



模板之用了就秃

WUST

So Like Coding? You Baldy

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0 Header

0.1 pbds

0.1.1 head

```
1 #include <bits/extc++.h>
2 #pragma comment(linker, "/STACK:102400000,102400000")
3 using namespace __gnu_pbds; // tree, gp_hash_table, trie
4 using namespace __gnu_cxx; // rope
5 tree<TYPE, null_type, less<>, rb_tree_tag, tree_order_statistics_node_update> tr;
6 // 可并堆
7 #include <ext/pb_ds/priority_queue.hpp>
8 using namespace __gnu_pbds;
9 __gnu_pbds::priority_queue<int, greater<int>, pairing_heap_tag> q[maxn];
10 //q[i].join(q[j]) 将j堆并入i
```

0.2 FastIO

0.2.1 快速 IO

```
1 // 适用于正负整数
2 template <class T>
3 inline bool scan(T &ret){
4     char c;
5     int sgn;
6     if (c = getchar(), c == EOF) return 0; //EOF
7     while (c != '-' && (c < '0' || c > '9')) c = getchar();
8     sgn = (c == '-') ? -1 : 1;
9     ret = (c == '-') ? 0 : (c - '0');
10    while (c = getchar(), c >= '0' && c <= '9') ret = ret * 10 + (c - '0');
11    ret *= sgn;
12    return 1;
13 }
14
15 template <class T>
16 inline void write(int x, int digit=10) {
17     static int sta[35];
18     int top = 0;
19     do {
20         sta[top++] = x % digit, x /= digit;
21     } while (x);
22     while (top) putchar(sta[--top] + 48);
23 }
24
25 inline int read() {
26     int x = 0;
27     char ch = getchar();
28     while (ch > '9' || ch < '0') ch = getchar();
29     while (ch >= '0' && ch <= '9') {
30         x = x * 10 + ch - '0';
31         ch = getchar();
32     }
33     return x;
34 }
```

0.2.2 快速幂

```
1 // 精确快速乘
2 ll mul(ll a, ll b, ll mod) {
3     a %= mod; b %= mod;
4     ll res = 0;
5     while (b > 0) {
6         if (b & 1) {
7             res = (res + a);
8             if (res >= mod) res -= mod;
9         }
10        a = (a + a);
11        if (a >= mod) a -= mod;
12        b >>= 1;
13    }
14    return res;
15 }
16
17 // O(1)快速乘
18 ll mul(ll a, ll b, ll M) { return (a * b - (ll)((long double)a * b / M) * M + M) % M; }
19
20 //int128
21 ll ans = ((__int128) a * b) % p;
22
23 // 10进制快速幂, 直接读入%s,c 预处理字符串len
24 char c[1000005];
25 int len;
26 ll qp(ll a) {
27     len--;
28     a %= mod;
29     ll s = a;
30     ll res = 1;
31     while (len >= 0) {
32         ll cur = s;
33         for (int i = 1; i <= c[len] - '0'; ++i) {
34             res = res * s % mod;
35         }
36         for (int i = 1; i < 10; ++i) {
37             cur = cur * s % mod;
38         }
39         s = cur;
40         len--;
41     }
42     return res;
43 }
```

0.2.3 Python 读入

```
1 // python一行读入
2 a,b = map(int, input().split())
3
4 a = []
5 for i in input().split():
6     a.append(int(i))
7
8
9 f.write('{0} {1}\n'.format(1, 2))
```

0.2.4 特殊 IO

```
1 // 代替gets
2 scanf("%[^\n]%*c", ss)
3
4 // 高精分数输出
5 void print(ll x, ll y, int need) { // need 表示小数位数
6     int a[need];
7     printf("%lld.", x / y);
8     ll t = x % y;
9     for (int i = 0; i < need; i++) {
10         t *= 10;
11         a[i] = t / y;
12         t = t % y;
13     }
14     if (t * 10 / y >= 5) a[need - 1]++;
15     int i = need - 1;
16     while (a[i] == 10) {
17         a[i] = 0;
18         a[--i]++;
19     }
20     for (int i = 0; i < need; i++) putchar(a[i]);
21     putchar('\n');
22 }
23
24 // type mx
25 // int 10
26 // int64 19
27 // int128 38
28 template<class T, int g = 10>
29 void print(T x) {
30     vector<char> a(38);
31     if (x < 0) {x = -x; putchar('-');}
32     if (x == 0) {putchar('0'); return;}
33     int tot = 0;
34     while (x > 0) {
35         a[tot++] = x % g;
36         x /= 10;
37     }
38     for (int i = tot - 1; i >= 0; --i) putchar('0' + a[i]);
39 }
```

0.3 header

```
1 // Editor -> Live Templates
2 // add template group acm
3 // add template main
4 // C++ Declaration
5
6 #include <bits/stdc++.h>
7 #define ll long long
8 #define x first
9 #define y second
10 #define sz size()
11 #define all(x) x.begin(), x.end()
12 using namespace std;
13
14 typedef pair<int, int> pii;
15 typedef vector<int> vi;
16 typedef vector<long long> vl;
```



```
17
18 template <class T>
19 inline bool scan(T &ret){
20     char c;
21     int sgn;
22     if (c = getchar(), c == EOF) return 0; //EOF
23     while (c != '-' && (c < '0' || c > '9')) c = getchar();
24     sgn = (c == '-') ? -1 : 1;
25     ret = (c == '-') ? 0 : (c - '0');
26     while (c = getchar(), c >= '0' && c <= '9') ret = ret * 10 + (c - '0');
27     ret *= sgn;
28     return 1;
29 }
30
31 const ll mod = 1e9+7;
32 const int maxn = $MAXN$;
33 const int inf = 0x3f3f3f3f;
34
35 ll qp(ll x, ll n, ll mod = ::mod) {
36     ll res = 1; x %= mod;
37     while (n > 0) {
38         if (n & 1) res = res * x % mod;
39         x = x * x % mod;
40         n >>= 1;
41     }
42     return res;
43 }
44
45 int main(int argc, char* argv[]) {
46     $END$
47     return 0;
48 }
49
50 // C++ Expression debug
51 freopen("data.in", "r", stdin);
52 freopen("data.out", "w", stdout);
53 clock_t ST = clock();
54 cerr << "time: " << ((clock()-ST)*1000.0 / CLOCKS_PER_SEC) << "ms" << endl;
55
56
57 // C++ Expression tkase
58 int T;
59 scanf("%d", &T);
60 for (int kase = 1; kase <= T; ++kase) {
61     $END$
62 }
```

0.3.1 comp

```
1 // 1 create directory comp
2 // 2 create directory comp/test
3 // 3 create duipai.cpp
4 // add text
5
6 #include<bits/stdc++.h>
7 using namespace std;
8 int main(){
9     int i;
```

```
10     for (i=1;;i++){
11         printf("The result of No. %d Case is: ",i);
12         system("python3 rand.py");
13         system("./std < test/data.in > test/std.out");
14         system("./my < test/data.in > test/my.out");
15         if (system("diff test/std.out test/my.out")){
16             printf("Wrong Answer\n");
17             return 0;
18         }
19         else printf("Accepted\n");
20     }
21     return 0;
22 }
23
24
25 // 4 create duipai.sh
26 #!/bin/bash
27 g++ std.cpp -o std
28 g++ my.cpp -o my
29 python3 rand.py
30 ./a.out
31
32 // 5 create rand.py
33 # coding=utf-8
34 from random import randint, choice, shuffle
35 # with open("../cmake-build-debug/data.in", "w") as f:
36 with open("test/data.in", "w") as f:
37     n = randint(1, 10)
38     m = randint(1, 10)
39     f.write(f"{n} {m}")
40
41 // 6 terminal: g++ duipai.cpp
42 // 7 terminal: sudo chmod 777 duipai.sh
43 // 8 add my.cpp and std.cpp
44 // 9 ./duipai.sh
```

0.3.2 vim

```
1 syntax on
2 set nu
3 set tabstop=4
4 set shiftwidth=4
5 set background=dark
6
7 map <C-A> ggVG"+y
8 map <F5> :call Run()<CR>
9 func! Run()
10     exec "w"
11     exec "!g++ -Wall % -o %<"
12     exec "!./%<"
13 endfunc
```

1 Math

1.1 群论

1.1.1 普通筛

```
1 int prim[maxn], tot;
2 bool v[maxn];
3 void init() {
4     for (int i = 2; i < maxn; ++i) {
5         if (!v[i]) prim[++tot] = i;
6         for (int j = i + i; j < maxn && j <= (ll)i * i; j += i) v[j] = 1;
7     }
8 }
```

1.1.2 线性筛

```
1 int mn[maxn], prim[maxprime];
2 int tot = 0;
3 void init() {
4     for (int i = 2; i < maxn; ++i) {
5         if (!mn[i]) prim[++tot] = i, mn[i] = i;
6         for (int j = 1; j <= tot && i * prim[j] < maxn; ++j) {
7             mn[i * prim[j]] = prim[j];
8             if (!(i % prim[j])) break;
9         }
10    }
11 }
12
13 // 素数个数表
14 // 10 4
15 // 1e5 9592 (1e4)
16 // 2e5 17984 (2e4)
17 // 5e5 41538 (5e4)
18 // 1e6 78498 (8e4)
19 // 2e6 148933 (1.5e5)
20 // 1e7 664579 (7e5)
21 // 3e7 1857859 (2e6)
22 // 1e8 5761455 (6e6)
```

1.1.3 Pollard Rho

```
1 typedef pair<ll, ll> PLL;
2 namespace Factor {
3     const int N = 1010000;
4     ll C, fac[10010], a[1001000];
5     int cnt, prime[N], p[N], psize, _cnt;
6     ll _e[100], _pr[100];
7     vector<ll> d;
8
9     // 快速乘
10    inline ll mul(ll a, ll b, ll p) {
11        if (p <= 1000000000) return a * b % p;
12        else if (p <= 1000000000000ll) return (((a * (b >> 20) % p) << 20) + (a * (b &
13            ((1 << 20) - 1)))) % p;
14        else {
15            ll d = (ll) floor(a * (long double) b / p + 0.5);
16            ll ret = (a * b - d * p) % p;
17        }
18    }
```

```

16         if (ret < 0) ret += p;
17         return ret;
18     }
19 }
20
21 // 素数筛
22 void prime_table() {
23     int i, j, tot, t1;
24     for (i = 1; i <= psize; i++) p[i] = i;
25     for (i = 2, tot = 0; i <= psize; i++) {
26         if (p[i] == i) prime[++tot] = i;
27         for (j = 1; j <= tot && (t1 = prime[j] * i) <= psize; j++) {
28             p[t1] = prime[j];
29             if (i % prime[j] == 0) break;
30         }
31     }
32 }
33
34 void init(int ps) {
35     psize = ps;
36     prime_table();
37 }
38
39 // 快速幂
40 ll powl(ll a, ll n, ll p) {
41     ll ans = 1;
42     for (; n; n >>= 1) {
43         if (n & 1) ans = mul(ans, a, p);
44         a = mul(a, a, p);
45     }
46     return ans;
47 }
48
49 bool witness(ll a, ll n) {
50     int t = 0;
51     ll u = n - 1;
52     for (; ~u & 1; u >>= 1) t++;
53     ll x = powl(a, u, n), _x = 0;
54     for (; t; t--) {
55         _x = mul(x, x, n);
56         if (_x == 1 && x != 1 && x != n - 1) return 1;
57         x = _x;
58     }
59     return _x != 1;
60 }
61
62 // n 是否为素数
63 bool miller(ll n) {
64     if (n < 2) return 0;
65     if (n <= psize) return p[n] == n;
66     if (~n & 1) return 0;
67     for (int j = 0; j <= 7; j++) if (witness(rand() % (n - 1) + 1, n)) return 0;
68     return 1;
69 }
70
71 ll gcd(ll a, ll b) {
72     ll ret = 1;
73     while (a != 0) {
74         if ((~a & 1) && (~b & 1)) ret <<= 1, a >>= 1, b >>= 1;

```

```

75         else if (~a & 1) a >>= 1;
76         else if (~b & 1) b >>= 1;
77         else {
78             if (a < b) swap(a, b);
79             a -= b;
80         }
81     }
82     return ret * b;
83 }
84
85 // 求 n 的一个素因子
86 ll rho(ll n) {
87     for (;;) {
88         ll X = rand() % n, Y, Z, T = 1, *lY = a, *lX = lY;
89         int tmp = 20;
90         C = rand() % 10 + 3;
91         X = mul(X, X, n) + C;
92         *(lY++) = X;
93         lX++;
94         Y = mul(X, X, n) + C;
95         *(lY++) = Y;
96         for (; X != Y;) {
97             ll t = X - Y + n;
98             Z = mul(T, t, n);
99             if (Z == 0) return gcd(T, n);
100            tmp--;
101            if (tmp == 0) {
102                tmp = 20;
103                Z = gcd(Z, n);
104                if (Z != 1 && Z != n) return Z;
105            }
106            T = Z;
107            Y = *(lY++) = mul(Y, Y, n) + C;
108            Y = *(lY++) = mul(Y, Y, n) + C;
109            X = *(lX++);
110        }
111    }
112 }
113
114 void _factor(ll n) {
115     for (int i = 0; i < cnt; i++) {
116         if (n % fac[i] == 0) n /= fac[i], fac[cnt++] = fac[i];
117     }
118     if (n <= psize) {
119         for (; n != 1; n /= p[n]) fac[cnt++] = p[n];
120         return;
121     }
122     if (miller(n)) fac[cnt++] = n;
123     else {
124         ll x = rho(n);
125         _factor(x);
126         _factor(n / x);
127     }
128 }
129
130 void dfs(ll x, int dep) {
131     if (dep == _cnt) d.push_back(x);
132     else {
133         dfs(x, dep + 1);

```

```

134         for (int i = 1; i <= _e[dep]; i++) dfs(x * _pr[dep], dep + 1);
135     }
136 }
137
138 void norm() {
139     sort(fac, fac + cnt);
140     _cnt = 0;
141     for (int i = 0; i < cnt; ++i)
142         if (i == 0 || fac[i] != fac[i - 1]) _pr[_cnt] = fac[i], _e[_cnt++] = 1;
143         else _e[_cnt - 1]++;
144 }
145
146 vector<ll> getd() {
147     d.clear();
148     dfs(1, 0);
149     return d;
150 }
151
152 // 返回所有因子
153 vector<ll> factor(ll n) {
154     cnt = 0;
155     _factor(n);
156     norm();
157     return getd();
158 }
159
160 // 返回所有素因子
161 vector<PLL> factorG(ll n) {
162     cnt = 0;
163     _factor(n);
164     norm();
165     vector<PLL> d;
166     for (int i = 0; i < _cnt; ++i) d.push_back(make_pair(_pr[i], _e[i]));
167     return d;
168 }
169
170 // a 是否为 p 的原根
171 bool is_primitive(ll a, ll p) {
172     vector<PLL> D = factorG(p - 1);
173     for (int i = 0; i < (int) D.size(); ++i) if (powl(a, (p - 1) / D[i].first, p)
174         == 1) return 0;
175     return 1;
176 }
177
178 // a 关于 g 的阶
179 int findorder(ll a, ll p) {
180     vector<PLL> D = factorG(p - 1);
181     int t = p - 1;
182     for (int i = 0; i < (int) D.size(); ++i) {
183         while (t % D[i].first == 0 && powl(a, t / D[i].first, p) == 1) t /= D[i].
184             first;
185     }
186     return t;
187 }

```

1.1.4 解乘法逆元

```

1 void exgcd(ll a, ll b, ll c, ll d, ll &x, ll &y) {

```

```

2     ll z = (a + b - 1) / b;
3     if (z <= c / d) {
4         x = z;
5         y = 1;
6         return;
7     }
8     a -= (z - 1) * b; c -= (z - 1) * d;
9     exgcd(d, c, b, a, y, x);
10    x += (z - 1) * y;
11 }
12
13 // 求  $a/b \bmod p = x \bmod p$ , 优先 b 小
14 pair<ll, ll> invInv(ll p, ll x) {
15     ll b, y;
16     exgcd(p, x, p, x - 1, b, y);
17     return {b * x - p * y, b};
18 }

```

1.1.5 EulerPhi

```

1 // 计算欧拉phi函数, phi(n)且与n互素的正整数个数
2
3 // 单点欧拉  $O(\sqrt{n})$ 
4 ll euler(ll n) {
5     ll rea = n;
6     for (ll i = 2; i * i <= n; i++)
7         if (n % i == 0) {
8             rea = rea - rea / i;
9             while (n % i == 0) n /= i;
10        }
11    if (n > 1) rea = rea - rea / n;
12    return rea;
13 }
14
15
16 // maxn 为根号最值
17 // maxprime 为 maxn 内素数个数
18 // 素数线筛 + 单点求值
19 int prim[maxprime], cnt;
20 bool v[maxn];
21 void init() {
22     v[0] = v[1] = 1;
23     for (int i = 2; i < maxn; i++) {
24         if (!v[i]) prim[++cnt] = i;
25         for (int j = 1; j <= cnt && i * prim[j] < maxn; j++) {
26             v[i * prim[j]] = 1;
27             if (!(i % prim[j])) break;
28         }
29     }
30 }
31 ll phi(ll n) {
32     ll rea = n;
33     for (ll i = 1; prim[i] * prim[i] <= n; i++)
34         if (n % prim[i] == 0) {
35             rea = rea - rea / prim[i];
36             while (n % prim[i] == 0) n /= prim[i];
37         }
38     if (n > 1) rea = rea - rea / n;

```

```

39     return rea;
40 }
41
42
43 // 线筛素数+欧拉函数
44 int phi[maxn], prim[maxprime], cnt;
45 bool v[maxn];
46 void init() {
47     phi[1] = 1;
48     for (int i = 2; i < maxn; ++i) {
49         if (!v[i]) prim[++cnt] = i, phi[i] = i - 1;
50         for (int j = 1; j <= cnt; ++j) {
51             if (i * prim[j] >= maxn) break;
52             v[i * prim[j]] = 1;
53             if (i % prim[j] == 0) {
54                 phi[i * prim[j]] = phi[i] * prim[j]; break;
55             } else phi[i * prim[j]] = phi[i] * phi[prim[j]];
56         }
57     }
58 }

```

1.1.6 求素因子

```

1 vector<pair<ll, int>> getFactors(ll x) {
2     vector<pair<ll, int>> fact;
3     for (int i = 1; prim[i] <= x / prim[i]; i++) {
4         if (x % prim[i] == 0) {
5             fact.emplace_back(prim[i], 0);
6             while (x % prim[i] == 0) fact.back().second++, x /= prim[i];
7         }
8     }
9     if (x != 1) fact.emplace_back(x, 1);
10    return fact;
11 }

```

1.1.7 gcd

```

1 ll gcd(ll a, ll b) {while(b ^= a ^= b ^= a % b); return a;}
2
3 ll gcd(ll a, ll b) { return b == 0 ? a : gcd(b, a % b); }
4
5 ll exgcd(ll a, ll b, ll &x, ll &y) {
6     ll d = a;
7     if (b) d = exgcd(b, a % b, y, x), y -= x * (a / b);
8     else x = 1, y = 0;
9     return d;
10 }
11
12 // ax + by = c
13 // x = x + k*dx
14 // y = y - k*dx
15 // 当x和y都非负时返回1, x, y即为当前最小非负整数解 (优先x)
16 bool solve(ll a, ll b, ll c, ll &x, ll &y, ll &dx, ll &dy) {
17     x = y = dx = dy = 0;
18     if (a == 0 && b == 0) return 0;
19     ll x0, y0;
20     ll d = exgcd(a, b, x0, y0);
21     if (c % d != 0) return 0;

```



```
22     dx = b / d, dy = a / d;
23     x = (x0 % dx * ((c / d) % dx) % dx + dx) % dx;
24     y = (c - a * x) / b;
25     // 删掉这一句返回x的最小非负整数解
26     if (y < 0) return 0;
27     return 1;
28 }
```

1.1.8 区间筛

```
1  bool f[maxlen];
2  bool sieve[maxn];
3
4  // maxn 至少为 sqrt(R), 预处理
5  void init() {
6      for (int i = 2; i < maxn; i++) sieve[i] = true;
7      for (int i = 2; i * i < maxn; i++) {
8          if (sieve[i]) {
9              for (int j = i * 2; j < maxn; j += i) {
10                 sieve[j] = false;
11             }
12         }
13     }
14 }
15 // 计算 [L,R] 素性, f[i] 为 1 表示 i+L 为素数
16 void cal(ll L, ll R) {
17     int len = R - L + 1;
18     for (int i = 0; i < len; i++) f[i] = true;
19     if (1 - L >= 0) f[1 - L] = false;
20     for (ll i = 2; i * i < R; i++) {
21         if (sieve[i]) {
22             for (ll j = max(1ll * 2, (L - 1 + i) / i) * i; j <= R; j += i) f[j - L] =
                false;
23         }
24     }
25 }
```

1.1.9 欧拉降幂

```
1  const int maxn = 1e7+50;
2
3  int prim[maxn], vis[maxn];
4  int tot, phi[maxn];
5  struct node {
6      ll res;
7      bool v;
8  };
9
10 node qpow(ll A, ll B, ll C) {
11     ll re = 1;
12     bool flag = true;
13     while (B) {
14         if (B & 1) {
15             if ((re * A) >= C) flag = 0;
16             re = re % C;
17         }
18         B = B >> 1;
19         if (B) {
```

```
20         if (A >= C) flag = 0;
21         A %= C;
22         if ((A * A) >= C) flag = 0;
23         A %= C;
24     }
25 }
26 return node{re, flag};
27 }
28
29 void init(int n) {
30     phi[1] = 1;
31     for (int i = 2; i <= n; i++) {
32         if (!vis[i]) {
33             prim[++tot] = i;
34             phi[i] = i - 1;
35         }
36         for (int j = 1; j <= tot && prim[j] * i <= n; j++) {
37             vis[i * prim[j]] = 1;
38             if (i % prim[j] == 0) {
39                 phi[i * prim[j]] = phi[i] * prim[j];
40                 break;
41             } else phi[i * prim[j]] = phi[i] * (prim[j] - 1);
42         }
43     }
44 }
45 }
46
47 inline ll Euler(ll x) {
48     if (x < maxn) return phi[x];
49     return 0;
50 }
51 }
52
53 node f(ll a, ll k, ll p) {
54     if (p == 1) return node{0, 0};
55     if (k == 0) return node{a % p, a < p};
56     ll ep = Euler(p);
57     node tmp = f(a, k - 1, ep);
58     if (__gcd(a, p) == 1) return qpow(a, tmp.res, p);
59     if (!tmp.v) {
60         tmp.res += ep;
61     }
62     return qpow(a, tmp.res, p);
63 }
64
65 int main() {
66     ll a, k, p;
67     init(1e7+2);
68     int T;
69     scanf("%d", &T);
70     for (int kase = 1; kase <= T; ++kase) {
71         // k次a次方模p的值
72         scanf("%lld%lld%lld", &a, &k, &p);
73         if (k == 0) printf("%lld\n", 1 % p);
74         else printf("%lld\n", f(a, k - 1, p).res);
75     }
76     return 0;
77 }
```

1.1.10 dirichlet

```

1  #define poly vector<int>
2
3  // mod need to be prim
4  namespace dirichlet {
5      int inv[maxn], cnt[maxn];
6      bool vis[maxn];
7      void init(int n) {
8          cnt[1] = 1;
9          for (int i = 2; i <= n; ++i) {
10             if (!vis[i]) {
11                 cnt[i] = 1;
12                 for (int j = 2, k = i * j; k <= n; ++j, k += i) {
13                     vis[k] = true;
14                     if (!cnt[k] && cnt[j]) cnt[k] = cnt[j] + 1;
15                 }
16             }
17         }
18         inv[1] = 1;
19         // no need to n (1e6 need 128)
20         for (int i = 2; i <= n; ++i) inv[i] = 1ll * (mod - mod / i) * inv[mod % i] %
            mod;
21     }
22
23     // no need to init, nlogn
24     poly mul(poly a, poly b) {
25         int n = max(a.size(), b.size()) - 1;
26         while ((int)a.size() <= n) a.push_back(0);
27         while ((int)b.size() <= n) b.push_back(0);
28         poly d(n + 1);
29         for (int i = 1; i <= n; ++i) {
30             for (int j = 1, k = i; k <= n; k += i, ++j) d[k] = (d[k] + 1ll * a[i] * b[j]
                ) % mod;
31         }
32         return d;
33     }
34
35     // nlogn
36     poly qp(poly a, int k) {
37         k %= mod;
38         if (k < 0) k += mod;
39         int n = a.size() - 1;
40         poly g(n + 1);
41         poly tf = a;
42         for (int i = 1; i <= n; ++i) a[i] = 1ll * a[i] * cnt[i] % mod * k % mod;
43         for (int i = 1; i <= n; ++i) {
44             int v = g[i];
45             g[i] = 1ll * g[i] * inv[cnt[i]] % mod + (i == 1);
46             for (int j = 2, t = i + i; t <= n; ++j, t += i)
47                 g[t] = (g[t] + 1ll * g[i] * a[j]) % mod;
48             for (int j = 2, t = i + i; t <= n; ++j, t += i)
49                 g[t] = (g[t] + 1ll * (mod - tf[j]) * v) % mod;
50         }
51         return g;
52     }
53
54     // nlogn
55     inline poly ln(const poly& a) {

```

```

56     int n = a.size() - 1;
57     poly b(n + 1);
58     for (int i = 2; i <= n; ++i) b[i] = 1ll * a[i] * cnt[i] % mod;
59     for (int i = 2; i <= n; ++i) {
60         for (int j = 2, k = i * j; k <= n; ++j, k += i) b[k] = (b[k] - 1ll * b[i] *
61             a[j]) % mod;
62         b[i] = 1ll * b[i] * inv[cnt[i]] % mod;
63         if (b[i] < 0) b[i] += mod;
64     }
65     return b;
66 }
67 // nlogn
68 inline poly exp(poly a) {
69     int n = a.size() - 1;
70     poly b(n + 1);
71     for (int i = 2; i <= n; ++i) a[i] = 1ll * a[i] * cnt[i] % mod;
72     b[1] = 1;
73     for (int i = 1; i <= n; ++i) {
74         b[i] = 1ll * b[i] * inv[cnt[i]] % mod;
75         for (int j = 2, k = i * j; k <= n; ++j, k += i) b[k] = (b[k] + 1ll * b[i] *
76             a[j]) % mod;
77     }
78     return b;
79 }
80 using namespace dirichlet;

```

1.2 同余

1.2.1 逆元

```

1  /*
2  1. 费马小定理
3  条件: mod为素数
4  */
5  ll inv(ll x) { return qp(x, mod - 2); }
6
7  /*
8  2. 扩展欧几里得
9  条件: gcd(a, mod) == 1
10  如果 gcd(a, mod) != 1 返回 -1
11  */
12  ll inv(ll a, ll p) {
13      ll g, x, y;
14      g = exgcd(a, p, x, y);
15      return g == 1 ? (x + p) % p : -1;
16  }
17
18  /*
19  3. 公式
20  a/b % mod = c
21  -> a % (b * mod) / b = c
22  */
23
24  /*
25  4. 逆元打表
26  p 是模
27  p 要求是奇素数

```

```
28 */
29 ll inv[maxn];
30
31 void getinv(int n, ll p) {
32     inv[1] = 1;
33     for (int i = 2; i <= n; i++) inv[i] = (p - p / i) * inv[p % i] % p;
34 }
35
36 // log逆元
37 ll dlog(ll g, ll b, ll p) {
38     ll m = sqrt(p - 1);
39     map<ll, ll> powers;
40     for (long j = 0; j < m; j++) powers[qp(g, j, p)] = j;
41     long gm = qp(g, -m + 2 * (p - 1), p);
42     for (int i = 0; i < m; i++) {
43         if (powers[b]) return i * m + powers[b];
44         b = b * gm % p;
45     }
46     return -1;
47 }
```

1.2.2 中国剩余定理

```
1 void exgcd(ll a, ll b, ll &x, ll &y) {
2     if (b) exgcd(b, a % b, y, x), y -= x * (a / b);
3     else x = 1, y = 0;
4 }
5
6 // x % m[i] = a[i]
7 ll crt(const vector<ll>& m, const vector<ll>& a) {
8     ll M = 1, ans = 0;
9     int n = a.size();
10    for (int i = 0; i < n; i++) M *= m[i];
11    for (int i = 0; i < n; i++) {
12        ll x = 0, y = 0, Mi = M / m[i];
13        exgcd(Mi, m[i], x, y);
14        ans = (ans + Mi % M * x % M * a[i] % M + M) % M;
15    }
16    if (ans < 0) ans += M;
17    return ans;
18 }
```

1.2.3 扩展中国剩余定理

```
1 ll exgcd(ll a, ll b, ll &x, ll &y) {
2     ll d = a;
3     if (b) d = exgcd(b, a % b, y, x), y -= x * (a / b);
4     else x = 1, y = 0;
5     return d;
6 }
7
8 // x % m[i] = a[i]
9 ll excrt(const vector<ll> &m, const vector<ll> &a) {
10    ll M = m[0], R = a[0];
11    int n = a.size();
12    for (int i = 1; i < n; i++) {
13        ll x = 0, y = 0, d = exgcd(M, m[i], x, y);
14        if ((R - a[i]) % d) return -1;
```

```
15 //      x = mul((R - a[i]) / d, x, m[i]);
16      x = (R - a[i]) / d * x % m[i];
17      R -= M * x;
18      M = M / d * m[i];
19      R %= M;
20  }
21  if (R < 0) R += M;
22  return R;
23 }
```

1.2.4 BSGS

```
1 // 求  $a^x = b \pmod p$ 
2 int BSGS(int a, int b, int p) {
3     map<int, int> hash;
4     b %= p;
5     int t = (int) sqrt(p) + 1;
6     for (int j = 0; j < t; j++) {
7         int val = 1ll * b * qp(a, j, p) % p;
8         hash[val] = j;
9     }
10    a = qp(a, t, p);
11    if (a == 0) return b == 0 ? 1 : -1;
12    for (int i = 0; i <= t; i++) {
13        int val = qp(a, i, p);
14        int j = hash.find(val) == hash.end() ? -1 : hash[val];
15        if (j >= 0 && i * t - j >= 0) return i * t - j;
16    }
17    return -1;
18 }
```

1.2.5 exBSGS

```
1 unordered_map<int, int> Hash;
2
3 int exBSGS(int a, int b, int p) {
4     a %= p, b %= p;
5     if (b == 1) return 0;
6     if (!b && !a) return 1;
7     if (!a) return -1;
8     if (!b) {
9         int ret = 0, d;
10        while ((d = __gcd(a, p)) != 1) {
11            ++ret, p /= d;
12            if (p == 1) return ret;
13        }
14        return -1;
15    }
16    int ret = 0, A = a, B = b, P = p, C = 1, d;
17    while ((d = __gcd(A, P)) != 1) {
18        if (B % d) return -1;
19        P /= d, B /= d;
20        C = 1ll * C * (A / d) % P;
21        ++ret;
22        if (C == B) return ret;
23    }
24    Hash.clear();
25    int f = 1, t = sqrt(P) + 1;
```

```

26     for (int i = 0; i < t; i++) {
27         Hash[1ll * f * B % P] = i;
28         f = 1ll * f * A % P;
29     }
30     int tf = f;
31     f = 1ll * f * C % P;
32     for (int i = 1; i <= t; i++) {
33         if (Hash.find(f) != Hash.end()) return ret + i * t - Hash[f];
34         f = 1ll * f * tf % P;
35     }
36     return -1;
37 }

```

1.2.6 模素数二次同余方程

```

1 // 要求模为素数, 输入n, mod, 返回  $x^2 \bmod n$ , 可解任意一元二次方程
2
3 bool Legendre(ll a, ll p) {
4     return qp(a, p - 1 >> 1, p) == 1;
5 }
6
7 ll modsqr(ll a, ll p) {
8     ll x;
9     ll i, k, b;
10    if (p == 2) x = a % p;
11    else if (p % 4 == 3) x = qp(a, p + 1 >> 2, p);
12    else {
13        for (b = 1; Legendre(b, p); ++b);
14        i = p - 1 >> 1;
15        k = 0;
16        do {
17            i >>= 1;
18            k >>= 1;
19            if (!((1LL * qp(a, i, p) * qp(b, k, p) + 1) % p)) k += p - 1 >> 1;
20        } while (!(i & 1));
21        x = 1ll * qp(a, i + 1 >> 1, p) * qp(b, k >> 1, p) % p;
22    }
23    return min(x, p - x);
24    // if(p-x<x) x=p-x;
25    // if(x==p-x) printf("%d\n",x);
26    // else printf("%d %d\n",x,p-x);
27 }

```

1.3 线代

1.3.1 线性基

```

1 struct Base {
2     #define TYPE ll
3     static const int len = 64;
4     bool rel; int sz;
5     TYPE a[len];
6
7     void init() {
8         rel = sz = 0;
9         memset(a, 0, sizeof a);
10    }
11 }

```

```
12     TYPE &operator[](int x) {
13         return a[x];
14     }
15
16     TYPE operator[](int x) const {
17         return a[x];
18     }
19
20
21     void ins(ll x) {
22         for(int i = 63; i >= 0; i --) {
23             if((x >> i) & 1) {
24                 if(!d[i]) return void(d[i] = x);
25                 x ^= d[i];
26             }
27         }
28     }
29
30     void insert(TYPE t) {
31         for (int i = len - 1; i >= 0; --i) {
32             if (!(t >> i & 1)) continue;
33             if (a[i]) t ^= a[i];
34             else {
35                 for (int j = 0; j < i; ++j) if (t >> j & 1) t ^= a[j];
36                 for (int j = i+1; j < len; ++j) if (a[j] >> i & 1) a[j] ^= t;
37                 a[i] = t;
38                 ++sz;
39                 return;
40             }
41         }
42         rel = true;
43     }
44
45     bool check(TYPE x) {
46         for (int i = len - 1; i >= 0; i--)
47             if ((x >> i) & 1) {
48                 if (a[i]) x ^= a[i];
49                 else return false;
50             }
51         return true;
52     }
53
54     TYPE mx() {
55         TYPE res = 0;
56         for (int i = len - 1; i >= 0; --i) {
57             if ((res ^ (a[i])) > res) res ^= a[i];
58         }
59         return res;
60
61     // vector<TYPE> v;
62     // void basis() {for (int i = 0; i < len; ++i) if (a[i]) v.push_back(a[i]);}
63     // TYPE k_th(TYPE k) {
64     //     k -= rel;
65     //     if(k >= (((TYPE)1) << sz)) return -1;
66     //     TYPE ans = 0;
67     //     for(int i = 0; i < (int)v.size(); i ++) if(k & (((TYPE)1) << i)) ans ^= v[i];
68     //     return ans;
69     // }
70     // void init()
71     // {
```



```

71 //      sz = 0;
72 //      for(int i = 0; i < len; i ++) if(a[i])
73 //          for(int j = 0; j < i; j ++)
74 //              if(a[i] & (1ll << j)) a[i] ^= a[j];
75 //      for(int i = 0; i < len; i ++) if(a[i]) a[sz ++] = a[i];
76 //  }
77 friend Base intersection(const Base &a, const Base &b) {
78     Base ans = {}, c = b, d = b;
79     for (int i = 0; i < len; i++) {
80         TYPE x = a[i];
81         if (!x)continue;
82         int j = i;
83         TYPE T = 0;
84         for (; j >= 0; --j) {
85             if ((x >> j) & 1)
86                 if (c[j]) {
87                     x ^= c[j];
88                     T ^= d[j];
89                 }
90             else break;
91         }
92         if (!x)ans[i] = T;
93         else {
94             c[j] = x;
95             d[j] = T;
96         }
97     }
98     return ans;
99 }
100
101 #undef TYPE
102 };
103
104 // 前缀线性基
105 struct LinearBasis {
106     int f[20], g[20];
107
108     void ins(int x, int idx) {
109         for (int i = 19; ~i; i--) {
110             if ((x >> i) & 1) {
111                 if (f[i]) {
112                     if (g[i] <= idx) {
113                         x ^= f[i];
114                         f[i] ^= x;
115                         swap(g[i], idx);
116                     }
117                     else x ^= f[i];
118                 } else {
119                     f[i] = x;
120                     g[i] = idx;
121                     break;
122                 }
123             }
124         }
125     }
126
127     int query(int l) {
128         int res = 0;
129         for (int i = 19; ~i; i--)

```

```

130         if (g[i] >= 1)
131             res = max(res, res ^ f[i]);
132     return res;
133 }
134 } base[maxn];

```

1.3.2 高斯消元

```

1  int n;
2  double b[maxn], c[maxn][maxn];
3  //c: 系数矩阵, b: 常数
4
5  void gauss() {
6      for (int i = 1; i <= n; i++) {
7          for (int j = i; j <= n; j++)
8              if (fabs(c[j][i]) > eps) {
9                  for (int k = 1; k <= n; k++) swap(c[i][k], c[j][k]);
10                 swap(b[i], b[j]);
11             }
12         for (int j = 1; j <= n; j++) {
13             if (i == j) continue;
14             double rate = c[j][i] / c[i][i];
15             for (int k = i; k <= n; k++) c[j][k] -= c[i][k] * rate;
16             b[j] -= b[i] * rate;
17         }
18     }
19 }

```

1.3.3 BM

```

1  //Berlekamp-Massey
2  typedef vector<int> VI;
3  namespace linear_seq
4  {
5      #define rep(i,a,n) for (int i=a;i<n;i++)
6      #define SZ(x) ((int)(x).size())
7      #define pb(x) push_back(x)
8      const ll mod=1e9+7;
9      ll powmod(ll a,ll b){ll res=1;a%=mod; assert(b>=0); for(;b>=>1){if(b&1)res=res*a%mod;a=a*a%mod;}return res;}
10     const int N=10010;
11     ll res[N],base[N],_c[N],_md[N];
12     vector<int> Md;
13     void mul(ll *a,ll *b,int k)
14     {
15         rep(i,0,k+k) _c[i]=0;
16         rep(i,0,k) if (a[i]) rep(j,0,k) _c[i+j]=(_c[i+j]+a[i]*b[j])%mod;
17         for (int i=k+k-1;i>=k;i--) if (_c[i])
18             rep(j,0,SZ(Md)) _c[i-k+Md[j]]=( _c[i-k+Md[j]]-_c[i]*_md[Md[j]])%mod;
19         rep(i,0,k) a[i]=_c[i];
20     }
21     int solve(ll n,VI a,VI b){
22         ll ans=0,pnt=0;
23         int k=SZ(a);
24         assert(SZ(a)==SZ(b));
25         rep(i,0,k) _md[k-1-i]=-a[i];_md[k]=1;
26         Md.clear();
27         rep(i,0,k) if (_md[i]!=0) Md.push_back(i);

```

```

28     rep(i,0,k) res[i]=base[i]=0;
29     res[0]=1;
30     while ((1ll<<pnt)<=n) pnt++;
31     for (int p=pnt;p>=0;p--) {
32         mul(res,res,k);
33         if ((n>>p)&1) {
34             for (int i=k-1;i>=0;i--) res[i+1]=res[i];res[0]=0;
35             rep(j,0,SZ(Md)) res[Md[j]]=(res[Md[j]]-res[k]*_md[Md[j]])%mod;
36         }
37     }
38     rep(i,0,k) ans=(ans+res[i]*b[i])%mod;
39     if (ans<0) ans+=mod;
40     return ans;
41 }
42 VI BM(VI s){
43     VI C(1,1),B(1,1);
44     int L=0,m=1,b=1;
45     rep(n,0,SZ(s)){
46         ll d=0;
47         rep(i,0,L+1) d=(d+(1ll)C[i]*s[n-i])%mod;
48         if(d==0) ++m;
49         else if(2*L<=n){
50             VI T=C;
51             ll c=mod-d*powmod(b,mod-2)%mod;
52             while (SZ(C)<SZ(B)+m) C.pb(0);
53             rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
54             L=n+1-L; B=T; b=d; m=1;
55         } else {
56             ll c=mod-d*powmod(b,mod-2)%mod;
57             while (SZ(C)<SZ(B)+m) C.pb(0);
58             rep(i,0,SZ(B)) C[i+m]=(C[i+m]+c*B[i])%mod;
59             ++m;
60         }
61     }
62     return C;
63 }
64 int gao(VI a,ll n)
65 {
66     VI c=BM(a);
67     c.erase(c.begin());
68     rep(i,0,SZ(c)) c[i]=(mod-c[i])%mod;
69     return solve(n,c,VI(a.begin(),a.begin()+SZ(c)));
70 }
71 };//linear_seq::gao(VI{},n-1)

```

1.3.4 exBM

```

1 // given first m items init[0..m-1] and coefficents trans[0..m-1] or
2 // given first 2 *m items init[0..2m-1], it will compute trans[0..m-1]
3 // for you. trans[0..m] should be given as that
4 //     init[m] = sum_{i=0}^{m-1} init[i] * trans[i]
5 struct LinearRecurrence
6 {
7     using int64 = long long;
8     using vec = std::vector<int64>;
9
10     static void extand(vec& a, size_t d, int64 value = 0)
11     {

```

```

12     if (d <= a.size()) return;
13     a.resize(d, value);
14 }
15 static vec BerlekampMassey(const vec& s, int64 mod)
16 {
17     std::function<int64(int64)> inverse = [&](int64 a) {
18         return a == 1 ? 1 : (int64)(mod - mod / a) * inverse(mod % a) % mod;
19     };
20     vec A = {1}, B = {1};
21     int64 b = s[0];
22     for (size_t i = 1, m = 1; i < s.size(); ++i, m++)
23     {
24         int64 d = 0;
25         for (size_t j = 0; j < A.size(); ++j)
26         {
27             d += A[j] * s[i - j] % mod;
28         }
29         if (!(d %= mod)) continue;
30         if (2 * (A.size() - 1) <= i)
31         {
32             auto temp = A;
33             extend(A, B.size() + m);
34             int64 coef = d * inverse(b) % mod;
35             for (size_t j = 0; j < B.size(); ++j)
36             {
37                 A[j + m] -= coef * B[j] % mod;
38                 if (A[j + m] < 0) A[j + m] += mod;
39             }
40             B = temp, b = d, m = 0;
41         }
42         else
43         {
44             extend(A, B.size() + m);
45             int64 coef = d * inverse(b) % mod;
46             for (size_t j = 0; j < B.size(); ++j)
47             {
48                 A[j + m] -= coef * B[j] % mod;
49                 if (A[j + m] < 0) A[j + m] += mod;
50             }
51         }
52     }
53     return A;
54 }
55 static void exgcd(int64 a, int64 b, int64& g, int64& x, int64& y)
56 {
57     if (!b)
58         x = 1, y = 0, g = a;
59     else
60     {
61         exgcd(b, a % b, g, y, x);
62         y -= x * (a / b);
63     }
64 }
65 static int64 crt(const vec& c, const vec& m)
66 {
67     int n = c.size();
68     int64 M = 1, ans = 0;
69     for (int i = 0; i < n; ++i) M *= m[i];
70     for (int i = 0; i < n; ++i)

```

```

71     {
72         int64 x, y, g, tm = M / m[i];
73         exgcd(tm, m[i], g, x, y);
74         ans = (ans + tm * x * c[i] % M) % M;
75     }
76     return (ans + M) % M;
77 }
78 static vec ReedsSloane(const vec& s, int64 mod)
79 {
80     auto inverse = [](int64 a, int64 m) {
81         int64 d, x, y;
82         exgcd(a, m, d, x, y);
83         return d == 1 ? (x % m + m) % m : -1;
84     };
85     auto L = [](const vec& a, const vec& b) {
86         int da = (a.size() > 1 || (a.size() == 1 && a[0])) ? a.size() - 1 : -1000;
87         int db = (b.size() > 1 || (b.size() == 1 && b[0])) ? b.size() - 1 : -1000;
88         return std::max(da, db + 1);
89     };
90     auto prime_power = [&](const vec& s, int64 mod, int64 p, int64 e) {
91         // linear feedback shift register mod p^e, p is prime
92         std::vector<vec> a(e), b(e), an(e), bn(e), ao(e), bo(e);
93         vec t(e), u(e), r(e), to(e, 1), uo(e), pw(e + 1);
94         ;
95         pw[0] = 1;
96         for (int i = pw[0] = 1; i <= e; ++i) pw[i] = pw[i - 1] * p;
97         for (int64 i = 0; i < e; ++i)
98         {
99             a[i] = {pw[i]}, an[i] = {pw[i]};
100            b[i] = {0}, bn[i] = {s[0] * pw[i] % mod};
101            t[i] = s[0] * pw[i] % mod;
102            if (t[i] == 0)
103            {
104                t[i] = 1, u[i] = e;
105            }
106            else
107            {
108                for (u[i] = 0; t[i] % p == 0; t[i] /= p, ++u[i])
109                ;
110            }
111        }
112        for (size_t k = 1; k < s.size(); ++k)
113        {
114            for (int g = 0; g < e; ++g)
115            {
116                if (L(an[g], bn[g]) > L(a[g], b[g]))
117                {
118                    ao[g] = a[e - 1 - u[g]];
119                    bo[g] = b[e - 1 - u[g]];
120                    to[g] = t[e - 1 - u[g]];
121                    uo[g] = u[e - 1 - u[g]];
122                    r[g] = k - 1;
123                }
124            }
125            a = an, b = bn;
126            for (int o = 0; o < e; ++o)
127            {
128                int64 d = 0;
129                for (size_t i = 0; i < a[o].size() && i <= k; ++i)

```

```

130         {
131             d = (d + a[o][i] * s[k - i]) % mod;
132         }
133         if (d == 0)
134         {
135             t[o] = 1, u[o] = e;
136         }
137         else
138         {
139             for (u[o] = 0, t[o] = d; t[o] % p == 0; t[o] /= p, ++u[o])
140                 ;
141             int g = e - 1 - u[o];
142             if (L(a[g], b[g]) == 0)
143             {
144                 extend(bn[o], k + 1);
145                 bn[o][k] = (bn[o][k] + d) % mod;
146             }
147             else
148             {
149                 int64 coef = t[o] * inverse(to[g], mod) % mod * pw[u[o] -
150                     uo[g]] % mod;
151                 int m = k - r[g];
152                 extend(an[o], ao[g].size() + m);
153                 extend(bn[o], bo[g].size() + m);
154                 for (size_t i = 0; i < ao[g].size(); ++i)
155                 {
156                     an[o][i + m] -= coef * ao[g][i] % mod;
157                     if (an[o][i + m] < 0) an[o][i + m] += mod;
158                 }
159                 while (an[o].size() && an[o].back() == 0) an[o].pop_back();
160                 for (size_t i = 0; i < bo[g].size(); ++i)
161                 {
162                     bn[o][i + m] -= coef * bo[g][i] % mod;
163                     if (bn[o][i + m] < 0) bn[o][i + m] += mod;
164                 }
165                 while (bn[o].size() && bn[o].back() == 0) bn[o].pop_back();
166             }
167         }
168     }
169     return std::make_pair(an[0], bn[0]);
170 };
171
172 std::vector<std::tuple<int64, int64, int>> fac;
173 for (int64 i = 2; i * i <= mod; ++i)
174 {
175     if (mod % i == 0)
176     {
177         int64 cnt = 0, pw = 1;
178         while (mod % i == 0) mod /= i, ++cnt, pw *= i;
179         fac.emplace_back(pw, i, cnt);
180     }
181 }
182 if (mod > 1) fac.emplace_back(mod, mod, 1);
183 std::vector<vec> as;
184 size_t n = 0;
185 for (auto&& x : fac)
186 {
187     int64 mod, p, e;

```

```

188         vec a, b;
189         std::tie(mod, p, e) = x;
190         auto ss = s;
191         for (auto&& x : ss) x %= mod;
192         std::tie(a, b) = prime_power(ss, mod, p, e);
193         as.emplace_back(a);
194         n = std::max(n, a.size());
195     }
196     vec a(n), c(as.size()), m(as.size());
197     for (size_t i = 0; i < n; ++i)
198     {
199         for (size_t j = 0; j < as.size(); ++j)
200         {
201             m[j] = std::get<0>(fac[j]);
202             c[j] = i < as[j].size() ? as[j][i] : 0;
203         }
204         a[i] = crt(c, m);
205     }
206     return a;
207 }
208
209 LinearRecurrence(const vec& s, const vec& c, int64 mod) : init(s), trans(c), mod(
    mod), m(s.size()) {}
210 LinearRecurrence(const vec& s, int64 mod, bool is_prime = true) : mod(mod)
211 {
212     vec A;
213     if (is_prime)
214         A = BerlekampMassey(s, mod);
215     else
216         A = ReedsSloane(s, mod);
217     if (A.empty()) A = {0};
218     m = A.size() - 1;
219     trans.resize(m);
220     for (int i = 0; i < m; ++i)
221     {
222         trans[i] = (mod - A[i + 1]) % mod;
223     }
224     std::reverse(trans.begin(), trans.end());
225     init = {s.begin(), s.begin() + m};
226 }
227 int64 calc(int64 n)
228 {
229     if (mod == 1) return 0;
230     if (n < m) return init[n];
231     vec v(m), u(m << 1);
232     int msk = !!n;
233     for (int64 m = n; m > 1; m >>= 1) msk <<= 1;
234     v[0] = 1 % mod;
235     for (int x = 0; msk; msk >>= 1, x <<= 1)
236     {
237         std::fill_n(u.begin(), m * 2, 0);
238         x |= !(n & msk);
239         if (x < m)
240             u[x] = 1 % mod;
241         else
242         { // can be optimized by fft/ntt
243             for (int i = 0; i < m; ++i)
244             {
245                 for (int j = 0, t = i + (x & 1); j < m; ++j, ++t)

```

```
246         {
247             u[t] = (u[t] + v[i] * v[j]) % mod;
248         }
249     }
250     for (int i = m * 2 - 1; i >= m; --i)
251     {
252         for (int j = 0, t = i - m; j < m; ++j, ++t)
253         {
254             u[t] = (u[t] + trans[j] * u[i]) % mod;
255         }
256     }
257 }
258 v = {u.begin(), u.begin() + m};
259 }
260 int64 ret = 0;
261 for (int i = 0; i < m; ++i)
262 {
263     ret = (ret + v[i] * init[i]) % mod;
264 }
265 return ret;
266 }
267
268 vec init, trans;
269 int64 mod;
270 int m;
271 };
```

1.4 组合数学

1.4.1 Lucas

```
1  const int maxn = 1e6 + 10;
2
3  ll fac[maxn], inv[maxn], facinv[maxn];
4
5  void init() {
6      fac[0] = inv[0] = facinv[0] = 1;
7      fac[1] = inv[1] = facinv[1] = 1;
8      for (int i = 2; i < maxn; i++) {
9          fac[i] = fac[i - 1] * i % mod;
10         inv[i] = mod - mod / i * inv[mod % i] % mod;
11         facinv[i] = facinv[i - 1] * inv[i] % mod;
12     }
13 }
14
15 ll C(int n, int k) {
16     if (k > n || k < 0) return 0;
17     return fac[n] * facinv[k] % mod * facinv[n - k] % mod;
18 }
19
20 ll lucas(ll n, ll m) {
21     ll res = 1;
22     while (n && m) {
23         res = res * C(n % mod, m % mod) % mod;
24         n /= mod;
25         m /= mod;
26     }
27     return res;
28 }
```


1.4.2 exLucas

```
1  typedef long long ll;
2
3  ll p, n, m;
4
5  ll exgcd(ll a, ll b, ll &x, ll &y)
6  {
7      if(!b) { x = 1; y = 0; return a; }
8      ll res = exgcd(b, a % b, x, y), t;
9      t = x, x = y, y = t - a / b * y;
10     return res;
11 }
12
13 ll qp(ll a, ll n, ll mod)
14 {
15     ll ans = 1, base = a;
16     for(; n; n >>= 1, (base *= base) %= mod) if(n & 1) (ans *= base) %= mod;
17     return ans;
18 }
19
20 ll fac(ll n, ll a, ll b)
21 {
22     if(!n) return 1;
23     ll res = 1;
24     for(ll i = 2; i <= b; i++)
25         if(i % a) (res *= i) %= b;
26     res = qp(res, n / b, b);
27     for(ll i = 2; i <= n % b; i++)
28         if(i % a) (res *= i) %= b;
29     return res * fac(n / a, a, b) % b;
30 }
31
32 ll inv(ll n, ll mod)
33 {
34     ll x, y;
35     exgcd(n, mod, x, y);
36     return (x += mod) > mod ? x - mod : x;
37 }
38
39 ll CRT(ll b, ll mod) { return b * inv(p / mod, mod) % p * (p / mod) % p; }
40
41 ll C(ll n, ll m, ll a, ll b)
42 {
43     ll up = fac(n, a, b), d1 = fac(m, a, b), d2 = fac(n - m, a, b);
44     ll k = 0;
45     for(ll i = n; i; i /= a) k += i / a;
46     for(ll i = m; i; i /= a) k -= i / a;
47     for(ll i = n - m; i; i /= a) k -= i / a;
48     return up * inv(d1, b) % b * inv(d2, b) % b * qp(a, k, b) % b;
49 }
50
51 ll exlucas(ll n, ll m)
52 {
53     ll res = 0, tmp = p, b;
54     int lim = sqrt(p) + 5;
55     for(int i = 2; i <= lim; i++) if(tmp % i == 0)
56     {
57         b = 1;
```

```
58     while(tmp % i == 0) tmp /= i, b *= i;
59     (res += CRT(C(n, m, i, b), b)) %= p;
60 }
61 if(tmp > 1) (res += CRT(C(n, m, tmp, tmp), tmp)) %= p;
62 return res;
63 }
```

1.4.3 递推组合数

```
1 // ---
2 // $0 \leq m \leq n \leq 1000$
3 // ---
4 const int maxn = 1010;
5 ll C[maxn][maxn];
6 void init() {
7     C[0][0] = 1;
8     for (int i = 1; i < maxn; i++)
9     {
10         C[i][0] = 1;
11         for (int j = 1; j <= i; j++) C[i][j] = (C[i - 1][j - 1] + C[i - 1][j]) % mod;
12     }
13 }
14
15 // ---
16 // $0 \leq m \leq n \leq 10^5$, 模$p$为素数
17 // ---
18 const int maxn = 100010;
19 ll f[maxn];
20 ll inv[maxn]; // 阶乘的逆元
21 void CalFact() {
22     f[0] = 1;
23     for (int i = 1; i < maxn; i++) f[i] = (f[i - 1] * i) % p;
24     inv[maxn - 1] = qp(f[maxn - 1], p - 2);
25     for (int i = maxn - 2; ~i; i--) inv[i] = inv[i + 1] * (i + 1) % p;
26 }
27 ll C(int n, int m) { return f[n] * inv[m] % p * inv[n - m] % p; }
```

1.4.4 小模数组合数

```
1 // ---
2 // $p$小$n$, $m$大
3 // ---
4
5 const int NICO = 100000+10;
6 const int MOD = 99991;
7 ll f[NICO];
8
9 ll Lucas(ll a, ll k)
10 {
11     ll res = 1;
12     while(a && k)
13     {
14         ll a1 = a % MOD;
15         ll b1 = k % MOD;
16         if(a1 < b1) return 0;
17         res = res * f[a1] * qp(f[b1] * f[a1 - b1] % MOD, MOD - 2) % MOD;
18         a /= MOD;
19         k /= MOD;
20     }
```

```
20     }
21     return res;
22 }
23
24 void init()
25 {
26     f[0] = 1;
27     for(int i=1;i<=MOD;i++)
28     {
29         f[i] = f[i-1]*i%MOD;
30     }
31 }
32
33 int main()
34 {
35     init();
36     cout << Lucas(5,2) << endl;
37 }
```

1.4.5 大模数组合数

```
1  // ---
2  // $n$, $m$ 小 $p$ 大
3  // ---
4  map<int, ll> m;
5
6  const int MOD = 1e9+7;
7  void fun(int n, int k) {
8      for (int i = 2; i <= sqrt(n * 1.0); i++) {
9          while (n % i == 0) {
10             n /= i;
11             m[i] += k;
12         }
13     }
14     if (n > 1) {
15         m[n] += k;
16     }
17 }
18
19 ll C(ll a, ll b) {
20     if (a < b || a < 0 || b < 0)
21         return 0;
22     m.clear();
23     ll ret = 1;
24     b = min(a - b, b);
25     for (int i = 0; i < b; i++) {
26         fun(a - i, 1);
27     }
28     for (int i = b; i >= 1; i--) {
29         fun(i, -1);
30     }
31     for (__typeof(m.begin()) it = m.begin(); it != m.end(); it++) {
32         if ((*it).second != 0) {
33             ret *= qp((*it).first, (*it).second);
34             ret %= MOD;
35         }
36     }
37     return ret;
}
```

```
38 }
39
40 int main(int argc, char *argv[])
41 {
42     ll a, b;
43     while (scanf("%lld%lld", &a, &b) != EOF) {
44         printf("%lld\n", C(a, b));
45     }
46     return 0;
47 }
```

1.5 多项式

1.5.1 FFT

```
1 // maxn 至少是大于m+n的2次方数
2 // m+n 2e5 maxn (1<<18)+50
3 // m+n 4e5 maxn (1<<19)+50
4 // m+n 2e6 maxn (1<<21)+50
5
6 // 普通fft
7 #include <bits/stdc++.h>
8
9 using namespace std;
10 typedef long long ll;
11
12 struct Complex {
13     double r, i;
14
15     Complex(double r, double i) : r(r), i(i) {}
16     Complex() {}
17     inline Complex operator+(const Complex &rhs) const { return Complex(r + rhs.r, i +
18         rhs.i); }
19     inline Complex operator-(const Complex &rhs) const { return Complex(r - rhs.r, i -
20         rhs.i); }
21     inline Complex operator*(const Complex &rhs) const { return Complex(r * rhs.r - i *
22         rhs.i, r * rhs.i + i * rhs.r); }
23     inline void operator/=(const double &x) { r /= x, i /= x; }
24     inline void operator*=(const Complex &rhs) { *this = Complex(r * rhs.r - i * rhs.i,
25         r * rhs.i + i * rhs.r); }
26     inline void operator+=(const Complex &rhs) { r += rhs.r, i += rhs.i; }
27 };
28
29 const int maxn = 4e6 + 6;
30 #define PI 3.14159265354
31 int pos[maxn];
32
33 void init(const int &n) {
34     for (int i = 0, j = 0; i < n; ++i) {
35         pos[i] = j;
36         for (int l = n >> 1; (j ^= 1) < 1; l >>= 1);
37     }
38 }
39
40 void transform(Complex *a, const int &n, bool inverse) {
41     for (int i = 0; i < n; ++i) if (i > pos[i]) std::swap(a[i], a[pos[i]]);
42     for (int l = 2; l <= n; l <= 1) {
43         int m = l / 2;
44         for (int i = 0; i < n; i += m) {
45             for (int j = 0; j < m; ++j) {
46                 Complex u = a[i + j], v = a[i + j + m];
47                 if (inverse) {
48                     Complex t1 = u + v, t2 = u - v;
49                     u = (t1 + t2 * Complex(0, 1)) * 0.5;
50                     v = (t1 - t2 * Complex(0, 1)) * 0.5;
51                 } else {
52                     Complex t1 = u + v, t2 = u - v;
53                     u = (t1 + t2 * Complex(0, 1)) * 0.5;
54                     v = (t1 - t2 * Complex(0, 1)) * 0.5;
55                 }
56                 a[i + j] = u;
57                 a[i + j + m] = v;
58             }
59         }
60     }
61 }
```

```

40     Complex omega = {cos(2 * PI / l), inverse ? -sin(2 * PI / l) : sin(2 * PI / l)
41     };
42     for (Complex *p = a; p != a + n; p += l) {
43         Complex x = {1, 0};
44         for (int i = 0; i < m; ++i, x *= omega) {
45             Complex t = x * p[m + i];
46             p[m + i] = p[i] - t;
47             p[i] += t;
48         }
49     }
50 }
51
52 void dft(Complex *a, const int &n) { transform(a, n, 0); }
53
54 void idft(Complex *a, const int &n) {
55     transform(a, n, 1);
56     for (int i = 0; i < n; ++i) a[i] /= n;
57 }
58
59 Complex A[maxn], B[maxn], C[maxn];
60
61 void FFT(int n, int m) { // len(A), len(B)
62     int cnt = 1;
63     while (cnt <= (n + m)) cnt <<= 1;
64     init(cnt);
65     dft(A, cnt);
66     dft(B, cnt);
67     for (int i = 0; i < cnt; i++) C[i] = A[i] * B[i];
68     idft(C, cnt);
69     for (int i = 0; i <= n + m; i++) C[i].r = ll(C[i].r + 0.01);
70 }
71
72 int main() {
73     int n, m, tem;
74     cin >> n >> m;
75     for (int i = 0; i <= n; i++) scanf("%d", &tem), A[i].r = tem;
76     for (int i = 0; i <= m; i++) scanf("%d", &tem), B[i].r = tem;
77     FFT(n, m);
78     for (int i = 0; i <= n + m; i++) printf("%lld ", ll(C[i].r));
79     cout << "\n";
80 }
81
82 // 两次变换
83 #define ld double
84 const ld PI = acos(-1);
85 struct cplx {
86     ld a, b;
87     cplx(ld a = 0, ld b = 0) : a(a), b(b) {}
88     const cplx operator+(const cplx &c) const { return cplx(a + c.a, b + c.b); }
89     const cplx operator-(const cplx &c) const { return cplx(a - c.a, b - c.b); }
90     const cplx operator*(const cplx &c) const { return cplx(a * c.a - b * c.b, a * c.b
91         + b * c.a); }
92     const cplx operator/(const ld &x) const { return cplx(a / x, b / x); }
93     const cplx conj() const { return cplx(a, -b); }
94 };
95 int rev[maxn];
96 cplx w[maxn];
97 cplx f[maxn];

```

```

97 void prepare(int &n) {
98     int sz = __builtin_ctz(n);
99     for (int i = 1; i < n; ++i) rev[i] = (rev[i] >> 1) >> 1 | ((i & 1) << (sz - 1));
100    w[0] = 0, w[1] = 1, sz = 1;
101    while (1 << sz < n) {
102        cplx w_n = cplx(cosl(2 * PI / (1 << (sz + 1))), sinl(2 * PI / (1 << (sz + 1))))
103        ;
104        for (int i = 1 << (sz - 1); i < (1 << sz); ++i) {
105            w[i << 1] = w[i], w[i << 1 | 1] = w[i] * w_n;
106        }
107        ++sz;
108    }
109    void fft(cplx *a, int n) {
110        for (int i = 1; i < n - 1; ++i) if (i < rev[i]) swap(a[i], a[rev[i]]);
111        for (int h = 1; h < n; h <= 1) {
112            for (int s = 0; s < n; s += h << 1) {
113                for (int i = 0; i < h; ++i) {
114                    cplx &u = a[s + i], &v = a[s + i + h], t = v * w[h + i];
115                    v = u - t, u = u + t;
116                }
117            }
118        }
119    }
120    template<class T>
121    vector<T> multiply(const vector<T> &a, const vector<T> &b) {
122        int n = a.size(), m = b.size(), sz = 1;
123        while (sz < n + m - 1) sz <= 1;
124        prepare(sz);
125        for (int i = 0; i < sz; ++i) f[i] = cplx(i < n ? a[i] : 0, i < m ? b[i] : 0);
126        fft(f, sz);
127        for (int i = 0; i <= (sz >> 1); ++i) {
128            int j = (sz - i) & (sz - 1);
129            cplx x = (f[i] * f[i] - (f[j] * f[j]).conj()) * cplx(0, -0.25);
130            f[j] = x, f[i] = x.conj();
131        }
132        fft(f, sz);
133        vector<T> c(n + m - 1);
134        for (int i = 0; i < n + m - 1; ++i) {
135            c[i] = ((T) (f[i].a / sz + 0.3));
136        }
137        return c;
138    }
139    #undef ld

```

1.5.2 NTT

```

1  const static int N = 4e6 + 6;
2  typedef long long ll;
3  ll pos[N];
4  const ll mod = 998244353, root = 3;
5  inline ll fastpow(ll a, ll b) {
6      ll ans = 1;
7      for (; b; a=a%mod,b>>=1) if(b&1)ans=ans*a%mod;
8      return ans;
9  }
10 inline void exgcd(ll a,ll b,ll &g,ll &x,ll &y) {
11     if (!b) g=a,x=1,y=0;

```

```

12     else exgcd(b,a%b,g,y,x),y-=x*(a/b);
13 }
14 inline ll inv(ll a) {
15     ll g,x,y;
16     exgcd(a,mod,g,x,y);
17     return (x%mod+mod)%mod;
18 }
19 void init(const int &n) {
20     for (int i = 0,j=0; i < n; ++i) {
21         pos[i]=j;for (int l = n >> 1; (j ^= 1) < 1; l >>= 1);
22     }
23 }
24 void transform(ll *a, const int &n, bool inverse) {
25     for (int i=0; i<n;++i) if(i>pos[i]) swap(a[i],a[pos[i]]);
26     for (int l=2; l<=n;l<=1) {
27         int m=l/2;ll omega=fastpow(inverse?inv(root):root,(mod-1)/l);
28         for (ll *p=a;p!=a+n;p+=l) {
29             ll x=1;
30             for (int i=0;i<m;++i,x=x*omega%mod) {
31                 ll t=x*p[m+i]%mod;
32                 p[m+i]=(p[i]-t+mod)%mod;(p[i]+=t)%=mod;
33             }
34         }
35     }
36 }
37 void dft(ll *a, const int &n) {
38     transform(a,n,0);
39 }
40 void idft(ll *a, const int &n) {
41     const ll INV=inv(n);
42     transform(a,n,1);
43     for (int i=0;i<n;i++) a[i]=a[i]*INV % mod;
44 }
45 const int maxn=4e6+6;
46 ll a[maxn],b[maxn],c[maxn],pic[maxn],pec[maxn],plc[maxn],ppc[maxn];
47 int INV[maxn];
48 void poly_inv(int n,ll *a,ll *b) {
49     if(n==1) {b[0]=inv(a[0]);return;}
50     poly_inv((n+1)/2,a,b);
51     int cnt=1;while(cnt<=n*2) cnt<=1;init(cnt);
52     copy(a,a+n,pic);fill(pic+n,pic+cnt,0);fill(b+n,b+cnt,0);dft(pic,cnt);dft(b,cnt);
53     for(int i=0;i<cnt;i++) (b[i]*=(2ll-pic[i]*b[i])%mod)%=mod;
54     for(int i=0;i<cnt;i++) b[i]=(b[i]+mod)%mod;
55     idft(b,cnt);fill(b+n,b+cnt,0);
56 }
57 void poly_ln(int n,ll *a,ll *b) { //G'=F'/F
58     poly_inv(n,a,b);
59     int cnt=1;while(cnt<=n*2-3) cnt<=1;init(cnt);
60     for(int i=0;i<n-1;i++) plc[i]=a[i+1]*(i+1)%mod;
61     fill(plc+n-1,plc+cnt,0);fill(b+n,b+cnt,0);dft(plc,cnt);dft(b,cnt);
62     for(int i=0;i<cnt;i++) b[i]=plc[i]*b[i]%mod;
63     idft(b,cnt);for(int i=n-1;i>=1;i--) b[i]=b[i-1]*INV[i]%mod;b[0]=0;
64     fill(b+n,b+cnt,0);
65 }
66 void poly_exp(int n,ll *a,ll *b) {
67     if(n==1) {b[0]=1;return;}
68     poly_exp((n+1)/2,a,b);poly_ln(n,b,pec);
69     int cnt=1;while(cnt<=n*2-2) cnt<=1;init(cnt);
70     for(int i=0;i<n;i++) pec[i]=(-pec[i]+a[i])%mod;pec[0]+=1;

```

```

71     dft(pec,cnt);dft(b,cnt);
72     for(int i=0;i<cnt;i++) b[i]=(b[i]*pec[i])%mod;
73     idft(b,cnt);fill(b+n,b+cnt,0);
74 }
75 void poly_pow(int n,ll *a,ll *b,ll k) {
76     poly_ln(n,a,ppc);
77     for(int i=0;i<n;i++) ppc[i]*=k;
78     poly_exp(n,ppc,b);
79 }
80 char buff[maxn];
81 int main(){
82     ios::sync_with_stdio(false);
83     cin.tie(0);cout.tie(0);
84     int n;cin>>n>>buff;
85     INV[1]=1;for(int i=2;i<n;++i) INV[i]=(mod-mod/i)*INV[mod%i]%mod;
86     ll k=0;
87     for(int i=0;buff[i];i++) k=(k*10+buff[i]-'0')%mod;
88     for(int i=0;i<n;i++) cin>>a[i];
89     poly_pow(n,a,b,k);
90     for(int i=0;i<n;i++) cout<<b[i]<<' ';
91 }

```

1.5.3 MTT

```

1  const int maxn = 262144;
2  int mod = 1e9+7;
3
4  namespace MTT {
5      struct comp {
6          double x, y;
7          comp(double x_ = 0, double y_ = 0): x(x_), y(y_) {}
8      };
9      inline comp operator + (const comp& a, const comp& b) { return comp(a.x + b.x, a.y
10         + b.y); }
11     inline comp operator - (const comp& a, const comp& b) { return comp(a.x - b.x, a.y
12         - b.y); }
13     inline comp operator * (const comp& a, const comp& b) { return comp(a.x * b.x - a.y
14         * b.y, a.x * b.y + a.y * b.x); }
15     inline comp conj(const comp& a) { return comp(a.x, -a.y); }
16
17     const double PI = acos(-1);
18     comp w[maxn + 5];
19     int bitrev[maxn + 5];
20     int N, L;
21
22     void fft(vector<comp>& a) {
23         for (int i = 0; i < N; ++i) if (i < bitrev[i]) swap(a[i], a[bitrev[i]]);
24         for (int i = 2, lyc = N >> 1; i <= N; i <= 1, lyc >>= 1) {
25             for (int j = 0; j < N; j += i) {
26                 int pl = j, pr = j + (i >> 1), p = 0;
27                 for (int k = 0; k < (i >> 1); ++k) {
28                     comp tmp = a[pr] * w[p];
29                     a[pr] = a[pl] - tmp, a[pl] = a[pl] + tmp;
30                     ++ pl, ++ pr, p += lyc;
31                 }
32             }
33         }
34     }
35 }

```



```

32
33 inline void fft_prepare() {
34     for (int i = 0; i < N; ++i) bitrev[i] = bitrev[i >> 1] >> 1 | ((i & 1) << (L -
        1));
35     for (int i = 0; i < N; ++i) w[i] = comp(cos(2 * PI * i / N), sin(2 * PI * i / N
        ));
36 }
37
38 // if max ans is n+m, n=n+m+1
39 void init(int n) {
40     L = 0;
41     while ((1 << L) < n) ++ L;
42     N = 1 << L;
43     fft_prepare();
44 }
45
46 inline vector<int> conv(vector<int> x, vector<int> y) {
47     vector<comp> a(N + 5), b(N + 5);
48     vector<comp> dfta(N + 5), dftb(N + 5), dftc(N + 5), dftd(N + 5);
49     for (int i = 0; i < (int)x.size(); ++i) {
50         if ((x[i] %= mod) < 0) x[i] += mod;
51         a[i] = comp(x[i] & 32767, x[i] >> 15);
52     }
53     for (int i = 0; i < (int)y.size(); ++i) {
54         if ((y[i] %= mod) < 0) y[i] += mod;
55         b[i] = comp(y[i] & 32767, y[i] >> 15);
56     }
57     fft(a), fft(b);
58     for (int i = 0; i < N; ++i) {
59         int j = (N - i) & (N - 1);
60         comp da = (a[i] + conj(a[j])) * comp(0.5, 0);
61         comp db = (a[i] - conj(a[j])) * comp(0, -0.5);
62         comp dc = (b[i] + conj(b[j])) * comp(0.5, 0);
63         comp dd = (b[i] - conj(b[j])) * comp(0, -0.5);
64         dfta[j] = da * dc;
65         dftb[j] = da * dd;
66         dftc[j] = db * dc;
67         dftd[j] = db * dd;
68     }
69     for (int i = 0; i < N; ++i) a[i] = dfta[i] + dftb[i] * comp(0, 1);
70     for (int i = 0; i < N; ++i) b[i] = dftc[i] + dftd[i] * comp(0, 1);
71     fft(a), fft(b);
72     vector<int> z(N);
73     for (int i = 0; i < N; ++i) {
74         int da = (ll)(a[i].x / N + 0.5) % mod;
75         int db = (ll)(a[i].y / N + 0.5) % mod;
76         int dc = (ll)(b[i].x / N + 0.5) % mod;
77         int dd = (ll)(b[i].y / N + 0.5) % mod;
78         z[i] = (da + ((ll)(db + dc) << 15) + ((ll)dd << 30)) % mod;
79         if (z[i] < 0) z[i] += mod;
80     }
81     while (!z.empty() && z.back() == 0) z.pop_back();
82     return z;
83 }
84 }
85
86 int main(int argc, char* argv[]) {
87     int n, m;
88     R(n, m);

```

```

89     VI a(n + 5), b(m + 5);
90     MTT::init(n + m + 1);
91     for (int i = 0; i <= n; ++i) R(a[i]);
92     for (int i = 0; i <= m; ++i) R(b[i]);
93     auto res = MTT::conv(a, b);
94     while (res.size() < (n + m + 1)) res.push_back(0);
95     W(res);
96     return 0;
97 }

```

1.5.4 FWT

```

1 //FWT 完后需要先模一遍
2 template<typename T>
3 void fwt(ll a[], int n, T f) {
4     for (int d = 1; d < n; d *= 2)
5         for (int i = 0, t = d * 2; i < n; i += t)
6             for (int j = 0; j < d; j++)
7                 f(a[i + j], a[i + j + d]);
8 }
9
10 void AND(ll& a, ll& b) { a += b; }
11 void OR(ll& a, ll& b) { b += a; }
12 void XOR (ll& a, ll& b) {
13     ll x = a, y = b;
14     a = (x + y) % mod;
15     b = (x - y + mod) % mod;
16 }
17 void rAND(ll& a, ll& b) { a -= b; }
18 void rOR(ll& a, ll& b) { b -= a; }
19 void rXOR(ll& a, ll& b) {
20     static ll INV2 = (mod + 1) / 2;
21     ll x = a, y = b;
22     a = (x + y) * INV2 % mod;
23     b = (x - y + mod) * INV2 % mod;
24 }
25
26 //FWT 子集卷积 i与j=0, i或j为k
27 a[popcount(x)][x] = A[x]
28 b[popcount(x)][x] = B[x]
29 fwt(a[i]) fwt(b[i])
30 c[i + j][x] += a[i][x] * b[j][x]
31 rfwt(c[i])
32 ans[x] = c[popcount(x)][x]

```

1.5.5 杜教筛

```

1 const int maxn=3e6+5;
2
3 unordered_map<int,ll> Mmu;
4 unordered_map<int,ll> Mphi;
5 ll mu[maxn],phi[maxn],prim[maxn],pcnt,mn[maxn];
6 void get_T() {
7     phi[1]=mu[1]=1;
8     for(int i=2;i<maxn;i++) {
9         if(!mn[i]) {
10             mn[i]=i;prim[++pcnt]=i;
11             phi[i]=i-1;mu[i]=-1;

```

```

12     }
13     for(int j=1;j<=pcnt&& i*prim[j]<maxn;j++) {
14         mn[i*prim[j]]=prim[j];
15         if(i%prim[j]==0) {
16             phi[i*prim[j]]=phi[i]*prim[j];
17             break;
18         }
19         mu[i*prim[j]]=-mu[i];
20         phi[i*prim[j]]=phi[i]*(prim[j]-1);
21     }
22 }
23 for(int i=1;i<maxn;i++)
24     mu[i]+=mu[i-1], phi[i]+=phi[i-1];
25 }
26 ll Smu(int n) {
27     if(n==1) return 1;
28     if(n<maxn) return mu[n];
29     if(Mmu.find(n)!=Mmu.end()) return Mmu[n];
30     ll ans=0;
31     for(int l=2,r;l<=n;l=r+1) {
32         r=min(n,n/(n/l));
33         ans+=Smu(n/l)*(r-l+1);
34     }
35     return Mmu[n]=1-ans;
36 }
37 ll Sphi(int n) {
38     if(n==1) return 1;
39     if(n<maxn) return phi[n];
40     if(Mphi.find(n)!=Mphi.end()) return Mphi[n];
41     ll ans=0;
42     for(int l=2,r;l<=n;l=r+1) {
43         r=min(n,n/(n/l));
44         ans+=Sphi(n/l)*(r-l+1);
45     }
46     return Mphi[n]=1ll*n*(n+1)/2-ans;
47 }
48 int main() {
49     get_T();
50     int T;cin>>T;while(T--) {
51         int n;cin>>n;
52         cout<<Sphi(n)<<' ' <<Smu(n)<<endl;
53     }
54     return 0;
55 }

```

1.5.6 min25

```

1  const int maxn=1e6+5;
2  int primer[maxn],pcnt;
3  bool insp[maxn];
4  ll fsum1[maxn],fsum2[maxn];
5  const int mod=1e9+7,INV6=166666668;
6  void sieve() {
7      for(int i=2;i<maxn;i++) {
8          if(!insp[i]) {
9              primer[++pcnt]=i;
10             fsum1[pcnt]=(fsum1[pcnt-1]+i)%mod;
11             fsum2[pcnt]=(fsum2[pcnt-1]+1ll*i*i)%mod;

```

```

12     }
13     for(int j=1;j<=pcnt&&primer[j]*i<maxn;j++) {
14         insp[i*primer[j]]=1;
15         if(i%primer[j]==0) break;
16     }
17 }
18 }
19
20 // Let g(j,m) be the sum of f'(i) for all i which is not greater
21 // than m and is prime or the min pri-factor of i > p[j]
22 // We got that g(j,m)=g(j-1,m)-f'(p[j])*( g(j-1,m/p[j])-\sum_{i=1}^{j-1}f'(p[i]) )
23 // g(0,m) = \sum_{i=2}^m f'(j)
24 // Here f'(j) is a function whose value equals f(j) when j is prime
25 // and f'(j) is an acompletely multiplicative function
26 // We can use DP to calc g(j,{n/1,n/2,...,n/n}) (sqrt(n) situation in total)
27 // So g(j,m) equals to the sum of f(i) for all primer i which not greater than m
28 // when p[j]*p[j] > m
29 int id1[maxn],id2[maxn],tot=0;
30 ll val[maxn],n,g1[maxn],g2[maxn];
31 //index of n/1,n/2...n/n
32 inline int getid(ll m) {
33     if(m<=n/m) return id1[m];
34     return id2[n/m];
35 }
36 void init() {
37     for(ll l=1,r;l<=n;l=r+1) {
38         r=n/(n/l);ll w=n/l;
39         val[++tot]=w;
40         if(w<=n/w) id1[w]=tot;
41         else id2[n/w]=tot;
42         w%=mod;
43         g1[tot]=w*(w+1)/2%mod;
44         g2[tot]=w*(w+1)%mod*(2ll*w+1)%mod*INV6%mod;
45         g1[tot]=(g1[tot]-1+mod)%mod;
46         g2[tot]=(g2[tot]-1+mod)%mod;
47     }
48     // Start DP
49     for(int i=1;1ll*primer[i]*primer[i]<=n;i++) {
50         for(int j=1;j<=tot&&1ll*primer[i]*primer[i]<=val[j];j++) {
51             int k=getid(val[j]/primer[i]);
52             g1[j]=(g1[j]-1ll*primer[i]*(g1[k]-fsum1[i-1]+mod)%mod+mod)%mod;
53             g2[j]=(g2[j]-1ll*primer[i]*primer[i]%mod*(g2[k]-fsum2[i-1]+mod)%mod+mod)%mod;
54         }
55     }
56 }
57 ll getS(ll m,int k) {
58     if(m<primer[k]) return 0;
59     ll res=(g2[getid(m)]-g1[getid(m)]-fsum2[k-1]+fsum1[k-1])%mod;
60     for(int i=k;1ll*primer[i]*primer[i]<=m;i++) {
61         for(ll pie=primer[i],pie1=pie*primer[i];pie1<=m;pie=pie1,pie1*=primer[i]){
62             ll t1=pie%mod,t2=pie1%mod;
63             res=(res+t1*(t1-1)%mod*getS(m/pie,i+1)%mod+t2*(t2-1)%mod)%mod;
64         }
65     }
66     return (res+mod)%mod;
67 }
68 int main() {
69     sieve();cin>>n;

```

```

70     init();
71     // for(int i=1;i<=10;i++) cout<<g1[i]<<endl;
72     // cout<<g1[1]<<endl;
73     cout<<(getS(n,1)+1)%mod<<endl;
74     return 0;
75 }

```

1.6 Others

1.6.1 公式

- 约数定理: 若 $n = \prod_{i=1}^k p_i^{a_i}$, 则
 - 约数个数 $f(n) = \prod_{i=1}^k (a_i + 1)$
 - 约数和 $g(n) = \prod_{i=1}^k (\sum_{j=0}^{a_i} p_i^j)$
- 小于 n 且互素的数之和为 $n\varphi(n)/2$
- 若 $\gcd(n, i) = 1$, 则 $\gcd(n, n-i) = 1 (1 \leq i \leq n)$
- 错排公式: $D(n) = (n-1)(D(n-2) + D(n-1)) = \sum_{i=2}^n \frac{(-1)^k n!}{k!} = \lfloor \frac{n!}{e} + 0.5 \rfloor$
- 部分错排公式: $n + m$ 个数中 m 个数必须错排求排列数
 - $1 \text{ dp}[i] = n * \text{dp}[i-1] + (i-1) * (\text{dp}[i-1] + \text{dp}[i-2]);$
 - $2 \text{ dp}[0] = n!;$
 - $3 \text{ dp}[1] = n * n!;$
 - $\text{dp}[m]$ 为所求解
- 海伦公式: $S = \sqrt{p(p-a)(p-b)(p-c)}$, 其中 $p = \frac{(a+b+c)}{2}$
- 求 $C(n, k)$ 中素因子 P 的个数: 把 n 转化为 P 进制, 并记它每个位上的和为 $S1$ 把 $n-k, k$ 做同样的处理, 得到 $S2, S3$ 则答案为: $\frac{S2+S3-S1}{P-1}$
- 威尔逊定理: $p \text{ is prime} \Rightarrow (p-1)! \equiv -1 \pmod{p}$
- 欧拉定理: $\gcd(a, n) = 1 \Rightarrow a^{\varphi(n)} \equiv 1 \pmod{n}$
- 欧拉定理推广: $\gcd(n, p) = 1 \Rightarrow a^n \equiv a^{n \% \varphi(p)} \pmod{p}$
- 模的幂公式: $a^n \pmod{m} = \begin{cases} a^n \pmod{m} & n < \varphi(m) \\ a^{n \% \varphi(m) + \varphi(m)} \pmod{m} & n \geq \varphi(m) \end{cases}$
- 素数定理: 对于不大于 n 的素数个数 $\pi(n)$, $\lim_{n \rightarrow \infty} \pi(n) = \frac{n}{\ln n}$
- 位数公式: 正整数 x 的位数 $N = \log_{10}(n) + 1$
- 斯特灵公式 $n! \approx \sqrt{2\pi n} (\frac{n}{e})^n$
- 设 $a > 1, m, n > 0$, 则 $\gcd(a^m - 1, a^n - 1) = a^{\gcd(m, n)} - 1$
- 设 $a > b, \gcd(a, b) = 1$, 则 $\gcd(a^m - b^m, a^n - b^n) = a^{\gcd(m, n)} - b^{\gcd(m, n)}$

$$G = \gcd(C_n^1, C_n^2, \dots, C_n^{n-1}) = \begin{cases} n, & n \text{ is prime} \\ 1, & n \text{ has multy prime factors} \\ p, & n \text{ has single prime factor } p \end{cases}$$

$$\gcd(\text{Fib}(m), \text{Fib}(n)) = \text{Fib}(\gcd(m, n))$$

- 求和公式:

- $\sum k = \frac{n(n+1)}{2}$
- $\sum 2k - 1 = n^2$
- $\sum k^2 = \frac{n(n+1)(2n+1)}{6}$
- $\sum (2k-1)^2 = \frac{n(4n^2-1)}{3}$

- (e) $\sum k^3 = \left(\frac{n(n+1)}{2}\right)^2$
 (f) $\sum (2k-1)^3 = n^2(2n^2-1)$
 (g) $\sum k^4 = \frac{n(n+1)(2n+1)(3n^2+3n-1)}{30}$
 (h) $\sum k^5 = \frac{n^2(n+1)^2(2n^2+2n-1)}{12}$
 (i) $\sum k(k+1) = \frac{n(n+1)(n+2)}{3}$
 (j) $\sum k(k+1)(k+2) = \frac{n(n+1)(n+2)(n+3)}{4}$
 (k) $\sum k(k+1)(k+2)(k+3) = \frac{n(n+1)(n+2)(n+3)(n+4)}{5}$

18. 若 $\gcd(m, n) = 1$, 则:

- (a) 最大不能组合的数为 $m * n - m - n$
 (b) 不能组合数个数 $N = \frac{(m-1)(n-1)}{2}$

19. $(n+1)lcm(C_n^0, C_n^1, \dots, C_n^{n-1}, C_n^n) = lcm(1, 2, \dots, n+1)$

20. 若 p 为素数, 则 $(x+y+\dots+w)^p \equiv x^p + y^p + \dots + w^p \pmod{p}$

21. 卡特兰数: 1, 1, 2, 5, 14, 42, 132, 429, 1430, 4862, 16796, 58786, 208012

$$h(0) = h(1) = 1, h(n) = \frac{(4n-2)h(n-1)}{n+1} = \frac{C_{2n}^n}{n+1} = C_{2n}^n - C_{2n}^{n-1}$$

22. 伯努利数: $B_n = -\frac{1}{n+1} \sum_{i=0}^{n-1} C_{n+1}^i B_i$

$$\sum_{i=1}^n i^k = \frac{1}{k+1} \sum_{i=1}^{k+1} C_{k+1}^i B_{k+1-i} (n+1)^i$$

23. 二项式反演:

$$f_n = \sum_{i=0}^n (-1)^i \binom{n}{i} g_i \Leftrightarrow g_n = \sum_{i=0}^n (-1)^i \binom{n}{i} f_i$$

$$f_n = \sum_{i=0}^n \binom{n}{i} g_i \Leftrightarrow g_n = \sum_{i=0}^n (-1)^{n-i} \binom{n}{i} f_i$$

24. 莫比乌斯反演:

(a) 令 $f(d) = \sum_{i=1}^n \sum_{j=1}^m [\gcd(i, j) = d]$

(b) $F(n) = \sum_{n|d} f(d) = \lfloor \frac{N}{n} \rfloor \lfloor \frac{M}{n} \rfloor$

(c) 有 $f(n) = \sum_{n|d} \mu(\lfloor \frac{d}{n} \rfloor) F(d)$

(d) $\phi(n) = \sum_{d|n} d * \mu(n/d)$

25. 2 的 n 次方, 在 pow 时可以精确输出最大 2^{1023} , pow(2,1023)

26. FFT 常用素数

$r \cdot 2^k + 1$	r	k	g
3	1	1	2
5	1	2	2
17	1	4	3
97	3	5	5
193	3	6	5
257	1	8	3
7681	15	9	17
12289	3	12	11
40961	5	13	3
65537	1	16	3
786433	3	18	10
5767169	11	19	3
7340033	7	20	3
23068673	11	21	3
104857601	25	22	3
167772161	5	25	3
469762049	7	26	3
998244353	119	23	3
1004535809	479	21	3
2013265921	15	27	31
2281701377	17	27	3
3221225473	3	30	5
75161927681	35	31	3
77309411329	9	33	7
206158430209	3	36	22
2061584302081	15	37	7
2748779069441	5	39	3
6597069766657	3	41	5
39582418599937	9	42	5
79164837199873	9	43	5
263882790666241	15	44	7
1231453023109121	35	45	3
1337006139375617	19	46	3
3799912185593857	27	47	5
4222124650659841	15	48	19
7881299347898369	7	50	6
31525197391593473	7	52	3
180143985094819841	5	55	6
1945555039024054273	27	56	5
4179340454199820289	29	57	3

1.7 高数

1.7.1 拉格朗日插值

```

1  const int maxn = 1e5 + 10;
2  int x[maxn], y[maxn];
3  int qp(int a, int n) {
4      ll ans = 1, base = a;
5      for(; n; (base *= base) %= mod, n >>= 1) if(n & 1) (ans *= base) %= mod;
6      return ans;
7  }
8
9  int lagrange(int n, int *x, int *y, int xi) {
10     int ans = 0;
11     for(int i = 0; i <= n; i++) {
12         int s1 = 1, s2 = 1;
13         for(int j = 0; j <= n; j++) if(i != j) {

```

```
14         s1 = 1ll * s1 * (xi - x[j]) % mod;
15         s2 = 1ll * s2 * (x[i] - x[j]) % mod;
16     }
17     ans = (1ll * ans + 1ll * y[i] * s1 % mod * qp(s2, mod - 2) % mod) % mod;
18 }
19 return (ans + mod) % mod;
20 }
```

1.7.2 拉格朗日插值 (连续取值)

```
1 int x[maxn], y[maxn];
2 int s1[maxn], s2[maxn], ifac[maxn];
3
4 //如果x的取值是连续一段, 可以做到O(n)求解
5 int lagrange(int n, int *x, int *y, int xi) {
6     int ans = 0;
7     s1[0] = (xi - x[0]) % mod, s2[n + 1] = 1;
8     for(int i = 1; i <= n; i++) s1[i] = 1ll * s1[i - 1] * (xi - x[i]) % mod;
9     for(int i = n; i >= 0; i--) s2[i] = 1ll * s2[i + 1] * (xi - x[i]) % mod;
10    ifac[0] = ifac[1] = 1;
11    for(int i = 2; i <= n; i++) ifac[i] = -1ll * mod / i * ifac[mod % i] % mod;
12    for(int i = 2; i <= n; i++) ifac[i] = 1ll * ifac[i] * ifac[i - 1] % mod;
13    for(int i = 0; i <= n; i++)
14        (ans += 1ll * y[i] * (i == 0 ? 1 : s1[i - 1]) % mod * s2[i + 1] % mod * ifac[i]
15         % mod * (((n - i) & 1) ? -1 : 1) * ifac[n - i] % mod) %= mod;
16    return (ans + mod) % mod;
17 }
```

1.7.3 辛普森积分

```
1 typedef double ld;
2 ld a, b, c, d;
3 inline ld f(ld x) { return (c * x + d) / (a * x + b); }
4 inline ld simpson(ld l, ld r, ld fl, ld fr, ld fm) { return (r - l) * (fl + fr + 4 * fm
5     ) / 6; }
6 inline ld asr(ld l, ld r, ld eps, ld fl, ld fr, ld fm) {
7     ld flm = f((l + (l + r) / 2) / 2), frm = f((r + (l + r) / 2) / 2);
8     ld L = simpson(l, (l + r) / 2, fl, fm, flm), R = simpson((l + r) / 2, r, fm, fr,
9         frm), A = simpson(l, r, fl, fr, fm);
10    if (fabs(L + R - A) <= 15 * eps) return L + R + (L + R - A) / 15;
11    return asr(l, (l + r) / 2, eps / 2, fl, fm, flm) + asr((l + r) / 2, r, eps / 2, fm,
12        fr, frm);
13 }
14
15 inline ld cal(ld l, ld r) { return asr(l, r, 1e-10, f(l), f(r), f((l + r) / 2)); }
```


2 Graph Theory

2.1 路径

2.1.1 Dijkstra

```
1  const int maxn = 1e5 + 10;
2  const int inf = 0x3f3f3f3f;
3
4  int head[maxn], dis[maxn], cnt, n;
5
6  struct Edge { int nex,to,w; }edge[20*maxn];
7
8  void add(int u,int v,int w)
9  {
10     edge[++cnt].nex=head[u];
11     edge[cnt].w=w;
12     edge[cnt].to=v;
13     head[u]=cnt;
14 }
15
16 void dijkstra(int s)
17 {
18     priority_queue<pair<int, int>, vector<pair<int, int> >, greater<pair<int, int> > >
19         que;
20     memset(dis, 0x3f, sizeof dis);
21     que.push({0, s}); dis[s] = 0;
22     while(!que.empty())
23     {
24         auto f = que.top(); que.pop();
25         int u = f.second, d = f.first;
26         if(d != dis[u]) continue;
27         for(int i = head[u]; ~i; i = edge[i].nex)
28         {
29             int v = edge[i].to, w = edge[i].w;
30             if(dis[u] + w < dis[v])
31             {
32                 dis[v] = dis[u] + w;
33                 que.push({dis[v], v});
34             }
35         }
36     }
```

2.1.2 Euler Path

```
1  int S[N << 1], top;
2  Edge edges[N << 1];
3  set<int> G[N];
4
5  void DFS(int u) {
6      S[top++] = u;
7      for (int eid: G[u])
8      {
9          int v = edges[eid].get_other(u);
10         G[u].erase(eid);
11         G[v].erase(eid);
12         DFS(v);
13         return;
```

```

14     }
15 }
16
17 void fleury(int start)
18 {
19     int u = start;
20     top = 0; path.clear();
21     S[top++] = u;
22     while (top)
23     {
24         u = S[--top];
25         if (!G[u].empty())
26             DFS(u);
27         else path.push_back(u);
28     }
29 }

```

2.1.3 K shortest Path(Astar)

```

1 // 有向图 k 短路
2 const int N = 1010, M = 10010, inf = 1e9+50;
3 int n, m;
4 int g[N], h[N], v[M<<1], w[M<<1], nxt[M<<1], ed, d[N], vis[N], ans[N];
5 typedef pair<int, int> P;
6 priority_queue<P, vector<P>, greater<P>> Q;
7 void add(int x, int y, int z) {
8     v[++ed] = x, w[ed] = z, nxt[ed] = g[y], g[y] = ed;
9     v[++ed] = y, w[ed] = z, nxt[ed] = h[x], h[x] = ed;
10 }
11
12 int KthShortest(int S, int T, int k) {
13     int x;
14     for (int i = 1; i <= k; ++i) ans[i] = -1;
15     for (int i = 1; i <= n; ++i) d[i] = inf;
16     Q.push(P(d[T] = 0, T));
17     while (!Q.empty()) {
18         P t = Q.top(); Q.pop();
19         if (d[t.second] < t.first) continue;
20         for (int i = g[x = t.second]; i; i = nxt[i]) {
21             if (d[x] + w[i] < d[v[i]]) Q.push(P(d[v[i]] = d[x] + w[i], v[i]));
22         }
23     }
24     if (d[S] < inf) Q.push(P(d[S], S));
25     while (!Q.empty()) {
26         P t = Q.top(); Q.pop(); vis[x = t.second] ++;
27         if (x == T && vis[T] <= k) ans[vis[T]] = t.first;
28         if (vis[T] > k) break;
29         if (vis[x] <= k) for (int i = h[x]; i; i = nxt[i]) {
30             Q.push(P(t.first - d[x] + d[v[i]] + w[i], v[i]));
31         }
32     }
33     return ans[k];
34 }

```

2.1.4 K shortest Path(可持久化可并堆)

```

1 #include <bits/stdc++.h>
2 #include<ext/pb_ds/priority_queue.hpp>

```

```
3
4 using namespace std;
5
6 const int N = '';
7 const int M = '';
8 const int logM = 20;
9 const int inf = 0x3f3f3f3f;
10
11 int n, m, k, S, T;
12
13 struct Edge{ int nex, to, w; };
14
15 struct Graph
16 {
17     int head[N], cnt;
18     Edge edge[M];
19     void init(int n) { for(int i = 0; i <= n; i++) head[i] = 0; cnt = 0; }
20     void addedge(int u, int v, int val) { edge[++ cnt].nex = head[u], edge[cnt].to = v,
        edge[cnt].w = val, head[u] = cnt; }
21 }g, rg;
22
23 int dis[N];
24
25 void dijkstra()
26 {
27     priority_queue<pair<int, int>, vector<pair<int, int> >, greater<pair<int, int> > >
        que;
28     memset(dis, inf, sizeof dis);
29     que.push({0, T}); dis[T] = 0;
30     const int *head = rg.head; const Edge *edge = rg.edge;
31     while(!que.empty())
32     {
33         auto f = que.top(); que.pop();
34         int u = f.second, d = f.first;
35         if(d != dis[u]) continue;
36         for(int i = head[u]; i; i = edge[i].nex)
37         {
38             int v = edge[i].to, w = edge[i].w;
39             if(dis[u] + w < dis[v]) { dis[v] = dis[u] + w; que.push({dis[v], v}); }
40         }
41     }
42 }
43
44 bool tree_edge[M], vis[N];
45 int fa[N], st[N], top;
46
47 void dfs(int u)
48 {
49     vis[u] = true;
50     st[++ top] = u;
51     for(int i = rg.head[u]; i; i = rg.edge[i].nex)
52     {
53         int v = rg.edge[i].to;
54         if(!vis[v] && dis[v] == dis[u] + rg.edge[i].w)
55         {
56             fa[v] = u;
57             tree_edge[i] = true;
58             dfs(v);
59         }
60     }
```

```

60     }
61 }
62
63 namespace LT
64 {
65     int son[M * logM][2];
66     int ht[M * logM], val[M * logM], id[M * logM];
67     int tot;
68
69     int newnode(int _val, int _id, int _dis = 0)
70     {
71         int now = ++ tot;
72         val[now] = _val, id[now] = _id;
73         ht[now] = _dis, son[now][0] = son[now][1] = 0;
74         return now;
75     }
76
77     int _copy(int ori)
78     {
79         int now = ++tot;
80         val[now] = val[ori], id[now] = id[ori];
81         ht[now] = ht[ori], son[now][0] = son[ori][0], son[now][1] = son[ori][1];
82         return now;
83     }
84
85     int merge(int a, int b)
86     {
87         if(!a || !b) return a | b;
88         if(val[a] > val[b]) swap(a, b);
89         int now = _copy(a);
90         son[now][1] = merge(son[now][1], b);
91         if(ht[son[now][0]] < ht[son[now][1]]) swap(son[now][0], son[now][1]);
92         ht[now] = ht[son[now][1]] + 1;
93         return now;
94     }
95
96     void insert(int &rt, int val, int id) { rt = merge(newnode(val, id), rt); }
97 }
98
99 int rt[M];
100
101 void build_heap()
102 {
103     for(int i = 1; i <= top; i ++)
104     {
105         int u = st[i];
106         rt[u] = rt[fa[u]];
107         for(int i = g.head[u]; i; i = g.edge[i].nex)
108         {
109             int v = g.edge[i].to;
110             if(!tree_edge[i] && dis[v] != inf) LT::insert(rt[u], dis[v] - dis[u] + g.
                edge[i].w, v);
111         }
112     }
113 }
114
115 int solve(int k)
116 {
117     if(k == 1) return dis[S];

```

```

118     __gnu_pbds::priority_queue<pair<int, int>, greater<pair<int, int> > > que;
119     que.push({dis[S] + LT::val[rt[S]], rt[S]});
120     while(!que.empty())
121     {
122         pair<int, int> f = que.top(); que.pop();
123         if((--k) == 1) return f.first;
124         int v = f.first, u = f.second;
125         int lc = LT::son[u][0], rc = LT::son[u][1], o = LT::id[u];
126         if(rt[o]) que.push({v + LT::val[rt[o]], rt[o]});
127         if(lc) que.push({v + LT::val[lc] - LT::val[u], lc});
128         if(rc) que.push({v + LT::val[rc] - LT::val[u], rc});
129     }
130     return -1;
131 }
132
133 void init()
134 {
135     g.init(n), rg.init(n);
136     memset(rt, 0, sizeof rt);
137     memset(tree_edge, 0, sizeof tree_edge);
138     top = LT::tot = 0;
139 }
140
141 void getans()
142 {
143     //input S-T
144     init();
145     dijkstra();
146     dfs(T);
147     build_heap();
148     cout << solve(k);
149 }

```

2.2 生成树

2.2.1 Matrix Tree

```

1  const int N = 305;
2  const int mod = 1e9 + 7;
3
4  int n, m, a[N][N];
5
6  int Gauss(int n) {
7      int ans = 1;
8      for (int i = 1; i <= n; i++) {
9          for (int k = i + 1; k <= n; k++) {
10             while (a[k][i]) {
11                 int d = a[i][i] / a[k][i];
12                 for (int j = i; j <= n; j++) {
13                     a[i][j] = (a[i][j] - 1LL * d * a[k][j] % mod + mod) % mod;
14                 }
15                 std::swap(a[i], a[k]);
16                 ans = -ans;
17             }
18         }
19         ans = 1LL * ans * a[i][i] % mod;
20     }
21     return (ans % mod + mod) % mod;
22 }

```

```

23
24 int main() {
25     scanf("%d%d", &n, &m);
26     for (int i = 1; i <= m; i++) {
27         int u, v;
28         scanf("%d%d", &u, &v);
29         a[u][v]--, a[v][u]--;
30         a[u][u]++, a[v][v]++;
31     }
32     printf("%d\n", Gauss(n - 1));
33     return 0;
34 }

```

2.2.2 Steiner Tree

```

1  /*BZOJ:4774
2  无向图G从1-n进行编号, 选择一些边, 使对于 $1 \leq i \leq d$ , i号点和n-i+1号点连通, 最小化选出的所有
   边权值和。
3  1. 枚举子树形态  $dp[S][i] = \min(dp[s] + dp[S \setminus s])$ 
4  2. 按照边进行松弛  $dp[S][i] = \min(dp[S][j] + w[j][i])$ 
5  其中 $S$ 为选取的子集,  $s$ 和 $S \setminus s$ 为 $S$ 的状态划分。第二类转移方程可以通过跑一次最短
   路进行松弛。
6  本题需要再做一次子集dp, 因为不成对的点可能不连通。
7  */
8  #include <bits/stdc++.h>
9
10 using namespace std;
11
12 const int maxn = 1e4 + 10;
13 const int inf = 0x3f3f3f3f;
14
15 int head[maxn], cnt;
16 struct Edge {int nex, to, w;} edge[maxn << 1];
17
18 void add(int u, int v, int w)
19 {
20     edge[cnt].nex = head[u];
21     edge[cnt].to = v;
22     edge[cnt].w = w;
23     head[u] = cnt++;
24 }
25
26 int f[1 << 10][maxn], ans[20];
27 bool in[maxn];
28
29 queue<int> que;
30
31 void spfa(int S)
32 {
33     while(!que.empty())
34     {
35         int u = que.front(); que.pop();
36         in[u] = false;
37         for(int i = head[u]; ~i; i = edge[i].nex)
38         {
39             int v = edge[i].to;
40             if(f[S][v] > f[S][u] + edge[i].w)
41                 f[S][v] = f[S][u] + edge[i].w;

```

```

42         f[S][v] = f[S][u] + edge[i].w;
43         if(!in[v]) que.push(v), in[v] = true;
44     }
45 }
46 }
47 }
48
49 int Steiner_Tree(int n, int d)
50 {
51     memset(f, 0x3f, sizeof f);
52     for(int i = 1; i <= d; i++)
53         f[1 << (i - 1)][i] = f[1 << (d + i - 1)][n - i + 1] = 0;
54     int lim = 1 << (d << 1);
55     for(int S = 1; S < lim; S++)
56     {
57         for(int i = 1; i <= n; i++)
58         {
59             for(int s = (S - 1) & S; s; s = (s - 1) & S)
60                 f[S][i] = min(f[S][i], f[s][i] + f[S ^ s][i]);
61             if(f[S][i] != inf) que.push(i), in[i] = true;
62         }
63         spfa(S);
64     }
65     lim = 1 << d;
66     memset(ans, 0x3f, sizeof ans);
67     for(int S = 1; S < lim; S++)
68         for(int i = 1; i <= n; i++)
69             ans[S] = min(ans[S], f[S ^ (S << d)][i]);
70     for(int S = 1; S < lim; S++)
71         for(int s = (S - 1) & S; s; s = (s - 1) & S)
72             ans[S] = min(ans[S], ans[s] + ans[S ^ s]);
73     return ans[lim - 1] == inf ? -1 : ans[lim - 1];
74 }
75
76 int main()
77 {
78     int n, m, d, u, v, w;
79     scanf("%d%d%d", &n, &m, &d);
80     memset(head, 0xff, sizeof head);
81     while(m--)
82     {
83         scanf("%d%d%d", &u, &v, &w);
84         add(u, v, w);
85         add(v, u, w);
86     }
87     printf("%d\n", Steiner_Tree(n, d));
88     return 0;
89 }

```

2.2.3 最小树形图

```

1  const int INF = 0x3f3f3f3f;
2  const int maxn = 10000;
3  const int maxm = 10000;
4
5  struct Edge{int u,v,cost; } edge[maxm];
6
7  int pre[maxn], id[maxn], vis[maxn], in[maxn];

```

```

8
9 int zhuliu(int root, int n, int m)
10 {
11     int res=0, u, v;
12     for(;;)
13     {
14         for(int i=0; i<n; i++) in[i] = INF;
15         for(int i=0; i<m; i++) if(edge[i].u != edge[i].v && edge[i].cost < in[edge[i].v
16             ])
17         {
18             pre[edge[i].v] = edge[i].u;
19             in[edge[i].v] = edge[i].cost;
20         }
21         for(int i=0; i<n; i++) if(i != root && in[i] ==INF) return -1;
22         int tn=0;
23         memset(id, 0xff, sizeof id);
24         memset(vis, 0xff, sizeof vis);
25         in[root] = 0;
26         for(int i=0; i<n;i++)
27         {
28             res += in[i];
29             v = i;
30             while( vis[v] != i && id[v] == -1 && v!= root) vis[v] = i, v = pre[v];
31             if(v != root && id[v] == -1)
32             {
33                 for(int u = pre[v]; u != v; u = pre[u]) id[u] = tn;
34                 id[v] = tn++;
35             }
36             if(tn == 0) break;
37             for(int i=0; i<n; i++) if(id[i] == -1) id[i] = tn++;
38             for(int i=0; i<m; )
39             {
40                 v = edge[i].v;
41                 edge[i].u = id[edge[i].u];
42                 edge[i].v = id[edge[i].v];
43                 if(edge[i].u != edge[i].v) edge[i++].cost -= in[v];
44                 else swap(edge[i], edge[--m]);
45             }
46             n = tn;
47             root = id[root];
48         }
49         return res;
50     }

```

2.3 连通性

2.3.1 割点

```

1 const int maxn = 1e4 + 10;
2
3 vector<int> edge[maxn];
4 int n, dfn[maxn], low[maxn], cnt = 0;
5 bool vis[maxn], cut[maxn];
6
7 void Tarjan(int u, int fa)
8 {
9     dfn[u] = low[u] = ++cnt;
10    vis[u] = true;

```



```

11     int children = 0;
12     for (int i = 0; i < edge[u].size(); i++)
13     {
14         int v = edge[u][i];
15         if (v != fa && vis[v])
16             low[u] = min(low[u], dfn[v]);
17         else if (!vis[v])
18         {
19             Tarjan(v, u);
20             children++;
21             low[u] = min(low[u], low[v]);
22             if (fa == -1 && children > 1) //若u是根节点且子节点数大于1
23                 cut[u] = true;        //u是割点
24             else if (fa != -1 && low[v] >= dfn[u]) //若u不是根节点且v不能访问到u的父
                节点
25                 cut[u] = true;        //u是割点
26         }
27     }
28 }

```

2.3.2 桥

```

1  const int maxn = 1e4 + 10;
2
3  vector<int> edge[maxn];
4  int n, dfn[maxn], low[maxn], father[maxn], cnt = 0;
5  bool bridge[maxn][maxn];
6
7  void Tarjan(int u, int fa)
8  {
9      dfn[u] = low[u] = ++cnt;
10     for (int i = 0; i < edge[u].size(); i++)
11     {
12         int v = edge[u][i];
13         if (!dfn[v]) //未访问节点v
14         {
15             Tarjan(v, u);
16             low[u] = min(low[u], low[v]);
17             if (low[v] > dfn[u]) //节点v到达祖先必须经过(u,v)
18                 bridge[u][v] = bridge[v][u] = true; // (u,v)是桥
19         }
20         else if (fa != v) //u的父节点不是v, (u,v)不存在重边
21             low[u] = min(low[u], dfn[v]);
22     }
23 }

```

2.3.3 强连通分量

```

1  const int maxn=1000+10;
2
3  vector<int> edge[maxn];
4
5  int dfn[maxn], low[maxn];
6  int stack[maxn], index, tot;
7  int belong[maxn], inde[maxn], outde[maxn], scc;
8  bool vis[maxn];
9
10 void add(int u, int v)

```

```
11 {
12     edge[u].push_back(v);
13     edge[v].push_back(u);
14 }
15
16 void Tarjan(int u)
17 {
18     dfn[u] = low[u] = ++tot;
19     stack[++index] = u;
20     vis[u] = true;
21     int v;
22     for(int i = 0; i < edge[u].size(); i++)
23     {
24         v = edge[u][i];
25         if(!dfn[v])
26         {
27             Tarjan(v);
28             low[u] = min(low[v], low[u]);
29         }
30         else if(vis[v]) low[u] = min(low[v], dfn[u]);
31     }
32     if(dfn[u] == low[u])
33     {
34         scc++;
35         do
36         {
37             v = stack[index--];
38             vis[v] = false;
39             belong[v] = scc;
40         } while(v != u);
41     }
42 }
```

2.3.4 点双联通分量

```
1  const int maxn = 10000 + 10;
2
3  struct Edge{ int u, v; };
4  vector<int> G[maxn], bcc[maxn];
5
6  int dfn[maxn], low[maxn], bccno[maxn], idx, bcc_cnt, bridge;
7  bool iscut[maxn];
8
9  stack<Edge> st;
10
11 void dfs(int u, int pre)
12 {
13     dfn[u] = low[u] = ++idx;
14     int child = 0;
15     for(auto v : G[u])
16     {
17         if(v == pre) continue;
18         if(!dfn[v])
19         {
20             child++;
21             st.push({u, v});
22             dfs(v, u);
23             low[u] = min(low[u], low[v]);
```

```

24         if(low[v] >= dfn[u])
25         {
26             iscut[u] = true;
27             bcc[++bcc_cnt].clear();
28             Edge x;
29             do
30             {
31                 x = st.top(); st.pop();
32                 if(bccno[x.u] != bcc_cnt) { bcc[bcc_cnt].push_back(x.u); bccno[x.u]
                    = bcc_cnt; }
33                 if(bccno[x.v] != bcc_cnt) { bcc[bcc_cnt].push_back(x.v); bccno[x.v]
                    = bcc_cnt; }
34             } while(x.u != u || x.v != v);
35         }
36         if(low[v] > dfn[u]) ++ bridge;
37     }
38     else if(dfn[v] < dfn[u])
39     {
40         st.push({u, v});
41         low[u] = min(low[u], dfn[v]);
42     }
43 }
44 if(pre < 0 && child == 1) iscut[u] = 0;
45 }
46
47 void find_bcc(int n)
48 {
49     memset(dfn, 0, sizeof dfn);
50     memset(iscut, 0, sizeof iscut);
51     memset(bccno, 0, sizeof bccno);
52     for(int i = 1; i <= bcc_cnt; i++) bcc[i].clear();
53     idx = bcc_cnt = bridge = 0;
54     for(int i = 0; i < n; i++) if(!dfn[i]) dfs(i, -1);
55 }

```

2.3.5 边双联通分量

```

1  const int maxn = 10000 + 10;
2
3  int low[maxn], dfn[maxn], head[maxn], cnt, idx;
4  int cutEdge[maxn << 2];
5  struct Edge { int nex, v; } edge[maxn << 2];
6
7  void add(int u, int v) { edge[cnt].nex = head[u], edge[cnt].v = v, head[u] = cnt++; }
8
9  void dfs(int u, int pre)
10 {
11     low[u] = dfn[u] = ++idx;
12     for(int i = head[u]; ~i; i = edge[i].nex)
13     {
14         int v = edge[i].v;
15         if(v == pre) continue;
16         if(!dfn[v])
17         {
18             dfs(v, u);
19             low[u] = min(low[u], low[v]);
20             if(low[v] > dfn[u]) cutEdge[i] = cutEdge[i ^ 1] = 1;
21         }

```

```

22         else if(dfn[v] < dfn[u]) low[u] = min(low[u], dfn[v]);
23     }
24 }
25
26 int n, m;
27 vector<int> edge[maxn];
28 int dfn[maxn], low[maxn], idx;
29 int st[maxn], stsz;
30 // 标号特性: 子节点标号大于父节点
31 int inWhichGroup[maxn], groupNow, groupRt[maxn];
32 int fa[maxn];
33
34 void dfs(int u, int fa) {
35     ::fa[u] = fa;
36     dfn[u] = low[u] = ++idx, st[++stsz] = u;
37     int firstToVisFa = 1;
38     for (auto v : edge[u]) {
39         if (v != fa || !firstToVisFa) {
40             if (!dfn[v]) {
41                 dfs(v, u);
42                 low[u] = min(low[u], low[v]);
43             } else low[u] = min(low[u], dfn[v]);
44             } else firstToVisFa = 0;
45     }
46     if (dfn[u] == low[u]) {
47         ++groupNow;
48         groupRt[groupNow] = u;
49         do {
50             inWhichGroup[st[stsz]] = groupNow;
51             } while (st[stsz--] != u);
52     }
53 }

```

2.4 图匹配

1. 二分图中的最大匹配数 = 最小点覆盖数
2. 最小路径覆盖 = 最小路径覆盖 = $|G|$ - 最大匹配数
3. 二分图最大独立集 = 顶点数 - 最小点覆盖
4. 二分图的最大团 = 补图的最大独立集

2.4.1 Hungary Algorithm

```

1  const int maxn = 150;
2
3  int n;
4  int edge[maxn][maxn];
5  int linker[maxn];
6  bool vis[maxn];
7
8  bool path(int u)
9  {
10     for (int v = 1; v <= n; v++)
11     {
12         if (edge[u][v] && !vis[v])
13         {
14             vis[v] = true;
15             if (linker[v] == -1 || path(linker[v]))

```

```
16         {
17             linker[v] = u;
18             return true;
19         }
20     }
21 }
22 return false;
23 }
24
25 int hungary()
26 {
27     int res = 0;
28     memset(linker, 0xff, sizeof(linker));
29     for (int i = 1; i <= n; i++)
30     {
31         memset(vis, false, sizeof(vis));
32         res += path(i);
33     }
34     return res;
35 }
```

2.4.2 Hopcroft-karp Algorithm

```
1 //复杂度O(n^0.5*m)
2 struct Hopcroft {
3     #define maxn 100005
4     #define maxm 100005
5     #define INF 0x3f3f3f3f
6     struct Edge { int v, next; } edge[maxm];
7     int nx, cnt, dis;
8     int first[maxn];
9     int xlink[maxn], ylink[maxn];
10    int dx[maxn], dy[maxn];
11    int vis[maxn];
12
13    void init(int n) {
14        cnt = 0;
15        for (int i = 0; i <= n; ++i) first[i] = ylink[i] = xlink[i] = -1;
16        nx = n + 1;
17    }
18
19    void add_edge(int u, int v) {
20        edge[cnt].v = v, edge[cnt].next = first[u], first[u] = cnt++;
21    }
22
23    int bfs() {
24        queue<int> q;
25        dis = INF;
26        for (int i = 0; i < nx; ++i) dx[i] = dy[i] = -1;
27        for (int i = 0; i < nx; i++) {
28            if (xlink[i] == -1) {
29                q.push(i);
30                dx[i] = 0;
31            }
32        }
33        while (!q.empty()) {
34            int u = q.front();
35            q.pop();
```

```

36         if (dx[u] > dis) break;
37         for (int e = first[u]; e != -1; e = edge[e].next) {
38             int v = edge[e].v;
39             if (dy[v] == -1) {
40                 dy[v] = dx[u] + 1;
41                 if (ylink[v] == -1) dis = dy[v];
42                 else {
43                     dx[ylink[v]] = dy[v] + 1;
44                     q.push(ylink[v]);
45                 }
46             }
47         }
48     }
49     return dis != INF;
50 }
51
52 int find(int u) {
53     for (int e = first[u]; e != -1; e = edge[e].next) {
54         int v = edge[e].v;
55         if (!vis[v] && dy[v] == dx[u] + 1) {
56             vis[v] = 1;
57             if (ylink[v] != -1 && dy[v] == dis) continue;
58             if (ylink[v] == -1 || find(ylink[v])) {
59                 xlink[u] = v, xlink[v] = u;
60                 return 1;
61             }
62         }
63     }
64     return 0;
65 }
66
67 int maxmatch() {
68     int ans = 0;
69     for (int i = 0; i < nx; ++i) vis[i] = 0;
70     while (bfs()) {
71         for (int i = 0; i < nx; ++i) vis[i] = 0;
72         for (int i = 0; i < nx; i++)
73             if (xlink[i] == -1)
74                 ans += find(i);
75     }
76     return ans;
77 }
78
79 #undef maxn
80 #undef maxm
81 } solve;

```

2.4.3 二分图多重匹配

```

1  const int maxn = 1e2 + 5; // 左边最大点数
2  const int maxm = 1e2 + 5; // 右边最大点数
3  int graph[maxn][maxm], vis[maxm]; // 图G和增广路访问标记
4  int match[maxm][maxn]; // 左边元素与右边元素第n次匹配
5  int nx, ny, m; // 左边点数, 右边点数, 边数
6  int vol[maxm]; // 右边点多重匹配可容纳值
7  int cnt[maxm]; // 右边点已匹配值
8
9  bool find_path(int u) // 找增广路

```

```

10 {
11     for (int i = 0; i < ny; i++)//注意, 这里节点是从0开始编号, 题目有时是从1开始编号
12     {
13         if (graph[u][i] && !vis[i])//不在增广路
14         {
15             vis[i] = 1;//放进增广路
16             if (cnt[i] < vol[i])//如果当前已匹配数量小于可容纳量, 则直接匹配
17             {
18                 match[i][cnt[i]++] = u;
19                 return true;
20             }
21             for (int j = 0; j < cnt[i]; j++)
22             {
23                 if (find_path(match[i][j]))//如果先前已匹配右边的点能另外找到增广路, 则
24                     //此点仍可匹配
25                 {
26                     match[i][j] = u;
27                     return true;
28                 }
29             }
30         }
31     }
32     return false;
33 }
34 int max_match()//计算多重匹配的最大匹配数
35 {
36     int res = 0;
37     memset(match, -1, sizeof(match));
38     memset(cnt, 0, sizeof(cnt));
39     for (int i = 0; i < nx; i++)
40     {
41         memset(vis, 0, sizeof(vis));
42         if (find_path(i)) res++;
43     }
44     return res;
45 }
46
47 bool all_match()//判断左边的点是否都与右边的点匹配了
48 {
49     memset(cnt, 0, sizeof(cnt));
50     for (int i = 0; i < nx; i++)
51     {
52         memset(vis, 0, sizeof(vis));
53         if (!find_path(i)) return false;
54     }
55     return true;
56 }

```

2.4.4 二分图最大权匹配 (KM 算法)

```

1  int n;
2  int cost[maxn][maxn];
3  int lx[maxn], ly[maxn], match[maxn], slack[maxn];
4  int Prev[maxn];
5  bool vy[maxn];
6  const int INF = 0x3f3f3f3f;
7

```

```

8 void augment(int root) {
9     fill(vy + 1, vy + n + 1, false);
10    fill(slack + 1, slack + n + 1, INF);
11    int py;
12    match[py = 0] = root;
13    do {
14        vy[py] = true;
15        int x = match[py], yy;
16        int delta = INF;
17        for (int y = 1; y <= n; y++) {
18            if (!vy[y]) {
19                if (lx[x] + ly[y] - cost[x][y] < slack[y])
20                    slack[y] = lx[x] + ly[y] - cost[x][y], Prev[y] = py;
21                if (slack[y] < delta) delta = slack[y], yy = y;
22            }
23        }
24        for (int y = 0; y <= n; y++) {
25            if (vy[y])
26                lx[match[y]] -= delta, ly[y] += delta;
27            else
28                slack[y] -= delta;
29        }
30        py = yy;
31    } while (match[py] != -1);
32    do {
33        int pre = Prev[py];
34        match[py] = match[pre], py = pre;
35    } while (py);
36 }
37
38 int KM() {
39     for (int i = 1; i <= n; i++) {
40         lx[i] = ly[i] = 0;
41         match[i] = -1;
42         for (int j = 1; j <= n; j++) lx[i] = max(lx[i], cost[i][j]);
43     }
44     int answer = 0;
45     for (int root = 1; root <= n; root++) augment(root);
46     for (int i = 1; i <= n; i++) answer += lx[i], answer += ly[i];
47     return answer;
48 }

```

2.4.5 一般图匹配带花树

```

1 //一般图匹配, 带花树算法
2 const int maxn = 1000 + 10;
3
4 struct Edmond
5 {
6     vector<int> edge[maxn];
7     queue<int> que;
8
9     int n, pre[maxn], type[maxn], link[maxn], nex[maxn], vis[maxn];
10
11     void init(int n)
12     {
13         this->n = n;
14         for (int i = 0; i <= n; i++) edge[i].clear();

```



```
15     memset(link, 0, sizeof(link));
16 }
17
18 void add(int u, int v)
19 {
20     edge[u].push_back(v);
21     edge[v].push_back(u);
22 }
23
24 int Find(int x)
25 {
26     return x == pre[x] ? x : pre[x] = Find(pre[x]);
27 }
28
29 void combine(int x, int lca)    //如果找到奇环, 对当前点x和找到的
30 {
31     while (x != lca)
32     {
33         int u = link[x], v = nex[u];
34         if (Find(v) != lca) nex[v] = u;
35         if (type[u] == 1) type[u] = 2, que.push(u);
36         pre[Find(x)] = Find(u);
37         pre[Find(u)] = Find(v);
38         x = v;
39     }
40 }
41
42 void contrack(int x, int y)
43 {
44     int lca = x;
45     memset(vis, 0, sizeof(vis));
46     for (int i = x; i; i = nex[link[i]])
47     {
48         i = Find(i);
49         vis[i] = 1;
50     }
51     for (int i = y; i; i = nex[link[i]])
52     {
53         i = Find(i);
54         if (vis[i])
55         {
56             lca = i;
57             break;
58         }
59     }
60     if (lca != Find(x)) nex[x] = y;
61     if (lca != Find(y)) nex[y] = x;
62     combine(x, lca);
63     combine(y, lca);
64 }
65
66 void bfs(int s)
67 {
68     memset(type, 0, sizeof(type));
69     memset(nex, 0, sizeof(nex));
70     for (int i = 1; i <= n; i++) pre[i] = i;
71     while (!que.empty()) que.pop();
72     que.push(s);
73     type[s] = 2;
```

```

74     while (!que.empty())
75     {
76         int x = que.front();
77         que.pop();
78         for (int i = 0; i < edge[x].size(); i++)
79         {
80             int y = edge[x][i];
81             if (Find(x) == Find(y) || link[x] == y || type[y] == 1) continue;
82             if (type[y] == 2) contrack(x, y);
83             else if (link[y])
84             {
85                 nex[y] = x;
86                 type[y] = 1;
87                 type[link[y]] = 2;
88                 que.push(link[y]);
89             } else
90             {
91                 nex[y] = x;
92                 int pos = y, u = nex[pos], v = link[u];
93                 while (pos)
94                 {
95                     link[pos] = u;
96                     link[u] = pos;
97                     pos = v;
98                     u = nex[pos];
99                     v = link[u];
100                 }
101                 return;
102             }
103         }
104     }
105 }
106
107 int maxmatch()
108 {
109     for (int i = 1; i <= n; i++) if (!link[i]) bfs(i);
110     int ans = 0;
111     for (int i = 1; i <= n; i++) if (link[i]) ans++;
112     return ans / 2;
113 }
114 }ans;

```

2.5 网络流

2.5.1 Dinic

```

1  const int MAX_V = 1000 + 10;
2  const int INF = 0x3f3f3f3f;
3
4  //用于表示边的结构体 (终点, 流量, 反向边)
5  struct edge{int to, cap, rev;};
6
7  vector<edge> G[MAX_V]; //图的邻接表表示
8  int level[MAX_V]; //顶点到源点的距离标号
9  int iter[MAX_V]; //当前弧
10
11 void add(int from, int to, int cap)
12 {
13     G[from].push_back((edge){to, cap, (int)G[to].size()});

```

```
14     G[to].push_back((edge){from, 0, (int)G[from].size() - 1});
15 }
16
17 // 计算从源点出发的距离标号
18 void bfs(int s)
19 {
20     memset(level, -1, sizeof(level));
21     queue<int> que;
22     level[s] = 0;
23     que.push(s);
24     while(!que.empty())
25     {
26         int v = que.front(); que.pop();
27         for(int i = 0; i < G[v].size(); i++)
28         {
29             edge &e = G[v][i];
30             if(e.cap > 0 && level[e.to] < 0)
31             {
32                 level[e.to] = level[v] + 1;
33                 que.push(e.to);
34             }
35         }
36     }
37 }
38
39 // 通过DFS寻找增广路
40 int dfs(int v, int t, int f)
41 {
42     if(v == t) return f;
43     for(int &i = iter[v]; i < G[v].size(); i++)
44     {
45         edge &e = G[v][i];
46         if(e.cap > 0 && level[v] < level[e.to])
47         {
48             int d = dfs(e.to, t, min(f, e.cap));
49             if(d > 0)
50             {
51                 e.cap -= d;
52                 G[e.to][e.rev].cap += d;
53                 return d;
54             }
55         }
56     }
57     return 0;
58 }
59
60 // 求解从s到t的最大流
61 int max_flow(int s, int t)
62 {
63     int flow = 0;
64     for(;;)
65     {
66         bfs(s);
67         if(level[t] < 0) return flow;
68         memset(iter, 0, sizeof(iter));
69         int f;
70         while((f = dfs(s,t,INF)) > 0) flow += f;
71     }
72 }
```

2.5.2 ISAP

```
1 struct Edge {
2     int from, to, cap, flow;
3     Edge(int u, int v, int c, int f) : from(u), to(v), cap(c), flow(f) {}
4 };
5
6 bool operator<(const Edge& a, const Edge& b) {
7     return a.from < b.from || (a.from == b.from && a.to < b.to);
8 }
9
10 struct ISAP {
11     int n, m, s, t;
12     vector<Edge> edges;
13     vector<int> G[maxn];
14     bool vis[maxn];
15     int d[maxn];
16     int cur[maxn];
17     int p[maxn];
18     int num[maxn];
19
20     void AddEdge(int from, int to, int cap) {
21         edges.push_back(Edge(from, to, cap, 0));
22         edges.push_back(Edge(to, from, 0, 0));
23         m = edges.size();
24         G[from].push_back(m - 2);
25         G[to].push_back(m - 1);
26     }
27
28     bool BFS() {
29         memset(vis, 0, sizeof(vis));
30         queue<int> Q;
31         Q.push(t);
32         vis[t] = 1;
33         d[t] = 0;
34         while (!Q.empty()) {
35             int x = Q.front();
36             Q.pop();
37             for (int i = 0; i < G[x].size(); i++) {
38                 Edge& e = edges[G[x][i] ^ 1];
39                 if (!vis[e.from] && e.cap > e.flow) {
40                     vis[e.from] = 1;
41                     d[e.from] = d[x] + 1;
42                     Q.push(e.from);
43                 }
44             }
45         }
46         return vis[s];
47     }
48
49     void init(int n) {
50         this->n = n;
51         for (int i = 0; i < n; i++) G[i].clear();
52         edges.clear();
53     }
54
55     int Augment() {
56         int x = t, a = INF;
57         while (x != s) {
```

```

58     Edge& e = edges[p[x]];
59     a = min(a, e.cap - e.flow);
60     x = edges[p[x]].from;
61 }
62 x = t;
63 while (x != s) {
64     edges[p[x]].flow += a;
65     edges[p[x] ^ 1].flow -= a;
66     x = edges[p[x]].from;
67 }
68 return a;
69 }
70
71 int Maxflow(int s, int t) {
72     this->s = s;
73     this->t = t;
74     int flow = 0;
75     BFS();
76     memset(num, 0, sizeof(num));
77     for (int i = 0; i < n; i++) num[d[i]]++;
78     int x = s;
79     memset(cur, 0, sizeof(cur));
80     while (d[s] < n) {
81         if (x == t) {
82             flow += Augment();
83             x = s;
84         }
85         int ok = 0;
86         for (int i = cur[x]; i < G[x].size(); i++) {
87             Edge& e = edges[G[x][i]];
88             if (e.cap > e.flow && d[x] == d[e.to] + 1) {
89                 ok = 1;
90                 p[e.to] = G[x][i];
91                 cur[x] = i;
92                 x = e.to;
93                 break;
94             }
95         }
96         if (!ok) {
97             int m = n - 1;
98             for (int i = 0; i < G[x].size(); i++) {
99                 Edge& e = edges[G[x][i]];
100                 if (e.cap > e.flow) m = min(m, d[e.to]);
101             }
102             if (--num[d[x]] == 0) break;
103             num[d[x] = m + 1]++;
104             cur[x] = 0;
105             if (x != s) x = edges[p[x]].from;
106         }
107     }
108     return flow;
109 }
110 };

```

2.5.3 MCMF

```

1 const int maxn = 10000 + 10;
2 const int inf = 0x3f3f3f3f;

```

```

3
4 struct Edge { int from, to, cap, flow, cost; };
5
6 struct MCMF
7 {
8     int n, m;
9     vector<Edge> edges;
10    vector<int> G[maxn];
11    bool inq[maxn];
12    int dis[maxn], path[maxn], a[maxn];
13
14    void init(int n)
15    {
16        this->n = n;
17        for(int i = 0; i <= n; i++)
18            G[i].clear();
19        edges.clear();
20    }
21
22    void addEdge(int from, int to, int cap, int cost)
23    {
24        edges.push_back(Edge{from, to, cap, 0, cost});
25        edges.push_back(Edge{to, from, 0, 0, -cost});
26        m = edges.size();
27        G[from].push_back(m - 2);
28        G[to].push_back(m - 1);
29    }
30
31    bool Bellman_Ford(int s, int t, int& flow, int& cost)
32    {
33        for(int i = 0; i <= n; i++) dis[i] = inf;
34        memset(inq, 0, sizeof inq);
35        dis[s] = 0, inq[s] = true, path[s] = 0, a[s] = inf;
36        queue<int> Q;
37        Q.push(s);
38        while(!Q.empty())
39        {
40            int u = Q.front(); Q.pop();
41            inq[u] = false;
42            for(int i = 0; i < G[u].size(); i++)
43            {
44                Edge& e = edges[G[u][i]];
45                if(e.cap > e.flow && dis[e.to] > dis[u] + e.cost)
46                {
47                    dis[e.to] = dis[u] + e.cost;
48                    path[e.to] = G[u][i];
49                    a[e.to] = min(a[u], e.cap - e.flow);
50                    if(!inq[e.to])
51                    {
52                        Q.push(e.to);
53                        inq[e.to] = true;
54                    }
55                }
56            }
57        }
58        if(dis[t] == inf) return false;    //求最小费用最大流
59        //if(1ll * dis[t] * a[t] > 0) return false; 求可行流最小费用, 因此当费用增量大于0时不继续增加流量
60        flow += a[t];

```

```

61     cost += dis[t] * a[t];
62     for(int u = t; u != s; u = edges[path[u]].from)
63     {
64         edges[path[u]].flow += a[t];
65         edges[path[u] ^ 1].flow -= a[t];
66     }
67     return true;
68 }
69
70 int mincostMaxFlow(int s, int t)
71 {
72     int flow = 0, cost = 0;
73     while(Bellman_Ford(s, t, flow, cost));
74     return cost;
75 }
76 };

```

2.5.4 Trick

建模技巧

二分图带权最大独立集。给出一个二分图，每个结点上有一个正权值。要求选出一些点，使得这些点之间没有边相连，且权值和最大。

解：在二分图的基础上添加源点 S 和汇点 T ，然后从 S 向所有 X 集合中的点连一条边，所有 Y 集合中的点向 T 连一条边，容量均为该点的权值。 X 结点与 Y 结点之间的边的容量均为无穷大。这样，对于图中的任意一个割，将割中的边对应的结点删掉就是一个符合要求的解，权和为所有权减去割的容量。因此，只需要求出最小割，就能求出最大权和。

公平分配问题。把 m 个任务分配给 n 个处理器。其中每个任务有两个候选处理器，可以任选一个分配。要求所有处理器中，任务数最多的那个处理器所分配的任务数尽量少。不同任务的候选处理器集 $\{p_1, p_2\}$ 保证不同。

解：本题有一个比较明显的二分图模型，即 X 结点是任务， Y 结点是处理器。二分答案 x ，然后构图，首先从源点 S 出发向所有的任务结点引一条边，容量等于 1，然后从每个任务结点出发引两条边，分别到达它所能分配到的两个处理器结点，容量为 1，最后从每个处理器结点出发引一条边到汇点 T ，容量为 x ，表示选择该处理器的任务不能超过 x 。这样网络中的每个单位流量都是从 S 流到一个任务结点，再到处理器结点，最后到汇点 T 。只有当网络中的总流量等于 m 时才意味着所有任务都选择了一个处理器。这样，我们通过 $O(\log m)$ 次最大流便算出了答案。

区间 k 覆盖问题。数轴上有一些带权值的左闭右开区间。选出权和尽量大的一些区间，使得任意一个数最多被 k 个区间覆盖。

解：本题可以用最小费用流解决，构图方法是把每个数作为一个结点，然后对于权值为 w 的区间 $[u, v)$ 加边 $u \rightarrow v$ ，容量为 1，费用为 $-w$ 。再对所有相邻的点加边 $i \rightarrow i+1$ ，容量为 k ，费用为 0。最后，求最左点到最右点的最小费用最大流即可，其中每个流量对应一组互不相交的区间。如果数值范围太大，可以先进行离散化。

最大闭合子图。给定带权图 G （权值可正可负），求一个权和最大的点集，使得起点在该点集中的任意弧，终点也在该点集中。

解：新增附加源 s 和附加汇 t ，从 s 向所有正权点引一条边，容量为权值；从所有负权点向汇点引一条边，容量为权值的相反数。求出最小割以后， $S - \{s\}$ 就是最大闭合子图。

最大密度子图。给出一个无向图，找一个点集，使得这些点之间的边数除以点数的值（称为子图的密度）最大。

解：如果两个端点都选了，就必然要选边，这就是一种推导。如果把每个点和每条边都看成新图中的结点，可以把问题转化为最大闭合子图。

无源汇有上下界可行流：附加源 S 和汇 T ；对于边 (u, v, min, max) ，记 $d[u]- = min, d[v]+ = max$ ，并添加弧 $(u, v, max - min)$ ；对于流量不平衡的点 u ，设多余流量为 W ，如果 $W > 0$ ，添加弧 $S \rightarrow u : W$ ，否则若 $W < 0$ ，添加弧 $u \rightarrow T : -W$ ，求改造后的网络 $S - T$ 最大流即可，当且仅当所有附加弧满载时原图有可行流。

有源汇有上下界可行流: 建 $t \rightarrow s$, 容量为 inf , 然后和无源汇相同。

有源汇有上下界最大/最小流: 与上面相同, 跑完可行流 $S \rightarrow T$ 后去掉边 $t \rightarrow s$, 最大流为加 $s \rightarrow t$, 最小流为 $G[s][t].\text{cap} - \text{maxflow}(t, s)$ 。

2.5.5 Stoer Wagner

```

1  #define INF 100000000
2  bool vis[maxn], com[maxn];
3  int mp[maxn][maxn], w[maxn], s, t;
4
5  int maxadj(int n, int v) {
6      int CUT = 0;
7      memset(vis, 0, sizeof vis);
8      memset(w, 0, sizeof w);
9      for (int i = 0; i < n; ++i) {
10         int num = 0, mx = -INF;
11         for (int j = 0; j < v; ++j) {
12             if (!com[j] && !vis[j] && w[j] > mx) {
13                 mx = w[j];
14                 num = j;
15             }
16         }
17         vis[num] = 1;
18         s = t;
19         t = num;
20         CUT = w[t];
21         for (int j = 0; j < v; ++j) {
22             if (!com[j] && !vis[j]) w[j] += mp[num][j];
23         }
24     }
25     return CUT;
26 }
27
28 int stoer(int v) {
29     int mincut = INF;
30     int n = v;
31     memset(com, 0, sizeof com);
32     for (int i = 0; i < v - 1; ++i) {
33         int cut;
34         s = 0, t = 0;
35         cut = maxadj(n, v);
36         n--;
37         if (cut < mincut) mincut = cut;
38         com[t] = 1;
39         for (int j = 0; j < v; ++j) {
40             if (!com[j]) {
41                 mp[j][s] += mp[j][t];
42                 mp[s][j] += mp[t][j];
43             }
44         }
45     }
46     return mincut;
47 }

```

2.5.6 ZKW 费用流

```

1  struct MCMF

```



```

2 {
3     int last[maxn], dis[maxn], cnt, ans;
4     int s, t;
5     bool vis[maxn];
6     struct edge { int from, to, cap, w, op, nex; } e[500000 + 10];
7
8     void init(int S, int T)
9     {
10         s = S, t = T;
11         cnt = 0, ans = 0;
12         memset(vis, 0, sizeof vis);
13         memset(dis, 0, sizeof dis);
14         memset(last, 0, sizeof last);
15     }
16
17     void add(int u, int v, int cap, int cost)
18     {
19         e[++ cnt] = { u, v, cap, cost, cnt + 1, last[u] };
20         last[u] = cnt;
21         e[++ cnt] = { v, u, 0, -cost, cnt - 1, last[v] };
22         last[v] = cnt;
23     }
24
25     int dfs(int x, int maxf)
26     {
27         if(x == t || maxf == 0) return maxf;
28         int ret = 0;
29         vis[x] = 1;
30         for(int i = last[x]; i; i = e[i].nex)
31             if(e[i].cap && dis[e[i].to] + e[i].w == dis[x] && !vis[e[i].to])
32             {
33                 int f = dfs(e[i].to, min(e[i].cap, maxf - ret));
34                 ans += f * e[i].w;
35                 e[i].cap -= f;
36                 e[e[i].op].cap += f;
37                 ret += f;
38                 if(ret == maxf) break;
39             }
40         return ret;
41     }
42
43     bool change()
44     {
45         int mn = inf;
46         for(int i = 0; i <= t; i ++)
47             if(vis[i])
48                 for(int j = last[i]; j; j = e[j].nex)
49                     if(!vis[e[j].to] && e[j].cap) mn = min(mn, -dis[i] + e[j].w + dis[e[j].to]);
50         if(mn == inf) return false;
51         for(int i = 0; i <= t; i ++) if(vis[i]) dis[i] += mn;
52         return true;
53     }
54
55     void zkw()
56     {
57         do
58         {
59             for(int i = 0; i <= t; i ++) vis[i] = 0;

```

```

60         while(dfs(s, inf)) for(int i = 0; i <= t; i++) vis[i] = 0;
61     }
62     while(change());
63 }
64 }ans;

```

2.6 Others

2.6.1 拓扑排序

```

1  const int maxn = 1e5 + 10;
2
3  vector<int> edge[maxn];
4  int indegree[maxn];
5
6  void add(int u, int v)
7  {
8      edge[u].push_back(v);
9      indegree[v]++;
10 }
11
12 void Toposort(int n)
13 {
14     queue<int> que;
15     for (int i = 1; i <= n; i++)
16         if (!indegree[i]) que.push(i);    //将图中没有前驱, 即入度为0的点加入队列
17     while (!que.empty())
18     {
19         int u = que.front();
20         que.pop();
21         indegree[u] = -1;    //从图中删去此顶点
22         for (int i = 0; i < edge[u].size(); i++)
23         {
24             int v = edge[u][i];
25             indegree[v]--;    //删去图中以u为尾的弧
26             if (!indegree[v]) que.push(v);    //将新增的当前入度为0的点压入队列中
27         }
28     }
29 }

```

2.6.2 2-SAT

```

1  /*2-SAT连边含义: 选A必选B
2      点$x_i$表示选, $x_i'$表示不选
3      1. 必选$x_i$, 等价于$x_i=1$: $x_i \rightarrow x_i$
4      2. 必不选$x_i$, 等价于$x_i=0$, $x_i \rightarrow x_i'$
5      3. $x_i$与$x_j$中至少选择一个, 等价于$x_i \vee x_j=1$, 连边$x_i \rightarrow x_j$, $x_j \rightarrow x_i$
6      4. $x_i$与$x_j$不都选, 等价于$x_i \wedge x_j=0$, 连边$x_i \rightarrow x_j'$, $x_j \rightarrow x_i'$
7      5. $x_i$与$x_j$情况相同, 等价于$x_i \oplus x_j=0$, 连边$x_i \rightarrow x_j$, $x_i' \rightarrow x_j'$, $x_j \rightarrow x_i$,
8          $x_j' \rightarrow x_i'$
9      6. $x_i$与$x_j$情况相反, 等价于$x_i \oplus x_j=1$, 连边$x_i \rightarrow x_j'$, $x_i' \rightarrow x_j$, $x_j \rightarrow x_i'$,
9          $x_j' \rightarrow x_i$
10 */
11 const int maxn = 2e6 + 10;
12
13 namespace twosat {
14     int n;

```

```
15     int low[maxn], dfn[maxn], color[maxn], cnt, scc_cnt;
16     bool instack[maxn];
17
18     vector<int> g[maxn];
19     stack<int> st;
20
21     void init(int _n) {
22         n = _n;
23         cnt = scc_cnt = 0;
24         for (int i = 0; i <= n * 2; ++i) {
25             dfn[i] = 0;
26             g[i].clear();
27         }
28     }
29
30     void Tarjan(int u) {
31         low[u] = dfn[u] = ++cnt;
32         st.push(u);
33         instack[u] = true;
34         for (const auto &v : g[u]) {
35             if (!dfn[v]) Tarjan(v), low[u] = min(low[u], low[v]);
36             else if (instack[v]) low[u] = min(low[u], dfn[v]);
37         }
38         if (low[u] == dfn[u]) {
39             ++scc_cnt;
40             do {
41                 color[u] = scc_cnt;
42                 u = st.top();
43                 st.pop();
44                 instack[u] = false;
45             } while (low[u] != dfn[u]);
46         }
47     }
48
49     inline void add(int a, int b) { g[a].push_back(b); }
50
51     inline void AND(int a, int b, int c) {
52         if (c == 1) add(a, a + n), add(b, b + n);
53         else add(a, b + n), add(b, a + n);
54     }
55
56     inline void OR(int a, int b, int c) {
57         if (c == 0) add(a + n, a), add(b + n, b);
58         else add(a + n, b), add(b + n, a);
59     }
60
61     inline void XOR(int a, int b, int c) {
62         if (c == 0) add(a, b), add(a + n, b + n), add(b, a), add(b + n, a + n);
63         else add(a, b + n), add(a + n, b), add(b, a + n), add(b + n, a);
64     }
65
66     bool TWO_SAT() {
67         for (int i = 1; i <= (n << 1); i++) if (!dfn[i]) Tarjan(i);
68         for (int i = 1; i <= n; i++)
69             if (color[i] == color[i + n]) return false;
70         for (int i = 1; i <= n; i++)
71             printf("%d ", color[i] > color[i + n]);
72         return true;
73     }
```

74 }

2.6.3 差分约束系统

```
1 //以$x_i-x_j y$为约束条件, 建图求最短路后得到的是最大解。所有的解都不大于且尽可能逼近
   $dis[x0]$
2 //最短路对应最大解, 最长路对应最小解
3
4 const int maxn = 1000 + 10;
5 const int inf = 0x3f3f3f3f;
6
7 struct Edge
8 {
9     int nex, to, w;
10 } edge[10 * maxn];
11
12 int head[maxn], cnt, dis[maxn], n;
13 bool vis[maxn];
14
15 void init()
16 {
17     cnt = 0;
18     memset(head, 0xff, sizeof head);
19 }
20
21 void add(int u, int v, int w)
22 {
23     edge[cnt].nex = head[u];
24     edge[cnt].to = v;
25     edge[cnt].w = w;
26     head[u] = ++cnt;
27 }
28
29 void spfa(int u)
30 {
31     int u, v, w;
32     for (int i = 1; i <= n; i++) dis[i] = inf, vis[i] = false;
33     dis[u] = 0;
34     queue<int> que;
35     que.push(u);
36     vis[u] = true;
37     while (!que.empty())
38     {
39         u = que.front();
40         que.pop();
41         vis[u] = false;
42         for (int i = head[u]; ~i; i = edge[i].nex)
43         {
44             v = edge[i].v, w = edge[i].w;
45             if (dis[u] + w < dis[v])
46             {
47                 dis[v] = dis[u] + w;
48                 if (!vis[v])
49                 {
50                     que.push(v);
51                     vis[v] = true;
52                 }
53             }
54         }
55     }
56 }
```

```
54     }
55 }
56 }
```

2.6.4 支配树

```
1  const int N = 2e5 + 10;
2
3  int n, m;
4
5  struct G
6  {
7      vector<int> edge[N];
8      inline void add(int u, int v) { edge[u].push_back(v); }
9  }a, b, c, d;
10
11 int dfn[N], id[N], fa[N], cnt;
12
13 void dfs(int u)
14 {
15     dfn[u] = ++ cnt; id[cnt] = u;
16     int len = a.edge[u].size();
17     for(auto v : a.edge[u]) if(!dfn[v]) { fa[v] = u; dfs(v); }
18 }
19
20 int semi[N], idom[N], belong[N], val[N];
21
22 int find(int x)
23 {
24     if(x == belong[x]) return x;
25     int tmp = find(belong[x]);
26     if(dfn[semi[val[belong[x]]]] < dfn[semi[val[x]]]) val[x] = val[belong[x]];
27     return belong[x] = tmp;
28 }
29
30 void tarjan()
31 {
32     for(int i = cnt; i > 1; i --)
33     {
34         int u = id[i];
35         for(auto v : b.edge[u])
36         {
37             if(!dfn[v]) continue;
38             find(v);
39             if(dfn[semi[val[v]]] < dfn[semi[u]]) semi[u] = semi[val[v]];
40         }
41         c.add(semi[u], u);
42         belong[u] = fa[u];
43         u = fa[u];
44         for(auto v : c.edge[u])
45         {
46             find(v);
47             if(semi[val[v]] == u) idom[v] = u;
48             else idom[v] = val[v];
49         }
50     }
51     for(int i = 2; i <= cnt; i ++){
52 }
```

```
53         int u = id[i];
54         if(idom[u] != semi[u]) idom[u] = idom[idom[u]];
55     }
56 }
57
58 int ans[N];
59
60 void dfs_ans(int u)
61 {
62     ans[u] = 1;
63     for(auto v : d.edge[u]) dfs_ans(v), ans[u] += ans[v];
64 }
65
66 void solve()
67 {
68     int u, v;
69     scanf("%d%d", &n, &m);
70     while(m --)
71     {
72         scanf("%d%d", &u, &v);
73         a.add(u, v);
74         b.add(v, u);
75     }
76     for(int i = 1; i <= n; i++) semi[i] = belong[i] = val[i] = i;
77     dfs(1);
78     tarjan();
79     for(int i = 2; i <= n; i++) d.add(idom[i], i);
80     dfs_ans(1);
81     for(int i = 1; i <= n; i++) printf("%d ", ans[i]);
82 }
```

2.6.5 Stable Matching Problem

```
1  const int maxn = 1000 + 10;
2
3  int pre[maxn][maxn], order[maxn][maxn], nex[maxn];
4  int hus[maxn], wife[maxn];
5  queue<int> que;
6
7  void engage(int man, int woman) {
8      int m = hus[woman];
9      if (m) wife[m] = 0, q.push(m);
10     wife[man] = woman;
11     hus[woman] = man;
12 }
13
14 int solve() {
15     for (int i = 1; i <= n; i++) {
16         for (int j = 1; j <= n; j++)
17             scanf("%d", &pre[i][j]);
18         nex[i] = 1;
19         wife[i] = 0;
20         que.push(i);
21     }
22     for (int i = 1; i <= n; i++) {
23         for (int j = 1; j <= n; j++) {
24             int x;
25             scanf("%d", &x);
```

```

26         order[i][x] = j;
27     }
28     hus[i] = 0;
29 }
30
31 while (!que.empty()) {
32     int man = que.front();
33     que.pop();
34     int woman = pre[man][nex[man]++];
35     if (!hus[woman]) engage(man, woman);
36     else if (order[woman][man] < order[woman][hus[woman]]) engage(man, woman);
37     else que.push(man);
38 }
39 }

```

2.6.6 一般图最大团

```

1  #define u64 unsigned long long
2  #define i64 long long
3
4  const u64 BITCOUNT = sizeof(u64) * 8;
5
6  u64 count_trailing_zeroes(u64 a) {
7      if (a == 0ull)
8          return BITCOUNT;
9      return __builtin_ctzll(a);
10 }
11 u64 disable_bit(u64 a, u64 bit) { return a & ~(1ull << bit); }
12 u64 popcount(u64 a) { return __builtin_popcountll(a); }
13
14 map<u64, u64> max_clique_cache;
15
16 u64 max_clique(u64 mask, vector<u64> const &graph_matrix) { // 最大独立集传补图即可
17     if (max_clique_cache.find(mask) != max_clique_cache.end())
18         return max_clique_cache[mask];
19     u64 a = count_trailing_zeroes(mask);
20     if (a == BITCOUNT) return 0;
21     u64 res1 = max_clique(disable_bit(mask, a), graph_matrix);
22     u64 res2 = max_clique(mask & disable_bit(graph_matrix[a], a), graph_matrix) | (1ull
        << a);
23     u64 res = popcount(res1) > popcount(res2) ? res1 : res2;
24     max_clique_cache[mask] = res;
25     return res;
26 }
27
28 int main() {
29     vector<u64> M;
30     int n;
31     while (scanf("%d", &n) != EOF) {
32         if (n == 0) break;
33         M.clear();
34         M.resize(n);
35         max_clique_cache.clear();
36         for (int i = 0; i < n; ++i) {
37             for (int j = 0; j < n; ++j) {
38                 int x; scanf("%d", &x);
39                 if (x == 1) M[i] |= 1ll << j;
40             }

```

```
41     }
42     printf("%d\n", popcount(max_clique((1ll << n) - 1, M)));
43 }
44 }
```


3 DataStructure

3.1 SegmentTreeDS

3.1.1 SGTB

```
1  int n, a[maxn];
2
3  struct node {
4      int p, t, se;
5
6      node() {}
7
8      node(int p, int t, int se) : p(p), t(t), se(se) {}
9
10     inline friend node combineMax(node a, node b) {
11         node c;
12         if (a.p < b.p) {
13             c.p = b.p;
14             c.t = b.t;
15             c.se = max(a.p, b.se);
16         } else if (a.p > b.p) {
17             c.p = a.p;
18             c.t = a.t;
19             c.se = max(a.se, b.p);
20         } else {
21             c.p = a.p;
22             c.t = a.t + b.t;
23             c.se = max(a.se, b.se);
24         }
25         return c;
26     }
27
28     inline friend node combineMin(node a, node b) {
29         node c;
30         if (a.p > b.p) {
31             c.p = b.p;
32             c.t = b.t;
33             c.se = min(a.p, b.se);
34         } else if (a.p < b.p) {
35             c.p = a.p;
36             c.t = a.t;
37             c.se = min(a.se, b.p);
38         } else {
39             c.p = a.p;
40             c.t = a.t + b.t;
41             c.se = min(a.se, b.se);
42         }
43         return c;
44     }
45
46     inline friend node operator+(node a, int p) {
47         return node(a.p + p, a.t, a.se + p);
48     }
49 };
50
51 struct SGTB {
52     #define inf 1e9
53     node mx[maxn << 2], mi[maxn << 2];
```

```

54     int tag[maxn << 2];
55     ll s[maxn << 2];
56 # define ls (x<<1)
57 # define rs (x<<1|1)
58
59     inline void up(int x) {
60         mx[x] = combineMax(mx[ls], mx[rs]);
61         mi[x] = combineMin(mi[ls], mi[rs]);
62         s[x] = s[ls] + s[rs];
63     }
64
65     // a = max(a, t)
66     inline void pushmax(int x, int l, int r, int p) {
67         s[x] += 1ll * mi[x].t * (p - mi[x].p);
68         mi[x].p = p;
69         mx[x].p = max(mx[x].p, p);
70         if (mi[x].p == mx[x].p) {
71             mi[x].se = inf, mx[x].se = -inf;
72             mi[x].t = mx[x].t = r - l + 1;
73             s[x] = 1ll * mi[x].p * (r - l + 1);
74         } else mx[x].se = max(mx[x].se, p);
75     }
76
77     // a = min(a, t)
78     inline void pushmin(int x, int l, int r, int p) {
79         s[x] += 1ll * mx[x].t * (p - mx[x].p);
80         mx[x].p = p;
81         mi[x].p = min(mi[x].p, p);
82         if (mi[x].p == mx[x].p) {
83             mi[x].se = inf, mx[x].se = -inf;
84             mi[x].t = mx[x].t = r - l + 1;
85             s[x] = 1ll * mi[x].p * (r - l + 1);
86         } else mi[x].se = min(mi[x].se, p);
87     }
88
89     inline void pushtag(int x, int l, int r, int p) {
90         tag[x] += p;
91         s[x] += 1ll * (r - l + 1) * p;
92         mx[x] = mx[x] + p, mi[x] = mi[x] + p;
93     }
94
95     inline void down(int x, int l, int r) {
96         int mid = l + r >> 1;
97         if (tag[x]) {
98             pushtag(ls, l, mid, tag[x]);
99             pushtag(rs, mid + 1, r, tag[x]);
100            tag[x] = 0;
101        }
102        if (mx[ls].p > mx[x].p && mx[ls].se < mx[x].p) pushmin(ls, l, mid, mx[x].p);
103        if (mx[rs].p > mx[x].p && mx[rs].se < mx[x].p) pushmin(rs, mid + 1, r, mx[x].p);
104        if (mi[ls].p < mi[x].p && mi[ls].se > mi[x].p) pushmax(ls, l, mid, mi[x].p);
105        if (mi[rs].p < mi[x].p && mi[rs].se > mi[x].p) pushmax(rs, mid + 1, r, mi[x].p);
106    }
107
108    inline void build(int x, int l, int r) {
109        tag[x] = 0;
110        if (l == r) {

```

```

111         mx[x].p = mi[x].p = s[x] = a[l], mx[x].t = mi[x].t = 1, mx[x].se = -inf, mi
           [x].se = inf;
112         return;
113     }
114     int mid = l + r >> 1;
115     build(ls, l, mid);
116     build(rs, mid + 1, r);
117     up(x);
118 }
119
120 // add p to [L, R]
121 inline void edt(int x, int l, int r, int L, int R, int p) {
122     if (L <= l && r <= R) {
123         pushtag(x, l, r, p);
124         return;
125     }
126     down(x, l, r);
127     int mid = l + r >> 1;
128     if (L <= mid) edt(ls, l, mid, L, R, p);
129     if (R > mid) edt(rs, mid + 1, r, L, R, p);
130     up(x);
131 }
132
133 // ai=min(ai, p) in [L, R]
134 inline void edtmin(int x, int l, int r, int L, int R, int p) {
135     if (mx[x].p <= p) return;
136     if (L <= l && r <= R && mx[x].se < p) {
137         pushmin(x, l, r, p);
138         return;
139     }
140     down(x, l, r);
141     int mid = l + r >> 1;
142     if (L <= mid) edtmin(ls, l, mid, L, R, p);
143     if (R > mid) edtmin(rs, mid + 1, r, L, R, p);
144     up(x);
145 }
146
147 // ai=max(ai, p) in [L, R]
148 inline void edtmax(int x, int l, int r, int L, int R, int p) {
149     if (mi[x].p >= p) return;
150     if (L <= l && r <= R && mi[x].se > p) {
151         pushmax(x, l, r, p);
152         return;
153     }
154     down(x, l, r);
155     int mid = l + r >> 1;
156     if (L <= mid) edtmax(ls, l, mid, L, R, p);
157     if (R > mid) edtmax(rs, mid + 1, r, L, R, p);
158     up(x);
159 }
160
161 inline int gmax(int x, int l, int r, int L, int R) {
162     if (L <= l && r <= R) return mx[x].p;
163     down(x, l, r);
164     int mid = l + r >> 1, ret = -inf;
165     if (L <= mid) ret = max(ret, gmax(ls, l, mid, L, R));
166     if (R > mid) ret = max(ret, gmax(rs, mid + 1, r, L, R));
167     return ret;
168 }

```

```

169
170     inline int gmin(int x, int l, int r, int L, int R) {
171         if (L <= l && r <= R) return mi[x].p;
172         down(x, l, r);
173         int mid = l + r >> 1, ret = inf;
174         if (L <= mid) ret = min(ret, gmin(ls, l, mid, L, R));
175         if (R > mid) ret = min(ret, gmin(rs, mid + 1, r, L, R));
176         return ret;
177     }
178
179     inline ll gsum(int x, int l, int r, int L, int R) {
180         if (L <= l && r <= R) return s[x];
181         down(x, l, r);
182         int mid = l + r >> 1;
183         ll ret = 0;
184         if (L <= mid) ret += gsum(ls, l, mid, L, R);
185         if (R > mid) ret += gsum(rs, mid + 1, r, L, R);
186         return ret;
187     }
188
189     inline void debug(int x, int l, int r) {
190         printf("%d %d %d [%d %d %d] [%d %d %d] %lld\n", x, l, r, mx[x].p, mx[x].t, mx[x]
            .se, mi[x].p, mi[x].t, mi[x].se,
191             s[x]);
192         if (l == r) return;
193         int mid = l + r >> 1;
194         debug(ls, l, mid);
195         debug(rs, mid + 1, r);
196     }
197
198 # undef ls
199 # undef rs
200 # undef inf
201 } T;

```

3.1.2 离散化区间

```

1 // 原题1e5个区间有2e5个端点, 离散化出来4e5个区间
2 // 然后线段树需要4e5*4=16e5的大小
3 // 注意三个数组要开离散化数量的四倍, 如果不需要sz可以不用这个数组。
4 int val[maxn << 4];
5 int lpos[maxn << 2], rpos[maxn << 2], tot, sz[maxn << 2];
6 vector<int> xpos;
7 sort(xpos.begin(), xpos.end());
8 xpos.erase(unique(xpos.begin(), xpos.end()), xpos.end());
9 tot = 1;
10 lpos[1] = rpos[1] = xpos[0];
11 sz[1] = 1;
12 for (int i = 1; i < xpos.size(); ++i) {
13     if (xpos[i] - xpos[i - 1] != 1) {
14         lpos[++tot] = xpos[i - 1] + 1;
15         rpos[tot] = xpos[i] - 1;
16         sz[tot] = rpos[tot] - lpos[tot] + 1;
17     }
18     ++tot;
19     lpos[tot] = rpos[tot] = xpos[i];
20     sz[tot] = 1;
21 }

```

```

22 le = lower_bound(lpos + 1, lpos + 1 + tot, p[i].x) - lpos;
23 re = upper_bound(rpos + 1, rpos + 1 + tot, p[i].y) - rpos - 1;

```

3.1.3 动态区间最大子段和

```

1 namespace ST {
2     struct node{
3         ll ans,ls,rs,sum;
4     }xx[maxn << 2];
5     inline void pushdown(int x){
6         xx[x].sum=xx[x<<1].sum+xx[x<<1|1].sum;
7         xx[x].ls=max(xx[x<<1].ls,xx[x<<1].sum+xx[x<<1|1].ls);
8         xx[x].rs=max(xx[x<<1|1].rs,xx[x<<1|1].sum+xx[x<<1].rs);
9         xx[x].ans=max(xx[x<<1].ans,max(xx[x<<1|1].ans,xx[x<<1].rs+xx[x<<1|1].ls));
10        return;
11    }
12    inline void build(int k,int l,int r){
13        if(l==r){
14            xx[k].ls=xx[k].rs=xx[k].ans=xx[k].sum=0;
15            return;
16        }
17        int mid=l+r>>1;
18        build(k<<1,l,mid),build(k<<1|1,mid+1,r);
19        pushdown(k);
20        return;
21    }
22    inline void change(int k,int l,int r,int x,int y,int w){ // 1, 1, n
23        if(x<=l&&r<=y){
24            xx[k].ls += w;
25            xx[k].rs += w;
26            xx[k].ans += w;
27            xx[k].sum += w;
28        } // xx[k].ls=xx[k].rs=xx[k].ans=xx[k].sum=w;
29        return;
30    }
31    int mid=l+r>>1;
32    if(x<=mid) change(k<<1,l,mid,x,y,w);
33    if(mid<y) change(k<<1|1,mid+1,r,x,y,w);
34    pushdown(k);
35    return;
36 }
37 inline node query(int k,int l,int r,int x,int y){
38     if(x<=l&&r<=y) {
39         return xx[k];
40     }
41     int mid=l+r>>1;
42     if(x<=mid&&!(mid<y)) return query(k<<1,l,mid,x,y);
43     else if(!(x<=mid)&&mid<y) return query(k<<1|1,mid+1,r,x,y);
44     else{
45         node st,t1=query(k<<1,l,mid,x,y),t2=query(k<<1|1,mid+1,r,x,y);
46         st.sum=t1.sum+t2.sum;
47         st.ls=max(t1.ls,t1.sum+t2.ls);
48         st.rs=max(t2.rs,t2.sum+t1.rs);
49         st.ans=max(t1.ans,max(t2.ans,t1.rs+t2.ls));
50         return st;
51     }
52 }
53 }

```

3.1.4 动态开点权值线段树

```

1  int root[100005];
2  int ls[1800000], rs[1800000], sum[1800000];
3  int sz = 0;
4
5  void insert(int &k, int l, int r, int val){
6      if (!k) k = ++sz;
7      if (l == r) {
8          sum[k] = 1;
9          return;
10     }
11     int mid = (l + r) >> 1;
12     if (val <= mid) insert(ls[k], l, mid, val);
13     else insert(rs[k], mid + 1, r, val);
14     sum[k] = sum[ls[k]] + sum[rs[k]];
15 }
16
17 int query(int k, int l, int r, int rank) {
18     if (l == r) return l;
19     int mid = (l + r) >> 1;
20     if (sum[ls[k]] >= rank) return query(ls[k], l, mid, rank);
21     else return query(rs[k], mid + 1, r, rank - sum[ls[k]]);
22 }
23 int merge(int x, int y)
24 {
25     if (!x) return y;
26     if (!y) return x;
27     ls[x] = merge(ls[x], ls[y]);
28     rs[x] = merge(rs[x], rs[y]);
29     sum[x] = sum[ls[x]] + sum[rs[x]];
30     return x;
31 }
32 insert(root[i], 1, n, a[i]);
33 query(root[p], 1, n, x);

```

3.1.5 扫描线

```

1  // 范用型扫描线, del储存上界+1, add储存下界, 先del后add即可
2  struct node {
3      int lpos, rpos, linepos;
4      bool operator < (const node& oth) const {
5          return linepos < oth.linepos;
6      }
7  };
8  vector<node> add, del;
9  int delpos = 0;
10 int res = 0;
11 for (int addpos = 0; addpos < add.size(); ++addpos) {
12     while (delpos < del.size() && del[delpos].linepos <= add[addpos].linepos) {
13         up(del[delpos].lpos, del[delpos].rpos, -1);
14         delpos ++;
15     }
16     up(add[addpos].lpos, add[addpos].rpos, 1);
17     res = max(res, val[1]);
18 }
19
20 // 求面积并

```

```

21 #define maxn 222
22 #define tmp (st<<1)
23 #define mid ((l+r)>>1)
24 #define lson l,mid,tmp
25 #define rson mid+1,r,tmp|1
26 using namespace std;
27 int cnt[maxn<<2];
28 double sum[maxn<<2];
29 double x[maxn];
30 struct Seg{
31     double h,l,r;
32     int s;
33     Seg(){}
34     Seg(double a,double b,double c,int d):l(a),r(b),h(c),s(d){}
35     bool operator<(const Seg &cmp)const{
36         return h<cmp.h;
37     }
38 }ss[maxn];
39 void push_up(int st,int l,int r){
40     if(cnt[st])sum[st]=x[r+1]-x[l];
41     else if(l==r)sum[st]=0;
42     else sum[st]=sum[tmp]+sum[tmp|1];
43 }
44 void update(int L,int R,int c,int l,int r,int st){
45     if(L<=l&&r<=R){
46         cnt[st]+=c;
47         push_up(st,l,r);
48         return ;
49     }
50     if(L<=mid)update(L,R,c,lson);
51     if(R>mid)update(L,R,c,rson);
52     push_up(st,l,r);
53 }
54 int main(){
55     int n,tot=1,m;
56     while(scanf("%d",&n)&&n){
57         double a,b,c,d;
58         m=0;
59         while(n--){
60             scanf("%lf%lf%lf%lf",&a,&b,&c,&d);
61             x[m]=a;
62             ss[m++]=Seg(a,c,b,1);
63             x[m]=c;
64             ss[m++]=Seg(a,c,d,-1);
65         }
66         sort(x,x+m);
67         sort(ss,ss+m);
68         double ans=0;
69         for(int i=0;i<m;++i){
70             int l=lower_bound(x,x+m,ss[i].l)-x;
71             int r=lower_bound(x,x+m,ss[i].r)-x-1;
72             update(l,r,ss[i].s,0,m-1,1);
73             ans+=sum[l]*(ss[i+1].h-ss[i].h);
74         }
75         printf("Test case # %d\nTotal explored area: %.2lf\n",tot++,ans);
76     }
77     return 0;
78 }
79

```

```
80 // 面积交
81 #include<bits/stdc++.h>
82 #define maxn 100005
83 #define lson l,mid,rt<<1
84 #define rson mid+1,r,rt<<1|1
85 #define pb push_back
86 using namespace std;
87
88 double tree[maxn<<2],tree2[maxn<<2];
89 int lazy[maxn<<2];
90 vector<double>ve;
91
92 struct seg{
93     double l,r,h;
94     int flag;
95     seg(){}
96     seg(double _l,double _r,double _h,int _flag){l=_l,r=_r,h=_h,flag=_flag;}
97     bool operator<(const seg &b)const{return h<b.h;}
98 }s[maxn];
99
100 void push_up(int l,int r,int rt){
101     if(lazy[rt]) tree[rt]=ve[r]-ve[l-1];
102     else if(l==r) tree[rt]=0;
103     else tree[rt]=tree[rt<<1]+tree[rt<<1|1];
104 }
105
106 void push_up2(int l,int r,int rt){
107     if(lazy[rt]>1) tree2[rt]=ve[r]-ve[l-1];
108     else if(l==r) tree2[rt]=0;
109     else if(lazy[rt]==1) tree2[rt]=tree[rt<<1]+tree[rt<<1|1];
110     else tree2[rt]=tree2[rt<<1]+tree2[rt<<1|1];
111 }
112
113 void build(int l,int r,int rt){
114     tree[rt]=0,lazy[rt]=0;
115     if(l==r) return;
116     int mid=l+r>>1;
117     build(lson);
118     build(rson);
119 }
120
121 void add(int L,int R,int v,int l,int r,int rt){
122     if(L<=l&&R>=r){
123         lazy[rt]+=v;
124         push_up(l,r,rt);
125         push_up2(l,r,rt);
126         return;
127     }
128     int mid=l+r>>1;
129     if(L<=mid) add(L,R,v,lson);
130     if(R>mid) add(L,R,v,rson);
131     push_up(l,r,rt);
132     push_up2(l,r,rt);
133 }
134
135 int getid(double x){ return lower_bound(ve.begin(),ve.end(),x)-ve.begin()+1;}
136
137 int main(){
138     int n;
```



```

139     int Case=1;
140     int T;
141     scanf("%d",&T);
142     while(T--){
143         scanf("%d",&n);
144         ve.clear();
145         int tot=0;
146         double x1,y1,x2,y2;
147         for(int i=1;i<=n;i++){
148             scanf("%lf %lf %lf %lf",&x1,&y1,&x2,&y2);
149             ve.pb(x1),ve.pb(x2);
150             s[++tot]=seg(x1,x2,y1,1);
151             s[++tot]=seg(x1,x2,y2,-1);
152         }
153         sort(ve.begin(),ve.end());
154         ve.erase(unique(ve.begin(),ve.end()),ve.end());
155         sort(s+1,s+tot+1);
156         int N=ve.size();
157         build(1,N,1);
158         double ans=0;
159         for(int i=1;i<tot;i++){
160             int L=getid(s[i].l);
161             int R=getid(s[i].r)-1;
162             add(L,R,s[i].flag,1,N,1);
163             ans+=tree2[1]*(s[i+1].h-s[i].h);
164         }
165         printf("%.2f\n",ans);
166     }
167 }
168
169 // 求周长并
170 #include<bits/stdc++.h>
171 #define maxn 100005
172 #define lson l,mid,rt<<1
173 #define rson mid+1,r,rt<<1|1
174 #define pb push_back
175 using namespace std;
176
177 int tree[maxn<<2];
178 int lazy[maxn<<2];
179 vector<int>ve[2];
180 int k;
181
182 struct seg{
183     int l,r,h;
184     int flag;
185     seg(){}
186     seg(int _l,int _r,int _h,int _flag){l=_l,r=_r,h=_h,flag=_flag;}
187     bool operator<(const seg &b)const{return h<b.h;}
188 }s[maxn];
189
190 void push_up(int l,int r,int rt){
191     if(lazy[rt]) tree[rt]=ve[k][r]-ve[k][l-1];
192     else if(l==r) tree[rt]=0;
193     else tree[rt]=tree[rt<<1]+tree[rt<<1|1];
194 }
195
196 void build(int l,int r,int rt){
197     tree[rt]=0,lazy[rt]=0;

```

```
198     if(l==r) return;
199     int mid=l+r>>1;
200     build(lson);
201     build(rson);
202 }
203
204 void add(int L,int R,int v,int l,int r,int rt){
205     if(L<=l&&R>=r){
206         lazy[rt]+=v;
207         push_up(l,r,rt);
208         return;
209     }
210     int mid=l+r>>1;
211     if(L<=mid) add(L,R,v,lson);
212     if(R>mid) add(L,R,v,rson);
213     push_up(l,r,rt);
214 }
215
216 int getid(int x){return lower_bound(ve[k].begin(),ve[k].end(),x)-ve[k].begin()+1;}
217
218 int main(){
219     int n;
220     while(~scanf("%d",&n)){
221         ve[0].clear();
222         ve[1].clear();
223         int x1,y1,x2,y2;
224         for(int i=1;i<=n;i++){
225             scanf("%d %d %d %d",&x1,&y1,&x2,&y2);
226             ve[0].pb(x1),ve[0].pb(x2);
227             ve[1].pb(y1),ve[1].pb(y2);
228             s[i]=seg(x1,x2,y1,1);
229             s[i+n]=seg(x1,x2,y2,-1);
230             s[i+n+n]=seg(y1,y2,x1,1);
231             s[i+n+n+n]=seg(y1,y2,x2,-1);
232         }
233         int ans=0;
234         int pos=1;
235         for(k=0;k<2;k++){
236             sort(ve[k].begin(),ve[k].end());
237             ve[k].erase(unique(ve[k].begin(),ve[k].end()),ve[k].end());
238             sort(s+pos,s+pos+n+n);
239             int N=ve[k].size();
240             build(1,N,1);
241             int pre=0;
242             for(int i=pos;i<pos+n+n;i++){
243                 int L=getid(s[i].l);
244                 int R=getid(s[i].r)-1;
245                 add(L,R,s[i].flag,1,N,1);
246                 ans+=abs(tree[1]-pre);
247                 pre=tree[1];
248             }
249             pos+=n+n;
250         }
251         printf("%d\n",ans);
252     }
253 }
```

3.2 HLD

3.2.1 HLD

```

1  #include <bits/stdc++.h>
2  #define ll long long
3  using namespace std;
4  /*
5  node 计算点权, path 下放后计算边权, edge 根据边的编号计算边权
6  work 中没有build需手动写
7  sz[] 数组, 以x为根的子树节点个数
8  top[] 数组, 当前节点的所在链的顶端节点
9  son[] 数组, 重儿子
10 deep[] 数组, 当前节点的深度
11 fa[] 数组, 当前节点的父亲
12 idx[] 数组, 树中每个节点剖分后的新编号
13 rnk[] 数组, idx的逆, 表示线段上中当前位置表示哪个节点
14 */
15
16 const int maxn = 1e5+5;
17
18 int sz[maxn], top[maxn], son[maxn], deep[maxn], fa[maxn], idx[maxn], rnk[maxn];
19 int tot;
20 int n, le, re;
21 ll k;
22
23 struct HLD {
24 #define type int
25
26     struct edge {
27         int a, b;
28         type v;
29
30         edge(int _a, int _b, type _v = 0) : a(_a), b(_b), v(_v) {}
31     };
32
33     struct node {
34         int to;
35         type w;
36
37         node() {}
38
39         node(int _to, type _w) : to(_to), w(_w) {}
40     };
41
42     vector<int> mp[maxn];
43     vector<edge> e;
44
45     void init(int _n) {
46         n = _n;
47         for (int i = 0; i <= n; i++) mp[i].clear();
48         e.clear();
49         e.push_back(edge(0, 0));
50     }
51
52     void add_edge(int a, int b, type v = 0) {
53 //         e.push_back(edge(a,b,v));
54         mp[a].push_back(b);
55         mp[b].push_back(a);

```

```

56     }
57
58     void dfs1(int x, int pre, int h) {
59         int i, to;
60         deep[x] = h;
61         fa[x] = pre;
62         sz[x] = 1;
63         for (i = 0; i < (int) (mp[x].size()); i++) {
64             to = mp[x][i];
65             if (to == pre) continue;
66             dfs1(to, x, h + 1);
67             sz[x] += sz[to];
68             if (son[x] == -1 || sz[to] > sz[son[x]]) son[x] = to;
69         }
70     }
71
72     void dfs2(int x, int tp) {
73         int i, to;
74         top[x] = tp;
75         idx[x] = ++tot;
76         rnk[idx[x]] = x;
77         if (son[x] == -1) return;
78         dfs2(son[x], tp);
79         for (i = 0; i < (int) (mp[x].size()); i++) {
80             to = mp[x][i];
81             if (to != son[x] && to != fa[x]) dfs2(to, to);
82         }
83     }
84
85     void work(int _rt = 1) {
86         memset(son, -1, sizeof son);
87         tot = 0;
88         dfs1(_rt, 0, 0);
89         dfs2(_rt, _rt);
90     }
91
92     int LCA(int x, int y) {
93         while (top[x] != top[y]) {
94             if (deep[top[x]] < deep[top[y]]) swap(x, y);
95             x = fa[top[x]];
96         }
97         if (deep[x] > deep[y]) swap(x, y);
98         return x;
99     }
100
101     void modify_node(int x, int y, type val) {
102         while (top[x] != top[y]) {
103             if (deep[top[x]] < deep[top[y]]) swap(x, y);
104             le = idx[top[x]], re = idx[x];
105             k = val;
106             update(1, 1, n);
107             x = fa[top[x]];
108         }
109         if (deep[x] > deep[y]) swap(x, y);
110         le = idx[x], re = idx[y];
111         k = val;
112         update(1, 1, n);
113     }
114

```

```

115     type query_node(int x, int y) {
116         type res = 0;
117         while (top[x] != top[y]) {
118             if (deep[top[x]] < deep[top[y]]) swap(x, y);
119             le = idx[top[x]], re = idx[x];
120             res += query(1, 1, n);
121             x = fa[top[x]];
122         }
123         if (deep[x] > deep[y]) swap(x, y);
124         le = idx[x], re = idx[y];
125         res += query(1, 1, n);
126         return res;
127     }
128
129     //path
130     void init_path()
131     {
132         v[idx[rt]]=0;
133         for(int i=1;i<n;i++)
134         {
135             if(deep[e[i].a]<deep[e[i].b]) swap(e[i].a,e[i].b);
136             a[idx[e[i].a]]=e[i].v;
137         }
138         build(n);
139     }
140     void modify_edge(int id, type val) {
141         if (deep[e[id].a] > deep[e[id].b]) {
142             le = idx[e[id].a], re = idx[e[id].a];
143             k = val;
144             update(1, 1, n);
145         } else {
146             le = idx[e[id].b], re = idx[e[id].b];
147             k = val;
148             update(1, 1, n);
149         }
150     }
151
152     void modify_path(int x, int y, type val) {
153         while (top[x] != top[y]) {
154             if (deep[top[x]] < deep[top[y]]) swap(x, y);
155             le = idx[top[x]], re = idx[x];
156             k = val;
157             update(1, 1, n);
158             x = fa[top[x]];
159         }
160         if (deep[x] > deep[y]) swap(x, y);
161         if (x != y) {
162             le = idx[x] + 1, re = idx[y];
163             k = val;
164             update(1, 1, n);
165         }
166     }
167
168     type query_path(int x, int y) {
169         type res = 0;
170         while (top[x] != top[y]) {
171             if (deep[top[x]] < deep[top[y]]) swap(x, y);
172             le = idx[top[x]], re = idx[x];
173             res += query(1, 1, n);

```

```

174         x = fa[top[x]];
175     }
176     if (deep[x] > deep[y]) swap(x, y);
177     if (x != y) {
178         le = idx[x] + 1, re = idx[y];
179         res += query(1, 1, n);
180     }
181     return res;
182 }
183
184 #undef type
185 } hld;

```

3.3 RMQ

3.3.1 RMQbyIndex

```

1  int pmax(int x, int y) { return a[x] > a[y] ? x : y; }
2  int pmin(int x, int y) { return a[x] < a[y] ? x : y; }
3  void init() {
4      for (int i = 1; i <= n; i++) {
5          maxx[i][0] = minn[i][0] = i;
6      }
7      for (int j = 1; 1 << (j - 1) <= n; j++) {
8          for (int i = 1; i + (1 << j) - 1 <= n; i++) {
9              int t = 1 << (j - 1);
10             maxx[i][j] = pmax(maxx[i][j - 1], maxx[i + t][j - 1]);
11             minn[i][j] = pmin(minn[i][j - 1], minn[i + t][j - 1]);
12         }
13     }
14 }
15 int query(int l, int r) {
16     int j = 0;
17     while ((1 << (j + 1)) <= r - l + 1) j++;
18     int i = r - (1 << j) + 1;
19     // return pmax(maxx[l][j], maxx[i][j]);
20     return pmin(minn[l][j], minn[i][j]);
21 }

```

3.3.2 RMQinNM

```

1  // 二维RMQ
2  int v[302][302];
3  int maxx[302][302][9][9], minn[302][302][9][9];
4
5  void RMQ(int n, int m) {
6      int i, j, ii, jj;
7      for (i = 1; i <= n; i++) {
8          for (j = 1; j <= m; j++) {
9              maxx[i][j][0][0] = minn[i][j][0][0] = v[i][j];
10         }
11     }
12     for (ii = 0; (1 << ii) <= n; ii++) {
13         for (jj = 0; (1 << jj) <= m; jj++) {
14             if (ii + jj) {
15                 for (i = 1; i + (1 << ii) - 1 <= n; i++) {
16                     for (j = 1; j + (1 << jj) - 1 <= m; j++) {
17                         if (ii) {

```

```

18             minn[i][j][ii][jj] = min(minn[i][j][ii - 1][jj], minn[i +
19             (1 << (ii - 1))][j][ii - 1][jj]);
20             maxx[i][j][ii][jj] = max(maxx[i][j][ii - 1][jj], maxx[i +
21             (1 << (ii - 1))][j][ii - 1][jj]);
22         } else {
23             minn[i][j][ii][jj] = min(minn[i][j][ii][jj - 1], minn[i][j
24             + (1 << (jj - 1))][ii][jj - 1]);
25             maxx[i][j][ii][jj] = max(maxx[i][j][ii][jj - 1], maxx[i][j
26             + (1 << (jj - 1))][ii][jj - 1]);
27         }
28     }
29 }
30
31 int query(int x1, int y1, int x2, int y2) {
32     int k1 = 0;
33     while ((1 << (k1 + 1)) <= x2 - x1 + 1) k1++;
34     int k2 = 0;
35     while ((1 << (k2 + 1)) <= y2 - y1 + 1) k2++;
36     x2 = x2 - (1 << k1) + 1;
37     y2 = y2 - (1 << k2) + 1;
38     return max(max(maxx[x1][y1][k1][k2], maxx[x1][y2][k1][k2]), max(maxx[x2][y1][k1][k2]
39     ], maxx[x2][y2][k1][k2]));
40 // return min(min(minn[x1][y1][k1][k2], minn[x1][y2][k1][k2]), min(minn[x2][y1][k1][k2],
41 minn[x2][y2][k1][k2]));
42 }

```

3.4 MO

3.4.1 分块

```

1 // 非预处理数组版
2 inline int belong(int x) { return (x - 1) / block + 1; }
3 inline int lpos(int x) { return 1 + (x - 1) * block; }
4 inline int rpos(int x) { return min(n, x * block); }
5 int sz = (n - 1) / block + 1;
6
7 // 预处理版, maxn大于1e6已经不可能处理了
8 const int maxb = 1005;
9 int n, m;
10 int belong[maxn], lpos[maxb], rpos[maxb];
11 int val[maxn], lazy[maxb];
12 int block;
13
14 scanf("%d", &n);
15 block = sqrt(n);
16 for (int i = 1; i <= n; ++i) {
17     scanf("%d", &val[i]);
18     belong[i] = (i - 1) / block + 1;
19 }
20 int sz = (n - 1) / block + 1;
21 for (int i = 1; i <= sz; ++i) {
22     lpos[i] = 1 + (i - 1) * block;
23     rpos[i] = i * block;
24 }
25 rpos[sz] = n;

```

3.4.2 带修莫队

```
1  #include <bits/stdc++.h>
2  #define ll long long
3  using namespace std;
4  const int maxn = 50005;
5
6  struct MO {
7      int l, r, id, oppre;
8  }q[maxn];
9
10 int n, m, col[maxn], block, belong[maxn], colpre[maxn];
11 int changepos[maxn], changepre[maxn], changenow[maxn];
12 int vis[maxn * 20];
13 int ans;
14 int res[maxn];
15 bool cmp(const MO& a, const MO& b) {
16     if (belong[a.l] != belong[b.l]) return a.l < b.l;
17     if (belong[a.r] != belong[b.r]) return a.r < b.r;
18     return a.oppre < b.oppre;
19 }
20 void add(int x) {}
21
22 void del(int x) {}
23
24 void unmodify(int pos, int now) {
25     if (q[pos].l <= changepos[now] && changepos[now] <= q[pos].r) {
26         del(changenow[now]);
27         add(changepre[now]);
28     }
29     col[changepos[now]] = changepre[now];
30 }
31
32 void modify(int pos, int now) {
33     if (q[pos].l <= changepos[now] && changepos[now] <= q[pos].r) {
34         del(changepre[now]);
35         add(changenow[now]);
36     }
37     col[changepos[now]] = changenow[now];
38 }
39
40 int main() {
41     scanf("%d%d", &n, &m);
42     block = pow(n, 0.66666);
43     for (int i = 1; i <= n; ++i) {
44         scanf("%d", &col[i]);
45         colpre[i] = col[i];
46         belong[i] = i / block + 1;
47     }
48     char s[2];
49     int t = 0, t2 = 0;
50     for (int i = 1; i <= m; ++i) {
51         scanf("%s", s);
52         if (s[0] == 'Q') {
53             ++t;
54             scanf("%d%d", &q[t].l, &q[t].r);
55             q[t].oppre = t2;
56             q[t].id = t;
57         } else {
```



```

58         ++t2;
59         scanf("%d%d", &changeupos[t2], &changenow[t2]);
60         changepre[t2] = colpre[changeupos[t2]];
61         colpre[changeupos[t2]] = changenow[t2];
62     }
63 }
64 sort(q + 1, q + 1 + t, cmp);
65 int l = 1, r = 0, now = 0;
66 for (int i = 1; i <= t; ++i) {
67     while(r < q[i].r) add(col[++r]);
68     while(r > q[i].r) del(col[r--]);
69     while(l < q[i].l) del(col[l++]);
70     while(l > q[i].l) add(col[--l]);
71     while (now < q[i].oppre) modify(i, ++now);
72     while (now > q[i].oppre) unmodify(i, now--);
73     res[q[i].id] = ans;
74 }
75 for (int i = 1; i <= t; ++i) printf("%d\n", res[i]);
76 return 0;
77 }

```

3.4.3 序列莫队

```

1  // const int maxn = 50005;
2
3  struct MO {
4      int l, r, id;
5  }q[maxn];
6
7  int n, m, col[maxn], block, belong[maxn];
8  int vis[maxn * 10];
9  ll res[maxn], ans;
10 bool cmp(const MO& a, const MO& b) { return belong[a.l] == belong[b.l] ? a.r < b.r : a.
    l < b.l; }
11 void add(int x) {
12     vis[x] ++;
13     ans += 1ll * x * (vis[x] * vis[x] - (vis[x] - 1) * (vis[x] - 1));
14 }
15
16 void del(int x) {
17     vis[x] --;
18     ans -= 1ll * x * ((vis[x] + 1) * (vis[x] + 1) - vis[x] * vis[x]);
19 }
20
21 int main() {
22     scanf("%d%d", &n, &m);
23     block = sqrt(n);
24     for (int i = 1; i <= n; ++i) {
25         scanf("%d", &col[i]);
26         belong[i] = i / block + 1;
27     }
28     for (int i = 1; i <= m; ++i) {
29         scanf("%d%d", &q[i].l, &q[i].r);
30         q[i].id = i;
31     }
32     sort(q + 1, q + 1 + m, cmp);
33     int l = 1, r = 0;
34     for (int i = 1; i <= m; ++i) {

```

```

35     while(r < q[i].r) add(col[++r]);
36     while(r > q[i].r) del(col[r--]);
37     while(l < q[i].l) del(col[l++]);
38     while(l > q[i].l) add(col[--l]);
39     res[q[i].id] = ans;
40 }
41 for (int i = 1; i <= m; ++i) printf("%lld\n", res[i]);
42 return 0;
43 }

```

3.4.4 弹飞绵羊

```

1  /*
2  每个装置设定初始弹力系数ki，当绵羊达到第i个装置时，它会往后弹ki步，
3  达到第i+ki个装置，若不存在第i+ki个装置，则绵羊被弹飞。
4  绵羊想知道当它从第i个装置起步时，被弹几次后会被弹飞。
5  为了使得游戏更有趣，Lostmonkey可以修改某个弹力装置的弹力系数，任何时候弹力系数均为正整
    数。
6  */
7
8  int n, m;
9  int belong[maxn], lpos[maxn], rpos[maxn];
10 int val[maxn], nxt[maxn], k[maxn], lst[maxn];
11 int block;
12
13 void update(int pos) {
14     int llim = lpos[belong[pos]], rlim = rpos[belong[pos]];
15     for (int i = pos; i >= llim; --i) {
16         if (val[i] + i > rlim) {
17             k[i] = 1;
18             nxt[i] = val[i] + i;
19             if (val[i] + i > n) lst[i] = i;
20             else lst[i] = lst[nxt[i]];
21         } else {
22             k[i] = 1 + k[val[i] + i];
23             nxt[i] = nxt[val[i] + i];
24             lst[i] = lst[val[i] + i];
25         }
26     }
27 }
28
29 void init() {
30     for (int i = n; i >= 1; --i) {
31         int rlim = rpos[belong[i]];
32         if (val[i] + i > rlim) {
33             k[i] = 1;
34             nxt[i] = val[i] + i;
35             if (val[i] + i > n) lst[i] = i;
36             else lst[i] = lst[nxt[i]];
37         } else {
38             k[i] = 1 + k[val[i] + i];
39             nxt[i] = nxt[val[i] + i];
40             lst[i] = lst[val[i] + i];
41         }
42     }
43 }
44
45 int query(int pos) {

```

```

46     int res = 0;
47     while (pos <= n) {
48         res += k[pos];
49         if (nxt[pos] > n) printf("%d ", lst[pos]);
50         pos = nxt[pos];
51     }
52     return res;
53 }
54
55 int main(int argc, char* argv[]) {
56     scanf("%d%d", &n, &m);
57     block = sqrt(n) * 1.6 + 1;
58     for (int i = 1; i <= n; ++i) {
59         scanf("%d", &val[i]);
60         belong[i] = (i - 1) / block + 1;
61     }
62     int sz = (n - 1) / block + 1;
63     for (int i = 1; i <= sz; ++i) {
64         lpos[i] = 1 + (i - 1) * block;
65         rpos[i] = i * block;
66     }
67     rpos[sz] = n;
68     init();
69     while (m--) {
70         int op;
71         scanf("%d", &op);
72         if (op == 1) {
73             int pos;
74             scanf("%d", &pos);
75             printf("%d\n", query(pos));
76         } else {
77             int pos, kl;
78             scanf("%d%d", &pos, &kl);
79             val[pos] = kl;
80             update(pos);
81         }
82     }
83     return 0;
84 }

```

3.4.5 树莫队

```

1 // rnk保存欧拉序
2 int sz[maxn], top[maxn], son[maxn], deep[maxn], fa[maxn], idx[maxn], ed[maxn], rnk[maxn
   *2];
3 int tot, n, m;
4 vector<int> edge[maxn];
5 int val[maxn];
6 vector<int> xpos;
7
8 inline void dfs1(int u, int pre, int h) {
9     deep[u] = h;
10    fa[u] = pre;
11    sz[u] = 1;
12    for (auto to : edge[u]) {
13        if (to == pre) continue;
14        dfs1(to, u, h + 1);
15        sz[u] += sz[to];

```

```

16         if (son[u] == 0 || sz[to] > sz[son[u]]) son[u] = to;
17     }
18 }
19
20 inline void dfs2(int u, int tp) {
21     top[u] = tp;
22     idx[u] = ++tot, rnk[tot] = u;
23     if (son[u] == 0) {
24         ed[u] = ++tot, rnk[tot] = u;
25         return;
26     }
27     dfs2(son[u], tp);
28     for (auto to : edge[u]) {
29         if (to != son[u] && to != fa[u]) dfs2(to, to);
30     }
31     ed[u] = ++tot, rnk[tot] = u;
32 }
33
34 inline int LCA(int x, int y) {
35     while (top[x] != top[y]) {
36         if (deep[top[x]] < deep[top[y]]) swap(x, y);
37         x = fa[top[x]];
38     }
39     if (deep[x] > deep[y]) swap(x, y);
40     return x;
41 }
42
43 int belong[maxn*2], block;
44 int res[maxn], ans;
45 // 每个点是否访问 (欧拉序去重)
46 int vis[maxn];
47 // 标记数组
48 int pre[maxn];
49 struct MO {
50     int l, r, id, lca;
51     bool operator < (const MO& oth) const {
52         return belong[l] == belong[oth.l] ? r < oth.r : belong[l] < belong[oth.l];
53     }
54 }q[maxm];
55
56 inline void add(int x) {
57     pre[x] ++;
58     if (pre[x] == 1) ans ++;
59 }
60
61 inline void del(int x) {
62     pre[x] --;
63     if (pre[x] == 0) ans --;
64 }
65
66 inline void deal(int x) {
67     vis[x] ? del(val[x]) : add(val[x]);
68     vis[x] = !vis[x];
69 }
70
71 int main(int argc, char* argv[]) {
72     scanf("%d%d", &n, &m);
73     block = sqrt(n);
74     xpos.resize(n + 1);

```

```

75     for (int i = 1; i <= n; ++i) {
76         scanf("%d", &val[i]);
77         xpos[i] = val[i];
78     }
79     sort(xpos.begin(), xpos.end());
80     xpos.erase(unique(xpos.begin(), xpos.end()), xpos.end());
81     for (int i = 1; i <= n; ++i) val[i] = lower_bound(xpos.begin(), xpos.end(), val[i])
        - xpos.begin();
82     // 欧拉序长度为n两倍所以分块要分两倍大小
83     for (int i = 1; i <= n * 2; ++i) {
84         belong[i] = (i - 1) / block + 1;
85     }
86     for (int i = 1, u, v; i < n; ++i) {
87         scanf("%d%d", &u, &v);
88         edge[u].push_back(v);
89         edge[v].push_back(u);
90     }
91     // 树剖预处理lca
92     dfs1(1, 0, 0);
93     dfs2(1, 1);
94     for (int i = 1, x, y; i <= m; ++i) {
95         scanf("%d%d", &x, &y);
96         if (idx[x] > idx[y]) swap(x, y);
97         int _lca = LCA(x, y);
98         q[i].id = i;
99         if (_lca == x) q[i].l = idx[x], q[i].r = idx[y], q[i].lca = 0;
100        else q[i].l = ed[x], q[i].r = idx[y], q[i].lca = _lca;
101    //      cerr << q[i].l << " " << q[i].r << " " << q[i].id << " " << q[i].lca << endl;
102    }
103    sort(q + 1, q + 1 + m);
104    int l = 1, r = 0;
105    for (int i = 1; i <= m; ++i) {
106        while(r < q[i].r) deal(rnk[++r]);
107        while(r > q[i].r) deal(rnk[r--]);
108        while(l < q[i].l) deal(rnk[l++]);
109        while(l > q[i].l) deal(rnk[--l]);
110        if (q[i].lca) deal(q[i].lca);
111        res[q[i].id] = ans;
112        if (q[i].lca) deal(q[i].lca);
113    }
114    for (int i = 1; i <= m; ++i) {
115        printf("%d\n", res[i]);
116    }
117    return 0;
118 }

```

3.5 VirtualTree

3.5.1 VirtualTree

```

1  const int pow2 = 19;
2  int n;
3  vector<int> adj0[maxn], adj1[maxn];
4  int st[maxn << 1][pow2 + 1], dep[maxn], euler[maxn], euler_clock;
5  int stk[maxn], fa0[maxn];
6  vector<int> cache;
7  void link0(int u, int v) { adj0[u].emplace_back(v); adj0[v].emplace_back(u); }
8  void link1(int u, int v) { adj1[u].emplace_back(v), cache.push_back(u); }
9  void clearAll() {

```

```

10     for (int i = 1; i <= n; ++i) {
11         adj0[i].clear();
12         adj1[i].clear();
13     }
14     euler_clock = 0;
15 }
16 void clearVT() { for (auto i : cache) adj1[i].clear(); cache.clear(); }
17 void dfs0(int u, int p) {
18     fa0[u] = p;
19     dep[u] = dep[p] + 1;
20     st[++euler_clock][0] = u;
21     euler[u] = euler_clock;
22     for (const auto& v : adj0[u]) if (v != p) {
23         dfs0(v, u);
24         st[++euler_clock][0] = u;
25     }
26 }
27 inline bool cmp(int u, int v) {return dep[u] < dep[v];}
28 inline int upper(int u, int v) {return cmp(u, v) ? u : v;}
29 void lca_init() {
30     for (int i = 0; i != 31 - __builtin_clz(euler_clock); ++i)
31         for (int j = 1; j + (1 << (i + 1)) <= euler_clock; ++j)
32             st[j][i + 1] = upper(st[j][i], st[j + (1 << i)][i]);
33 }
34 inline int lca(int u, int v) {
35     if (u == v) return u;
36     u = euler[u];
37     v = euler[v];
38     if (u > v) swap(u, v);
39     int temp = 31 - __builtin_clz(++v - u);
40     return upper(st[u][temp], st[v - (1 << temp)][temp]);
41 }
42 // build 后 stk[1] 是该树的根节点, 且为有根树
43 void build(vector<int>& key) {
44     sort(key.begin(), key.end(), [&] (int u, int v) { return euler[u] < euler[v]; });
45     key.erase(unique(key.begin(), key.end()), key.end());
46     int top = 0;
47     for (auto u : key) {
48         if (!top) {
49             stk[++top] = u;
50             continue;
51         }
52         int p = lca(u, stk[top]);
53         while (euler[p] < euler[stk[top]]) {
54             if (euler[p] >= euler[stk[top - 1]]) {
55                 link1(p, stk[top]);
56                 if (stk[--top] != p) stk[++top] = p;
57                 break;
58             }
59             link1(stk[top - 1], stk[top]);
60             --top;
61         }
62         stk[++top] = u;
63     }
64     while (top > 1) {
65         link1(stk[top - 1], stk[top]);
66         --top;
67     }
68 }

```

```
69
70 int f[maxn];
71 int res;
72 int vis[maxn];
73 void dfs1(int u) {
74     fa1[u] = p;
75     len[u] = dep[u] - dep[p];
76     for (auto v : adj1[u]) {
77         dfs1(v);
78     }
79 }
80
81 int main(int argc, char* argv[]) {
82     scanf("%d", &n);
83     for (int i = 1, u, v; i < n; ++i) {
84         scanf("%d%d", &u, &v);
85         link0(u, v);
86     }
87     dfs0(1, 0);
88     lca_init();
89     int m; scanf("%d", &m);
90     for (int i = 0; i < m; ++i) {
91         int sz; scanf("%d", &sz);
92         vector<int> key(sz);
93         for (int j = 0; j < sz; ++j) {
94             scanf("%d", &key[j]);
95             vis[key[j]] = 1;
96         }
97         build(key);
98         res = 0;
99         dfs1(stk[1]);
100        printf("%d\n", res);
101        for (int j = 0; j < sz; ++j) vis[key[j]] = 0;
102        clearVT();
103    }
104    return 0;
105 }
```

3.6 PersistentDS

3.6.1 主席树区间 k 大

```
1 // const int maxn = 100005;
2 int n, m;
3 int a[maxn];
4 int root[maxn];
5 int cnt = 0;
6 vector<int> b;
7 struct node {
8     int l, r, val;
9 }p[maxn * 40];
10
11 void update(int l, int r, int pre, int &now, int pos) {
12     p[now = ++cnt] = p[pre];
13     p[now].val++;
14     if (l == r) {
15         return;
16     }
17     int mid = l + r >> 1;
```

```

18     if (pos <= mid) update(l, mid, p[pre].l, p[now].l, pos);
19     else update(mid + 1, r, p[pre].r, p[now].r, pos);
20 }
21
22 /*
23 void build(int pre, int &now, int pos) {
24     now = pre;
25     for (auto i : a[pos]) {
26         update(1, n, pre, now, i);
27     }
28 }
29 */
30
31 int query(int l, int r, int x, int y, int k) {
32     if (l == r) return b[l - 1];
33     int mid = l + r >> 1;
34     int temp = p[p[y].l].val - p[p[x].l].val;
35     if (k <= temp) return query(l, mid, p[x].l, p[y].l, k);
36     return query(mid + 1, r, p[x].r, p[y].r, k - temp);
37 }
38
39 int main(int argc, char *argv[])
40 {
41     while (scanf("%d%d", &n, &m) != EOF) {
42         b.clear();
43         cnt = 0;
44         for (int i = 1; i <= n; ++i) scanf("%d", &a[i]), b.push_back(a[i]);
45         sort(b.begin(), b.end());
46         b.erase(unique(b.begin(), b.end()), b.end());
47         for (int i = 1; i <= n; ++i) {
48             update(1, b.size(), root[i - 1], root[i], lower_bound(b.begin(), b.end(), a
49                 [i]) - b.begin() + 1);
50         }
51         int L, R, k;
52         while (m--) {
53             scanf("%d%d%d", &L, &R, &k);
54             printf("%d\n", query(1, b.size(), root[L - 1], root[R], k));
55         }
56     }
57     return 0;
58 }

```

3.6.2 动态森林

```

1  const int maxn = 10010;
2  struct Splay {
3      int ch[maxn][2], fa[maxn], tag[maxn];
4      void clear(int x) { ch[x][0] = ch[x][1] = fa[x] = tag[x] = 0; }
5      int getch(int x) { return ch[fa[x]][1] == x; }
6      int isroot(int x) { return ch[fa[x]][0] != x && ch[fa[x]][1] != x; }
7      void pushdown(int x) {
8          if (tag[x]) {
9              if (ch[x][0]) swap(ch[ch[x][0]][0], ch[ch[x][0]][1]), tag[ch[x][0]] ^= 1;
10             if (ch[x][1]) swap(ch[ch[x][1]][0], ch[ch[x][1]][1]), tag[ch[x][1]] ^= 1;
11             tag[x] = 0;
12         }
13     }
14     void update(int x) {

```



```

15         if (!isroot(x)) update(fa[x]);
16         pushdown(x);
17     }
18     void rotate(int x) {
19         int y = fa[x], z = fa[y], chx = getch(x), chy = getch(y);
20         fa[x] = z;
21         if (!isroot(y)) ch[z][chy] = x;
22         ch[y][chx] = ch[x][chx ^ 1];
23         fa[ch[x][chx ^ 1]] = y;
24         ch[x][chx ^ 1] = y;
25         fa[y] = x;
26     }
27     void splay(int x) {
28         update(x);
29         for (int f = fa[x]; f = fa[x], !isroot(x); rotate(x))
30             if (!isroot(f)) rotate(getch(x) == getch(f) ? f : x);
31     }
32     void access(int x) {
33         for (int f = 0; x; f = x, x = fa[x]) splay(x), ch[x][1] = f;
34     }
35     void makeroot(int x) {
36         access(x);
37         splay(x);
38         swap(ch[x][0], ch[x][1]);
39         tag[x] ^= 1;
40     }
41     int find(int x) {
42         access(x);
43         splay(x);
44         while (ch[x][0]) x = ch[x][0];
45         splay(x);
46         return x;
47     }
48 } st;
49
50 // 动态森林
51 namespace LCT {
52     void clear(int n) { for (int i = 0; i <= n; ++i) st.clear(i); }
53     bool isConnect(int u, int v) { return st.find(u) == st.find(v); }
54     bool add(int u, int v) {
55         if (isConnect(u, v)) return false;
56         st.makeroot(u), st.fa[u] = v;
57         return true;
58     }
59     // 无法判断是否存在直接相连的边
60     void del(int u, int v) {
61         st.makeroot(u);
62         st.access(v);
63         st.splay(v);
64         if (st.ch[v][0] == u && !st.ch[u][1]) st.ch[v][0] = st.fa[u] = 0;
65     }
66 }
67 using namespace LCT;

```

3.7 Tree

3.7.1 LCA

```

1 // const int maxn = 1e5 + 10;

```

```

2
3 // 普通倍增lca
4 int n, dep[maxn], fa[maxn][30];
5 vector<int> edge[maxn];
6
7 void dfs(int u, int pre) {
8     dep[u] = dep[pre] + 1, fa[u][0] = pre;
9     for(int i = 1; (1 << i) <= n; i++)
10         fa[u][i] = fa[fa[u][i - 1]][i - 1];
11     for(auto v : edge[u]) if(v != pre) dfs(v, u);
12 }
13
14 int LCA(int u, int v) {
15     if(dep[u] < dep[v]) swap(u, v);
16     int d = dep[u] - dep[v];
17     for(int i = 0; (1 << i) <= d; i++)
18         if((1 << i) & d) u = fa[u][i];
19     if(u == v) return u;
20     for(int i = 20; i >= 0; i--)
21         if(fa[u][i] != fa[v][i])
22             u = fa[u][i], v = fa[v][i];
23     return fa[u][0];
24 }

```

3.7.2 点分治

```

1 int n, k;
2
3 // 清零 head 和 tot
4 const int maxm = maxn * 2;
5 int ver[maxn], Next[maxn], head[maxn], edge[maxm];
6 int tot;
7 void addEdge(int u, int v, int w){
8     ver[++tot]=v;
9     Next[tot]=head[u];
10    head[u]=tot;
11    edge[tot]=w;
12 }
13
14 int sz[maxn], vis[maxn];
15 int rt, mxsz, has;
16
17 void getrt(int u, int pre) {
18     sz[u] = 1;
19     int mxnow = 0;
20     for (int i = head[u]; i; i = Next[i]) {
21         int v = ver[i];
22         if (v == pre || vis[v]) continue;
23         getrt(v, u);
24         sz[u] += sz[v];
25         mxnow = max(mxnow, sz[v]);
26     }
27     mxnow = max(mxnow, has - sz[u]);
28     if (mxnow < mxsz) {
29         mxsz = mxnow, rt = u;
30     }
31 }
32

```

```
33 int dl[maxn], r;
34 int val[maxn];
35
36 void getdis(int u, int pre) {
37     dl[r++] = val[u];
38     for (int i = head[u]; i; i = Next[i]) {
39         int v = ver[i];
40         if (v == pre || vis[v]) continue;
41         val[v] = val[u] + edge[i];
42         getdis(v, u);
43     }
44 }
45
46 ll cal(int u, int pre) {
47     r = 0;
48     val[u] = pre;
49     getdis(u, 0);
50     ll sum = 0;
51     sort(dl, dl + r);
52     r--;
53     int l = 0;
54     while (l < r) {
55         if (dl[l] + dl[r] > k) r--;
56         else sum += r - l, l++;
57     }
58     return sum;
59 }
60
61 ll res = 0;
62 void dfs(int u) {
63     res += cal(u, 0);
64     vis[u] = 1;
65     for (int i = head[u]; i; i = Next[i]) {
66         int v = ver[i];
67         if (vis[v]) continue;
68         res -= cal(v, edge[i]);
69         has = sz[v];
70         mxsz = 0x3f3f3f3f;
71         getrt(v, 0);
72         dfs(rt);
73     }
74 }
75
76 int main(int argc, char* argv[]) {
77     while (scanf("%d%d", &n, &k) != EOF && (n || k)) {
78         tot = 0; memset(head, 0, sizeof head);
79         memset(vis, 0, sizeof vis);
80         res = 0;
81         for (int i = 1, u, v, w; i < n; ++i) {
82             scanf("%d%d%d", &u, &v, &w);
83             addEdge(u, v, w);
84             addEdge(v, u, w);
85         }
86         mxsz = 0x3f3f3f3f;
87         has = n;
88         getrt(1, 0);
89         dfs(rt);
90         printf("%lld\n", res);
91     }
```

```

92     return 0;
93 }

```

3.8 Splay

```

1  /*
2  1. 插入x数
3  2. 删除x数(若有多个相同的数, 因只删除一个)
4  3. 查询x数的排名(若有多个相同的数, 因输出最小的排名)
5  4. 查询排名为x的数
6  5. 求x的前驱(前驱定义为小于x, 且最大的数)
7  6. 求x的后继(后继定义为大于x, 且最小的数)
8  */
9
10 const int N = 1e5 + 7;
11
12 struct Splay {
13     int ch[N][2], fa[N], val[N], cnt[N], size[N], tol, root;
14     inline bool chk(int x) {
15         return ch[fa[x]][1] == x;
16     }
17     inline void pushup(int x) {
18         size[x] = size[ch[x][0]] + size[ch[x][1]] + cnt[x];
19     }
20     void rotate(int x) {
21         int y = fa[x], z = fa[y], k = chk(x), w = ch[x][k ^ 1];
22         ch[y][k] = w; fa[w] = y;
23         ch[z][chk(y)] = x; fa[x] = z;
24         ch[x][k ^ 1] = y; fa[y] = x;
25         pushup(y); pushup(x);
26     }
27     void splay(int x, int goal = 0) {
28         while (fa[x] != goal) {
29             int y = fa[x], z = fa[y];
30             if (z != goal) {
31                 if (chk(x) == chk(y)) rotate(y);
32                 else rotate(x);
33             }
34             rotate(x);
35         }
36         if (!goal) root = x;
37     }
38     void insert(int x) {
39         int cur = root, p = 0;
40         while (cur && val[cur] != x) {
41             p = cur;
42             cur = ch[cur][x > val[cur]];
43         }
44         if (cur) {
45             cnt[cur]++;
46         } else {
47             cur = ++tol;
48             if (p) ch[p][x > val[p]] = cur;
49             ch[cur][0] = ch[cur][1] = 0;
50             fa[cur] = p; val[cur] = x;
51             cnt[cur] = size[cur] = 1;
52         }
53         splay(cur);

```

```
54     }
55     void find(int x) {
56         int cur = root;
57         while (ch[cur][x > val[cur]] && x != val[cur])
58             cur = ch[cur][x > val[cur]];
59         splay(cur);
60     }
61     int kth(int k) {
62         int cur = root;
63         while (1) {
64             if (ch[cur][0] && k <= size[ch[cur][0]])
65                 cur = ch[cur][0];
66             else if (k > size[ch[cur][0]] + cnt[cur])
67                 k -= size[ch[cur][0]] + cnt[cur], cur = ch[cur][1];
68             else
69                 break;
70         }
71         return cur;
72     }
73     int pre(int x) {
74         find(x);
75         if (val[root] < x) return root;
76         int cur = ch[root][0];
77         while (ch[cur][1]) cur = ch[cur][1];
78         return cur;
79     }
80     int succ(int x) {
81         find(x);
82         if (val[root] > x) return root;
83         int cur = ch[root][1];
84         while (ch[cur][0]) cur = ch[cur][0];
85         return cur;
86     }
87     void del(int x) {
88         int last = pre(x), nxt = succ(x);
89         splay(last); splay(nxt, last);
90         int del = ch[nxt][0];
91         if (cnt[del] > 1)
92             cnt[del]--, splay(del);
93         else
94             ch[nxt][0] = 0;
95     }
96     int getrk(int x) {
97         find(x);
98         return size[ch[root][0]];
99     }
100 } splay;
101
102 int n;
103
104 int main() {
105     //freopen("in.txt", "r", stdin);
106     splay.insert(0x3f3f3f3f);
107     splay.insert(0xcfcfcfcf);
108     read(n);
109     while (n--) {
110         int opt, x;
111         read(opt, x);
112         if (opt == 1) splay.insert(x);
```

```

113     else if (opt == 2) splay.del(x);
114     else if (opt == 3) print(splay.getrk(x));
115     else if (opt == 4) print(splay.val[splay.kth(x + 1)]);
116     else if (opt == 5) print(splay.val[splay.pre(x)]);
117     else print(splay.val[splay.succ(x)]);
118 }
119 flush();
120 return 0;
121 }

```

3.9 Others

3.9.1 BITinNM

```

1 struct Fenwick_Tree {
2 #define type int
3     type bit[maxn][maxn];
4     int n, m;
5     void init(int _n, int _m) {
6         n = _n;
7         m = _m;
8         mem(bit, 0);
9     }
10    int lowbit(int x) { return x & (-x); }
11    void update(int x, int y, type v) {
12        int i, j;
13        for (i = x; i <= n; i += lowbit(i)) {
14            for (j = y; j <= m; j += lowbit(j)) {
15                bit[i][j] += v;
16            }
17        }
18    }
19    type get(int x, int y) {
20        type i, j, res = 0;
21        for (i = x; i > 0; i -= lowbit(i)) {
22            for (j = y; j > 0; j -= lowbit(j)) {
23                res += bit[i][j];
24            }
25        }
26        return res;
27    }
28    type query(int x1, int x2, int y1, int y2) {
29        x1--;
30        y1--;
31        return get(x2, y2) - get(x1, y2) - get(x2, y1) + get(x1, y1);
32    }
33 #undef type
34 } tr;
35
36 // 二维区间前缀和写法 (非树状数组)
37 inline void range_add(int xa, int ya, int xb, int yb) { add(xa, ya, 1), add(xa, yb + 1,
38     -1), add(xb + 1, ya, -1), add(xb + 1, yb + 1, 1); }
39 inline ll range_ask(int xa, int ya, int xb, int yb){ return ask(xb, yb) - ask(xb, ya -
40     1) - ask(xa - 1, yb) + ask(xa - 1, ya - 1); }
41 inline void build() {
42     // 预处理出每个点的单点值
43     for (int i = 1; i < n + 5; ++i) {
44         for (int j = 1; j < m + 5; ++j) {
45             st[i][j] += st[i - 1][j] + st[i][j - 1] - st[i - 1][j - 1];

```

```

44     }
45 }
46 // 再求一次处理出每个点的前缀和
47 for (int i = 1; i < n + 5; ++i) {
48     for (int j = 1; j < m + 5; ++j) {
49         if (st[i][j] > 1) st[i][j] = 1;
50         st[i][j] += st[i - 1][j] + st[i][j - 1] - st[i - 1][j - 1];
51     }
52 }
53 }
54
55 // 二维树状数组区间加与求和
56 ll t1[maxn][maxn], t2[maxn][maxn], t3[maxn][maxn], t4[maxn][maxn];
57 void add(ll x, ll y, ll z){
58     for(int X = x; X <= n; X += X & -X)
59         for(int Y = y; Y <= m; Y += Y & -Y){
60             t1[X][Y] += z;
61             t2[X][Y] += z * x;
62             t3[X][Y] += z * y;
63             t4[X][Y] += z * x * y;
64         }
65 }
66 ll ask(ll x, ll y){
67     ll res = 0;
68     for(int i = x; i; i -= i & -i)
69         for(int j = y; j; j -= j & -j)
70             res += (x + 1) * (y + 1) * t1[i][j]
71                 - (y + 1) * t2[i][j]
72                 - (x + 1) * t3[i][j]
73                 + t4[i][j];
74     return res;
75 }
76
77 // 区间加, 询问单点: 直接维护前缀差分数组, 求单点=普通求前缀和

```

3.9.2 静态区间 k 大划分树

```

1 // const int maxn = 100010;
2 int tree[20][maxn];
3 // 读入sorted并排序, 赋值给tree的第0层
4 int sorted[maxn];
5 int toleft[20][maxn];
6 // 保存左子树的和
7 // ll sum[20][maxn];
8
9 // 1, n, 0
10 void build(int l, int r, int dep) {
11     if (l == r) return;
12     // sum[dep][0] = 0;
13     toleft[dep][0] = 0;
14     int mid = l + r >> 1;
15     int same = mid - l + 1;
16     for (int i = l; i <= r; ++i) {
17         if (tree[dep][i] < sorted[mid]) same--;
18     }
19     int lpos = l, rpos = mid + 1;
20     for (int i = l; i <= r; ++i) {
21         // sum[dep][i] = sum[dep][i - 1];

```

```

22     if (tree[dep][i] < sorted[mid]) {
23         // sum[dep][i] += tree[dep][i];
24         tree[dep + 1][lpos++] = tree[dep][i];
25     }
26     else if (tree[dep][i] == sorted[mid] && same > 0) {
27         // sum[dep][i] += tree[dep][i];
28         tree[dep + 1][lpos++] = tree[dep][i];
29         same--;
30     } else tree[dep + 1][rpos++] = tree[dep][i];
31     toleft[dep][i] = toleft[dep][l - 1] + lpos - 1;
32 }
33 build(l, mid, dep + 1);
34 build(mid + 1, r, dep + 1);
35 }
36
37 //(1~k-1)的数的和, 注意每次查询前初始化
38 // ll res = 0;
39
40 // L = 1, R = n,  dep = 0, l,r是查询区间
41 int query(int L, int R, int dep, int k) {
42     if (l == r) return tree[dep][l];
43     int mid = (L + R) >> 1;
44     int cnt = toleft[dep][r] - toleft[dep][l - 1];
45     if (cnt >= k) {
46         int newl = L + toleft[dep][l - 1] - toleft[dep][L - 1];
47         int newr = newl + cnt - 1;
48         return query(L, mid, newl, newr, dep + 1, k);
49     } else {
50         int newr = r + toleft[dep][R] - toleft[dep][r];
51         int newl = newr - (r - l - cnt);
52         // res += sum[dep][r] - sum[dep][l - 1];
53         return query(mid + 1, R, newl, newr, dep + 1, k - cnt);
54     }
55 }
56
57
58 scan(n), scan(m);
59 for (int i = 1; i <= n; ++i) {
60     scan(sorted[i]);
61     tree[0][i] = sorted[i];
62 }
63 sort(sorted + 1, sorted + 1 + n);
64 build(1, n, 0);
65 int l, r, k;
66 while (m--) {
67     scan(l), scan(r), scan(k);
68     printf("%d\n", query(1, n, l, r, 0, k));
69 }

```

3.9.3 二叉堆

```

1  template<class T = int>
2  struct BinaryHeap {
3      vector<T> Heap;
4      int sz;
5      BinaryHeap(int s = 0) { Heap.resize(s), sz = 0; }
6      void init(int s) { Heap.resize(s), sz = 0; }
7      bool less(T a, T b) { return a > b; }

```



```
8     void up(int p) {
9         while (p > 1) {
10             if (less(Heap[p / 2], Heap[p])) {
11                 swap(Heap[p], Heap[p / 2]);
12                 p /= 2;
13             } else break;
14         }
15     }
16     void down(int p) {
17         int s = p * 2;
18         while (s <= sz) {
19             if (s < sz && less(Heap[s], Heap[s + 1])) s++;
20             if (less(Heap[p], Heap[s])) {
21                 swap(Heap[s], Heap[p]);
22                 p = s, s = p * 2;
23             } else break;
24         }
25     }
26     void insert(int val) {
27         Heap[++sz] = val;
28         up(sz);
29     }
30     void removeTop() {
31         Heap[1] = Heap[sz--];
32         down(1);
33     }
34     void remove(int pos) {
35         Heap[pos] = Heap[sz--];
36         up(pos), down(pos);
37     }
38     int getTop() {
39         assert(sz > 0);
40         return Heap[1];
41     }
42 };
```

4 String

4.1 KMP

4.1.1 KMP

```

1 // nxt[0]表示失配到完全不匹配
2 int nxt[maxm];
3
4 void getNext(char *s, int len) {
5     int i = 0, j = -1;
6     nxt[i] = j;
7     while (i < len) {
8         if (j == -1 || s[i] == s[j]) nxt[++i] = ++j;
9         else j = nxt[j];
10    }
11 }
12
13 // a为原串, b为模式串, 下标从0开始, 找第一个出现模式串的位置 (起点为1), 找不到返回-1
14 int KMP(char *a, char *b, int n, int m) {
15     getNext(b, m);
16     int i = 0, j = 0;
17     while (i < n && j < m) {
18         if (j == -1 || a[i] == b[j]) ++i, ++j;
19         else j = nxt[j];
20    }
21    return j == m ? i - m + 1 : -1;
22 }

```

4.1.2 exKMP

```

1 const int maxn = 1e5 + 10;
2 int nex[maxn], extend[maxn];
3
4 //预处理计算Next数组
5 void getNext(char *str)
6 {
7     int i = 0, j, po, len = strlen(str);
8     nex[0] = len;    //初始化nex[0]
9     while (str[i] == str[i + 1] && i + 1 < len) i++;    //计算nex[1]
10    nex[1] = i;
11    po = 1;    //初始化po的位置
12    for (int i = 2; i < len; i++)
13    {
14        if (nex[i - po] + i < nex[po] + po) //第一种情况, 可以直接得到nex[i]的值
15            nex[i] = nex[i - po];
16        else //第二种情况, 要继续匹配才能得到nex[i]的值
17        {
18            j = nex[po] + po - i;
19            if (j < 0) j = 0;    //如果i>po+nex[po],则要从头开始匹配
20            while (i + j < len && str[j] == str[j + i]) j++;
21            nex[i] = j;
22            po = i;    //更新po的位置
23        }
24    }
25 }
26
27 void EXKMP(char *s1, char *s2)
28 {

```

```
29     int i = 0, j, po, len = strlen(s1), l2 = strlen(s2);
30     getNext(s2);
31     while (s1[i] == s2[i] && i < l2 && i < len) i++;
32     extend[0] = i;
33     po = 0;
34     for (int i = 1; i < len; i++)
35     {
36         if (nex[i - po] + i < extend[po] + po)
37             extend[i] = nex[i - po];
38         else
39         {
40             j = extend[po] + po - i;
41             if (j < 0) j = 0;
42             while (i + j < len && j < l2 && s1[j + i] == s2[j]) j++;
43             extend[i] = j;
44             po = i;
45         }
46     }
47 }
```

4.2 Trie

4.2.1 Trie

```
1  const int maxn = 2e6 + 10;
2
3  int trie[maxn][30], tot;
4  bool flag[maxn];
5
6  void insert_ch(char *str)
7  {
8      int len = strlen(str);
9      int root = 0;
10     for (int i = 0; i < len; i++)
11     {
12         int id = str[i] - 'a';
13         if (!trie[root][id]) trie[root][id] = ++tot;
14         root = trie[root][id];
15     }
16     flag[root] = true;
17 }
18
19 bool find_ch(char *str)
20 {
21     int len = strlen(str);
22     int root = 0;
23     for (int i = 0; i < len; i++)
24     {
25         int id = str[i] - 'a';
26         if (!trie[root][id]) return false;
27         root = trie[root][id];
28     }
29     return true;
30 }
```

4.2.2 Persistence Trie

```
1  const int maxn = 1e5 + 10;
2
3  int a[maxn], rt[maxn], n;
4
5  struct Trie
6  {
7      int tot;
8      int child[maxn * 32][2], sum[maxn * 32];
9      int insert(int x, int val)
10     {
11         int tmp, y;
12         tmp = y = ++tot;
13         for(int i = 30; i >= 0; --i)
14             {
15                 child[y][0] = child[x][0];
16                 child[y][1] = child[x][1];
17                 sum[y] = sum[x] + 1;
18                 int t = val >> i & 1;
19                 x = child[x][t];
20                 child[y][t] = ++tot;
21                 y = child[y][t];
22             }
23         sum[y] = sum[x] + 1;
24         return tmp;
25     }
26     int query(int l, int r, int val)
27     {
28         int tmp = 0;
29         for(int i = 30; i >= 0; --i)
30             {
31                 int t = val >> i & 1;
32                 if(sum[child[r][t^1]] - sum[child[l][t^1]]) tmp += (1<<i), r = child[r][t^1], l = child[l][t^1];
33                 else r = child[r][t], l = child[l][t];
34             }
35         return tmp;
36     }
37 }trie;
```

4.2.3 01Trie

```
1  struct Trie {
2      int tree[maxn*20][2], tot;
3      int flag[maxn*20];
4
5      void insert_ch(int x) {
6          int root = 0;
7          flag[0]++;
8          for (int i = 30; i >= 0; --i) {
9              int id = (x >> i) & 1;
10             if (!tree[root][id]) {
11                 tree[root][id] = ++tot;
12                 tree[tree[root][id]][0] = tree[tree[root][id]][1] = 0;
13                 flag[tree[root][id]] = flag[tree[tree[root][id]][0]] = flag[tree[tree[root][id]][1]] = 0;
14             }
15             root = tree[root][id];
16             flag[root]++;
17         }
18     }
```

```
17     }
18 }
19
20 void del(int x) {
21     int root = 0;
22     flag[0]--;
23     for (int i = 30; i >= 0; --i) {
24         int id = (x >> i) & 1;
25         assert(tree[root][id]);
26         if (flag[tree[root][id]] == 1) {
27             flag[tree[root][id]] = 0;
28             tree[root][id] = 0;
29             return;
30         }
31         root = tree[root][id];
32         flag[root]--;
33     }
34 }
35
36 int find_ch(int x, int flag = 0) { // flag 0 最小异或值, 1 最大异或值
37     int root = 0;
38     int res = 0;
39     for (int i = 30; i >= 0; --i) {
40         int id = ((x >> i) & 1);
41         if (flag) id = !id;
42         if (tree[root][id]) {
43             root = tree[root][id];
44             res = res << 1 | id;
45         } else {
46             root = tree[root][!id];
47             res = res << 1 | (!id);
48         }
49     }
50     return res;
51 }
52
53 void init() {
54     tree[0][0] = tree[0][1] = 0;
55     tot = 0;
56 }
57 };
```

4.3 Manachar

4.3.1 Manacher

```
1 const int maxn = 1e5 + 10;
2
3 char s[maxn];
4
5 char tmp[maxn << 1];
6 int Len[maxn << 1];
7
8 int init(char *str)
9 {
10     int len = strlen(str);
11     tmp[0] = '@';
12     for (int i = 1; i <= 2 * len; i += 2)
13     {
```

```
14         tmp[i] = '#';
15         tmp[i + 1] = str[i / 2];
16     }
17     tmp[2 * len + 1] = '#';
18     tmp[2 * len + 2] = '$';
19     tmp[2 * len + 3] = 0;
20     return 2 * len + 1;
21 }
22
23 int manacher(char *str)
24 {
25     int mx = 0, ans = 0, pos = 0;
26     int len = init(str);
27     for (int i = 1; i <= len; i++)
28     {
29         if (mx > i) Len[i] = min(mx - i, Len[2 * pos - i]);
30         else Len[i] = 1;
31         while (tmp[i - Len[i]] == tmp[i + Len[i]]) Len[i]++;
32         if (Len[i] + i > mx) mx = Len[i] + i, pos = i;
33     }
34 }
```

4.4 Aho-Corasick Automation

4.4.1 AC Automation

```
1 class AC_automation
2 {
3 public:
4     int trie[maxn][26], cnt;
5     int tag[maxn];
6     int fail[maxn], num[maxn], res[maxn], in[maxn], Map[maxn];
7
8     void init()
9     {
10         memset(trie, 0, sizeof trie);
11         memset(tag, 0, sizeof tag);
12         memset(fail, 0, sizeof fail);
13         cnt = 0;
14     }
15
16     void insert(char *str, int id)
17     {
18         int root = 0;
19         for (int i = 0; str[i]; i++)
20         {
21             int id = str[i] - 'a';
22             if (!trie[root][id]) trie[root][id] = ++cnt;
23             root = trie[root][id];
24         }
25         if (!tag[root]) tag[root] = id;
26         Map[id] = tag[root];
27     }
28
29     void build()
30     {
31         queue<int> que;
32         for (int i = 0; i < 26; i++) if (trie[0][i]) que.push(trie[0][i]);
33         while (!que.empty())
```

```

34     {
35         int k = que.front();
36         que.pop();
37         for (int i = 0; i < 26; i++)
38         {
39             if (trie[k][i])
40             {
41                 fail[trie[k][i]] = trie[fail[k]][i];
42                 que.push(trie[k][i]);
43                 in[fail[trie[k][i]]] ++;
44             } else trie[k][i] = trie[fail[k]][i];
45         }
46     }
47 }
48
49 void toposort()
50 {
51     queue<int> que;
52     for(int i = 1; i <= cnt; i ++) if(in[i] == 0) que.push(i);
53     while(!que.empty())
54     {
55         int u = que.front(); que.pop();
56         res[tag[u]] = num[u];
57         int v = fail[u]; in[v] --;
58         num[v] += num[u];
59         if(in[v] == 0) que.push(v);
60     }
61 }
62
63 void query(char *str, int n)
64 {
65     int u = 0, len = strlen(s);
66     for(int i = 0; i < len; i ++)
67         u = trie[u][str[i] - 'a'], num[u] ++;
68     toposort();
69     for(int i = 1; i <= n; i ++) printf("%d\n", res[Map[i]]);
70 }
71 } AC;

```

4.5 Suffix Array

4.5.1 Suffix Array

```

1 char s[maxn];
2 int sa[maxn], t[maxn], t2[maxn], c[maxn], n;
3
4 //build_sa(n + 1, 130), sa, height下标从1开始,rk下标从0开始
5 void build_sa(int n, int m)
6 {
7     int *x = t, *y = t2;
8     for(int i = 0; i < m; i++) c[i] = 0;
9     for(int i = 0; i < n; i++) c[x[i] = s[i]]++;
10    for(int i = 1; i < m; i++) c[i] += c[i - 1];
11    for(int i = n - 1; i >= 0; i--) sa[--c[x[i]]] = i;
12    for(int k = 1; k <= n; k <= 1)
13    {
14        int p = 0;
15        for(int i = n - k; i < n; i++) y[p++] = i;
16        for(int i = 0; i < n; i++) if(sa[i] >= k) y[p++] = sa[i] - k;

```

```

17     for(int i = 0; i < m; i++) c[i] = 0;
18     for(int i = 0; i < n; i++) c[x[y[i]]]++;
19     for(int i = 0; i < m; i++) c[i] += c[i - 1];
20     for(int i = n - 1; i >= 0; i--) sa[--c[x[y[i]]]] = y[i];
21     swap(x, y);
22     p = 1; x[sa[0]] = 0;
23     for(int i = 1; i < n; i++)
24         x[sa[i]] = y[sa[i - 1]] == y[sa[i]] && y[sa[i - 1] + k] == y[sa[i] + k] ? p
                - 1 : p++;
25     if(p >= n) break;
26     m = p;
27 }
28 }
29
30 int rk[maxn], height[maxn];
31
32 void getHeight()
33 {
34     for(int i = 1; i <= n; i++) rk[sa[i]] = i;
35     for(int i = 0, k = 0; i < n; i++)
36     {
37         if(k) k--;
38         int j = sa[rk[i] - 1];
39         while(s[i + k] == s[j + k]) k++;
40         height[rk[i]] = k;
41     }
42 }
43
44 int dp[maxn][20];
45
46 void RMQ()
47 {
48     for(int i = 1; i <= n; i++) dp[i][0] = height[i];
49     for(int j = 1; (1 << j) < maxn; j++)
50         for(int i = 1; i + (1 << j) - 1 <= n; i++)
51             dp[i][j] = min(dp[i][j - 1], dp[i + (1 << (j - 1))][j - 1]);
52 }
53
54 int query(int l, int r)
55 {
56     int k = 0;
57     while((1 << (k + 1)) <= r - l + 1) k++;
58     return min(dp[l][k], dp[r - (1 << k) + 1][k]);
59 }
60
61 int lcp(int x, int y)
62 {
63     x = rk[x], y = rk[y];
64     if(x > y) swap(x, y);
65     return query(x + 1, y);
66 }

```

4.5.2 SA badcw

```

1 namespace SA {
2     const int maxn = 2e5 + 10;
3     int t1[maxn], t2[maxn], c[maxn];
4     int Rank[maxn], height[maxn];

```



```

5   int RMQ[maxn];
6   int mm[maxn];
7   int sa[maxn];
8   int best[25][maxn];
9   bool cmp(int *r, int a, int b, int l) {
10      return r[a] == r[b] && r[a + l] == r[b + l];
11  }
12  void da(char str[], int sa[], int Rank[], int height[], int n, int m) {
13      n++;
14      int i, j, p, *x = t1, *y = t2;
15      for (i = 0; i < m; i++) c[i] = 0;
16      for (i = 0; i < n; i++) c[x[i] = str[i]]++;
17      for (i = 1; i < m; i++) c[i] += c[i - 1];
18      for (i = n - 1; i >= 0; i--) sa[--c[x[i]]] = i;
19      for (j = 1; j <= n; j <= 1) {
20          p = 0;
21          for (i = n - j; i < n; i++) y[p++] = i;
22          for (i = 0; i < n; i++) if (sa[i] >= j) y[p++] = sa[i] - j;
23          for (i = 0; i < m; i++) c[i] = 0;
24          for (i = 0; i < n; i++) c[x[y[i]]]++;
25          for (i = 1; i < m; i++) c[i] += c[i - 1];
26          for (i = n - 1; i >= 0; i--) sa[--c[x[y[i]]]] = y[i];
27          swap(x, y);
28          p = 1;
29          x[sa[0]] = 0;
30          for (i = 1; i < n; i++)
31              x[sa[i]] = cmp(y, sa[i - 1], sa[i], j) ? p - 1 : p++;
32          if (p >= n) break;
33          m = p;
34      }
35      int k = 0;
36      n--;
37      for (i = 0; i <= n; i++) Rank[sa[i]] = i;
38      for (i = 0; i < n; i++) {
39          if (k) k--;
40          j = sa[Rank[i] - 1];
41          while (str[i + k] == str[j + k]) k++;
42          height[Rank[i]] = k;
43      }
44  }
45  void initRMQ(int n) {
46      for (int i = 1; i <= n; i++)
47          RMQ[i] = height[i];
48      mm[0] = -1;
49      for (int i = 1; i <= n; i++)
50          mm[i] = ((i & (i - 1)) == 0) ? mm[i - 1] + 1 : mm[i - 1];
51      for (int i = 1; i <= n; i++) best[0][i] = i;
52      for (int i = 1; i <= mm[n]; i++)
53          for (int j = 1; j + (1 << i) - 1 <= n; j++) {
54              int a = best[i - 1][j];
55              int b = best[i - 1][j + (1 << (i - 1))];
56              if (RMQ[a] < RMQ[b]) best[i][j] = a;
57              else best[i][j] = b;
58          }
59  }
60  int askRMQ(int a, int b) {
61      int t;
62      t = mm[b - a + 1];
63      b -= (1 << t) - 1;

```

```

64     a = best[t][a];
65     b = best[t][b];
66     return RMQ[a] < RMQ[b] ? a : b;
67 }
68 int lcp(int a, int b) {
69     a = Rank[a];
70     b = Rank[b];
71     if (a > b) swap(a, b);
72     //cout << askRMQ(a + 1, b) << endl;
73     return height[askRMQ(a + 1, b)];
74 }
75 void preprocess(char *str, int n, int m) {
76     da(str, sa, Rank, height, n, m);
77     initRMQ(n);
78 }
79 }

```

4.6 PalindromicTree

4.6.1 PalindromicTree

```

1  const int maxn = 2e6+6;
2  const int N = 26;
3  const int mod = 51123987;
4
5  struct Palindromic_Tree {
6  //     vector<pair<int, int> > next[maxn];
7      int next[maxn][N]; //next 指针, next 指针和字典树类似, 指向的串为当前串两端加上同一个
                          // 字符构成
8      int fail[maxn]; //fail 指针, 失配后跳转到 fail 指针指向的节点
9      int cnt[maxn]; //表示节点 i 表示的本质不同的串的个数 (建树时求出的不是完全的, 最后
                      // count() 函数跑一遍以后才是正确的)
10     int num[maxn]; //表示以节点 i 表示的最长回文串的最右端点为回文串结尾的回文串个数
11     int len[maxn]; //len[i] 表示节点 i 表示的回文串的长度 (一个节点表示一个回文串)
12     int S[maxn]; //存放添加的字符
13     int last; //指向新添加一个字母后所形成的最长回文串表示的节点。
14     int n; //表示添加的字符个数。
15     int p; //表示添加的节点个数。
16     //0 向前加, 1 向后加字符
17     //int last[2];
18     //int lpos, rpos;
19
20     int newnode(int l) { //新建节点
21 //         next[p].clear();
22         for (int i = 0; i < N; ++i) next[p][i] = 0;
23         cnt[p] = 0;
24         num[p] = 0;
25         len[p] = 1;
26         return p++;
27     }
28
29     void init() { //初始化
30         n = last = p = 0;
31         newnode(0);
32         newnode(-1);
33         S[n] = -1; //开头放一个字符集中没有的字符, 减少特判
34         fail[0] = 1;
35         // lpos 为字符串最大长度
36         // last[0] = last[1] = 0;

```

```

37         // lpos = 100000, rpos = lpos - 1;
38         // S[lpos - 1] = S[rpos + 1] = -1;
39     }
40
41     int get_fail(int x) { //和KMP一样, 失配后找一个尽量最长的
42         // op 0 向前, 1 向后
43         // if (op == 0) while (S[lpos + len[x] + 1] != S[lpos]) x = fail[x];
44         // else while(S[rpos - len[x] - 1] != S[rpos]) x = fail[x];
45         while (S[n - len[x] - 1] != S[n]) x = fail[x];
46         return x;
47     }
48
49     // int find(int u, int c) {
50     //     vector<pair<int, int> > & x = next[u];
51     //     int sz = x.size();
52     //     for(int i = 0; i < sz; ++i) {
53     //         if(x[i].first == c) return x[i].second;
54     //     }
55     //     return 0;
56     // }
57
58     int add(int c) {
59         // 注意清空左右字符
60         // if (op == 0) S[--lpos] = c, S[lpos - 1] = -1;
61         // else S[++rpos] = c, S[rpos + 1] = -1;
62         S[++n] = c;
63         int cur = get_fail(last); //通过上一个回文串找这个回文串的匹配位置
64         // int x = find(cur, c);
65         // if (!x) {
66         if (!next[cur][c]) { //如果这个回文串没有出现过, 说明出现了一个新的本质不同的回
            文串
67             int now = newnode(len[cur] + 2); //新建节点
68             // x = now;
69             // fail[now] = find(get_fail(fail[cur]), c);
70             // next[cur].emplace_back(make_pair(c, now));
71             fail[now] = next[get_fail(fail[cur])][c]; //和AC自动机一样建立fail指针, 以便
            失配后跳转
72             next[cur][c] = now;
73             num[now] = num[fail[now]] + 1;
74         }
75         // last = x;
76         // 修改最终长度
77         // if (len[last[op]] == rpos - lpos + 1) last[op ^ 1] = last[op];
78         last = next[cur][c];
79         cnt[last]++;
80         return num[last];
81     }
82
83     void count() {
84         for (int i = p - 1; i >= 0; --i) cnt[fail[i]] += cnt[i];
85         //父亲累加儿子的cnt, 因为如果fail[v]=u, 则u一定是v的子回文串!
86     }
87 } solve;
88
89 char s[maxn];
90
91 // 求相交回文串数量
92 ll a[maxn], b[maxn];
93 int main() {

```

```

94     solve.init();
95     int n;
96     scanf("%d", &n);
97     scanf("%s", s);
98     for (int i = 0; i < n; ++i) a[i] = solve.add(s[i] - 'a');
99     solve.init();
100    for (int i = n - 1; i >= 0; --i) b[i] = (b[i + 1] + solve.add(s[i] - 'a')) % mod;
101    ll res = (b[0] * (b[0] - 1) / 2) % mod;
102    for (int i = 0; i < n; ++i) res = ((res - (a[i] * b[i + 1]) + mod) % mod) % mod;
103    printf("%lld\n", res);
104    return 0;
105 }

```

4.7 Hash

4.7.1 hash

```

1 // hash常用素数
2 // 61, 83, 113, 151, 211
3 // 91815541, 38734667, 68861641
4 // 917120411, 687840301, 386910137, 515880193
5 // 1222827239, 1610612741
6
7 typedef unsigned long long ull;
8 struct mhash {
9     // 自然溢出无模数 805306457
10    ull base[maxn];
11    ull hash_index[maxn];
12    ull seed; //31, 131
13    void inithash(ull seedt = 31) {
14        base[0] = 1;
15        seed = seedt;
16        for (int i = 1; i < maxn; ++i) base[i] = base[i - 1] * seed;
17    }
18    void H(char *p, int n) { // from 1 to n
19        hash_index[0] = 0;
20        for (int i = 1; i <= n; ++i) hash_index[i] = hash_index[i - 1] * seed + p[i] -
21            'a';
22    }
23    ull gethash(int s, int e) {
24        return hash_index[e] - hash_index[s - 1] * base[e - s + 1];
25    }
26 };
27 // 26个素数, 解决加法hash
28 int prime[] = {34183, 13513, 152993, 13591, 19687, 350869, 111187, 766091, 769297,
29     633469, 752273, 298651, 617191, 880421, 136067,
30     1408397, 726899, 458921, 2133701, 2599847, 2730947, 4696343, 10267237,
31     18941059, 34078909, 69208409};

```

4.7.2 doubleHash

```

1 namespace Hash{
2
3     template<class __A, class __B>
4     class Hash{
5     private:
6         static const int size=2000000;

```

```

7     __B *hash; __A *0; int sz;
8 public:
9     Hash(int hash_size=size){ sz=hash_size;
10         hash=(__B *)malloc(sizeof(__B)*sz);
11         0=(__A *)malloc(sizeof(__A)*sz);
12         memset(0,0xff,sizeof(__A)*sz);
13     }~Hash(){free(0);free(hash);}
14     __B &operator [](const __A &_0){
15         int loc=_0%sz;
16         while(~0[loc]&&0[loc]!=_0){
17             ++loc;
18             if(loc>sz)loc=0;
19         }if(!~0[loc])0[loc]=_0;
20         return hash[loc];
21     }
22     void clear(){memset(0,0xff,sizeof(__A)*sz);}
23 };
24
25 struct StringDoubleHashResult{
26     int32_t *H1,*H2,c_len,len;
27     StringDoubleHashResult(int32_t sz=0){
28         len=sz; c_len=0; //cur_len;
29         if(len<=0){
30             H1=H2=0;
31             return;
32         }
33         H1=(int32_t *)malloc(sizeof(int32_t)*sz);
34         H2=(int32_t *)malloc(sizeof(int32_t)*sz);
35     }
36     ~StringDoubleHashResult(){}
37     void clear(){free(H1);free(H2);len=0;H1=H2=0;}
38     void resize(int new_len){
39         int32_t *T1=(int32_t *)malloc(sizeof(int32_t)*new_len);
40         int32_t *T2=(int32_t *)malloc(sizeof(int32_t)*new_len);
41         for(int i=0;i<c_len;++i)T1[i]=H1[i],T2[i]=H2[i];
42         free(H1);free(H2); H1=T1; H2=T2; len=new_len;
43     }
44     void erase(int ers_len){//erase suffix
45         c_len-=ers_len;
46         if(c_len<0)c_len=0;
47     }
48     //erase prefix not better than reculc
49 };
50
51 namespace hash_random{
52     const int mod_tot=5;
53     const int mod[]={1000000009,1000000007,998244353,917120411,515880193};
54 };
55
56 class StringDoubleHash{
57 private:
58     static const int enable_random=1;
59     int32_t sz,HA1,HA2;
60     long long B,C;
61     int32_t *H1,*H2;
62 public:
63     StringDoubleHash(int32_t SZ=2e6+5,int32_t ha1=-1,int32_t ha2=-1,int32_t b=-1,
64         int32_t c=-1){
65         sz=SZ;

```

```

65         if(enable_random){
66             std::mt19937 rnd(time(0)+19990630);
67             int z1= rnd() % hash_random::mod_tot;
68             int z2= (z1 +rnd()%(hash_random::mod_tot - 1) + 1) % hash_random::
                mod_tot;
69             if(ha1<0)ha1=hash_random::mod[z1];
70             if(ha2<0)ha2=hash_random::mod[z2];
71             if(b<0)b=rnd()%114514+23333;
72             if(c<0)c=rnd()%1919810+23333;
73         } else {
74             if(ha1<0)ha1=1e9+7;
75             if(ha2<0)ha2=1e9+9;
76             if(b<0)b=114514;
77             if(c<0)c=1919810;
78         }
79         HA1=ha1; HA2=ha2; B=b; C=c;
80         //cerr<<HA1<<" "<<HA2<<" "<<B<<" "<<C<<endl;
81         H1=(int32_t *)malloc(sizeof(int32_t)*sz);
82         H2=(int32_t *)malloc(sizeof(int32_t)*sz);
83         init_hash_val();
84     }
85     ~StringDoubleHash(){free(H1);free(H2);}
86     void init_hash_val(){
87         H1[0]=H2[0]=1;
88         for(int32_t i=1;i<sz;++i){
89             H1[i]=(H1[i-1]*B)%HA1;
90             H2[i]=(H2[i-1]*B)%HA2;
91         }
92     }
93     template <class _Tp>
94     StringDoubleHashResult culc_hash(const _Tp &s,int32_t len,int32_t tot_len=-1){
95         if(tot_len<0)tot_len=len;
96         StringDoubleHashResult R(tot_len);
97         if(len<=0)return R;
98         R.H1[0]=(s[0]+C)%HA1;
99         R.H2[0]=(s[0]+C)%HA2;
100        for(int32_t i=1;i<len;++i){
101            R.H1[i]=(R.H1[i-1]*B+s[i]+C)%HA1;
102            R.H2[i]=(R.H2[i-1]*B+s[i]+C)%HA2;
103        }
104        R.c_len=len;
105        return R;
106    }
107    // s is the char* first, len is the append length
108    template <class _Tp>
109    void append(StringDoubleHashResult &R,const _Tp &s,int32_t len){
110        if(len<=0)return;
111        int t_len=R.len;
112        while(R.c_len+len>t_len)t_len<=1;
113        if(t_len>R.len)R.resize(t_len);
114        for(int32_t i=R.c_len;i<R.c_len+len;++i){
115            if(i==0){
116                R.H1[i]=(s[i-R.c_len]+C)%HA1;
117                R.H2[i]=(s[i-R.c_len]+C)%HA2;
118            } else {
119                R.H1[i]=(R.H1[i-1]*B+s[i-R.c_len]+C)%HA1;
120                R.H2[i]=(R.H2[i-1]*B+s[i-R.c_len]+C)%HA2;
121            }
122        }

```

```

123         R.c_len+=len;
124     }
125     void append(StringDoubleHashResult &R, char s){
126         int t_len=R.len;
127         while(R.c_len+1>t_len)t_len<=1;
128         if(t_len>R.len)R.resize(t_len);
129         for(int32_t i=R.c_len;i<R.c_len+1;++i){
130             if(i==0){
131                 R.H1[i]=(s+C)%HA1;
132                 R.H2[i]=(s+C)%HA2;
133             } else {
134                 R.H1[i]=(R.H1[i-1]*B+s+C)%HA1;
135                 R.H2[i]=(R.H2[i-1]*B+s+C)%HA2;
136             }
137         }
138         R.c_len+=1;
139     }
140     //return hash [l,r)
141     ll gethash(const StringDoubleHashResult &R, int32_t l,int32_t r){
142         if(l>r||l<0||r-->R.c_len)return -1;//fail
143         ll v1=l>0?R.H1[l-1]*(long long)H1[r-l+1]%HA1:0;
144         ll v2=l>0?R.H2[l-1]*(long long)H2[r-l+1]%HA2:0;
145         v1=R.H1[r]-v1; v2=R.H2[r]-v2;
146         if(v1<0)v1+=HA1; if(v2<0)v2+=HA2;
147         return v1<<32|v2;
148     }
149     //merge two hashes as one(s1+s2), but need s2's length
150     ll merge_hash(const long long &hs1,const long long &hs2,int lenr){
151         int32_t m1=hs1>>32,m2=hs1&0xffffffffLL;
152         int32_t m3=hs2>>32,m4=hs2&0xffffffffLL;
153         m1=m1*(long long)H1[lenr]%HA1+m3;
154         if(m1>=HA1)m1-=HA1;
155         m2=m2*(long long)H2[lenr]%HA2+m4;
156         if(m2>=HA2)m2-=HA2;
157         return (long long)m1<<32|m2;
158     }
159 };
160 };

```

4.7.3 二维 hash

```

1  #define ull unsigned long long
2  const int maxn = 1005;
3  ull hs[maxn][maxn];
4  char a[maxn][maxn];
5  int n, m;
6  ull base1 = 131, base2 = 13331;
7  ull pwb1[maxn] = {1}, pwb2[maxn] = {1};
8
9  void init() {
10     for (int i = 1; i < maxn; ++i) {
11         pwb1[i] = pwb1[i - 1] * base1;
12         pwb2[i] = pwb2[i - 1] * base2;
13     }
14 }
15
16 void Hash() {
17     for(int i=1;i<=n;i++)

```

```

18         for(int j=1;j<=m;j++)
19             hs[i][j]=hs[i][j-1]*base1+a[i][j] - 'a';
20     for(int i=1;i<=n;i++)
21         for(int j=1;j<=m;j++)
22             hs[i][j]+=hs[i-1][j]*base2;
23 }
24
25 // 右下角(i,j), 行列长度n,m
26 ull getHs(int i, int j, int lenn, int lenm) {
27     return hs[i][j] - hs[i - lenn][j] * pwb2[lenn] -
28         hs[i][j - lenm] * pwb1[lenm] +
29         hs[i - lenn][j - lenm] * pwb2[lenn] * pwb1[lenm];
30 }

```

4.7.4 树 hash 同构

```

1 // n=1e5的话base开2e6+9, 可以输出看到top不比n小即可
2 const int base = 2e6+9;
3 // vis大小要开到素数大小, turn表示当前树的编号, p是预处理数组
4 int vis[base + 1], top, turn, p[base + 1];
5 // 程序开头调用一次
6 void init() {
7     top = 0;
8     for (int i = 2; i <= base; ++i) {
9         if (!vis[i]) {
10             p[++top] = i;
11         }
12         for (int j = 1; j <= top && i * p[j] <= base; ++j) {
13             vis[i * p[j]] = 1;
14             if (i % p[j] == 0) break;
15         }
16     }
17     assert(top >= maxn);
18 }
19
20 vector<int> edge[maxn];
21 // h[x]表示x这棵子树的hash值, g[x]表示以x为根的hash值
22 int h[maxn], g[maxn], sz[maxn];
23
24 struct TreeHash {
25     int n;
26     // 如果树比较多, 在类内部开edge可能会炸内存, 可以改到外面做前向星
27     // 除了hs是答案其他都可以改到外部, 只有edge需要清零
28     // vector<int> edge[maxn];
29     // int h[maxn], g[maxn], sz[maxn];
30     vector<int> hs;
31
32     void init(int n_ = 0) {
33         n = n_;
34         hs.clear();
35     }
36
37     void dfs1(int u, int pre) {
38         sz[u] = 1;
39         h[u] = 1;
40         for (auto v : edge[u]) {
41             if (v == pre) continue;
42             dfs1(v, u);

```



```

43         h[u] = (h[u] + 1ll * h[v] * p[sz[v]] % mod) % mod;
44         sz[u] += sz[v];
45     }
46 }
47
48 void dfs2(int u, int pre, int V, int needres = 1) {
49     g[u] = (h[u] + 1ll * V * p[n - sz[u]] % mod) % mod;
50     if (needres) hs.push_back(g[u]);
51     for (auto v : edge[u]) {
52         if (v == pre) continue;
53         dfs2(v, u, (g[u] - 1ll * h[v] * p[sz[v]] % mod + mod) % mod);
54     }
55 }
56
57 void work(int needres = 1) {
58     // 无根树选一个不存在的点当pre即可, 当多棵无根树判重时需要sort
59     dfs1(1, 0);
60     dfs2(1, 0, 0, needres);
61     sort(hs.begin(), hs.end());
62 }
63 };
64
65 // 获取删掉某叶子节点后以与该叶子节点相邻点开头的hash值
66 // int res = (hs[edge[i][0]] - 2 + mod) % mod;

```

4.8 Suffix Automation

4.8.1 SAM

```

1  const int maxn = 2e4 + 10;
2
3  struct SuffixAutomation
4  {
5      int last, cnt;
6      int ch[maxn << 1][26], fa[maxn << 1], len[maxn << 1], pos[maxn << 1];
7      int sz[maxn << 1], a[maxn << 1], c[maxn << 1];
8
9      void init()
10     {
11         last = cnt = 1;
12         memset(ch[1], 0, sizeof ch[1]);
13         fa[1] = len[1] = 0;
14         a[1] = c[1] = 0;
15     }
16
17     int inline newnode(int idx)
18     {
19         ++cnt;
20         memset(ch[cnt], 0, sizeof ch[cnt]);
21         fa[cnt] = len[cnt] = sz[cnt] = a[cnt] = c[cnt] = 0;
22         pos[cnt] = idx;
23         return cnt;
24     }
25
26     void ins(int c)
27     {
28         int p = last, np = newnode(pos[last] + 1);
29         last = np, len[np] = len[p] + 1;
30         for(;; p && !ch[p][c]; p = fa[p]) ch[p][c] = np;

```

```

31     if(!p) fa[np] = 1;
32     else
33     {
34         int q = ch[p][c];
35         if(len[p] + 1 == len[q]) fa[np] = q;
36         else
37         {
38             int nq = newnode(pos[p] + 1);
39             len[nq] = len[p] + 1;
40             memcpy(ch[nq], ch[q], sizeof ch[q]);
41             fa[nq] = fa[q], fa[q] = fa[np] = nq;
42             for(; ch[p][c] == q; p = fa[p]) ch[p][c] = nq;
43         }
44     }
45     sz[np] = 1;
46 }
47
48 int solve(int n)
49 {
50     /* 求两个串的LCS:
51        对一个字符串建立SAM, 记录一个当前匹配的长度Len和当前节点v, 枚举另一个字符串
52        的每个字符;
53        如果p有字符v的转移边出边, 则使Len加一, 并使p转移到出边指向的节点上;
54        否则不断向父节点上跳, 直到当前节点有字符p的转移出边, 或者跳到根节点;
55    */
56     int p = 1, ans = 0, now_len = 0;
57     for(int i = 0; s2[i]; i++)
58     {
59         if(ch[p][s2[i] - 'a']) p = ch[p][s2[i] - 'a'], now_len++;
60         else
61         {
62             for(; p && !ch[p][s2[i] - 'a']; p = fa[p]);
63             if(p == 0) now_len = 0, p = 1;
64             else now_len = len[p] + 1, p = ch[p][s2[i] - 'a'];
65         }
66         ans = max(now_len, ans);
67     }
68
69     void Toposort()
70     {
71         long long ans = 0;
72         for(int i = 1; i <= cnt; i++) c[len[i]]++;
73         for(int i = 1; i <= cnt; i++) c[i] += c[i - 1];
74         for(int i = 1; i <= cnt; i++) a[c[len[i]]--] = i;
75         for(int i = cnt; i; i--) sz[fa[a[i]]] += sz[a[i]];
76     }
77 }sam;

```

4.9 Others

4.9.1 最小表示法

```

1 // 0起始
2 int Gao(char a[], int len) {
3     int i = 0, j = 1, k = 0;
4     while (i < len && j < len && k < len) {
5         int cmp = a[(j + k) % len] - a[(i + k) % len];
6         if (cmp == 0) k++;

```

```
7         else {
8             if (cmp > 0) j += k + 1;
9             else i += k + 1;
10            if (i == j) j ++;
11            k = 0;
12        }
13    }
14    return min(i, j);
15 }
```

4.9.2 Lyndon

```
1 void duval(char s[])
2 {
3     int n = strlen(s), i = 0;
4     vector<int> res; //Lyndon串的开头
5     while(i < n)
6     {
7         int j = i, k = i + 1;
8         while(j < n && s[j] <= s[k])
9         {
10             if(s[j] < s[k]) j = i;
11             else j ++;
12             k ++;
13         }
14         while(i <= j)
15         {
16             res.push_back(i);
17             i += k - j;
18         }
19     }
20 }
```

5 dp

5.1 BitDP

5.1.1 数位 dp 计和

```

1  #include <bits/stdc++.h>
2  #define ll long long
3  using namespace std;
4  const int mod = 998244353;
5  pair<ll, ll> dp[20][1<<10];
6  bool vis[20][1<<10];
7  int k;
8  int t[20];
9  ll base[20];
10
11 pair<ll, ll> dfs(int pos, int state, bool limit, bool lead) {
12     if (pos == -1) return __builtin_popcount(state) <= k ? make_pair(1, 0) : make_pair
        (0, 0);
13     if (!limit && !lead && vis[pos][state]) return dp[pos][state];
14     int up = limit ? t[pos] : 9;
15     pair<ll, ll> res = {0, 0};
16     for (int i = 0; i <= up; ++i) {
17         int n_s = state;
18         if (lead && i == 0) n_s = 0;
19         else n_s = state | (1 << i);
20         auto tmp = dfs(pos - 1, n_s, limit && i == t[pos], lead && i == 0);
21         ll pre = 1ll * i * base[pos] % mod;
22         (res.first += tmp.first) %= mod;
23         (res.second += tmp.second + pre * tmp.first) %= mod;
24     }
25     if (!limit && !lead) dp[pos][state] = res, vis[pos][state] = 1;
26     return res;
27 }
28
29 ll solve(ll x) {
30     int pos = 0;
31     do {
32         t[pos++] = x % 10;
33     } while (x /= 10);
34     return dfs(pos - 1, 0, true, true).second;
35 }
36
37 int main(int argc, char *argv[])
38 {
39     base[0] = 1;
40     for (int i = 1; i < 20; ++i) base[i] = base[i - 1] * 10;
41     ll l, r;
42     scanf("%lld%lld%d", &l, &r, &k);
43     printf("%lld\n", (solve(r) - solve(l - 1) + mod) % mod);
44     return 0;
45 }

```

5.1.2 两个数数位 dp

```

1  // 二进制数位dp, 求a 在 1~x 和 b 在 1~y, 满足 a & b > c || a ^ b < c 的对数
2  ll dp[maxn][2][2][2][2];
3  int a[maxn], b[maxn], c[maxn];
4

```

```

5
6 void cal(int *xt, ll x) {
7     int has = 0;
8     while (x) {
9         xt[has++] = x % 2;
10        x /= 2;
11    }
12 }
13
14 ll dfs(int pos, int o1, int o2, int lim1, int lim2) {
15     if (pos < 0) return 1;
16     ll &t = dp[pos][o1][o2][lim1][lim2];
17     if (t != -1) return t;
18     int up1 = o1 ? a[pos] : 1;
19     int up2 = o2 ? b[pos] : 1;
20     ll res = 0;
21     for (int i = 0; i <= up1; ++i) {
22         for (int j = 0; j <= up2; ++j) {
23             int t1 = i & j;
24             int t2 = i ^ j;
25             if (lim1 && t1 > c[pos]) continue;
26             if (lim2 && t2 < c[pos]) continue;
27             res += dfs(pos - 1, o1 && i == up1, o2 && j == up2, lim1 && t1 == c[pos],
28                        lim2 && t2 == c[pos]);
29         }
30     }
31     return t = res;
32 }
33 ll solve(ll x, ll y, ll z) {
34     memset(dp, -1ll, sizeof dp);
35     for (int i = 0; i < 33; ++i) a[i] = b[i] = c[i] = 0;
36     cal(a, x);
37     cal(b, y);
38     cal(c, z);
39     return dfs(32, 1, 1, 1, 1);
40 }
41
42 int main(int argc, char *argv[]) {
43     int T;
44     scanf("%d", &T);
45     ll x, y, z;
46     for (int kase = 1; kase <= T; ++kase) {
47         scanf("%lld%lld%lld", &x, &y, &z);
48         ll res = solve(x, y, z);
49         res -= max(0ll, y - z + 1);
50         res -= max(0ll, x - z + 1);
51         printf("%lld\n", x * y - res);
52     }
53     return 0;
54 }

```

5.2 Subsequence

5.2.1 MaxSum

```

1 // 传入序列a和长度n, 返回最大子序列和
2 int MaxSeqSum(int a[], int n)
3 {

```

```
4     int rt = 0, cur = 0;
5     for (int i = 0; i < n; i++)
6         cur += a[i], rt = max(cur, rt), cur = max(0, cur);
7     return rt;
8 }
```

5.2.2 LIS

```
1 // 简单写法(下标从0开始,只返回长度)
2 int dp[N];
3 int LIS(int a[], int n)
4 {
5     memset(dp, 0x3f, sizeof(dp));
6     for (int i = 0; i < n; i++) *lower_bound(dp, dp + n, a[i]) = a[i];
7     return lower_bound(dp, dp + n, INF) - dp;
8 }
9
10 // 小常数nlogn求序列用树状数组维护dp即可
11 // dp[i] = max(dp[j]) + 1 (j < i && a[j] < a[i])
```

5.2.3 LongestCommonIncrease

```
1 // 序列下标从1开始
2 int LCIS(int a[], int b[], int n, int m)
3 {
4     memset(dp, 0, sizeof(dp));
5     for (int i = 1; i <= n; i++)
6     {
7         int ma = 0;
8         for (int j = 1; j <= m; j++)
9         {
10             dp[i][j] = dp[i - 1][j];
11             if (a[i] > b[j]) ma = max(ma, dp[i - 1][j]);
12             if (a[i] == b[j]) dp[i][j] = ma + 1;
13         }
14     }
15     return *max_element(dp[n] + 1, dp[n] + 1 + m);
16 }
```

5.2.4 LCS

```
1 #include <stdio.h>
2 #include <string.h>
3
4 #define M 30005
5 #define SIZE 128
6 #define WORDMAX 3200
7 #define BIT 32
8
9 char s1[M], s2[M];
10 int nword;
11 unsigned int str[SIZE][WORDMAX];
12 unsigned int tmp1[WORDMAX], tmp2[WORDMAX];
13
14 void pre(int len)
15 {
16     int i, j;
```

```
17     memset(str, 0, sizeof(str));
18     for(i = 0; i < len; i ++){
19         str[s1[i]][i / BIT] |= 1 << (i % BIT);
20     }
21
22 void cal(unsigned int *a, unsigned int *b, char ch)
23 {
24     int i, bottom = 1, top;
25     unsigned int x, y;
26     for(i = 0; i < nword; i ++){
27         {
28             y = a[i];
29             x = y | str[ch][i];
30             top = (y >> (BIT - 1)) & 1;
31             y = (y << 1) | bottom;
32             if(x < y) top = 1;
33             b[i] = x & ((x - y) ^ x);
34             bottom = top;
35         }
36     }
37
38 int bitcnt(unsigned int *a)
39 {
40     int i, j, res = 0, t;
41     unsigned int b[5] = {0x55555555, 0x33333333, 0x0f0f0f0f, 0x00ff00ff, 0x0000ffff}, x
42     ;
43     for(i = 0; i < nword; i ++){
44         {
45             x = a[i];
46             t = 1;
47             for(j = 0; j < 5; j ++, t <= 1)
48                 x = (x & b[j]) + ((x >> t) & b[j]);
49             res += x;
50         }
51     }
52     return res;
53
54 void process()
55 {
56     int i, j, len1, len2;
57     unsigned int *a, *b, *t;
58     len1 = strlen(s1);
59     len2 = strlen(s2);
60     nword = (len1 + BIT - 1) / BIT;
61     pre(len1);
62     memset(tmp1, 0, sizeof(tmp1));
63     a = &tmp1[0];
64     b = &tmp2[0];
65     for(i = 0; i < len2; i ++){
66         {
67             cal(a, b, s2[i]);
68             t = a; a = b; b = t;
69         }
70     }
71     printf("%d\n", bitcnt(a));
72
73 int main()
74 {
75     while(scanf("%s%s", s1, s2) != EOF)
```

```

75     process();
76     return 0;
77 }

```

5.3 Others

问题 设 $f(i) = \min(y[k] - s[i] \times x[k]), k \in [1, i-1]$, 现在要求出所有 $f(i), i \in [1, n]$ 考虑两个决策 j 和 k , 如果 j 比 k 优, 则

$$y[j] - s[i] \times x[j] < y[k] - s[i] \times x[k]$$

化简得:

$$\frac{y_j - y_k}{x_j - x_k} < s_i$$

不等式左边是个斜率, 我们把它设为 $\text{slope}(j, k)$

我们可以维护一个单调递增的队列, 为什么呢?

因为如果 $\text{slope}(q[i-1], q[i]) > \text{slope}(q[i], q[i+1])$, 那么当前者成立时, 后者必定成立。即 $q[i]$ 决策优于 $q[i-1]$ 决策时, $q[i+1]$ 必然优于 $q[i]$, 因此 $q[i]$ 就没有存在的必要了。所以我们要维护递增的队列。

那么每次的决策点 i , 都要满足

$$\begin{cases} \text{slope}(q[i-1], q[i]) < s[i] \\ \text{slope}(q[i], q[i+1]) \geq s[i] \end{cases}$$

一般情况去二分这个 i 即可。

如果 $s[i]$ 是单调不降的, 那么对于决策 j 和 $k (j < k)$ 来说, 如果决策 k 优于决策 j , 那么对于 $i \in [k+1, n]$, 都存在决策 k 优于决策 j , 因此决策 j 就可以舍弃了。这样的话我们可以用单调队列进行优化, 可以少个 \log 。

单调队列滑动窗口最大值

```

1 // k为滑动窗口的大小, 数列下标从1开始, d为序列长度+1
2 deque<int> q;
3 for (int i = 0, j = 0; i + k <= d; i++)
4 {
5     while (j < i + k)
6     {
7         while (!q.empty() && a[q.back()] < a[j]) q.pop_back();
8         q.push_back(j++);
9     }
10    while (q.front() < i) q.pop_front();
11    // a[q.front()]为当前滑动窗口的最大值
12 }

```

5.3.1 矩阵快速幂

```

1 const int sz = 5;
2 struct Matrix {
3     ll a[sz][sz];
4     Matrix() { memset(a, 0, sizeof a); }
5     void pr() {
6         printf("*\n");
7         for (int i = 0; i < sz; ++i) {
8             for (int j = 0; j < sz; ++j) {
9                 printf("%lld ", a[i][j]);
10            }
11            printf("\n");
12        }
13    }
14    void tr() {
15        for (int i = 0; i < sz; ++i) {

```



```
16         for (int j = i + 1; j < sz; ++j) {
17             swap(a[i][j], a[j][i]);
18         }
19     }
20 }
21 } res, t1;
22
23 void init() {
24     ;
25 }
26
27 Matrix mul(Matrix a, Matrix b) {
28     Matrix res;
29     // assert(a.m == b.n)
30     for (int i = 0; i < sz; i++) // a.n
31         for (int j = 0; j < sz; j++) // b.m
32             for (int k = 0; k < sz; k++) // a.m, b.n
33                 (res.a[i][j] += a.a[i][k] * b.a[k][j] % mod) %= mod;
34     return res;
35 }
36
37 Matrix Pow(ll n) {
38     init();
39     //for(int i = 0; i < cur; i++) res.a[i][i] = 1;
40     while (n > 0) {
41         if (n & 1) res = mul(res, t1);
42         t1 = mul(t1, t1);
43         n >>= 1;
44     }
45     return res;
46 }
```

5.3.2 单调栈

```
1 // 求左边第一个比a[i]小的和右边最后一个不比a[i]小的位置
2 for (int i = 1; i <= n; i++) {
3     while (top && a[sta[top - 1]] >= a[i]) top--;
4     la[i] = (top == 0) ? 1 : sta[top - 1] + 1;
5     sta[top++] = i;
6 }
7 top = 0;
8 for (int i = n; i >= 1; i--) {
9     while (top && a[sta[top - 1]] >= a[i]) top--;
10    ra[i] = (top == 0) ? n : sta[top - 1] - 1;
11    sta[top++] = i;
12 }
```

5.3.3 单调队列

```
1 // 循环序列的最大子段和
2 int a[maxn];
3 int pre[maxn * 2];
4 int qu[maxn * 2];
5 int n, resl, resr, res, k;
6
7 int main(int argc, char* argv[]) {
8     int T;
9     scanf("%d", &T);
```

```
10     for (int kase = 1; kase <= T; ++kase) {
11         scanf("%d%d", &n, &k);
12         for (int i = 1; i <= n; ++i) {
13             scanf("%d", &a[i]);
14             pre[i] = pre[i - 1] + a[i];
15         }
16         for (int i = n + 1; i <= 2 * n; ++i) {
17             pre[i] = pre[i - 1] + a[i - n];
18         }
19         res = -0x3f3f3f3f;
20         resl = resr = -1;
21         int l = 1, r = 0;
22         for (int i = 1; i <= 2 * n; ++i) {
23             while (l <= r && pre[qu[r]] >= pre[i - 1]) r--;
24             qu[++r] = i - 1;
25             while (l <= r && qu[l] < i - k) l++;
26             int tmp = pre[i] - pre[qu[l]];
27             if (tmp > res) {
28                 res = tmp;
29                 resl = qu[l] + 1;
30                 resr = i;
31             }
32         }
33         if (resl > n) resl -= n;
34         if (resr > n) resr -= n;
35         printf("%d %d %d\n", res, resl, resr);
36     }
37     return 0;
38 }
```

6 Geometry

6.1 geo

```

1  #define mp make_pair
2  #define fi first
3  #define se second
4  #define pb push_back
5  typedef double db;
6  const db eps=1e-6;
7  const db pi=acos(-1);
8  int sign(db k){
9      if (k>eps) return 1; else if (k<-eps) return -1; return 0;
10 }
11 int cmp(db k1,db k2){return sign(k1-k2);}
12 int inmid(db k1,db k2,db k3){return sign(k1-k3)*sign(k2-k3)<=0;}// k3 在 [k1,k2] 内
13 struct point{
14     db x,y;
15     point operator + (const point &k1) const{return (point){k1.x+x,k1.y+y};}
16     point operator - (const point &k1) const{return (point){x-k1.x,y-k1.y};}
17     point operator * (db k1) const{return (point){x*k1,y*k1};}
18     point operator / (db k1) const{return (point){x/k1,y/k1};}
19     int operator == (const point &k1) const{return cmp(x,k1.x)==0&&cmp(y,k1.y)==0;}
20     // 逆时针旋转
21     point turn(db k1){return (point){x*cos(k1)-y*sin(k1),x*sin(k1)+y*cos(k1)};}
22     point turn90(){return (point){-y,x};}
23     bool operator < (const point k1) const{
24         int a=cmp(x,k1.x);
25         if (a==1) return 1; else if (a==0) return 0; else return cmp(y,k1.y)<0;
26     }
27     db abs(){return sqrt(x*x+y*y);}
28     db abs2(){return x*x+y*y;}
29     db dis(point k1){return ((*this)-k1).abs();}
30     point unit(){db w=abs(); return (point){x/w,y/w};}
31     void scan(){double k1,k2; scanf("%lf%lf",&k1,&k2); x=k1; y=k2;}
32     void print(){printf("%.11lf %.11lf\n",x,y);}
33     db getw(){return atan2(y,x);}
34     point getdel(){if (sign(x)==-1||(sign(x)==0&&sign(y)==-1)) return (*this)*(-1);
35         else return (*this);}
36     int getP() const{return sign(y)==1||(sign(y)==0&&sign(x)==-1);}
37 };
38 int inmid(point k1,point k2,point k3){return inmid(k1.x,k2.x,k3.x)&&inmid(k1.y,k2.y,k3.y);}
39 db cross(point k1,point k2){return k1.x*k2.y-k1.y*k2.x;}
40 db dot(point k1,point k2){return k1.x*k2.x+k1.y*k2.y;}
41 db rad(point k1,point k2){return atan2(cross(k1,k2),dot(k1,k2));}
42 // -pi -> pi
43 int compareangle (point k1,point k2){
44     return k1.getP()<k2.getP()||(k1.getP()==k2.getP()&&sign(cross(k1,k2))>0);
45 }
46 point proj(point k1,point k2,point q){ // q 到直线 k1,k2 的投影
47     point k=k2-k1; return k1+k*(dot(q-k1,k)/k.abs2());
48 }
49 point reflect(point k1,point k2,point q){return proj(k1,k2,q)*2-q;}
50 int clockwise(point k1,point k2,point k3){// k1 k2 k3 逆时针 1 顺时针 -1 否则 0
51     return sign(cross(k2-k1,k3-k1));
52 }
53 int checkLL(point k1,point k2,point k3,point k4){// 求直线 (L) 线段 (S)k1,k2 和 k3,k4
54     // 的交点

```

```

53     return cmp(cross(k3-k1,k4-k1),cross(k3-k2,k4-k2))!=0;
54 }
55 point getLL(point k1,point k2,point k3,point k4){
56     db w1=cross(k1-k3,k4-k3),w2=cross(k4-k3,k2-k3); return (k1*w2+k2*w1)/(w1+w2);
57 }
58 int intersect(db l1,db r1,db l2,db r2){
59     if (l1>r1) swap(l1,r1); if (l2>r2) swap(l2,r2); return cmp(r1,l2)!=-1&&cmp(r2,l1)
        !=-1;
60 }
61 int checkSS(point k1,point k2,point k3,point k4){
62     return intersect(k1.x,k2.x,k3.x,k4.x)&&intersect(k1.y,k2.y,k3.y,k4.y)&&
63     sign(cross(k3-k1,k4-k1))*sign(cross(k3-k2,k4-k2))<=0&&
64     sign(cross(k1-k3,k2-k3))*sign(cross(k1-k4,k2-k4))<=0;
65 }
66 db disSP(point k1,point k2,point q){
67     point k3=proj(k1,k2,q);
68     if (inmid(k1,k2,k3)) return q.dis(k3); else return min(q.dis(k1),q.dis(k2));
69 }
70 db disSS(point k1,point k2,point k3,point k4){
71     if (checkSS(k1,k2,k3,k4)) return 0;
72     else return min(min(disSP(k1,k2,k3),disSP(k1,k2,k4)),min(disSP(k3,k4,k1),disSP(k3,
        k4,k2)));
73 }
74 int onS(point k1,point k2,point q){return inmid(k1,k2,q)&&sign(cross(k1-q,k2-k1))==0;}
75 struct circle{
76     point o; db r;
77     void scan(){o.scan(); scanf("%lf",&r);}
78     int inside(point k){return cmp(r,o.dis(k));}
79 };
80 struct line{
81     // p[0]->p[1]
82     point p[2];
83     line(point k1,point k2){p[0]=k1; p[1]=k2;}
84     point& operator [] (int k){return p[k];}
85     int include(point k){return sign(cross(p[1]-p[0],k-p[0]))>0;}
86     point dir(){return p[1]-p[0];}
87     line push(){ // 向外 ( 左手边 ) 平移 eps
88         const db eps = 1e-6;
89         point delta=(p[1]-p[0]).turn90().unit()*eps;
90         return {p[0]-delta,p[1]-delta};
91     }
92 };
93 point getLL(line k1,line k2){return getLL(k1[0],k1[1],k2[0],k2[1]);}
94 int parallel(line k1,line k2){return sign(cross(k1.dir(),k2.dir()))==0;}
95 int sameDir(line k1,line k2){return parallel(k1,k2)&&sign(dot(k1.dir(),k2.dir()))==1;}
96 int operator < (line k1,line k2){
97     if (sameDir(k1,k2)) return k2.include(k1[0]);
98     return compareangle(k1.dir(),k2.dir());
99 }
100 int checkpos(line k1,line k2,line k3){return k3.include(getLL(k1,k2));}
101 vector<point> halfPlaneIS(vector<line> &l) {// 求半平面交 , 半平面是逆时针方向 , 输出按
    照逆时针, 点数小于等于2表示无法构成半平面交
102     sort(l.begin(), l.end());
103     deque<line> q;
104     for (int i = 0; i < (int)l.size(); ++i) {
105         if (i && sameDir(l[i], l[i - 1])) continue;
106         while (q.size() > 1 && !checkpos(q[q.size() - 2], q[q.size() - 1], l[i])) q.
            pop_back();
107         while (q.size() > 1 && !checkpos(q[1], q[0], l[i])) q.pop_front();

```

```

108     q.push_back(l[i]);
109 }
110 while (q.size() > 2 && !checkpos(q[q.size() - 2], q[q.size() - 1], q[0])) q.
    pop_back();
111 while (q.size() > 2 && !checkpos(q[1], q[0], q[q.size() - 1])) q.pop_front();
112 vector<point> ret;
113 for (int i = 0; i < (int)q.size(); ++i) ret.push_back(getLL(q[i], q[(i + 1) % q.
    size()]));
114 return ret;
115 }
116 // vector<line> getHL(vector<line> &L){ // 求半平面交 , 半平面是逆时针方向 , 输出按照逆
    时针
117 //     sort(L.begin(),L.end()); deque<line> q;
118 //     for (int i=0;i<(int)L.size();i++){
119 //         if (i&&sameDir(L[i],L[i-1])) continue;
120 //         while (q.size()>1&&!checkpos(q[q.size()-2],q[q.size()-1],L[i])) q.pop_back()
    ;
121 //         while (q.size()>1&&!checkpos(q[1],q[0],L[i])) q.pop_front();
122 //         q.push_back(L[i]);
123 //     }
124 //     while (q.size()>2&&!checkpos(q[q.size()-2],q[q.size()-1],q[0])) q.pop_back();
125 //     while (q.size()>2&&!checkpos(q[1],q[0],q[q.size()-1])) q.pop_front();
126 //     vector<line>ans; for (int i=0;i<q.size();i++) ans.push_back(q[i]);
127 //     return ans;
128 // }
129 db closepoint(vector<point>&A,int l,int r){ // 最近点对 , 先要按照 x 坐标排序
130     if (r-l<=5){
131         db ans=1e20;
132         for (int i=l;i<=r;i++) for (int j=i+1;j<=r;j++) ans=min(ans,A[i].dis(A[j]));
133         return ans;
134     }
135     int mid=l+r>>1; db ans=min(closepoint(A,l,mid),closepoint(A,mid+1,r));
136     vector<point>B; for (int i=l;i<=r;i++) if (abs(A[i].x-A[mid].x)<=ans) B.push_back(A
        [i]);
137     sort(B.begin(),B.end(),[](point k1,point k2){return k1.y<k2.y;});
138     for (int i=0;i<B.size();i++) for (int j=i+1;j<B.size()&&B[j].y-B[i].y<ans;j++) ans=
        min(ans,B[i].dis(B[j]));
139     return ans;
140 }
141 int checkposCC(circle k1,circle k2){// 返回两个圆的公切线数量
142     if (cmp(k1.r,k2.r)==-1) swap(k1,k2);
143     db dis=k1.o.dis(k2.o); int w1=cmp(dis,k1.r+k2.r),w2=cmp(dis,k1.r-k2.r);
144     if (w1>0) return 4; else if (w1==0) return 3; else if (w2>0) return 2;
145     else if (w2==0) return 1; else return 0;
146 }
147 vector<point> getCL(circle k1,point k2,point k3){ // 沿着 k2->k3 方向给出 , 相切给出两
    个
148     point k=proj(k2,k3,k1.o); db d=k1.r*k1.r-(k-k1.o).abs2();
149     if (sign(d)==-1) return {};
150     point del=(k3-k2).unit()*sqrt(max((db)0.0,d)); return {k-del,k+del};
151 }
152 vector<point> getCC(circle k1,circle k2){// 沿圆 k1 逆时针给出 , 相切给出两个
153     int pd=checkposCC(k1,k2); if (pd==0||pd==4) return {};
154     db a=(k2.o-k1.o).abs2(),cosA=(k1.r*k1.r+a-k2.r*k2.r)/(2*k1.r*sqrt(max(a,(db)0.0)));
155     db b=k1.r*cosA,c=sqrt(max((db)0.0,k1.r*k1.r-b*b));
156     point k=(k2.o-k1.o).unit(),m=k1.o+k*b,del=k.turn90()*c;
157     return {m-del,m+del};
158 }
159 vector<point> TangentCP(circle k1,point k2){// 沿圆 k1 逆时针给出

```

```

160     db a=(k2-k1.o).abs(),b=k1.r*k1.r/a,c=sqrt(max((db)0.0,k1.r*k1.r-b*b));
161     point k=(k2-k1.o).unit(),m=k1.o+k*b,del=k.turn90()*c;
162     return {m-del,m+del};
163 }
164 vector<line> TangentoutCC(circle k1,circle k2){
165     int pd=checkposCC(k1,k2); if (pd==0) return {};
166     if (pd==1){point k=getCC(k1,k2)[0]; return {(line){k,k}};}
167     if (cmp(k1.r,k2.r)==0){
168         point del=(k2.o-k1.o).unit().turn90().getdel();
169         return {(line){k1.o-del*k1.r,k2.o-del*k2.r},{k1.o+del*k1.r,k2.o+del*k2.r}};
170     } else {
171         point p=(k2.o*k1.r-k1.o*k2.r)/(k1.r-k2.r);
172         vector<point>A=TangentCP(k1,p),B=TangentCP(k2,p);
173         vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});
174         return ans;
175     }
176 }
177 vector<line> TangentinCC(circle k1,circle k2){
178     int pd=checkposCC(k1,k2); if (pd<=2) return {};
179     if (pd==3){point k=getCC(k1,k2)[0]; return {(line){k,k}};}
180     point p=(k2.o*k1.r+k1.o*k2.r)/(k1.r+k2.r);
181     vector<point>A=TangentCP(k1,p),B=TangentCP(k2,p);
182     vector<line>ans; for (int i=0;i<A.size();i++) ans.push_back((line){A[i],B[i]});
183     return ans;
184 }
185 vector<line> TangentCC(circle k1,circle k2){
186     int flag=0; if (k1.r<k2.r) swap(k1,k2),flag=1;
187     vector<line>A=TangentoutCC(k1,k2),B=TangentinCC(k1,k2);
188     for (line k:B) A.push_back(k);
189     if (flag) for (line &k:A) swap(k[0],k[1]);
190     return A;
191 }
192 db getarea(circle k1,point k2,point k3){
193     // 圆 k1 与三角形 k2 k3 k1.o 的有向面积交
194     point k=k1.o; k1.o=k1.o-k; k2=k2-k; k3=k3-k;
195     int pd1=k1.inside(k2),pd2=k1.inside(k3);
196     vector<point>A=getCL(k1,k2,k3);
197     if (pd1>=0){
198         if (pd2>=0) return cross(k2,k3)/2;
199         return k1.r*k1.r*rad(A[1],k3)/2+cross(k2,A[1])/2;
200     } else if (pd2>=0){
201         return k1.r*k1.r*rad(k2,A[0])/2+cross(A[0],k3)/2;
202     } else {
203         int pd=cmp(k1.r,disSP(k2,k3,k1.o));
204         if (pd<=0) return k1.r*k1.r*rad(k2,k3)/2;
205         return cross(A[0],A[1])/2+k1.r*k1.r*(rad(k2,A[0])+rad(A[1],k3))/2;
206     }
207 }
208 circle getcircle(point k1,point k2,point k3){
209     db a1=k2.x-k1.x,b1=k2.y-k1.y,c1=(a1*a1+b1*b1)/2;
210     db a2=k3.x-k1.x,b2=k3.y-k1.y,c2=(a2*a2+b2*b2)/2;
211     db d=a1*b2-a2*b1;
212     point o=(point){k1.x+(c1*b2-c2*b1)/d,k1.y+(a1*c2-a2*c1)/d};
213     return (circle){o,k1.dis(o)};
214 }
215 circle getScircle(vector<point> A){
216     random_shuffle(A.begin(),A.end());
217     circle ans=(circle){A[0],0};

```

```

218     for (int i=1;i<A.size();i++)
219         if (ans.inside(A[i])== -1){
220             ans=(circle){A[i],0};
221             for (int j=0;j<i;j++){
222                 if (ans.inside(A[j])== -1){
223                     ans.o=(A[i]+A[j])/2; ans.r=ans.o.dis(A[i]);
224                     for (int k=0;k<j;k++){
225                         if (ans.inside(A[k])== -1)
226                             ans=getcircle(A[i],A[j],A[k]);
227                     }
228                 }
229             }
230         }
231     db area(vector<point> A){ // 多边形用 vector<point> 表示 , 逆时针
232         db ans=0;
233         for (int i=0;i<A.size();i++) ans+=cross(A[i],A[(i+1)%A.size()]);
234         return ans/2;
235     }
236     int checkconvex(vector<point>A){
237         int n=A.size(); A.push_back(A[0]); A.push_back(A[1]);
238         for (int i=0;i<n;i++) if (sign(cross(A[i+1]-A[i],A[i+2]-A[i]))== -1) return 0;
239         return 1;
240     }
241     int contain(vector<point>A,point q){ // 2 内部 1 边界 0 外部
242         int pd=0; A.push_back(A[0]);
243         for (int i=1;i<A.size();i++){
244             point u=A[i-1],v=A[i];
245             if (onS(u,v,q)) return 1; if (cmp(u.y,v.y)>0) swap(u,v);
246             if (cmp(u.y,q.y)>0||cmp(v.y,q.y)<0) continue;
247             if (sign(cross(u-v,q-v))<0) pd^=1;
248         }
249         return pd<<1;
250     }
251     vector<point> ConvexHull(vector<point>A,int flag=1){ // flag=0 不严格 flag=1 严格
252         int n=A.size(); vector<point>ans(n*2);
253         sort(A.begin(),A.end()); int now=-1;
254         for (int i=0;i<A.size();i++){
255             while (now>0&&sign(cross(ans[now]-ans[now-1],A[i]-ans[now-1]))<flag) now--;
256             ans[++now]=A[i];
257         } int pre=now;
258         for (int i=n-2;i>=0;i--){
259             while (now>pre&&sign(cross(ans[now]-ans[now-1],A[i]-ans[now-1]))<flag) now--;
260             ans[++now]=A[i];
261         } ans.resize(now); return ans;
262     }
263     db convexDiameter(vector<point>A){
264         int now=0,n=A.size(); db ans=0;
265         for (int i=0;i<A.size();i++){
266             now=max(now,i);
267             while (1){
268                 db k1=A[i].dis(A[now%n]),k2=A[i].dis(A[(now+1)%n]);
269                 ans=max(ans,max(k1,k2)); if (k2>k1) now++; else break;
270             }
271         }
272         return ans;
273     }
274     vector<point> convexcut(vector<point>A,point k1,point k2){
275         // 保留 k1,k2,p 逆时针的所有点
276         int n=A.size(); A.push_back(A[0]); vector<point>ans;

```

```

277     for (int i=0;i<n;i++){
278         int w1=clockwise(k1,k2,A[i]),w2=clockwise(k1,k2,A[i+1]);
279         if (w1>=0) ans.push_back(A[i]);
280         if (w1*w2<0) ans.push_back(getLL(k1,k2,A[i],A[i+1]));
281     }
282     return ans;
283 }
284 int checkPoS(vector<point>A,point k1,point k2){
285     // 多边形 A 和直线 ( 线段 )k1->k2 严格相交 , 注释部分为线段
286     struct ins{
287         point m,u,v;
288         int operator < (const ins& k) const {return m<k.m;}
289     }; vector<ins>B;
290     //if (contain(A,k1)==2||contain(A,k2)==2) return 1;
291     vector<point>poly=A; A.push_back(A[0]);
292     for (int i=1;i<A.size();i++) if (checkLL(A[i-1],A[i],k1,k2)){
293         point m=getLL(A[i-1],A[i],k1,k2);
294         if (inmid(A[i-1],A[i],m)/*&&inmid(k1,k2,m)*/) B.push_back((ins){m,A[i-1],A[i]});
295     }
296     if (B.size()==0) return 0; sort(B.begin(),B.end());
297     int now=1; while (now<B.size()&&B[now].m==B[0].m) now++;
298     if (now==B.size()) return 0;
299     int flag=contain(poly,(B[0].m+B[now].m)/2);
300     if (flag==2) return 1;
301     point d=B[now].m-B[0].m;
302     for (int i=now;i<B.size();i++){
303         if (!(B[i].m==B[i-1].m)&&flag==2) return 1;
304         int tag=sign(cross(B[i].v-B[i].u,B[i].m+d-B[i].u));
305         if (B[i].m==B[i].u||B[i].m==B[i].v) flag+=tag; else flag+=tag*2;
306     }
307     //return 0;
308     return flag==2;
309 }
310 int checkinp(point r,point l,point m){
311     if (compareangle(l,r)){return compareangle(l,m)&&compareangle(m,r);}
312     return compareangle(l,m)||compareangle(m,r);
313 }
314 int checkPosFast(vector<point>A,point k1,point k2){ // 快速检查线段是否和多边形严格相交
315     if (contain(A,k1)==2||contain(A,k2)==2) return 1; if (k1==k2) return 0;
316     A.push_back(A[0]); A.push_back(A[1]);
317     for (int i=1;i+1<A.size();i++){
318         if (checkLL(A[i-1],A[i],k1,k2)){
319             point now=getLL(A[i-1],A[i],k1,k2);
320             if (inmid(A[i-1],A[i],now)==0||inmid(k1,k2,now)==0) continue;
321             if (now==A[i]){
322                 if (A[i]==k2) continue;
323                 point pre=A[i-1],ne=A[i+1];
324                 if (checkinp(pre-now,ne-now,k2-now)) return 1;
325             } else if (now==k1){
326                 if (k1==A[i-1]||k1==A[i]) continue;
327                 if (checkinp(A[i-1]-k1,A[i]-k1,k2-k1)) return 1;
328             } else if (now==k2||now==A[i-1]) continue;
329             else return 1;
330         }
331     }
332     return 0;
333 }
334 // 拆分凸包成上下凸壳 凸包尽量都随机旋转一个角度来避免出现相同横坐标
335 // 尽量特判只有一个点的情况 凸包逆时针

```



```

335 void getUDP(vector<point>A,vector<point>&U,vector<point>&D){
336     db l=1e100,r=-1e100;
337     for (int i=0;i<A.size();i++) l=min(l,A[i].x),r=max(r,A[i].x);
338     int wherel,wherer;
339     for (int i=0;i<A.size();i++) if (cmp(A[i].x,l)==0) wherel=i;
340     for (int i=A.size();i;i--) if (cmp(A[i-1].x,r)==0) wherer=i-1;
341     U.clear(); D.clear(); int now=wherel;
342     while (1){D.push_back(A[now]); if (now==wherer) break; now++; if (now>=A.size())
        now=0;}
343     now=wherel;
344     while (1){U.push_back(A[now]); if (now==wherer) break; now--; if (now<0) now=A.size()
        (-1);}
345 }
346 // 需要保证凸包点数大于等于 3,2 内部 ,1 边界 ,0 外部
347 int containCoP(const vector<point>&U,const vector<point>&D,point k){
348     db lx=U[0].x,rx=U[U.size()-1].x;
349     if (k==U[0]||k==U[U.size()-1]) return 1;
350     if (cmp(k.x,lx)==-1||cmp(k.x,rx)==1) return 0;
351     int where1=lower_bound(U.begin(),U.end(),(point){k.x,-1e100})-U.begin();
352     int where2=lower_bound(D.begin(),D.end(),(point){k.x,-1e100})-D.begin();
353     int w1=clockwise(U[where1-1],U[where1],k),w2=clockwise(D[where2-1],D[where2],k);
354     if (w1==1||w2==-1) return 0; else if (w1==0||w2==0) return 1; return 2;
355 }
356 // d 是方向 , 输出上方切点和下方切点
357 pair<point,point> getTangentCow(const vector<point> &U,const vector<point> &D,point d){
358     if (sign(d.x)<0||(sign(d.x)==0&&sign(d.y)<0)) d=d*(-1);
359     point whereU,whereD;
360     if (sign(d.x)==0) return mp(U[0],U[U.size()-1]);
361     int l=0,r=U.size()-1,ans=0;
362     while (l<r){int mid=l+r>>1; if (sign(cross(U[mid+1]-U[mid],d))<=0) l=mid+1,ans=mid
        +1; else r=mid;}
363     whereU=U[ans]; l=0,r=D.size()-1,ans=0;
364     while (l<r){int mid=l+r>>1; if (sign(cross(D[mid+1]-D[mid],d))>=0) l=mid+1,ans=mid
        +1; else r=mid;}
365     whereD=D[ans]; return mp(whereU,whereD);
366 }
367 // 先检查 contain, 逆时针给出
368 pair<point,point> getTangentCoP(const vector<point>&U,const vector<point>&D,point k){
369     db lx=U[0].x,rx=U[U.size()-1].x;
370     if (k.x<lx){
371         int l=0,r=U.size()-1,ans=U.size()-1;
372         while (l<r){int mid=l+r>>1; if (clockwise(k,U[mid],U[mid+1])==1) l=mid+1; else
            ans=mid,r=mid;}
373         point w1=U[ans]; l=0,r=D.size()-1,ans=D.size()-1;
374         while (l<r){int mid=l+r>>1; if (clockwise(k,D[mid],D[mid+1])==1) l=mid+1; else
            ans=mid,r=mid;}
375         point w2=D[ans]; return mp(w1,w2);
376     } else if (k.x>rx){
377         int l=1,r=U.size(),ans=0;
378         while (l<r){int mid=l+r>>1; if (clockwise(k,U[mid],U[mid-1])==1) r=mid; else
            ans=mid,l=mid+1;}
379         point w1=U[ans]; l=1,r=D.size(),ans=0;
380         while (l<r){int mid=l+r>>1; if (clockwise(k,D[mid],D[mid-1])==1) r=mid; else
            ans=mid,l=mid+1;}
381         point w2=D[ans]; return mp(w2,w1);
382     } else {
383         int where1=lower_bound(U.begin(),U.end(),(point){k.x,-1e100})-U.begin();
384         int where2=lower_bound(D.begin(),D.end(),(point){k.x,-1e100})-D.begin();
385         if ((k.x==lx&&k.y>U[0].y)||((where1&&clockwise(U[where1-1],U[where1],k)==1)){

```

```

386         int l=1,r=where1+1,ans=0;
387         while (l<r){int mid=l+r>>1; if (clockwise(k,U[mid],U[mid-1])==1) ans=mid,l=
            mid+1; else r=mid;}
388         point w1=U[ans]; l=where1,r=U.size()-1,ans=U.size()-1;
389         while (l<r){int mid=l+r>>1; if (clockwise(k,U[mid],U[mid+1])==1) l=mid+1;
            else ans=mid,r=mid;}
390         point w2=U[ans]; return mp(w2,w1);
391     } else {
392         int l=1,r=where2+1,ans=0;
393         while (l<r){int mid=l+r>>1; if (clockwise(k,D[mid],D[mid-1])==1) ans=mid,l=
            mid+1; else r=mid;}
394         point w1=D[ans]; l=where2,r=D.size()-1,ans=D.size()-1;
395         while (l<r){int mid=l+r>>1; if (clockwise(k,D[mid],D[mid+1])==1) l=mid+1;
            else ans=mid,r=mid;}
396         point w2=D[ans]; return mp(w1,w2);
397     }
398 }
399 }
400 struct P3{
401     db x,y,z;
402     P3 operator + (P3 k1){return (P3){x+k1.x,y+k1.y,z+k1.z};}
403     P3 operator - (P3 k1){return (P3){x-k1.x,y-k1.y,z-k1.z};}
404     P3 operator * (db k1){return (P3){x*k1,y*k1,z*k1};}
405     P3 operator / (db k1){return (P3){x/k1,y/k1,z/k1};}
406     db abs2(){return x*x+y*y+z*z;}
407     db abs(){return sqrt(x*x+y*y+z*z);}
408     P3 unit(){return (*this)/abs();}
409     int operator < (const P3 k1) const{
410         if (cmp(x,k1.x)!=0) return x<k1.x;
411         if (cmp(y,k1.y)!=0) return y<k1.y;
412         return cmp(z,k1.z)==-1;
413     }
414     int operator == (const P3 k1){
415         return cmp(x,k1.x)==0&&cmp(y,k1.y)==0&&cmp(z,k1.z)==0;
416     }
417     void scan(){
418         double k1,k2,k3; scanf("%lf%lf%lf",&k1,&k2,&k3);
419         x=k1; y=k2; z=k3;
420     }
421 };
422 P3 cross(P3 k1,P3 k2){return (P3){k1.y*k2.z-k1.z*k2.y,k1.z*k2.x-k1.x*k2.z,k1.x*k2.y-k1.
    y*k2.x};}
423 db dot(P3 k1,P3 k2){return k1.x*k2.x+k1.y*k2.y+k1.z*k2.z;}
424 //p=(3,4,5),l=(13,19,21),theta=85 ans=(2.83,4.62,1.77)
425 P3 turn3D(db k1,P3 l,P3 p){
426     l=l.unit(); P3 ans; db c=cos(k1),s=sin(k1);
427     ans.x=p.x*(l.x*l.x*(1-c)+c)+p.y*(l.x*l.y*(1-c)-l.z*s)+p.z*(l.x*l.z*(1-c)+l.y*s);
428     ans.y=p.x*(l.x*l.y*(1-c)+l.z*s)+p.y*(l.y*l.y*(1-c)+c)+p.z*(l.y*l.z*(1-c)-l.x*s);
429     ans.z=p.x*(l.x*l.z*(1-c)-l.y*s)+p.y*(l.y*l.z*(1-c)+l.x*s)+p.z*(l.x*l.x*(1-c)+c);
430     return ans;
431 }
432 typedef vector<P3> VP;
433 typedef vector<VP> VVP;
434 db Acos(db x){return acos(max(-(db)1,min(x,(db)1)));}
435 // 球面距离 , 圆心原点 , 半径 1
436 db Odist(P3 a,P3 b){db r=Acos(dot(a,b)); return r;}
437 db r; P3 rnd;
438 vector<db> solve(db a,db b,db c){
439     db r=sqrt(a*a+b*b),th=atan2(b,a);

```

```

440     if (cmp(c,-r)==-1) return {0};
441     else if (cmp(r,c)<=0) return {1};
442     else {
443         db tr=pi-Acos(c/r); return {th+pi-tr,th+pi+tr};
444     }
445 }
446 vector<db> jiao(P3 a,P3 b){
447     // dot(rd+x*cos(t)+y*sin(t),b) >= cos(r)
448     if (cmp(0dist(a,b),2*r)>0) return {0};
449     P3 rd=a*cos(r),z=a.unit(),y=cross(z,rd).unit(),x=cross(y,z).unit();
450     vector<db> ret = solve(-(dot(x,b)*sin(r)),-(dot(y,b)*sin(r)),-(cos(r)-dot(rd,b)));
451     return ret;
452 }
453 db norm(db x,db l=0,db r=2*pi){ // change x into [l,r)
454     while (cmp(x,l)==-1) x+=(r-l); while (cmp(x,r)>=0) x-=(r-l);
455     return x;
456 }
457 db disLP(P3 k1,P3 k2,P3 q){
458     return (cross(k2-k1,q-k1)).abs()/(k2-k1).abs();
459 }
460 db disLL(P3 k1,P3 k2,P3 k3,P3 k4){
461     P3 dir=cross(k2-k1,k4-k3); if (sign(dir.abs())==0) return disLP(k1,k2,k3);
462     return fabs(dot(dir.unit(),k1-k2));
463 }
464 VP getFL(P3 p,P3 dir,P3 k1,P3 k2){
465     db a=dot(k2-p,dir),b=dot(k1-p,dir),d=a-b;
466     if (sign(fabs(d))==0) return {};
467     return {(k1*a-k2*b)/d};
468 }
469 VP getFF(P3 p1,P3 dir1,P3 p2,P3 dir2){// 返回一条线
470     P3 e=cross(dir1,dir2),v=cross(dir1,e);
471     db d=dot(dir2,v); if (sign(abs(d))==0) return {};
472     P3 q=p1+v*dot(dir2,p2-p1)/d; return {q,q+e};
473 }
474 // 3D Convex Hull Template
475 db getV(P3 k1,P3 k2,P3 k3,P3 k4){ // get the Volume
476     return dot(cross(k2-k1,k3-k1),k4-k1);
477 }
478 db rand_db(){return 1.0*rand()/RAND_MAX;}
479 VP convexHull2D(VP A,P3 dir){
480     P3 x={(db)rand(),(db)rand(),(db)rand()}; x=x.unit();
481     x=cross(x,dir).unit(); P3 y=cross(x,dir).unit();
482     P3 vec=dir.unit()*dot(A[0],dir);
483     vector<point>B;
484     for (int i=0;i<A.size();i++) B.push_back((point){dot(A[i],x),dot(A[i],y)});
485     B=ConvexHull(B); A.clear();
486     for (int i=0;i<B.size();i++) A.push_back(x*B[i].x+y*B[i].y+vec);
487     return A;
488 }
489 namespace CH3{
490     VVP ret; set<pair<int,int> >e;
491     int n; VP p,q;
492     void wrap(int a,int b){
493         if (e.find({a,b})==e.end()){
494             int c=-1;
495             for (int i=0;i<n;i++) if (i!=a&&i!=b){
496                 if (c==-1||sign(getV(q[c],q[a],q[b],q[i]))>0) c=i;
497             }
498             if (c!=-1){

```

```

499         ret.push_back({p[a],p[b],p[c]});
500         e.insert({a,b}); e.insert({b,c}); e.insert({c,a});
501         wrap(c,b); wrap(a,c);
502     }
503 }
504 }
505 VVP ConvexHull3D(VP _p){
506     p=q=_p; n=p.size();
507     ret.clear(); e.clear();
508     for (auto &i:q) i=i+(P3){rand_db()*1e-4,rand_db()*1e-4,rand_db()*1e-4};
509     for (int i=1;i<n;i++) if (q[i].x<q[0].x) swap(p[0],p[i]),swap(q[0],q[i]);
510     for (int i=2;i<n;i++) if ((q[i].x-q[0].x)*(q[i].y-q[0].y)>(q[i].y-q[0].y)*(q
        [1].x-q[0].x)) swap(q[1],q[i]),swap(p[1],p[i]);
511     wrap(0,1);
512     return ret;
513 }
514 }
515 VVP reduceCH(VVP A){
516     VVP ret; map<P3,VP> M;
517     for (VP nowF:A){
518         P3 dir=cross(nowF[1]-nowF[0],nowF[2]-nowF[0]).unit();
519         for (P3 k1:nowF) M[dir].pb(k1);
520     }
521     for (pair<P3,VP> nowF:M) ret.pb(convexHull2D(nowF.se,nowF.fi));
522     return ret;
523 }
524 // 把一个面变成 ( 点 , 法向量 ) 的形式
525 pair<P3,P3> getF(VP F){
526     return mp(F[0],cross(F[1]-F[0],F[2]-F[0]).unit());
527 }
528 // 3D Cut 保留 dot(dir,x-p)>=0 的部分
529 VVP ConvexCut3D(VVP A,P3 p,P3 dir){
530     VVP ret; VP sec;
531     for (VP nowF: A){
532         int n=nowF.size(); VP ans; int dif=0;
533         for (int i=0;i<n;i++){
534             int d1=sign(dot(dir,nowF[i]-p));
535             int d2=sign(dot(dir,nowF[(i+1)%n]-p));
536             if (d1>=0) ans.pb(nowF[i]);
537             if (d1*d2<0){
538                 P3 q=getFL(p,dir,nowF[i],nowF[(i+1)%n])[0];
539                 ans.push_back(q); sec.push_back(q);
540             }
541             if (d1==0) sec.push_back(nowF[i]); else dif=1;
542             dif|=(sign(dot(dir,cross(nowF[(i+1)%n]-nowF[i],nowF[(i+1)%n]-nowF[i]))
                ==-1);
543         }
544         if (ans.size()>0&&dif) ret.push_back(ans);
545     }
546     if (sec.size()>0) ret.push_back(convexHull2D(sec,dir));
547     return ret;
548 }
549 db vol(VVP A){
550     if (A.size()==0) return 0; P3 p=A[0][0]; db ans=0;
551     for (VP nowF:A)
552         for (int i=2;i<nowF.size();i++)
553             ans+=abs(getV(p,nowF[0],nowF[i-1],nowF[i]));
554     return ans/6;
555 }

```

```
556 VVP init(db INF) {
557     VVP pss(6,VP(4));
558     pss[0][0] = pss[1][0] = pss[2][0] = {-INF, -INF, -INF};
559     pss[0][3] = pss[1][1] = pss[5][2] = {-INF, -INF, INF};
560     pss[0][1] = pss[2][3] = pss[4][2] = {-INF, INF, -INF};
561     pss[0][2] = pss[5][3] = pss[4][1] = {-INF, INF, INF};
562     pss[1][3] = pss[2][1] = pss[3][2] = {INF, -INF, -INF};
563     pss[1][2] = pss[5][1] = pss[3][3] = {INF, -INF, INF};
564     pss[2][2] = pss[4][3] = pss[3][1] = {INF, INF, -INF};
565     pss[5][0] = pss[4][0] = pss[3][0] = {INF, INF, INF};
566     return pss;
567 }
```

7 Others

7.1 mint 类

```
1  const int mod = 998244353;
2
3  struct mint {
4      int n;
5      mint(int n_ = 0) : n(n_) {}
6  };
7
8  mint operator+(mint a, mint b) { return (a.n += b.n) >= mod ? a.n - mod : a.n; }
9  mint operator-(mint a, mint b) { return (a.n -= b.n) < 0 ? a.n + mod : a.n; }
10 mint operator*(mint a, mint b) { return 1LL * a.n * b.n % mod; }
11 mint &operator+=(mint &a, mint b) { return a = a + b; }
12 mint &operator-=(mint &a, mint b) { return a = a - b; }
13 mint &operator*=(mint &a, mint b) { return a = a * b; }
14 ostream &operator<<(ostream &o, mint a) { return o << a.n; }
```

7.2 不重叠区间贪心

```
1  #include <bits/stdc++.h>
2  #define ll long long
3  using namespace std;
4
5  const int maxn = 5e5+5;
6  pair<int, int> a[maxn];
7  int main() {
8      int n;
9      cin >> n;
10     for (int i = 1; i <= n; ++i) {
11         cin >> a[i].second >> a[i].first;
12     }
13     sort(a + 1, a + 1 + n);
14     int res = 1;
15     int tmp = a[1].first;
16     // printf("%d %d\n", a[1].second, a[1].first);
17     for (int i = 2; i <= n; ++i) {
18         if (a[i].second > tmp) {
19             res++;
20             // printf("%d %d\n", a[i].second, a[i].first);
21             tmp = a[i].first;
22         }
23     }
24     printf("%d\n", res);
25     return 0;
26 }
```

7.3 BigInt 类

```
1  const double PI = acos(-1.0);
2  struct Complex{
3      double x,y;
4      Complex(double _x = 0.0,double _y = 0.0){
5          x = _x;
6          y = _y;
7      }
```

```

8     Complex operator-(const Complex &b)const{
9         return Complex(x - b.x,y - b.y);
10    }
11    Complex operator+(const Complex &b)const{
12        return Complex(x + b.x,y + b.y);
13    }
14    Complex operator*(const Complex &b)const{
15        return Complex(x*b.x - y*b.y,x*b.y + y*b.x);
16    }
17 };
18 void change(Complex y[],int len){
19     int i,j,k;
20     for(int i = 1,j = len/2;i<len-1;i++){
21         if(i < j)    swap(y[i],y[j]);
22         k = len/2;
23         while(j >= k){
24             j = j - k;
25             k = k/2;
26         }
27         if(j < k)    j+=k;
28     }
29 }
30 void fft(Complex y[],int len,int on){
31     change(y,len);
32     for(int h = 2;h <= len;h<=1){
33         Complex wn(cos(on*2*PI/h),sin(on*2*PI/h));
34         for(int j = 0;j < len;j += h){
35             Complex w(1,0);
36             for(int k = j;k < j + h/2;k++){
37                 Complex u = y[k];
38                 Complex t = w*y[k + h/2];
39                 y[k] = u + t;
40                 y[k + h/2] = u - t;
41                 w = w*wn;
42             }
43         }
44     }
45     if(on == -1){
46         for(int i = 0;i < len;i++){
47             y[i].x /= len;
48         }
49     }
50 }
51 class BigInt
52 {
53     #define Value(x, nega) ((nega) ? -(x) : (x))
54     #define At(vec, index) ((index) < vec.size() ? vec[(index)] : 0)
55     static int absComp(const BigInt &lhs, const BigInt &rhs)
56     {
57         if (lhs.size() != rhs.size())
58             return lhs.size() < rhs.size() ? -1 : 1;
59         for (int i = lhs.size() - 1; i >= 0; --i)
60             if (lhs[i] != rhs[i])
61                 return lhs[i] < rhs[i] ? -1 : 1;
62         return 0;
63     }
64     using Long = long long;
65     const static int Exp = 9;
66     const static Long Mod = 1000000000;

```

```
67     mutable std::vector<Long> val;
68     mutable bool nega = false;
69     void trim() const
70     {
71         while (val.size() && val.back() == 0)
72             val.pop_back();
73         if (val.empty())
74             nega = false;
75     }
76     int size() const { return val.size(); }
77     Long &operator[](int index) const { return val[index]; }
78     Long &back() const { return val.back(); }
79     BigInt(int size, bool nega) : val(size), nega(nega) {}
80     BigInt(const std::vector<Long> &val, bool nega) : val(val), nega(nega) {}
81
82 public:
83     friend std::ostream &operator<<(std::ostream &os, const BigInt &n)
84     {
85         if (n.size())
86         {
87             if (n.nega)
88                 putchar('-');
89             for (int i = n.size() - 1; i >= 0; --i)
90             {
91                 if (i == n.size() - 1)
92                     printf("%lld", n[i]);
93                 else
94                     printf("%0*lld", n.Exp, n[i]);
95             }
96         }
97         else
98             putchar('0');
99         return os;
100     }
101     friend BigInt operator+(const BigInt &lhs, const BigInt &rhs)
102     {
103         BigInt ret(lhs);
104         return ret += rhs;
105     }
106     friend BigInt operator-(const BigInt &lhs, const BigInt &rhs)
107     {
108         BigInt ret(lhs);
109         return ret -= rhs;
110     }
111     BigInt(Long x = 0)
112     {
113         if (x < 0)
114             x = -x, nega = true;
115         while (x >= Mod)
116             val.push_back(x % Mod), x /= Mod;
117         if (x)
118             val.push_back(x);
119     }
120     BigInt(const char *s)
121     {
122         int bound = 0, pos;
123         if (s[0] == '-')
124             nega = true, bound = 1;
125         Long cur = 0, pow = 1;
```



```

126         for (pos = strlen(s) - 1; pos >= Exp + bound - 1; pos -= Exp, val.push_back(cur
127             ), cur = 0, pow = 1)
128             for (int i = pos; i > pos - Exp; --i)
129                 cur += (s[i] - '0') * pow, pow *= 10;
129         for (cur = 0, pow = 1; pos >= bound; --pos)
130             cur += (s[pos] - '0') * pow, pow *= 10;
131         if (cur)
132             val.push_back(cur);
133     }
134     BigInt &operator=(const char *s){
135         BigInt n(s);
136         *this = n;
137         return n;
138     }
139     BigInt &operator=(const Long x){
140         BigInt n(x);
141         *this = n;
142         return n;
143     }
144     friend std::istream &operator>>(std::istream &is, BigInt &n){
145         string s;
146         is >> s;
147         n=(char*)s.data();
148         return is;
149     }
150     BigInt &operator+=(const BigInt &rhs)
151     {
152         const int cap = std::max(size(), rhs.size()) + 1;
153         val.resize(cap);
154         int carry = 0;
155         for (int i = 0; i < cap - 1; ++i)
156         {
157             val[i] = Value(val[i], nega) + Value(At(rhs, i), rhs.nega) + carry, carry =
158                 0;
159             if (val[i] >= Mod)
160                 val[i] -= Mod, carry = 1;
161             else if (val[i] < 0)
162                 val[i] += Mod, carry = -1;
163         }
164         if ((val.back() = carry) == -1) //assert(val.back() == 1 or 0 or -1)
165         {
166             nega = true, val.pop_back();
167             bool tailZero = true;
168             for (int i = 0; i < cap - 1; ++i)
169             {
170                 if (tailZero && val[i])
171                     val[i] = Mod - val[i], tailZero = false;
172                 else
173                     val[i] = Mod - 1 - val[i];
174             }
175             trim();
176             return *this;
177         }
178     friend BigInt operator-(const BigInt &rhs)
179     {
180         BigInt ret(rhs);
181         ret.nega ^= 1;
182         return ret;

```

```

183     }
184     BigInt &operator--(const BigInt &rhs)
185     {
186         rhs.nega ^= 1;
187         *this += rhs;
188         rhs.nega ^= 1;
189         return *this;
190     }
191     friend BigInt operator*(const BigInt &lhs, const BigInt &rhs)
192     {
193         int len=1;
194         BigInt ll=lhs,rr=rhs;
195         ll.nega = lhs.nega ^ rhs.nega;
196         while(len<2*lhs.size()||len<2*rhs.size())len<=1;
197         ll.val.resize(len),rr.val.resize(len);
198         Complex x1[len],x2[len];
199         for(int i=0;i<len;i++){
200             Complex nx(ll[i],0.0),ny(rr[i],0.0);
201             x1[i]=nx;
202             x2[i]=ny;
203         }
204         fft(x1,len,1);
205         fft(x2,len,1);
206         for(int i = 0 ; i < len; i++)
207             x1[i] = x1[i] * x2[i];
208         fft( x1 , len , -1 );
209         for(int i = 0 ; i < len; i++)
210             ll[i] = int( x1[i].x + 0.5 );
211         for(int i = 0 ; i < len; i++){
212             ll[i+1]+=ll[i]/Mod;
213             ll[i]%=Mod;
214         }
215         ll.trim();
216         return ll;
217     }
218     friend BigInt operator*(const BigInt &lhs, const Long &x){
219         BigInt ret=lhs;
220         bool negat = ( x < 0 );
221         Long xx = (negat) ? -x : x;
222         ret.nega ^= negat;
223         ret.val.push_back(0);
224         ret.val.push_back(0);
225         for(int i = 0; i < ret.size(); i++)
226             ret[i]*=xx;
227         for(int i = 0; i < ret.size(); i++){
228             ret[i+1]+=ret[i]/Mod;
229             ret[i] %= Mod;
230         }
231         ret.trim();
232         return ret;
233     }
234     BigInt &operator*=(const BigInt &rhs) { return *this = *this * rhs; }
235     BigInt &operator*=(const Long &x) { return *this = *this * x; }
236     friend BigInt operator/(const BigInt &lhs, const BigInt &rhs)
237     {
238         static std::vector<BigInt> powTwo{BigInt(1)};
239         static std::vector<BigInt> estimate;
240         estimate.clear();
241         if (absComp(lhs, rhs) < 0)

```

```

242         return BigInt();
243     BigInt cur = rhs;
244     int cmp;
245     while ((cmp = absComp(cur, lhs)) <= 0)
246     {
247         estimate.push_back(cur), cur += cur;
248         if (estimate.size() >= powTwo.size())
249             powTwo.push_back(powTwo.back() + powTwo.back());
250     }
251     if (cmp == 0)
252         return BigInt(powTwo.back().val, lhs.nega ^ rhs.nega);
253     BigInt ret = powTwo[estimate.size() - 1];
254     cur = estimate[estimate.size() - 1];
255     for (int i = estimate.size() - 1; i >= 0 && cmp != 0; --i)
256         if ((cmp = absComp(cur + estimate[i], lhs)) <= 0)
257             cur += estimate[i], ret += powTwo[i];
258     ret.nega = lhs.nega ^ rhs.nega;
259     return ret;
260 }
261 friend BigInt operator/(const BigInt &num, const Long &x){
262     bool negat = ( x < 0 );
263     Long xx = (negat) ? -x : x;
264     BigInt ret;
265     Long k = 0;
266     ret.val.resize( num.size() );
267     ret.nega = (num.nega ^ negat);
268     for(int i = num.size() - 1 ; i >= 0; i--){
269         ret[i] = ( k * Mod + num[i]) / xx;
270         k = ( k * Mod + num[i]) % xx;
271     }
272     ret.trim();
273     return ret;
274 }
275 bool operator==(const BigInt &rhs) const
276 {
277     return nega == rhs.nega && val == rhs.val;
278 }
279 bool operator!=(const BigInt &rhs) const { return nega != rhs.nega || val != rhs.
    val; }
280 bool operator>=(const BigInt &rhs) const { return !(*this < rhs); }
281 bool operator>(const BigInt &rhs) const { return !(*this <= rhs); }
282 bool operator<=(const BigInt &rhs) const
283 {
284     if (nega && !rhs.nega)
285         return true;
286     if (!nega && rhs.nega)
287         return false;
288     int cmp = absComp(*this, rhs);
289     return nega ? cmp >= 0 : cmp <= 0;
290 }
291 bool operator<(const BigInt &rhs) const
292 {
293     if (nega && !rhs.nega)
294         return true;
295     if (!nega && rhs.nega)
296         return false;
297     return (absComp(*this, rhs) < 0) ^ nega;
298 }
299 void swap(const BigInt &rhs) const

```

```
300     {
301         std::swap(val, rhs.val);
302         std::swap(nega, rhs.nega);
303     }
304 };
305 BigInt ba,bb;
306 int main(){
307     cin>>ba>>bb;
308     cout << ba + bb << '\n';//和
309     cout << ba - bb << '\n';//差
310     cout << ba * bb << '\n';//积
311     BigInt d;
312     cout << (d = ba / bb) << '\n';//商
313     cout << ba - d * bb << '\n';//余
314     return 0;
315 }
```

7.4 date

```
1  string dayOfWeek[] = {"Mo", "Tu", "We", "Th", "Fr", "Sa", "Su"};
2  // converts Gregorian date to integer (Julian day number)
3  int DateToInt (int m, int d, int y){
4      return
5          1461 * (y + 4800 + (m - 14) / 12) / 4 +
6          367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
7          3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
8          d - 32075;
9  }
10
11 // converts integer (Julian day number) to Gregorian date: month/day/year
12 void IntToDate (int jd, int &m, int &d, int &y){
13     int x, n, i, j;
14     x = jd + 68569;
15     n = 4 * x / 146097;
16     x -= (146097 * n + 3) / 4;
17     i = (4000 * (x + 1)) / 1461001;
18     x -= 1461 * i / 4 - 31;
19     j = 80 * x / 2447;
20     d = x - 2447 * j / 80;
21     x = j / 11;
22     m = j + 2 - 12 * x;
23     y = 100 * (n - 49) + i + x;
24 }
25 // converts integer (Julian day number) to day of week
26 string IntToDay (int jd){
27     return dayOfWeek[jd % 7];
28 }
```

7.5 Frac 类

```
1  struct Frac {
2      ll a, b;
3      void getJian() {
4          ll gcd = abs(__gcd(a, b));
5          a /= gcd;
6          b /= gcd;
7          if (b < 0) {
```

```
8         a = -a;
9         b = -b;
10    }
11 }
12 Frac(1l a_ = 1, 1l b_ = 1) {
13     a = a_;
14     b = b_;
15     getJian();
16 }
17 Frac operator + (const Frac& oth) {
18     1l bt = b * oth.b;
19     1l at = a * oth.b + oth.a * b;
20     return Frac(at, bt);
21 }
22 Frac operator * (const Frac& oth) {
23     a *= oth.a;
24     b *= oth.b;
25     getJian();
26     return *this;
27 }
28 bool operator < (const Frac& oth) const {
29     return a * oth.b < b * oth.a;
30 }
31 bool operator == (const Frac& oth) const {
32     return a * oth.b == b * oth.a;
33 }
34 bool operator <= (const Frac& oth) const {
35     return a * oth.b <= b * oth.a;
36 }
37 };
```

7.6 模拟退火 (最小圆覆盖)

```
1 const int maxn = 1e5 + 10;
2 const double eps = 1e-8;
3 const double delta = 0.98;
4 const double inf = 1e18;
5
6 struct Point { double x, y; } p[maxn];
7
8 double dis(Point A, Point B) { return sqrt((A.x - B.x) * (A.x - B.x) + (A.y - B.y) * (A
    .y - B.y)); };
9
10 double Simulate_Annea(int n)
11 {
12     Point S;
13     S.x = S.y = 0;
14     double t = 1000;
15     double res = inf;
16     while(t > eps)
17     {
18         int k = 0;
19         for(int i = 0; i < n; i++) if(dis(S, p[i]) > dis(S, p[k])) k = i;
20         double d = dis(S, p[k]);
21         res = min(res, d);
22         S.x += (p[k].x - S.x) / d * t;
23         S.y += (p[k].y - S.y) / d * t;
24         t *= delta;
```

```
25     }
26     return res;
27 }
28
29 int main()
30 {
31     int n;
32     scanf("%d", &n);
33     for(int i = 0; i < n; i++) scanf("%lf%lf", &p[i].x, &p[i].y);
34     printf("%.3f\n", Simulate_Annea(n));
35     return 0;
36 }
```

7.7 string 类

```
1  const int maxn = 1005;
2  struct String{
3      int nex[maxn];
4      char x[maxn];
5      int len;
6      int getLength() {
7          return len;
8      }
9      void getNext() {
10         int n = len, i = 0, j = -1;
11         nex[0] = -1;
12         while (i < n) {
13             if (j == -1 || x[i] == x[j]) nex[++i] = ++j;
14             else j = -1;
15         }
16     }
17     void input() {
18         scanf("%s", x);
19         len = strlen(x);
20     }
21     void inputAndCal() {
22         scanf("%s", x);
23         len = strlen(x);
24         getNext();
25     }
26     void show() {
27         printf("%s\n", x);
28     }
29     bool operator < (const String&oth) const {
30         return strcmp(x, oth.x) < 0;
31     }
32     char operator [] (const int a) const {
33         return x[a];
34     }
35     bool substring(String b) { //b is the substring of a
36         int m = len, n = b.getLength();
37         int i = 0, j = 0;
38         while (i < m && j < n) {
39             if (j == -1 || x[i] == b[j]) ++i, ++j;
40             else j = b.nex[j];
41             if (j == n) return true;
42         }
43         return false;
44     }
```

```
44     }
45 };
```

7.8 前缀异或和

```
1 ll xor_sum(ll n) {
2     ll t=n&3;
3     if (t&1) return t/2ull^1;
4     return t/2ull^n;
5 }
```

7.9 约瑟夫环第 k 个

```
1 ll kth(ll n, ll m, ll k) { // n个人, m间隔, 第k个出列的人
2     if (m == 1) return k;
3     ll res = (m - 1) % (n - k + 1);
4     for (ll i = n - k + 2, stp = 0; i <= n; i += stp, res += stp * m) {
5         if (res + m >= i) {
6             res = (res + m) % i;
7             i++;
8             stp = 0;
9         } else {
10            stp = (i - res - 2) / (m - 1);
11            if (i + stp > n) {
12                res += (n - (i - 1)) * m;
13                break;
14            }
15        }
16    }
17    return res + 1;
18 }
19
20 ll dieInXturn(int n, int k, int x) { // n个人, m间隔, 第k个人出列时间
21     ll tmp = 0;
22     while (n) {
23         x = (x + n) % n;
24         if (k > n) x += (k - x - 1 + n - 1) / n * n;
25         if ((x + 1) % k == 0) {
26             tmp += (x + 1) / k;
27             break;
28         } else {
29             if (k > n) {
30                 tmp += x / k;
31                 ll ttmp = x;
32                 x = x - (x / n + 1) * (x / k) + (x + n) / n * n - k;
33                 n -= ttmp / k;
34             } else {
35                 tmp += n / k;
36                 x = x - x / k;
37                 x += n - n / k * k;
38                 n -= n / k;
39             }
40         }
41     }
42 }
43 return tmp;
44 }
```

7.10 二分

```
1 // a为二分数组, x为需要查找的数, 返回最左端和最右端
2 pair<int, int> F(vector<int> a, int x) {
3     int l = 0, r = a.size() - 1;
4     int lres = -1;
5     while (l <= r) {
6         int mid = l + r >> 1;
7         int tt = a[mid];
8         if (tt >= x) {
9             r = mid - 1;
10        } else if (tt < x) {
11            l = mid + 1;
12        }
13    }
14    if (l >= a.size() || a[l] != x) return make_pair(-1, -1);
15    lres = l;
16    l = 0, r = a.size() - 1;
17    while (l <= r) {
18        int mid = l + r >> 1;
19        int tt = a[mid];
20        if (tt > x) {
21            r = mid - 1;
22        } else if (tt <= x) {
23            l = mid + 1;
24        }
25    }
26    return make_pair(lres, r);
27 }
```

7.11 猛男 IO 挂

```
1 const int LEN = 100000;
2 struct fastio {
3     int it, len;
4     char s[LEN + 5];
5     fastio() {
6         it = len = 0;
7     }
8     char get() {
9         if (it < len) return s[it++];
10        it = 0, len = fread(s, 1, LEN, stdin);
11        return len ? s[it++] : EOF;
12    }
13    bool notend() {
14        char c;
15        for (c = get(); c == ' ' || c == '\n'; c = get());
16        if (it) it--;
17        return c != EOF;
18    }
19    void put(char c) {
20        if (it == LEN) fwrite(s, 1, LEN, stdout), it = 0;
21        s[it++] = c;
22    }
23    void flush() {
24        fwrite(s, 1, it, stdout);
25    }
26 } buff, bufo;
```



```

27 inline int getint() {
28     char c;
29     int res = 0, sig = 1;
30     for (c = buff.get(); c < '0' || c > '9'; c = buff.get()) if (c == '-') sig = -1;
31     for (; c >= '0' && c <= '9'; c = buff.get()) res = res * 10 + (c - '0');
32     return sig * res;
33 }
34 inline ll getll() {
35     char c;
36     ll res = 0, sig = 1;
37     for (c = buff.get(); c < '0' || c > '9'; c = buff.get()) if (c == '-') sig = -1;
38     for (; c >= '0' && c <= '9'; c = buff.get()) res = res * 10 + (c - '0');
39     return sig * res;
40 }
41 inline void putint(int x, char suf) {
42     if (!x) bufo.put('0');
43     else {
44         if (x < 0) bufo.put('-'), x = -x;
45         int k = 0;
46         char s[15];
47         while (x) {
48             s[++k] = x % 10 + '0';
49             x /= 10;
50         }
51         for (; k; k--) bufo.put(s[k]);
52     }
53     bufo.put(suf);
54 }
55 inline void putll(ll x, char suf) {
56     if (!x) bufo.put('0');
57     else {
58         if (x < 0) bufo.put('-'), x = -x;
59         int k = 0;
60         char s[25];
61         while (x) {
62             s[++k] = x % 10 + '0';
63             x /= 10;
64         }
65         for (; k; k--) bufo.put(s[k]);
66     }
67     bufo.put(suf);
68 }
69 inline char get_char() {
70     char c;
71     for (c = buff.get(); c == ' ' || c == '\n'; c = buff.get());
72     return c;
73 }

```

7.12 贪心结论

- 1 // n 个区间, 挪到使得某个点被所有区间覆盖需要的最少步数时, 选择的点是所有区间端点的中位数 ($mid-mid+1$ 答案都是一样的)
- 2
- 3 // $2 * n$ 的格子填数, 使得列上的最大和最小的填充方法: 最大配最小, 次大配次小, 以此类推
- 4
- 5 // n 个数, 重排后使得相同位置上数不同的最大值为: 如果 $max \leq (sum - max)$ 则为 sum , 如果 $max > (sum - max)$ 则为 $2 * (sum - max)$
- 6

```

7 // 不重叠区间贪心
8 pair<int, int> a[maxn];
9 int main() {
10     int n;
11     cin >> n;
12     for (int i = 1; i <= n; ++i) {
13         cin >> a[i].second >> a[i].first;
14     }
15     sort(a + 1, a + 1 + n);
16     int res = 1;
17     int tmp = a[1].first;
18     // printf("%d %d\n", a[1].second, a[1].first);
19     for (int i = 2; i <= n; ++i) {
20         if (a[i].second > tmp) {
21             res++;
22             // printf("%d %d\n", a[i].second, a[i].first);
23             tmp = a[i].first;
24         }
25     }
26     printf("%d\n", res);
27     return 0;
28 }

```

7.13 builtin

```

1 __builtin_popcount(unsigned int n) // 1的个数
2 __builtin_parity(unsigned int n) // 奇数个1返回1, 偶数个返回0
3 __builtin_ctz(unsigned int n) // 判断n的二进制末尾后面0的个数
4 __builtin_clz(unsigned int n) // 返回前导0的个数

```

7.14 n 以内 k 因子的个数

```

1 // 返回1~n中k因子的个数
2 ll dig(ll n, ll k) {
3     if (k == 1) return n;
4     ll res = 0;
5     while (n > 0) {
6         res += n / k;
7         n /= k;
8     }
9     return res;
10 }

```

7.15 每个点左右两边最长不重子序列

```

1 int r = 1;
2 for (int i = 1; i <= n; ++i) {
3     while (r <= n && !vis[a[r]]) vis[a[r++]] = 1;
4     vis[a[i]] = 0;
5     R[i] = r - 1;
6 }
7 int l = n;
8 for (int i = n; i >= 1; --i) {
9     while (l >= 1 && !vis[a[l]]) vis[a[l--]] = 1;
10    vis[a[i]] = 0;
11    L[i] = l + 1;
12 }

```