2.x-k1.x*k2.z,k1.x*k2.y-k1.y*k2.x};}
407
     db dot(P3 k1,P3 k2){return k1.x*k2.x+k1.y*k2.y+k1.z*k2.z;}
408
     //p=(3,4,5),l=(13,19,21),theta=85 ans=(2.83,4.62,1.77)
409
410
     P3 turn3D(db k1,P3 1,P3 p){
411
         l=1.unit(); P3 ans; db c=cos(k1),s=sin(k1);
412
         ans.x = p.x * (1.x * 1.x * (1-c) + p.y * (1.x * 1.y * (1-c) - 1.z * s) + p.z * (1.x * 1.z * (1-c) + 1.y * s);
413
         ans.y = p.x * (1.x * 1.y * (1-c) + 1.z * s) + p.y * (1.y * 1.y * (1-c) + c) + p.z * (1.y * 1.z * (1-c) - 1.x * s);
414
         ans.z = p.x * (1.x * 1.z * (1-c) - 1.y * s) + p.y * (1.y * 1.z * (1-c) + 1.x * s) + p.z * (1.x * 1.x * (1-c) + c);
415
         return ans;
416 }
417 typedef vector<P3> VP;
418 typedef vector<VP> VVP;
419 db A\cos(db x){return a\cos(max(-(db)1,min(x,(db)1)));}
420 // 球面距离 , 圆心原点 , 半径 1
421 db Odist(P3 a,P3 b){db r=Acos(dot(a,b)); return r;}
422 db r; P3 rnd;
423
     vector<db> solve(db a,db b,db c){
424
         db r=sqrt(a*a+b*b),th=atan2(b,a);
425
          if (cmp(c,-r)==-1) return {0};
426
         else if (cmp(r,c) \le 0) return \{1\};
427
         else {
428
              db tr=pi-Acos(c/r); return {th+pi-tr,th+pi+tr};
429
         }
430 }
431
     vector<db> jiao(P3 a,P3 b){
432
         // dot(rd+x*cos(t)+y*sin(t),b) >= cos(r)
433
         if (cmp(Odist(a,b),2*r)>0) return {0};
434
         P3 rd=a*cos(r),z=a.unit(),y=cross(z,rnd).unit(),x=cross(y,z).unit();
435
         vector < db > ret = solve(-(dot(x,b)*sin(r)), -(dot(y,b)*sin(r)), -(cos(r)-dot(rd,b)));
436
         return ret;
437
438
     db norm(db x,db l=0,db r=2*pi){ // change x into [1,r)
439
         while (cmp(x,1)==-1) x+=(r-1); while (cmp(x,r)>=0) x-=(r-1);
440
         return x;
441 }
442
     db disLP(P3 k1,P3 k2,P3 q){
443
         return (cross(k2-k1,q-k1)).abs()/(k2-k1).abs();
444 }
     db disLL(P3 k1,P3 k2,P3 k3,P3 k4){
445
446
         P3 dir=cross(k2-k1,k4-k3); if (sign(dir.abs())==0) return disLP(k1,k2,k3);
447
         return fabs(dot(dir.unit(),k1-k2));
448
449
     VP getFL(P3 p,P3 dir,P3 k1,P3 k2){
450
          db a=dot(k2-p,dir),b=dot(k1-p,dir),d=a-b;
451
          if (sign(fabs(d))==0) return {};
452
         return {(k1*a-k2*b)/d};
453
454 VP getFF(P3 p1,P3 dir1,P3 p2,P3 dir2){// 返回一条线
```

```
455
         P3 e=cross(dir1,dir2),v=cross(dir1,e);
456
         db d=dot(dir2,v); if (sign(abs(d))==0) return {};
457
         P3 q=p1+v*dot(dir2,p2-p1)/d; return {q,q+e};
458
459
    // 3D Covex Hull Template
     db getV(P3 k1,P3 k2,P3 k3,P3 k4){ // get the Volume
460
         return dot(cross(k2-k1,k3-k1),k4-k1);
461
462 }
463
    db rand_db(){return 1.0*rand()/RAND_MAX;}
    VP convexHull2D(VP A,P3 dir){
         P3 x={(db)rand(),(db)rand(),(db)rand()}; x=x.unit();
465
466
         x=cross(x,dir).unit(); P3 y=cross(x,dir).unit();
467
         P3 vec=dir.unit()*dot(A[0],dir);
468
         vector<point>B;
469
         for (int i=0;i<A.size();i++) B.push_back((point){dot(A[i],x),dot(A[i],y)});</pre>
470
         B=ConvexHull(B); A.clear();
471
         for (int i=0;i<B.size();i++) A.push_back(x*B[i].x+y*B[i].y+vec);</pre>
472
         return A;
473 }
474
    namespace CH3{
475
         VVP ret; set<pair<int,int> >e;
476
         int n; VP p,q;
477
         void wrap(int a,int b){
             if (e.find({a,b})==e.end()){
478
479
                 int c=-1;
480
                 for (int i=0;i<n;i++) if (i!=a&&i!=b){</pre>
481
                     if (c==-1||sign(getV(q[c],q[a],q[b],q[i]))>0) c=i;
482
                 }
483
                 if (c!=-1){
484
                     ret.push_back({p[a],p[b],p[c]});
485
                     e.insert({a,b}); e.insert({b,c}); e.insert({c,a});
486
                     wrap(c,b); wrap(a,c);
487
                 }
             }
488
         }
489
         VVP ConvexHull3D(VP _p){
490
491
             p=q=_p; n=p.size();
492
             ret.clear(); e.clear();
493
             for (auto &i:q) i=i+(P3){rand_db()*1e-4,rand_db()*1e-4,rand_db()*1e-4};
494
             for (int i=1;i<n;i++) if (q[i].x<q[0].x) swap(p[0],p[i]),swap(q[0],q[i]);</pre>
495
             for (int i=2;i<n;i++) if ((q[i].x-q[0].x)*(q[1].y-q[0].y)>(q[i].y-q[0].y)*(q[1].x-q[0].x))
         swap(q[1],q[i]),swap(p[1],p[i]);
496
             wrap(0,1);
497
             return ret;
498
         }
499
    VVP reduceCH(VVP A){
500
501
         VVP ret; map<P3,VP> M;
502
         for (VP nowF:A){
             P3 dir=cross(nowF[1]-nowF[0],nowF[2]-nowF[0]).unit();
503
504
             for (P3 k1:nowF) M[dir].pb(k1);
505
506
         for (pair<P3,VP> nowF:M) ret.pb(convexHull2D(nowF.se,nowF.fi));
507
         return ret;
508
509
    // 把一个面变成 (点,法向量)的形式
    pair<P3,P3> getF(VP F){
510
511
         return mp(F[0],cross(F[1]-F[0],F[2]-F[0]).unit());
512 }
```

```
513 // 3D Cut 保留 dot(dir,x-p)>=0 的部分
514
    VVP ConvexCut3D(VVP A,P3 p,P3 dir){
515
         VVP ret; VP sec;
516
         for (VP nowF: A){
             int n=nowF.size(); VP ans; int dif=0;
517
518
             for (int i=0;i<n;i++){</pre>
                 int d1=sign(dot(dir,nowF[i]-p));
519
                 int d2=sign(dot(dir,nowF[(i+1)%n]-p));
520
                 if (d1>=0) ans.pb(nowF[i]);
521
522
                 if (d1*d2<0){
                     P3 q=getFL(p,dir,nowF[i],nowF[(i+1)%n])[0];
523
524
                     ans.push_back(q); sec.push_back(q);
525
                 }
526
                 if (d1==0) sec.push_back(nowF[i]); else dif=1;
527
                 dif|=(sign(dot(dir,cross(nowF[(i+1)%n]-nowF[i],nowF[(i+1)%n]-nowF[i])))==-1);
             }
528
529
             if (ans.size()>0&&dif) ret.push_back(ans);
         }
530
         if (sec.size()>0) ret.push_back(convexHull2D(sec,dir));
531
532
         return ret;
533 }
    db vol(VVP A){
534
535
         if (A.size()==0) return 0; P3 p=A[0][0]; db ans=0;
536
         for (VP nowF:A)
537
             for (int i=2;i<nowF.size();i++)</pre>
538
                 ans+=abs(getV(p,nowF[0],nowF[i-1],nowF[i]));
539
         return ans/6;
540
    }
    VVP init(db INF) {
541
542
         VVP pss(6, VP(4));
543
         pss[0][0] = pss[1][0] = pss[2][0] = {-INF, -INF, -INF};
544
         pss[0][3] = pss[1][1] = pss[5][2] = {-INF, -INF};
         pss[0][1] = pss[2][3] = pss[4][2] = {-INF, INF, -INF};
545
         pss[0][2] = pss[5][3] = pss[4][1] = {-INF, INF, INF};
546
         pss[1][3] = pss[2][1] = pss[3][2] = {INF, -INF, -INF};
547
548
         pss[1][2] = pss[5][1] = pss[3][3] = {INF, -INF, INF};
         pss[2][2] = pss[4][3] = pss[3][1] = {INF, INF, -INF};
549
550
         pss[5][0] = pss[4][0] = pss[3][0] = {INF, INF, INF};
551
         return pss;
552 }
     4.1.2 3D 凸包
    #include<stdio.h>
    #include<algorithm>
 3 #include<string.h>
 4 #include<math.h>
 5 #include<stdlib.h>
 6 using namespace std;
 7
    const int MAXN=1050;
 8
    const double eps=1e-8;
 9
10 \quad {\tt struct} \ {\tt Point}
11
    {
12
         double x,y,z;
13
         Point() {}
14
         Point(double xx,double yy,double zz):x(xx),y(yy),z(zz) {}
         //两向量之差
15
```

```
16
       Point operator -(const Point p1)
17
18
           return Point(x-p1.x,y-p1.y,z-p1.z);
19
        }
20
        //两向量之和
21
       Point operator +(const Point p1)
22
        {
23
           return Point(x+p1.x,y+p1.y,z+p1.z);
24
       }
25
       //叉乘
26
       Point operator *(const Point p)
27
28
           return Point(y*p.z-z*p.y,z*p.x-x*p.z,x*p.y-y*p.x);
29
       }
30
       Point operator *(double d)
31
32
           return Point(x*d,y*d,z*d);
33
       }
34
       Point operator / (double d)
35
        {
36
           return Point(x/d,y/d,z/d);
37
       }
       //点乘
38
39
       double operator ^(Point p)
40
41
           return (x*p.x+y*p.y+z*p.z);
42
       }
43
   };
   struct CH3D
44
45
   {
46
        struct face
47
        {
           //表示凸包一个面上的三个点的编号
48
49
           int a,b,c;
           //表示该面是否属于最终凸包上的面
50
51
           bool ok;
52
       };
53
       //初始顶点数
54
       int n;
        //初始顶点
55
       Point P[MAXN];
56
        //凸包表面的三角形数
57
58
        int num;
59
        //凸包表面的三角形
60
       face F[8*MAXN];
        //凸包表面的三角形
61
62
        int g[MAXN][MAXN];
63
        //向量长度
       double vlen(Point a)
64
65
        {
66
           return sqrt(a.x*a.x+a.y*a.y+a.z*a.z);
67
       }
68
       //叉乘
69
       Point cross(const Point &a,const Point &b,const Point &c)
70
           return Point((b.y-a.y)*(c.z-a.z)-(b.z-a.z)*(c.y-a.y),
71
72
                        (b.z-a.z)*(c.x-a.x)-(b.x-a.x)*(c.z-a.z),
73
                        (b.x-a.x)*(c.y-a.y)-(b.y-a.y)*(c.x-a.x)
74
                       );
```

```
75
         }
         //三角形面积*2
76
77
         double area(Point a,Point b,Point c)
78
         {
79
             return vlen((b-a)*(c-a));
80
         }
81
         //四面体有向体积*6
82
         double volume(Point a,Point b,Point c,Point d)
83
         {
84
             return (b-a)*(c-a)^(d-a);
85
         }
86
         //正: 点在面同向
87
         double dblcmp(Point &p,face &f)
88
89
             Point m=P[f.b]-P[f.a];
             Point n=P[f.c]-P[f.a];
90
91
             Point t=p-P[f.a];
92
             return (m*n)^t;
93
         }
94
         void deal(int p,int a,int b)
95
             int f=g[a][b];//搜索与该边相邻的另一个平面
96
97
             face add;
98
             if(F[f].ok)
99
             {
100
                 if(dblcmp(P[p],F[f])>eps)
101
                     dfs(p,f);
102
                 else
103
                 {
104
                     add.a=b;
105
                     add.b=a;
106
                     add.c=p;//这里注意顺序,要成右手系
107
                     add.ok=true;
108
                     g[p][b]=g[a][p]=g[b][a]=num;
109
                     F[num++]=add;
                 }
110
111
             }
112
        }
113
         void dfs(int p,int now)//递归搜索所有应该从凸包内删除的面
114
         {
             F[now].ok=0;
115
             deal(p,F[now].b,F[now].a);
116
117
             deal(p,F[now].c,F[now].b);
118
             deal(p,F[now].a,F[now].c);
         }
119
120
        bool same(int s,int t)
121
             Point &a=P[F[s].a];
122
             Point &b=P[F[s].b];
123
124
             Point &c=P[F[s].c];
125
             return fabs(volume(a,b,c,P[F[t].a]))<eps &&</pre>
126
                    fabs(volume(a,b,c,P[F[t].b]))<eps &&</pre>
127
                    fabs(volume(a,b,c,P[F[t].c]))<eps;</pre>
128
         }
129
         //构建三维凸包
130
         void create()
131
         {
132
             int i,j,tmp;
133
             face add;
```

```
134
             num=0;
             if(n<4)return;</pre>
135
136
             //*******
             //此段是为了保证前四个点不共面
137
138
             bool flag=true;
139
             for(i=1; i<n; i++)</pre>
                 if(vlen(P[0]-P[i])>eps)
140
                 {
141
                      swap(P[1],P[i]);
142
143
                      flag=false;
144
                      break;
145
                 }
146
             if(flag)return;
             flag=true;
147
             //使前三个点不共线
148
             for(i=2; i<n; i++)</pre>
149
                 if(vlen((P[0]-P[1])*(P[1]-P[i]))>eps)
150
151
152
                      swap(P[2],P[i]);
153
                      flag=false;
154
                      break;
                 }
155
             if(flag)return;
156
157
             flag=true;
158
             //使前四个点不共面
159
             for(int i=3; i<n; i++)</pre>
160
                 if(fabs((P[0]-P[1])*(P[1]-P[2])^(P[0]-P[i]))>eps)
161
                      swap(P[3],P[i]);
162
163
                      flag=false;
164
                      break;
165
166
             if(flag)return;
             //***********************
167
168
             for(i=0; i<4; i++)</pre>
169
             {
170
                 add.a=(i+1)\%4;
171
                 add.b=(i+2)%4;
                 add.c=(i+3)\%4;
172
173
                 add.ok=true;
174
                 if(dblcmp(P[i],add)>0)swap(add.b,add.c);
175
                 g[add.a] [add.b]=g[add.b] [add.c]=g[add.c] [add.a]=num;
176
                 F[num++]=add;
177
178
             for(i=4; i<n; i++)</pre>
179
                 for(j=0; j<num; j++)</pre>
180
                      if(F[j].ok&&dblcmp(P[i],F[j])>eps)
181
182
                          dfs(i,j);
183
                          break;
184
185
             tmp=num;
186
             for(i=num=0; i<tmp; i++)</pre>
187
                 if(F[i].ok)
188
                      F[num++]=F[i];
189
190
         }
191
         //表面积
         double area()
192
```

```
193
         {
194
             double res=0;
195
             if(n==3)
196
             {
197
                  Point p=cross(P[0],P[1],P[2]);
198
                  res=vlen(p)/2.0;
199
                  return res;
200
             }
201
             for(int i=0; i<num; i++)</pre>
202
                  res+=area(P[F[i].a],P[F[i].b],P[F[i].c]);
203
             return res/2.0;
204
         }
         //体积
205
206
         double volume()
207
208
             double res=0;
209
             Point tmp(0,0,0);
210
             for(int i=0; i<num; i++)</pre>
211
                  res+=volume(tmp,P[F[i].a],P[F[i].b],P[F[i].c]);
212
             return fabs(res/6.0);
         }
213
214
         //表面三角形个数
215
         int triangle()
216
         {
217
             return num;
         }
218
219
         //表面多边形个数
220
         int polygon()
221
222
             int i,j,res,flag;
223
             for(i=res=0; i<num; i++)</pre>
224
             {
225
                  flag=1;
226
                  for(j=0; j<i; j++)</pre>
227
                      if(same(i,j))
228
229
                          flag=0;
230
                          break;
231
                      }
232
                  res+=flag;
233
             }
234
             return res;
         }
235
236
         //三维凸包重心
237
         Point barycenter()
238
         {
239
             Point ans(0,0,0),o(0,0,0);
240
             double all=0;
             for(int i=0; i<num; i++)</pre>
241
242
             {
243
                  double vol=volume(o,P[F[i].a],P[F[i].b],P[F[i].c]);
244
                  ans=ans+(o+P[F[i].a]+P[F[i].b]+P[F[i].c])/4.0*vol;
245
                  all+=vol;
246
             }
247
             ans=ans/all;
248
             return ans;
249
         }
250
         //点到面的距离
251
         double ptoface(Point p,int i)
```

```
252
         {
253
             return fabs(volume(P[F[i].a],P[F[i].b],P[F[i].c],p)/vlen((P[F[i].b]-P[F[i].a])*(P[F[i].c]-P
         [F[i].a])));
254
         }
255 };
256
    CH3D hull;
257
    int main()
258 {
259
         while(~scanf("%d",&hull.n))
260
261
             for(int i=0; i<hull.n; i++)</pre>
262
                 scanf("\%lf\%lf",\&hull.P[i].x,\&hull.P[i].y,\&hull.P[i].z);\\
             hull.create();
263
264
             printf("%d\n",hull.polygon());
265
         }
266
         return 0;
267 }
```

5 String

5.1 KMP

```
5.1.1 KMP
```

```
// nxt[0]表示失配到完全不匹配
  int nxt[maxm];
3
   void getNext(char *s, int len) {
4
5
       int i = 0, j = -1;
6
       nxt[i] = j;
7
       while (i < len) {</pre>
8
           if (j == -1 || s[i] == s[j]) nxt[++i] = ++j;
9
           else j = nxt[j];
10
       }
11 }
12
13 // a为原串, b为模式串, 下标从0开始, 找第一个出现模式串的位置 (起点为1), 找不到返回-1
   int KMP(char *a, char *b, int n, int m) {
       getNext(b, m);
15
16
       int i = 0, j = 0;
       while (i < n && j < m) {</pre>
17
18
           if (j == -1 || a[i] == b[j]) ++i, ++j;
19
           else j = nxt[j];
20
       }
21
       return j == m ? i - m + 1 : -1;
22 }
   5.1.2 exKMP
1 const int maxn = 1e5 + 10;
2 int nex[maxn], extend[maxn];
3
4 //预处理计算Next数组
5 void getNext(char *str)
6 {
7
       int i = 0, j, po, len = strlen(str);
8
       nex[0] = len;
                      //初始化nex[0]
9
       while (str[i] == str[i + 1] && i + 1 < len) i++; //计算nex[1]
10
       nex[1] = i;
11
       po = 1; //初始化po的位置
12
       for (int i = 2; i < len; i++)</pre>
13
14
           if (nex[i - po] + i < nex[po] + po) //第一种情况, 可以直接得到nex[i]的值
15
              nex[i] = nex[i - po];
16
           else
                  //第二种情况,要继续匹配才能得到nex[i]的值
17
           {
              j = nex[po] + po - i;
18
19
              if (j < 0) j = 0;
                                  //如果i>po+nex[po],则要从头开始匹配
20
              while (i + j < len \&\& str[j] == str[j + i]) j++;
21
              nex[i] = j;
22
              po = i; //更新po的位置
23
           }
24
       }
25 }
27 void EXKMP(char *s1, char *s2)
28 {
```

```
29
        int i = 0, j, po, len = strlen(s1), 12 = strlen(s2);
30
        getNext(s2);
        while (s1[i] == s2[i] && i < 12 && i < len) i++;</pre>
31
32
        extend[0] = i;
33
        po = 0;
34
        for (int i = 1; i < len; i++)</pre>
35
            if (nex[i - po] + i < extend[po] + po)</pre>
36
37
                extend[i] = nex[i - po];
38
            else
39
            {
40
                j = extend[po] + po - i;
41
                if (j < 0) j = 0;
                while (i + j < len && j < 12 && s1[j + i] == s2[j]) j++;
42
43
                extend[i] = j;
44
                po = i;
45
            }
46
        }
47 }
    5.2 Trie
    5.2.1 Trie
   const int maxn = 2e6 + 10;
 2
 3
   int trie[maxn][30], tot;
 4 bool flag[maxn];
 5
 6 void insert_ch(char *str)
 7
 8
        int len = strlen(str);
 9
        int root = 0;
10
        for (int i = 0; i < len; i++)</pre>
11
12
            int id = str[i] - 'a';
13
            if (!trie[root][id]) trie[root][id] = ++tot;
14
            root = trie[root][id];
15
16
        flag[root] = true;
17
   }
18
19 bool find_ch(char *str)
20 {
21
        int len = strlen(str);
22
        int root = 0;
23
        for (int i = 0; i < len; i++)</pre>
24
25
            int id = str[i] - 'a';
26
            if (!trie[root][id]) return false;
27
            root = trie[root][id];
        }
28
29
        return true;
30 }
    5.2.2 Persistence Trie
 1 const int maxn = 1e5 + 10;
```

```
int a[maxn], rt[maxn], n;
3
4
5
   struct Trie
6
   {
7
       int tot;
8
       int child[maxn * 32][2], sum[maxn *32];
9
       int insert(int x, int val)
10
11
           int tmp, y;
12
           tmp = y= ++tot;
13
           for(int i = 30; i >= 0; --i)
14
               child[y][0] = child[x][0];
15
16
               child[y][1] = child[x][1];
17
               sum[y] = sum[x] + 1;
18
               int t = val >> i & 1;
19
               x = child[x][t];
20
               child[y][t] = ++tot;
21
               y = child[y][t];
22
23
           sum[y] = sum[x] + 1;
24
           return tmp;
25
26
       int query(int 1, int r, int val)
27
28
           int tmp = 0;
29
           for(int i =30; i >= 0; --i)
30
31
               int t = val >> i & 1;
32
               [1][t ^ 1];
33
               else r = child[r][t], l = child[l][t];
           }
34
35
           return tmp;
       }
36
37 }trie;
   5.2.3 01Trie
   struct Trie {
2
       int tree[maxn*20][2], tot;
3
       int flag[maxn*20];
4
5
       void insert_ch(int x) {
6
           int root = 0;
7
           flag[0]++;
8
           for (int i = 30; i >= 0; --i) {
9
               int id = (x >> i) & 1;
10
               if (!tree[root][id]) {
11
                   tree[root][id] = ++tot;
                   tree[tree[root][id]][0] = tree[tree[root][id]][1] = 0;
12
                   flag[tree[root][id]] = flag[tree[tree[root][id]][0]] = flag[tree[tree[root][id
13
       ]][1]] = 0;
14
               }
               root = tree[root][id];
15
16
               flag[root]++;
17
           }
```

```
}
18
19
        void del(int x) {
20
21
            int root = 0;
22
            flag[0]--;
23
            for (int i = 30; i >= 0; --i) {
24
                int id = (x >> i) & 1;
25
                assert(tree[root][id]);
26
                if (flag[tree[root][id]] == 1) {
27
                    flag[tree[root][id]] = 0;
28
                    tree[root][id] = 0;
29
                    return;
                }
30
31
                root = tree[root][id];
32
                flag[root]--;
            }
33
        }
34
35
36
        int find_ch(int x, int flag = 0) { // flag 0 最小异或值, 1 最大异或值
            int root = 0;
37
            int res = 0;
38
39
            for (int i = 30; i >= 0; --i) {
40
                int id = ((x >> i) & 1);
41
                if (flag) id = !id;
42
                if (tree[root][id]) {
43
                    root = tree[root][id];
44
                    res = res << 1 | id;
45
                } else {
46
                    root = tree[root][!id];
47
                    res = res << 1 | (!id);
48
49
            }
50
            return res;
51
        }
52
        void init() {
53
54
            tree[0][0] = tree[0][1] = 0;
55
            tot = 0;
56
        }
57 };
         Manachar
    5.3
    5.3.1 Manacher
1 const int maxn = 1e5 + 10;
2
3 char s[maxn];
4
5 char tmp[maxn << 1];</pre>
   int Len[maxn << 1];</pre>
7
   int init(char *str)
8
9
10
        int len = strlen(str);
        tmp[0] = '@';
11
        for (int i = 1; i <= 2 * len; i += 2)</pre>
12
13
14
            tmp[i] = '#';
```

```
15
            tmp[i + 1] = str[i / 2];
        }
16
17
        tmp[2 * len + 1] = '#';
18
        tmp[2 * len + 2] = '$';
19
        tmp[2 * len + 3] = 0;
20
        return 2 * len + 1;
21 }
22
23
   int manacher(char *str)
24
25
        int mx = 0, ans = 0, pos = 0;
26
        int len = init(str);
27
        for (int i = 1; i <= len; i++)</pre>
28
29
            if (mx > i) Len[i] = min(mx - i, Len[2 * pos - i]);
30
            else Len[i] = 1;
31
            while (tmp[i - Len[i]] == tmp[i + Len[i]]) Len[i]++;
32
            if (Len[i] + i > mx) mx = Len[i] + i, pos = i;
33
        }
34 }
```

5.4 Aho-Corasick Automation

5.4.1 AC Automation

```
1 class AC_automation
2 {
3
   public:
        int trie[maxn][26], cnt;
4
5
        int tag[maxn];
6
        int fail[maxn], num[maxn], res[maxn], in[maxn], Map[maxn];
7
8
        void init()
9
        {
10
            memset(trie, 0, sizeof trie);
            memset(tag, 0, sizeof tag);
11
            memset(fail, 0, sizeof fail);
12
13
            cnt = 0;
14
15
16
        void insert(char *str, int id)
17
18
            int root = 0;
19
            for (int i = 0; str[i]; i++)
20
            {
21
                int id = str[i] - 'a';
                if (!trie[root][id]) trie[root][id] = ++cnt;
22
23
                root = trie[root][id];
24
            }
25
            if(!tag[root]) tag[root] = id;
26
            Map[id] = tag[root];
27
        }
28
29
        void build()
30
31
            queue<int> que;
32
            for (int i = 0; i < 26; i++) if (trie[0][i]) que.push(trie[0][i]);</pre>
33
            while (!que.empty())
34
            {
```

```
35
                 int k = que.front();
36
                 que.pop();
37
                 for (int i = 0; i < 26; i++)</pre>
38
39
                     if (trie[k][i])
40
                     {
                         fail[trie[k][i]] = trie[fail[k]][i];
41
42
                         que.push(trie[k][i]);
                         in[fail[trie[k][i]]] ++;
43
44
                     } else trie[k][i] = trie[fail[k]][i];
                 }
45
46
            }
47
        }
48
49
        void toposort()
50
51
             queue<int> que;
52
             for(int i = 1; i <= cnt; i ++) if(in[i] == 0) que.push(i);</pre>
53
            while(!que.empty())
54
55
                 int u = que.front(); que.pop();
                 res[tag[u]] = num[u];
56
57
                 int v = fail[u]; in[v] --;
                 num[v] += num[u];
58
59
                 if(in[v] == 0) que.push(v);
60
            }
        }
61
62
63
        void query(char *str, int n)
64
65
             int u = 0, len = strlen(s);
66
             for(int i = 0; i < len; i ++)</pre>
67
                 u = trie[u][str[i] - 'a'], num[u] ++;
68
             toposort();
69
             for(int i = 1; i <= n; i ++) printf("%d\n", res[Map[i]]);</pre>
70
        }
71 } AC;
    5.5
          Suffix Array
    5.5.1 Suffix Array
    char s[maxn];
    int sa[maxn], t[maxn], t2[maxn], c[maxn], n;
 3
 4
   //build_sa(n + 1, 130), sa, height下标从1开始,rk下标从0开始
 5
    void build_sa(int n, int m)
 6
    {
 7
        int *x = t, *y = t2;
 8
        for(int i = 0; i < m; i++) c[i] = 0;</pre>
 9
        for(int i = 0; i < n; i++) c[x[i] = s[i]]++;
10
        for(int i = 1; i < m; i++) c[i] += c[i - 1];</pre>
11
        for(int i = n - 1; i \ge 0; i--) sa[--c[x[i]]] = i;
12
        for(int k = 1; k <= n; k <<= 1)</pre>
13
14
             int p = 0;
15
             for(int i = n - k; i < n; i++) y[p++] = i;
16
            for(int i = 0; i < n; i++) if(sa[i] >= k) y[p++] = sa[i] - k;
            for(int i = 0; i < m; i++) c[i] = 0;</pre>
17
```

```
18
            for(int i = 0; i < n; i++) c[x[y[i]]]++;</pre>
19
            for(int i = 0; i < m; i++) c[i] += c[i - 1];
20
            for(int i = n - 1; i \ge 0; i--) sa[--c[x[y[i]]]] = y[i];
            swap(x, y);
21
22
            p = 1; x[sa[0]] = 0;
23
            for(int i = 1; i < n; i++)</pre>
                x[sa[i]] = y[sa[i - 1]] == y[sa[i]] && y[sa[i - 1] + k] == y[sa[i] + k] ? p - 1 : p++;
24
25
            if(p >= n) break;
26
            m = p;
27
        }
28 }
29
30 int rk[maxn], height[maxn];
31
32 void getHeight()
33 {
34
        for(int i = 1; i <= n; i++) rk[sa[i]] = i;</pre>
35
        for(int i = 0, k = 0; i < n; i++)</pre>
36
        {
37
            if(k) k--;
38
            int j = sa[rk[i] - 1];
            while(s[i + k] == s[j + k]) k++;
39
40
            height[rk[i]] = k;
41
        }
42 }
43
44 int dp[maxn][20];
45
   void RMQ()
46
47
   {
48
        for(int i = 1; i <= n; i ++) dp[i][0] = height[i];</pre>
49
        for(int j = 1; (1 << j) < maxn; j ++)
50
            for(int i = 1; i + (1 << j) - 1 <= n; i ++)
                dp[i][j] = min(dp[i][j-1], dp[i+(1 << (j-1))][j-1]);
51
52 }
53
54 int query(int 1, int r)
55 {
56
        int k = 0;
        while((1 << (k + 1)) <= r - 1 + 1) k ++;</pre>
57
        return min(dp[l][k], dp[r - (1 << k) + 1][k]);</pre>
58
59 }
60
61 int lcp(int x, int y)
62 {
63
        x = rk[x], y = rk[y];
64
        if(x > y) swap(x, y);
65
        return query(x + 1, y);
66 }
```

5.6 PalindromicTree

5.6.1 PalindromicTree

```
1 const int maxn = 2e6+6;
2 const int N = 26;
3 const int mod = 51123987;
4
5 struct Palindromic_Tree {
```

```
vector<pair<int, int> > next[maxn];
7
       int next[maxn][N];//next指针, next指针和字典树类似, 指向的串为当前串两端加上同一个字符构成
       int fail[maxn]{};//fail指针,失配后跳转到fail指针指向的节点
8
9
       int cnt[maxn]{}; //表示节点i表示的本质不同的串的个数 (建树时求出的不是完全的, 最后count()函数跑一遍
       以后才是正确的)
10
       int num[maxn]{}; //表示以节点i表示的最长回文串的最右端点为回文串结尾的回文串个数
       int len[maxn]{};//len[i]表示节点i表示的回文串的长度(一个节点表示一个回文串)
11
12
       int S[maxn]{};//存放添加的字符
       int last{};//指向新添加一个字母后所形成的最长回文串表示的节点。
13
       int n{};//表示添加的字符个数。
14
15
       int p{};//表示添加的节点个数。
16
       //0向前加,1向后加字符
17
       //int last[2];
18
       //int lpos, rpos;
19
20
       int newnode(int 1) {//新建节点
21
   //
            next[p].clear();
22
          for (int i = 0; i < N; ++i) next[p][i] = 0;</pre>
23
          cnt[p] = 0;
24
          num[p] = 0;
25
          len[p] = 1;
26
          return p++;
      }
27
28
29
       void init() {//初始化
30
          n = last = p = 0;
31
          newnode(0);
32
          newnode(-1);
33
          S[n] = -1;//开头放一个字符集中没有的字符,减少特判
34
          fail[0] = 1;
35
          // lpos 为字符串最大长度
36
          // last[0] = last[1] = 0;
37
          // lpos = 100000, rpos = lpos - 1;
38
          // S[lpos - 1] = S[rpos + 1] = -1;
39
      }
40
       int get_fail(int x) {//和KMP一样, 失配后找一个尽量最长的
41
42
          // op 0 向前, 1 向后
43
          // if (op == 0) while (S[lpos + len[x] + 1] != S[lpos]) x = fail[x];
          // else while(S[rpos - len[x] - 1] != S[rpos]) x = fail[x];
44
          while (S[n - len[x] - 1] != S[n]) x = fail[x];
45
46
          return x;
47
       }
48
49 //
        int find(int u, int c) {
50 //
            vector<pair<int, int> > & x = next[u];
51 //
            int sz = x.size();
52 //
            for(int i = 0; i < sz; ++i) {
53 //
                if(x[i].first == c) return x[i].second;
54 //
55 //
            return 0;
56 //
        }
57
       int add(int c) {
58
59
          // 注意清空左右字符
60
          // if (op == 0) S[--lpos] = c, S[lpos - 1] = -1;
          // else S[++rpos] = c, S[rpos + 1] = -1;
61
62
          S[++n] = c;
63
          int cur = get_fail(last);//通过上一个回文串找这个回文串的匹配位置
```

```
64 //
              int x = find(cur, c);
65
   //
              if (!x) {
            if (!next[cur][c]) {//如果这个回文串没有出现过,说明出现了一个新的本质不同的回文串
66
67
                int now = newnode(len[cur] + 2);//新建节点
68 //
                  x = now;
69 //
                  fail[now] = find(get_fail(fail[cur]), c);
70 //
                  next[cur].emplace_back(make_pair(c, now));
                fail[now] = next[get_fail(fail[cur])][c];//和AC自动机一样建立fail指针,以便失配后跳转
71
72
                next[cur][c] = now;
73
                num[now] = num[fail[now]] + 1;
74
75 //
              last = x;
76
            // 修改最终长度
77
            // if (len[last[op]] == rpos - lpos + 1) last[op ^ 1] = last[op];
78
            last = next[cur][c];
79
            cnt[last]++;
80
            return num[last];
        }
81
82
        void count() {
83
            for (int i = p - 1; i >= 0; --i) cnt[fail[i]] += cnt[i];
84
            //父亲累加儿子的cnt,因为如果fail[v]=u,则u一定是v的子回文串!
85
        }
86
87 } solve;
88
89
    char s[maxn];
90
91
    // 求相交回文串数量
92
    11 a[maxn], b[maxn];
93
    int main() {
94
        solve.init();
        int n;
95
96
        scanf("%d", &n);
        scanf("%s", s);
97
98
        for (int i = 0; i < n; ++i) a[i] = solve.add(s[i] - 'a');</pre>
99
        solve.init();
        for (int i = n - 1; i \ge 0; --i) b[i] = (b[i + 1] + solve.add(s[i] - 'a')) % mod;
100
101
        ll res = (b[0] * (b[0] - 1) / 2) \% mod;
102
        for (int i = 0; i < n; ++i) res = ((res - (a[i] * b[i + 1]) + mod) % mod) % mod;
103
        printf("%lld\n", res);
104
        return 0;
105 }
    5.7 Hash
    5.7.1 hash
 1 // hash常用素数
 2 // 61, 83, 113, 151, 211
 3 // 91815541, 38734667, 68861641
 4 // 917120411, 687840301, 386910137, 515880193
 5 // 1222827239, 1610612741
 6
 7
    typedef unsigned long long ull;
 8
    struct mhash {
 9
        // 自然溢出无模数 805306457
 10
        ull base[maxn];
11
        ull hash_index[maxn];
        ull seed; //31, 131
12
```

```
13
        void inithash(ull seedt = 31) {
14
            base[0] = 1;
15
            seed = seedt;
16
            for (int i = 1; i < maxn; ++i) base[i] = base[i - 1] * seed;</pre>
17
        }
18
        void H(char *p, int n) { // from 1 to n
19
            hash_index[0] = 0;
20
            for (int i = 1; i <= n; ++i) hash_index[i] = hash_index[i - 1] * seed + p[i] - 'a';</pre>
21
        }
22
        ull gethash(int s, int e) {
23
            return hash_index[e] - hash_index[s - 1] * base[e - s + 1];
24
25
   };
26
27
   // 26个素数,解决加法hash
   int prime[] =
        {34183,13513,152993,13591,19687,350869,111187,766091,769297,633469,752273,298651,617191,880421,136067,1408397,7
```

5.7.2 doubleHash

```
namespace Hash{
 1
2
3
        template<class __A,class __B>
4
        class Hash{
        private:
5
6
            static const int size=2000000;
7
            __B *hash; __A *0; int sz;
 8
        public:
9
            Hash(int hash_size=size){ sz=hash_size;
10
                hash=(__B *)malloc(sizeof(__B)*sz);
11
                O=(__A *)malloc(sizeof(__A)*sz);
12
                memset(0,0xff,sizeof(__A)*sz);
13
            }~Hash(){free(0);free(hash);}
14
            __B &operator [](const __A &_0){
15
                int loc=_0%sz;
16
                while(~0[loc] &&0[loc]!=_0){
17
                     ++loc:
18
                     if(loc>sz)loc=0;
19
                }if(!~0[loc])0[loc]=_0;
20
                return hash[loc];
21
            }
22
            void clear(){memset(0,0xff,sizeof(__A)*sz);}
23
        };
24
        struct StringDoubleHashResult{
25
26
            int32_t *H1,*H2,c_len,len;
27
            StringDoubleHashResult(int32_t sz=0){
28
                len=sz; c_len=0; //cur_len;
29
                if(len <= 0){
30
                    H1=H2=0;
31
                    return;
32
                }
33
                H1=(int32_t *)malloc(sizeof(int32_t)*sz);
34
                H2=(int32_t *)malloc(sizeof(int32_t)*sz);
35
36
            ~StringDoubleHashResult(){}
            void clear(){free(H1);free(H2);len=0;H1=H2=0;}
37
```

```
38
             void resize(int new_len){
39
                 int32_t *T1=(int32_t *)malloc(sizeof(int32_t)*new_len);
40
                 int32_t *T2=(int32_t *)malloc(sizeof(int32_t)*new_len);
41
                 for(int i=0;i<c_len;++i)T1[i]=H1[i],T2[i]=H2[i];</pre>
42
                 free(H1);free(H2); H1=T1; H2=T2; len=new_len;
43
            }
44
             void erase(int ers_len){//erase suffix
45
                 c_len-=ers_len;
46
                 if(c_len<0)c_len=0;</pre>
47
            }
48
             //erase prefix not better than reculc
49
        };
50
51
        namespace hash_random{
52
             const int mod_tot=5;
53
             const int mod[]={1000000009,1000000007,998244353,917120411,515880193};
54
        };
55
56
        class StringDoubleHash{
57
        private:
             static const int enable_random=1;
58
59
             int32_t sz,HA1,HA2;
60
            long long B,C;
61
             int32_t *H1,*H2;
62
        public:
63
             StringDoubleHash(int32_t SZ=2e6+5,int32_t ha1=-1,int32_t ha2=-1,int32_t b=-1,int32_t c=-1){
64
                 sz=SZ;
65
                 if(enable_random){
66
                     std::mt19937 rnd(time(0)+19990630);
67
                     int z1= rnd() % hash_random::mod_tot;
                     int z2= (z1 +rnd()%(hash_random::mod_tot - 1) + 1) % hash_random::mod_tot;
68
69
                     if(ha1<0)ha1=hash_random::mod[z1];</pre>
                     if(ha2<0)ha2=hash_random::mod[z2];</pre>
70
71
                     if(b<0)b=rnd()%114514+23333;</pre>
72
                     if(c<0)c=rnd()%1919810+23333;</pre>
73
                 } else {
74
                     if(ha1<0)ha1=1e9+7;</pre>
75
                     if(ha2<0)ha2=1e9+9;
76
                     if(b<0)b=114514;</pre>
77
                     if(c<0)c=1919810;</pre>
                 }
78
79
                 HA1=ha1; HA2=ha2; B=b; C=c;
80
                 //cerr<<HA1<<" "<<HA2<<" "<<B<<" "<<C<<endl;
81
                 H1=(int32_t *)malloc(sizeof(int32_t)*sz);
82
                 H2=(int32_t *)malloc(sizeof(int32_t)*sz);
83
                 init_hash_val();
84
             ~StringDoubleHash(){free(H1);free(H2);}
85
86
             void init_hash_val(){
87
                 H1[0]=H2[0]=1;
88
                 for(int32_t i=1;i<sz;++i){</pre>
89
                     H1[i]=(H1[i-1]*B)%HA1;
90
                     H2[i]=(H2[i-1]*B)%HA2;
                 }
91
92
93
             template <class _Tp>
94
             StringDoubleHashResult culc_hash(const _Tp &s,int32_t len,int32_t tot_len=-1){
95
                 if(tot_len<0)tot_len=len;</pre>
96
                 StringDoubleHashResult R(tot_len);
```

```
97
                  if(len<=0)return R;</pre>
 98
                  R.H1[0] = (s[0]+C)%HA1;
 99
                  R.H2[0] = (s[0]+C)%HA2;
100
                  for(int32_t i=1;i<len;++i){</pre>
101
                      R.H1[i] = (R.H1[i-1]*B+s[i]+C)%HA1;
102
                      R.H2[i]=(R.H2[i-1]*B+s[i]+C)%HA2;
103
                  }
104
                  R.c_len=len;
105
                  return R;
106
              }
107
              // s is the char* first, len is the append length
108
              template <class _Tp>
109
              void append(StringDoubleHashResult &R,const _Tp &s,int32_t len){
110
                  if(len<=0)return;</pre>
111
                  int t_len=R.len;
112
                  while(R.c_len+len>t_len)t_len<<=1;</pre>
113
                  if(t_len>R.len)R.resize(t_len);
114
                  for(int32_t i=R.c_len;i<R.c_len+len;++i){</pre>
115
                      if(i==0){
                          R.H1[i]=(s[i-R.c_len]+C)%HA1;
116
117
                          R.H2[i]=(s[i-R.c_len]+C)%HA2;
118
119
                          R.H1[i]=(R.H1[i-1]*B+s[i-R.c_len]+C)%HA1;
120
                          R.H2[i]=(R.H2[i-1]*B+s[i-R.c_len]+C)%HA2;
121
                      }
122
                  }
123
                  R.c_len+=len;
124
125
              void append(StringDoubleHashResult &R, char s){
126
                  int t_len=R.len;
127
                  while(R.c_len+1>t_len)t_len<<=1;</pre>
128
                  if(t_len>R.len)R.resize(t_len);
129
                  for(int32_t i=R.c_len;i<R.c_len+1;++i){</pre>
130
                      if(i==0){
131
                          R.H1[i]=(s+C)%HA1;
132
                          R.H2[i]=(s+C)%HA2;
133
                      } else {
134
                          R.H1[i] = (R.H1[i-1]*B+s+C)%HA1;
135
                          R.H2[i]=(R.H2[i-1]*B+s+C)%HA2;
                      }
136
                  }
137
138
                  R.c_len+=1;
             }
139
              //return hash [1,r)
140
141
              11 gethash(const StringDoubleHashResult &R, int32_t 1,int32_t r){
142
                  if(l>r||1<0||r-->R.c_len)return -1;//fail
143
                  11 v1=1>0?R.H1[1-1]*(long long)H1[r-1+1]%HA1:0;
144
                  ll v2=1>0?R.H2[1-1]*(long long)H2[r-1+1]%HA2:0;
145
                  v1=R.H1[r]-v1; v2=R.H2[r]-v2;
                  if(v1<0)v1+=HA1; if(v2<0)v2+=HA2;</pre>
146
147
                  return v1<<32|v2;</pre>
148
              }
149
              //merge two hashes as one(s1+s2), but need s2's length
150
              ll merge_hash(const long long &hs1,const long long &hs2,int lenr){
151
                  int32_t m1=hs1>>32,m2=hs1&0xffffffffLL;
152
                  int32_t m3=hs2>>32,m4=hs2&0xffffffffLL;
153
                  m1=m1*(long long)H1[lenr]%HA1+m3;
154
                  if(m1>=HA1)m1-=HA1;
155
                  m2=m2*(long long)H2[lenr]%HA2+m4;
```

```
156
                 if(m2>=HA2)m2-=HA2;
157
                 return (long long)m1<<32|m2;</pre>
             }
158
159
         };
160 };
     5.7.3 二维 hash
 1 #define ull unsigned long long
 ^{2}
    const int maxn = 1005;
    ull hs[maxn][maxn];
 4 char a[maxn][maxn];
 5 int n, m;
 6 ull base1 = 131, base2 = 13331;
 7 ull pwb1[maxn] = {1}, pwb2[maxn] = {1};
 8
 9
    void init() {
10
         for (int i = 1; i < maxn; ++i) {</pre>
11
             pwb1[i] = pwb1[i - 1] * base1;
12
             pwb2[i] = pwb2[i - 1] * base2;
13
         }
    }
14
15
    void Hash() {
16
        for(int i=1;i<=n;i++)</pre>
17
18
             for(int j=1; j<=m; j++)</pre>
19
                 hs[i][j]=hs[i][j-1]*base1+a[i][j] - 'a';
20
         for(int i=1;i<=n;i++)</pre>
21
             for(int j=1; j<=m; j++)</pre>
22
                 hs[i][j]+=hs[i-1][j]*base2;
23 }
24
25 // 右下角(i,j), 行列长度n,m
    ull getHs(int i, int j, int lenn, int lenm) {
         return hs[i][j] - hs[i - lenn][j] * pwb2[lenn] -
27
28
                 hs[i][j - lenm] * pwb1[lenm] +
29
                 hs[i - lenn][j - lenm] * pwb2[lenn] * pwb1[lenm];
30 }
    5.7.4 树 hash 同构
 1 // n=1e5的话base开2e6+9, 可以输出看到top不比n小即可
    const int base = 2e6+9;
    // vis大小要开到素数大小, turn表示当前树的编号, p是预处理数组
 4 int vis[base + 1], top, turn, p[base + 1];
 5 // 程序开头调用一次
 6
    void init() {
 7
         top = 0;
 8
         for (int i = 2; i <= base; ++i) {</pre>
 9
             if (!vis[i]) {
10
                 p[++top] = i;
             }
11
12
             for (int j = 1; j \le top && i * p[j] \le base; ++j) {
13
                 vis[i * p[j]] = 1;
14
                 if (i % p[j] == 0) break;
15
             }
16
         }
17
         assert(top >= maxn);
```

```
18 }
19
20
  vector<int> edge[maxn];
   // h[x]表示x这棵子树的hash值, g[x]表示以x为根的hash值
   int h[maxn], g[maxn], sz[maxn];
23
  struct TreeHash {
24
25
       int n;
26
       // 如果树比较多,在类内部开edge可能会炸内存,可以改到外面做前向星
27
       // 除了hs是答案其他都可以改到外部,只有edge需要清零
28
       // vector<int> edge[maxn];
29
       // int h[maxn], g[maxn], sz[maxn];
30
       vector<int> hs;
31
32
       void init(int n_ = 0) {
33
           n = n_{j}
34
           hs.clear();
35
36
37
       void dfs1(int u, int pre) {
           sz[u] = 1;
38
           h[u] = 1;
39
40
           for (auto v : edge[u]) {
               if (v == pre) continue;
41
42
               dfs1(v, u);
43
               h[u] = (h[u] + 111 * h[v] * p[sz[v]] % mod) % mod;
               sz[u] += sz[v];
44
           }
45
       }
46
47
48
       void dfs2(int u, int pre, int V, int needres = 1) {
           g[u] = (h[u] + 111 * V * p[n - sz[u]] % mod) % mod;
49
50
           if (needres) hs.push_back(g[u]);
51
           for (auto v : edge[u]) {
               if (v == pre) continue;
52
53
               dfs2(v, u, (g[u] - 1]l * h[v] * p[sz[v]] % mod + mod) % mod);
54
           }
55
       }
56
57
       void work(int needres = 1) {
           // 无根树选一个不存在的点当pre即可, 当多棵无根树判重时需要sort
58
59
           dfs1(1, 0);
60
           dfs2(1, 0, 0, needres);
           sort(hs.begin(), hs.end());
61
62
       }
63 };
64
65 // 获取删掉某叶子节点后以与该叶子节点相邻点开头的hash值
66 // \text{ int res} = (hs[edge[i][0]] - 2 + mod) \% mod;
   5.8 Suffix Automation
   5.8.1 SAM
1 const int maxn = 2e4 + 10;
2
3 struct SuffixAutomation
4
   {
5
       int last, cnt;
```

```
6
        int ch[maxn << 1][26], fa[maxn << 1], len[maxn << 1], pos[maxn << 1];</pre>
7
       int sz[maxn << 1], a[maxn << 1], c[maxn << 1];</pre>
8
9
       void init()
10
       {
11
           last = cnt = 1;
12
           memset(ch[1], 0, sizeof ch[1]);
13
           fa[1] = len[1] = 0;
14
       }
15
       int inline newnode(int idx)
16
17
       {
18
           ++cnt;
19
           memset(ch[cnt], 0, sizeof ch[cnt]);
20
           fa[cnt] = len[cnt] = 0;
21
           pos[cnt] = idx;
22
           return cnt;
23
24
25
       void ins(int c)
26
27
           int p = last , np = newnode(pos[last] + 1);
28
           last = np, len[np] = len[p] + 1;
29
           for(; p && !ch[p][c]; p = fa[p]) ch[p][c] = np;
30
           if(!p) fa[np] = 1;
31
           else
32
           {
33
               int q = ch[p][c];
34
               if(len[p] + 1 == len[q]) fa[np] = q;
35
               else
36
               {
37
                   int nq = newnode(pos[p] + 1);
38
                   len[nq] = len[p] + 1;
39
                   memcpy(ch[nq], ch[q], sizeof ch[q]);
40
                   fa[nq] = fa[q], fa[q] = fa[np] = nq;
41
                   for(; ch[p][c] == q; p = fa[p]) ch[p][c] = nq;
42
               }
43
           }
44
           sz[np] = 1;
       }
45
46
47
       int solve(int n)
48
49
           /*求两个串的LCS:
               对一个字符串建立SAM,记录一个当前匹配的长度Len和当前节点v,枚举另一个字符串的每个字符;
50
               如果p有字符v的转移边出边,则使Len加一,并使p转移到出边指向的节点上;
51
                否则不断向父节点上跳,直到当前节点有字符p的转移出边,或者跳到根节点;
52
53
54
           int p = 1, ans = 0, now_len = 0;
           for(int i = 0; s2[i]; i ++)
55
56
57
               if(ch[p][s2[i] - 'a']) p = ch[p][s2[i] - 'a'], now_len ++;
58
               else
               {
59
60
                   for(;p && !ch[p][s2[i] -'a'] ; p = fa[p]) ;
61
                   if(p == 0) now_len = 0, p = 1;
62
                   else now_len = len[p] + 1, p = ch[p][s2[i] - 'a'];
63
64
               ans = max(now_len, ans);
```

```
65
            }
        }
66
67
        void Toposort()
68
69
70
            long long ans = 0;
71
            for(int i = 1; i <= cnt; i ++) c[len[i]] ++;</pre>
72
            for(int i = 1; i <= cnt; i ++) c[i] += c[i - 1];</pre>
73
            for(int i = 1; i <= cnt; i ++) a[c[len[i]] --] = i;</pre>
74
            for(int i = cnt; i; i --) sz[fa[a[i]]] += sz[a[i]];
75
76 }sam;
    5.9 Others
```

5.9.1 最小表示法

```
1 // 0起始
   int Gao(char a[], int len) {
3
        int i = 0, j = 1, k = 0;
        while (i < len && j < len && k < len) {
4
5
            int cmp = a[(j + k) % len] - a[(i + k) % len];
            if (cmp == 0) k++;
6
7
            else {
8
                if (cmp > 0) j += k + 1;
9
                else i += k + 1;
10
                if (i == j) j ++;
11
                k = 0;
12
            }
        }
13
14
        return min(i, j);
15 }
```

6 dp

6.1 BitDP

6.1.1 数位 dp 计和

```
1 #include <bits/stdc++.h>
2 #define 11 long long
3 using namespace std;
4 const int mod = 998244353;
5 pair<11, 11> dp[20][1<<10];
6 bool vis[20][1<<10];
7 int k;
8 int t[20];
9 11 base[20];
10
11 pair<11, 11> dfs(int pos, int state, bool limit, bool lead) {
        if (pos == -1) return _ builtin popcount(state) <= k ? make pair(1, 0) : make pair(0, 0);
        if (!limit && !lead && vis[pos][state]) return dp[pos][state];
13
14
        int up = limit ? t[pos] : 9;
       pair<11, 11> res = {0, 0};
15
        for (int i = 0; i <= up; ++i) {</pre>
16
            int n_s = state;
17
18
            if (lead && i == 0) n_s = 0;
19
            else n_s = state | (1 << i);</pre>
20
            auto tmp = dfs(pos - 1, n_s, limit && i == t[pos], lead && i == 0);
21
            11 pre = 111 * i * base[pos] % mod;
22
            (res.first += tmp.first) %= mod;
23
            (res.second += tmp.second + pre * tmp.first) %= mod;
24
        }
25
        if (!limit && !lead) dp[pos][state] = res, vis[pos][state] = 1;
        return res;
27 }
28
29 11 solve(11 x) {
30
        int pos = 0;
31
        do {
            t[pos ++] = x % 10;
32
33
        } while (x /= 10);
34
        return dfs(pos - 1, 0, true, true).second;
35 }
36
37 int main(int argc, char *argv[])
38 {
39
        base[0] = 1;
        for (int i = 1; i < 20; ++i) base[i] = base[i - 1] * 10;</pre>
40
41
42
        scanf("%lld%lld%d", &1, &r, &k);
        printf("%lld\n", (solve(r) - solve(l - 1) + mod) % mod);
43
44
        return 0;
45 }
    6.1.2 两个数数位 dp
1 // 二进制数位dp, 求a $\in$ 1~x 和 b $\in$ 1~y, 满足 $a & b > c || a ^ b < c$的对数
2 11 dp[maxn][2][2][2][2];
3 int a[maxn], b[maxn], c[maxn];
4
5
```

```
void cal(int *xt, ll x) {
7
        int has = 0;
8
        while (x) {
            xt[has++] = x % 2;
9
10
            x /= 2;
11
        }
12 }
13
    11 dfs(int pos, int o1, int o2, int lim1, int lim2) {
14
        if (pos < 0) return 1;</pre>
        11 &t = dp[pos][01][02][lim1][lim2];
16
17
        if (t != -1) return t;
18
        int up1 = o1 ? a[pos] : 1;
        int up2 = o2 ? b[pos] : 1;
19
20
        11 \text{ res} = 0;
        for (int i = 0; i <= up1; ++i) {</pre>
21
22
            for (int j = 0; j <= up2; ++j) {</pre>
23
                 int t1 = i & j;
24
                 int t2 = i ^ j;
25
                 if (lim1 && t1 > c[pos]) continue;
26
                 if (lim2 && t2 < c[pos]) continue;</pre>
27
                 res += dfs(pos - 1, o1 && i == up1, o2 && j == up2, lim1 && t1 == c[pos], lim2 && t2 ==
         c[pos]);
28
            }
29
30
        return t = res;
31 }
32
    11 solve(ll x, ll y, ll z) {
33
34
        memset(dp, -111, sizeof dp);
35
        for (int i = 0; i < 33; ++i) a[i] = b[i] = c[i] = 0;
36
        cal(a, x);
37
        cal(b, y);
38
        cal(c, z);
39
        return dfs(32, 1, 1, 1, 1);
40 }
41
42
    int main(int argc, char *argv[]) {
43
        int T;
        scanf("%d", &T);
44
45
        11 x, y, z;
46
        for (int kase = 1; kase <= T; ++kase) {</pre>
47
            scanf("%lld%lld%lld", &x, &y, &z);
48
            11 \text{ res} = \text{solve}(x, y, z);
49
            res -= max(011, y - z + 1);
50
            res -= max(011, x - z + 1);
51
            printf("%lld\n", x * y - res);
        }
52
53
        return 0;
54 }
    6.2
         Subsequence
```

6.2.1 MaxSum

```
    // 传入序列a和长度n, 返回最大子序列和
    int MaxSeqSum(int a[], int n)
    {
    int rt = 0, cur = 0;
```

```
5
        for (int i = 0; i < n; i++)</pre>
6
           cur += a[i], rt = max(cur, rt), cur = max(0, cur);
7
       return rt;
8
   }
   6.2.2 LIS
1 // 简单写法(下标从0开始,只返回长度)
   int dp[N];
3 int LIS(int a[], int n)
4 {
5
       memset(dp, 0x3f, sizeof(dp));
6
        for (int i = 0; i < n; i++) *lower_bound(dp, dp + n, a[i]) = a[i];</pre>
7
        return lower_bound(dp, dp + n, INF) - dp;
8
   }
9
10 // 小常数nlogn求序列用树状数组维护dp即可
11 // dp[i] = max(dp[j]) + 1 (j < i && a[j] < a[i])
   6.2.3 LongestCommonIncrease
1 // 序列下标从1开始
2
   int LCIS(int a[], int b[], int n, int m)
3
4
       memset(dp, 0, sizeof(dp));
5
       for (int i = 1; i <= n; i++)</pre>
6
7
           int ma = 0;
8
           for (int j = 1; j \le m; j++)
9
10
               dp[i][j] = dp[i - 1][j];
               if (a[i] > b[j]) ma = max(ma, dp[i - 1][j]);
11
12
               if (a[i] == b[j]) dp[i][j] = ma + 1;
13
           }
14
       }
15
       return *max_element(dp[n] + 1, dp[n] + 1 + m);
16 }
   6.2.4 LCS
   #include <stdio.h>
   #include <string.h>
3
4 #define M 30005
5 #define SIZE 128
6 #define WORDMAX 3200
7 #define BIT 32
8
9 char s1[M], s2[M];
10 int nword;
11 unsigned int str[SIZE][WORDMAX];
12 unsigned int tmp1[WORDMAX], tmp2[WORDMAX];
13
14 void pre(int len)
15 {
16
       int i, j;
17
       memset(str, 0, sizeof(str));
```

```
18
        for(i = 0; i < len; i ++)</pre>
19
            str[s1[i]][i / BIT] |= 1 << (i % BIT);
20 }
21
22 void cal(unsigned int *a, unsigned int *b, char ch)
23 {
24
        int i, bottom = 1, top;
25
        unsigned int x, y;
26
        for(i = 0; i < nword; i ++)</pre>
27
28
            y = a[i];
            x = y | str[ch][i];
29
30
            top = (y >> (BIT - 1)) & 1;
            y = (y \ll 1) \mid bottom;
31
            if(x < y) top = 1;
32
            b[i] = x & ((x - y) ^ x);
33
34
            bottom = top;
35
        }
36 }
37
38 int bitcnt(unsigned int *a)
39 {
40
        int i, j, res = 0, t;
41
        unsigned int b[5] = {0x555555555, 0x33333333, 0x0f0f0f0f, 0x000ff00ff, 0x0000ffff}, x;
        for(i = 0; i < nword; i ++)</pre>
42
43
            x = a[i];
44
45
            t = 1;
46
            for(j = 0; j < 5; j ++, t <<= 1)
47
                x = (x \& b[j]) + ((x >> t) \& b[j]);
48
            res += x;
49
        }
50
        return res;
51 }
52
53 void process()
54 {
55
        int i, j, len1, len2;
56
        unsigned int *a, *b, *t;
        len1 = strlen(s1);
57
58
        len2 = strlen(s2);
        nword = (len1 + BIT - 1) / BIT;
59
60
        pre(len1);
61
        memset(tmp1, 0, sizeof(tmp1));
        a = &tmp1[0];
62
63
        b = &tmp2[0];
64
        for(i = 0; i < len2; i ++)</pre>
65
        {
66
            cal(a, b, s2[i]);
67
            t = a; a = b; b = t;
68
69
        printf("%d\n", bitcnt(a));
70 }
71
72 int main()
73 {
74
        while(scanf("%s%s", s1, s2) != EOF)
            process();
75
76
        return 0;
```

77 }

6.3 Others

问题 设 $f(i) = \min(y[k] - s[i] \times x[k]), k \in [1, i-1]$, 现在要求出所有 $f(i), i \in [1, n]$ 考虑两个决策 j 和 k, 如果 j 比 k 优,则

$$y[j] - s[i] \times x[j] < y[k] - s[i] \times x[k]$$

化简得:

$$\frac{y_j - y_k}{x_i - x_k} < s_i$$

不等式左边是个斜率,我们把它设为 slope(j,k)

我们可以维护一个单调递增的队列,为什么呢?

因为如果 slope(q[i-1],q[i]) > slope(q[i],q[i+1]),那么当前者成立时,后者必定成立。即 q[i] 决策优于 q[i-1] 决策时,q[i+1] 必然优于 q[i],因此 q[i] 就没有存在的必要了。所以我们要维护递增的队列。

那么每次的决策点 i, 都要满足

$$\begin{cases} \operatorname{slope}(q[i-1], q[i]) < s[i] \\ \operatorname{slope}(q[i], q[i+1]) \ge s[i] \end{cases}$$

一般情况去二分这个 i 即可。

如果 s[i] 是单调不降的,那么对于决策 j 和 k(j < k) 来说,如果决策 k 优于决策 j,那么对于 $i \in [k+1,n]$,都存在决策 k 优于决策 j,因此决策 j 就可以舍弃了。这样的话我们可以用单调队列进行优化,可以少个 \log 。

单调队列滑动窗口最大值

```
1 // k为滑动窗口的大小,数列下标从1开始,d为序列长度+1
   deque<int> q;
3
   for (int i = 0, j = 0; i + k \le d; i++)
4
5
       while (j < i + k)
6
7
           while (!q.empty() && a[q.back()] < a[j]) q.pop_back();</pre>
8
           q.push_back(j++);
9
10
       while (q.front() < i) q.pop_front();</pre>
11
       // a[q.front()]为当前滑动窗口的最大值
12 }
```

6.3.1 矩阵快速幂

```
struct Matrix {
2
        int sz;
3
        // int n, m;
 4
        11 a[maxn] [maxn];
        Matrix(int sz_ = 0):sz(sz_) {
5
            memset(a, 0, sizeof a);
6
 7
        }
 8
        void pr() {
            printf("*\n");
9
10
            for(int i = 0; i < sz; ++i) {</pre>
11
                 for (int j = 0; j < sz; ++j) {
12
                     printf("%lld ", a[i][j]);
13
                 printf("\n");
14
15
            }
16
        void tr() {
17
```

```
18
            for (int i = 0; i < sz; ++i) {</pre>
19
                for (int j = i + 1; j < sz; ++j) {
20
                     swap(a[i][j], a[j][i]);
21
22
            }
23
        }
24
   }res, t1;
25
26
   void init() {
27
28 }
29
30 Matrix mul(Matrix a, Matrix b)
31 {
32
        Matrix res(a.sz);
        // if (a.m != b.n) return res;
33
34
        for(int i = 0; i < res.sz; i++) // a.n</pre>
35
            for(int j = 0; j < res.sz; j++) // b.m
                for(int k = 0; k < res.sz; k++) // a.m, b.n</pre>
36
37
                     (res.a[i][j] += a.a[i][k] * b.a[k][j] % mod) %= mod;
38
        return res;
39 }
40
41 Matrix pow(ll n)
42 {
43
        init();
        //for(int i = 0; i < cur; i++) res.a[i][i] = 1;
44
45
        while(n > 0) {
            if(n & 1) res = mul(res, t1);
46
47
            t1 = mul(t1, t1);
48
            n >>= 1;
49
        }
50
        return res;
51 }
```

7 Others

7.1 mint 类

```
const int mod = 998244353;
3
   struct mint {
     int n;
4
     mint(int n_ = 0) : n(n_) {}
5
6
  };
7
8 mint operator+(mint a, mint b) { return (a.n += b.n) >= mod ? a.n - mod : a.n; }
9 mint operator-(mint a, mint b) { return (a.n -= b.n) < 0 ? a.n + mod : a.n; }
10 mint operator*(mint a, mint b) { return 1LL * a.n * b.n % mod; }
11 mint &operator+=(mint &a, mint b) { return a = a + b; }
12 mint & operator = (mint &a, mint b) { return a = a - b; }
13 mint &operator*=(mint &a, mint b) { return a = a * b; }
14 ostream &operator << (ostream &o, mint a) { return o << a.n; }
```

7.2 不重叠区间贪心

```
1 #include <bits/stdc++.h>
2 #define 11 long long
3 using namespace std;
4
5 const int maxn = 5e5+5;
6 pair<int, int> a[maxn];
   int main() {
7
8
        int n;
9
        cin >> n;
10
        for (int i = 1; i <= n; ++i) {</pre>
11
            cin >> a[i].second >> a[i].first;
12
13
        sort(a + 1, a + 1 + n);
14
        int res = 1;
        int tmp = a[1].first;
15
          printf("%d %d\n", a[1].second, a[1].first);
16
17
        for (int i = 2; i <= n; ++i) {</pre>
18
            if (a[i].second > tmp) {
                  printf("%d %d\n", a[i].second, a[i].first);
20 //
21
                tmp = a[i].first;
            }
22
23
24
        printf("%d\n", res);
25
        return 0;
26 }
```

7.3 BigInt 类

```
8
        Complex operator-(const Complex &b)const{
 9
            return Complex(x - b.x,y - b.y);
10
11
        Complex operator+(const Complex &b)const{
12
            return Complex(x + b.x,y + b.y);
13
        }
14
        Complex operator*(const Complex &b)const{
15
            return Complex(x*b.x - y*b.y,x*b.y + y*b.x);
16
17
    };
    void change(Complex y[],int len){
18
19
        int i,j,k;
20
        for(int i = 1,j = len/2;i<len-1;i++){</pre>
21
            if(i < j)
                          swap(y[i],y[j]);
22
            k = len/2;
23
             while(j \ge k){
                 j = j - k;
24
                 k = k/2;
25
26
            }
27
            if(j < k)
                          j+=k;
28
        }
29 }
    void fft(Complex y[],int len,int on){
31
        change(y,len);
32
        for(int h = 2;h <= len;h<<=1){</pre>
33
             Complex wn(cos(on*2*PI/h),sin(on*2*PI/h));
34
             for(int j = 0; j < len; j += h){
35
                 Complex w(1,0);
36
                 for(int k = j; k < j + h/2; k++){
37
                     Complex u = y[k];
38
                     Complex t = w*y[k + h/2];
39
                     y[k] = u + t;
40
                     y[k + h/2] = u - t;
41
                     w = w*wn;
42
                 }
            }
43
        }
44
45
        if(on == -1){
46
            for(int i = 0;i < len;i++){</pre>
                 y[i].x /= len;
47
48
        }
49
50 }
   class BigInt
51
52 {
    #define Value(x, nega) ((nega) ? -(x) : (x))
    #define At(vec, index) ((index) < vec.size() ? vec[(index)] : 0)</pre>
54
        static int absComp(const BigInt &lhs, const BigInt &rhs)
55
56
57
             if (lhs.size() != rhs.size())
58
                 return lhs.size() < rhs.size() ? -1 : 1;</pre>
59
            for (int i = lhs.size() - 1; i >= 0; --i)
60
                 if (lhs[i] != rhs[i])
                     return lhs[i] < rhs[i] ? -1 : 1;</pre>
61
62
             return 0;
63
64
        using Long = long long;
65
        const static int Exp = 9;
66
        const static Long Mod = 1000000000;
```

```
67
         mutable std::vector<Long> val;
 68
         mutable bool nega = false;
 69
         void trim() const
 70
         {
 71
             while (val.size() && val.back() == 0)
 72
                 val.pop_back();
 73
             if (val.empty())
 74
                 nega = false;
 75
         }
 76
         int size() const { return val.size(); }
 77
         Long &operator[](int index) const { return val[index]; }
 78
         Long &back() const { return val.back(); }
 79
         BigInt(int size, bool nega) : val(size), nega(nega) {}
         BigInt(const std::vector<Long> &val, bool nega) : val(val), nega(nega) {}
 80
 81
 82
     public:
 83
         friend std::ostream &operator<<(std::ostream &os, const BigInt &n)</pre>
 84
 85
             if (n.size())
 86
             {
 87
                 if (n.nega)
                      putchar('-');
 88
 89
                 for (int i = n.size() - 1; i >= 0; --i)
 90
 91
                      if (i == n.size() - 1)
 92
                          printf("%lld", n[i]);
                      else
 93
 94
                          printf("%0*lld", n.Exp, n[i]);
                 }
 95
             }
 96
 97
             else
 98
                 putchar('0');
99
             return os;
100
         }
101
         friend BigInt operator+(const BigInt &lhs, const BigInt &rhs)
102
103
             BigInt ret(lhs);
104
             return ret += rhs;
105
         }
106
         friend BigInt operator-(const BigInt &lhs, const BigInt &rhs)
107
108
             BigInt ret(lhs);
109
             return ret -= rhs;
110
         }
111
         BigInt(Long x = 0)
112
         {
113
             if (x < 0)
114
                 x = -x, nega = true;
115
             while (x \ge Mod)
116
                 val.push_back(x % Mod), x /= Mod;
117
118
                 val.push_back(x);
119
         }
120
         BigInt(const char *s)
121
122
             int bound = 0, pos;
123
             if (s[0] == '-')
124
                 nega = true, bound = 1;
125
             Long cur = 0, pow = 1;
```

```
126
             for (pos = strlen(s) - 1; pos >= Exp + bound - 1; pos -= Exp, val.push_back(cur), cur = 0,
         pow = 1
127
                 for (int i = pos; i > pos - Exp; --i)
128
                      cur += (s[i] - '0') * pow, pow *= 10;
129
             for (cur = 0, pow = 1; pos >= bound; --pos)
130
                 cur += (s[pos] - '0') * pow, pow *= 10;
131
             if (cur)
132
                 val.push_back(cur);
133
         }
134
         BigInt &operator=(const char *s){
135
             BigInt n(s);
136
             *this = n;
137
             return n;
         }
138
139
         BigInt &operator=(const Long x){
140
             BigInt n(x);
141
             *this = n;
142
             return n;
143
         }
         friend std::istream &operator>>(std::istream &is, BigInt &n){
144
145
             string s;
             is >> s;
146
147
             n=(char*)s.data();
148
             return is;
149
         }
150
         BigInt &operator+=(const BigInt &rhs)
151
152
             const int cap = std::max(size(), rhs.size()) + 1;
153
             val.resize(cap);
             int carry = 0;
154
155
             for (int i = 0; i < cap - 1; ++i)</pre>
156
157
                 val[i] = Value(val[i], nega) + Value(At(rhs, i), rhs.nega) + carry, carry = 0;
158
                 if (val[i] >= Mod)
                      val[i] -= Mod, carry = 1;
159
160
                 else if (val[i] < 0)</pre>
161
                      val[i] += Mod, carry = -1;
162
             }
163
             if ((val.back() = carry) == -1) //assert(val.back() == 1 or 0 or -1)
164
165
                 nega = true, val.pop_back();
166
                 bool tailZero = true;
167
                 for (int i = 0; i < cap - 1; ++i)</pre>
168
169
                      if (tailZero && val[i])
170
                          val[i] = Mod - val[i], tailZero = false;
171
                      else
172
                          val[i] = Mod - 1 - val[i];
                 }
173
174
             }
175
             trim();
176
             return *this;
177
         }
178
         friend BigInt operator-(const BigInt &rhs)
179
180
             BigInt ret(rhs);
181
             ret.nega ^= 1;
182
             return ret;
183
         }
```

```
184
         BigInt &operator == (const BigInt &rhs)
185
186
              rhs.nega ^= 1;
187
              *this += rhs;
188
              rhs.nega ^= 1;
189
             return *this;
190
         }
         friend BigInt operator*(const BigInt &lhs, const BigInt &rhs)
191
192
193
              int len=1;
194
              BigInt ll=lhs,rr=rhs;
195
              11.nega = lhs.nega ^ rhs.nega;
196
              while(len<2*lhs.size()||len<2*rhs.size())len<<=1;</pre>
197
              ll.val.resize(len),rr.val.resize(len);
198
              Complex x1[len],x2[len];
199
              for(int i=0;i<len;i++){</pre>
200
                  Complex nx(ll[i],0.0),ny(rr[i],0.0);
201
                  x1[i]=nx;
202
                  x2[i]=ny;
203
             }
204
             fft(x1,len,1);
205
              fft(x2,len,1);
206
              for(int i = 0 ; i < len; i++)</pre>
207
                  x1[i] = x1[i] * x2[i];
208
              fft( x1 , len , -1 );
209
              for(int i = 0 ; i < len; i++)</pre>
210
                  11[i] = int(x1[i].x + 0.5);
211
              for(int i = 0 ; i < len; i++){</pre>
212
                  ll[i+1]+=ll[i]/Mod;
213
                  11[i]%=Mod;
214
215
             11.trim();
216
             return 11;
217
         }
218
         friend BigInt operator*(const BigInt &lhs, const Long &x){
219
             BigInt ret=lhs;
220
              bool negat = (x < 0);
221
             Long xx = (negat) ? -x : x;
222
             ret.nega ^= negat;
223
             ret.val.push_back(0);
224
              ret.val.push_back(0);
225
              for(int i = 0; i < ret.size(); i++)</pre>
226
                  ret[i]*=xx;
227
              for(int i = 0; i < ret.size(); i++){</pre>
228
                  ret[i+1]+=ret[i]/Mod;
229
                  ret[i] %= Mod;
230
             }
231
              ret.trim();
232
              return ret;
233
234
         BigInt &operator*=(const BigInt &rhs) { return *this = *this * rhs; }
235
         BigInt &operator*=(const Long &x) { return *this = *this * x; }
236
         friend BigInt operator/(const BigInt &lhs, const BigInt &rhs)
237
238
              static std::vector<BigInt> powTwo{BigInt(1)};
239
              static std::vector<BigInt> estimate;
240
              estimate.clear();
241
              if (absComp(lhs, rhs) < 0)</pre>
242
                  return BigInt();
```

```
243
             BigInt cur = rhs;
244
             int cmp;
245
             while ((cmp = absComp(cur, lhs)) <= 0)</pre>
246
             {
247
                 estimate.push_back(cur), cur += cur;
248
                 if (estimate.size() >= powTwo.size())
249
                      powTwo.push_back(powTwo.back() + powTwo.back());
250
             }
251
             if (cmp == 0)
252
                 return BigInt(powTwo.back().val, lhs.nega ^ rhs.nega);
253
             BigInt ret = powTwo[estimate.size() - 1];
254
             cur = estimate[estimate.size() - 1];
255
             for (int i = estimate.size() - 1; i >= 0 && cmp != 0; --i)
256
                 if ((cmp = absComp(cur + estimate[i], lhs)) <= 0)</pre>
257
                      cur += estimate[i], ret += powTwo[i];
             ret.nega = lhs.nega ^ rhs.nega;
258
259
             return ret;
260
         }
261
         friend BigInt operator/(const BigInt &num,const Long &x){
262
             bool negat = (x < 0);
263
             Long xx = (negat) ? -x : x;
264
             BigInt ret;
265
             Long k = 0;
266
             ret.val.resize( num.size() );
267
             ret.nega = (num.nega ^ negat);
268
             for(int i = num.size() - 1; i >= 0; i--){
269
                 ret[i] = (k * Mod + num[i]) / xx;
270
                 k = (k * Mod + num[i]) % xx;
271
             }
272
             ret.trim();
273
             return ret;
274
         }
275
         bool operator==(const BigInt &rhs) const
276
         {
277
             return nega == rhs.nega && val == rhs.val;
         }
278
279
         bool operator!=(const BigInt &rhs) const { return nega != rhs.nega || val != rhs.val; }
280
         bool operator>=(const BigInt &rhs) const { return !(*this < rhs); }</pre>
281
         bool operator>(const BigInt &rhs) const { return !(*this <= rhs); }</pre>
282
         bool operator<=(const BigInt &rhs) const</pre>
283
284
             if (nega && !rhs.nega)
285
                 return true;
286
             if (!nega && rhs.nega)
287
                 return false;
288
             int cmp = absComp(*this, rhs);
289
             return nega ? cmp >= 0 : cmp <= 0;</pre>
290
         }
291
         bool operator<(const BigInt &rhs) const</pre>
292
293
             if (nega && !rhs.nega)
294
                 return true;
295
             if (!nega && rhs.nega)
296
                 return false;
297
             return (absComp(*this, rhs) < 0) ^ nega;</pre>
298
         }
299
         void swap(const BigInt &rhs) const
300
         {
301
             std::swap(val, rhs.val);
```

```
302
             std::swap(nega, rhs.nega);
        }
303
304 };
305 BigInt ba,bb;
306 int main(){
307
         cin>>ba>>bb;
308
         cout << ba + bb << '\n';//和
         cout << ba - bb << '\n';//差
309
         cout << ba * bb << '\n';//积
310
311
        BigInt d;
        cout << (d = ba / bb) << '\n';//商
312
313
         cout << ba - d * bb << '\n';//余
314
        return 0;
315 }
     7.4 date
 1 string dayOfWeek[] = {"Mo", "Tu", "We", "Th", "Fr", "Sa", "Su"};
   // converts Gregorian date to integer (Julian day number)
    int DateToInt (int m, int d, int y){
 3
         return
 4
 5
             1461 * (y + 4800 + (m - 14) / 12) / 4 +
 6
             367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
 7
             3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
 8
             d - 32075;
 9 }
10
11 // converts integer (Julian day number) to Gregorian date: month/day/year
12 void IntToDate (int jd, int &m, int &d, int &y){
        int x, n, i, j;
13
14
        x = jd + 68569;
        n = 4 * x / 146097;
15
16
        x = (146097 * n + 3) / 4;
17
        i = (4000 * (x + 1)) / 1461001;
18
        x = 1461 * i / 4 - 31;
19
         j = 80 * x / 2447;
20
        d = x - 2447 * j / 80;
21
        x = j / 11;
22
        m = j + 2 - 12 * x;
23
        y = 100 * (n - 49) + i + x;
24 }
25 // converts integer (Julian day number) to day of week
26 string IntToDay (int jd){
        return dayOfWeek[jd % 7];
27
28 }
     7.5 Frac 类
    struct Frac {
        ll a, b;
 3
        void getJian() {
             11 \text{ gcd} = abs(\_gcd(a, b));
 4
 5
             a /= gcd;
 6
             b /= gcd;
 7
             if (b < 0) {
 8
                 a = -a;
                 b = -b;
```

10

```
}
        }
11
12
        Frac(ll a_ = 1, ll b_ = 1) {
13
            a = a_;
14
            b = b_;
15
            getJian();
16
        }
17
        Frac add(const Frac& oth) {
            ll bt = b * oth.b;
18
19
            11 at = a * oth.b + oth.a * b;
20
            return Frac(at, bt);
21
        }
22
        Frac multi(const Frac& oth) {
23
            a *= oth.a;
24
            b *= oth.b;
25
            getJian();
26
            return *this;
27
        }
28
        bool operator < (const Frac& oth) const {</pre>
29
            return a * oth.b < b * oth.a;</pre>
30
        }
31
        bool operator == (const Frac& oth) const {
32
            return a * oth.b == b * oth.a;
33
34
        bool operator <= (const Frac& oth) const {</pre>
35
            return a * oth.b <= b * oth.a;</pre>
36
        }
37 };
         模拟退火 (最小圆覆盖)
1 const int maxn = 1e5 + 10;
2 const double eps = 1e-8;
3 const double delta = 0.98;
4 const double inf = 1e18;
5
6 struct Point { double x, y; } p[maxn];
7
8
    double dis(Point A, Point B) { return sqrt((A.x - B.x) * (A.x - B.x) + (A.y - B.y) * (A.y - B.y));
        };
9
10
   double Simulate_Annea(int n)
11
12
        Point S;
13
        S.x = S.y = 0;
        double t = 1000;
14
15
        double res = inf;
16
        while(t > eps)
17
        {
18
19
            for(int i = 0; i < n; i ++) if(dis(S, p[i]) > dis(S, p[k])) k = i;
            double d = dis(S, p[k]);
20
21
            res = min(res, d);
22
            S.x += (p[k].x - S.x) / d * t;
23
            S.y += (p[k].y - S.y) / d * t;
24
            t *= delta;
25
        }
26
        return res;
```

```
27 }
28
29 int main()
30 {
31
        int n;
32
        scanf("%d", &n);
33
        for(int i = 0; i < n; i ++) scanf("%lf%lf", &p[i].x, &p[i].y);</pre>
34
        printf("%.3f\n", Simulate_Annea(n));
35
        return 0;
36 }
    7.7 string 类
 1 const int maxn = 1005;
    struct String{
 3
        int nex[maxn];
 4
        char x[maxn];
 5
        int len;
 6
        int getLength() {
 7
            return len;
 8
 9
        void getNext() {
10
            int n = len, i = 0, j = -1;
            nex[0] = -1;
11
12
            while (i < n) {</pre>
13
                 if (j == -1 \mid | x[i] == x[j]) nex[++i] = ++j;
14
                 else j = -1;
15
            }
16
        }
17
        void input() {
18
            scanf("%s", x);
            len = strlen(x);
19
20
21
        void inputAndCal() {
22
            scanf("%s", x);
23
            len = strlen(x);
24
            getNext();
25
        }
26
        void show() {
            printf("%s\n", x);
27
28
        }
29
        bool operator < (const String&oth) const {</pre>
30
            return strcmp(x, oth.x) < 0;</pre>
31
32
        char operator [] (const int a) const {
33
            return x[a];
34
        bool substring(String b) {//b is the substring of a
35
36
            int m = len, n = b.getLength();
37
            int i = 0, j = 0;
            while (i < m && j < n) \{
38
39
                 if (j == -1 || x[i] == b[j]) ++i, ++j;
                 else j = b.nex[j];
40
41
                 if (j == n) return true;
42
43
            return false;
44
        }
45 };
```

7.8 前缀异或和

```
ll xor_sum(ll n) {
1
2
        11 t=n&3;
3
        if (t&1) return t/2ull^1;
4
        return t/2ull^n;
5 }
    7.9 约瑟夫环第 k 个
1
    ll kth(ll n,ll m,ll k) {
2
        if(m==1) return k;
3
        11 res=(m-1)\%(n-k+1);
4
        for(11 i=n-k+2, stp=0; i\leq n; i+stp, res+stp*m) {
5
            if(res+m>=i) {
6
                res=(res+m)%i;
7
                i++;
8
                stp=0;
            }
9
            else {
10
11
                stp=(i-res-2)/(m-1);
12
                if(i+stp>n)
13
                {
14
                    res+=(n-(i-1))*m;
15
                    break;
16
                }
17
            }
18
        }
19
        return res+1;
20
    7.10 二分
 1 // a为二分数组, x为需要查找的数, 返回最左端和最右端
    pair<int, int> F(vector<int> a, int x) {
3
        int 1 = 0, r = a.size() - 1;
4
        int lres = -1;
5
        while (1 <= r) {</pre>
6
            int mid = 1 + r >> 1;
7
            int tt = a[mid];
8
            if (tt >= x) {
9
                r = mid - 1;
10
            } else if (tt < x) {</pre>
11
                1 = mid + 1;
12
13
        }
14
        if (l >= a.size() || a[l] != x) return make_pair(-1, -1);
15
        lres = 1;
16
        1 = 0, r = a.size() - 1;
17
        while (1 <= r) {</pre>
18
            int mid = 1 + r >> 1;
19
            int tt = a[mid];
20
            if (tt > x) {
21
                r = mid - 1;
22
            } else if (tt <= x) {</pre>
23
                l = mid + 1;
24
```

```
25
        }
26
        return make_pair(lres, r);
27 }
    7.11 猛男 IO 挂
1 const int LEN = 100000;
2
    struct fastio {
3
        int it, len;
4
        char s[LEN + 5];
5
        fastio() {
6
            it = len = 0;
7
        }
8
        char get() {
9
            if (it < len) return s[it++];</pre>
10
            it = 0, len = fread(s, 1, LEN, stdin);
11
            return len ? s[it++] : EOF;
12
        }
13
        bool notend() {
14
            char c;
            for (c = get(); c == ' ' || c == '\n'; c = get());
15
16
            if (it) it--;
17
            return c != EOF;
18
        }
19
        void put(char c) {
20
            if (it == LEN) fwrite(s, 1, LEN, stdout), it = 0;
21
            s[it++] = c;
22
        }
23
        void flush() {
24
            fwrite(s, 1, it, stdout);
25
26 } buff, bufo;
27 inline int getint() {
28
        char c;
29
        int res = 0, sig = 1;
        for (c = buff.get(); c < '0' || c > '9'; c = buff.get()) if (c == '-') sig = -1;
30
31
        for (; c >= '0' && c <= '9'; c = buff.get()) res = res * 10 + (c - '0');
32
        return sig * res;
33 }
34 inline 11 get11() {
35
        char c;
36
        11 \text{ res} = 0, \text{ sig} = 1;
37
        for (c = buff.get(); c < '0' || c > '9'; c = buff.get()) if (c == '-') sig = -1;
38
        for (; c >= '0' && c <= '9'; c = buff.get()) res = res * 10 + (c - '0');</pre>
39
        return sig * res;
40 }
41
   inline void putint(int x, char suf) {
42
        if (!x) bufo.put('0');
43
        else {
44
            if (x < 0) bufo.put('-'), x = -x;
45
            int k = 0;
46
            char s[15];
47
            while (x) {
48
                s[++k] = x \% 10 + '0';
49
                x /= 10;
50
51
            for (; k; k--) bufo.put(s[k]);
52
        }
```

```
bufo.put(suf);
53
54
   }
    inline void putll(ll x, char suf) {
55
56
        if (!x) bufo.put('0');
57
        else {
            if (x < 0) bufo.put('-'), x = -x;
58
59
            int k = 0;
60
            char s[25];
61
            while (x) {
62
                s[++k] = x % 10 + '0';
63
                x /= 10;
64
            }
65
            for (; k; k--) bufo.put(s[k]);
66
        }
67
        bufo.put(suf);
   }
68
69
    inline char get_char() {
        char c;
70
        for (c = buff.get(); c == ' ' || c == '\n'; c = buff.get());
71
72
        return c;
73 }
```

7.12 贪心结论

1 // n个区间, 挪到使得某个点被所有区间覆盖需要的最少步数时, 选择的点是所有区间端点的中位数 (mid~mid+1答案 都是一样的)

7.13 builtin

```
1 __builtin_popcount(unsigned int n) // 1的个数
2 __builtin_parity(unsigned int n) // 奇数个1返回1, 偶数个返回0
3 __builtin_ctz(unsigned int n) // 判断n的二进制末尾后面0的个数
4 __builtin_clz(unsigned int n) //返回前导0的个数
```

7.14 n 以内 k 因子的个数